# Limen Alpha Programmer's manual

# **Table of Contents**

1	A br	ief description of Limen Alpha	3
2	Mem	nory model	3
3	Regi	ster file	3
4	List	of instructions	. 4
5	Instr	ruction formats	5
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9	Introducing	5 6 8 9 9
6	Instr	ruction processing	
	6.1	Instruction phases	
7	Con	trol registers	. 11
	7.1 7.2 7.3	MSR - Machine Status Register	12
8	Inter	rupt processing	12
9	Core	es synchronization	13
	9.1 9.2 9.3	Mechanism to achieve synchronization	13
10	Prog	ıram example	14
	List	of Symbols and Abbreviations	15

# 1 A brief description of Limen Alpha

Limen Alpha is processor architecture with following characteristics:

- 16-bit width,
- Von Neumann type,
- RISC type,
- dual-core variant,
- can address up to 64 kDB.

# 2 Memory model

It's used physically addressed flat memory model - every address used in instruction is equal address that will appear on the address bus. Memory address space also contains mapped I/O device and is shared for both processor cores.

# 3 Register file

Architecture defines eight general purpose registers. However, register R0 represents always zero value, write to this register is ignored. There is also IP register that holds address of actual executed instruction and some control registers.

Register file 3

# 4 List of instructions

- SLU set if less unsigned
- SL set if less
- ADD addition
- LL load linked
- LD load
- SC store conditional
- ST store
- OR logical disjunction
- ORN logical disjunction not
- AND logical conjunction
- ANDN logical conjunction not
- XOR exclusive logical disjunction
- SLL shift left logical
- SRL shift right logical
- SRA shift right arithmetic

- SUB subtraction
- RTC register to control
- CTR control to register
- LI load immediate
- LIS load immediate scaled
- LIL load immediate low
- LIH load immediate high
- JNE jump if not equal
- JE jump if equal
- JL jump if less
- JLE jump if less or equal
- JG jump if greater
- JGE jump if greater or equal
- JWL jump with link

List of instructions 4

## **5** Instruction formats

Bits with indexes 13 to 15 of every instruction word determine its instruction format. Instructions that are in the same instruction format are often executed very similarly.

## 5.1 Introducing

The following concept [s]imm<x>, where <x> is natural number, represents numbers:

- imm<x> from 0 to  $2^x$  1,
- simm<x> from  $-2^{(x-1)}$  to  $2^{(x-1)}$  1.

imm<x> format and natural numbers of simm<x> format can be expressed in various radix:

- Decimal  $\langle digits(0-9) \rangle$  or  $\langle digits(0-9) \rangle$ d,
- Hexadecimal <digits(0-F or 0-f)>H or <digits(0-F or 0-f)>h,
- Octal  $\langle digits(0-7) \rangle O$  or  $\langle digits(0-7) \rangle O$ ,
- Binary  $\langle \text{digits}(0,1) \rangle$ B or  $\langle \text{digits}(0,1) \rangle$ b.

Negative numbers must be expressed only in one of decimal radix formats.

## 5.2 Arithmetic with immediate value

Instructions of arithmetic type with 5-bit signed immediate value. This value is in two's complement format.

#### Instruction word structure

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	fui	nc		S	imm	5			Ry			Rz	

#### List of instructions

fu	nc	mnemonic	C language
0	0	SLU Rz, Ry, simm5	Rz = Ry < simm5;
0	1	SL Rz, Ry, simm5	Rz = sRy < simm5;
1	0	Х	Χ
1	1	ADD Rz, Ry, simm5	Rz = Ry + simm5;

## 5.3 Transfer with immediate displacement

The only format that allows access to memory and perform synchronization matters. The final memory address is result of sum register value and 5-bit signed immediate value. This value represents displacement from register value.

## Instruction word structure

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	fu	nc		S	imm	5			Ry			Rz	

## List of instructions

```
func
         mnemonic
                                    C language
0 0
         LL Rz, Ry, simm5
                                    Rz = *(Ry + simm5);
                                    link_core = current_core;
                                    link_bit = 1;
   1
         LD Rz, Ry, simm5
                                    Rz = *(Ry + simm5);
  0
         SC Rz, Ry, simm5
                                    if (link_bit && current_core == link_core) {
                                        *(Ry + simm5) = Rz;
                                        Rz = 1;
                                        link_bit = 0;
                                    } else {
                                        Rz = 0;
                                    }
1 1
         ST Rz, Ry, simm5
                                    *(Ry + simm5) = Rz;
```

# 5.4 Logic with immediate value

This instruction format can perform logic operations with 4-bit immediate value, which is first of all extended with zeros to 16-bit width. Instruction format also includes logical and arithmetic shift instructions.

