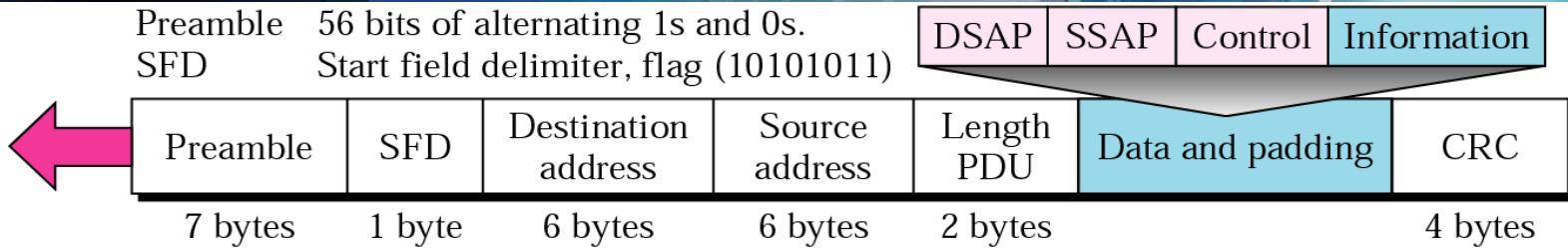


# Overview

- Switches / Ethernet Switches
- Network Layer
- Addressing, IPv4 Addresses & Routers
- CIDR
- NAT
- ARP



# 802.3 MAC Format



- 64-bit frame preamble (10101010) used to synchronize reception
  - 7 bit preamble (10101010) + 1 start flag (10101011)
- Maximum frame length: 1536 bytes
  - ⇒ max 1500 bytes payload
- Minimum frame length: 64 bytes
  - ⇒ min 46 bytes payload

\* Figure is courtesy of B. Forouzan

# Ethernet Standards

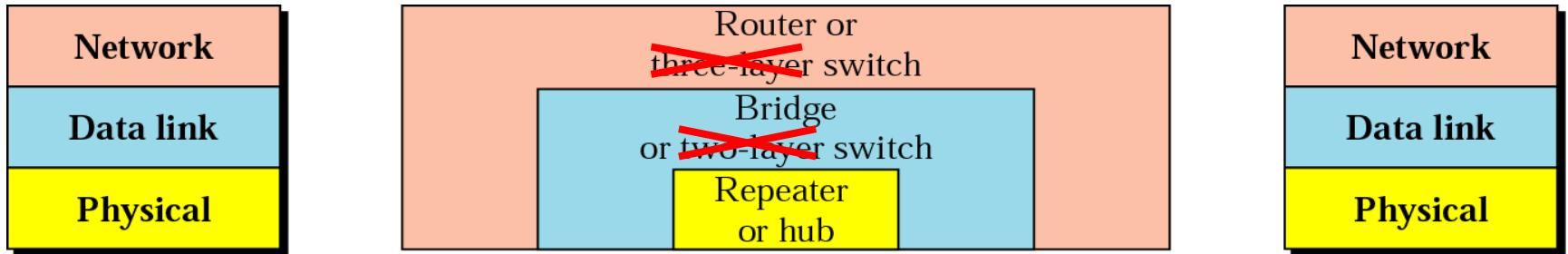
Name	Cable	Max. seg.	Nodes/seg.	Advantages
10Base5	Thick coax	500 m	100	Original cable; now obsolete
10Base2	Thin coax	185 m	30	No hub needed
10Base-T	Twisted pair	100 m	1024	Cheapest system
10Base-F	Fiber optics	2000 m	1024	Best between buildings

Name	Cable	Max. segment	Advantages
100Base-T4	Twisted pair	100 m	Uses category 3 UTP
100Base-TX	Twisted pair	100 m	Full duplex at 100 Mbps
100Base-FX	Fiber optics	2000 m	Full duplex at 100 Mbps; long runs

Name	Cable	Max. segment	Advantages
1000Base-SX	Fiber optics	550 m	Multimode fiber (50, 62.5 microns)
1000Base-LX	Fiber optics	5000 m	Single (10 $\mu$ ) or multimode (50, 62.5 $\mu$ )
1000Base-CX	2 Pairs of STP	25 m	Shielded twisted pair
1000Base-T	4 Pairs of UTP	100 m	Standard category 5 UTP



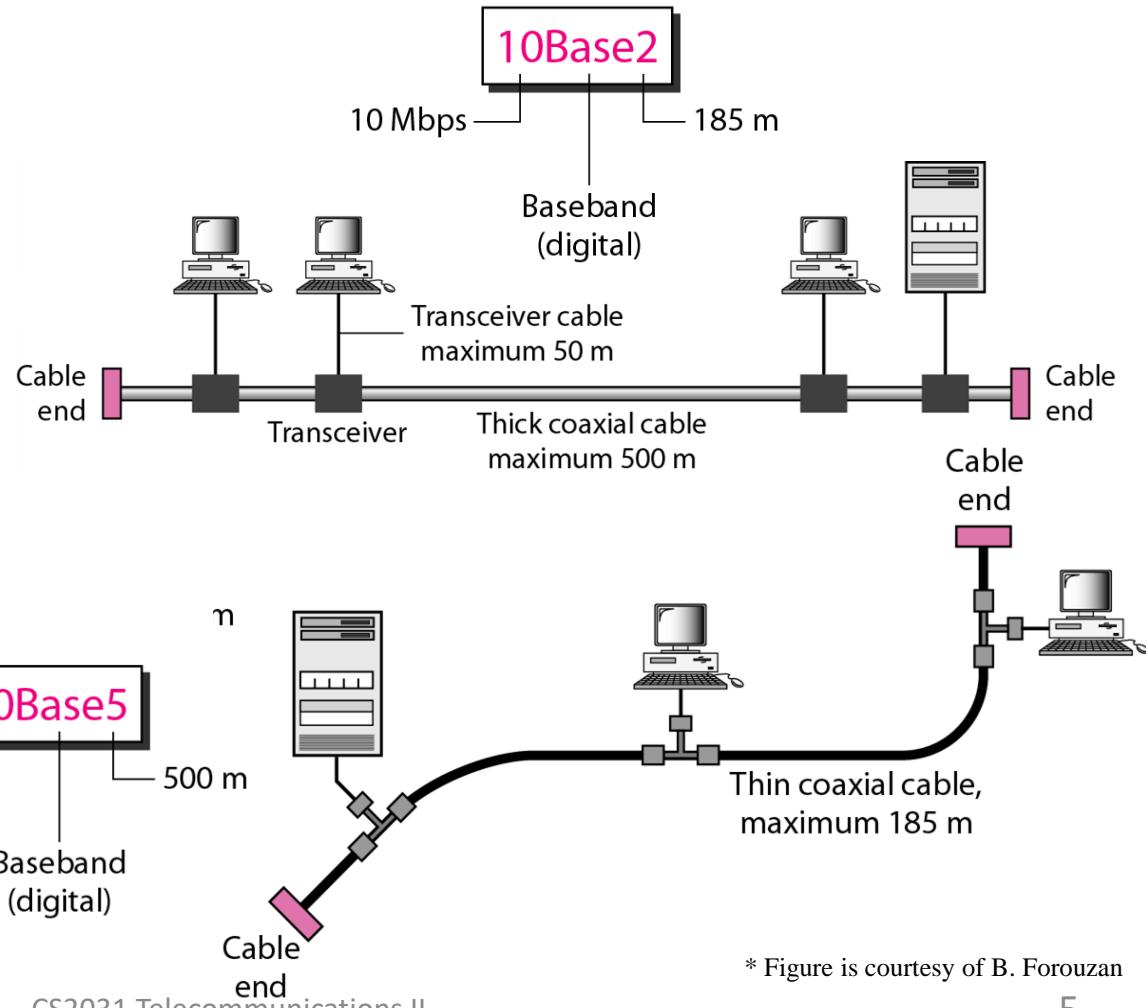
# Connecting Devices



- Physical Layer: Repeater or Hub
- Data Link Layer: Bridge or Layer-2 Switch
- Network Layer: Router or Layer-3 Switch

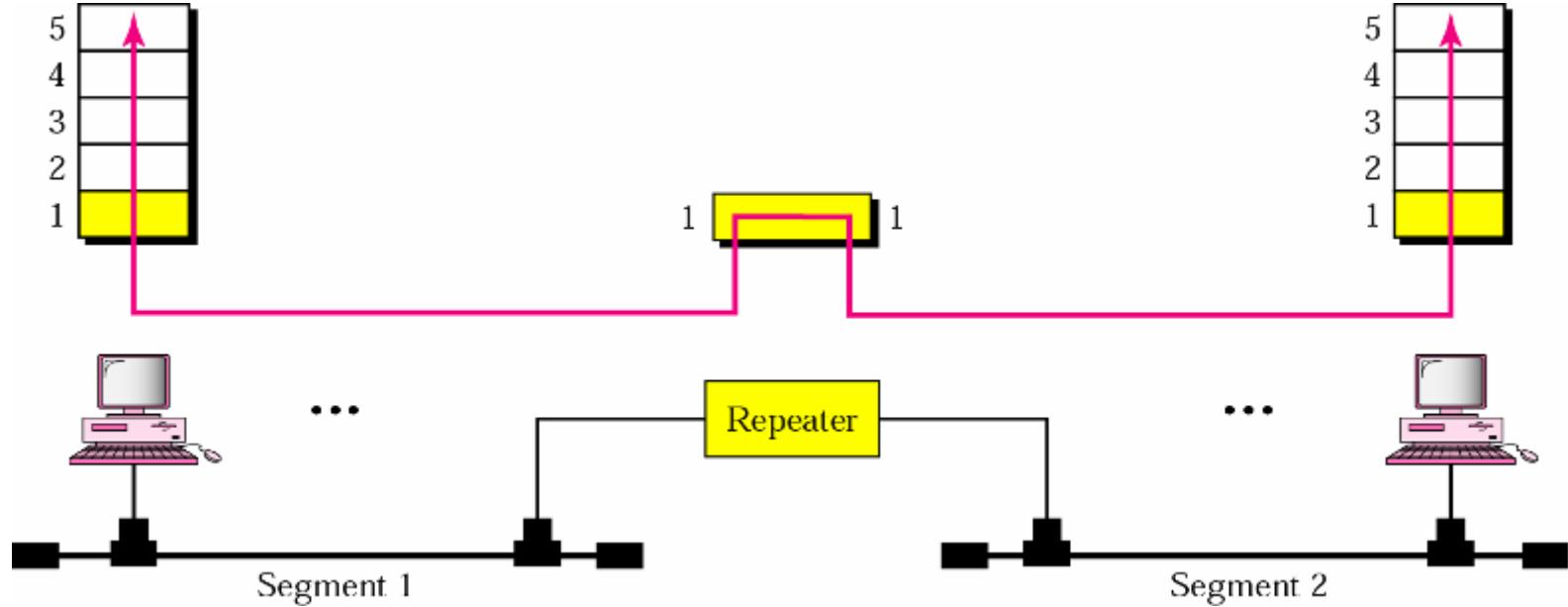
# 10Base5 & 10Base2

- Signal travels over cable & is picked up by all stations
- Used as backbone technology
- 10Base5: Stations linked into coaxial cable



