Towards The Reactive Web

Master Thesis

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Chapter 1

Introduction

The fast evolving web provides ever more data and functionalities, both in terms of volume and also complexity. Together with the ubiquitous access to the web through mobile devices and their push notifications, users get flooded with informations and, in turn, also produce increasingly more information. Governing the web's information flood is getting more difficult and even fairly impossible for human beings with the tools given.

It becomes important to the individual to be able to filter out personally important bits and pieces. Basic filters are often available but they do not allow smart filtering in any way. Also, apart from smart filtering of information, people should have the possibility to aggregate important information in their desired place and in a way it's most useful for themselves. Such an aggregation implies access to services that consume data and produce an output, be it an answer or storing the data. This means the user should get access to data and functionality services in a way that she can combine the possibilities in a suitable way to generated the most valuable output for her. Even if the web service access gets simpler these days, the average user is not able to weild them. The challenge to provide users with ways to handle these already simpler accessible services, called Web API's has received a notable amount of attention over the last few years.

Currently only few possibilities exist in the web to turn the web into reactive entity, which allows to process information as it arises.

Chapter 2

Related Work

2.1 Web Services / Web API's

Data and functionalities in the web were always accessible via web services, whatever this means. reengineering of services in the beginning with standardised REST API's we got to Web API, easy access and understandable

The fast evolving web has brought up a trend towards easy to master interfaces to services, the so called Web APIs. They do not only provide access to mere services but whole applications that allow access over Web APIs. These trending Web APIs benefit from a RESTful architecture which predominantly uses HTTP and thus relies on the most basic and powerful operations and the basis of the Web itself, the HTTP protocol.

2.2 Mashups

Mashups combine information and functionality of more than one web service in a single place. The mashing up of such web services allows data to be enhanced with new informations, processing / refinement of the information, or even ways to interact with them, e.g. through Google Maps. Simple functions from different sources can be combined into more powerful ones, which influence data and services in a way their founders eventually didn't even think of. Web service mashups have been developped ever since services in the web started to exist and were accessible in a more or less convenient way. We introduce Paul Rademacher as an example for how recent the invention of web service mashups are. He's one of the first inventors of such a web service mashup. In the same year after Google Maps came up in 2005, he invented a site[12, 3] that displayed Craigslist houses on a Google Map. With no Google Maps API at that time, he needed time and skills to reverse engineer Google Map's functionalities.

A large number of such "static" mashups were and are still developped. They are static in the way that they aggregate a fixed (and mostly low) number, of either data or functionality resources, to provide an enhanced resource in a specialized domain. Of course Mashups can be mashed up again, to provide even more sophisticated functionality and data. Some latest example Mashups, taken from the ProgrammableWeb[13] directory, are:

- Wifi and Plugs[5]: MapBox, Google Docs and Import.io API's used to display where Wi-Fi and plugs are available in London.

- MapLight[7]: GovTrack.us and OpenSecrets API's used to combine political results with financial contributions to show how capital contributions affect voting.
- Shared Count[15]: Facebook, LinkedIn, Pinterest and Twitter API's used to display informations about how well spread a URL is on social media sites.

In the past few years, research and development for platforms to allow users to flexibly mashup Web APIs got attention. With IFTTT and Zapier, two platforms have evolved out of this process. Users that register on those platforms are provided with a multitude of Web API functions that act as event triggers and such that are used to execute actions. The user is then free to combine these event triggers and actions in the way it suits best, creating helpful Web API mashups on their own.

2.3 Rule Languages

It turns out that Web API mashing up is not able to bring reactivity to the web. They are merely aggregations of services that only provide data or functions but no write possibilities such as web applications provide.

