# Group actions on models

## G. A. Polymath

ABSTRACT. In a nutshell, a set is *strongly syndetic* if the intersection of any finitely many of its translation is syndetic (a.k.a. generic). I do not know if the notion has been studied in the context presented here. To demostrate its interest I use it for a short proof of (a generalization of) Newelski's theorem on the diamter of the Lascar graph, see Theorem 12.

I only recently realized the connnections with topological dynamics (though, I see, it is written on the wall). I only adopted some terminologi and did not investigate further.

Theorem 14 shows that the phenomenon  $strongly \ syndetic = syndetic$  has some interest (it is reminiscent of forking = dividing).

In the last section I try to apply these notions to definable groups hoping to recover part of the stable group theory but I fail. I think I am missing something important (or commit errors somewhere).

#### 1. The two perspectives on the invariance of types

This section I review well-known matter and set the terminology.

Below  $\Delta \subseteq L_{xz}(\mathcal{U})$ ,  $\mathfrak{X} \subseteq \mathcal{U}^x$ , and  $\mathfrak{Z} \subseteq \mathcal{U}^z$  are some arbitrary nonempty sets (at some point we will require that  $\mathfrak{X}$  and  $\mathfrak{Z}$  are type-definable). We write  $L_{\Delta}(\mathfrak{Z})$  for the set of formulas  $\theta(x)$  that are Boolean combination of formulas  $\varphi(x)$  for some  $\varphi(x) \in \Delta$  and some  $b \in \mathfrak{Z}$ . Such formulas a called  $\Delta$ -formulas. A relatively  $\Delta$ -definable set is a set of the form  $\theta(\mathfrak{X})$  for some  $\Delta$ -formula  $\theta(x)$ . Subsets of  $L_{\Delta}(\mathfrak{Z})$  are called  $\Delta$ -types. We write  $S_{\Delta}(\mathfrak{Z})$  for the set of complete  $\Delta$ -types with parameters in  $\mathfrak{Z}$ . Beware that there may be other parameters hidden ion  $\Delta$ ).

**1 Assumption** Let *G* be a group that acts on  $\mathfrak{X}$  and on  $\mathfrak{Z}$  from the left. We require that for every  $\varphi(x;z) \in \Delta$  the set  $\varphi(\mathfrak{X};\mathfrak{Z})$  is invariant under the action of *G*.

Let  $\mathcal{D} \subseteq \mathcal{U}^z$ . We say that  $\mathcal{D}$  is invariant under the action of G, or G-invariant, if  $\mathcal{D} \cap \mathcal{Z}$  is fixed setwise by G. Yet in other words, if

is1. 
$$a \in \mathcal{D} \leftrightarrow ga \in \mathcal{D}$$
 for every  $a \in \mathcal{Z}$  and every  $g \in G$ .

A formula is invariant if the set it defines is invariant. We say that  $p(x) \subseteq L(\mathcal{U})$  is invariant under the action of G, or G-invariant, if for every formula  $\varphi(x;z) \in L$ 

it1. 
$$\varphi(x;a) \in p \Leftrightarrow \varphi(x;ga) \in p$$
 for every  $a \in \mathbb{Z}$  and every  $g \in G$ .

It should be evident that invariant under the action of  $\operatorname{Aut}(\mathcal{U}/A)$  coincides with invariant over A and Lascar invariant over A coincides with invariant under the action of  $\operatorname{Aut}(\mathcal{U}/A)$ .

1

We have just defined invariance using the subsets of  $\mathbb{Z}$  (externally) defined by p. Now we discuss invariance using the subsets of  $\mathfrak{X}$  that are in p.

An immediate consequence of Assumption 1 is that any G-translate of a  $\Delta$ -definable set is again  $\Delta$ -definable. In particular for every  $\Delta$ -formula  $\vartheta(x;\bar{b})$  and every  $g \in G$ 

$$g[\vartheta(\mathfrak{X};\bar{b})] = \vartheta(\mathfrak{X};g\bar{b}).$$

Therefore  $p(\mathbf{x}) \subseteq L_{\Delta}(\mathcal{Z})$  is invariant if

 $p(x) \vdash x \in \mathcal{D} \iff p(x) \vdash x \in g\mathcal{D}$  for every  $\Delta$ -definable  $\mathcal{D} \subseteq \mathcal{U}^x$  and  $g \in G$ ,

where by  $p(x) \vdash x \in \mathcal{D}$  we understand  $\vartheta(\mathfrak{X}) \subseteq \mathcal{D}$  for some  $\vartheta(x)$  that is conjunction of formulas in p(x).

