# CSU44004/CSU55004: FORMAL VERIFICATION

# Lecture 1: Module Overview & Introduction

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#### THE PERILS OF SOFTWARE

- → Software is now controlling critical machines:
  - → Transportation: cars (> 100M LoC [IEEE]), airplains, trains, spacecraft, ...
  - → Medical: pacemakers, MRI machines, ...
  - ightarrow Utilities: power grids, telephone centres, ...
  - $\rightarrow$  Finance: online banking, stock prices, ...
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  - → ...
- → BUT Software is very unreliable
  - → http://en.wikipedia.org/wiki/List\_of\_software\_bugs
  - → http://www.cs.tau.ac.il/~nachumd/horror.html













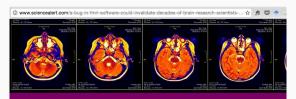


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# A bug in fMRI software could invalidate 15 years of brain research

#### This is huge.

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There could be a very serious problem with the past 15 years of research into human brain activity, with a new study suggesting that a bug in fMRI software could invalidate the results of some 40,000 papers.

That's massive, because functional magnetic resonance imaging (fMRI) is one of the best tools we have to measure brain activity, and if it's flawed, it means all those conclusions about what our brains look like during things like exercise, gaming, love, and drug addiction are wrong.

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#### TIMSORT

- → Hybrid Sorting Algorithm of MergeSort + InsertionSort
- → Used in Android JDK, Sun JDK, OpenJDK, Python, GNU Octave
- → The implementation uses an intermediate array of fixed size
- → The (wrong) assumption is that the array will never fill up
- → But it does fill up for carefully selected inputs arrays of size > 562 trillion

http://www.envisage-project.eu/

proving-android-java-and-python-sorting-algorithm-is-broken-and-how-to-fix-it/

# OpenJDK's java.utils.Collection.sort() is broken: The good, the bad and the worst case $^{\star}$

Stijn de Gouw<sup>1,2</sup>, Jurriaan Rot<sup>3,1</sup>, Frank S. de Boer<sup>1,3</sup>, Richard Bubel<sup>4</sup>, and Reiner Hähnle<sup>4</sup>

- CWI, Amsterdam, The Netherlands
- SDL, Amsterdam, The Netherlands
   Leiden University, The Netherlands
- <sup>4</sup> Technische Universität Darmstadt, Germany

Abstract. We investigate the correctases of TimSort, which is the main serting algorithm provided by the Jawa standard library. The goal is functional verification with mechanical proofs. During our verification strept we discovered a bug which causes the implementation to crash. We characterize the conditions under which the bug occurs, and from this we derive a bug-free version that does not componing the performance. We formally specify the new version and mechanically verify the absence of this bug with Keyl's atstace-of-the-art wrification tool for Jawa.



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- → What does it even mean for quicksort to be correct?
- → A number of ways to approach the problem.

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To prove correctness we neet to test all possible input-output pairs (infinite testsuite).

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That is, create a program av that inputs another program p and after finite time outputs false if p has a bug for some input, or true otherwise.

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NO! If we could create av we could solve the halting problem [Alan Turing 1936].

→ Deductive Verification: Can we have a mathematical proof system to prove correctness/incorrectness for all programs?

Create a system L of logical axioms and rules, such that for any program p we can write a **finite proof** that shows either

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If such a system L exists then we can create a fully automatic verification algorithm (simply systematically explore all logical derivations and eventually, in finite time, derive "p has a bug" or "p has no bug".)

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A software verification method that is:

- → Sound: If the verification method returns yes, then indeed the program under examination has no bug
- → Complete: If the verification method returns no, then indeed the program under examination has a bug
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Unfortunately we can only pick TWO. A verification method with all three properties is theoretically impossible.

Usually verification systems pick soundness and termination.

However not all is lost! Sound and terminating systems can prove the correctness of virtually every program we would care about.

The scientific community continuously pushes the limits of these systems.

- → Algorithmic verification: model checking, abstract interpretation, static analysis, type systems.
  - → create a model of the program in a decidable framework (finite state system, pushdown system)
  - → building the model may require some user input
  - → "push-button" verification of correctness

#### → Deductive verification:

- → create a correctness proof of the program in a logic (with axioms and logical rules)
- → constructing the proof likely to require user input
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In this module: we will use the second approach.

- → We will manually prove programs correct using pen and paper proofs.
- → We will use software that provide *some* automation to make this easier.

## Success stories start appearing from the mid-1990s.

→ Paris metro line 14, (1998,refinement approach – combination of algorithmic and deductive approach)

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http://www.methode-b.com/documentation b/ClearSy-Industrial Use of B.pdf
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- → Flight control software of A380: verified absence of run-time errors (2005, abstract interpretation) http://www.astree.ens.fr/
- → L4.verifiedmicro-kernel: a formally correct operating system kernel (2010, deductive verification) http://www.ertos.nicta.com.au/research/t4.verified/
- → SLAM: verifier for MS Windows drivers (2010, model checking) http://research.microsoft.com/en-us/projects/slam/
- → Facebook's Infer verifier: detects bugs in Android and iOS apps http://fbinfer.com/
- → And more...



# CSU44004/CSU55004 THEORY

Book: Logic in Computer Science: Modelling and Reasoning about Systems, 2nd Edition, by M. Huth & M. Ryan

- 1. Symbolic logic
  - → Natural deduction
  - → Propositional logic
  - → First-order Predicate logic
- 2. Correctness of imperative programs
  - → Floyd-Hoare logic
  - → Weakest Precondition calculus
  - → Loop invariants
  - → functional abstractions, etc.



# CSU44004/CSU55004: PRACTICE

## Exercises, Tutorials and Assignments:

- 1. Pencil&paper proofs: propositional, inductive, correctness proofs.
- Semi-automatic proofs using the Dafny verifier
   http://research.microsoft.com/en-us/projects/dafny/ http://rise4fun.com/Dafny

# CSU44004/CSU55004: OTHER

#### Marks:

- → 2% marks from attendance (random sampling)
- → 33% marks from coursework
- → 65% marks from annual exam (before Christmas)
- → Second attempt supplementals: 100% exam

Final exam: pencil&paper proofs of simple propositional properties, and program correctness properties.

You will need to install Visual Studio + Dafny (in Windows) or VStudio Code + Dafny (in Linux/OS X).

Three lectures per week - mix of lecture and in-class tutorial.

More information here:

www.scss.tcd.ie/Vasileios.Koutavas/teaching/cs4004-4504