Thus

Several different rule languages have been developped for different purposes over the last years and they vary grately in their purpose. A compilation of research on different emerging rule-based languages and technologies [9] gives an oveview over such efforts. We examined different existing rule languages with respect to a certain use case to identify its applicability for reactivity in the web. The use case is defined such that the rule needs to suffice the ECA paradigm:

- Event: Receipt of an Email
- Condition: Check for a certain sender
- Action: Store it remotely via a Web API

We defined an email event which the rule languages need to be able to process. The JSON representation of the given email event is shown in Listing 2.1.

```
{
  "eventname": "email",
  "body": {
    "sender": "sender@mail.com",
    "subject": "Important subject!",
    "textbody": "Hi User,\n\nThis is a lengthy mail body"
  }
}
```

Listing 2.1: Example E-Mail event expressed in JSON

2.3.1 RDF & XML

An early ECA Rule Language for XML repositories[8] was postulated in 2003 and was picked up by many researches afterwards. It was designed to react on insert and delete events within XML repositories and as an action change XML documents.

```
ON INSERT document( inbound_queue.xml )/mails/mail
IF $delta/sender[.= sender@mail.com ]
DO DELETE document( inbound_queue.xml )/mails/mail;
  LET $api = resource("www.webapi.com") IN
  INSERT ($api, newcontent,
    <content>New mail: {$delta/subject}</content>)
```

Listing 2.2: E-Mail Example rule expressed in RDF

Now apart from implementing a rules engine, we would also need to add an XML document event manager which interpretes and pushes events into the XML file *inbound_queue.xml*. Then again this instance would interprete the ouptuts of the ECA engine, which would theoretically manifest in other XML documents, and produce meaningful actions on remote hosts. This wouldn't be an architecture which has its focus on the solution of our use case and, as a result, add complexity and create an unnecessary overhead.

2.3.2 Notation 3

To make the lengthy RDF definitions smaller and more readable, Notation 3[1] was designed and announced in 2005. Through the implies operator(=>) an "event" can be connected to an "action", both expressed in RDF's subject, predicate, object notation, which makes the expression of ECA rules a complicated and not very intuitive task. A solution to our use case would look as follows:

```
{ ?x :event "email". ?x :sender "sender@mail.com" }
  => { :webapi :newcontent ?x }
```

Listing 2.3: E-Mail Example rule expressed in Notation 3

This language is used to express relations between entities and thus not really suitable for our use case, since we would require another interpreter to infer the actions. But concepts and ideas of the work that was done in these consortias could eventually still find influence into our solution.

2.3.3 XChange/Xcerpt

The rule language XChange[11] was the outcome of the REWERSE ([14], Reasoning on the Web with Rules and Semantics) project, which was funded by the EU and Switzerland. Their work influenced a number of future research. The language was designed to add reactive behaviour to a "static" web which is represented through XML resources. Thus we have action logics to alter such resources through insertions and deletions. Since we aim to utilize web API's for our rule language we need a more generic approach which adds flexibility in term of the API provided. But the thorough research done with the language XChange holds valuable concepts, especially in terms of temporal evet composition. This could be a rule according to our use case:

But XChange is designed to access other resources in an action and thus provides powerful tools:

```
TRANSACTION
  in {
    resource { "http://www.webapi.com"},
    newcontents {{
      insert newcontent { var Mail }
    }}
  }
ON
  xchange:event {{
    xchange:sender { "http://mailserver.com" },
    var Mail -> email {{
      sender { "sender@mail.com" }
    }}
  }}
END
```

Listing 2.4: E-Mail Example rule expressed in XChange

```
TRANSACTION
  [...]
ON
  [...]
FROM
  in {
    resource { "http://www.weather.com"},
    temperatures {{
      var T -> temperature {{
        datetime { "2013-10-20-08:00:00AM" }
      }}
    }}
  }
END
```

Listing 2.5: XChange Rule accesses remote resource

2.3.4 JSON Rules

In 2008 *JSON Rules* [6] was introduced as a language to easily react on specific DOM tree compositions. The usage of JavaScript allowed them to provide simple functions which could be called directly by the actions, thus abstracting functionality from the language. This key concept found influence into our language as it allows different layers of abstractions. Through this it is possible to provide generic functions for expert user as well as very limited functions with only few possibilities for parameterization to be used by unexperienced persons. A drawback of this language is its binding to DOM tree events, where we would want to react on any events happening in the world. Also the temporal composition to complex events is not a subject of their work and needs further attention.

```
{
    "id": 0,
    "conditions": [
        {
            "type": "email",
            "constraints": [
                {
                    "propertyName": "sender",
                    "operator": "EQ",
```

```
                    "restriction": {
                        "type": "String",
                        "value": "sender@mail.com"
                    }
                },
                {
                    "bind": "$S",
                    "propertyName": "subject"
                }
            ]
        }
    ],
    "actions": [
        "webapi('addcontent', $S)"
    ]
}
```

