



## Chapter 6: Database Design Using the E-R Model

Database System Concepts, 7<sup>th</sup> Ed.

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### Outline

- Overview of the Design Process
- The Entity-Relationship Model
- Complex Attributes
- Mapping Cardinalities
- Primary Key
- Removing Redundant Attributes in Entity Sets
- Reducing ER Diagrams to Relational Schemas
- Extended E-R Features
- Entity-Relationship Design Issues
- Alternative Notations for Modeling Data
- Other Aspects of Database Design



## Outline

- Extended E-R Features
- Entity-Relationship Design Issues
- Alternative Notations for Modeling Data
- Other Aspects of Database Design



## Design Phases

- Initial phase -- characterize fully the data needs of the prospective database users.
- Second phase -- choosing a data model
  - Applying the concepts of the chosen data model
  - Translating these requirements into a conceptual schema of the database.
  - A fully developed conceptual schema indicates the functional requirements of the enterprise.
    - Describe the kinds of operations (or transactions) that will be performed on the data.



## Design Phases (Cont.)

- Final Phase -- Moving from an abstract data model to the implementation of the database
  - Logical Design – Deciding on the database schema. Database design requires that we find a “good” collection of relation schemas.
    - Business decision – What attributes should we record in the database?
    - Computer Science decision – What relation schemas should we have and how should the attributes be distributed among the various relation schemas?
  - Physical Design – Deciding on the physical layout of the database



## Design Alternatives

- In designing a database schema, we must ensure that we avoid two major pitfalls:
  - Redundancy: a bad design may result in repeat information.
    - Redundant representation of information may lead to **data inconsistency** among the various copies of information when updating the information.
  - Incompleteness: a bad design may make certain aspects of the enterprise difficult or impossible to model.
- Avoiding bad designs is not enough. There may be a large number of good designs from which we must choose.



## Design Approaches

- Entity Relationship Model (covered in this chapter)
  - Models an enterprise as a collection of *entities* and *relationships*
    - Entity: a “thing” or “object” in the enterprise that is distinguishable from other objects
      - Described by a set of *attributes*
    - Relationship: an association among several entities
  - Represented diagrammatically by an *entity-relationship diagram*:
- Normalization Theory (Chapter 7)
  - Formalize what designs are bad, and test for them



## Outline of the ER Model



## ER model -- Database Modeling

- The ER data model was developed to facilitate database design by allowing specification of an **enterprise schema** that represents the overall logical structure of a database.
- The ER data model employs three basic concepts:
  - entity sets,
  - relationship sets,
  - attributes.
- The ER model also has an associated diagrammatic representation, the **ER diagram**, which can express the overall logical structure of a database graphically.



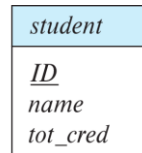
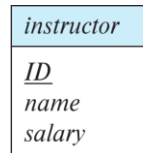
## Entity Sets

- An **entity** is an object that exists and is distinguishable from other objects.
  - Example: specific person, company, event, plant
- An **entity set** is a set of entities of the same type that share the same properties.
  - Example: set of all persons, companies, trees, holidays
- An entity is represented by a set of attributes; i.e., descriptive properties possessed by all members of an entity set.
  - Example:  
*instructor* = (*ID*, *name*, *salary*)  
*course* = (*course\_id*, *title*, *credits*)
- A subset of the attributes form a **primary key** of the entity set; i.e., uniquely identifying each member of the set.



## Representing Entity sets in ER Diagram

- Entity sets can be represented graphically as follows:
  - Rectangles represent entity sets.
  - Attributes listed inside entity rectangle
  - Underline indicates primary key attributes



## Relationship Sets

- A **relationship** is an association among several entities

Example:

44553 (Peltier)	<u>advisor</u>	22222 (Einstein)
student entity	relationship set	instructor entity

- A **relationship set** is a mathematical relation among  $n \geq 2$  entities, each taken from entity sets

$$\{(e_1, e_2, \dots, e_n) \mid e_1 \in E_1, e_2 \in E_2, \dots, e_n \in E_n\}$$

where  $(e_1, e_2, \dots, e_n)$  is a relationship

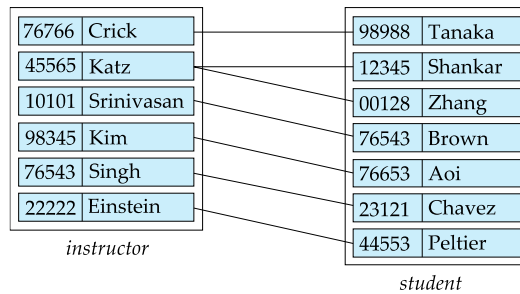
- Example:

$$(44553, 22222) \in \text{advisor}$$



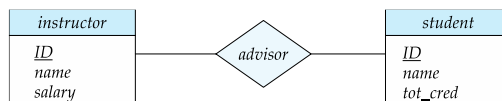
## Relationship Sets (Cont.)

- Example: we define the relationship set *advisor* to denote the associations between students and the instructors who act as their advisors.
- Pictorially, we draw a line between related entities.



## Representing Relationship Sets via ER Diagrams

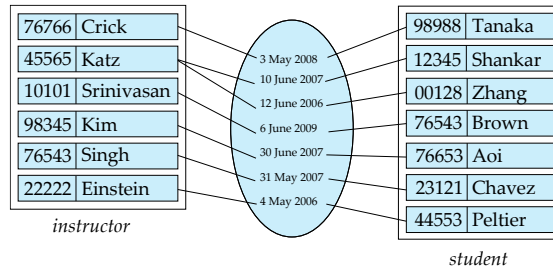
- Diamonds represent relationship sets.



