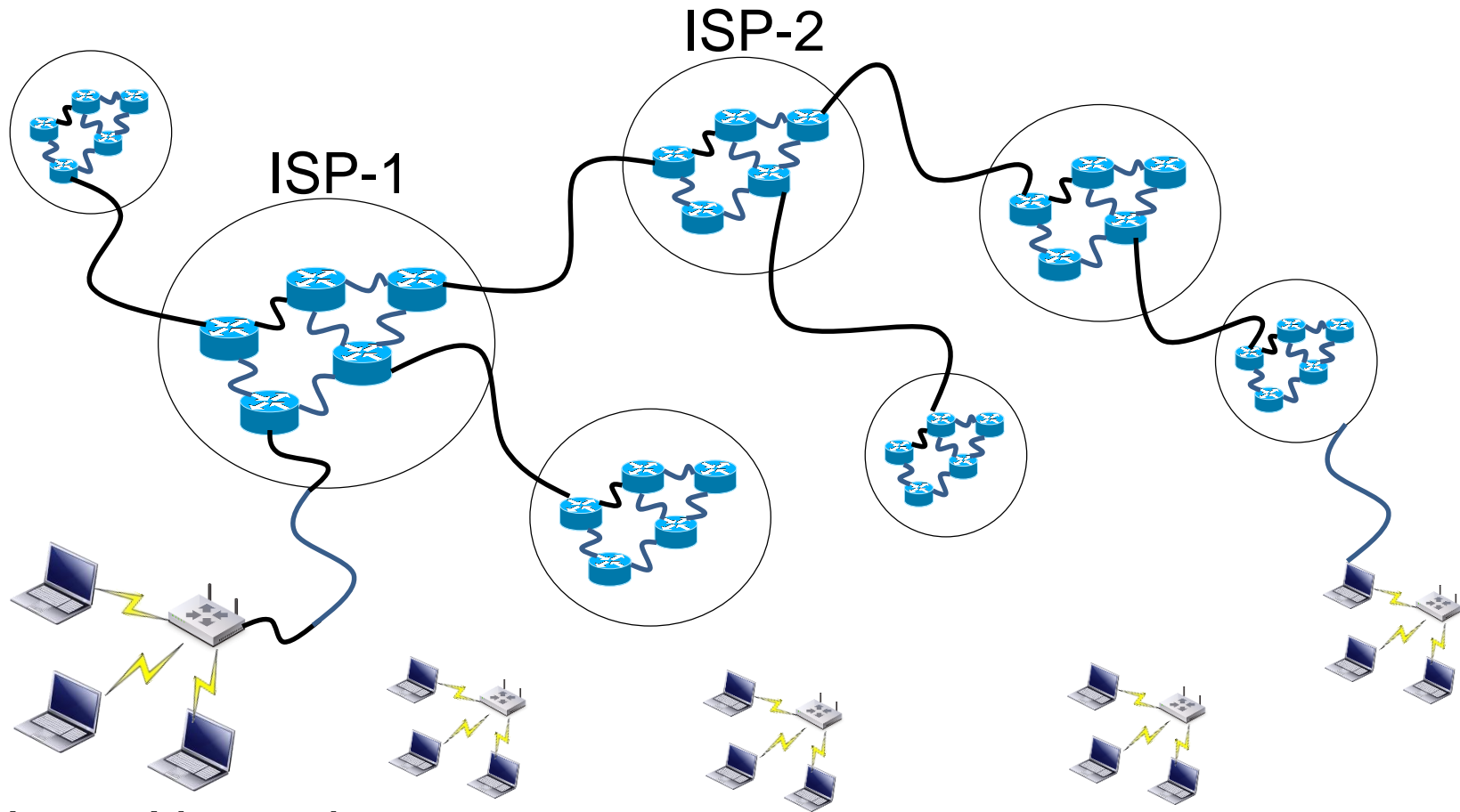


# CS2031

## Telecommunications II

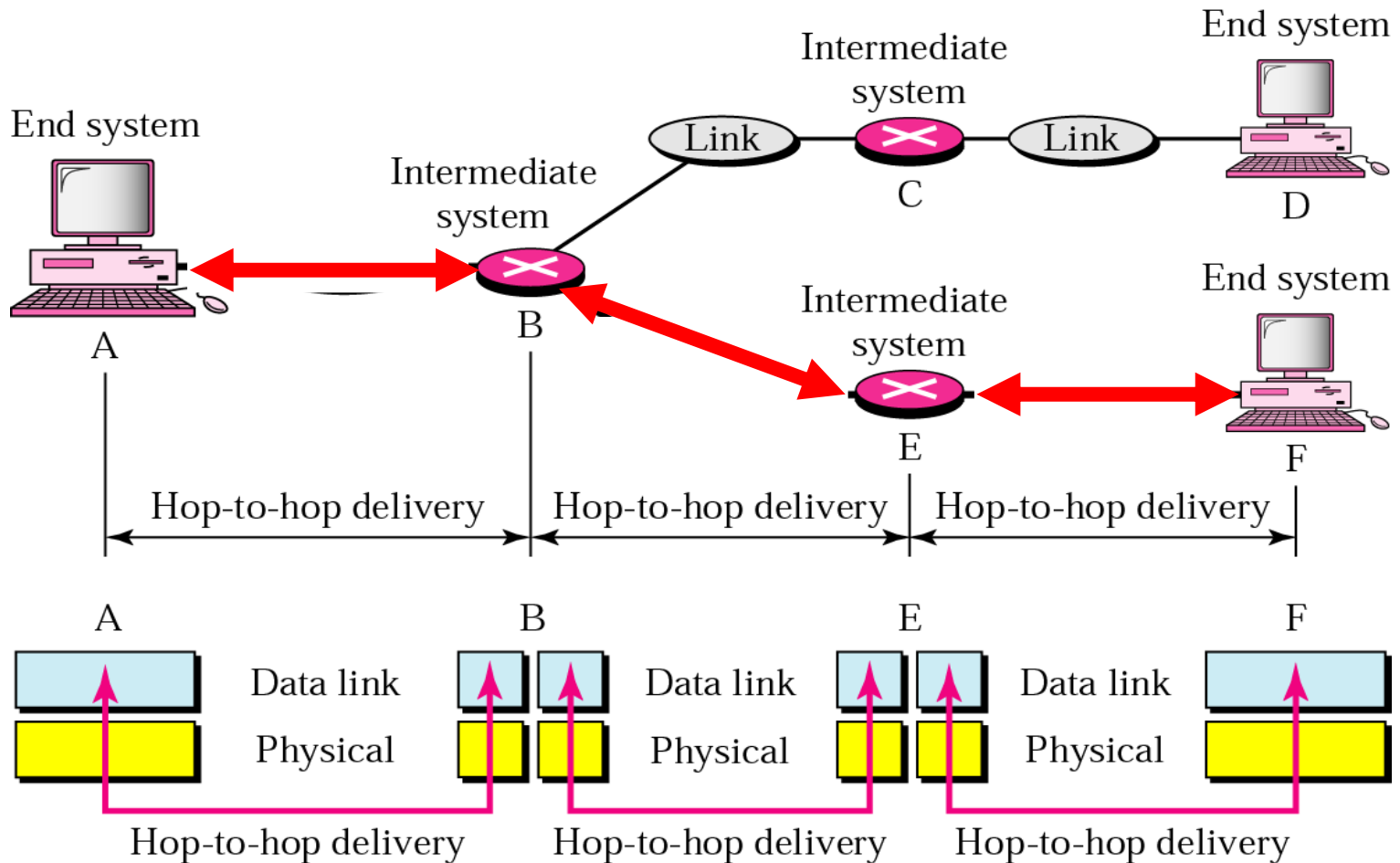
### Error Detection and Correction

# Internet = Network of Networks



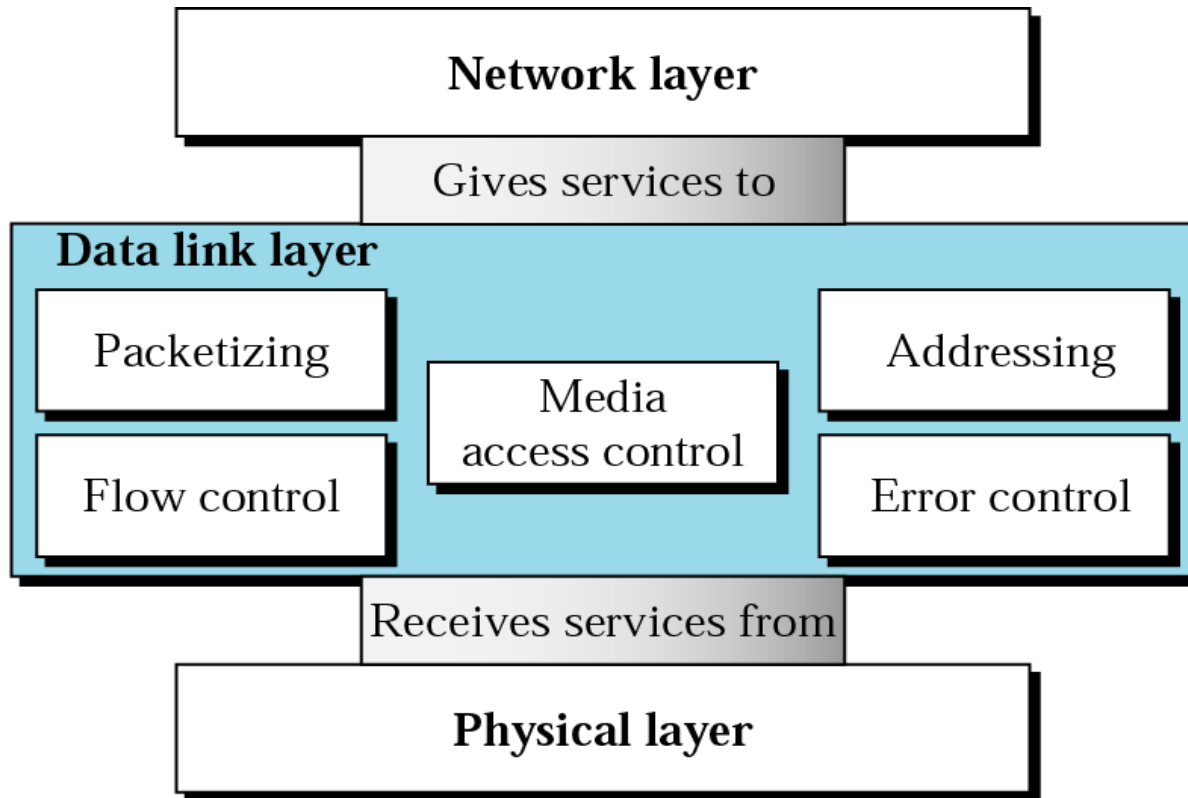
Home Network

# Link Layer



\* Figure is courtesy of B. Forouzan

# Duties of the Link Layer



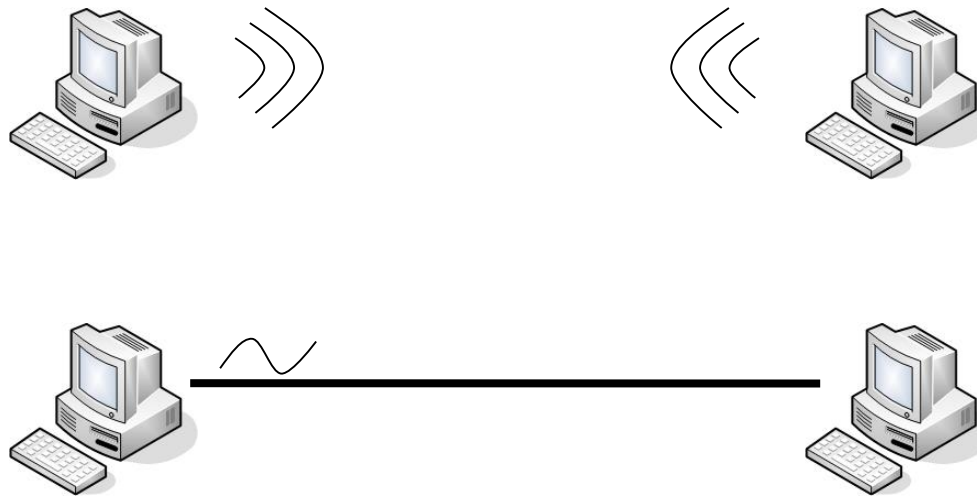
The link layer is responsible for transmitting frames from one station to the next.

\* Figure is courtesy of B. Forouzan

# Errors in Transmissions

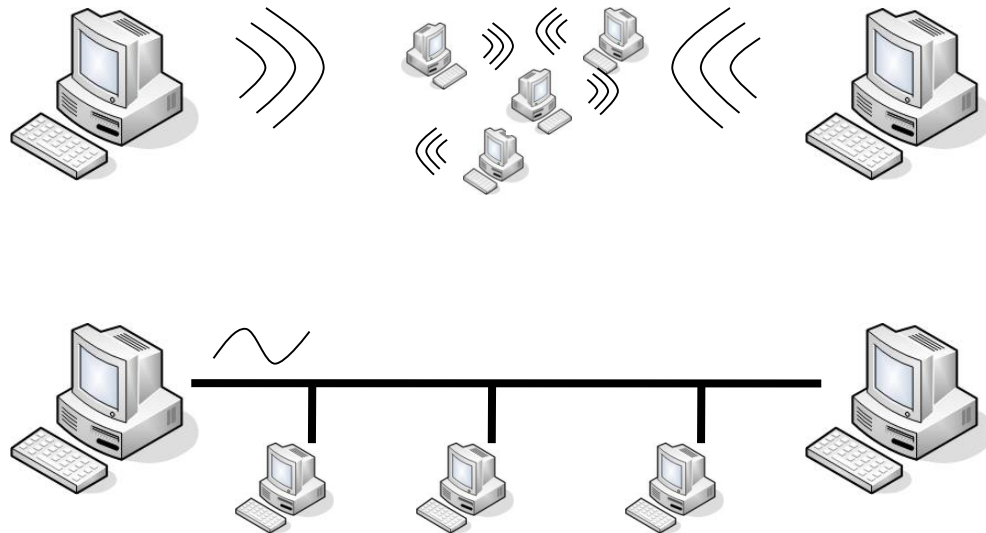
- Causes for Errors
- Types of Errors
- Detection of Errors
- Correction of Errors