Listing 2.6: E-Mail Example rule in JSON Rules

2.3.5 Kinetics Rule Language (KRL)

A most recent (2011) open-source development is the Kinetic Rules Engine together with the Kinetics Rule Language [16]. It is built for the purpose of adding reactivity to the cloud. The language is based on declarative syntax, enriched with imparative elements. But it is a tedious task to get into a whole new language and their caveats. *authorization?*

```
rule store_mail {
  select when mail newmail
  sender re#sender@mail.com#
  subject re#*# setting(subj)
  http:post("http://www.webapi.com/newcontent")
  with params = {
    "text": subj
  }
}
```

Listing 2.7: E-Mail Example rule in KRL

2.3.6 (Reaction) RuleML

The basis of *RuleML* [2] is datalog, a language in the intersection of SQL and Prolog. In 2012 the *Reaction RuleML* [10] language incorporated several different types of rules into the RuleML syntax, to establish a uniform syntax and interchangability of rules. *Reaction RuleML* is a valuable resource in terms of manifold research that has been done in the domain of rule languages, but the syntax is not user-friendly.

R2ML allows usage for RuleML together with many other dialects. Really!?

```
<Rule style="active">
  <on>
    <Event>
      <Atom>
        <Rel per="value">mail</Rel>
```

```
6           <Var>sender</Var>
7           <Var>subject</Var>
8         </Atom>
9       </Event>
10    </on>
11    <if>
12      <Atom>
13        <op><Rel>equals</Rel></op>
14        <Var>sender</Var>
15        <Ind>sender@mail.com</Ind>
16      </Atom>
17    </if>
18    <do>
19      <Atom>
20        <oid><Ind uri="http://webapi.com"/></oid>
21        <Rel>newcontent</Rel>
22        <Var>subject</Var>
23      </Atom>
24    </do>
25  </Rule>
```

Most of the examined rule languages are designed for the interchangability of rules between different service providers. We do not attempt to jump into this domain but we rather pick up important concepts to manifest web API's as first class citizens of our rule language. This allows the ad-hoc design and implementation of reactive rules between existing web API's without the need for their cooperation in setting up their endpoint in a special way.

Chapter 3

Conceptual Model for Reactive Web Systems

3.1 The Web as Event Producer

3.2 From Web Events to Web Actions

3.2.1 Engine / Rules

In the last section we showed how mashups create additional value for the web by combining several WebAPI's. But it turned out, that such mashups are closed systems, which ofen only allow little degree of parametrization. To get past such limitations and define a conceptual model for reactive web systems, it is necessary to define a

existing rule languages, rule engines,

Existing ECA systems all act on local data. Looking at (Wikipedia...) their definition is actions on local data. This does only add reactivity to these systems and not to the Web per se.

Such systems are merely event sinks which add fairly any value to the Web, except for the individual users and the system itself.