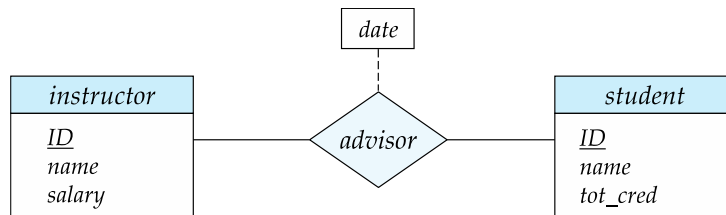


## Relationship Sets (Cont.)

- An attribute can also be associated with a relationship set.
- For instance, the *advisor* relationship set between entity sets *instructor* and *student* may have the attribute *date* which tracks when the student started being associated with the advisor



## Relationship Sets with Attributes

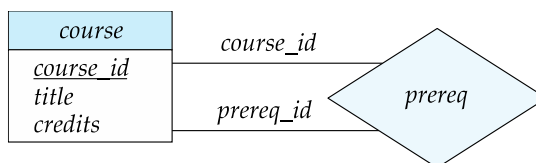






## Roles

- Entity sets of a relationship need not be distinct
  - Each occurrence of an entity set plays a “role” in the relationship
- The labels “*course\_id*” and “*prereq\_id*” are called **roles**.



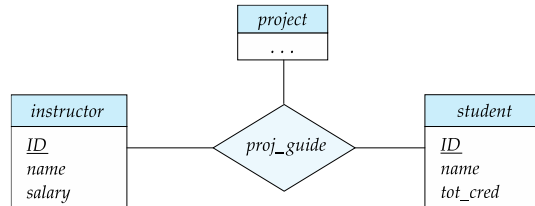
## Degree of a Relationship Set

- Binary relationship
  - involve two entity sets (or degree two).
  - most relationship sets in a database system are binary.
- Relationships between more than two entity sets are rare. Most relationships are binary. (More on this later.)
  - Example: *students* work on research *projects* under the guidance of an *instructor*.
  - relationship *proj\_guide* is a ternary relationship between *instructor*, *student*, and *project*



## Non-binary Relationship Sets

- Most relationship sets are binary
- There are occasions when it is more convenient to represent relationships as non-binary.
- E-R Diagram with a Ternary Relationship



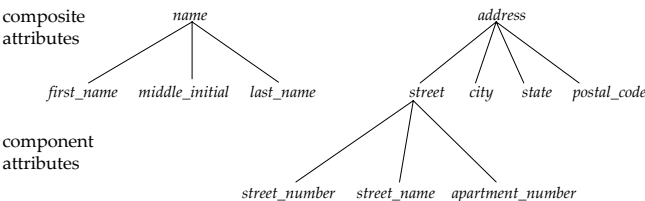
## Complex Attributes

- Attribute types:
  - **Simple** and **composite** attributes.
  - **Single-valued** and **multivalued** attributes
    - Example: multivalued attribute: *phone\_numbers*
  - **Derived** attributes
    - Can be computed from other attributes
    - Example: age, given date\_of\_birth
- **Domain** – the set of permitted values for each attribute

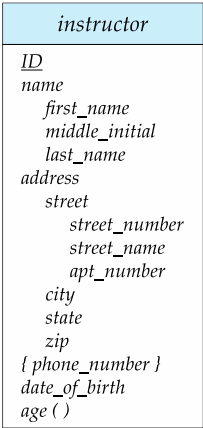


# Composite Attributes

- Composite attributes allow us to divided attributes into subparts (other attributes).



# Representing Complex Attributes in ER Diagram



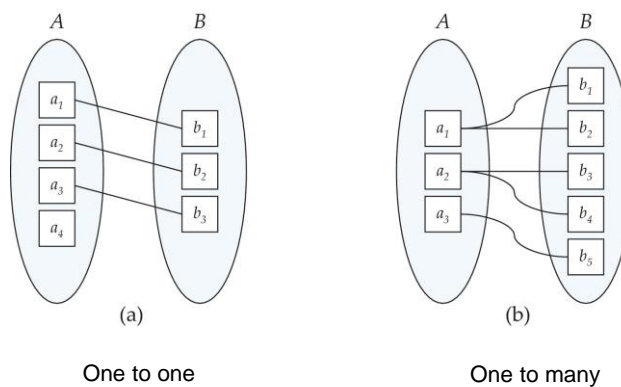


## Mapping Cardinality Constraints

- Express the number of entities to which another entity can be associated via a relationship set.
- Most useful in describing binary relationship sets.
- For a binary relationship set the mapping cardinality must be one of the following types:
  - One to one
  - One to many
  - Many to one
  - Many to many



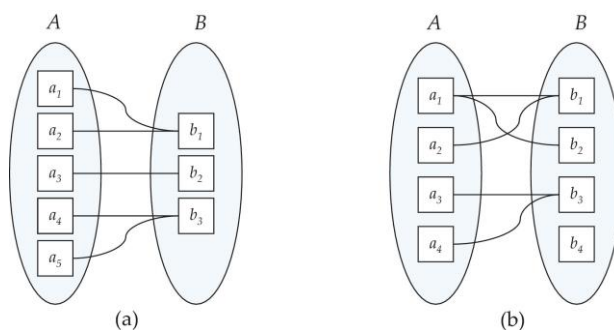
## Mapping Cardinalities



Note: Some elements in  $A$  and  $B$  may not be mapped to any elements in the other set



## Mapping Cardinalities



Many to one

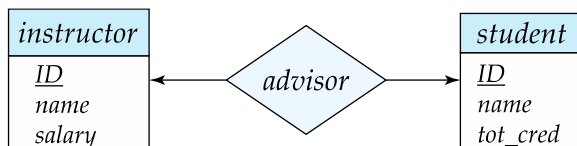
Many to many

Note: Some elements in A and B may not be mapped to any elements in the other set



