# Group actions on models

ABSTRACT. A set is *strongly generic* if the intersection of any finitely many of its translations is generic.

Theorem 18 shows that the condition *strongly generic* = *generic* is robust. It might be of some interest (it is reminiscent of *forking* = *dividing*). Does it have interesting examples?

Proposition 26 looks interesting. Can it be strengthened to obtain a type-definable group?

The connections with topological dynamics have not been explored :(

### 1. The dual perspective on invariance

The notions in this section are well-known but sometimes the terminology differs.

In this chapter  $\Delta \subseteq L_{xz}(\mathcal{U})$ . Let  $\mathcal{Z} \subseteq \mathcal{U}^z$ . We write  $\Delta(\mathcal{Z})$  for the set of formulas of the form  $\varphi(x;b)$  for some  $\varphi(x;z) \in \Delta$  and some  $b \in \mathcal{Z}$ . We write  $\Delta^{\pm}(\mathcal{Z})$  for the set of formulas in  $\Delta(\mathcal{Z})$  or negation thereof. Furthermore, we write  $S_{\Delta}(\mathcal{Z})$  for the set of complete  $\Delta^{\pm}(\mathcal{Z})$ -types.

We write  $\Delta^{\mathbb{B}}(\mathbb{Z})$  for the set of Boolean combinations of formulas in  $\Delta(\mathbb{Z})$ .

**1 Assumption** Let G be a group that acts on some sets  $\mathfrak{X} \subseteq \mathcal{U}^x$  and  $\mathfrak{Z} \subseteq \mathcal{U}^z$ . We require that for every  $\varphi(x;z) \in \Delta$  the set  $\varphi(\mathfrak{X};\mathfrak{Z})$  is invariant under the action of G. For convenience, we will assume that G is the identity outside  $\mathfrak{X}$  and  $\mathfrak{Z}$ .

When  $p(x) \subseteq \Delta^{\mathbb{B}}(\mathcal{Z})$  and  $\mathfrak{D} \subseteq \mathcal{U}^x$  we write  $p(x) \vdash x \in \mathfrak{D}$  if the inclusion  $\psi(\mathfrak{X}) \subseteq \mathfrak{D}$  for some  $\psi(x)$  that is conjunctions of formulas in p(x).

Let  $\mathcal{D} \subseteq \mathcal{U}^z$ . We say that  $\mathcal{D}$  is invariant under the action of G, or G-invariant, if  $\mathcal{D}$  is fixed setwise by G. That is,  $g\mathcal{D} = \mathcal{D}$  for every  $g \in G$ . Yet in other words, if

is1. 
$$a \in \mathcal{D} \leftrightarrow ga \in \mathcal{D}$$
 for every  $a \in \mathcal{Z}$  and every  $g \in G$ .

A formula is *G*-invariant if the set it defines is *G*-invariant. We say that  $p(x) \subseteq \Delta^{B}(\mathbb{Z})$  is invariant under the action of *G*, or *G*-invariant, if for every  $\Delta^{B}(\mathbb{Z})$ -formula  $\vartheta(x;\bar{a})$ .

it1. 
$$\vartheta(x;\bar{a}) \in p \Leftrightarrow \vartheta(x;g\bar{a}) \in p$$
 for every  $g \in G$ .

It should be clear that invariant under the action of Aut(U/A) coincides with invariant over A and Lascar invariant over A coincides with invariant under the action of Autf(U/A).

Note that p(x) is G-invariant exactly when the sets  $\mathcal{D}_{p,\vartheta} = \{b : \vartheta(x;b) \in p\} \subseteq \mathcal{U}^z$  are. Now we would like to discuss invariance using sets  $\mathcal{D} \subseteq \mathcal{U}^x$ .

An immediate consequence of the invariance of  $\varphi(\mathfrak{X}; \mathbb{Z})$  is that any G-translate of a  $\Delta^{\mathbb{B}}(\mathbb{Z})$ -definable set is again  $\Delta^{\mathbb{B}}(\mathbb{Z})$ -definable. In particular for every  $\Delta^{\mathbb{B}}(\mathbb{Z})$ -formula  $\vartheta(x; \bar{b})$  and every  $g \in G$ 

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$$g[\vartheta(\mathbf{X};\bar{b})] = \vartheta(\mathbf{X};g\bar{b}).$$

Therefore, a type  $p(x) \subseteq \Delta^{B}(\mathcal{Z})$  is *G*-invariant if

$$p(x) \vdash x \in \mathcal{D} \iff p(x) \vdash x \in g \cdot \mathcal{D} \text{ for every } \Delta^{\mathsf{B}}(\mathcal{Z})\text{-definable } \mathcal{D} \subseteq \mathcal{X} \text{ and } g \in G.$$