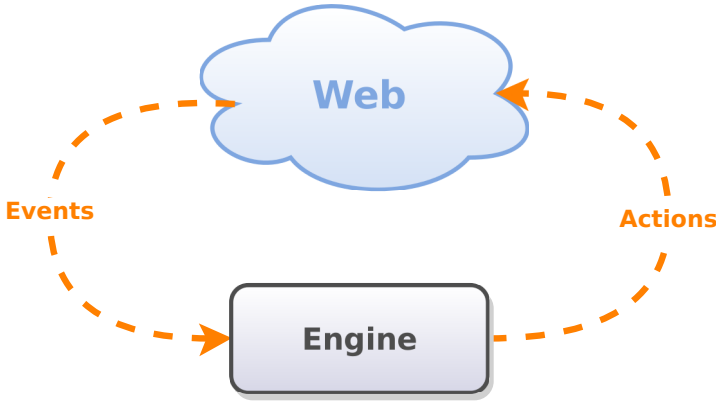

Figure 3.1: Reactor ;) Conceptual model for reactive web systems

- from web to events

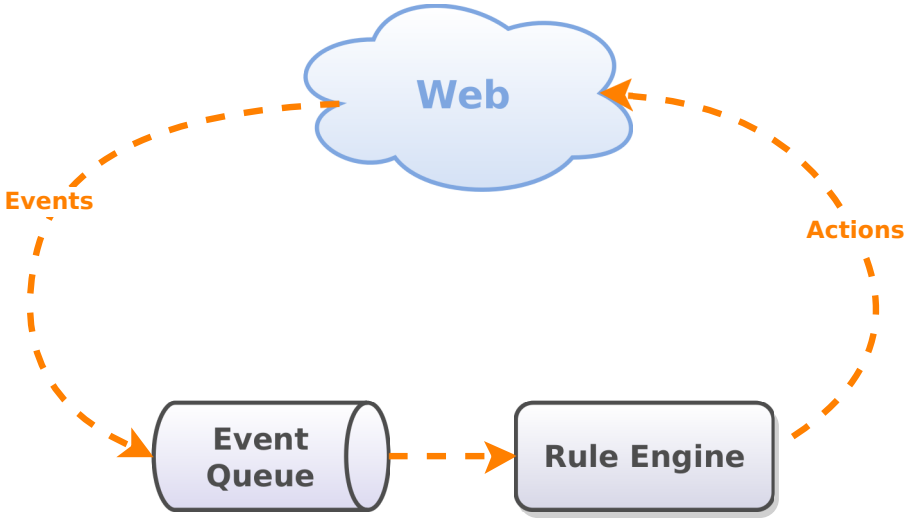

Figure 3.2: Conceptual model for reactive Web Systems

- from events to rules
- from rules to actions
- from actions to the Web
- from concept to engine

Own Rules

```
on mail
if sender="sender@mail.com"
do webapi->newcontent(subject)
```

Would be translated into:

```
{
  "event": "mail",
  "conditions": [
    { "sender": "sender@mail.com" },
  ],
  "actions": [
    {
      "api": "webapi",
      "method": "newcontent",
      "arguments": {
        "text": "$X.subject"
      }
    }
  ]
}
```

```
on weather->tempRaisesAbove(20)
do probinder->addContent(temp)
```

```
on emailyak->newMail
if FromAddress="dominic.bosch.db@gmail.com"
do probinder->newContent(TextBody)
```

```
on probinder->unreadContent
if serviceId=32
do probinder->markread(id),
   probinder->createContent(id, title, tab_name)
```

```
function call(args) {
  require('needle').post(
    'https://probinder.com/service/'
      + args.service + '/' + args.method,
    args.data,
    args.credentials
  );
};

function newContent(txt){
  call({
    service: '27',
    method: 'save',
    data: {
      companyId: '961',
      context: '17930',
      text: txt
    }
  });
}

on mail
do probinder->createContent(subject)

on mail
do probinder->call("27","save",
     ["961", "17930", subject]
   )

on probinder->unread
if serviceId=32
do probinder->setRead(id),
   probinder->makeFileEntry(service, id)
```

```
    "event": "emailyak->newMail",
    "condition": { "FromAddress": "dominic.bosch.db@gmail.com"},
    "actions": [
      {
        "module": "probinder->newContent",
        "arguments": {
          "content": "Received from EmailYak: $X.TextBody"
        }
      }
    ]
```

```
function newMail(callback) {
  needle.get('https://api.emailyak.com/v1/'+key+'/json/get/new/email/',
    function (error, response, body){
        var mails = JSON.parse(body).Emails;
        for(var i = 0; i < mails.length; i++) callback(mails[i]);
    }
  );
}
```

```
{

  "event": "emailyak->newMail",
  "ToAddressList": "test@mscliveweb.simpleyak.com",
  "FromAddress": "dominic.bosch.db@gmail.com",
  "TextBody": "Lengthy body [...]",
  "Subject": "Fwd: test subject",
  [...]

}
```

Chapter 4

The "XYZ" Prototype System

The "XYZ" prototype system is the realisation of a reactive web system. It was developped during the research for this thesis and acts as a platform for feasibility studies of certain use cases.

