# Array Access Analysis by Ilvmpy

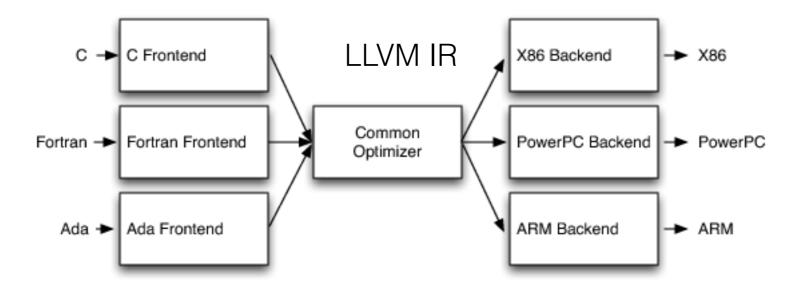
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### LLVM infrastructure

Three Phase Design of LLVM infrastructure





## LLVM Intermediate Representation (IR)

#### Features:

- Three address form
- Static single assignment
- low-level RISC like instruction set with high-level type information

#### Example:

```
int max(int a, int b) {
    if (a > b)
        return a;
    else
        return b;
}
```

```
define i32 @max(i32 %a, i32 %b) #0 {
 %1 = alloca i32, align 4
 %2 = alloca i32, align 4
 %3 = alloca i32, align 4
 store i32 %a, i32* %2, align 4
 store i32 %b, i32* %3, align 4
 %4 = load i32 * %2, align 4
 %5 = load i32* %3, align 4
 %6 = icmp \ sqt \ i32 \ %4, \ %5
 br i1 %6, label %7, label %9
; <label>:7
                                                    ; preds = %0
 %8 = load i32 * %2, align 4
  store i32 %8, i32* %1
 br label %11
; <label>:9
                                                    ; preds = %0
  %10 = load i32* %3, align 4
  store i32 %10, i32* %1
  br label %11
; <label>:11
                                                    ; preds = %9, %7
 %12 = load i32* %1
  ret i32 %12
```

## LLVM API

- Features:
  - C++ implementation
  - templates and pointers are everywhere

```
for (Function::iterator b = F.begin(), be = F.end(); b != be; ++b) {
    for (BasicBlock::iterator i = b->begin(), ie = b->end(); i != ie; ++i) {
        if (CallInst* callInst = dyn_cast<CallInst>(&*i)) {
            if (callInst->getCalledFunction() == targetFunc)
            ++callCounter;
        }
    }
}
```

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```

#### LLVMPY

## Array access analysis of GPU kernel program

- Kernel Program
  - blockldx
  - threadIdx

```
__global__ void add( int *a, int *b, int *c ) {
        int tid = threadIdx.x; // handle the data at this index
        if (tid < N)
        c[tid] = a[tid] + b[tid];
}
```

- Array Access Analysis
  - Extract the index
    - tid -> threadIdx.x
  - Calculate the range
    - thread block |- 0 -|- 1 -|- 2 -|...
    - threads |-0-|-1-|-2-|-3-|-4-|-5-|-6-|-7-|-8-|-9-|10|11|...
    - array |-0-|-1-|-2-|-3-|-4-|-5-|-6-|-7-|-8-|-9-|10|11|...

## Implementation

- Algorithm (Array index analysis)
  - locating array access instruction 'getelementptr'
  - construct prefix expression by define-use chain
  - construct infix expression from prefix expression
- Algorithm (Induction variable analysis)
  - locating store to induction variable instructions
  - construct expression
- Algorithm (Constrain analysis)
  - locating the basic blocks contain array access
  - construct Execution Path Tree from control flow graph
  - extract branch conditions to constrains

## Test

#### MatrixMul

```
Array access analysis result=======
A = a + wA * ty + tx
B = b + wB * ty + tx
C = wB * 32 * by + 32 * bx + wB * ty + tx
induction variables analysis result=====
a = wA * 32 * blockIdx.y
a = a + 32
b = 32 * blockIdx_x
b = b + 32 * wB
constrains analysis result========
full execution path:
Start From Basic Block 0 : [[0]]
Start From Basic Block 1 : [[1, 0], [1, 7, 6, 3, 2, 1], [1, 7, 6, 3, 5, 4, 3]]
Start From Basic Block 2 : [[2, 1, 0], [2, 1, 7, 6, 3, 2], [2, 1, 7, 6, 3, 2, 5, 4, 3]]
Start From Basic Block 3: [[3, 5, 4, 3], [3, 2, 1, 0], [3, 2, 1, 7, 6, 3]]
Start From Basic Block 4: [[4, 3, 5, 4], [4, 3, 2, 1, 0], [4, 3, 2, 1, 7, 6, 3]]
Start From Basic Block 5 : [[5, 4, 3, 5], [5, 4, 3, 2, 1, 0], [5, 4, 3, 2, 1, 7, 6, 3]]
Start From Basic Block 6: [[6, 3, 5, 4, 3], [6, 3, 2, 1, 0], [6, 3, 2, 1, 7, 6]]
Start From Basic Block 7: [[7, 6, 3, 2, 1, 7], [7, 6, 3, 5, 4, 3], [7, 6, 3, 2, 1, 0]]
Start From Basic Block 8: [[8, 1, 0], [8, 1, 7, 6, 3, 2, 1], [8, 1, 7, 6, 3, 5, 4, 3]]
execution path of a,b:[2, 1, 0]
execution path of c:[8, 1, 0]
a,b constrains: a \leq wA * 32 * blockIdx.y + wA - 1
c constrains: a > wA * 32 * blockIdx.y + wA - 1
```

#### • Reference:

• [1] Chris Lattner. The Architecture of Open Source Applications: LLVM.