A set  $\mathcal{D} \subseteq \mathcal{X}$  is **generic** under the action of G, or G-generic for short, if finitely many G-translates of  $\mathcal{D}$  cover  $\mathcal{X}$ ; we say n-G-generic if  $\leq n$  translates suffices. Dually, we say that  $\mathcal{D}$  is **persistent** under the action of G, or G-persistent for short, if the intersection of any finitely many G-translates of  $\mathcal{D}$  is nonempty; we say n-G-persistent when the request is limited to  $\leq n$  translates. When  $\mathcal{X}$  and  $\mathcal{X}$  are not clear from the context, we say that these notions are **relative** to  $\mathcal{X}$  and  $\mathcal{X}$ .

The same properties may be attributed to formulas (as these are identified with the set they define). When these properties are attributed to a type p(x), we understand that they hold for every conjunction of formulas in p(x).

The terminology above is non-standard. In [1] the authors write *quasi-non-dividing* for *persistent* under the action of Aut(U/A). Their terminology has good motivations, but it would be a mouthful if adapted to our context. In topological dynamics similar notions have been introduced with different terminology: *syndetic* corresponds to *generic* and *thick* corresponds to *persistent*.

**2 Example** If  $p(x) \subseteq L(\mathcal{U})$  is finitely satisfiable in A then p(x) is persistent under the action of  $G = \operatorname{Aut}(\mathcal{U}/A)$  relative to any  $\mathcal{X} \supseteq A^x$ . In fact, the same  $a \in A^x$  that satisfies  $\varphi(x)$  also satisfies every  $\operatorname{Aut}(\mathcal{U}/A)$ -translate of  $\varphi(x)$ .

Notation: for  $\mathfrak{D} \subseteq \mathcal{U}^x$  and  $H \subseteq G$  we write  $H\mathfrak{D}$  for  $\{h\mathfrak{D} : h \in H\}$ .

In this chapter many proofs require some juggling with negations as epitomized by the following fact.

- 3 Fact The following are equivalent
  - 1.  $\mathcal{D}$  is not *G*-generic
  - 2.  $\neg \mathbb{D}$  is *G*-persistent.

**Proof.** Immediate by spelling out the definitions

- 1. there are no finite  $H \subseteq G$  such that  $\mathfrak{X} \subseteq \cup H\mathfrak{D}$ .
- 2.  $\emptyset \neq \mathfrak{X} \cap (\cap H \neg \mathfrak{D})$  for every finite  $H \subseteq G$ .

Define the following type

$$\gamma_G(x) = \{\vartheta(x) \in \Delta^B(\mathcal{Z}) : \vartheta(x) \text{ is } G\text{-generic}\}$$

**4 Corollary** Let  $p(x) \subseteq \Delta^{\mathsf{B}}(\mathcal{Z})$  be such that  $\gamma_G(x) \cup p(x)$  is finitely satisfiable in  $\mathcal{X}$ . Then p(x) is G-persistent.

**Proof.** Let  $\vartheta(x)$  be a conjunction of formulas in p(x). As  $\gamma_G(x)$  is finitely satisfiable in  $\vartheta(\mathfrak{X})$ , it cannot be that  $\neg \vartheta(x)$  is *G*-generic. From Fact 3, we obtain that  $\vartheta(x)$  is *G*-persistent.

The converse implication holds for complete types.

- **5 Theorem** Let  $p(x) \in S_{\Delta}(\mathcal{Z})$ . Then the following are equivalent
  - 1. p(x) is *G*-invariant and finitely satisfiable in X
  - 2.  $p(\mathbf{x}) \vdash \gamma_G(\mathbf{x})$
  - 3. p(x) is *G*-persistent.

**Proof.**  $1\Rightarrow 2$ . Let  $H\subseteq G$  be finite such that  $\mathfrak{X}\subseteq \cup H\mathfrak{D}$ . By completeness and finite satisfiability,  $p(x)\vdash x\in \cup H\mathfrak{D}$ . Again by completeness,  $p(x)\vdash x\in h\mathfrak{D}$  for some  $h\in H$ . Finally, by invariance,  $p(x)\vdash x\in \mathfrak{D}$ .

2⇒3. Let  $\mathbb D$  be defined by a conjunction of formulas in p(x). If  $\mathbb D$  is not G-persistent then, by Fact 3,  $\neg \mathbb D$  is G-generic. By 2,  $p(x) \vdash x \notin \mathbb D$ , a contradiction.