## Representing Cardinality Constraints in ER Diagram

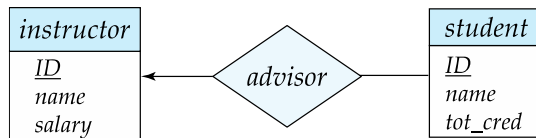
- We express cardinality constraints by drawing either a directed line ( $\rightarrow$ ), signifying "one," or an undirected line ( $\text{—}$ ), signifying "many," between the relationship set and the entity set.
- One-to-one relationship between an *instructor* and a *student* :
  - A student is associated with at most one *instructor* via the relationship *advisor*
  - A *student* is associated with at most one *department* via *stud\_dept*





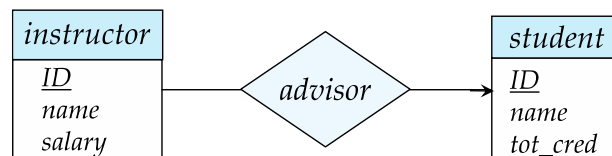
## One-to-Many Relationship

- one-to-many relationship between an *instructor* and a *student*
  - an instructor is associated with several (including 0) students via *advisor*
  - a student is associated with at most one instructor via *advisor*,



## Many-to-One Relationships

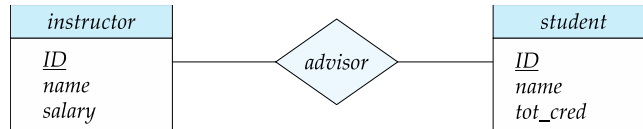
- In a many-to-one relationship between an *instructor* and a *student*,
  - an instructor is associated with at most one student via *advisor*,
  - and a student is associated with several (including 0) instructors via *advisor*





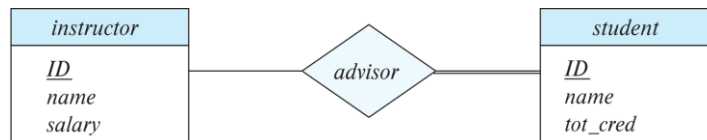
## Many-to-Many Relationship

- An instructor is associated with several (possibly 0) students via *advisor*
- A student is associated with several (possibly 0) instructors via *advisor*



## Total and Partial Participation

- **Total participation** (indicated by double line): every entity in the entity set participates in at least one relationship in the relationship set



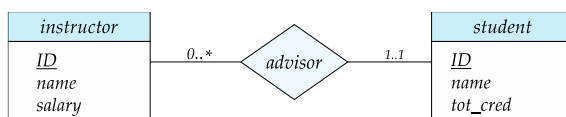
participation of *student* in *advisor* relation is total

- every *student* must have an associated instructor
- **Partial participation**: some entities may not participate in any relationship in the relationship set
  - Example: participation of *instructor* in *advisor* is partial



## Notation for Expressing More Complex Constraints

- A line may have an associated minimum and maximum cardinality, shown in the form  $l..h$ , where  $l$  is the minimum and  $h$  the maximum cardinality
  - A minimum value of 1 indicates total participation.
  - A maximum value of 1 indicates that the entity participates in at most one relationship
  - A maximum value of \* indicates no limit.
- Example

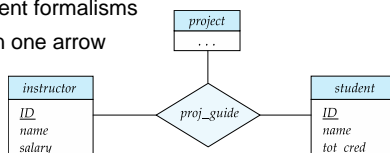


- Instructor can advise 0 or more students. A student must have 1 advisor; cannot have multiple advisors



## Cardinality Constraints on Ternary Relationship

- We allow at most one arrow out of a ternary (or greater degree) relationship to indicate a cardinality constraint
- For example, an arrow from *proj\_guide* to *instructor* indicates each student has at most one guide for a project
- If there is more than one arrow, there are two ways of defining the meaning.
  - For example, a ternary relationship  $R$  between  $A$ ,  $B$  and  $C$  with arrows to  $B$  and  $C$  could mean
    1. Each  $A$  entity is associated with a unique entity from  $B$  and  $C$  or
    2. Each pair of entities from  $(A, B)$  is associated with a unique  $C$  entity, and each pair  $(A, C)$  is associated with a unique  $B$
  - Each alternative has been used in different formalisms
  - To avoid confusion we outlaw more than one arrow







## Primary Key

- Primary keys provide a way to specify how entities and relations are distinguished. We will consider:
  - Entity sets
  - Relationship sets.
  - Weak entity sets



## Primary key for Entity Sets

- By definition, individual entities are distinct.
- From database perspective, the differences among them must be expressed in terms of their attributes.
- The values of the attribute values of an entity must be such that they can uniquely identify the entity.
  - No two entities in an entity set are allowed to have exactly the same value for all attributes.
- A key for an entity is a set of attributes that suffice to distinguish entities from each other



## Primary Key for Relationship Sets

- To distinguish among the various relationships of a relationship set we use the individual primary keys of the entities in the relationship set.
  - Let  $R$  be a relationship set involving entity sets  $E_1, E_2, \dots, E_n$
  - The **set of attributes** of  $R$  consists of the union of the primary keys of entity sets  $E_1, E_2, \dots, E_n$ , which describes an individual relationship in  $R$ .
  - If the relationship set  $R$  has attributes  $a_1, a_2, \dots, a_m$  associated with it, then **the set of attributes** of  $R$  also includes the attributes  $a_1, a_2, \dots, a_m$
- The choice of the primary key for a relationship set depends on the mapping cardinality of the relationship set.
  - **the union of the primary keys of entity sets  $E_1, E_2, \dots, E_n$  forms a superkey for  $R$**
- Example: relationship set “advisor”.
  - The primary key consists of *instructor.ID* and *student.ID*



## Choice of Primary key for Binary Relationship

- Many-to-Many relationships. The preceding union of the primary keys is a minimal superkey and is chosen as the primary key.
- One-to-Many relationships. The primary key of the “Many” side is a minimal superkey and is used as the primary key.
- Many-to-one relationships. The primary key of the “Many” side is a minimal superkey and is used as the primary key.
- One-to-one relationships. The primary key of either one of the participating entity sets forms a minimal superkey, and either one can be chosen as the primary key.



