

# Information Theory: Entropy

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# Definition

Let  $X$  be a discrete random variable with probability distribution  $p(x)$ .

## Entropy $H_X$

$$H_X = - \sum_{x \in X} p(x) \log_2 p(x) = \mathbb{E} \log_2 \left[ \frac{1}{p(X)} \right].$$

Here we define  $0 \log_2 0 = 0$ .

In a sense, the entropy of a random variable shows how “uncertain” the event is.

# Examples

## Entropy of a fair coin

A fair coin has entropy  $\frac{1}{2} \log(2) + \frac{1}{2} \log(2) = 1$ .

## Entropy of a fair $m$ -sided die

A  $m$ -sided die has entropy  $\log m$ .

## Entropy of $n$ fair coin flips

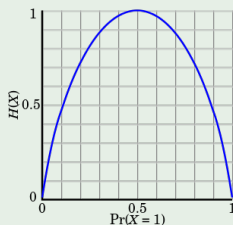
Flipping  $n$  coins produce  $2^n$  uniformly distributed possibilities, and hence, the entropy is  $n$ .

# Examples

## Entropy of a Bernoulli trial

If  $X$  is a random variable taking values between 0 and 1, where  $p(0) = p$  and  $p(1) = 1 - p$ , its entropy is

$$H(p) = -p \log p - (1 - p) \log(1 - p).$$



# Examples

## Entropy of an unfair dice

An unfair dice with four faces and

$$p(1) = 1/2, p(2) = 1/4, p(3) = 1/8, p(4) = 1/8$$

has entropy  $7/4$ , smaller than the one of the corresponding fair dice, which is 2. (This dice is less uncertain than the fair dice.)

# Motivation [TODO]

Define a function  $H$  that takes in a random variable and outputs an integer, such that:

- If a random variable  $X$  takes  $n$  values, then  $H(X)$  is maximized if  $X$  is uniform.
- Entropy is additive, in the sense that: if  $X$  takes  $x_i$  with probability  $p_i$ ,  $Y$  takes  $x_i$  with probability  $q_i$ , and  $Z$  takes  $x_i$  with probability  $\alpha p_i + \beta q_i$  for all  $i$ , then

$$H(Z) = \alpha H(X) + \beta H(Y).$$

Then  $H(X) = -\sum p_i \log p_i$  is the only possible function (up to the base of the logarithm).

# Properties (TPM 15.7.1)

- 1  $H(X) \leq \log |S|.$
- 2  $H(X) \leq H(X, Y) \leq H(X) + H(Y).$
- 3  $H(X|Y, Z) \leq H(X|Y).$

## Entropy is subadditive (TPM 15.7.2)

Let  $X = (X_1, \dots, X_n)$  be a random variable taking values in the set  $S = S_1 \times S_2 \times \dots \times S_n$ , where each of the coordinates  $X_i$  of  $X$  is a random variable taking values in  $S_i$ . Then

$$H(X) \leq \sum_{i=1}^n H(X_i).$$

(This is just induction on the property  $H(X, Y) \leq H(X) + H(Y).$ )

## TPM 15.7.3

Let  $S$  be a family of subsets of  $\{1, 2, \dots, n\}$  and let  $p_i$  denote the fraction of sets in  $S$  that contain  $i$ . Then

$$|S| \leq 2^{\sum H(p_i)}.$$

### Proof

Let  $X = (X_1, \dots, X_n)$  take elements in  $S$  with equal probability. Then  $H(X) \leq \sum H(X_i)$  implies  $\log |S| \leq \sum H(p_i)$ .



## TPM 15.7.4

Let  $X = (X_1, \dots, X_n)$  taking values in  $S = S_1 \times \dots \times S_n$ , where each  $X_i$  takes values in  $S_i$ . For an index set  $I \subseteq N = \{1, 2, \dots, n\}$  let  $X(I)$  denote  $(X_i)_{i \in I}$ . If  $T$  is a family of subsets of  $N$  and each  $i \in N$  belongs to at least  $k$  members of  $T$ , then

$$kH(X) \leq \sum_{G \in T} H(X(G)).$$

### Proof

Use induction. If there is  $G \in T$  where  $G = N$  we are done. Otherwise, we prove

$$H(X(G \cup G')) + H(X(G \cap G')) \leq H(X(G)) + H(X(G')).$$

## TPM 15.7.5

Take  $F \subseteq S_1 \times S_2 \times \cdots \times S_n$ . Let  $\mathcal{I} = \{I_1, I_2, \dots, I_m\}$  be a collection of index sets (i.e. subsets of  $N$ ), and suppose that each element  $i \in N$  belongs to at least  $k$  members of  $\mathcal{I}$ . For each  $1 \leq i \leq m$  let  $F_i$  be the set of all projections of the members of  $F$  on  $I_i$ . Then

$$|F|^k \leq \prod_{i=1}^m |F_i|.$$

### Proof

Let  $X = (X_1, \dots, X_n)$  take elements of  $F$  with equal probability. Then  $kH(X) \leq \sum H(X(I_i))$  implies  $k \log |F| \leq \sum \log |F_i|$ .

## Corollaries of TPM 15.7.5

Take  $I_i = N - \{i\}$  for all  $i$ , and notice that the volume of a set can be approximated by a fine enough aligned boxes.

### TPM 15.7.6

Let  $B$  be a measurable body in the  $n$ -dimensional Euclidean space, let  $\text{Vol}(B)$  denote its ( $n$ -dimensional) volume, and let  $\text{Vol}(B_i)$  denote the  $(n - 1)$ -dimensional volume of the projection of  $B$  on the hyperplane spanned by all coordinates besides the  $i$ -th one.

Then

$$(\text{Vol}(B))^{n-1} \leq \prod_{i=1}^n \text{Vol}(B_i).$$

# Corollaries of TPM 15.7.5

Take  $S_i = \{0, 1\}$  in TPM 15.7.5.

## TPM 15.7.7

Let  $F$  and  $\mathcal{I} = \{I_1, I_2, \dots, I_m\}$  be a collection of subsets of  $N$ . Suppose that each element of  $N$  belongs to at least  $k$  members of  $\mathcal{I}$ . For each  $1 \leq i \leq m$ , define  $F_i = \{f \cap I_i : f \in F\}$ . Then

$$|F|^k \leq \prod_{i=1}^m |F_i|.$$

# Corollaries of TPM 15.7.5

## TPM 15.7.8

Let  $F$  be a family of graphs on the labeled set of vertices  $\{1, \dots, t\}$ , and suppose that for any two members of  $F$  there is a triangle contained in both of them. Then

$$|F| < \frac{1}{4} 2^{\binom{t}{2}}.$$

## Proof

Take  $N$  to be the set of all edges of  $K_t$  and  $\mathcal{I}$  be edges of all possible  $K_{\lfloor \frac{t}{2} \rfloor} \cup K_{\lceil \frac{t}{2} \rceil}$ .