 $3\Rightarrow 1$ . First note that *G*-persistent types are finitely satisfiable in  $\mathfrak{X}$ . Now, suppose p(x) is not *G*-invariant. Then, by completeness,  $p(x) \vdash \varphi(x;b) \land \neg \varphi(x;gb)$  for some  $g \in G$ . Clearly  $\varphi(x;b) \land \neg \varphi(x;gb)$  is not 2-*G*-persistent as it is inconsistent with its *g*-translate.

- **6 Corollary** The following are equivalent for every  $\Delta^{B}(\mathcal{Z})$ -definable set  $\mathfrak{D}$ 
  - 1.  $\gamma_G(x) \vdash x \in \mathcal{D}$
  - 2.  $p(x) \vdash x \in \mathcal{D}$  for every *G*-persistent  $p(x) \in S_{\Delta}(\mathcal{Z})$ .

**Proof.**  $1\Rightarrow 2$ . This is an immediate consequence of Theorem 5.

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2⇒1. Suppose \gamma_G(x) \not\vdash x \in \mathcal{D}. Then there is a type p(x) \in S_{\Delta}(\mathcal{Z}) consistent with \gamma_G(x) \cup \{x \notin \mathcal{D}\}. By Corollary 4 p(x) is G-persistent. Then ¬2.
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The theorem yields a necessary condition for the existence of *G*-invariant global  $\Delta^{B}(\mathcal{Z})$ -types.

- **7 Corollary** If there exists a *G*-invariant type  $p(x) \in S_{\Delta}(\mathbb{Z})$  finitely satisfiable in  $\mathfrak{X}$  then for every  $\Delta^{\mathsf{B}}(\mathbb{Z})$ -definable set  $\mathfrak{D}$ 
  - 1.  $\mathcal{D}$  and  $\neg \mathcal{D}$  are not both *G*-generic
  - 2. if  $\mathfrak{D}$  is *G*-generic then it is *G*-persistent
  - 3.  $\gamma_G(x)$  is finitely satisfiable in X.

**Proof.** Clearly, 1 and 2 are equivalent by Fact 3 and follow from 3. Finally, 3 is an immediate consequence of 2 of Theorem 5.  $\Box$ 

The following theorem gives a necessary and sufficient condition for the existence of global G-invariant  $\Delta^B(\mathbb{Z})$ -type. Ideally, we would like to have that every G-persistent  $\Delta^B(\mathbb{Z})$ -type extends to a global persistent type. Unfortunately this is not true in general (it requires stronger assumptions, see Section 4). A set  $\mathbb{D}$  is G-wide if every finite cover of  $\mathbb{D}$  by  $\Delta^B(\mathbb{Z})$ -definable sets contains a G-persistent set. In [1] a similar property is called *quasi-non-forking*. Our use of the term *wide* is consistent with [2], though we apply it to a narrow context. A type is G-wide if every conjunction of formulas in the type is G-wide.

- **8 Theorem** Let  $\mathcal{D}$  be a  $\Delta^{\mathbb{B}}(\mathbb{Z})$ -definable set. Then the following are equivalent
  - 1.  $\gamma_G(x)$  is finitely satisfied in  $\mathfrak{D} \cap \mathfrak{X}$
  - 2. there exists a *G*-persistent type  $p(x) \in S_{\Delta}(\mathcal{Z})$  containing  $x \in \mathcal{D}$
  - 3.  $\mathfrak{D}$  is *G*-wide.

**Proof.** 1 $\Rightarrow$ 2. By Corollary 4, it suffices to pick any  $p(x) \in S_{\Delta}(\mathcal{Z})$  containing  $\gamma_G(x)$  and finitely satisfied in  $\mathfrak{D} \cap \mathfrak{X}$ .

 $2 \Rightarrow 1$ . By Theorem 5.

2⇒3. Let  $\mathcal{C}_1, \ldots, \mathcal{C}_n$  be  $\Delta^{\mathbb{B}}(\mathcal{Z})$ -definable sets that cover  $\mathfrak{D}$ . Pick p(x) as in 2. By completeness,  $p(x) \vdash x \in \mathcal{C}_i$  for some i. Then, by Theorem 5,  $\neg \mathcal{C}_i$  is not G-generic. Therefore, by Fact 3,  $\mathcal{C}_i$  is G-persistent.