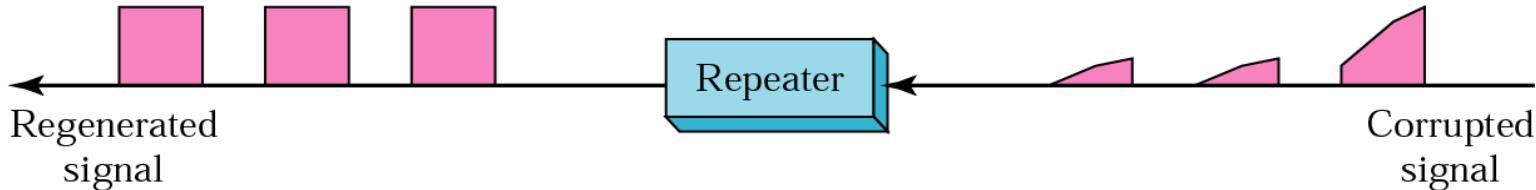
\* Figure is courtesy of B. Forouzan

# Repeater

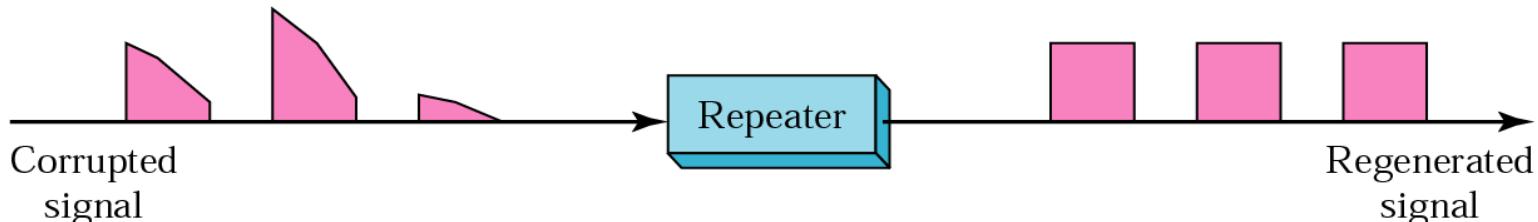


- A repeater connects segments of a LAN
- A repeater forwards every frame; it has no filtering capability

# Function of a Repeater



a. Right-to-left transmission.

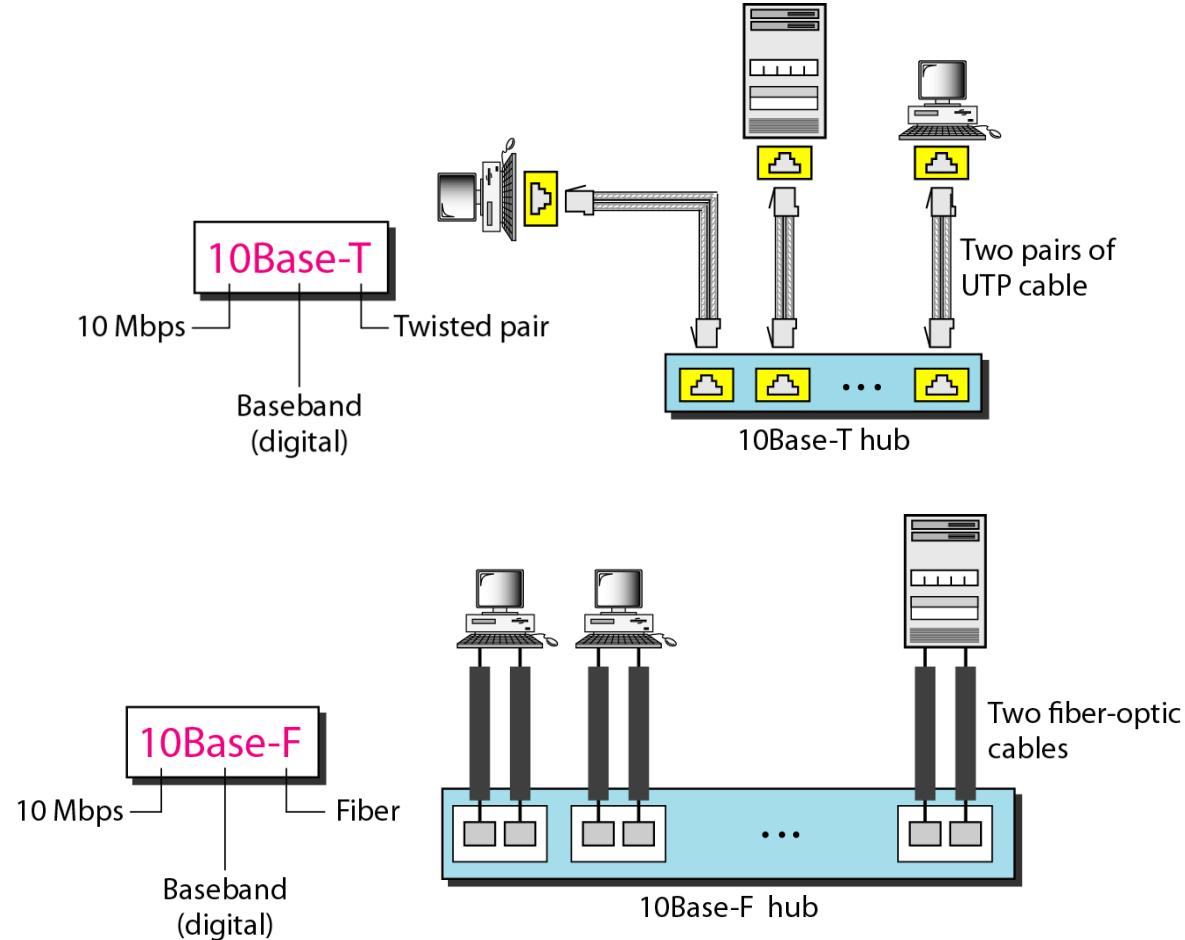


b. Left-to-right transmission.

- A repeater is a regenerator,  
not an amplifier!

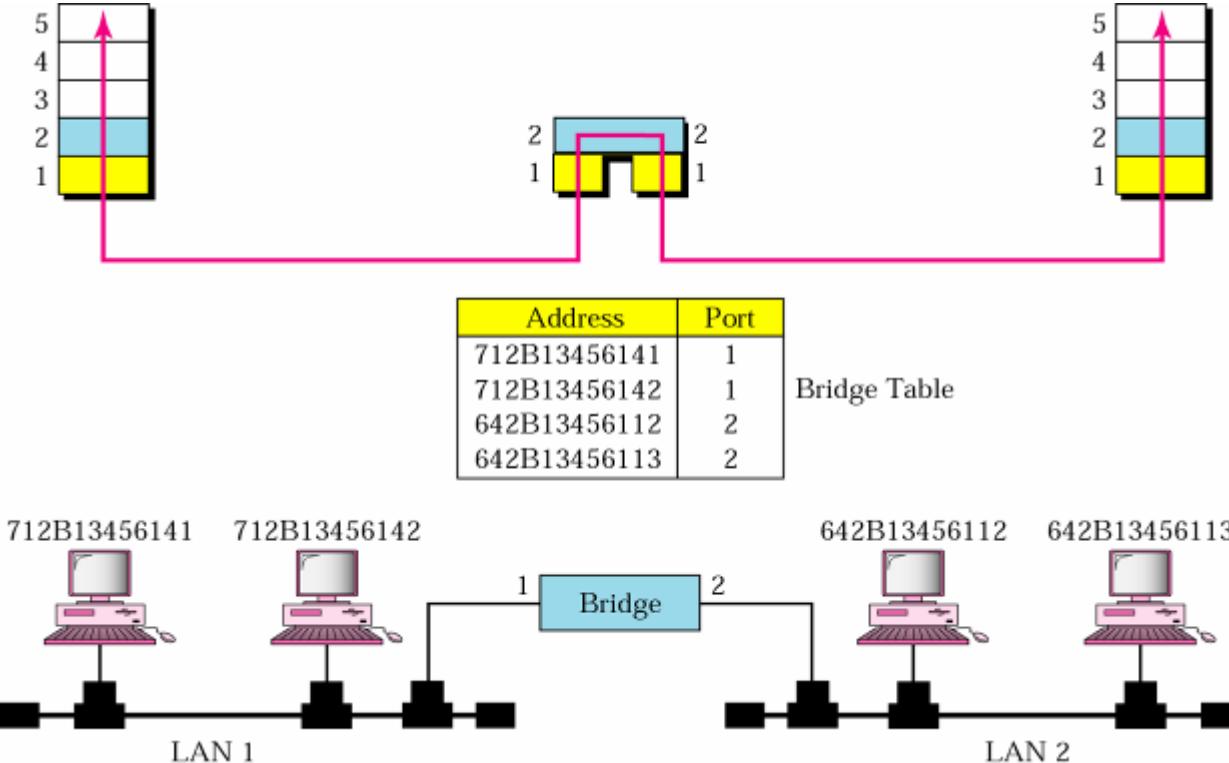
# 10Base-T & 10Base-F

- Hub replicates traffic to connected stations
- Each station has its own connection to hub
- Every station hears all traffic!