## Weak Entity Sets

- Consider a *section* entity, which is uniquely identified by a *course\_id*, *semester*, *year*, and *sec\_id*.
- Clearly, section entities are related to course entities. Suppose we create a relationship set *sec\_course* between entity sets *section* and *course*.
- Note that the information in *sec\_course* is redundant, since *section* already has an attribute *course\_id*, which identifies the course with which the section is related.
- One option to deal with this redundancy is to get rid of the relationship *sec\_course*; however, by doing so the relationship between *section* and *course* becomes implicit in an attribute, which is not desirable.

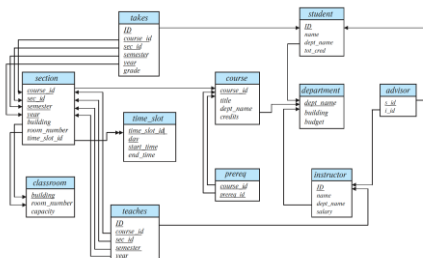


Figure 2.9 Schema diagram for the university database.



## Weak Entity Sets (Cont.)

- An alternative way to deal with this redundancy is to not store the attribute *course\_id* in the *section* entity and to only store the remaining attributes *section\_id*, *year*, and *semester*.
  - However, the entity set *section* then does not have enough attributes to identify a particular *section* entity uniquely
- To deal with this problem, we treat the relationship *sec\_course* as a special relationship that provides extra information, in this case, the *course\_id*, required to identify *section* entities uniquely.
- A **weak entity set** is one whose existence is dependent on another entity, called its **identifying entity**
- Instead of associating a primary key with a weak entity, we use the identifying entity, along with extra attributes called **discriminator** to uniquely identify a weak entity.



## Weak Entity Sets (Cont.)

- An entity set that is not a weak entity set is termed a **strong entity set**.
- Every weak entity must be associated with an identifying entity; that is, the weak entity set is said to be **existence dependent** on the identifying entity set.
- The identifying entity set is said to **own** the weak entity set that it identifies.
- The relationship associating the weak entity set with the identifying entity set is called the **identifying relationship**.
- Note that the relational schema we eventually create from the entity set *section* does have the attribute *course\_id*, for reasons that will become clear later, even though we have dropped the attribute *course\_id* from the entity set *section*.



## Expressing Weak Entity Sets

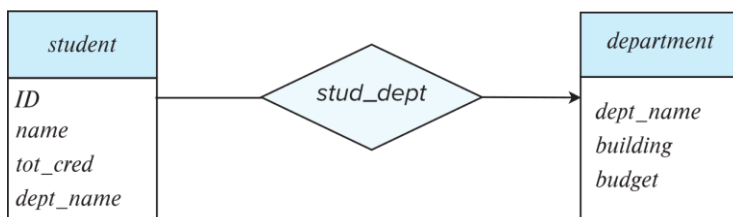
- In E-R diagrams, a weak entity set is depicted via a double rectangle.
- We underline the discriminator of a weak entity set with a dashed line.
- The relationship set connecting the weak entity set to the identifying strong entity set is depicted by a double diamond.
- Primary key for *section* – (*course\_id*, *sec\_id*, *semester*, *year*)



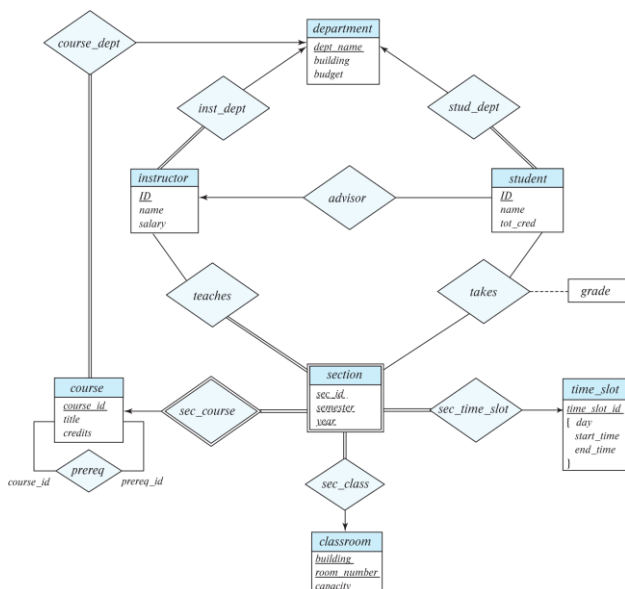


## Redundant Attributes

- Suppose we have entity sets:
  - student*, with attributes: *ID*, *name*, *tot\_cred*, *dept\_name*
  - department*, with attributes: *dept\_name*, *building*, *budget*
- We model the fact that each student has an associated department using a relationship set *stud\_dept*
- The attribute *dept\_name* in *student* replicates information present in the relationship and is therefore redundant
  - and needs to be removed.
- BUT: when converting back to tables, in some cases the attribute gets reintroduced, as we will see later.



## E-R Diagram for a University Enterprise





## Reduction to Relation Schemas



## Reduction to Relation Schemas

- Entity sets and relationship sets can be expressed uniformly as *relation schemas* that represent the contents of the database.
- A database which conforms to an E-R diagram can be represented by a collection of schemas.
- For each entity set and relationship set there is a unique schema that is assigned the name of the corresponding entity set or relationship set.
- Each schema has a number of columns (generally corresponding to attributes), which have unique names.