 $3\Rightarrow 2$ . Let p(x) be maximal among the  $\Delta^B(\mathcal{Z})$ -types that are finitely satisfiable in  $\mathcal{X}\cap \mathcal{D}$  and are such that  $\vartheta(\mathcal{U}^x)$  is G-wide for every  $\vartheta(x)$  that is conjunction of formulas in p(x). We claim that p(x) is a complete  $\Delta^B(\mathcal{Z})$ -type. Suppose for a contradiction that  $\vartheta(x)$ ,  $\neg \vartheta(x) \notin p$ . By maximality there is some formula  $\psi(x)$ , a conjunction of formulas in p(x), and some  $\Delta^B(\mathcal{Z})$ -definable sets  $\mathcal{C}_1, \ldots, \mathcal{C}_n$  that cover both  $\psi(\mathcal{U}^x) \cap \vartheta(\mathcal{U}^x)$  and  $\psi(\mathcal{U}^x) \setminus \vartheta(\mathcal{U}^x)$  and such that no  $\mathcal{C}_i$  is G-persistent. As  $\mathcal{C}_1, \ldots, \mathcal{C}_n$  cover  $\psi(\mathcal{U}^x)$  this is a contradiction. It is only left to show that p(x) is G-invariant. This follows from completeness and Theorem 5.

**9 Corollary** Let  $\mathcal{D}$  be a G-wide  $\Delta^{\mathsf{B}}(\mathcal{Z})$ -definable set. Then  $\mathcal{D} \cap g \cdot \mathcal{D}$  is G-wide for every  $g \in G$ .

**Proof.** Let  $p(x) \in S_{\Delta}(\mathcal{Z})$  be a *G*-persistent type such that  $p(x) \vdash x \in \mathcal{D}$ . By *G*-invariance  $p(x) \vdash x \in g \cdot \mathcal{D}$ .

- **10 Fact** Let  $\mathcal{D}$  be a  $\Delta^{\mathsf{B}}(\mathcal{Z})$ -definable set. The following are equivalent
  - 1.  $\mathcal{D}$  is G-wide;
  - 2. every finite cover of  $\mathfrak{D}$  by  $\Delta^{\pm}(\mathfrak{Z})$ -definable sets contains a G-persistent set.

Let  $q(x) \subseteq \Delta^{\mathsf{B}}(\mathcal{Z})$  be *G*-invariant. We say that q(x) is *G*-prime if for every  $\Delta^{\mathsf{B}}(\mathcal{Z})$ -definable set  $\mathcal{D}$  and every  $g \in G$  if  $q(x) \vdash x \in (\mathcal{D} \cup g \cdot \mathcal{D})$  then  $q(x) \vdash x \in \mathcal{D}$ .

**11 Proposition** The type  $\gamma_G(x)$  is *G*-prime.

**Proof.** Suppose  $\gamma_G(x) \vdash x \in (\mathcal{D} \cup g \cdot \mathcal{D})$ . Assume that  $\gamma_G(x)$  is consistent, otherwise the claim is trivial. Let  $p(x) \in S_{\Delta}(\mathcal{Z})$  be G-persistent. We claim that  $p(x) \vdash x \in \mathcal{D}$ . The proposition follows from the claim by Corollary 6. By completeness  $p(x) \vdash x \in \mathcal{D}$  or  $p(x) \vdash x \in g \cdot \mathcal{D}$ . If the latter occurs,  $p(x) \vdash x \in \mathcal{D}$  follows from invariance.  $\square$ 

## 2. Strong genericity

Unfortunately, *G*-genericy is not preserved under intersection. To obtain closure under intersection, we need to push the concept to a higher level of complexity.

A set  $\mathcal{D} \subseteq \mathcal{U}^x$  is strongly *G*-generic if for every finite  $H \subseteq G$  the set  $\cap H \mathcal{D}$  is generic (recall that  $H \mathcal{D}$  stands for  $\{h \cdot \mathcal{D} : h \in H\}$ ). Dually, we say that  $\mathcal{D}$  is

weakly *G*-persistent if for some finite  $H \subseteq G$  the set  $\cup H \mathcal{D}$  is persistent. Again, the same properties may be attributed to formulas and types.

**12 Lemma** The intersection of two strongly *G*-generic sets is strongly *G*-generic.

**Proof.** We may assume that all sets mentioned below are subsets of  $\mathfrak{X}$ . Let  $\mathfrak{D}$  and  $\mathfrak{C}$  be strongly G-generic and let  $K \subseteq G$  be an arbitrary finite set. It suffices to prove that  $\mathfrak{B} = \cap K(\mathfrak{C} \cap \mathfrak{D})$  is G-generic. Clearly  $\mathfrak{B} = \mathfrak{C}' \cap \mathfrak{D}'$ , where  $\mathfrak{C}' = \cap K\mathfrak{C}$  and  $\mathfrak{D}' = \cap K\mathfrak{D}$ . Note that  $\mathfrak{C}'$  and  $\mathfrak{D}'$  are both strongly G-generic. In particular  $\mathfrak{X} = \bigcup H\mathfrak{D}'$  for some finite  $H \subseteq G$ . Now, from

As  $\mathfrak{C}'$  is strongly G-generic,  $\cap H \mathfrak{C}'$  is G-generic. Therefore  $\cup H \mathfrak{B}$  is also G-generic. The G-genericity of  $\mathfrak{B}$  follows.