4.1 Architecture

"XYZ" consists of a queue in which all incoming events are pushed, and an engine that picks the events from the end of the queue whenever it is idle. Since "XYZ"'s core functionalty is the communication with resources in the web, the architecture bases on HTTP protocol in several parts. For example the events are meant to be retrieved completely via HTTP, the user interface is a webpage which posts requests to the system and most actions are also meant to be HTTP requests, or at least using them to gather information.

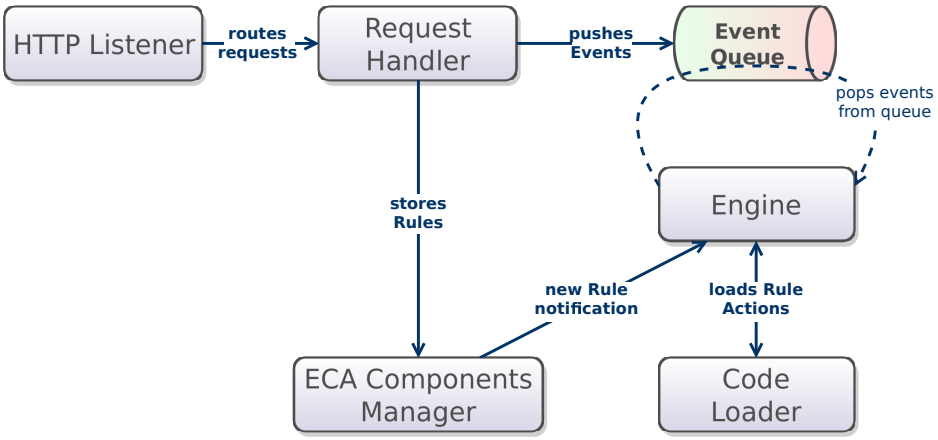

Figure 4.1: "XYZ" Architecture

4.2 Event Triggers

Event Gathering is the E in ECA and without one of these letters such a system would not run. It is of utmost importance to find as much as possible ways to get data into a system.