A set  $\mathcal{D} \subseteq \mathcal{X}$  is syndetic under the action of G, or G-syndetic for short, if finitely many G-translates of  $\mathcal{D}$  cover  $\mathcal{X}$ ; we say n-G-syndetic if  $\leq n$  translates suffices. Dually, we say that  $\mathcal{D}$  is thick under the action of G, or G-thick for short, if the intersection of any finitely many G-translates of  $\mathbb{D}$  is nonempty; we say n-G-thick when the request is limited to  $\leq n$  translates. We may drop reference to G when it is clear from the context.

The same properties may be attributed to formulas (as these are identified with the set they define). When these properties are attributed to a type p(x), we understand that they hold for every conjunction of formulas in p(x).

The terminology is taken from topological dynamics (I hope correctly). Model theorists say generic for syndetic. In [1] the authors write quasi-non-dividing for thick when  $G = Aut(\mathcal{U}/A)$ .

**2 Example** If  $p(x) \subseteq L(\mathcal{U})$  is finitely satisfiable in A then p(x) is thick (in any  $\mathfrak{X} \supseteq$  $A^x$ ) under the action of Aut( $\mathcal{U}/A$ ). In fact, the same  $a \in A^x$  that satisfies  $\varphi(x)$  also satisfies every Aut( $\mathcal{U}/A$ )-translate of  $\varphi(x)$ .

In this chapter many proofs require some juggling with negations.

- 3 Fact (Assume 1) The following are equivalent
  - 1.  $\mathcal{D}$  is not *n*-syndetic
  - 2.  $\neg \mathbb{D}$  is *n*-thick.

**Proof.** Immediate by spelling out the definitions

1. there are no 
$$g_1, \ldots, g_n \in G$$
 such that  $\mathfrak{X} \subseteq \bigcup_{i=1}^n g_i \mathfrak{D}$   
2.  $\varnothing \neq \mathfrak{X} \cap \bigcap_{i=1}^n \neg g_i \mathfrak{D}$  for every  $g_1, \ldots, g_n \in G$ .

**4 Theorem** (Assume 1) Let  $p(x) \in S_{\Delta}(\mathcal{Z})$  be finitely satisfiable in  $\mathcal{X}$ . Then the following are equivalent

- 1. p(x) is invariant
- 2.  $p(x) \vdash x \in \mathcal{D}$  for every syndetic relatively  $\Delta$ -definable set  $\mathcal{D}$
- 3. p(x) is thick.

**Proof.**  $1\Rightarrow 2$ . Let  $g_1\mathcal{D}, \ldots, g_n\mathcal{D}$  be translations of  $\mathcal{D}$  that cover  $\mathcal{X}$ . Negate 2. By completeness,  $p(x) \vdash x \notin \mathcal{D}$ . Hence, from invariance we obtain

$$p(\mathbf{x}) \vdash \mathbf{x} \notin \bigcup_{i=1}^n g_i \mathbb{D}.$$

This contradicts the finite satisfiability of p(x) in X.

2⇒3. Let  $\mathfrak{D}$  be defined by a conjunction of formulas in p(x). If  $\mathfrak{D}$  is not thick then, by Fact 3,  $\neg \mathfrak{D}$  is syndetic. By 2,  $p(x) \vdash x \notin \mathfrak{D}$ , a contradiction.

 $3\Rightarrow 1$ . If p(x) is not invariant then, by completeness,  $p(x) \vdash \varphi(x;b) \land \neg \varphi(x;gb)$  for some  $g \in G$ . Clearly  $\varphi(x;b) \land \neg \varphi(x;gb)$  is not thick as it is inconsistent with its g-translate.

5 Remark In the theorem above, 2 and 3 can be replaced by

2'.  $p(x) \vdash x \in \mathcal{D}$  for every 2-syndetic relatively  $\Delta$ -definable set  $\mathcal{D}$ 

3'. p(x) is 2-thick.

The theorem yields an immediate necessary condition for the existence of an invariant global  $\Delta$ -type.

- **6 Corollary** (Assume 1) If there exists an invariant global Δ-type then for every  $\Delta$ -definable set  $\mathfrak{D}$ 
  - 1.  $\mathcal{D}$  and  $\neg \mathcal{D}$  cannot be both syndetic
  - 2. if  $\mathcal{D}$  is syndetic than it is thick.