Define the following type

$${}^{\mathbf{s}}\gamma_G(x) = \{\vartheta(x) \in \Delta^{\mathbf{B}}(\mathbb{Z}) : \vartheta(x) \text{ is strong } G\text{-generic}\}$$

**13 Corollary** The type  ${}^{s}\gamma_{G}(x)$  is finitely satisfiable in X, strongly G-generic, and G-invariant.

**Proof.** The strong G-genericity is an immediate consequence of Lemma 12. The finite satisfiability is a consequence of G-genericity. As for invariance, note that any translate of a strongly G-generic formula is also strongly G-generic.

**14 Corollary** Let  $p(x) \subseteq L(\mathcal{U})$  be such that  ${}^s\gamma(x) \cup p(x)$  is finitely satisfiable in  $\mathcal{X}$ . Then p(x) is weakly G-persistent.

**Proof.** Similar to Corollary 4. Let  $\vartheta(x)$  be a conjunction of formulas in p(x). As  ${}^{s}\gamma(\mathfrak{X})$  is finitely satisfiable in  $\vartheta(\mathfrak{X})$ , it cannot be that  $\neg\vartheta(x)$  is strongly *G*-generic. From Fact 3, we obtain that  $\neg\vartheta(x)$  not being strongly *G*-generic is equivalent to  $\vartheta(x)$  being weakly *G*-persistent.

#### 3. The diameter of a Lascar type

As an application we prove an interesting theorem of Newelski's on Lascar types. Let  $\mathcal{L}(a/A)$ , the set of tuples with the same Lascar strong type as a over A. This set is the union of a chain of type-definable sets of the form  $\{x: d_A(a,x) \leq n\}$ , where  $d_A$  is the distance in the Lascar graph. In this section we prove that  $\mathcal{L}(a/A)$  is type-definable (if and) only this chain is finite. In other words, only if the connected component of a in the Lascar graph has finite diameter.

It is convenient to address the problem in more general terms. We work under Assumption 1 and also assume that G acts transitively on  $\mathfrak{X}$  i.e.  $Ga = \mathfrak{X}$  for every  $a \in \mathfrak{X}$ . Let  $K \subseteq G$  be a set of generators that is

- 1. symmetric i.e. it contains the unit and is closed under inverse
- 2. conjugacy invariant i.e.  $g K g^{-1} = K$  for every  $g \in G$

We define a discrete metric on  $\mathfrak{X}$ . For  $a,b \in \mathfrak{X}$  let d(a,b) be the minimal n such that  $a \in K^n b$ . This defines a metric which is G-invariant by 2. The diameter of a set  $\mathfrak{C} \subseteq \mathfrak{X}$  is the supremum of d(a,b) for  $a,b \in \mathfrak{C}$ .

We are interested in sufficient conditions for  $\mathcal{X}$  to have finite diameter. The notions introduced in Section 2 offer some hint.

**15 Proposition** If X has a weakly persistent subset of finite diameter, then X itself has finite diameter.

**Proof.** Let  $\mathcal{C} \subseteq \mathcal{X}$  be a weakly persistent set of diameter n. Let  $H \subseteq G$  be finite such that  $\cup H$   $\mathcal{C}$  is persistent. We claim that also  $\cup H$   $\mathcal{C}$  has finite diameter. Let  $a \in \mathcal{C}$  be arbitrary. Let m be larger than d(ha,ka) for all  $h,k,\in H$ . Now, let hb and kc, for some  $h,k,\in H$  and  $b,c\in \mathcal{C}$ , be two arbitrary elements of  $\cup H$   $\mathcal{C}$ . As  $h\mathcal{C}$  and  $k\mathcal{C}$  have the same diameter of  $\mathcal{C}$ ,

$$d(hb, kc) \leq d(hb, ha) + d(ha, ka) + d(ka, kc)$$
  
$$< n + m + n.$$

This proves that  $\cup$  H  $\mathfrak{C}$  has finite diameter. Therefore, without loss of generality, we may assume that  $\mathfrak{C}$  itself is persistent.

By the transitivity of the action, any two elements of  $\mathfrak{X}$  are of the form ha, ka for some  $h, k \in G$  and some  $a \in \mathfrak{C}$ . By persistency, there are  $c \in \mathfrak{C} \cap h\mathfrak{C}$  and  $d \in \mathfrak{C} \cap k\mathfrak{C}$ . Then

$$d(ha, ka) \leq d(ha, c) + d(c, d) + d(d, ka)$$
  
$$\leq n + n + n.$$

Therefore the diameter of  $\mathcal{X}$  does not exceed 3n.