## Representing Entity Sets

- A strong entity set reduces to a schema with the same attributes

*student*(ID, name, tot\_cred)

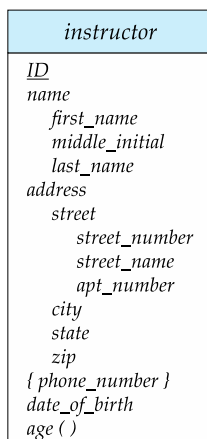
- A weak entity set becomes a table that includes a column for the primary key of the identifying strong entity set

*section* ( course\_id, sec\_id, sem, year )

- Example



## Representation of Entity Sets with Composite Attributes



- Composite attributes are flattened out by creating a separate attribute for each component attribute
  - Example: given entity set *instructor* with composite attribute *name* with component attributes *first\_name* and *last\_name* the schema corresponding to the entity set has two attributes *name\_first\_name* and *name\_last\_name*
    - Prefix omitted if there is no ambiguity (*name\_first\_name* could be *first\_name*)
- Ignoring multivalued attributes, extended instructor schema is
  - *instructor*(ID, *first\_name*, *middle\_initial*, *last\_name*, *street\_number*, *street\_name*, *apt\_number*, *city*, *state*, *zip\_code*, *date\_of\_birth*)



## Representation of Entity Sets with Multivalued Attributes

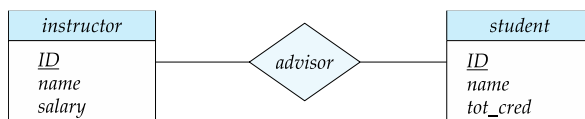
- A multivalued attribute  $M$  of an entity  $E$  is represented by a separate schema  $EM$
- Schema  $EM$  has attributes corresponding to the primary key of  $E$  and an attribute corresponding to multivalued attribute  $M$
- Example: Multivalued attribute *phone\_number* of *instructor* is represented by a schema:  
 $inst\_phone = (\underline{ID}, phone\_number)$
- Each value of the multivalued attribute maps to a separate tuple of the relation on schema  $EM$ 
  - For example, an *instructor* entity with primary key 22222 and phone numbers 456-7890 and 123-4567 maps to two tuples:  
 $(22222, 456-7890)$  and  $(22222, 123-4567)$
- What about “*time\_slot*” entity set having multivalued composite attributes in E-R diagram for a university enterprise ?
  - Note: A missing arrow head of the **sec\_time\_slot** relationship of ER diagram in Fig. 6.15 of the textbook (pp. 264)
  - Read the last paragraph at the page 266 and all the paragraphs before section 6.7.3 at the page 267 of the textbook.



## Representing Relationship Sets

- A many-to-many relationship set is represented as a schema with attributes for the primary keys of the two participating entity sets, and any descriptive attributes of the relationship set.
- Example: schema for relationship set *advisor*

$advisor = (\underline{s\_id}, i\_id)$

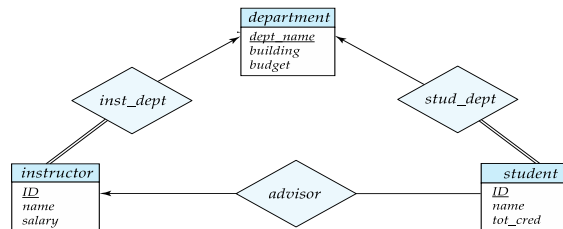






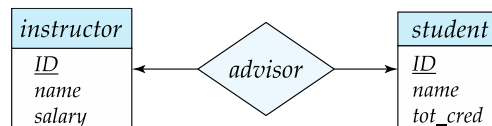
## Redundancy of Schemas

- Many-to-one and one-to-many relationship sets that are total on the many-side can be represented by adding an extra attribute to the “many” side, containing the primary key of the “one” side
- Example: Instead of creating a schema for relationship set *inst\_dept*, add an attribute *dept\_name* to the schema arising from entity set *instructor*
- Example
  - instructor*(ID, name, salary, *dept\_name*)
  - student*(ID, name, tot\_cred, *dept\_name*)



## Redundancy of Schemas (Cont.)

- For one-to-one relationship sets, either side can be chosen to act as the “many” side
  - That is, an extra attribute can be added to either of the tables corresponding to the two entity sets



- If participation is *partial* on the “many” side, replacing a schema by an extra attribute in the schema corresponding to the “many” side could result in null values



## Redundancy of Schemas (Cont.)

- The schema corresponding to a relationship set linking a weak entity set to its identifying strong entity set is redundant.
- Example: The *section* schema already contains the attributes that would appear in the *sec\_course* schema



## Extended E-R Features



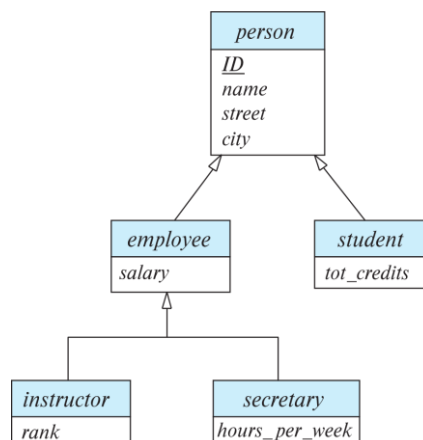
## Specialization

- Top-down design process; we designate sub-groupings within an entity set that are distinctive from other entities in the set.
- These sub-groupings become lower-level entity sets that have attributes or participate in relationships that do not apply to the higher-level entity set.
- Depicted by a *triangle* component labeled ISA (e.g., *instructor* “is a” *person*).
- **Attribute inheritance** – a lower-level entity set inherits all the attributes and relationship participation of the higher-level entity set to which it is linked.