**Proof.** By Fact 3, 1 and 2 are equivalent; 1 is an immediate consequenc of 2 of Theorem 4.  $\Box$ 

The following theorem gives a necessary and sufficient condition for the existence of global invariant  $\Delta$ -type. Ideally, we would like to prove that every thick  $\Delta$ -type extends to a global persitent type. Unfortunately this is not true – we need a stronger property. A  $\Delta$ -definable set  ${\mathbb D}$  is hereditarely thick if every finite cover of  ${\mathbb D}$  by  $\Delta$ -definable sets contains a thick set. A type is hereditarely thick if every conjunction of formulas in the type is hereditarely thick.

⚠ I do not know the term used in topological dynamics for hereditarely thick – the terminology is provisional. In [1] a related property is called *quasi-non-forking*.

**7 Theorem** (Assume 1) Let  $q(x) \subseteq L(\mathcal{U})$ . Then the following are equivalent

- 1. q(x) extends to an invariant type  $p(x) \in S_{\Delta}(\mathcal{Z})$  finitely satisfiable in  $\mathcal{X}$
- 2. q(x) is hereditarely thick.

**Proof.**  $1\Rightarrow 2$ . Let  $\vartheta(x)$  be a conjunction of formulas in q(x). Suppose  $\mathfrak{C}_1, \ldots, \mathfrak{C}_n$  cover  $\vartheta(\mathfrak{U}^x)$  and pick p(x) as in 1. By completeness,  $p(x) \vdash x \in \mathfrak{C}_i$  for some i. Then, by Theorem 4,  $\neg \mathfrak{C}_i$  is not syndetic. Therefore, by Fact 3,  $\mathfrak{C}_i$  is thick.

2⇒1. Let p(x) be maximal among the  $\Delta$ -types that contain q(x) and are such that  $\vartheta(\mathcal{U}^x)$  is hereditarely thick for every  $\vartheta(x)$  that is conjunction of formulas in p(x). We claim that p is a complete  $\Delta$ -type. Suppose for a contradiction that  $\vartheta(x)$ ,  $\neg\vartheta(x) \notin p$ . By maximality there is some formula  $\psi(x)$ , a conjunction of formulas in p(x) and

some  $\mathcal{C}_1, \ldots, \mathcal{C}_n$  that cover both  $\psi(\mathcal{U}^x) \cap \vartheta(\mathcal{U}^x)$  and  $\psi(\mathcal{U}^x) \setminus \vartheta(\mathcal{U}^x)$  and such that no  $\mathcal{C}_i$  is thick. As  $\mathcal{C}_1, \ldots, \mathcal{C}_n$  cover  $\psi(\mathcal{U}^x)$  this is a contradiction. It is only left to show that p(x) is finitely satisfiable in  $\mathcal{X}$  and invariant. Finite satisfiability follows from thickness. From completeness and Theorem 4 we obtain invariance.

## 2. Strong syndeticty

Unfortunatelly, genericy is not preserved under intersection. To obtain closure under intersection, we need to push the concept to a higher level of complexity. This is only required in Section 3.

A set  $\mathcal{D} \subseteq \mathcal{U}^x$  is strongly syndetic if the intersection of  $\mathcal{D}$  with any of its translations  $\mathcal{D}$  is syndetic. Dually, we say that  $\mathcal{D}$  is weakly thick if the union of  $\mathcal{D}$  with one of its translations is peristent. Again, the same properties may be attributed to formulas and types.

Notation: for  $\mathcal{B} \subseteq \mathcal{X}$  and  $H \subseteq G$  we write  $H \mathcal{B}$  for  $\{h\mathcal{B} : h \in H\}$ .

8 Lemma (Assume 1) The intersection of strongly syndetic sets is strongly syndetic.

**Proof.** We may assume that all sets mentioned below are subsets of  $\mathfrak{X}$ . Let  $\mathfrak{D}$  and  $\mathfrak{C}$  be strongly syndetic and let  $K \subseteq G$  be an arbitrary finite set. It suffices to prove that  $\mathfrak{B} = \cap K \ (\mathfrak{C} \cap \mathfrak{D})$  is syndetic. Clearly  $\mathfrak{B} = \mathfrak{C}' \cap \mathfrak{D}'$ , where  $\mathfrak{C}' = \cap K \ \mathfrak{C}$  and  $\mathfrak{D}' = \cap K \ \mathfrak{D}$ . Note that  $\mathfrak{C}'$  and  $\mathfrak{D}'$  are both strongly syndetic. In particular  $\mathfrak{X} = \cup H \ \mathfrak{D}'$  for some finite  $H \subseteq G$ . Now, from