# Terminal to Terminal Comms



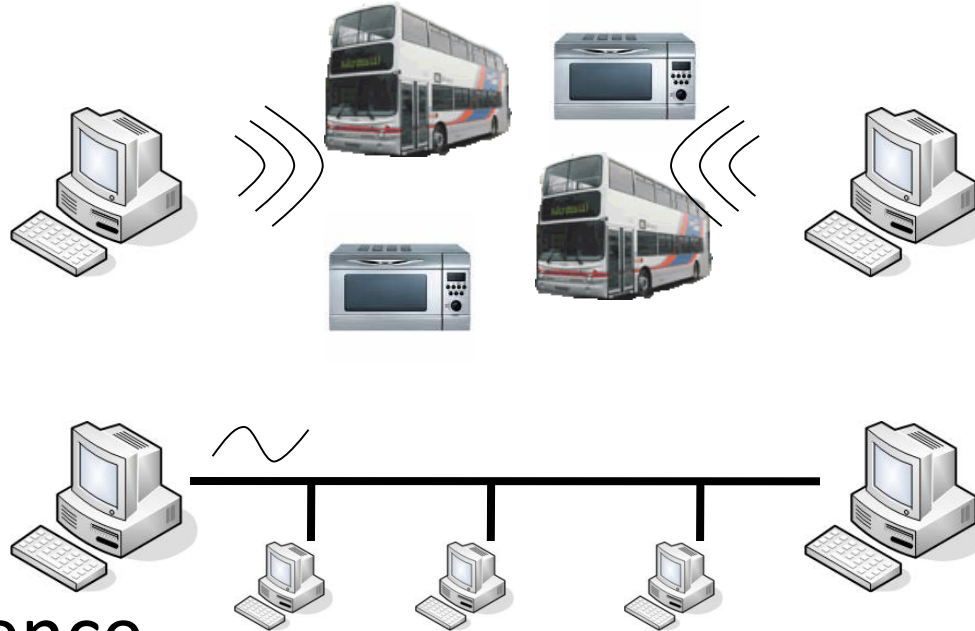
- Either over dedicated or shared medium

# Causes for Errors



- Interference
  - Collision with communication from other nodes
  - Electrical interference from third parties
  - Thermal interference

# Causes for Errors



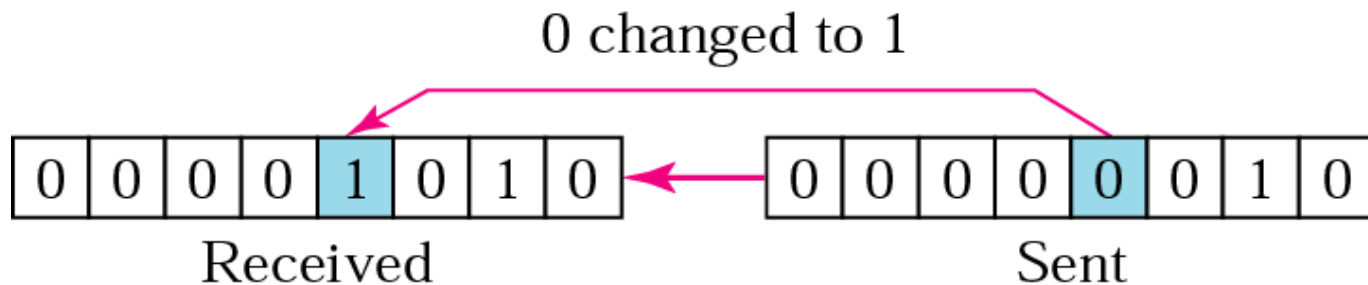
- Interference

- Collision with communication from other nodes
- Electrical interference from third parties
- Thermal interference



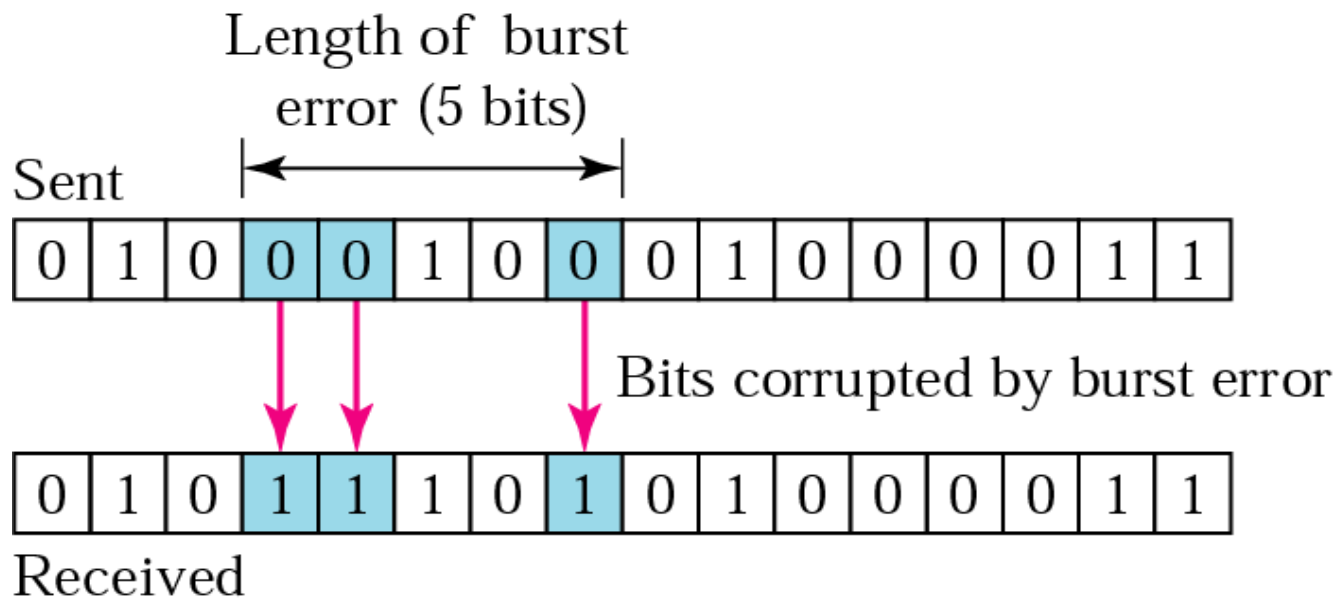
# Types of Errors: Single-Bit Error

In a **single-bit error**, only **one bit** in the data unit has changed.