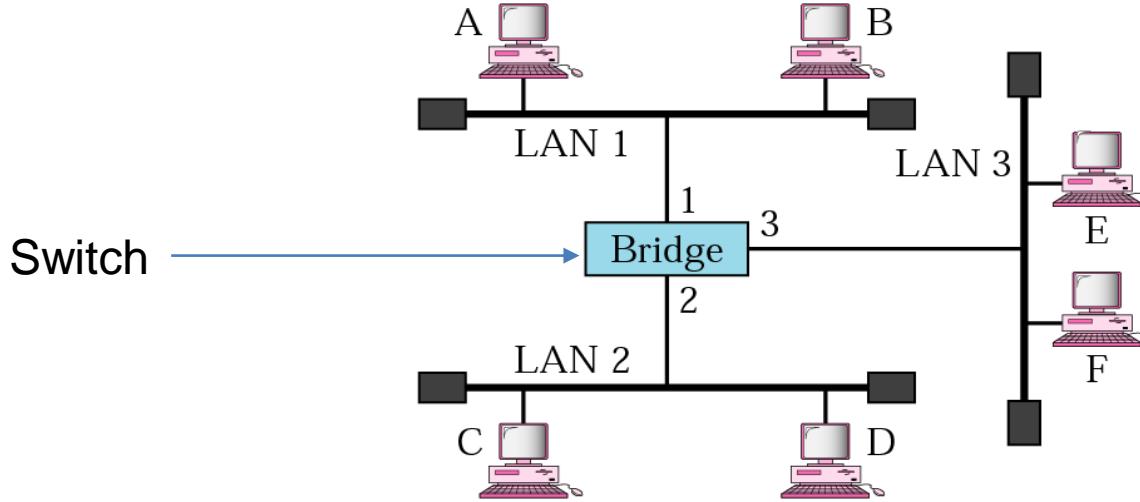
\* Figure is courtesy of B. Forouzan

# Functions of a Bridge



- Read all frames transmitted on one LAN and accept those address to any station on the other LAN
- Using MAC protocol for second LAN, retransmit each frame
- Do the same the other way round

# Learning Bridges



Address	Port

a. Original

Address	Port
A	1

b. After A sends a frame to D

Address	Port
A	1
E	3

c. After E sends a frame to A

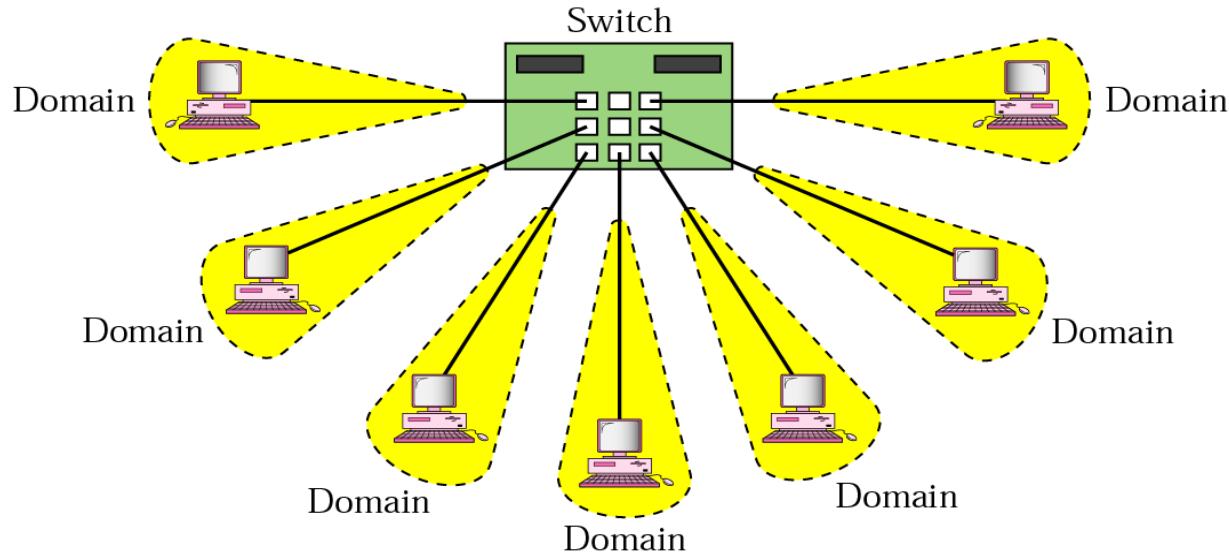
Address	Port
A	1
E	3
B	1

d. After B sends a frame to C

- Initially bridge forwards frames on all segments except incoming port
- Learns addresses from frames that pass through



# Switched Ethernet



- **Switch delivers packets to individual machines**
  - Without affecting communication with other machines
- **Collisions only occur on individual links**

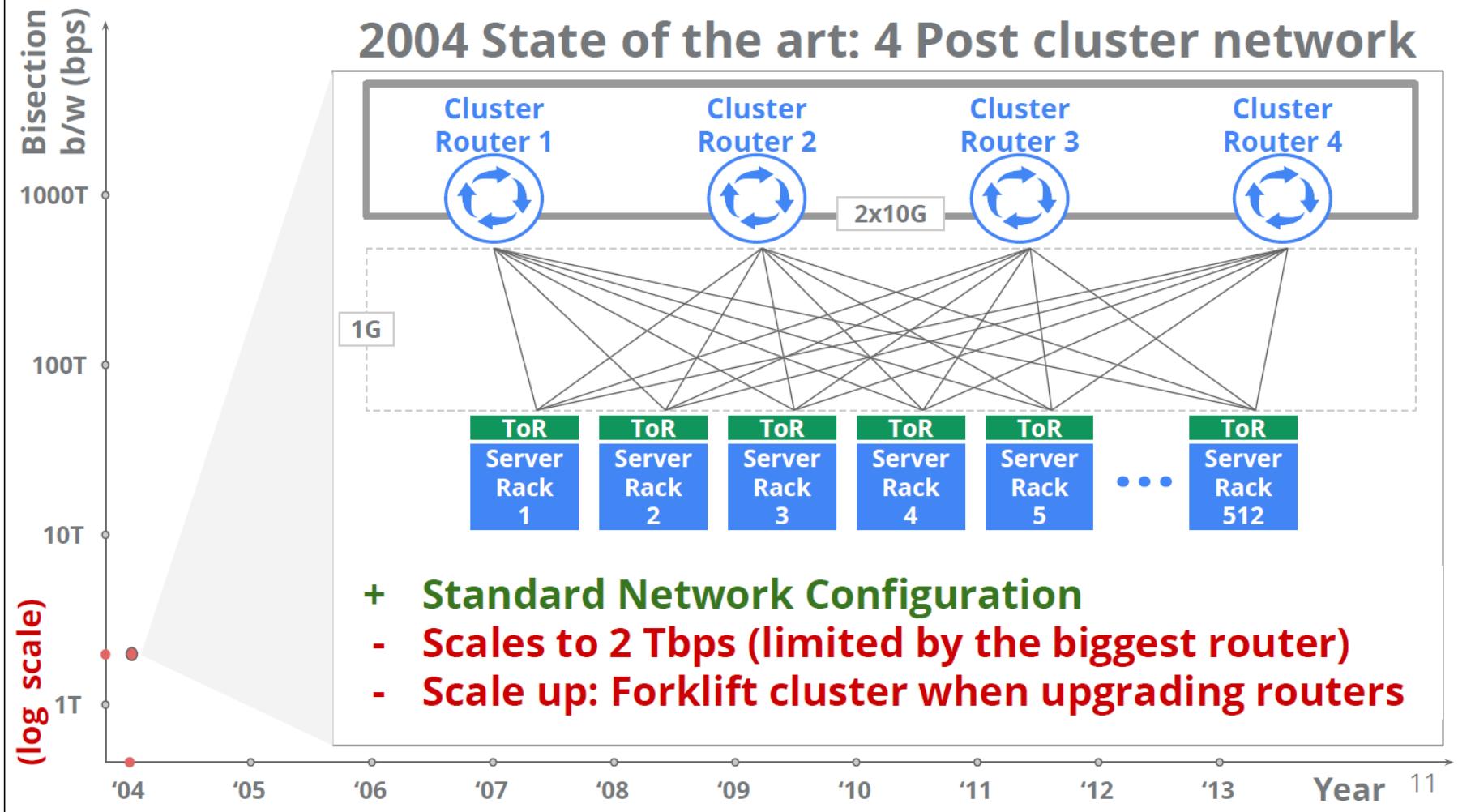
\* Figure is courtesy of B. Forouzan

# Switch



\* Figure is courtesy of Cisco & Reddit

# Google's Infrastructure



# Types of Layer 2 Switches

- **Store-and-Forward** switch
  - Accepts frame on input line
  - Buffers it briefly,
  - Then routes it to appropriate output line
  - Delay between sender and receiver
  - Boosts integrity of network
- **Cut-Through** switch
  - Takes advantage of destination address appearing at beginning of frame
  - Switch begins repeating frame onto output line as soon as it recognizes destination address
  - Highest possible throughput
  - Risk of propagating bad frames
    - Switch unable to check CRC prior to retransmission