## Instruction word structure

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0		func			imı	m4			Ry			Rz	

## List of instructions

	fu	nc	mnemonic	C language
0	0	0	OR Rz, Ry, imm4	Rz = Ry   imm4;
0	0	1	ORN Rz, Ry, imm4	Rz = Ry   ~imm4;
0	1	0	AND Rz, Ry, imm4	Rz = Ry & imm4;
0	1	1	ANDN Rz, Ry, imm4	Rz = Ry & ~imm4;
1	0	0	XOR Rz, Ry, imm4	Rz = Ry ^ imm4;
1	0	1	SLL Rz, Ry, imm4	Rz = Ry << imm4;
1	1	0	SRL Rz, Ry, imm4	Rz = Ry >> imm4;
1	1	1	SRA Rz, Ry, imm4	Rz = sRy >> imm4;

# 5.5 Arithmetic, logic and control with registers

Instruction format can perform all previous logical and arithmetic instructions. Moreover there is a subtraction. This instruction format uses only registers as operands. There are also instructions for access to control registers.\*

## Instruction word structure

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1		fu	nc			Rx		F	Ry/C	i		Rz	

## List of instructions

		fu	nc	mnemonic	C language
0	0	0	0	OR Rz, Ry, Rx	$Rz = Ry \mid Rx;$
0	0	0	1	ORN Rz, Ry, Rx	$Rz = Ry \mid \sim Rx;$
0	0	1	0	AND Rz, Ry, Rx	Rz = Ry & Rx;
0	0	1	1	ANDN Rz, Ry, Rx	$Rz = Ry \& \sim Rx;$
0	1	0	0	XOR Rz, Ry, Rx	$Rz = Ry ^Rx;$
0	1	0	1	SLL Rz, Ry, Rx	Rz = Ry << (Rx & 0x000F);
0	1	1	0	SRL Rz, Ry, Rx	Rz = Ry >> (Rx & 0x000F);
0	1	1	1	SRA Rz, Ry, Rx	Rz = sRy >> (Rx & 0x000F);
1	0	0	0	SLU Rz, Ry, Rx	Rz = Ry < Rx;
1	0	0	1	SL Rz, Ry, Rx	Rz = sRy < sRx;
1	0	1	0	SUB Rz, Ry, Rx	Rz = Ry - Rx;
1	0	1	1	ADD Rz, Ry, Rx	Rz = Ry + Rx;
1	1	0	0	Х	X
1	1	0	1	Х	X
1	1	1	0	RTC Rx, Ci	Ci = Rx
1	1	1	1	CTR Rz, Ci	Rz = Ci

<sup>\*</sup> When accessing to control registers, unused register index must have zero value.

## 5.6 Load immediate value

Instructions that load 8-bit unsigned immediate to registers. There are two mods: 1) load 8-bit immediate direct/scaled to register, 2) merge 8-bit immediate with lower/higher part of register.

#### Instruction word structure

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0				imı	n8				fui	nc		Rz	

## List of instructions

fu	nc	mnemonic	C language
0	0	LI Rz, imm8	Rz = imm8;
0	1	LIS Rz, imm8	Rz = imm8 << 8;
1	0	LIL Rz, imm8	Rz = (Rz & 0xFF00) + imm8;
1	1	LIH Rz, imm8	Rz = (imm8 << 8) + (Rz & 0x00FF);

## 5.7 Conditional jump with immediate displacement

Provides the only way to realize conditional branching. The final address of success jump is displacement of actual instruction address (IP register value). For this reason, instruction word contains 7-bit signed immediate value. Conditional branch is always result of compare register value and zero.

#### Instruction word structure

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
1	0	1			si	mm	7				Ry			func		

## List of instructions

	fu	nc	mnemonic	C language
0	0	0	Х	Χ
0	0	1	Х	Χ
0	1	0	JNE Ry, simm7	if(Ry != 0) IP += simm7;
0	1	1	JE Ry, simm7	if(Ry == 0) IP += simm7;
1	0	0	JL Ry, simm7	if(sRy < 0) IP $+=$ simm7;
1	0	1	JLE Ry, simm7	if(sRy <= 0) IP += simm7;
1	1	0	JG Ry, simm7	if(sRy $>$ 0) IP += simm7;
1	1	1	JGE Ry, simm7	if(sRy >= 0) IP += simm7;

## 5.8 Unconditional jump with immediate displacement

Includes only one instruction - instruction for realize unconditional jumps and calling functions in program. The first operand is register into which is saved return address (incremented IP register value). Using register R0 if don't needed to save return address. The second operand is 10-bit signed immediate in which range can be jump performed.

#### Instruction word structure

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0					simı	m10						Rz	

## List of instructions

mnemonic	C language
JWL Rz, simm10	Rz = IP + 1; $IP += simm10$ ;

## 5.9 Unconditional jump with registers

Another format for implement unconditional jumps and calling functions. However, opposed to previous format, the jump address is register value that is used as a pointer, save return address is also possible.

