Operating Systems

Xinu – Internals List and Queue Manipulation

Unified Lists in Xinu

• The process manager handles processes

1

- A process moves among the lists frequently
- At any time, a process is only in one list
- Rather than store all the information about a process, the process manager is free to store only the process ID (PID) or Thread ID (TID) in a list
 - So when we refer to to putting a process on a list, it really means putting the PID there – the process control block need not move
- Unified implementation means that not every subsystem uses all the list features

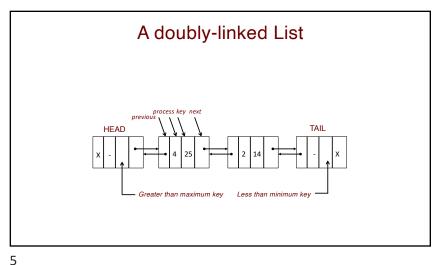
Linked Lists in the OS

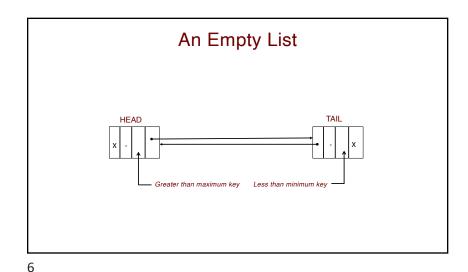
- Manipulating lists of processes is an important operation in the OS
 - A process's lifecycle consists of moving between, and in, queues and lists
- Xinu implements a unified approach to list management
 - All list management uses this common infrastructure
 - Common functions to create a new list, insert an element at the end of the list, insert or remove from the middle, remove an item from the front
- We can think of this code as being sequential
 - Protected with mutual exclusion as necessary

2

List Properties

- All lists are doubly linked each node points to predecessor and successor
- Each node stores a key as well as a process ID, even though a key is not used in a FIFO list
- Each list has head and tail nodes; the head and tail nodes have the same memory layout as other nodes
- Non-FIFO lists are ordered in descending order; the key in the head node is greater than the maximum valid key value, and the key value in the tail node is less than the minimum valid key





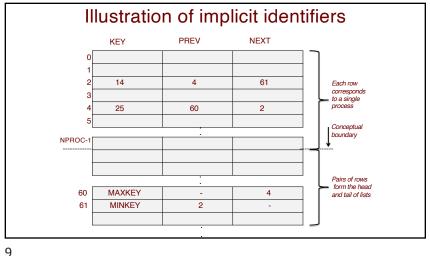
A compact list structure

- Compact memory use is important
 - in general, but in embedded systems in particular
- Two ideas related to process specific representation
- Relative pointers
 - Given that there is some (small) fixed number of processes (NPROC)
 - One might use a pointer in this situation, which is 4 bytes
 - For NPROCS < 62, we only need 6 bits
 - Allocate the nodes in a contiguous array and use the array index as a "pointer"

A compact list structure

- Implicit data structure based on the fact that a process can only be in one list, we can use the list position to indicate the
- To omit the PID, use an array and use the ith element of the array for process i
- To put process 3 in a linked list, insert node 3 into the list
- The relative address of a node is the same as the ID of a process being stored

- List membership is implicit



Notes

- The table contains NQENT entries
- There is an implicit divider at NPROCS-1
- · Between queuetab[NPROC] and queuetab[NQENT-1] are the heads and tails of lists
- The default NPROC + 4 + NSEM + NSEM allocates enough space for
 - each process

10

- head and tail entries for the ready and sleep lists
- a head and tail entry for each of the NSEM semaphores in the system
- · This can be changed at compile time

```
Implementation of the Queue Data
            Structure
```

```
/* queue.h - firstid, firstkey, isempty, lastkey, nonempty
/* Queue structure declarations, constants, and inline functions */
/* Default # of queue entries: 1 per process plus 2 for ready list plus*/
     2 for sleep list plus 2 per semaphore */
#ifndef NQENT
#define NQENT (NPROC + 4 + NSEM + NSEM)
#endif
                      /* null value for qnext or qprev index
#define MAXKEY 0x7FFFFFFF /* max key that can be stored in queue
#define MINKEY 0x80000000 /* min key that can be stored in queue */
```

```
/* one per process plus two per list */
struct gentry {
                     /* key on which the queue is ordered */
  int32 qkey;
  qid16
          qnext;
                     /* index of next process or tail */
  qid16 qprev;
                     /* index of previous process or head */
};
extern struct gentry queuetab[];
/* Inline queue manipulation functions */
#define queuehead(q) (q)
\#define queuetail(q) ((q) + 1)
#define firstid(q) (queuetab[queuehead(q)].qnext)
#define lastid(q) (queuetab[queuetail(q)].qprev)
#define isempty(q) (firstid(q) >= NPROC)
#define nonempty(q)(firstid(q) < NPROC)
#define firstkey(q) (queuetab[firstid(q)].qkey)
#define lastkey(q) (queuetab[ lastid(q)].qkey)
```

```
/* assumes interrupts are disabled */
/* Inline to check queue id */

#define isbadqid(x)(((int32)(x) < 0) || (int32)(x) >= NQENT-1)

/* Queue function prototypes */

pid32 getfirst(qid16);
pid32 getlast(qid16);
pid32 getitem(pid32);
pid32 enqueue(pid32, qid16);
pid32 dequeue(qid16);
status insert(pid32, qid16, int);
status insertd(pid32, qid16, int);
qid16 newqueue(void);
```

Basic Functions to Extract A Process From A List

13

FIFO Queue Manipulation

```
{
  int tail, prev;    /* tail & previous node indexes */
  if (isbadqid(q) || isbadpid(pid)) {
    return SYSERR;
  }

  tail = queuetail(q);
  prev = queuetab[tail].qprev;

  queuetab[pid].qnext = tail;    /* insert just before tail node */
  queuetab[pid].qprev = prev;
  queuetab[prev].qnext = pid;
  queuetab[tail].qprev = pid;
  return pid;
}
```

17

Manipulation of Priority Queues

```
/* runs through items in a queue*/
  int16 curr:
  int16 prev;
                        /* holds previous node index */
  if (isbadqid(q) || isbadpid(pid)) {
  return SYSERR;
  curr = firstid(q);
  while (queuetab[curr].qkey >= key) {
      curr = queuetab[curr].qnext;
/* insert process between curr node and previous node */
  prev = queuetab[curr].qprev; /* get index of previous node */
  queuetab[pid].qnext = curr;
  queuetab[pid].qprev = prev;
  queuetab[pid].qkey = key;
  queuetab[prev].qnext = pid;
  queuetab[curr].qprev = pid;
  return OK;
```

21

List Initialization