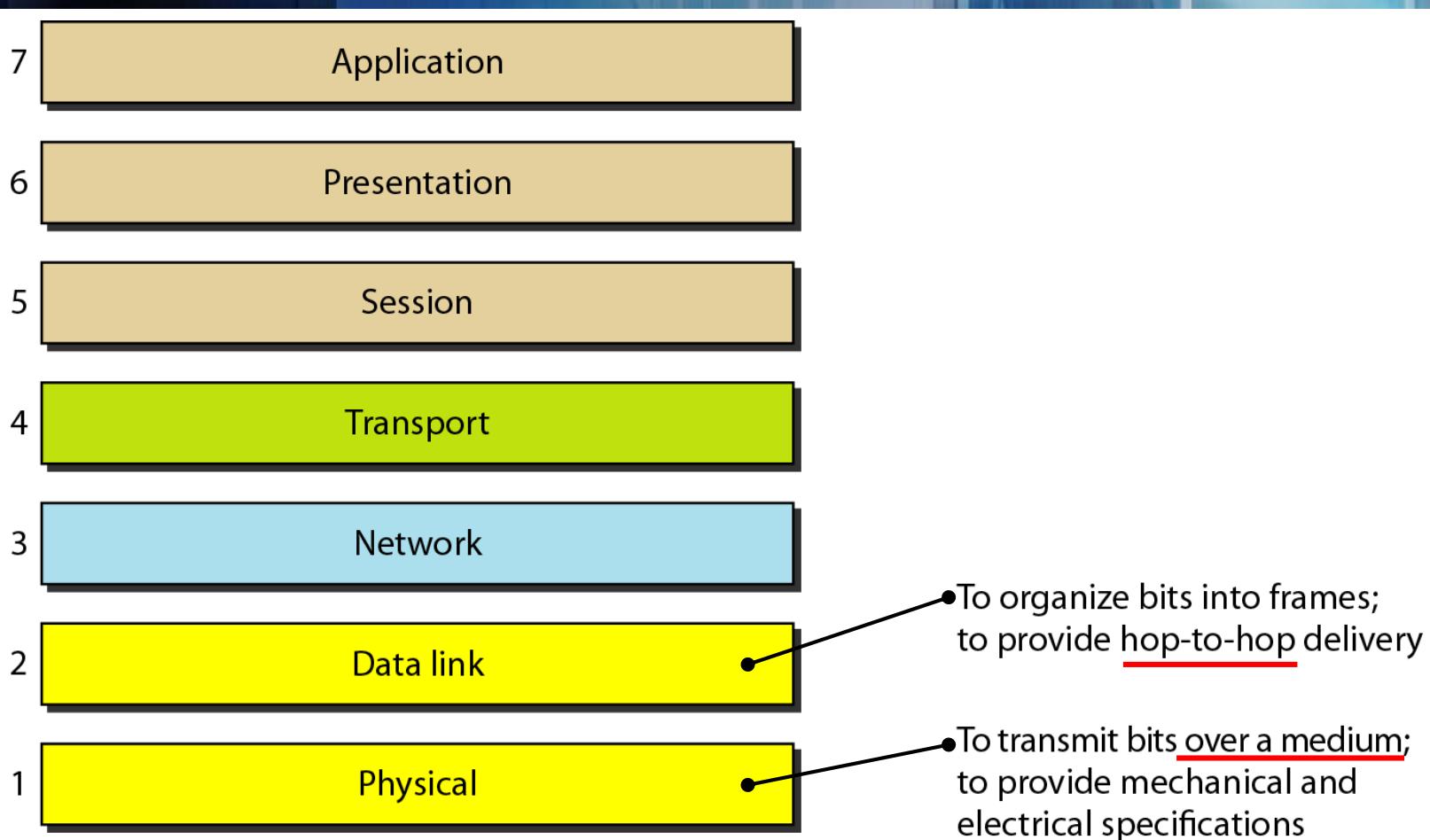
# CS2031

## Telecommunications II

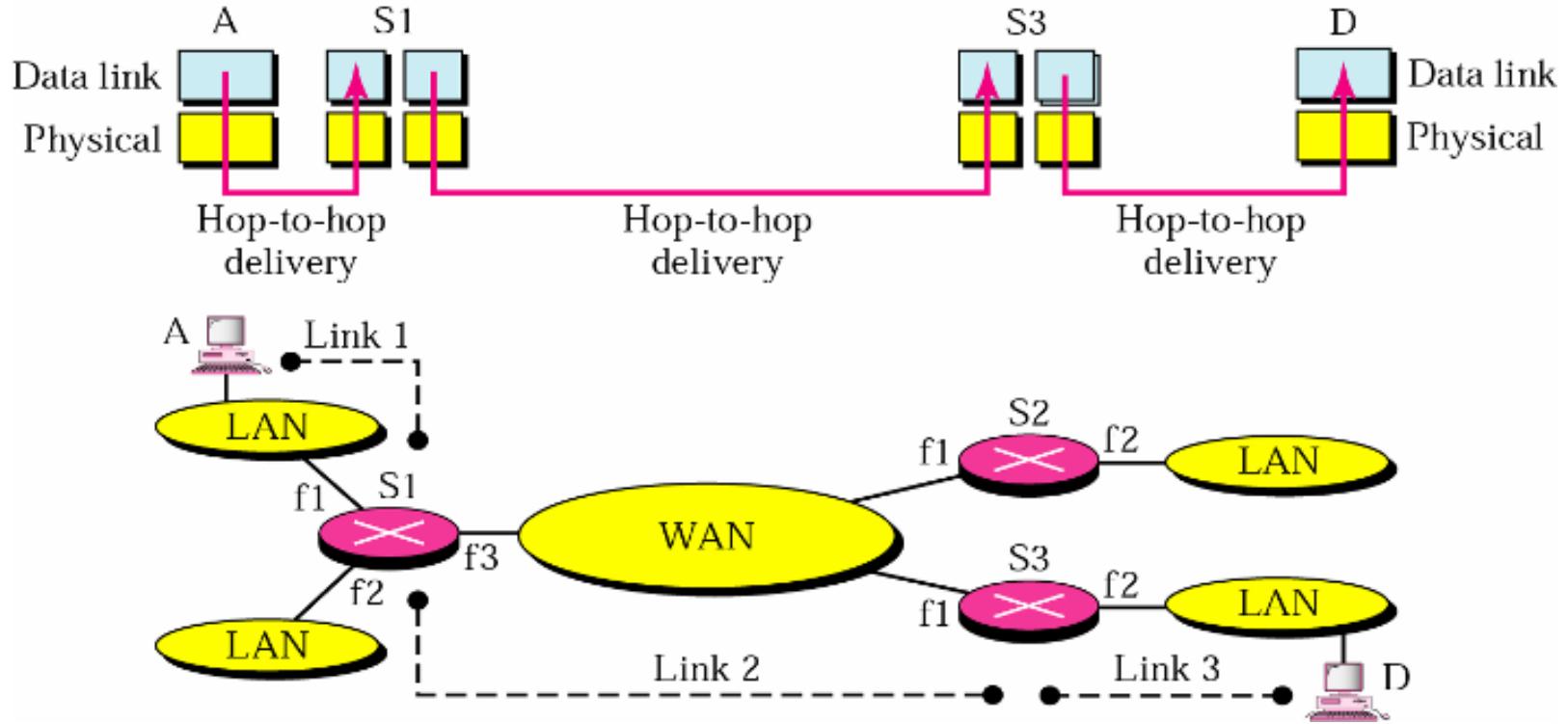
### Network Layer



# OSI Model



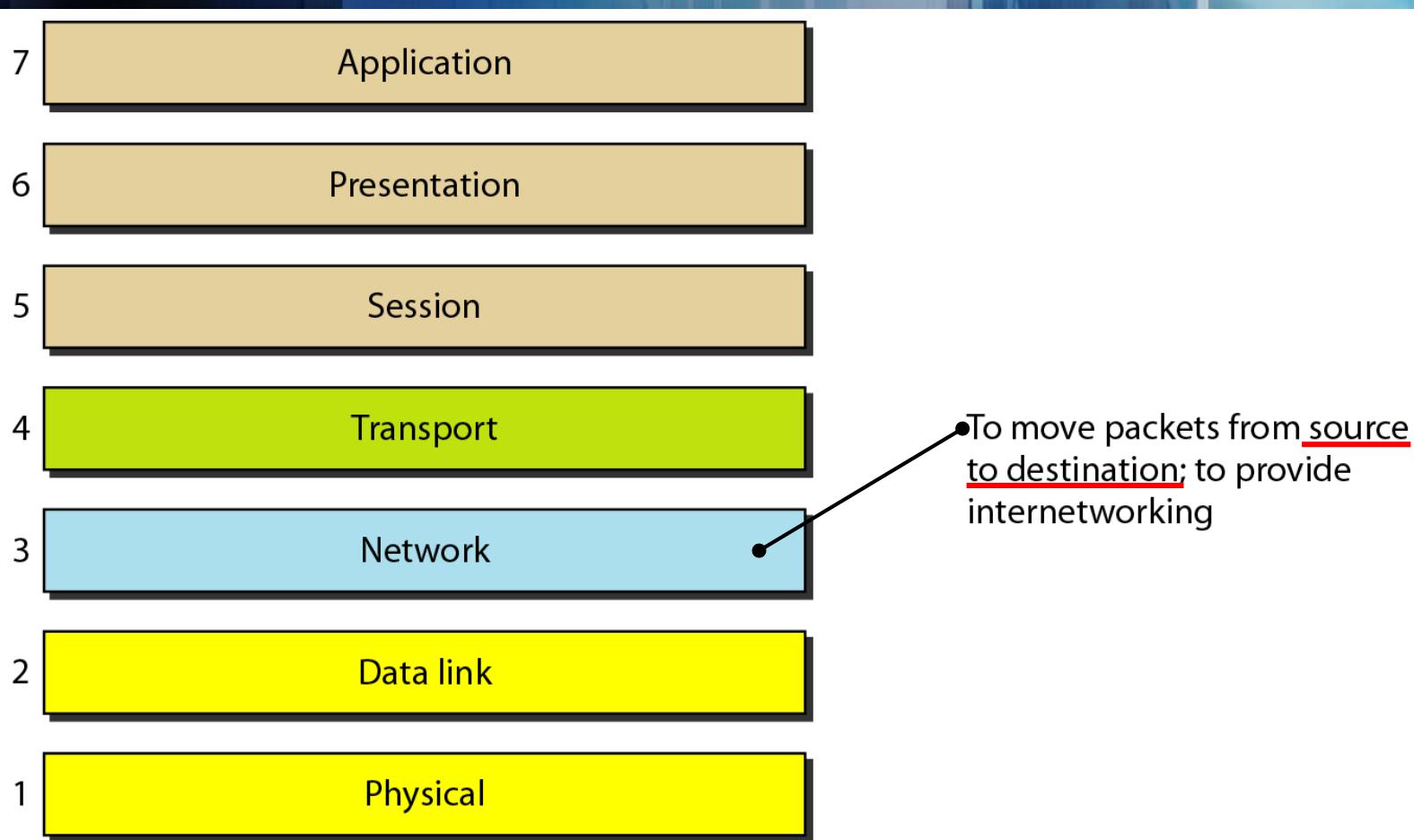
# Data Link Layer



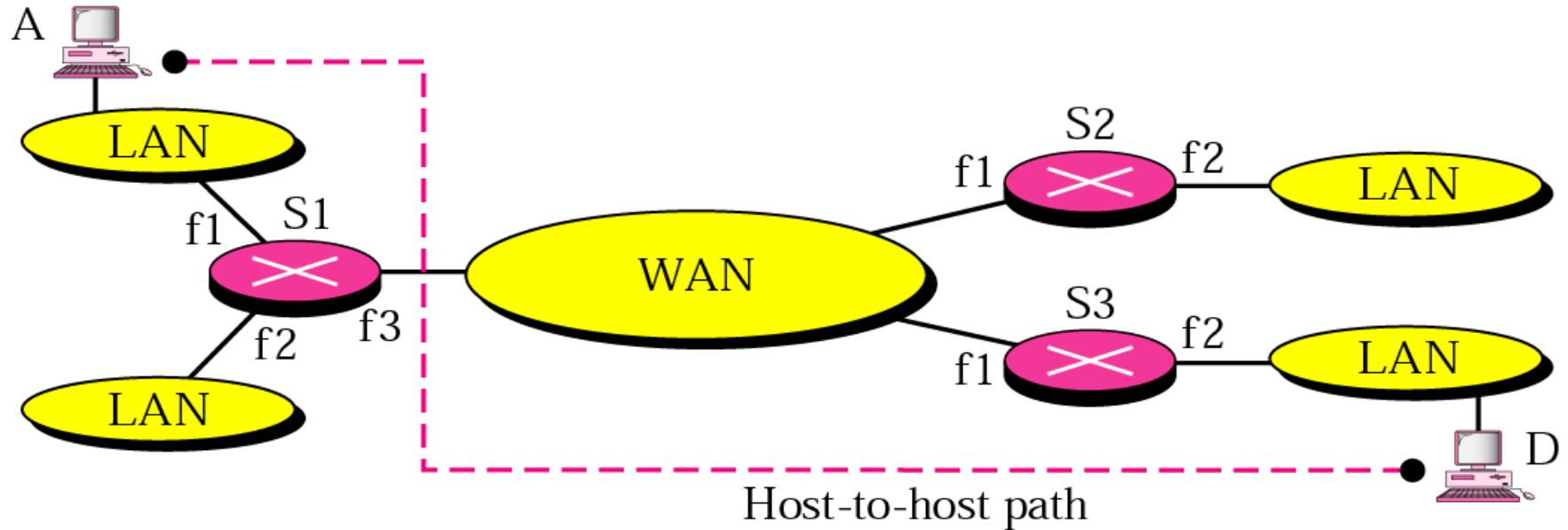
- Hop-to-hop communication

\* Figure is courtesy of B. Forouzan

# OSI Model

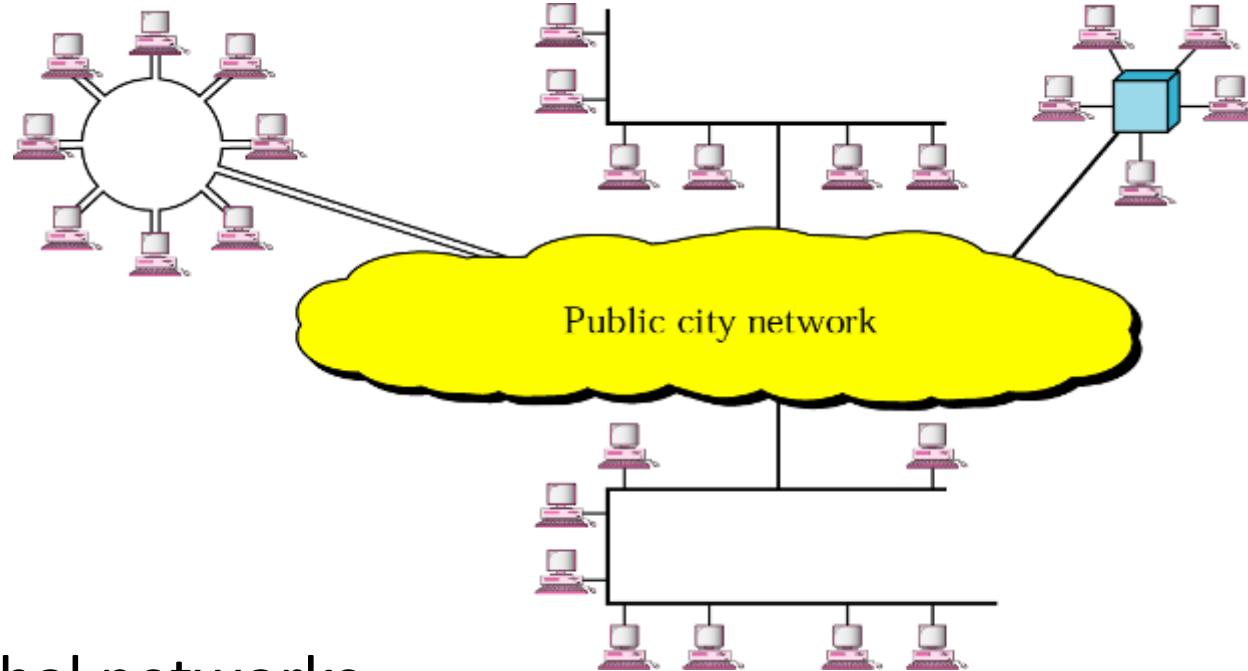


# Network Layer



- Responsible for source-to-destination delivery of a packet across multiple network links.

# Problems faced by Network Layers



- Global networks
  - Various different architectures
  - Various formats for hardware addresses
  - Various maximum transmission units (MTUs)

# Paul Baran, 1964

- Paul Barran, On Distributed Communication Networks, IEEE Transactions on Communication Systems, Volume 12, No. 1, March 1964