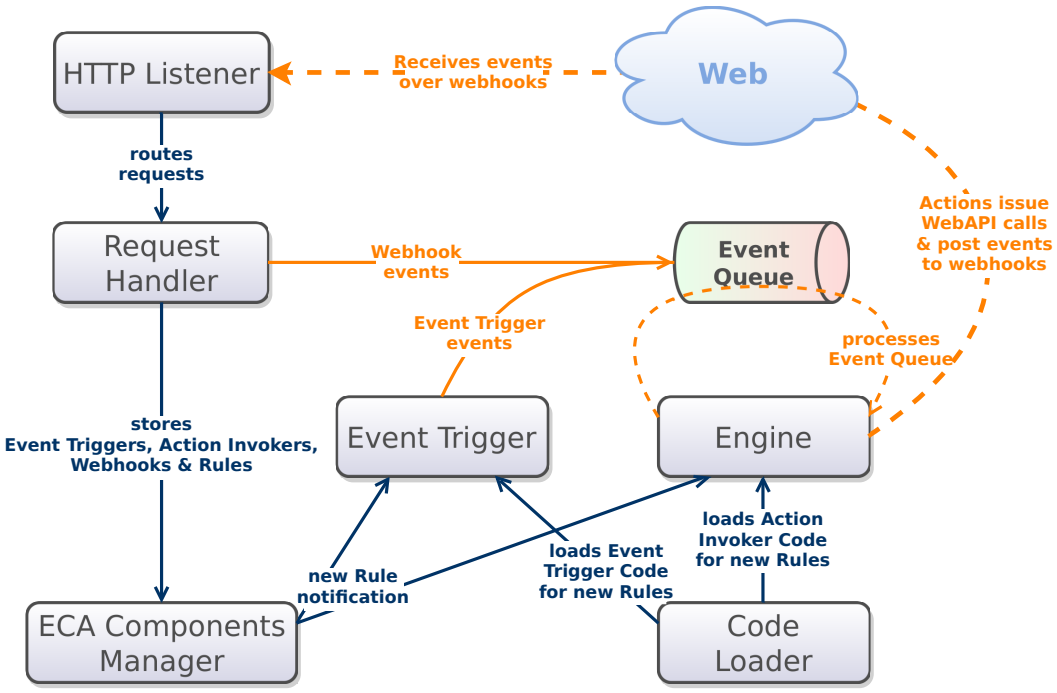

Figure 4.2: "XYZ" Architecture

4.2.1 Polling

4.2.2 Webhooks

4.3 Conditions

4.4 Actions

4.4.1 Aditional Information

4.5 ECA Rules

4.6 Engine

4.7 Asynchronous Closures

Often, optimization approaches and programming language concepts require special attention to avoid common pitfalls. When closures are used as asynchronous functions, developers need to be very careful not to end up with race conditions.

Looking at an example of sequential code execution in Figure 4.3, we see that function execution of `fA` is halted until function `fB` is finished. If `fB` happens to be a latency-driven I/O operation the completion of `fA` could be deferred for a relatively long time. While the application waits for the completion of the I/O operation, some remaining operations in `fA` could eventually already be executed without causing any race conditions.

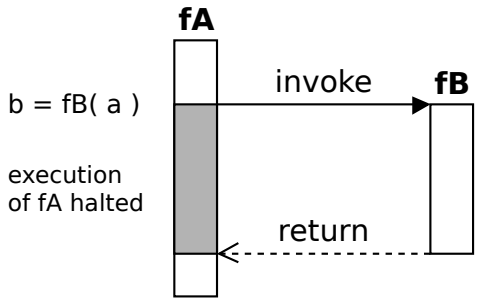

Figure 4.3: Synchronous Function Call

Asynchronous code execution, as shown in Figure 4.4, allows non-blocking and thus scalable applications. Non-blocking operations are a remedy for optimzed resource allocation and open up ways to overcome previously described unnecessary resource bindings. Processing any kind of latency-driven I/O operation asynchronously (e.g. filesystem access and socket communication) exploits resources that would otherwise be bound while waiting for completion. Such operations are processed and completed whenever required resources are available.

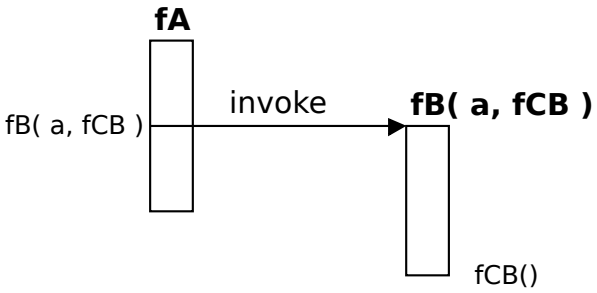

Figure 4.4: Asynchronous Function Call

Often other operations depend on the completion of asynchronous operations, hence their execution needs to be deferred. This necessary code execution deferral is achieved through the use of callback functions, denoted `fCB` in Figure 4.4. Any code placed in a callback function, which is assigned to an asynchronous operation, is only executed after the respective asynchronous operation completed. This allows stacking of functions and operations upon each other which automatically results in a flexible and event-driven application.

Now we take closures into this asynchronous context, as defined in ECMAScript[4], which is the base for widely-spread script languages like JavaScript, JScript and ActionScript. Closures in ECMAScript[4] are defined such as they have access to the context of the function they were created in. This is shown in Figure 4.5 where c from `fA`'s context is accessible from within `fB`, assuming that `fB` was created in `fA` and not only invoked from there. Using asynchronous closures it becomes evident, that the context in the invoking function can change while the closure is still computing and eventually referencing the outer context, thus causing race conditions. This will be most obvious in a loop that immediately invokes `fB` several times, as shown in Figure 4.6. In such a setup c will have different values in the same part of different invocations of `fB`. This might be a very well hidden pitfall since often developers will take care that the context will not change during execution of `fA`. But it is likely that the context will change if `fA` is invoked again while `fB` is still running, which is a common behaviour in event-driven architectures.