**16 Theorem** Suppose that  $\mathfrak{X}$  and the sets  $\mathfrak{X}_n = K^n a$ , for some  $a \in \mathfrak{X}$ , are type-definable. Then  $\mathfrak{X}$  has finite diameter.

**Proof.** By Proposition 15, it suffices to prove that  $\mathcal{X}_n$  is weakly persistent. By Corollary 14 it suffices to show that for some n the type  ${}^{\mathrm{s}}\gamma_G(x)$  is finitely satisfied in  $\mathcal{X}_n$ . Suppose not. Let  $\psi_n(x) \in {}^{\mathrm{s}}\gamma_G$  be a formula that is not satisfied in  $\mathcal{X}_n$ . Then the type  $p(x) = \{\psi_n(x) : n \in \omega\}$  is finitely satisfied in  $\mathcal{X}$ . From the type-definablity of  $\mathcal{X}$  it follows that p(x) has a realization in  $\mathcal{X}$ . As this realization belongs to some  $\mathcal{X}_n$  we contradict the definition of  $\psi_n(x)$ .

**17 Example** Let  $K \subseteq \operatorname{Aut}(\mathcal{U}/A)$  be the set of automorphisms that fix a model containing A. Then the group G generated by K is  $\operatorname{Autf}(\mathcal{U}/A)$  and  $G \cdot a = \mathcal{X}$  is  $\mathcal{L}(a/A)$ . Let  $\Delta = L_{xz}$  and  $\mathcal{Z} = \mathcal{U}^z$ . Then d(a,b) coincides with the distance in the Lascar graph. The sets  $K^n \cdot a = \{x : d(x,a) \le n\}$  are type definable. Then from Theorem 16 it follows that  $\mathcal{L}(a/A)$  is type definable (if and) only if it has a finite diameter.

### 4. A tamer landscape

Under suitable assumptions some notion introduced in this chapter coalesce, and we are left with a tamer landscape. In general, we have the following.

#### 18 Theorem Assume

1. *G*-persistent  $\Delta^{\mathbb{B}}(\mathbb{Z})$ -definable sets are *G*-wide.

Then the following hold

- 1. *G*-persistent  $\Delta^{B}(\mathbb{Z})$ -definable sets are *G*-wide
- 2. *G*-generic  $\Delta^{B}(\mathbb{Z})$ -definable sets are closed under intersection
- 3. *G*-generic  $\Delta^{B}(\mathcal{I})$ -definable sets are strongly *G*-generic
- 4. weakly persistent  $\Delta^{B}(\mathbb{Z})$ -definable sets are *G*-persistent.

## **Proof.** Clearly $2 \Leftrightarrow 3 \Leftrightarrow 4$ .

1⇒2. Let  $\mathfrak C$  and  $\mathfrak D$  be G-generic  $\Delta^{\mathbb B}(\mathcal Z)$ -definable sets. Suppose for a contradiction that  $\mathfrak C \cap \mathfrak D$  is not G-generic. Then  $\neg (\mathfrak C \cap \mathfrak D)$  is G-persistent. By 1 and Theorem 8 there is a G-invariant global  $\Delta^{\mathbb B}(\mathcal Z)$ -type p(x) containing  $x \notin \mathfrak C \cap \mathfrak D$ . By completeness either  $p(x) \vdash x \notin \mathfrak C$  or  $p(x) \vdash x \notin \mathfrak D$ . This is a contradiction because by Theorem 5  $p(x) \vdash x \in \mathfrak C$  and  $p(x) \vdash x \in \mathfrak D$ .

19 Question Is 1 in Theorem 18 equivalent to 2-4?

**20 Remark** Assume 2-4 in Theorem 18 hold. Then the types  $\gamma_G(x)$  and  ${}^s\gamma_G(x)$  coincide, and therefore *G*-invariant global types exist. It is also worth mentioning that every positive Boolean combination of *G*-generic sets is *G*-generic.

### 5. The action of normal subgroups

Let  $H \leq G$ . The following is an immediate consequence of normality.

**21 Remark** For every  $\mathfrak{D} \subseteq \mathcal{U}^x$  and every  $g \in G$ 

$$\mathfrak{D}$$
 is  $H$ -foo  $\Leftrightarrow g \cdot \mathfrak{D}$  is  $H$ -foo,

where *foo* can be replaced by *generic, invariant, persistent, wide.* In particular, the type  $\gamma_H(x)$  is *G*-invariant.