- Introduced Concepts of
  - Distributed Networks
  - Routing  
(hot-potato-routing)
  - Packet-Switching  
(message-block)

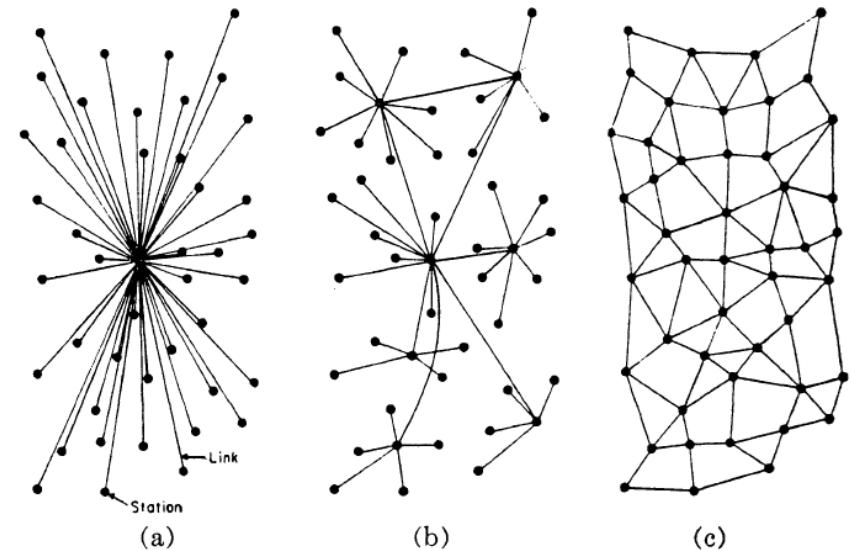
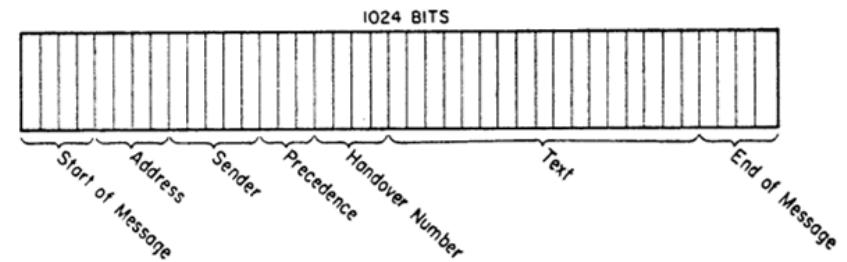


Fig. 1—(a) Centralized. (b) Decentralized. (c) Distributed networks.



AT&T's comment back then was that packet-switching will never be useful.

# HOST-HOST Communication, 1970

- Paul Barran, [On Distributed Communication Networks](#), IEEE Transactions on Communication Systems, Volume 12, No. 1, March 1964
- Introduced Concepts of
  - Distributed Networks
  - Routing  
(hot-potato-routing)
  - Packet-Switching  
(message-block)
- Name worth mentioning: Bob Kahn

