

Below several pairs of predicates are given. For each pair do the following:

- Identify the atomic assertions common to each pair of assertions and assign a variable to each of these assertions.
- Find logical expressions for each sentence in terms of the variables.
- Determine whether the first can be logically deduced from the second, and whether the second can be logically deduced from the first. Explain your answers.

- if  $n$  is prime or  $n + 2$  is prime, then  $n^2 + 2$  is prime or  $n^2 - 2$  is prime
  - $n^2 + 2$  is non-prime and  $n^2 - 2$  is non-prime implies  $n$  is non-prime and  $n + 2$  is non-prime.

$A := n$  is prime.

$B := n + 2$  is prime.

$C := n^2 + 2$  is prime.

$D := n^2 - 2$  is prime.

$$(a) (A \vee B) \Rightarrow (C \vee D)$$

$$(b) \neg C \wedge \neg D \Rightarrow \neg A \wedge \neg B$$

$$(A \vee B) \Rightarrow (C \vee D) \leftrightarrow \neg(C \vee D) \Rightarrow \neg(A \vee B) \text{ (Modus Tollens)}$$

$$\neg C \wedge \neg D \Rightarrow \neg A \wedge \neg B \leftrightarrow \neg(C \vee D) \Rightarrow \neg(A \vee B) \text{ (DeMorgan's rule)}$$

Therefore the two statements are logically equivalent.

- For all real numbers  $x$ , there is a real number  $y$  such that  $y^2 + y + 10x = 0$  or  $x \leq 9$  and there is a real number  $z$  such that  $z^2 + 2z + 15x = 0$ .
  - For all real numbers  $x$ ,  $x \leq 9$  or there is both a real number  $y$  such that  $y^2 + y + 10x = 0$  and a real number  $z$  such that  $z^2 + 2z + 15x = 0$ .

$$A(x, y) := y^2 + y + 10x = 0$$

$$B(x) := x \leq 9$$

$$C(x, z) := z^2 + 2z + 15x = 0$$

$$(a) (\forall x \in \mathbb{R}, \exists y \in \mathbb{R}, (A(x, y) \vee B(x))) \wedge (\exists z \in \mathbb{R}, C(x, z))$$

$$(b) (\forall x \in \mathbb{R}, B(x)) \vee (\exists y \in \mathbb{R}, \exists z \in \mathbb{R}, (A(x, y) \wedge C(x, z)))$$

(b) cannot imply (a) as if  $A = T, B = T, C = F$  then  $T \vee (T \wedge F) = T \vee F = T \not\Rightarrow (T \vee T) \wedge F = T \wedge F = F$ . However (a) implies (b) as if  $(\forall x \in \mathbb{R}, \exists y \in \mathbb{R}, (A(x, y) \vee B(x))) \wedge (\exists z \in \mathbb{R}, C(x, z))$  is true then  $(A \wedge B)$  is true. Since it is an or statement we can choose  $\forall x \in \mathbb{R} B(x)$  to be true. Since the highest level logical operator in (b) is an or, then  $B(x) = T$  satisfies the statement.

- $f(x) > y$  and  $g(y) > x$  implies  $f(g(y)) > y$  and  $g(f(x)) > x$ .

(b)  $f(g(y)) \leq y$  implies  $f(x) \leq y$ , and  $g(f(x)) \leq x$  implies  $g(y) \leq x$ .

$$\begin{aligned} A(x, y) &:= f(x) > y \\ B(x, y) &:= g(y) > x \\ C(y) &:= f(g(y)) > y \\ D(x) &:= g(f(x)) > x \end{aligned}$$

(a)  $(A(x, y) \wedge B(x, y)) \Rightarrow (C(y) \wedge D(x))$

(b)  $(\neg C(y) \Rightarrow \neg A(x, y)) \wedge (\neg D(x) \Rightarrow \neg B(x, y))$

(a)  $\neg(A \wedge B) \vee (C \wedge D)$  (Implication definition)

(a)  $((\neg A \vee \neg B) \vee C) \wedge ((\neg A \vee \neg B) \vee D)$  (DeMorgan's rule)

(a)  $(\neg A \vee (\neg B \vee C)) \wedge (\neg A \vee (\neg B \vee D))$  (or associativity)

(a)  $(A \Rightarrow (B \Rightarrow C)) \wedge (A \Rightarrow (B \Rightarrow D))$  (b)  $(A \Rightarrow C) \wedge (B \Rightarrow D)$  (Contrapositive)

a does not logically imply b as in the case where  $A$  is true and the rest are false then  $(T \wedge F) \Rightarrow (F \wedge F) = F \Rightarrow F = T \neq (T \Rightarrow (F \wedge F)) \wedge (F \Rightarrow (F \wedge F)) = (T \Rightarrow F) \wedge (F \Rightarrow F) = F \wedge T = F$ . However b implies a, as if  $(A \Rightarrow C) \wedge (B \Rightarrow D) = T$ ,  $(A \Rightarrow C) = T$ ,  $(B \Rightarrow D) = T$ , then since we already know  $(A \Rightarrow C) \Rightarrow (A \Rightarrow (B \Rightarrow C))$  from problem 2, then we have  $T \wedge (A \Rightarrow T) = T \wedge (\neg A \vee T) = T \wedge T = T$ .

4. In this pair of assertions,  $S, T, V$ , and  $W$  are all sets.

(a)  $S \subseteq T$  if and only if  $S \subseteq V$ , or  $S \subseteq T$  if and only if  $W \subseteq T$

(b)  $S \subseteq T$  if and only if  $(S \subseteq T \text{ or } W \subseteq T)$ .

$$\begin{aligned} A(S, T) &:= S \subseteq T \\ B(S, V) &:= S \subseteq V \\ C(T, W) &:= W \subseteq T \end{aligned}$$

(a)  $A(S, T) \Leftrightarrow B(S, V) \vee A(S, T) \Leftrightarrow C(T, W)$

(b)  $A(S, T) \Leftrightarrow (B(S, V) \vee C(T, W))$

b doesn't imply a since if  $A = F, B = T, C = F$ , then  $T \Leftrightarrow F \vee F \Leftrightarrow F = F \vee T = T \neq F \Leftrightarrow (T \vee F) = F \Leftrightarrow T = F$ . However it does work the other way around, if we take (b) to be true then we have  $A \Leftrightarrow (B \vee C)$ . Since we have an or statement, then we can just take  $B$  out of it, and we're left with the statement  $A \Leftrightarrow B$ , which if substituted into a as true, then we get  $T \vee A \Leftrightarrow C = T$ .