# Types of Errors: Burst Error

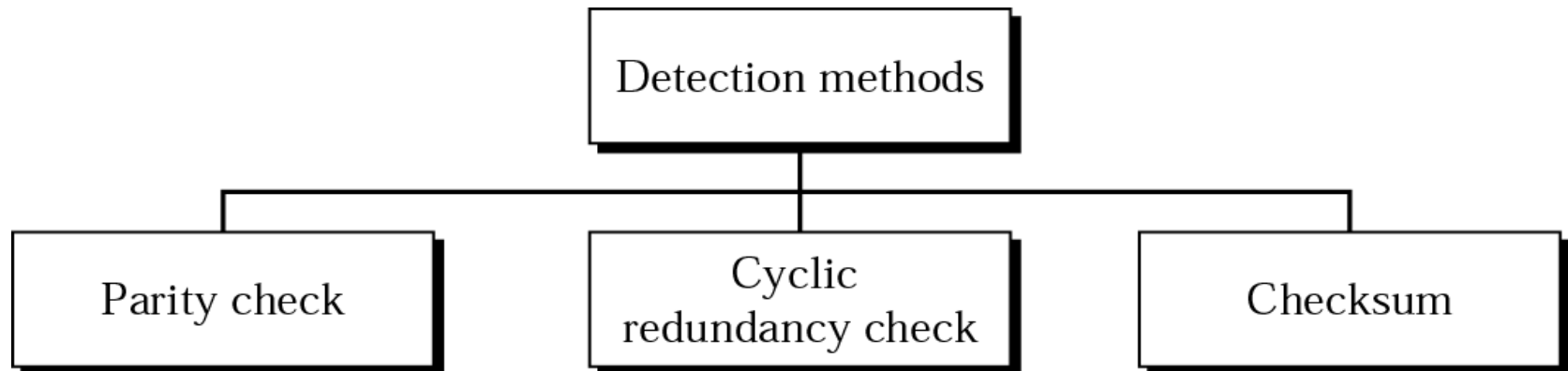
A **burst error** means that **2 or more bits** in the data unit have changed



\* Figure is courtesy of B. Forouzan

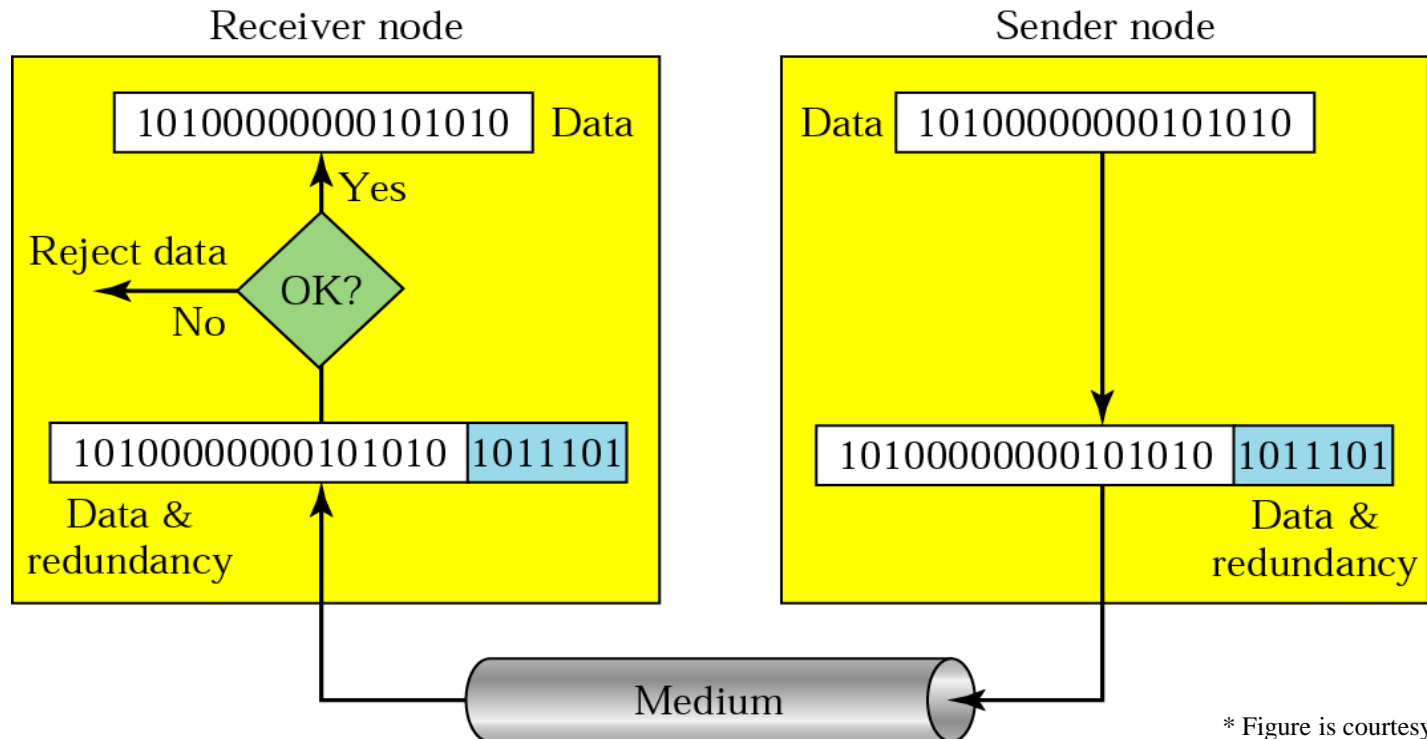
# Detection of Errors

- Redundancy



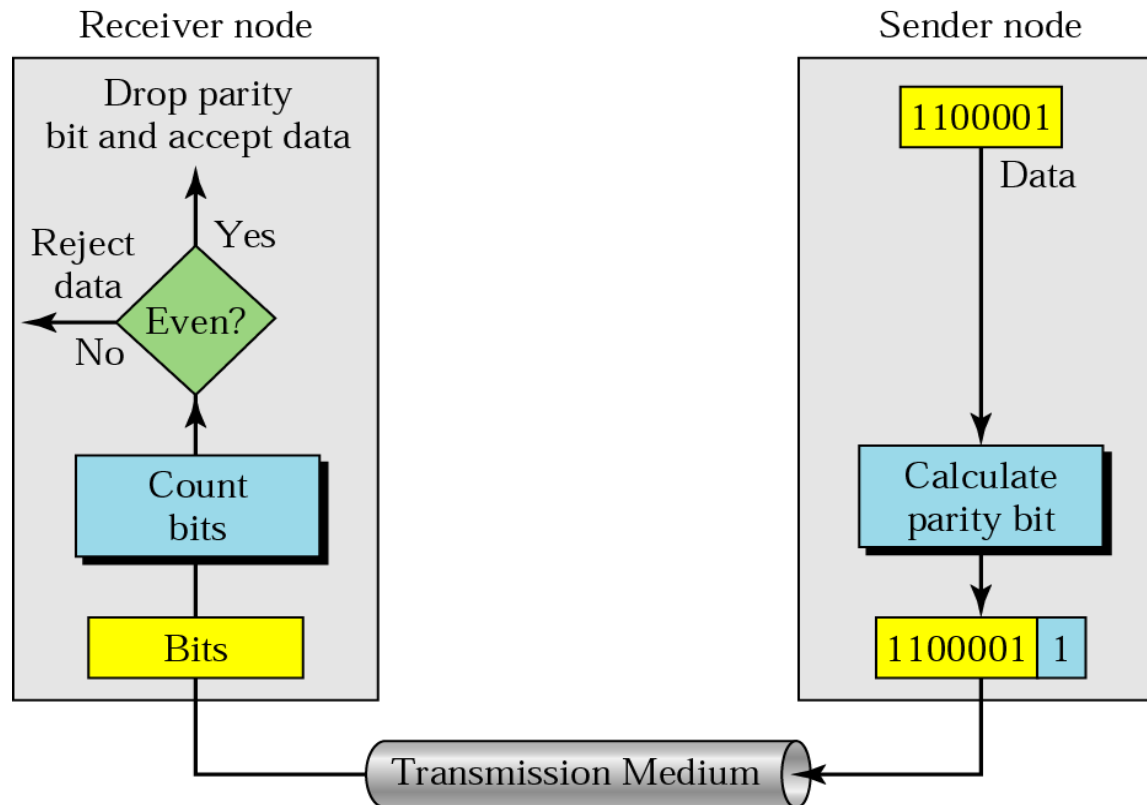
# Redundancy

Error detection uses the concept of redundancy, which means adding extra bits for detecting errors at the destination



\* Figure is courtesy of B. Forouzan

# Even-Parity Concept



A parity bit is added to every data unit so that the total number of 1s is even (or odd for odd-parity).

\* Figure is courtesy of B. Forouzan

## Even-Parity: Example - Sender

- Assume you want to send the following:

1110111    1101111    1110010    1101100    1100100

- The following bits are actually sent:

11101110    11011110    11100100    11011000    11001001

## Even-Parity: Example - Receiver

11101110 11011110 11100100 11011000 11001001

6

6

4

4

4

- The receiver counts the 1s in each character and comes up with even numbers (6, 6, 4, 4, 4). The data are accepted.

11111110 11011110 11101100 11011000 11001001

7

6

5

4

4

- The receiver counts the 1s in each character and comes up with even and odd numbers (7, 6, 5, 4, 4). The receiver knows that the data are corrupted, discards them, and asks for retransmission.

