

Wednesday, February 7, 2024

**1. Recurrence Relations.** Give the big- $\Theta$  bound for each of the following recurrence relations. You may assume the Master Theorem gives a big- $\Theta$  bound (the proof is similar to the big- $O$  proof shown in class) and that  $T(1) = O(1)$ . Justify each of your answers.

- a)  $T(n) = 4T(n/2) + 42n$
- b)  $T(n) = T(n-1) + n^3$
- c)  $T(n) = 7T(n/2) + n^2 + \log n$
- d)  $T(n) = T(n/3) + 5$

**2. Recurrence Relations.** What are the running times of each of these algorithms in big- $O$  notation? Which is the fastest?

- a) Algorithm A solves problems by dividing them into five subproblems of half the size, recursively solving each subproblem, and then combining the solutions in linear time.
- b) Algorithm B solves problems of size  $n$  by recursively solving two subproblems of size  $n-1$  and then combining the solutions in constant time.
- c) Algorithm C solves problems of size  $n$  by dividing them into eight subproblems of size  $\frac{n}{4}$ , recursively solving each subproblem, and then combining the solutions in  $O(n^3)$  time.

**3. Merge.** A  $k$ -way merge operation. Suppose you have  $k$  sorted arrays, each with  $n$  elements, and you want to combine them into a single sorted array of  $kn$  elements.

- a) Here's one strategy: Using the merge procedure, merge the first two arrays, then merge in the third, then merge in the fourth, and so on. What is the time complexity of this algorithm, in terms of  $k$  and  $n$ ?
- b) Give a more efficient solution to this problem that runs in  $O(nk \log k)$  using divide-and-conquer.

**4. Array Rotations.** Consider a rotation operation that takes an array and moves its last element to the beginning. After  $n$  rotations, an array  $[a_0, a_1, a_2, \dots, a_{m-1}]$  of size  $m$  where  $0 < n \leq m$ , will become:

$$[a_{m-n}, a_{m-n+1}, \dots, a_{m-2}, a_{m-1}, a_0, a_1, \dots, a_{m-n-1}]$$

Notice how  $a_{m-1}$  is adjacent to  $a_0$  in the middle of the new array. For example, two rotations on the array  $[1, 2, 3, 4, 5]$  will yield  $[4, 5, 1, 2, 3]$ .

You are given a list of unique integers `nums`, which was previously sorted in ascending order, but has now been rotated an unknown number of times. Find the number of rotations in  $O(\log n)$  time. (*Hint: consider Binary Search.*)

**5. More Recurrence Relations.** Give the big- $\Theta$  bound for each of the following recurrence relations. You may assume that  $T(1) = O(1)$ . Justify each of your answers.

a)  $T(n) = 2T(2n/3) + T(n/3) + n^2$

b)  $T(n) = 3T(n/4) + n \log n$