Those event-driven context overwrites can be taken care of by shielding the closure from context changes, as shown in Figure 4.7. To shield the closure form context changes, closure

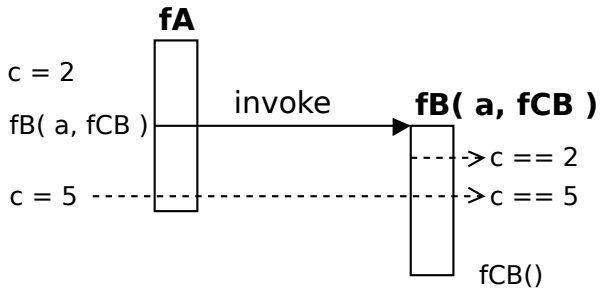

Figure 4.5: Closure Scope and referenced context

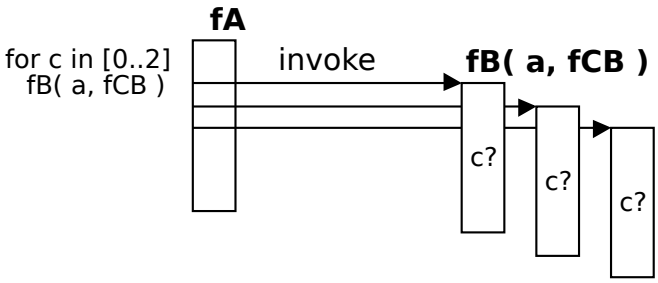

Figure 4.6: Closure context changes in a loop

`fB` needs to create another closure `fC` and return it to `fA`. The argument passed to `fB` is the context (c in Figure 4.7) that might change but requires to be persistent for one invocation. `fC` has now c as a fixed context, which can't be overwritten anymore. Now the only thing left is `fC` needs to be invoked and it will retain the original context. This implementation is necessary when the closure acts as a callback function for asynchronous operations, to preserve the original context in case it is required within the callback function.