## Specialization Example

- **Overlapping** – *employee* and *student*
- **Disjoint** – *instructor* and *secretary*
- Total and partial





## Representing Specialization via Schemas

- Method 1:
  - Form a schema for the higher-level entity
  - Form a schema for each lower-level entity set, include primary key of higher-level entity set and local attributes

schema	attributes
person	ID, name, street, city
student	ID, tot_cred
employee	ID, salary

- Drawback: getting information about, an *employee* requires accessing two relations, the one corresponding to the low-level schema and the one corresponding to the high-level schema



## Representing Specialization as Schemas (Cont.)

- Method 2:
  - Form a schema for each entity set with all local and inherited attributes

schema	attributes
person	ID, name, street, city
student	ID, name, street, city, tot_cred
employee	ID, name, street, city, salary

- Drawback: *name*, *street* and *city* may be stored redundantly for people who are both students and employees



## Generalization

- **A bottom-up design process** – combine a number of entity sets that share the same features into a higher-level entity set.
- Specialization and generalization are simple inversions of each other; they are represented in an E-R diagram in the same way.
- The terms specialization and generalization are used interchangeably.



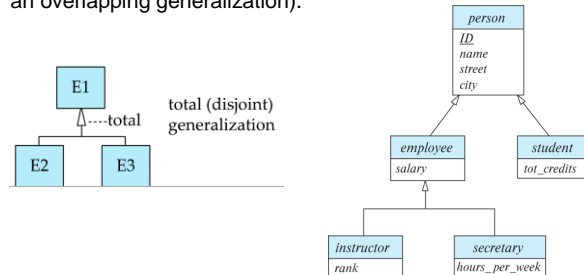
## Completeness constraint

- **Completeness constraint** -- specifies whether or not an entity in the higher-level entity set must belong to at least one of the lower-level entity sets within a generalization.
  - **total**: an entity must belong to one of the lower-level entity sets
  - **partial**: an entity need not belong to one of the lower-level entity sets



## Completeness constraint (Cont.)

- Partial generalization is the default.
- We can specify total generalization in an ER diagram by adding the keyword **total** in the diagram and drawing a dashed line from the keyword to the corresponding hollow arrow-head to which it applies (for a total generalization), or to the set of hollow arrow-heads to which it applies (for an overlapping generalization).

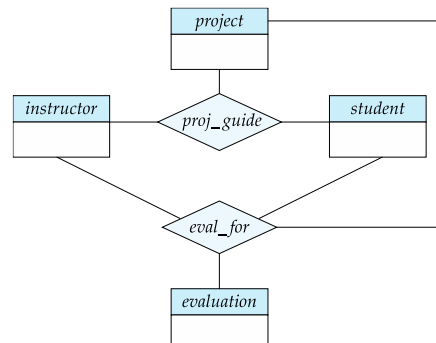


- The *student* generalization is total: All student entities must be either graduate or undergraduate. Because the higher-level entity set arrived at through generalization is generally composed of only those entities in the lower-level entity sets, the completeness constraint for a generalized higher-level entity set is usually total



## Aggregation

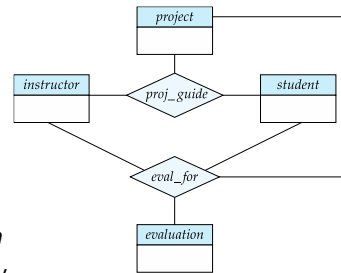
- Consider the ternary relationship *proj\_guide*, which we saw earlier
- Suppose we want to record evaluations of a student by a guide on a project





## Aggregation (Cont.)

- Relationship sets *eval\_for* and *proj\_guide* represent overlapping information
  - Every *eval\_for* relationship corresponds to a *proj\_guide* relationship
  - However, some *proj\_guide* relationships may not correspond to any *eval\_for* relationships
    - So we can't discard the *proj\_guide* relationship

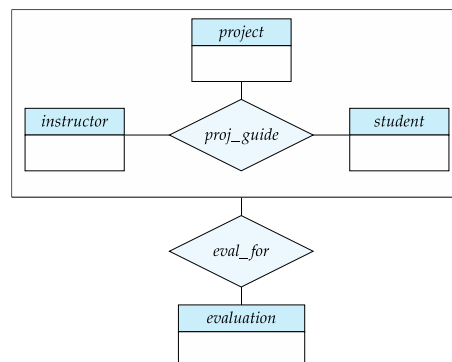


- Eliminate this redundancy via *aggregation*
  - Treat relationship as an abstract entity
  - Allows relationships between relationships
  - Abstraction of relationship into new entity



## Aggregation (Cont.)

- Eliminate this redundancy via *aggregation* without introducing redundancy, the following diagram represents:
  - A student is guided by a particular instructor on a particular project
  - A student, instructor, project combination may have an associated evaluation





## Reduction to Relational Schemas

- To represent aggregation, create a schema containing
  - Primary key of the aggregated relationship,
  - The primary key of the associated entity set
  - Any descriptive attributes
- In our example:
  - The schema *eval\_for* is:  
*eval\_for* (*s\_ID*, *project\_id*, *i\_ID*, *evaluation\_id*)
  - The schema *proj\_guide* is redundant.



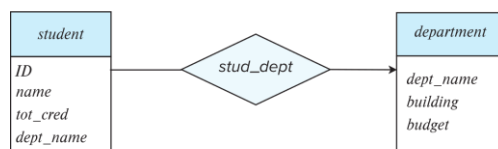
## Design Issues



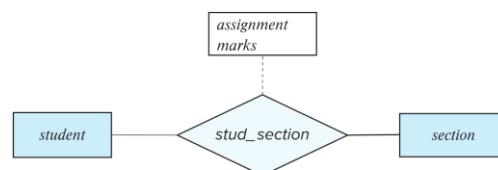


## Common Mistakes in E-R Diagrams

- Example of erroneous E-R diagrams



(a) Incorrect use of attribute

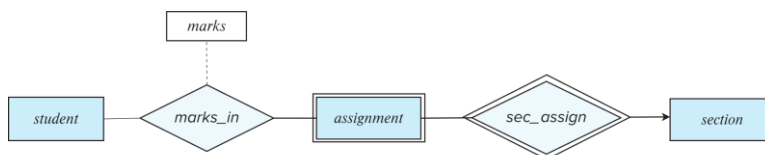


(b) Erroneous use of relationship attributes

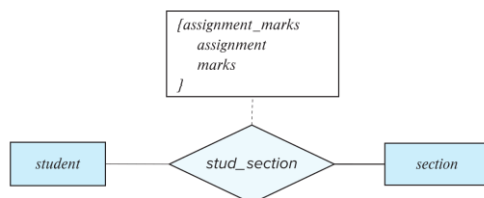


## Common Mistakes in E-R Diagrams (Cont.)

- Correct versions of the E-R diagram of previous slide



(c) Correct alternative to erroneous E-R diagram (b)