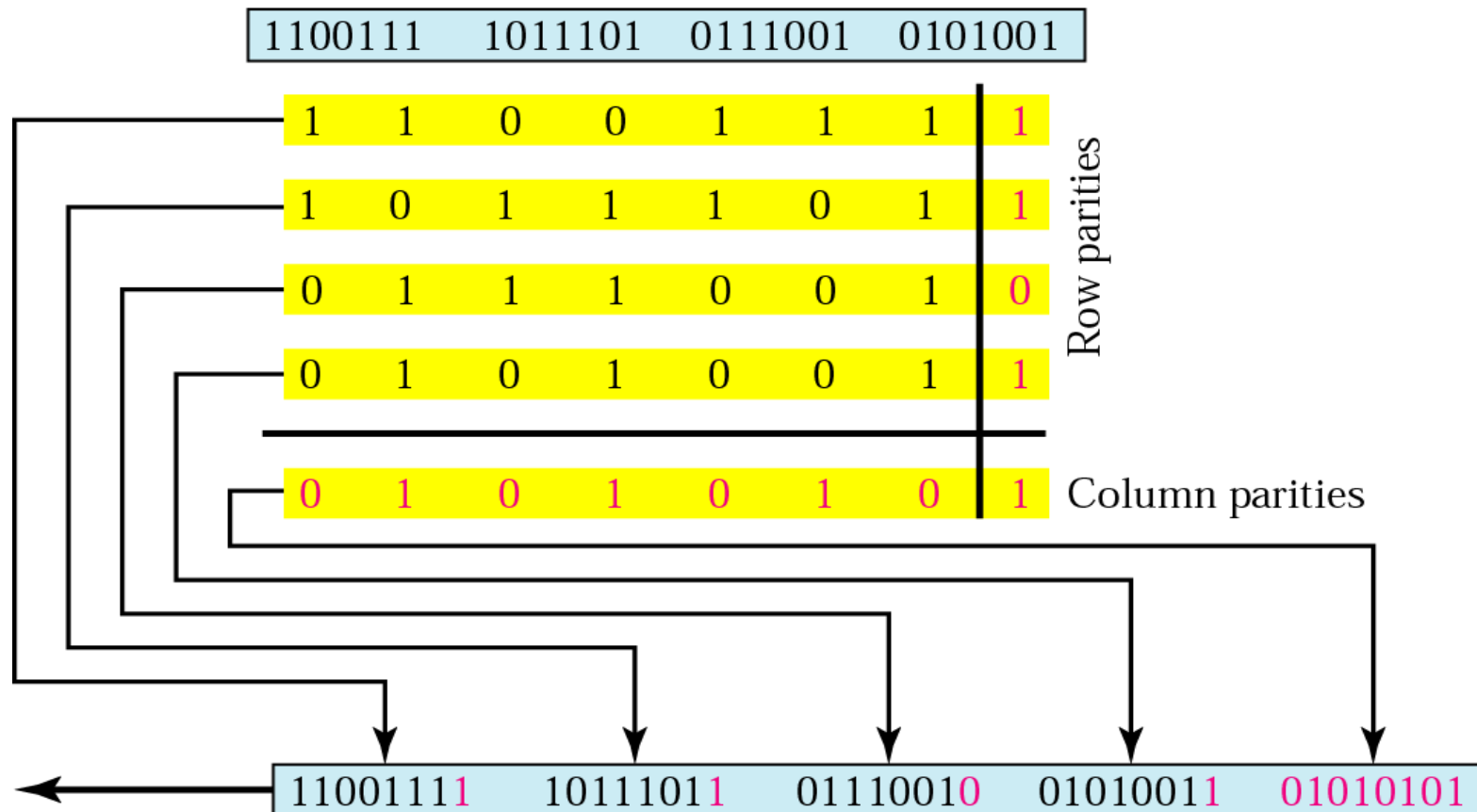
# Simple Parity Check

- Can detect all single-bit errors
- Can detect burst errors only if the total number of errors in each data unit is odd



# Two-Dimensional Parity Check

In two-dimensional parity check, a block of bits is divided into rows and a redundant row of bits is added to the whole block.



\* Figure is courtesy of B. Forouzan

## Example: 2D-Parity Check

Suppose the following block is sent:

10101001 00111001 11011101 11100111 10101010

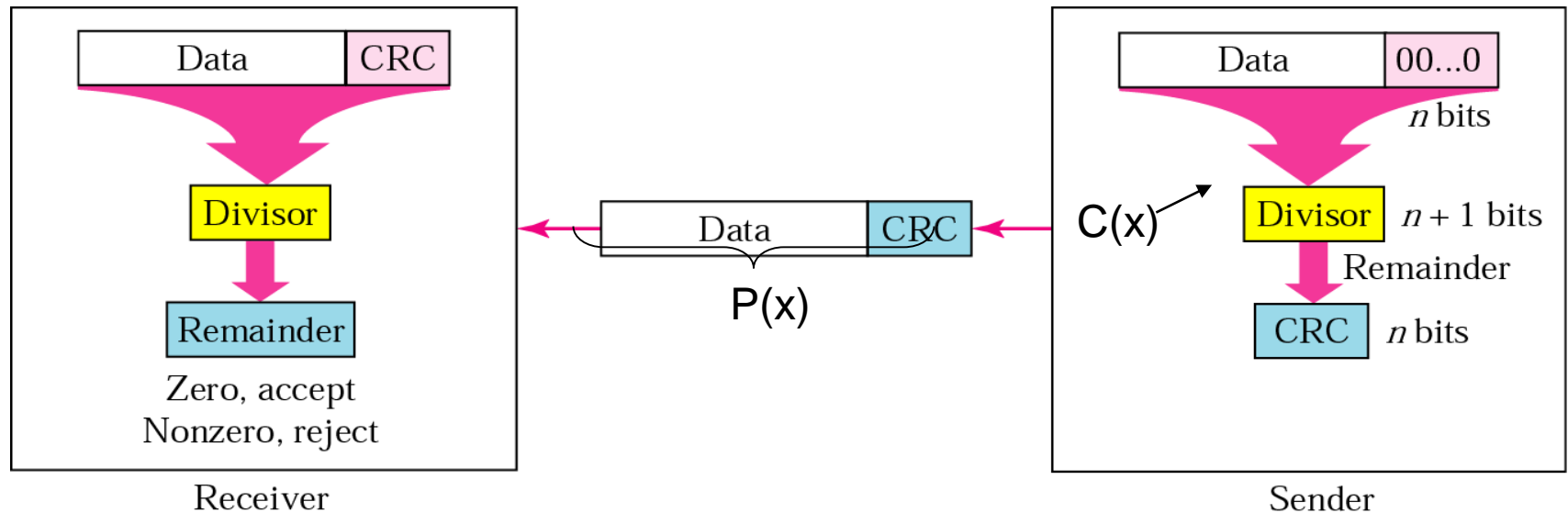
However, it is hit by a burst noise of length 8, and some bits are corrupted.

10100011 10001001 11011101 11100111 10101010

When the receiver checks the parity bits, some of the bits do not follow the even-parity rule and the whole block is discarded.

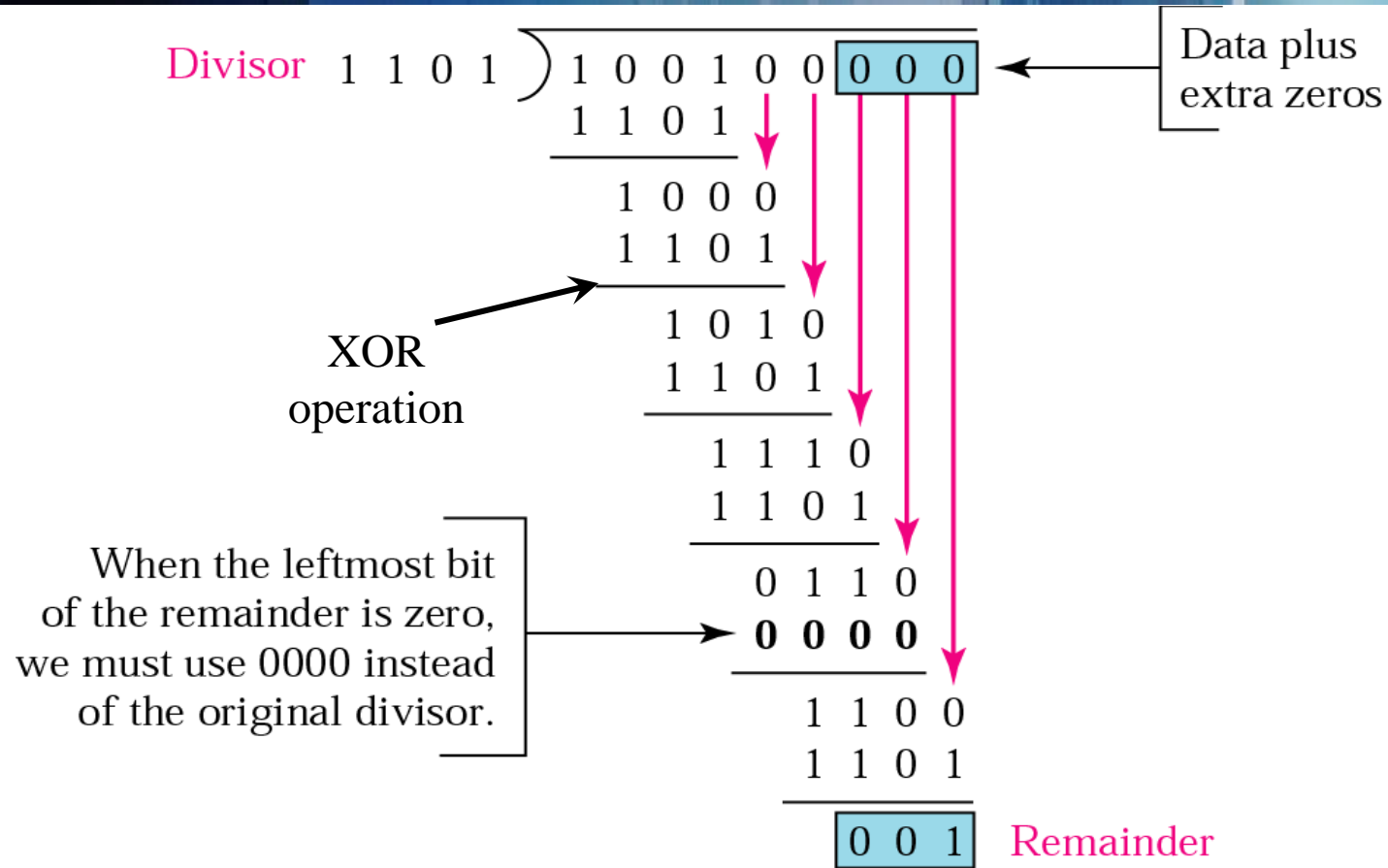
10100011 10001001 11011101 11100111 10101010

# Cyclic Redundancy Check (CRC)



- $P(x)$  divided by  $C(x) = 0$
- $(P(x) + \text{remainder})$  divided by  $C(x)$  should be  $\neq 0$

