# Quantum Aspects of Cryptography

Assignments 10 and 11—BQP $\neq$ QMA, yet PRU exist, Self-testing, QFHE and Verification (topics from Lectures 22 to 27) Version:  $\alpha.1$ —April 25, 2025

**Instructions.** Same as those in previous assignments (including the updates since Assignment 5) except that *May 6*, *2025* for this course is a *hard deadline* for all assignments—no unused extension can be availed to extend beyond that.

- 1. If your name is *Alice* and you're submitting answers to *Assignment 10*, use Alice10.pdf as your filename when submitting.
- 2. Submit your assignment using the appropriate link below and add the submission date to this Google spreadsheet.<sup>1</sup>

Please let me know if you spot a mistake or if something is unclear or feels suspicious.

<sup>&</sup>lt;sup>1</sup>The system automatically adds a date-time stamp when you upload it to the OneDrive folder but I didn't want to spend time writing a script to fetch this data into Google sheets.

# Assignment 10—BQP=QMA yet PRUs exist and QFHE

- Submission link: OneDrive link for Assignment 10.
- Due: Tuesday, May 6, 2025 (no extensions)

### A. BQP=QMA yet PURs exist

This is based on by Kretschmer (v5 on arXiv) [3]. While there seems to be a gap in one of the steps, let us review what we covered in class anyway.

**Exercise 1.** The following argument shows that QMA can distinguish PRS from Haar random states. What is wrong with the argument?

QMA protocol:

Arthur holds many copies of  $|\psi\rangle$ .

Merlin gives the circuit C that produces  $|\psi\rangle$  to Arthur.

Arthur does a swap test between  $|\psi\rangle$  and  $C|0^n\rangle$ .

If  $|\psi\rangle$  is PRS, such a circuit C exists and Merlin can convince Arthur that the state is PRS and not Haar random.

### **Exercise 2.** Pre-requisites. State the following.

- 1. From §2.2 Probability
  - (a) Lemma 5 (Baye's decision rule); also briefly explain how this shows one can't do better than Baye's rule for guessing
  - (b) Lemma 6 (Borel-Cantelli); also briefly explain how this result can be seen as a criterion under which at most finitely many events occur with probability 1 (taking  $X_i$  to denote whether event i occurs or not).
- 2. From §2.3 Quantum Information: Fact 7, Fact 8, Lemma 9
- 3. From §2.4 Haar measure and Concentration:
  - (a) Definition of  $\mathbb{U}(N)^M$ ,  $\mu_N^M$
  - (b) Definition of *L*-Lipschitz
  - (c) Theorem 10
- 4. From §2.5 Complexity theory
  - (a) Definition of a language, a promise problem and the  $Dom(\Pi)$  notation for a promise problem  $\Pi$
  - (b) Definition 11 (Promise BQP) and 11 (PromiseQMA)
- 5. From §2.6 Quantum Oracles,
  - (a) explain how controlled- $\mathcal{U}$  can be viewed as  $\mathbb{I} \oplus \mathcal{U}$ .
  - (b) explain how the 'query cost' is defined for queries  $\mathcal{U} = \{\mathcal{U}_n\}_{n \in \mathbb{N}}$ .
- 6. From §2.7 Cryptography
  - (a) Definition 16 (Pseudorandom unitary transformations)
  - (b) Briefly explain why, in this work,  $n(\kappa)$  the number of qubits on which the pseudorandom unitary acts, is taken to be at least  $\omega(\log \kappa)$  where  $\kappa$  is the security parameter.

Exercise 3. (2x) This is based on §5 Pseudorandomness from a quantum oracle.

- 1. From §5.1, define the language  $\mathcal C$
- 2. From §5.2 PromiseBQP = PromiseQMA relative to  $(\mathcal{U}, \mathcal{C})$ 
  - (a) Lemma 28 and its proof
  - (b) Lemma 29 and its proof

Exercise 4. (10x) State and prove Theorem 30; identify the step that has a gap in its justification (we discussed this in class).

### Exercise 5. From §5.3 Pseudorandom unitaries

- 1. State the construction for the PRU ensemble used by the author
- 2. State Lemma 31
- 3. (2x) Prove Lemma 31

### B. QFHE

The following refers to statements in Appendix A of KLVY '22[2].

### Exercise 6. Answer the following

- 1. Explain informally what is meant by homomorphic encryption.
- 2. Write down Definition A1.
- 3. Explain briefly the connection between your informal explanation and the formal definition.

Answer the following, semi-formally, about Mahadev's encryption scheme (as described in Appendix A of [2]).

**Exercise 7.** Explain the encryption procedure Mahadev uses.

**Exercise 8.** Explain how the Eval function is implemented for Clifford gates

**Exercise 9.** Explain why this strategy fails for the Toffoli gate, T, by explicitly writing down the extra factors that appear when one moves a Toffoli past the Pauli pad.

**Exercise 10.** Explain what additional properties of trapdoor claw-free functions are assumed by Mahadev for implementing Eval for T.

**Exercise 11** (3x). Explain how Eval is implemented for T gates.

# Assignment 11-Self-testing and Verification

- Submission link: OneDrive link for Assignment 11.
- Due: Tue, May 6, 2025 (no extensions)

# A. Self-testing

The following refers to statements in Section 4 of the review paper by Šupić and Bowles '20 [4].