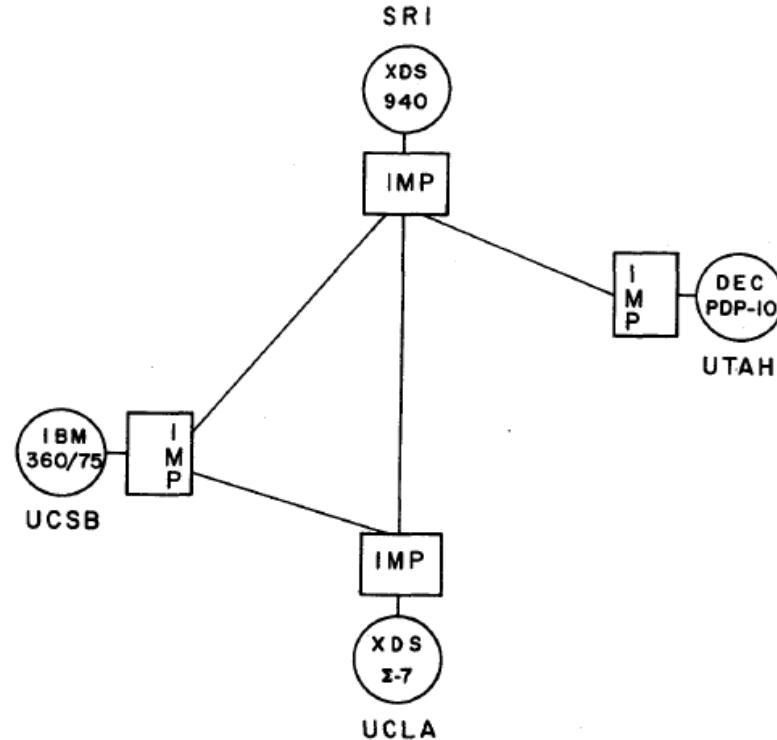
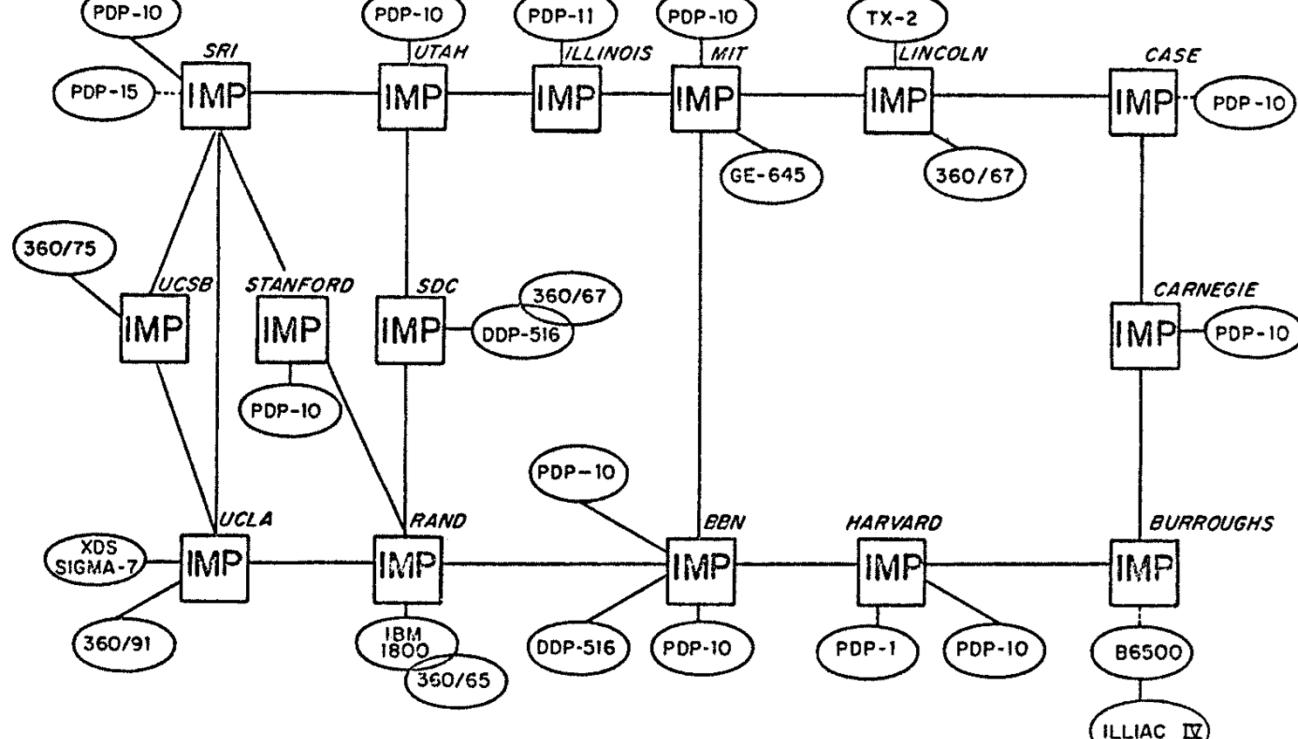


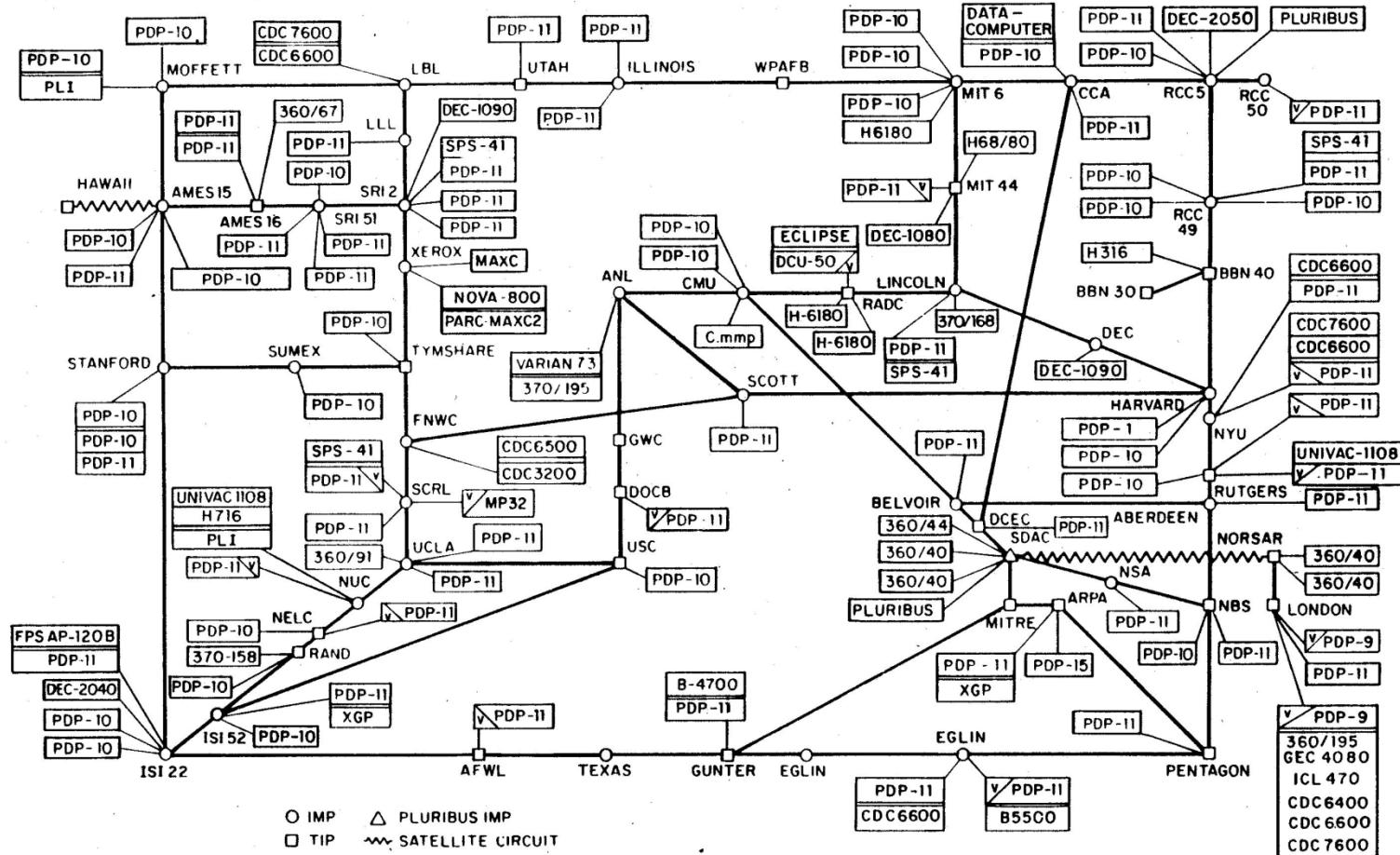
Figure 1—Initial network configuration

\* C. Stephen Carr, Stephen D. Crocker, and Vinton G. Cerf, [HOST-HOST Communication Protocol in the ARPA Network](#), Spring Joint Computer Conference, pages 589-597, Atlantic City, NJ, USA, May 1970

# ARPANET 1971



# ARPANET 1977



(PLEASE NOTE THAT WHILE THIS MAP SHOWS THE HOST POPULATION OF THE NETWORK ACCORDING TO THE BEST INFORMATION OBTAINABLE, NO CLAIM CAN BE MADE FOR ITS ACCURACY.)