Recall that if  $\gamma_H(x)$  is finitely satisfiable in  $\mathfrak{X}$ , then H-generic sets are H-wide, cf. Theorem 8. As it happens, we can slightly strengthen this fact.

**22 Proposition** Assume that  $\gamma_H(x)$  is consistent. Let  $\mathcal{D}$  be a  $\Delta^B(\mathbb{Z})$ -definable set. Then if  $\mathcal{D}$  is G-generic it is also H-wide.

**Proof.** Let  $p(x) \in S_{\Delta}(\mathbb{Z})$  be consistent with  $\gamma_H(x)$ . As  $\mathbb{D}$  is G-generic, by completeness  $p(x) \vdash x \in g \cdot \mathbb{D}$  for some  $g \in G$ . Equivalently,  $g^{-1} \cdot p \vdash x \in \mathbb{D}$ . As p(x) is H-persistent, by Remark 21  $g^{-1} \cdot p(x)$  is also H-persistent. Then the proposition follows from Theorem 8.

### 6. Definable groups

In this section we assume that  $\mathcal{Z}$  and  $\mathcal{X}$  are type-definable over some set of parameters A. Moreover we assume that  $\mathcal{Z}$  is a group that act on  $\mathcal{X}$ . The group operations and the group action are assumed definable over A. We use the symbol  $\cdot$  for both the group multiplication and the group action.

Let  $\Psi \subseteq L_x(\mathcal{U})$  be some small set of formulas. In this section  $\Delta$  contains formulas  $\varphi(x;z)$  of the form  $\psi(z^{-1} \cdot x)$  for  $\psi(x) \in \Psi$ . The sets  $\varphi(\mathfrak{X};\mathcal{Z})$  are  $\mathcal{Z}$ -invariant. We write 1 for the identity of  $\mathcal{Z}$ . If  $\varphi(\mathfrak{X};1) \in \Delta^{\mathsf{B}}(1)$  then  $\varphi(\mathfrak{X};g) = g \cdot \varphi(\mathfrak{X};1)$ .

The following auxiliary structure is useful. Let  $\mathcal{U}^{\Delta} = \langle \mathbf{X}; \mathbf{Z} \rangle$  be a 2-sorted structure whose signature  $L^{\Delta}$  contains relation symbols for every formula  $\varphi(\mathbf{x}; z) \in \Delta$ . As there is little risk of confusion, these relations symbols are also denoted by  $\varphi(\mathbf{x}; z)$ . In  $L^{\Delta}$  we include also symbols for the formulas

$$\zeta_n(z_1,\ldots,z_n,z'_1,\ldots,z'_n) = \exists w \bigwedge_{i=1}^n w \cdot z_i = z'_i$$

Finally, note that as  $\mathcal{X}$  and  $\mathcal{Z}$  are assumed to be type-definable,  $\mathcal{U}^{\Delta}$  is a saturated  $L^{\Delta}$ -structure. In this section  $G = \operatorname{Aut}(\mathcal{U}^{\Delta})$ .

To each  $h \in \mathbb{Z}$  we associate the  $L^{\Delta}$ -automorphism  $\langle a; g \rangle \mapsto \langle h \cdot a; h \cdot g \rangle$ . Therefore  $\mathbb{Z}$  is, up to isomorphism, a subgroup of G. In fact, it is a normal subgroup.

**23 Remark** It is worth noticing that automorphisms of  $\mathcal{U}^{\Delta}$  do not preserve the group operations nor the group action. However the formulas  $\zeta_n(z_1,\ldots,z_n,z'_1,\ldots,z'_n)$  ensure that for every automorphism in  $f \in G$  and finite set  $B \subseteq \mathcal{Z}$ , there is an automorphism in  $g \in \mathcal{Z}$  that coincides with f on g. By saturation this is true also for any small set g. Therefore for every  $g \in G$  there is a  $g \in \mathcal{Z}$  such that  $g \in \mathcal{Z}$  holds for every  $g \in \mathcal{Z}$ . We conclude that for  $g \in \mathcal{Z}$  definable sets the notions of generic and persistent under the actions of  $g \in \mathcal{Z}$  and  $g \in \mathcal{Z}$  coincide.

We also consider the action some other normal subgroup  $H \subseteq G$ .

- **24 Proposition** Let  $\mathfrak D$  be a  $\Delta^{\!B}$ -definable set. Assume that  $\gamma_H(x)$  is consistent. Then 1 $\Rightarrow$ 2 holds, where
  - 1.  $\mathcal{D}$  is  $\mathbb{Z}$ -generic
  - 2.  $g \cdot D$  is *H*-wide for every  $g \in \mathbb{Z}$ .