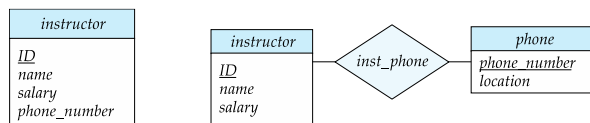


(d) Correct alternative to erroneous E-R diagram (b)



## Entities vs. Attributes

- Use of entity sets vs. attributes



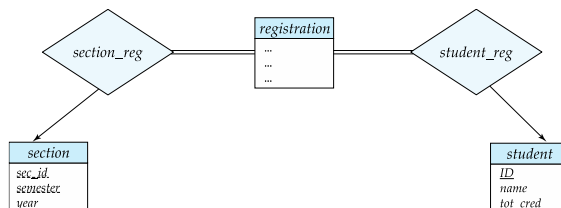
- Use of phone as an entity allows extra information about phone numbers (plus multiple phone numbers)



## Entities vs. Relationship sets

- Use of entity sets vs. relationship sets

Possible guideline is to designate a relationship set to describe an action that occurs between entities



- Placement of relationship attributes

For example, attribute date as attribute of advisor or as attribute of student



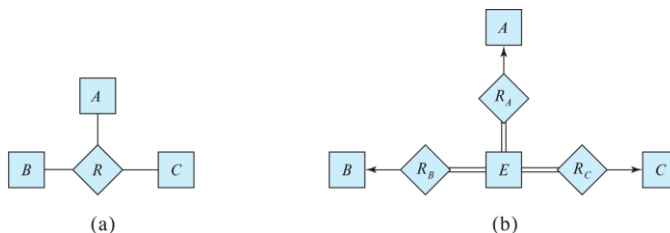
## Binary Vs. Non-Binary Relationships

- Although it is possible to replace any non-binary ( $n$ -ary, for  $n > 2$ ) relationship set by a number of distinct binary relationship sets, a  $n$ -ary relationship set shows more clearly that several entities participate in a single relationship.
- Some relationships that appear to be non-binary may be better represented using binary relationships
  - For example, a ternary relationship *parents*, relating a child to his/her father and mother, is best replaced by two binary relationships, *father* and *mother*
    - Using two binary relationships allows partial information (e.g., only mother being known)
  - But there are some relationships that are naturally non-binary
    - Example: *proj\_guide*



## Converting Non-Binary Relationships to Binary Form

- In general, any non-binary relationship can be represented using binary relationships by creating an artificial entity set.
  - Replace  $R$  between entity sets  $A$ ,  $B$  and  $C$  by an entity set  $E$ , and three relationship sets:
    1.  $R_A$ , relating  $E$  and  $A$
    2.  $R_B$ , relating  $E$  and  $B$
    3.  $R_C$ , relating  $E$  and  $C$
  - Create an identifying attribute for  $E$  and add any attributes of  $R$  to  $E$
  - For each relationship  $(a_i, b_i, c_i)$  in  $R$ , create
    1. a new entity  $e_i$  in the entity set  $E$
    2. add  $(e_i, a_i)$  to  $R_A$
    3. add  $(e_i, b_i)$  to  $R_B$
    4. add  $(e_i, c_i)$  to  $R_C$





## Converting Non-Binary Relationships (Cont.)

- Also need to translate constraints
  - Translating all constraints may not be possible
  - There may be instances in the translated schema that cannot correspond to any instance of  $R$ 
    - Exercise: *add constraints to the relationships  $R_A$ ,  $R_B$  and  $R_C$  to ensure that a newly created entity corresponds to exactly one entity in each of entity sets  $A$ ,  $B$  and  $C$*
  - We can avoid creating an identifying attribute by making  $E$  a weak entity set (described shortly) identified by the three relationship sets