NAMES SHOWN ARE IMP NAMES, NOT (NECESSARILY) HOST NAMES



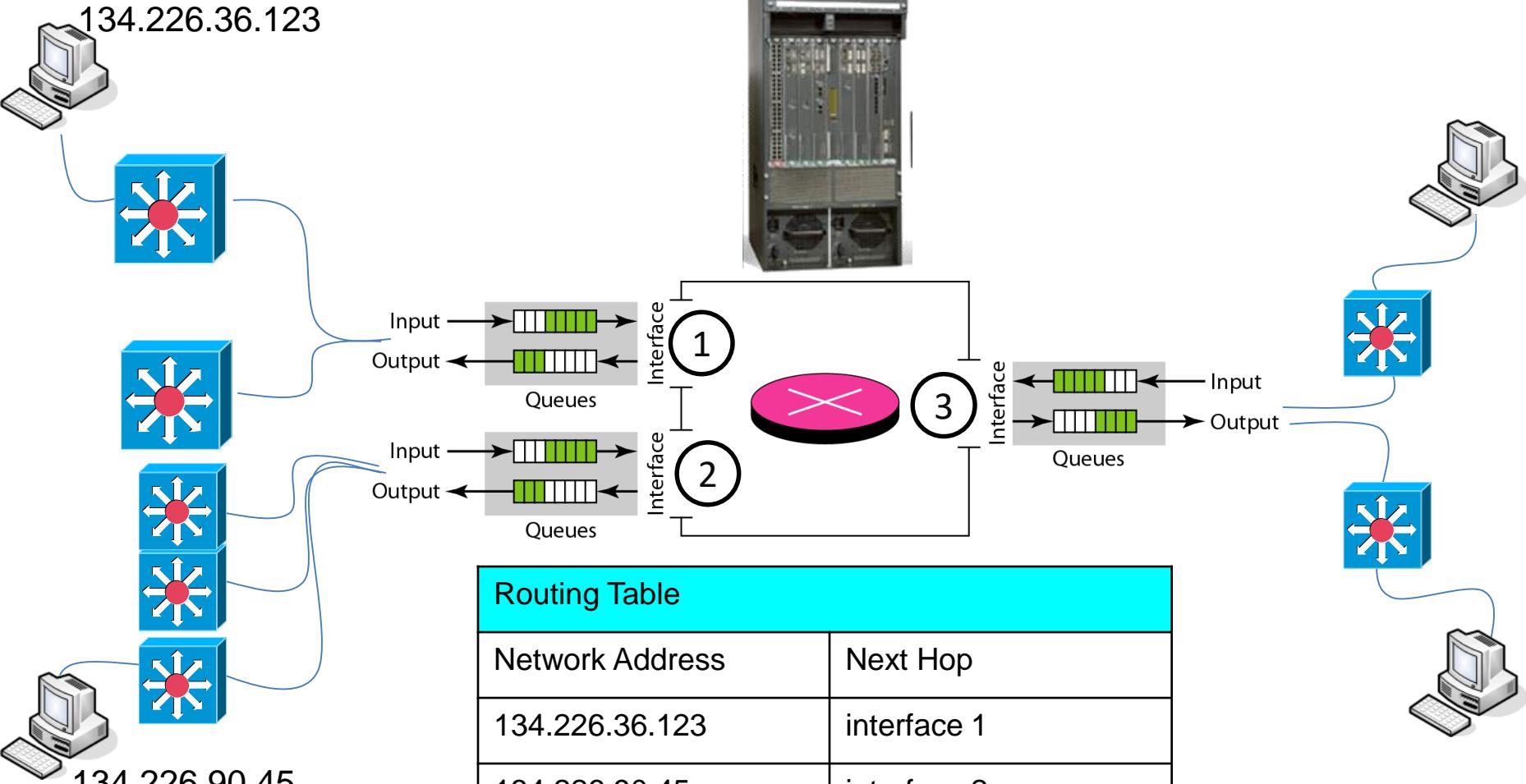
# IP Addresses

10000000 00001011 00000011 00011111

**128.11.3.31**

- 32-bit number
  - 4.294.967.296 addresses
- IP addresses are unique and universal
  - with some exceptions
- Dotted decimal notation:
  - Bytes of binary notation represented as decimal separated by dot

# Task of Routers



# Routers

- One Main Interest  
**Forwarding Packets**
- Important Aspects  
**Queue Length**  
**Routing Table**

Destination	Gateway	Interface
IP Range <sub>1</sub>	G <sub>1</sub>	IF <sub>1</sub>
IP Range <sub>2</sub>	G <sub>2</sub>	IF <sub>2</sub>



\* Figure is courtesy of Cisco

# Everything's a Router

**Active Routes:**

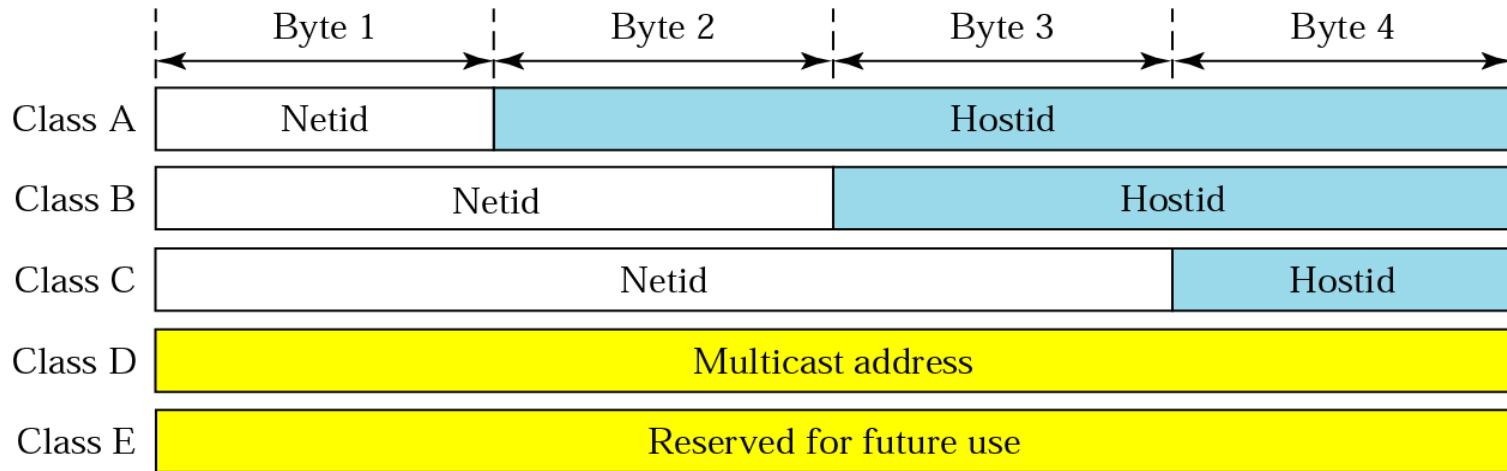
Network	Destination	Netmask	Gateway	Interface	Metric
	0.0.0.0	0.0.0.0	192.168.192.1	192.168.192.37	25
	127.0.0.0	255.0.0.0	On-link	127.0.0.1	306
	127.0.0.1	255.255.255.255	On-link	127.0.0.1	306
127.255.255.255	255.255.255.255		On-link	127.0.0.1	306
	192.168.21.0	255.255.255.0	On-link	192.168.21.1	276
	192.168.21.1	255.255.255.255	On-link	192.168.21.1	276
	192.168.21.255	255.255.255.255	On-link	192.168.21.1	276
	192.168.111.0	255.255.255.0	On-link	192.168.111.1	276
	192.168.111.1	255.255.255.255	On-link	192.168.111.1	276
192.168.111.255	255.255.255.255		On-link	192.168.111.1	276
	192.168.150.0	255.255.255.0	On-link	192.168.150.1	276
	192.168.150.1	255.255.255.255	On-link	192.168.150.1	276
192.168.150.255	255.255.255.255		On-link	192.168.150.1	276
	192.168.192.0	255.255.255.0	On-link	192.168.192.37	281
	192.168.192.37	255.255.255.255	On-link	192.168.192.37	281
192.168.192.255	255.255.255.255		On-link	192.168.192.37	281



# 4.294.967.296 Possible Nodes???



# Network ID and Host ID

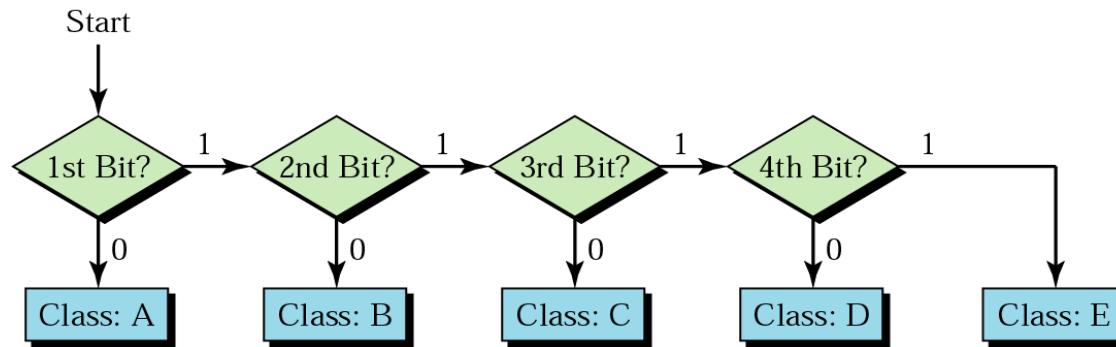


- Network ID: Used to find a particular network
- Host ID: Identifies individual nodes

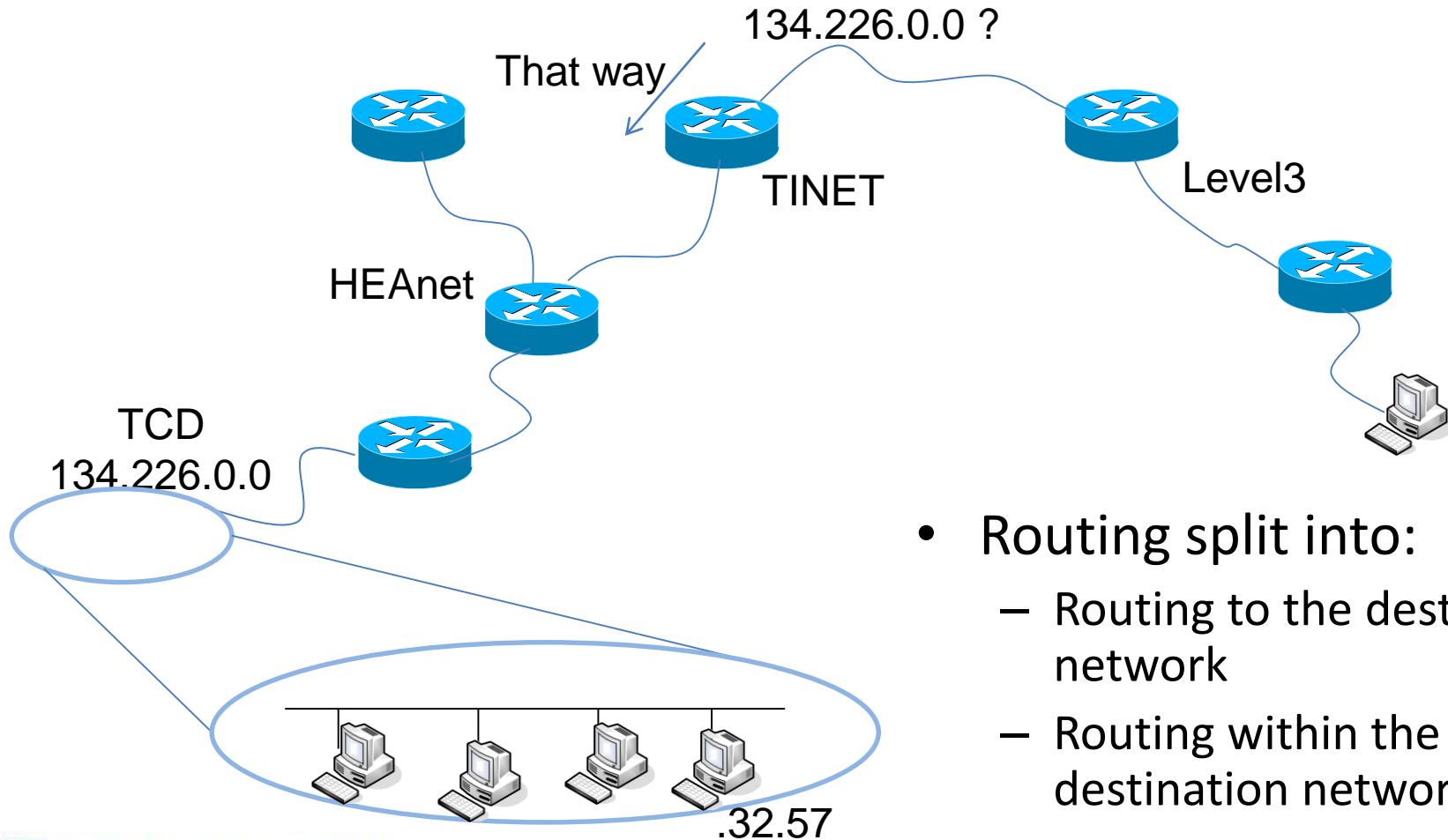
# Classes in Binary Notation

	First byte	Second byte	Third byte	Fourth byte
Class A	<b>0</b>			
Class B	<b>10</b>			
Class C	<b>110</b>			
Class D	<b>1110</b>			
Class E	<b>1111</b>			