$$\begin{array}{rcl}
\cup H \, \mathfrak{B} &=& \cup H \Big[ \mathfrak{C}' \, \cap \, \mathfrak{D}' \Big] \\
\cup H \, \mathfrak{B} &\supseteq & \cup H \Big[ \big( \cap H \, \mathfrak{C}' \big) \, \cap \, \mathfrak{D}' \Big] \\
&=& \big( \cap H \, \mathfrak{C}' \big) \, \cap \, \big( \cup H \, \mathfrak{D}' \big) \\
&=& \cap H \, \mathfrak{C}'
\end{array}$$

As  $\mathfrak{C}'$  is strongly syndetic,  $\cap H \mathfrak{C}'$  is syndetic. Therefore  $\cup H \mathfrak{B}$  is also syndetic. The syndeticty of  $\mathfrak{B}$  follows.

**9 Corollary** (Assume 1) Let  $q(x) = \{\vartheta(x) \in L_{\varphi}(\mathcal{U}) : \vartheta(x) \text{ strongly syndetic}\}$ . Then q(x) is finitely satisfiable in  $\mathcal{X}$ , strongly syndetic, and invariant.

**Proof.** Strong syndeticty is an immediate consequence of Lemma 8. Finite satisfiability follows easily from syndeticty. As for invariance, note that any translate of a strongly syndetic formula is also strongly syndetic.

**10 Corollary** (Assume 1) Let q(x) be as in Corollary 9. Let  $p(x) \subseteq L(\mathcal{U})$  be such that  $p(x) \cup q(x)$  is finitely satisfied in  $\mathcal{X}$ . Then p(x) is weakly thick.

**Proof.** Let  $\vartheta(x) \in p$ . As q(x) is finitely satisfiable in  $\vartheta(\mathcal{U}^x)$ , we cannot have that  $\neg \vartheta(x)$  is strongly syndetic. From Fact 3, we obtain that  $\neg \vartheta(\mathcal{U}^x)$  non strongly syndetic is equivalent to  $\vartheta(x)$  weakly thick.

#### 3. The diameter of a Lascar type

Recall that  $\mathcal{L}(a/A)$ , the Lascar strong type of  $a \in \mathcal{U}^x$ , is the union of a chain of type-definable sets of the form  $\{x: d_A(a,x) \leq n\}$ . In this section we prove that  $\mathcal{L}(a/A)$  is type-definable (if and) only this chain is finite. In other words, only if the connected component of a in the Lascar graph has finite diameter. It is convenient to address the problem in more general terms.

Let *G* be a normal subgroup of Aut( $\mathcal{U}$ ). Assume *G* is the union of a countable chain of sets  $\langle G_n : n \in \omega \rangle$  with the following properties

- 1. every  $G_n$  is symmetric i.e. it contains the unit and is closed under inverse
- 2. every  $G_n$  is conjugancy invariant i.e.  $g G_n g^{-1} = G_n$  for every  $g \in G$
- 3.  $G_nG_m \subseteq G_{n+m}$  for every  $n, m \in \omega$ .

Assume G acts transitively on  $\mathfrak{X}$  i.e.,  $Ga = \mathfrak{X}$  for every  $a \in \mathfrak{X}$ . We define a discrete metric on  $\mathfrak{X}$ . For  $a,b \in \mathfrak{X}$  let d(a,b) be the minimal n such that  $a \in G_nb$ . This defines a metric by 1 and 3. By 2, this metric in G-invariant. The diameter of a set  $\mathfrak{C} \subseteq \mathfrak{X}$  is the supremum of d(a,b) for  $a,b \in \mathfrak{C}$ .

We are interested in sufficient conditions for  $\mathfrak X$  to have finite diameter. The notions introduced in Section 2 offer some hints.