# CRC: Sender

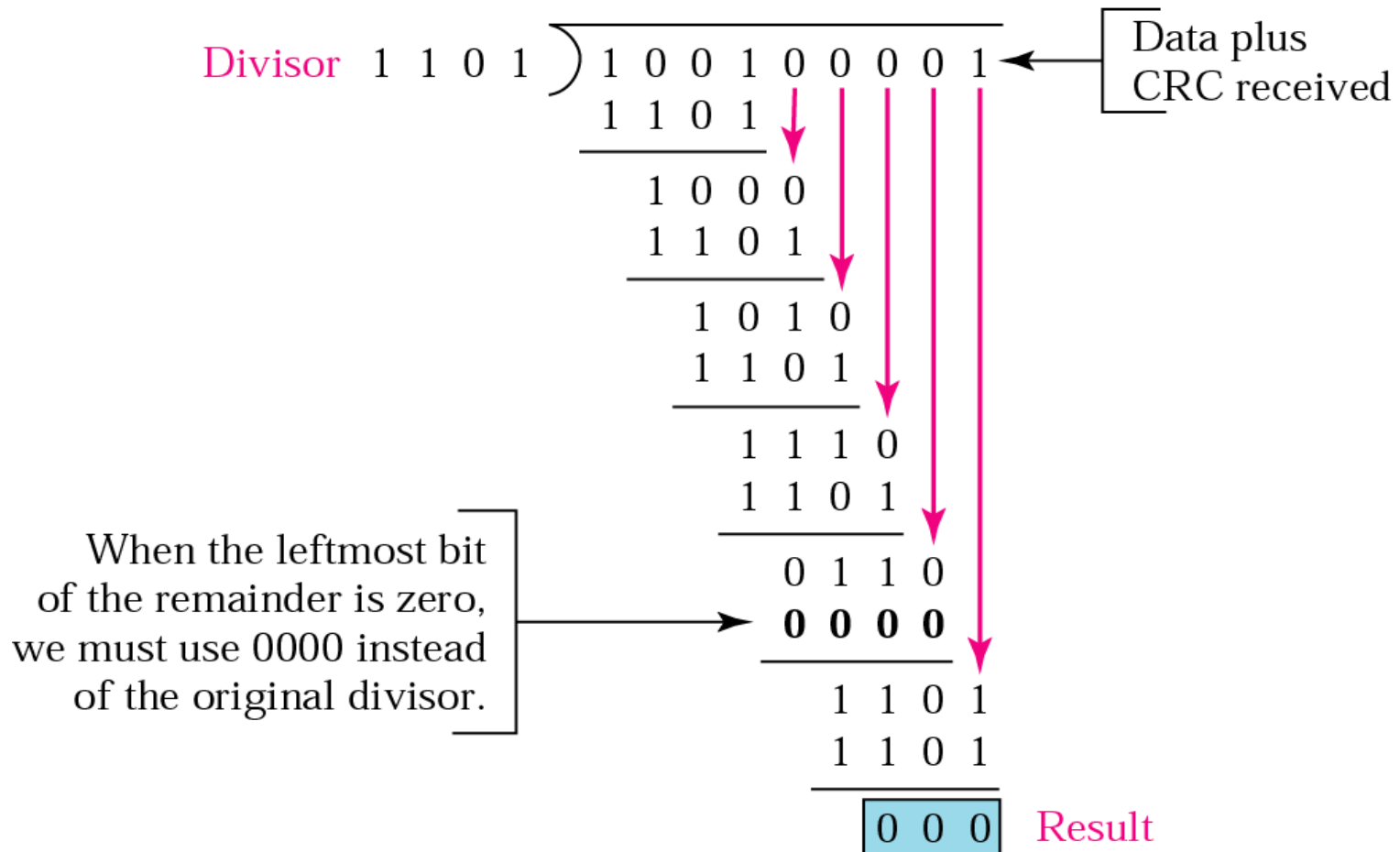


Data transmitted to receiver: 1 0 0 1 0 0 0 0 1

Data CRC

\* Figure is courtesy of B. Forouzan

# CRC: Receiver

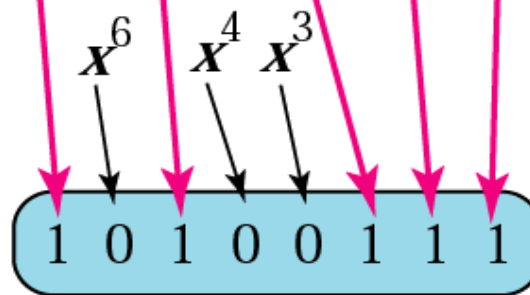


\* Figure is courtesy of B. Forouzan

# Polynomial Notation

Polynomial

$$x^7 + x^5 + x^2 + x + 1$$



Divisor

- Rules for selecting divisor:
  - It should not be divisible by  $x$
  - It should be divisible by  $x+1$

\* Figure is courtesy of B. Forouzan

# Polynomials

- We cannot choose  $x$  (binary 10) or  $x^2 + x$  (binary 110) as polynomial because both are divisible by  $x$ .
- However, we can choose  $x + 1$  (binary 11) because it is not divisible by  $x$ , but is divisible by  $x + 1$ . We can also choose  $x^2 + 1$  (binary 101) because it is divisible by  $x + 1$  (binary division).

# Standard Polynomials

Name	Polynomial	Application
<b>CRC-8</b>	$x^8 + x^2 + x + 1$	<b>ATM header</b>
<b>CRC-10</b>	$x^{10} + x^9 + x^5 + x^4 + x^2 + 1$	<b>ATM AAL</b>
<b>CRC-16</b>	$x^{16} + x^{12} + x^5 + 1$	<b>HDLC</b>
<b>CRC-32</b>	$x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^8 + x^7 + x^5 + x^4 + x^2 + x + 1$	<b>LANs</b>



## CRC Performance

- Can detect all burst errors that effect an odd number of bits
- Can detect all burst errors of the length less than or equal to the degree of the polynomial
- Can detect with a very high probability burst errors of a length greater than the degree of the polynomial.

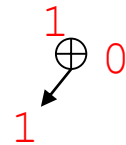
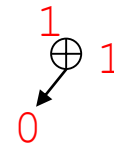
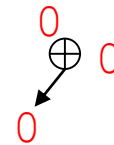
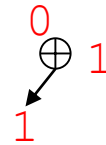
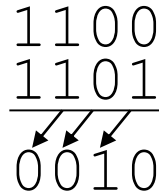
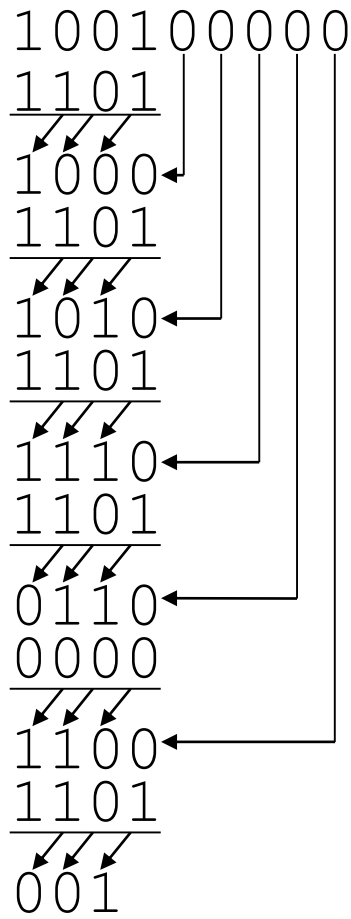
# CRC-12 Performance

## The CRC-12

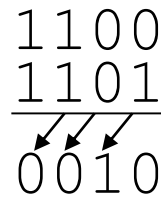
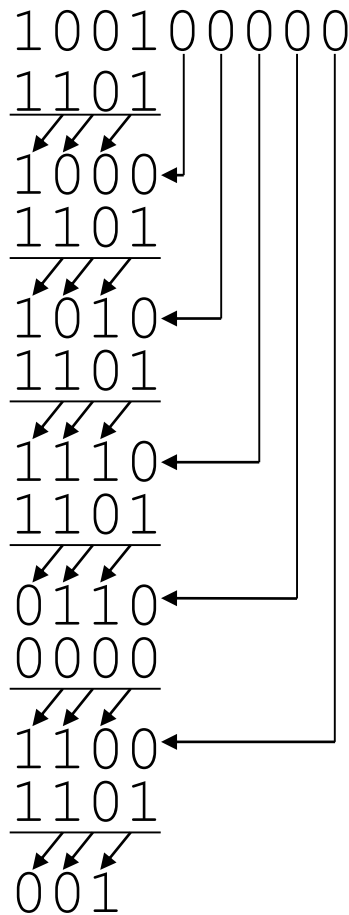
$$x^{12} + x^{11} + x^3 + x + 1$$

which has a degree of 12, will detect all burst errors affecting an odd number of bits, will detect all burst errors with a length less than or equal to 12, and will detect, 99.97 percent of the time, burst errors with a length of 12 or more.

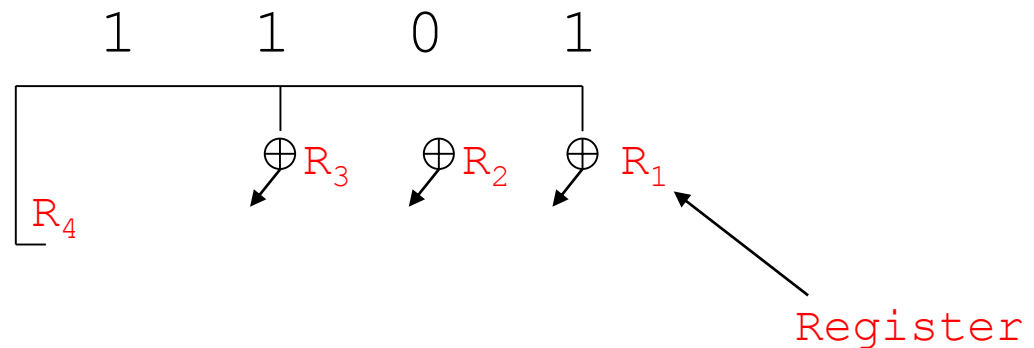
# CRC Calculation



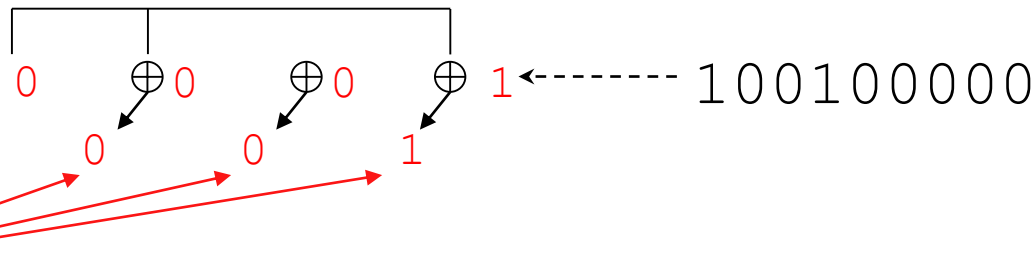
# CRC Calculation



Representation of divisor:

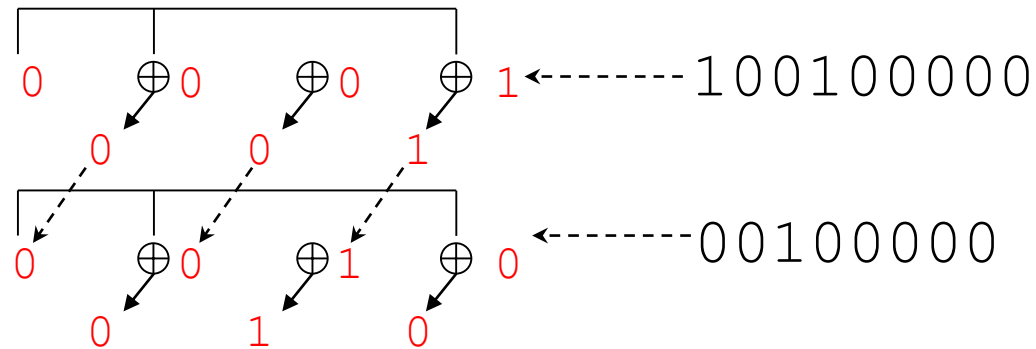


# CRC Calculation

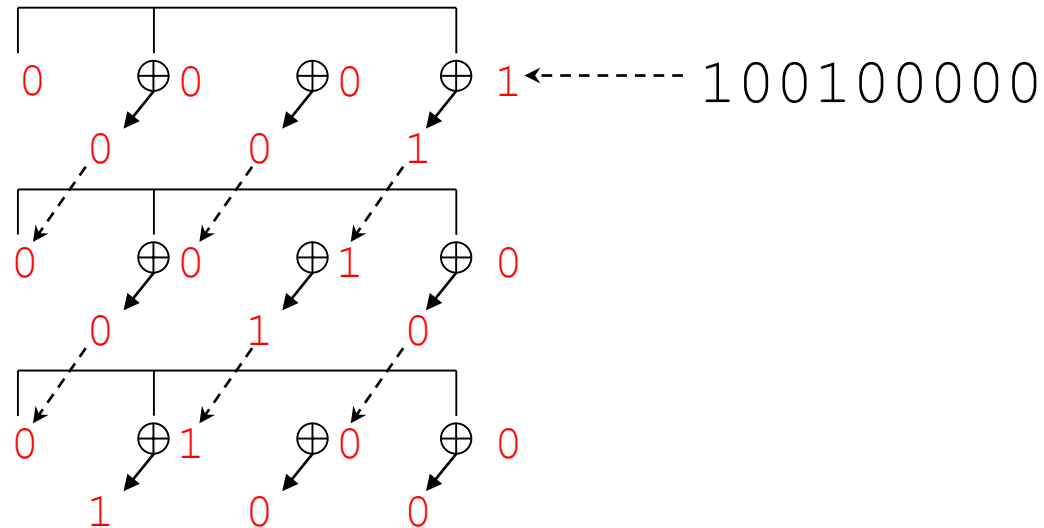


New content  
for registers  
R<sub>4</sub>, R<sub>3</sub>, R<sub>2</sub>

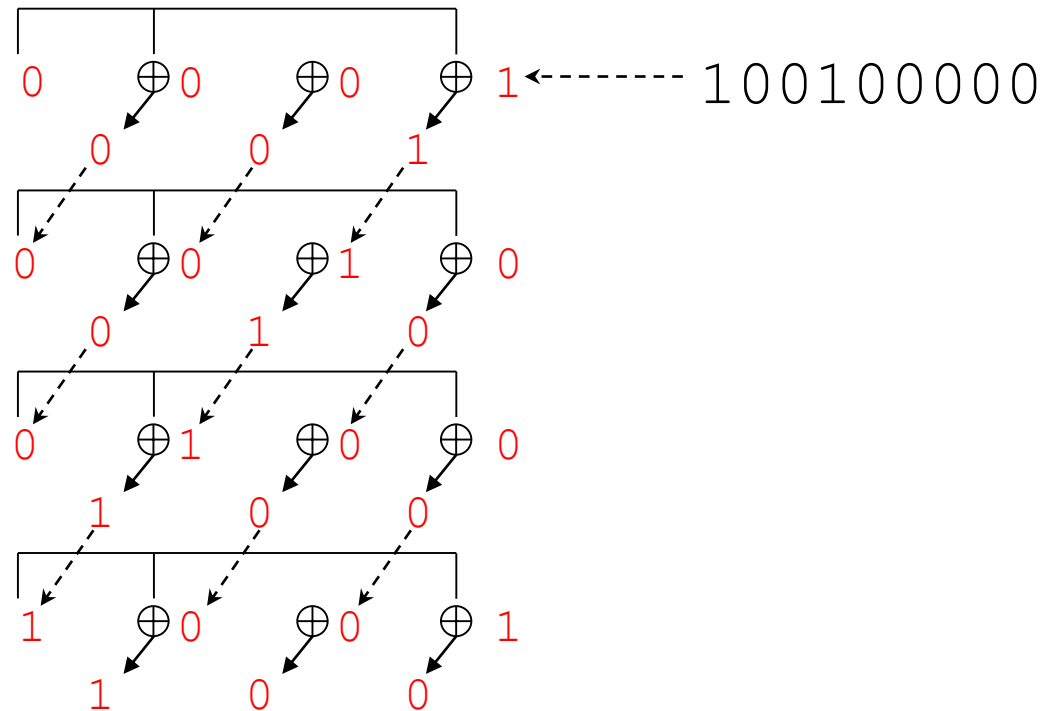
# CRC Calculation



# CRC Calculation

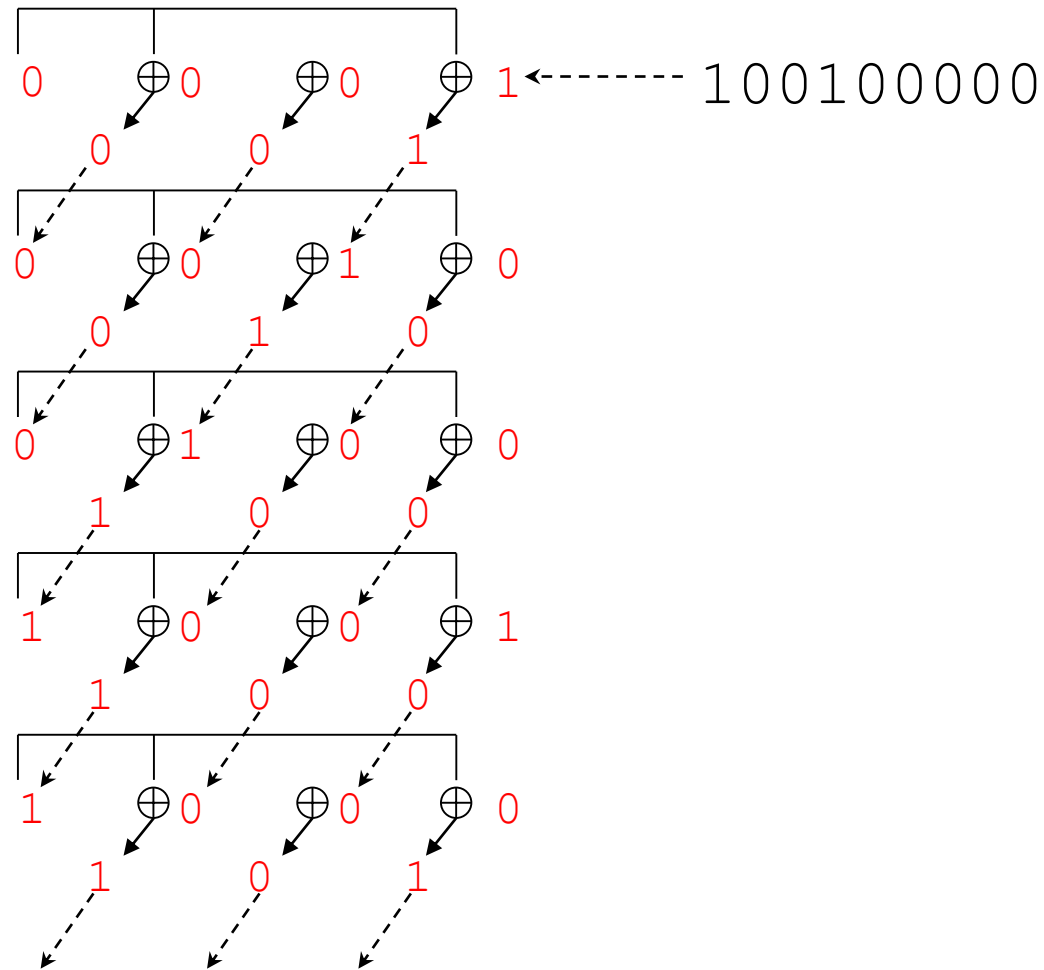


# CRC Calculation

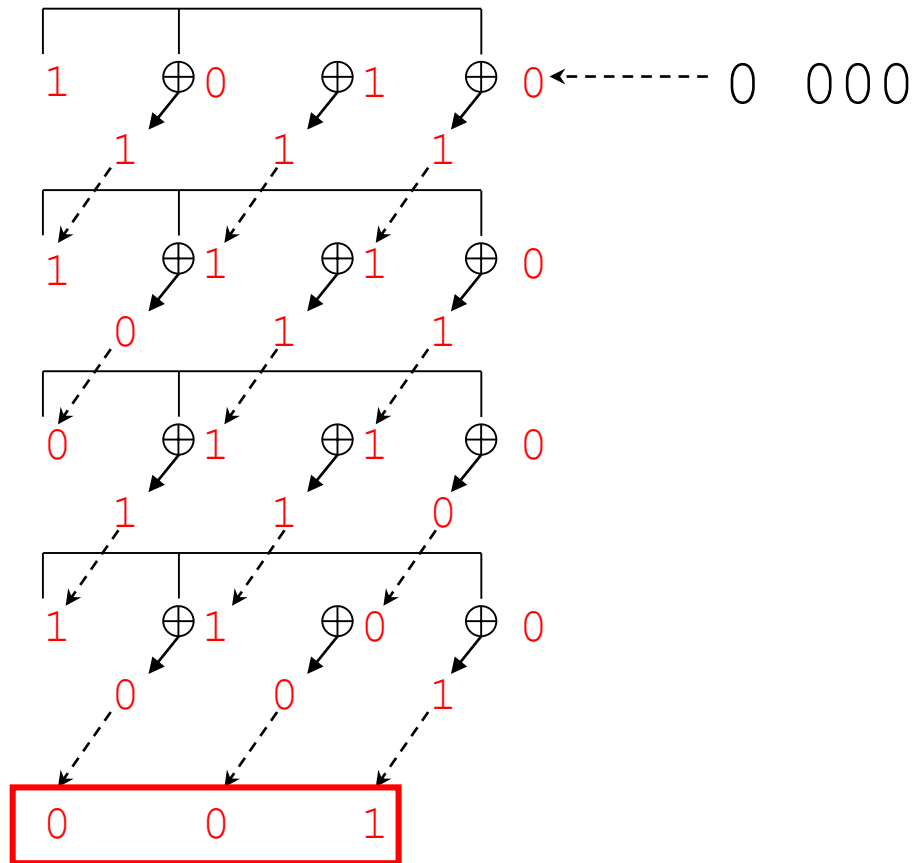




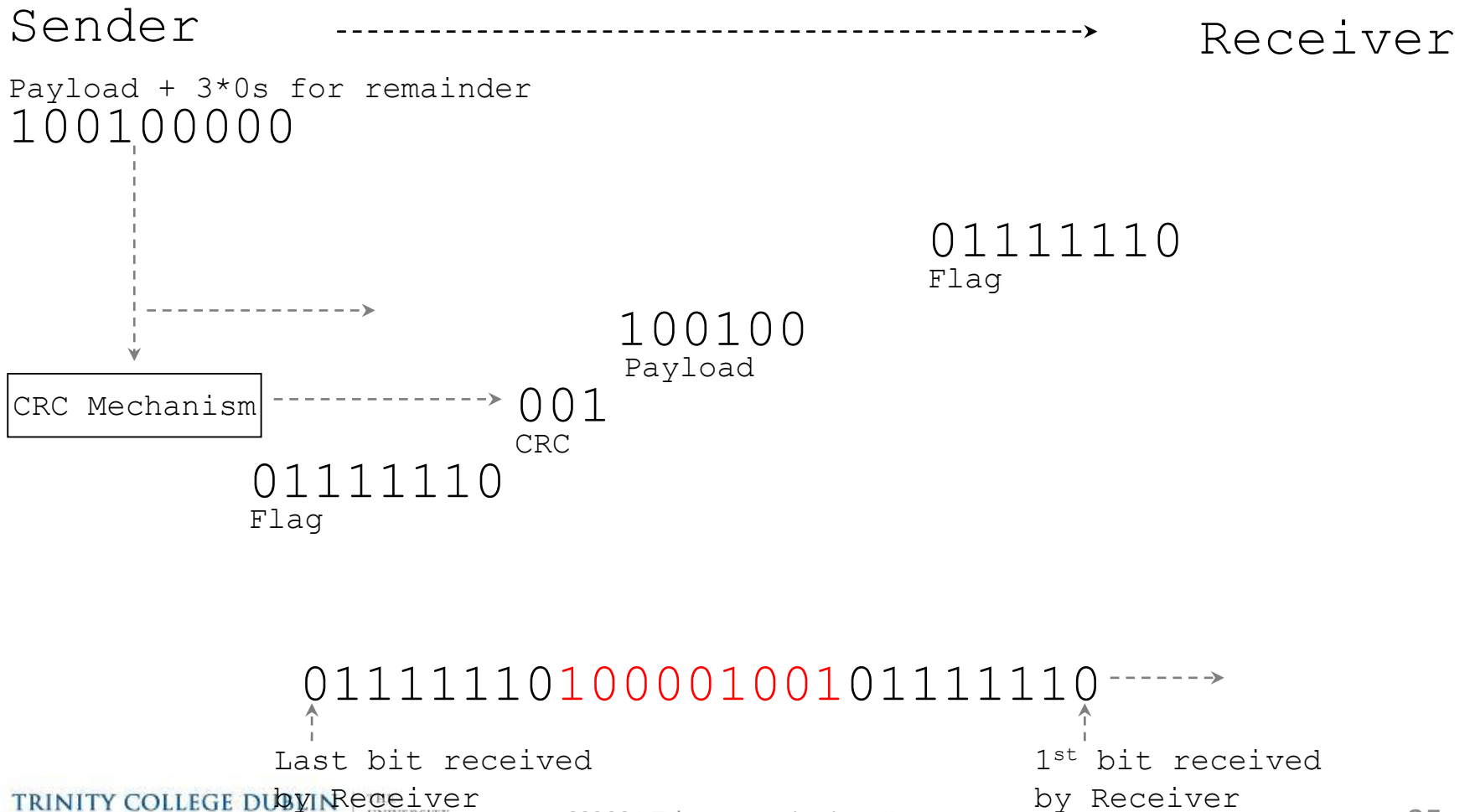
# CRC Calculation



# CRC Calculation



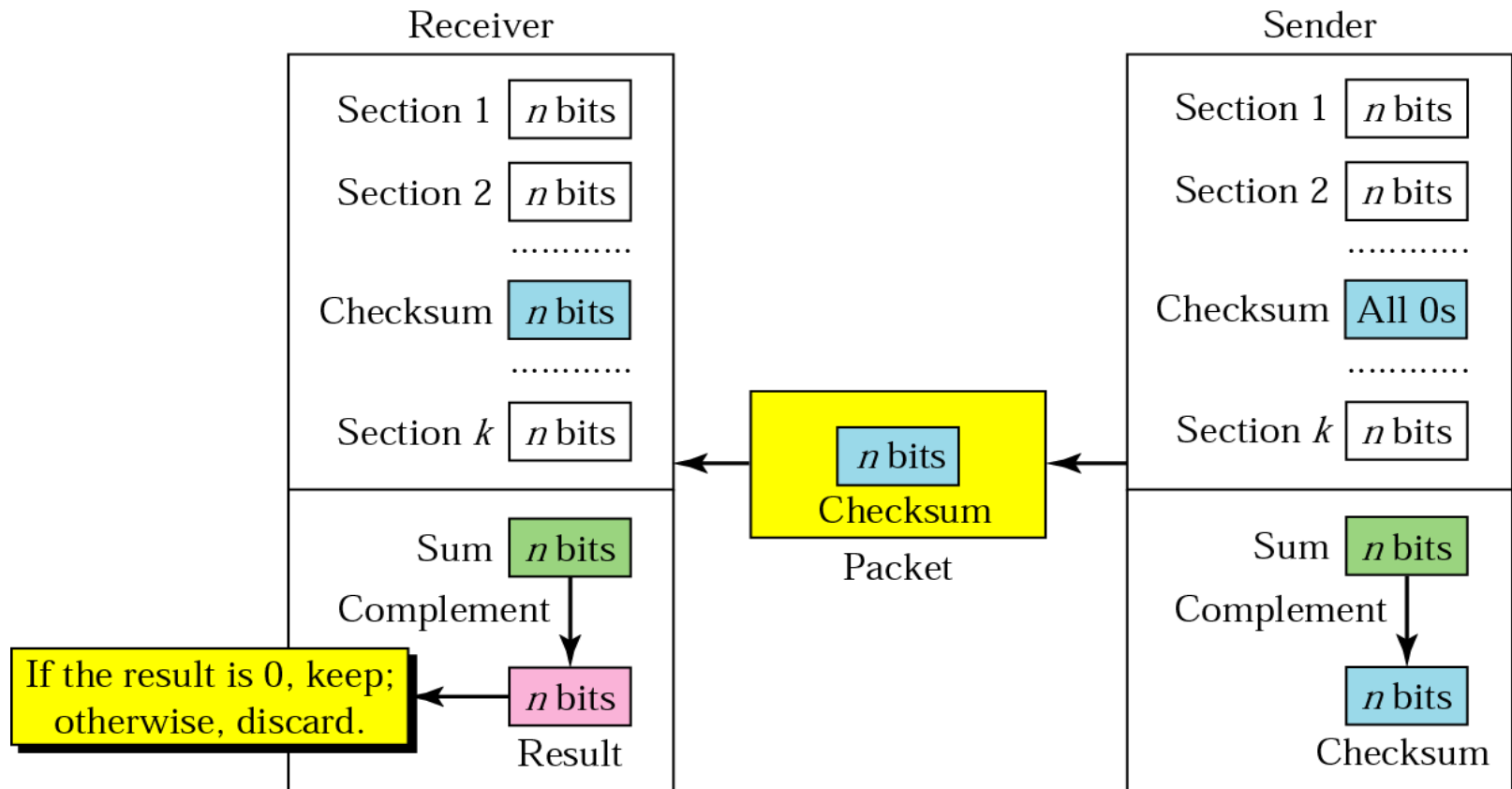
# CRC Calculation



# Example from a Linux box

```
wlan0      Link encap:Ethernet  HWaddr 00:0b:81:89:56:ca  
           inet addr:192.168.192.12  Bcast:192.168.192.255  Mask:255.255.255.0  
           inet6 addr: fe80::20b:81ff:fe89:56ca/64 Scope:Link  
           UP BROADCAST RUNNING MULTICAST  MTU:1500  Metric:1  