**Exercise 12.** Semi-formally, describe what is meant by self-testing.

Exercise 13. Setting up the general strategy

- 1. Explain how  $A_x$  and  $B_y$  are defined in terms of projectors and why this is without loss of generality.
- 2. Show Eq (20), i.e. these observables are Herimitian and unitary.

### Exercise 14. Bell operator.

- 1. Write down the expression for  $\beta_{CHSH}$  and explain how  $\langle A_x B_y \rangle$  are to be interpreted.
- 2. For what choices of  $\{A_x\}_x$  and  $\{B_y\}_y$  and state  $|\psi\rangle$  do we get  $\beta_{CHSH}$  to be  $2\sqrt{2}$  (don't need to prove it if you don't feel like it)

### Exercise 15. SoS and anti-commutation

- 1. Explain how the SoS decomposition technique can be used to deduce that  $P_{\lambda} | \psi \rangle = 0$  for all  $\lambda$  where  $\beta_Q \mathbb{I} \mathcal{B} = \sum_{\lambda} P_{\lambda}^{\dagger} P_{\lambda}$ ,  $\mathcal{B}$  is a Bell operator (see Eq (26)) and  $\beta_Q$  is the maximum quantum value (achieved by  $|\psi\rangle$ ).
- 2. Use Eq (30) and the argument above, to deduce that  $\frac{(A_0\pm A_1)}{\sqrt{2}}\,|\psi\rangle=B_{0/1}\,|\psi\rangle$ .
- 3. Using your answer above, show that  $B_0$  and  $B_1$  anti-commute (simply expand  $\{B_0, B_1\} | \psi \rangle$  in terms of  $A_0 \pm A_1$ ).

### Exercise 16. Swap Isometry

- 1. Write down the Swap Isometry  $\Phi$  as shown in Figure 4 of the paper.
- 2. Show that this indeed acts the Swap isometry when  $Z_A, Z_B$  are taken to be the Pauli z operators and  $X_A, X_B$  are taken to be the Pauli x operators, and  $|\psi\rangle_{AB}$  is a 2-qubit state.
- 3. Suppose  $Z_A = \frac{1}{\sqrt{2}} \left( A_0 + A_1 \right), Z_B = B_0, X_A = \frac{1}{\sqrt{2}} (A_0 A_1), X_B = B_1.$ 
  - (a) Show that  $\{Z_A, X_A\} = 0$
  - (b) Also show that  $\{Z_B, X_B\} | \psi \rangle = 0$
  - (c)  $Z_A |\psi\rangle = Z_B |\psi\rangle$  and  $X_A |\psi\rangle = X_B |\psi\rangle$
- 4. (5x) For simplicity, suppose  $\Phi$  is an isometry even for the choice of Zs and Xs we made above. Now, establish that

$$\Phi[|\psi\rangle] = \sum_{i,j\in\{0,1\}} |ij\rangle_{A'B'} \otimes \underbrace{\left(\frac{1}{4}X_A^i(\mathbb{I} + (-1)^i Z_A)X_B^j(\mathbb{I} - (-1)^j Z_B)\right)}_{=:\hat{f}_{ij}} |\psi\rangle_{AB}.$$

- 5. Prove that  $\hat{f}_{01} | \psi \rangle = \hat{f}_{10} | \psi \rangle = 0$ .
- 6. Prove that  $\hat{f}_{11} | \psi \rangle = \hat{f}_{00} | \psi \rangle$ .
- 7. Explain how existence of  $\Phi$  shows that one can self-test the Bell state  $\frac{|00\rangle+|11\rangle}{\sqrt{2}}$  (be consistent with your answer to ??).

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### **B.** Verification

This last part is based on Grilo '20 [1].

**Exercise 17.** Background/Self-testing. State the following.

- 1. Lemma 1 and also explain the magic square game
- 2. Pauli Braiding Test (as written in Figure 1)
- 3. Theorem 2

Exercise 18. Background/Local Hamiltonian problem. State the following.

- 1. Describe what is meant by a local hamiltonian problem
- 2. Definition 3 (XZ Local Hamiltonian)
- 3. Lemma 4 (QMA completeness of XZ local hamiltonians)
- 4. Lemma 5 (amplifying soundness/completeness of XZ Local Hamiltonians)
- 5. Definition 6 (What is meant by non-local games for Local Hamiltonians)

**Exercise 19.** Write down the protocol in Figure 2 and intuitively (to the extent reasonable) explain what is going on.

Exercise 20. (3x) State and prove Lemma 7 (bound on max winning probability of semi-honest strategies).

Exercise 21. (3x) State and prove Lemma 8 (bound on general strategies)

Exercise 22. (3x) State and prove Theorem 9 (uses Lemma 7 and Lemma 8 to satisfy Definition 6; then amplifies the soundness/completeness gap)

Exercise 23. State Corollary 9.

# References

- [1] Alex B. Grilo. A simple protocol for verifiable delegation of quantum computation in one round, 2020.
- [2] Yael Kalai, Alex Lombardi, Vinod Vaikuntanathan, and Lisa Yang. Quantum advantage from any non-local game, 2022.
- [3] William Kretschmer. Quantum pseudorandomness and classical complexity. 2021.
- [4] Ivan Šupić and Joseph Bowles. Self-testing of quantum systems: a review. Quantum, 4:337, September 2020.