Under the assumption of stability and with  $H = \operatorname{Autf}(\mathcal{U}^{\Delta})$  a stronger claim obtains – the consistency of  $\gamma_H(\mathbf{x})$  is guaranteed, and also the converse implication holds

**Proof.** Let g be given. If  $\mathfrak{D}$  is  $\mathfrak{Z}$ -generic, then so is  $g \cdot \mathfrak{D}$ . Then  $g \cdot \mathfrak{D}$  is, a fortiori, G-generic. Therefore 2 follows from Proposition 22.

We write  $(g)_H$  for the *H*-orbit of *g*, that is, the set  $\{f(g): f \in H\}$ .

- **25 Proposition** Let  $\vartheta(x; z_1, ..., z_n)$  be a Boolean combination of formulas  $\varphi_i(x; z_i)$  for some  $\varphi_i(x; z) \in \Delta$ . Then for every  $h_i \in (g_i)_H$  the following are equivalent
  - 1.  $\vartheta(\mathbf{x}; g_1, \dots, g_n)$  is *H*-wide
  - 2.  $\vartheta(\mathbf{x}; h_1, \dots, h_n)$  is *H*-wide.

**Proof.** Let  $f_i \in H$  be such that  $h_i \in f_i(g_i)$ . Without loss of generality we can assume that only conjunctions occur in  $\vartheta(x; g_1, \ldots, g_n)$ . Let  $\mathfrak{C}_i = \varphi_i(\mathfrak{X}; 1)$ . Then 1 says that  $\mathfrak{C} = g_1 \cdot \mathfrak{C}_1 \cap \cdots \cap g_n \cdot \mathfrak{C}_n$  is H-wide. Let  $f_i \in H$  be such that  $h_i = f_i(g_i)$ . Then, by Corollary 9 also the intersection of the sets  $f_i[\mathfrak{C}]$  is H-wide. A fortiori the intersection of the sets  $f_i[g_i \cdot \mathfrak{C}_i]$  is H-wide. Then  $1 \Rightarrow 2$  follows from Remark 23. By symmetry, this proves the equivalence.

If  $A \subseteq \mathcal{I}$ , we write  $\langle A \rangle$  for the subgroup generated by A.

**26 Proposition** Let  $\vartheta(x; z_1, ..., z_n)$  be a Boolean combination of formulas  $\varphi_i(x; z_i)$  for some  $\varphi_i(x; z) \in \Delta$ . Let  $g \in \mathbb{Z}$  be arbitrary. Assume that  $\vartheta(x; 1, ..., 1)$  is H-wide. Then  $\vartheta(x; h_1, ..., h_n)$  is H-wide for every

$$h_1,\ldots,h_n \in \left\langle \bigcup_{g\in\mathcal{Z}} (g)_H^{-1} \cdot (g)_H \right\rangle.$$

**Proof.** We proceed by induction on the number of factors of the form  $a^{-1} \cdot b$ , for some  $a, b \in (g)_H$ , that occur in  $h_1, \ldots, h_n$ . Without loss of generality we can assume that only conjunctions occur in  $\vartheta(x; g_1, \ldots, g_n)$ . Let  $\mathfrak{C}_i = \varphi_i(\mathfrak{X}; 1)$ . Assume inductively that  $h_1 \cdot \mathfrak{C}_1 \cap \cdots \cap h_n \cdot \mathfrak{C}_n$  is H-wide. Pick two arbitrary  $a, b \in (g)_H$ . By Remark 21

$$a \cdot \mathcal{C}_1 \cap a \cdot h_1^{-1} \cdot h_2 \cdot \mathcal{C}_2 \cap \dots \cap a \cdot h_1^{-1} \cdot h_n \cdot \mathcal{C}_n$$
 is *H*-wide.

By Proposition 25, we can replace  $a \cdot \mathcal{C}_1$  by  $b \cdot \mathcal{C}_1$ . Then, again by Remark 21

$$h_1 \cdot a^{-1} \cdot b \cdot \mathcal{C}_1 \cap h_2 \cdot \mathcal{C}_2 \cap \dots \cap h_n \cdot \mathcal{C}_n$$
 is *H*-wide.

**27 Question** Can we replace the group  $\langle \bigcup_{g \in \mathbb{Z}} (g)_H^{-1} \cdot (g)_H \rangle$  in Proposition 26 with an *H*-invariant group?

#### 7. Notes and references

- [1] Artem Chernikov and Itay Kaplan, Forking and dividing in NTP<sub>2</sub> theories, J. Symbolic Logic 77 (2012), 1–20.
- [2] Ehud Hrushovski, *Stable group theory and approximate subgroups*, J. Amer. Math. Soc. **25** (2012), no. 1, 189–243.