#### Instruction word structure

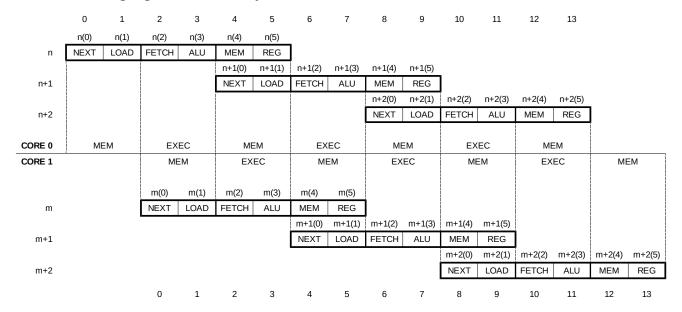
15	5 14	1	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1		1				Χ					Ry			Rz	

#### List of instructions

```
\begin{tabular}{lll} \begin{
```

# 6 Instruction processing

Each instruction takes six clock signals and is overlapped in two phases with a following instruction. So every four clock signals is completed one instruction on the processor core. A pair of processor cores are running together without any stalls.



## 6.1 Instruction phases

- NEXT write new instruction address to IP register
- LOAD read instruction from memory
- FETCH fetch instruction from memory bus
- ALU perform arithmetic or logic operation
- MEM access to memory / control registers
- REG write result to register

# 7 Control registers

Control registers are auxiliary core registers that provides way to use some extra functionality of core.

index	mnemonic	access	internal structure
0	MSR	RW	ID
1	SIP	RW	addr16
2	PRNG	R	rand16
3 - 7	Χ		

Control registers 11

## 7.1 MSR - Machine Status Register

ID - Interrupt Disable

#### Reset value

ID - 1, others - 0

## 7.2 SIP - Saved Instruction Pointer

When interrupt occurs, the return address is saved here.

## 7.3 PRNG - Pseudorandom number generator

Every clock is generated a new pseudorandom number, which can be read from this register. All program writes to this register are ignored.

#### Reset value

Core 0 - 0xFFFF, Core 1 - 0xDDD4

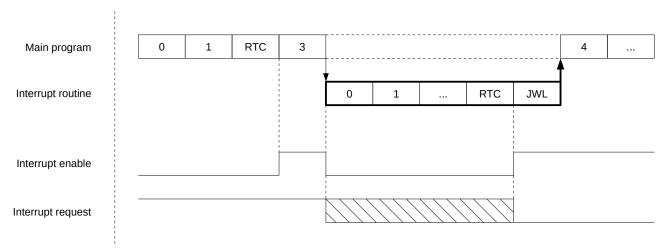
## Algorithm

For generate pseudorandom numbers is used linear-feedback shift register Galois type.

Polynomial used: 
$$x^{16} + x^{15} + x^{14} + x^{12} + x^{11} + x^9 + x^8 + x^7 + x^5 + x^2 + 1$$

# 8 Interrupt processing

One instruction that follows RTC instruction can not be interrupted and always will be executed. After accept interrupt, further interrupts are automatically disabled.



# 9 Cores synchronization

Synchronization mechanism is used when core requires exclusive access to shared data. Before program gets exclusive access to this data, must use some form of synchronization primitives. The very access to this data is done in "critical section".

## 9.1 Mechanism to achieve synchronization

To realize synchronization of access to shared data, there are two instructions LL and SC. Indeed, this instructions provide a way to implement higher level constructions for synchronization.

#### 9.2 Behavior of LL and SC instructions

In the processor there are two bit registers called Link bit and Link core. When LL instruction is executed, Link bit is set to one and Link core is set to core's index which executed this instruction. When SC instruction is executed, it will be indicate whether it was successful or not with write value to register (1 - success, 0 - fail). SC will be successful only if Link bit is set to one and Link core is equal to index of core that is executing this instruction. If this condition is met, write to memory is approved and performed and Link bit is set to zero. Otherwise write to memory is dismissed and Link bit doesn't change.

## 9.3 Lock implementation example

This example uses Test and test-and-set method with active waiting.

# 10 Program example

```
; == Call procedure ==
; test call procedure via register value as pointer
;address mnemonic
                              binary
                                                   C like
                                                     void main() {
                                                    do {
          LI R4, 0
                              1000000000000100
                                                   ; R4 = 0;
1
          LIH R4, 4
                              1000000010011100
                                                     R4 = 1024;
2
          SRL R5, R4, 3
                              0101100011100101
                                                     R5 = R4 >> 3;
3
          JWL R6, R5
                                                     <R5>(); //128
                              1110000000101110
4
          LIL R4, 0
                              100000000010100
                                                    R4 &= 0xFF00;
5
          JWL R0, R4
                              1110000000100000
                                                     } while(1);
                                                     void 128() {
128 :
          SUB R4, R4, R5
                              0111010101100100
                                                     R4 -= R5;
129:
          SRL R4, R4, 1
                                                     R4 >>= 1;
                              0101100001100100
          LIH R4, 0
                              100000000011100
130 :
                                                     R4 &= 0 \times 00 FF;
131:
          JWL R0, R4
                              1110000000100000
                                                     //"JMP R4" (192)
192 :
          JWL R0, R6
                              1110000000110000
                                                     return;
```

The program is only for demonstrate the possibilities. Used algorithm is not intended to be optimized.

Program example 14

# **List of Symbols and Abbreviations**

DB Double byte

IP Instruction pointer

Rz, Ry, Rx General purpose register indexes

LC Link core

AC Actual core

LB Link bit

Ci Control register index

DEF Define