           RX packets:292 errors:0 dropped:374 overruns:0 frame:0  
           TX packets:199 errors:0 dropped:2 overruns:0 carrier:0  
           collisions:0 txqueuelen:1000  
           RX bytes:47787 (46.6 KiB)  TX bytes:26749 (26.1 KiB)
```

# Checksum



\* Figure is courtesy of B. Forouzan

# Checksum II

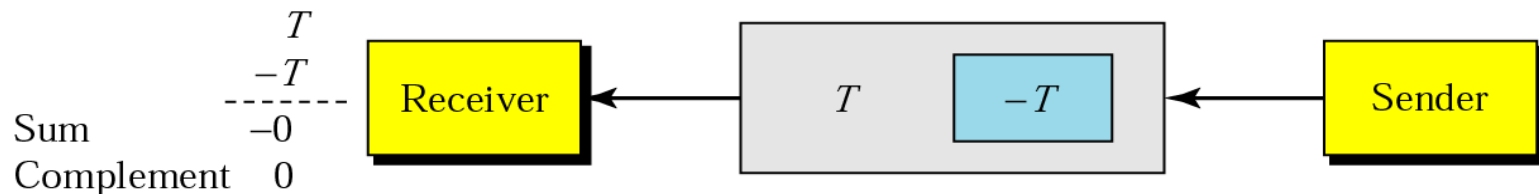
## Sender:

The unit is divided into  $k$  sections, each of  $n$  bits.

All sections are added using one's complement to get the sum.

The sum is complemented and becomes the checksum.

The checksum is sent with the data.



## Receiver:

The unit is divided into  $k$  sections, each of  $n$  bits.

All sections are added using one's complement to get the sum.

The sum is complemented.

If the result is zero, the data are accepted: otherwise, rejected.