**11 Proposition** If X has a weakly thick subset of finite diameter, then X itself has finite diameter.

**Proof.** Let  $\mathcal{C} \subseteq \mathcal{X}$  be a weakly thick set of diameter n. Let  $H \subseteq G$  be finite such that  $\cup H$   $\mathcal{C}$  is thick. We claim that also  $\cup H$   $\mathcal{C}$  has finite diameter. Let  $a \in \mathcal{C}$  be arbitrary. Let m be larger than d(ha, ka) for all  $h, k, \in H$ . Now, let hb and kc, for some  $h, k, \in H$  and  $b, c \in \mathcal{C}$ , be two arbitrary elements of  $\cup H$   $\mathcal{C}$ . As  $h\mathcal{C}$  and  $k\mathcal{C}$  have the same diameter of  $\mathcal{C}$ ,

$$d(hb, kc) \leq d(hb, ha) + d(ha, ka) + d(ka, kc)$$
  
$$< n + m + n.$$

This proves that  $\cup$  H  $\mathfrak{C}$  has finite diameter. Therefore, without loss of generality, we may assume that  $\mathfrak{C}$  itself is thick.

By the transitivity of the action, any two elements of  $\mathfrak{X}$  are of the form ha, ka for some  $h, k \in G$  and some  $a \in \mathfrak{C}$ . By percistency, there are  $c \in \mathfrak{C} \cap h\mathfrak{C}$  and  $d \in \mathfrak{C} \cap k\mathfrak{C}$ . Then

$$d(ha, ka) \leq d(ha, c) + d(c, d) + d(d, ka)$$
  
$$\leq n + n + n.$$

Therefore the diameter of  $\mathcal{X}$  does not exceed 3n.

**12 Theorem** Suppose that  $\mathfrak{X}$  and the sets  $\mathfrak{X}_n = G_n a$ , for some  $a \in \mathfrak{X}$ , are type-definable. Then  $\mathfrak{X}$  has finite diameter.

**Proof.** By Proposition 11, it suffices to prove that  $\mathcal{X}_n$  is weakly thick. Define

$$q(x) = \{\vartheta(x) \in L(\mathcal{U}) : \vartheta(x) \text{ strongly syndetic}\}.$$

By Corollary 10, with L for  $L_{\varphi}$ , it suffices to prove that for some n the type q(x) is finitely satisfied in  $\mathfrak{X}_n$ . Suppose not. Let  $\psi_n(x) \in q$  be a formula that is not satisfied in  $\mathfrak{X}_n$ . The type  $\{\psi_n(x) : n \in \omega\}$  is finitely satisfied in  $\mathfrak{X}$ . Then it has a realization in  $\mathfrak{X}$ . As this realization belongs to some  $\mathfrak{X}_n$  this contradicts the definition of  $\psi_n(x)$ .

**13 Example** Let  $G_1 \subseteq \operatorname{Aut}(\mathcal{U}/A)$  be the set of automorphisms that fix a model containing A. Thenthe group G generated by  $G_1$  is  $\operatorname{Autf}(\mathcal{U}/A)$  and  $G \cdot a = \mathcal{X}$  is  $\mathcal{L}(a/A)$ . If we let  $G_1$  be the set of automorphisms that fix a model  $M \supseteq A$  and  $G_n = G_1^n$ , then d(a,b) concides with the dinstance in the Lascar graph. It is not difficult to see that the sets  $G_n \cdot a$  are type definable. Then from Theorem 12 it follows that  $\mathcal{L}(a/A)$  is type definable (if and) only if it has a finite diameter.

## 4. A simplified landscape

Under suitable assumptions – e.g. the sability of  $\varphi(x;z)$  – some of the notions introduced above coalesce and we are left with cleaner theory. We prove the following theorem.

- **14 Theorem** (Assume 1) The following are equivalent
  - 1. thick  $\Delta$ -definable sets are hereditarely thick
  - 2. syndetic  $\Delta$ -definable sets are strongly syndetic
  - 3. syndetic  $\Delta$ -definable sets are closed under intersection
  - 4. weakly persisent  $\Delta$ -definable sets are thick.

**Proof.**  $1\Rightarrow 2$ . It suffices to prove that syndetic sets are closed under intersection. Let  $\mathbb{C}$  and  $\mathbb{D}$  be syndetic  $\Delta$ -definable sets. Suppose for a contradiction that  $\mathbb{C}\cap\mathbb{D}$  is not syndetic. By 1 and Theorem 7 there is an invariant global  $\Delta$ -type p(x) containing  $x \in \neg \mathbb{C} \cup \neg \mathbb{D}$ . By completeness either  $p(x) \vdash x \in \neg \mathbb{C}$  or  $p(x) \vdash x \in \neg \mathbb{D}$ . This is a contradiction because, by Theorem 4,  $p(x) \vdash x \in \mathbb{C}$  and  $p(x) \vdash x \in \mathbb{D}$ .