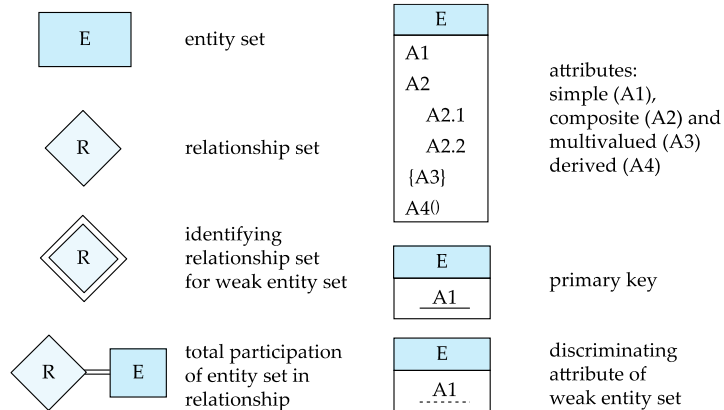


## E-R Design Decisions

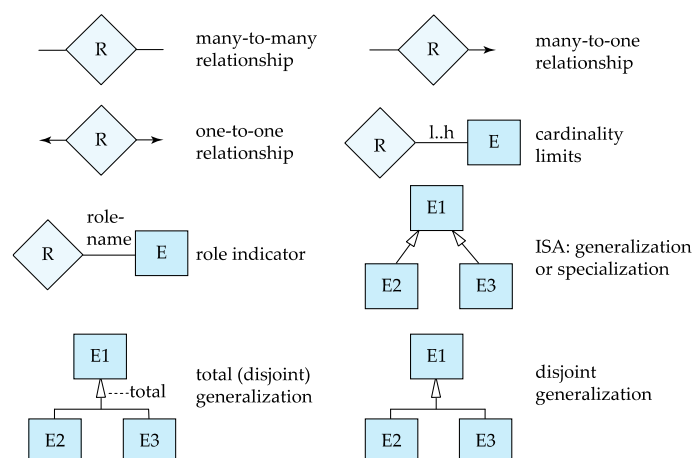
- The use of an attribute or entity set to represent an object.
- Whether a real-world concept is best expressed by an entity set or a relationship set.
- The use of a ternary relationship versus a pair of binary relationships.
- The use of a strong or weak entity set.
- The use of specialization/generalization – contributes to modularity in the design.
- The use of aggregation – can treat the aggregate entity set as a single unit without concern for the details of its internal structure.



## Summary of Symbols Used in E-R Notation



## Symbols Used in E-R Notation (Cont.)

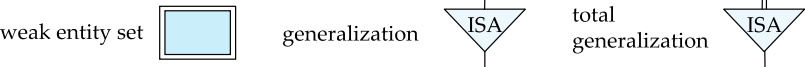
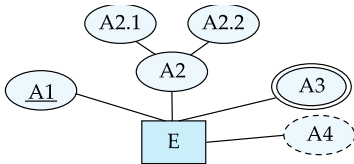




# Alternative ER Notations

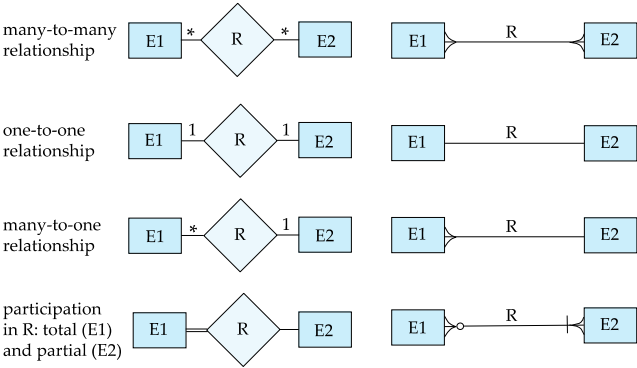
- Chen, IDE1FX, ...

entity set E with  
simple attribute A1,  
composite attribute A2,  
multivalued attribute A3,  
derived attribute A4,  
and primary key A1



# Alternative ER Notations

Chen IDE1FX (Crows feet notation)





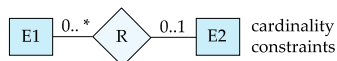
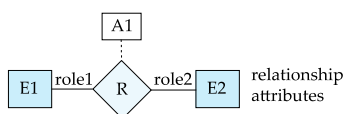
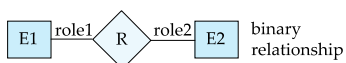
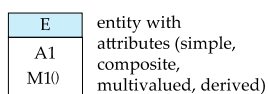
## UML

- **UML:** Unified Modeling Language
- UML has many components to graphically model different aspects of an entire software system
- UML Class Diagrams correspond to E-R Diagram, but several differences.

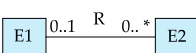
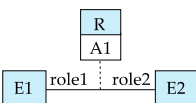
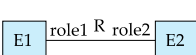
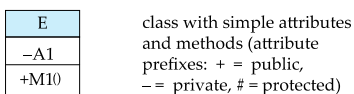


## ER vs. UML Class Diagrams

### ER Diagram Notation



### Equivalent in UML

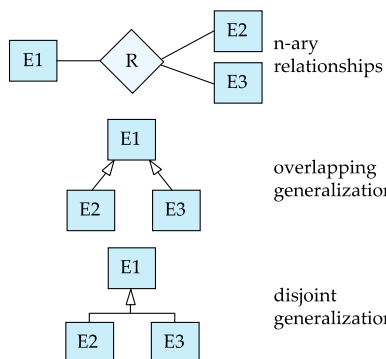


\* Note reversal of position in cardinality constraint depiction

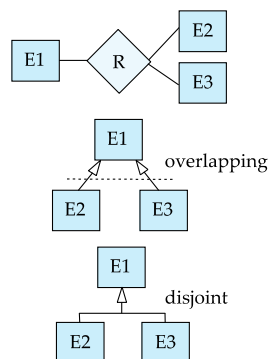


## ER vs. UML Class Diagrams

### ER Diagram Notation



### Equivalent in UML



\* Generalization can use merged or separate arrows independent of disjoint/overlapping



## UML Class Diagrams (Cont.)

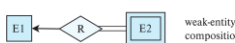
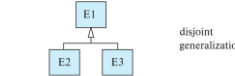
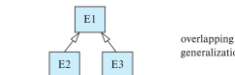
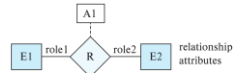
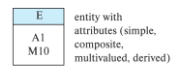
- Binary relationship sets are represented in UML by just drawing a line connecting the entity sets. The relationship set name is written adjacent to the line.
- The role played by an entity set in a relationship set may also be specified by writing the role name on the line, adjacent to the entity set.
- The relationship set name may alternatively be written in a box, along with attributes of the relationship set, and the box is connected, using a dotted line, to the line depicting the relationship set.



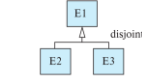
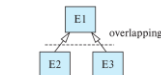
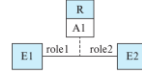
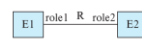
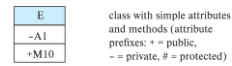


## ER vs. UML Class Diagrams

### ER Diagram Notation



### Equivalent in UML



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## Other Aspects of Database Design

- Functional Requirements
- Data Flow, Workflow
- Schema Evolution

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## End of Chapter 6