\* Figure is courtesy of B. Forouzan

## Example: Checksum

Sender:

10101001

00111001

-----

Sum            11100010

Checksum      **00011101**

The data that is send:

10101001   00111001   **00011101**

## Example: Checksum

Sender:

10101001

00111001

-----

Sum 11100010

Checksum **00011101**

The data that is send:

10101001 00111001 **00011101**

Receiver:

10101001

00111001

00011101

Sum 11111111

Complement **00000000**

Complement: **00000000**

means that the frame is OK.



# Example: Checksum

Sender:

10101001

00111001

-----

Sum

11100010

Checksum

**00011101**

Receiver:

10101111

11111001

00011101

Partial Sum

**1** 11000101

Carry

**1**

Sum

11000110

Complement

**00111001**

The data that is send:

10101001 00111001 **00011101**

Complement: **00111001**

means that the frame is corrupted.

# Sample TCP / IP Packet

0000	00 07 e9 7c 22 fc 00 11 93 85 e0 c4 08 00 45 00
0010	00 2c db 26 40 00 3f 06 0e 77 86 e2 20 37 86 e2
0020	24 33 01 bd 12 3f 3d fa 0f b6 a8 6f 87 c0 50 18
0030	bc 40 8a 7c 00 00 85 00 00 00 00 00

...|".....E.  
 .,.&@.?..w.. 7..  
 \$3...?=....o..P.  
 .@.|.....

## Ethernet Header:

src addr: 00 07 e9 7c 22 fc  
 dest addr: 00 11 93 85 e0 c4

## IP Header:

src addr: 134.226.36.55  
 dest addr: 134.226.36.51

## TCP Header:

src port: 445  
 dest port: 4671

## NetBios Information

## IP Header Checksum

The IP header is generally 20 byte and can be divided into units of 2 bytes/16 bits to calculate the checksum

# Summary: Detection of Errors

- Parity Check
- Cyclic Redundancy Check (CRC)
- Checksum

# Correction of Errors

- Error Correction through Retransmission
  - Parity, CRC, Checksum determine validity
  - If not valid, discard and wait for sender to retransmit
- Forward Error Correction
  - Determine the corrupted bit or bits at the receiver

# Hamming Code

11	10	9	8	7	6	5	4	3	2	1
d	d	d	$r_8$	d	d	d	$r_4$	d	$r_2$	$r_1$

- Redundancy bits distributed throughout data bits
- Individual redundancy bits work as parity bits for specific data bits
  - e.g.  $r_1$  is the parity bit for all odd numbers
 

3 = binary 001 $r_1$

7 = binary 011 $r_1$

5 = binary 010 $r_1$

9 = binary 100 $r_1$

\* Figure is courtesy of B. Forouzan

# Redundancy Bits Calculation

$r_1$  will take care of these bits.

11		9		7		5		3		1
d	d	d	$r_8$	d	d	d	$r_4$	d	$r_2$	$r_1$

$r_2$  will take care of these bits.

11	10			7	6			3	2	
d	d	d	$r_8$	d	d	d	$r_4$	d	$r_2$	$r_1$

$r_4$  will take care of these bits.

				7	6	5	4			
d	d	d	$r_8$	d	d	d	$r_4$	d	$r_2$	$r_1$

$r_8$  will take care of these bits.

11	10	9	8							
d	d	d	$r_8$	d	d	d	$r_4$	d	$r_2$	$r_1$

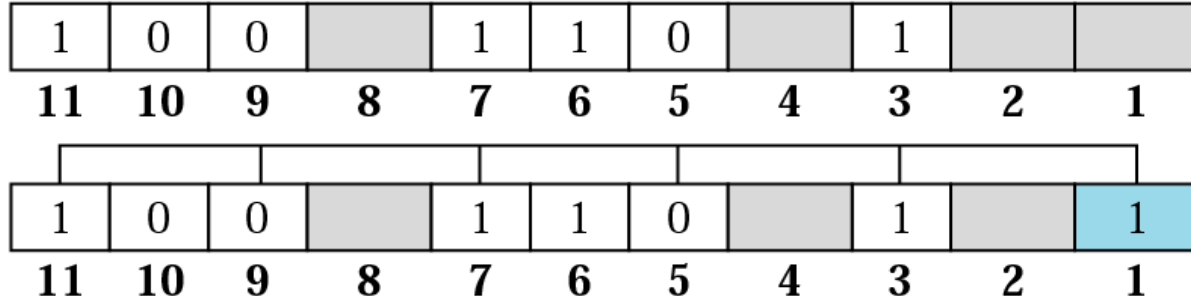
\* Figure is courtesy of B. Forouzan

# Redundancy Bit Calculation

1	0	0		1	1	0		1		
11	10	9	8	7	6	5	4	3	2	1

**Data:**  
**1 0 0 1 1 0 1**

# Redundancy Bit Calculation



**Data:**  
**1 0 0 1 1 0 1**

\* Figure is courtesy of B. Forouzan



# Redundancy Bit Calculation

**Data:**  
**1 0 0 1 1 0 1**

1	0	0		1	1	0		1		
11	10	9	8	7	6	5	4	3	2	1

Adding  $r_1$

1	0	0		1	1	0		1		1
11	10	9	8	7	6	5	4	3	2	1

Adding  $r_2$

1	0	0		1	1	0		1	0	1
11	10	9	8	7	6	5	4	3	2	1

# Redundancy Bit Calculation

**Data:**  
**1 0 0 1 1 0 1**

1	0	0		1	1	0		1		
11	10	9	8	7	6	5	4	3	2	1

Adding  $r_1$

1	0	0		1	1	0		1		1
11	10	9	8	7	6	5	4	3	2	1

Adding  $r_2$

1	0	0		1	1	0		1	0	1
11	10	9	8	7	6	5	4	3	2	1

Adding  $r_4$

1	0	0		1	1	0	0	1	0	1
11	10	9	8	7	6	5	4	3	2	1

# Redundancy Bit Calculation

**Data:**  
**1 0 0 1 1 0 1**

1	0	0		1	1	0		1		
11	10	9	8	7	6	5	4	3	2	1

Adding  $r_1$

1	0	0		1	1	0		1		1
11	10	9	8	7	6	5	4	3	2	1

Adding  $r_2$

1	0	0		1	1	0		1	0	1
11	10	9	8	7	6	5	4	3	2	1

Adding  $r_4$

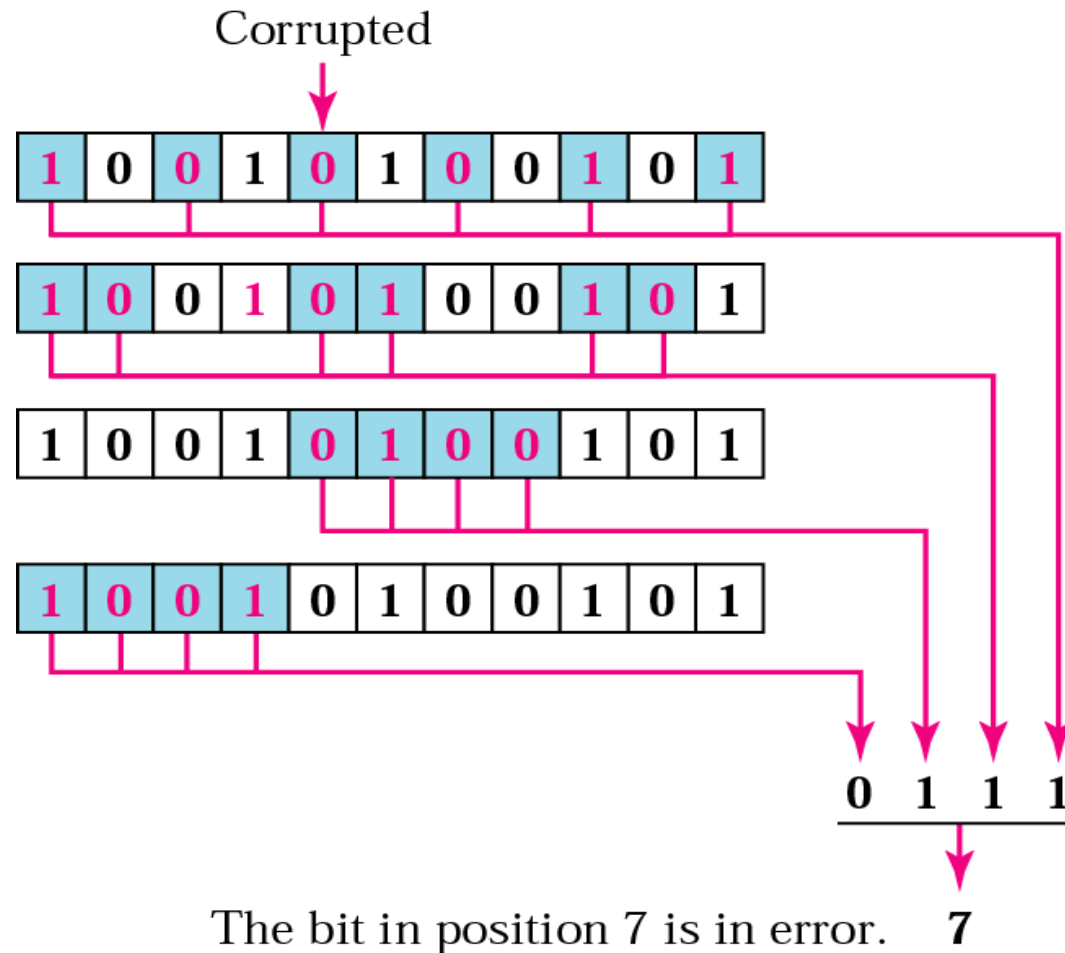
1	0	0		1	1	0	0	1	0	1
11	10	9	8	7	6	5	4	3	2	1

Adding  $r_8$

1	0	0	1	1	1	0	0	1	0	1
11	10	9	8	7	6	5	4	3	2	1

**Code:**  
**1 0 0 1 1 1 0 0 1 0 1**

# Error Detection using Hamming Code



\* Figure is courtesy of B. Forouzan

# Data and Redundancy Bits

Number of data bits $m$	Number of redundancy bits $r$	Total bits $m + r$
1	2	3
2	3	5
3	3	6
4	3	7
5	4	9
6	4	10
7	4	11

\* Figure is courtesy of B. Forouzan

# Summary

- Types of Errors
  - Single-Bit & Burst Errors
- Detection of Errors
  - Parity Check / 2D Parity Check
  - CRC
  - Checksum
- Correction of Errors
  - Error Correction by Retransmission
  - Forward Error Correction – Hamming Code



That's all  
folks