2⇔3⇔4. Clear.

4⇒1. Let  $q(x) = \{\vartheta(x) \in L_{\varphi}(\mathfrak{U}) : \vartheta(x) \text{ syndetic}\}$ . By 2 this is the same type defined in Corollary 10. Therefore, any completion of q(x) is, by 4, thick. Let  $\mathfrak{D}$  be a thick Δ-definable set. By Theorems 4 and 7 it suffices to show that  $\mathfrak{D}$  is consistent with q(x). Suppose not, then  $q(x) \vdash x \in \neg \mathfrak{D}$ . Therefore, by 3,  $\neg \mathfrak{D}$  is syndetic. This is a contradiction by Fact 3.

### 5. Definable groups

In this section we set  $G = \mathbb{Z}$  and require that  $\mathbb{Z}$  and  $\mathbb{X}$  are type-definable over A. We assume that the group operations and the group action are definable over A. We use the symbol  $\cdot$  for both the group multiplication and the group action. Clearly,  $\mathbb{Z}$  also acts on itself by left multiplication.

In this section we deal with the actions of two groups:  $\mathcal{Z}$  and  $\operatorname{Aut}(\mathcal{U}/A)$ . Syndetic and thick only refer to the action of  $\mathcal{Z}$ . We will always be explicit about invariance.

Let  $\psi(x;y) \in L(A)$ . We write  $\varphi(x;z;y)$  for the formula  $\psi(z^{-1} \cdot x;y)$ . In this section  $\Delta$  contains the formulas  $\varphi(x;z;a)$  where a ranges over the realizations of a

given  $q(y) \in S(A)$ . Note that  $\varphi(\mathfrak{X}; \mathcal{Z}; a)$  is invariant under the action of  $\mathcal{Z}$  and  $\varphi(\mathfrak{X}; g; a) = g \cdot \varphi(\mathfrak{X}; 1; a)$  for every a. Therefore the above G and  $\Delta$  satisfy Assumption 1.

Let 1 be the identity of  $\mathbb{Z}$  which, for simplicity, we assume is a constant of L.

**15 Fact** There are some formulas  $\gamma(y)$ ,  $\pi(y) \in L(A)$  such that, for every  $a \in \mathcal{U}^y$ 

$$\gamma(a) \Leftrightarrow \varphi(x;1;a)$$
 is syndetic

$$\pi(a) \Leftrightarrow \varphi(x;1;a)$$
 is thick.

Similar claims holds for  $\neg \varphi(x; 1; a)$ .

**Proof.** By an easy argument of compactness, there is an n such that if  $\varphi(x;1;a)$  is syndetic then it is also n-syndetic. Then

$$\gamma(y) = \exists z_1, \dots, z_n \ \forall x \bigvee_{i=1}^n \varphi(x; z_i; y)$$

The other claims follow easily.

As  $\varphi(x;g;a)$  is syndetic/thick if and only if  $\varphi(x;1;a)$  is such,  $\operatorname{Aut}(\mathcal{U}/A)$  maps syndetic/thick sets of the form  $\varphi(X;g;a)$  to syndetic/thick sets of the same form.

**16 Fact** If  $p(x) \in S_{\Delta}(\mathcal{Z})$  is thick (equivalently,  $\mathcal{Z}$ -invariant) then it is invariant over A – i.e. under the action of  $\operatorname{Aut}(\mathcal{U}/A)$ .

**Proof.** By Theorem 4, p(x) is  $\mathbb{Z}$ -invariant. Therefore  $\varphi(x;g;a) \in p$  if and only if  $\varphi(x;1;a) \in p$ . As  $\varphi(x;1;a) \leftrightarrow \varphi(x;1;fa)$  for every  $f \in \operatorname{Aut}(\mathbb{U}/A)$ , invariance over A follows.

**17 Assumption** We assume that the three equivalent conditions in Theorem 14 hold for  $G = \mathcal{Z}$ , X, and  $\Delta$  introduced in this section.

Under this assumption there is a unique maximal syndetic type  $q(x) \subseteq L_{\Delta}(z)$ . The types  $p(x) \in S_{\Delta}(z)$  that extend q(x) are thick, hence invariant under the action of both z and Aut(U/A).

## 6. Notes and references

[1] Artem Chernikov and Itay Kaplan, Forking and dividing in NTP<sub>2</sub> theories, J. Symbolic Logic 77 (2012), 1–20.