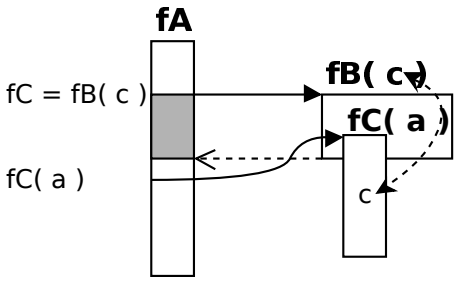

Figure 4.7: Closure context shielding

4.8 Node.js

4.8.1 Benchmarking JavaScript vs. Java

```
1 /*
2  * BenchmarkingDeferred.java
3  */
4 import java.util.concurrent.ScheduledExecutorService;
5 import java.util.concurrent.Executors;
6 import java.util.concurrent.TimeUnit;
```

```
7  import java.util.HashMap;
8
9  public class BenchmarkingDeferred {
10
11   private static Runtime runtime = Runtime.getRuntime();
12   private static final ScheduledExecutorService worker =
13     Executors.newSingleThreadScheduledExecutor();
14
15   private static void deferFunctionCall( int numScopeVars, int delay,
         String scopeId ) {
16     HashMap<String, String> mapVars = new HashMap<String, String>();
17     for( int i = 0; i < numScopeVars; i++ ) {
18       mapVars.put( "id" + i, "12345678" ); // 8 bytes per stored scope
             variable
19     }
20     Object context = new TimeoutContext( "TimeoutFunction" );
21     Runnable task = new RunnableCallbackFunction( mapVars, context );
22     worker.schedule( task, delay, TimeUnit.SECONDS );
23   }
24
25   public static void main( String[] args ) {
26     long startTime, stopTime;
27     int numVars = 10, firstArg = 0;
28     firstArg = Integer.parseInt( args[0] );
29     numVars = Integer.parseInt( args[1] );
30     int j = 0, numFuncs = 1 << firstArg;
31
32     startTime = System.nanoTime();
33     while( j++ < numFuncs) {
34       deferFunctionCall( numVars, numFuncs * 10, numFuncs + "(" + j + ")"
             );
35     }
36     stopTime = System.nanoTime();
37
38     // [...] benchmark system out
39
40     worker.shutdownNow();
41   }
42 }
43
44 /*
45  * RunnableCallbackFunction.java
46  */
47 import java.util.HashMap;
48
49 /*
50  * The Callback function instance.
51  */
52 public class RunnableCallbackFunction implements Runnable {
53
54   // The hashhmap is used to store variables and their value as the scope
55   private HashMap<String, String> mapScope;
56   private Object context;
57
58   public RunnableCallbackFunction( HashMap<String, String> scope, Object
         context ) {
59     this.mapScope = scope;
60     this.context = context;
61   }
62
63   // If this is executing, we didn't wait long enough and the
64   // benchmark time is compromised
```

```
65    public void run() {
66      System.out.println( mapScope.toString() );
67    }
68
69  }
70
71  /*
72   * TimeoutContext.java
73   */
74  public class TimeoutContext {
75    private long idleTimeout = 1;
76    private long idlePrev;
77    private long idleNext;
78    private long idleStart = 140000505;
79    private String onTimeout = null;
80    private boolean repeat = false;
81
82    public TimeoutContext( String cb ) {
83      this.onTimeout = cb;
84    }
85  }
```

Listing 4.1: Closure Benchmarking: Java Code

```
1  /*
2  The function deferral measurements in node.js
3  */
4
5  var deferredFunction = function ( numScopeVars, delay, scopeId ) {
6    var scope = {};
7    for ( var i = 0; i < numScopeVars; i++ ) {
8      scope[ "id" + i ] = "12345678"; // 8 bytes per stored scope variable
9    }
10   setTimeout( function () {
11     // If this is executed we didn't wait long enough
12     console.log( JSON.stringify( scope, null, '  ' ) );
13   }, delay );
14 }
15
16 var numOfFunctions,
17     numOfScopeVars = process.argv[ 3 ];
18
19 numOfFunctions = Math.pow( 2,  process.argv[ 2 ] );
20
21 var time = process.hrtime();
22 for (var i = 0; i < numOfFunctions; i++) {
23   deferredFunction( numOfScopeVars, 1000 * numOfFunctions, numOfFunctions
         + "(" + i + ")" );
24 };
25 var diff = process.hrtime( time );
26
27 var mem = process.memoryUsage();
28 // [...] benchmark system out
29 process.exit( 0 );
```

Listing 4.2: Closure Benchmarking: JavaScript Code

"XYZ" Name discussion

Examples for XYZ:

- **GEEK** (General-Purpose ECA Engine Kernel)
- **DECADE** (Dynamic ECA Demo Engine)
- **WECAST** (WebAPI ECA Service Trigger)
- **RECAST** (Reactive ECA Service Trigger)
- **PECAN** (Productive ECA eNgine)
- **ICECAP** (Inet-Service Calls through ECA Paradigm)

Chapter 5

Discussion & Results

5.1 Future Work

We have seen that the ECA approach is already a powerful one to make the web reactive. A future improvement of this could be to adopt Complex Event Processing (CEP). This would mean that several events could be stored in a rule and be evaluated in terms of time constraints. Through this more complex events can be created as a result of several atomic events which would lead into semantically more complex events. A change in paradigm will result in an approach where events are not just processed when they are entering the system and evaluated against rules, but these events would need to be stored for quite a long time. Also the rules will not all be checked for each event but they are subject to a scheduler. It can be decided when and how often a rule is evaluated and all events will be checked at these point in times, whether they are candidates for firing the rule. A relational database will be needed in order to search through the timestamps

Appendix A

Rules

A.1 Wow

A.1.1 Binder voila

```
'use strict';
var express = require('express');
var qs = require('querystring');
var engine = require('./ecainference');
```

Appendix B

Rules Too

B.1 Wow

B.1.1 Binder voila

Appendix C

Benchmarking

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