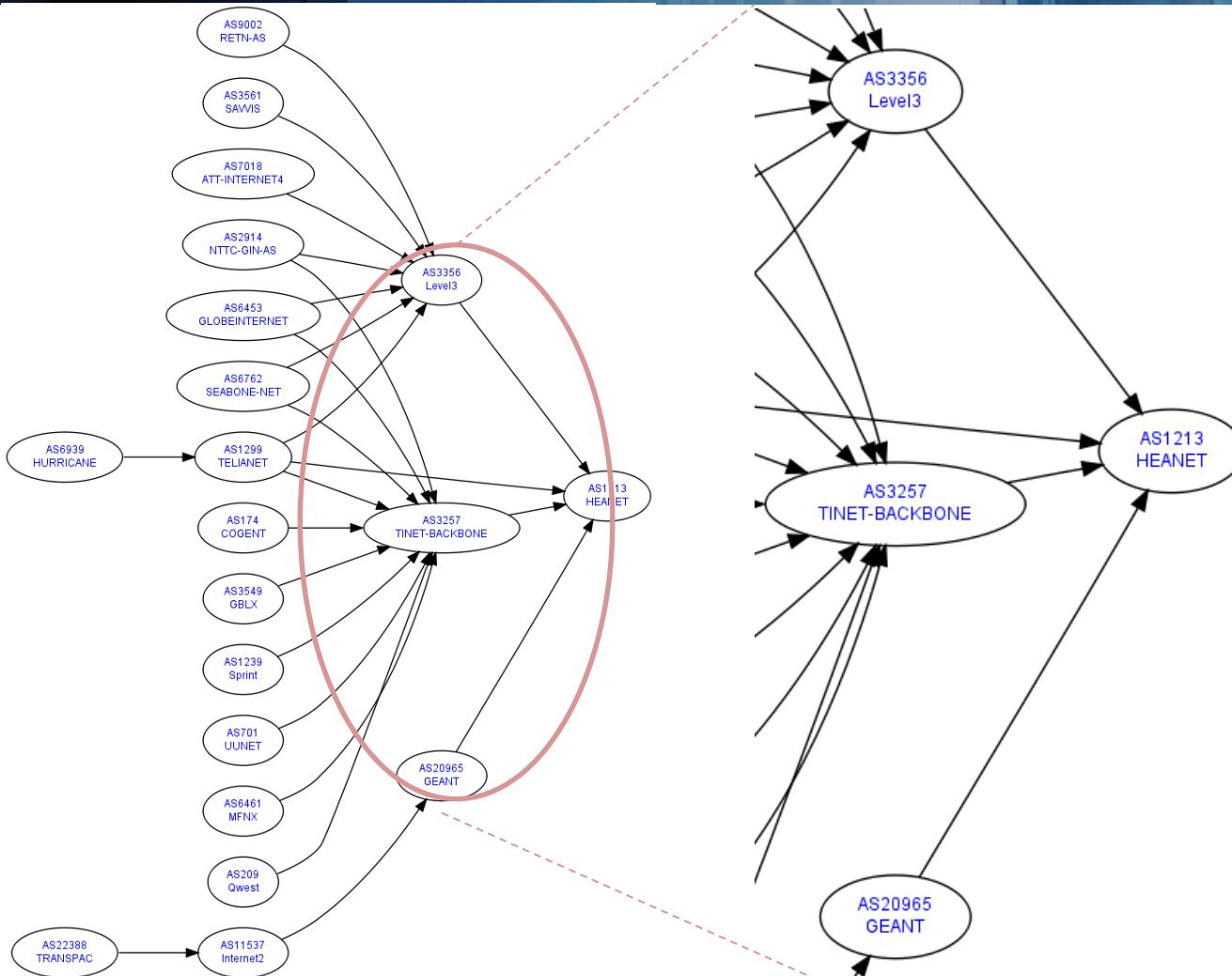
- Decision Process:



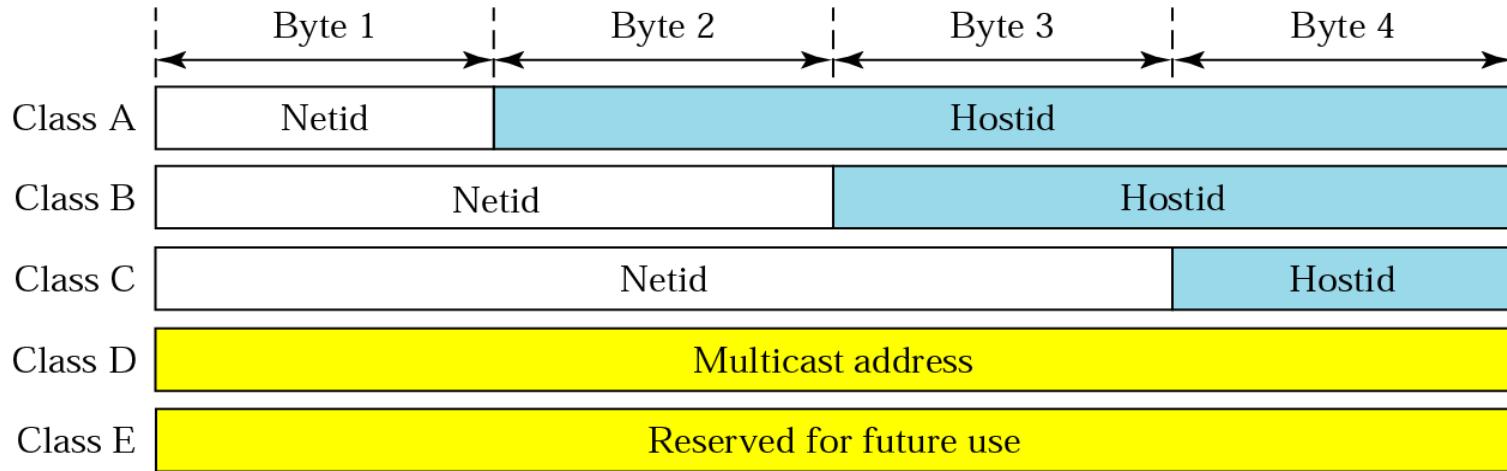
# Network IDs and Host IDs



# AS1213 - HEANET

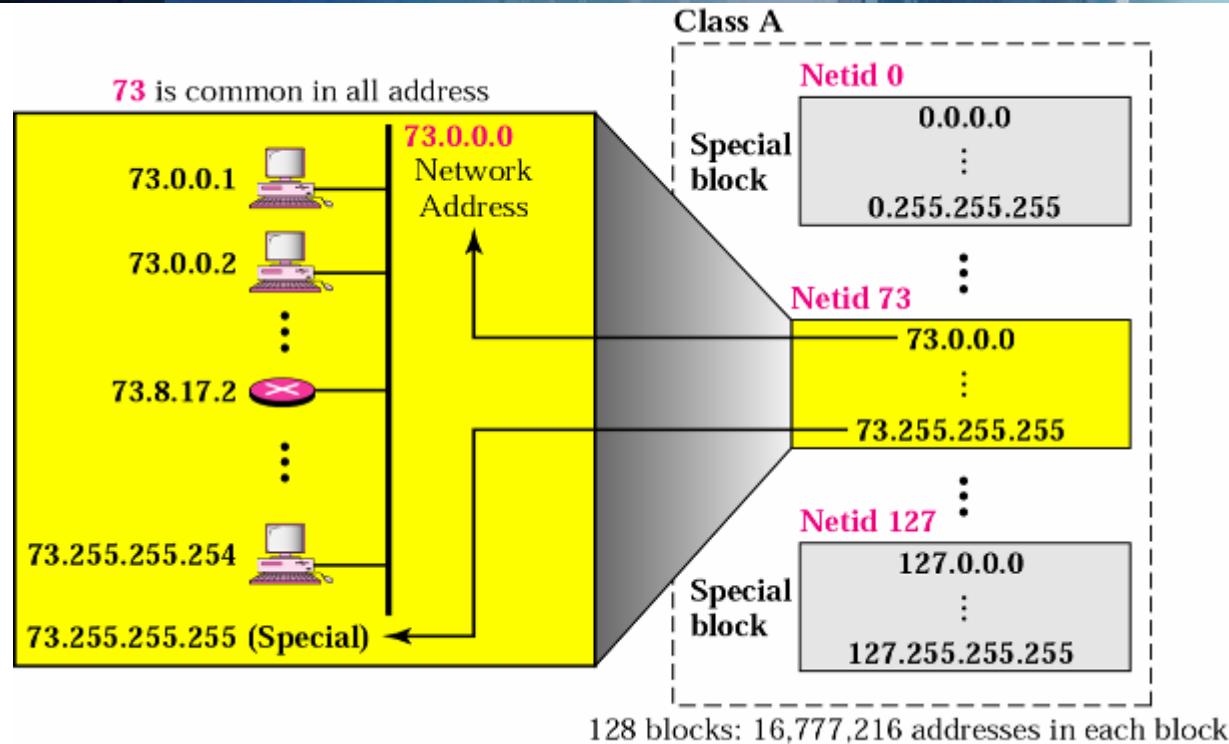


# Classful Addresses



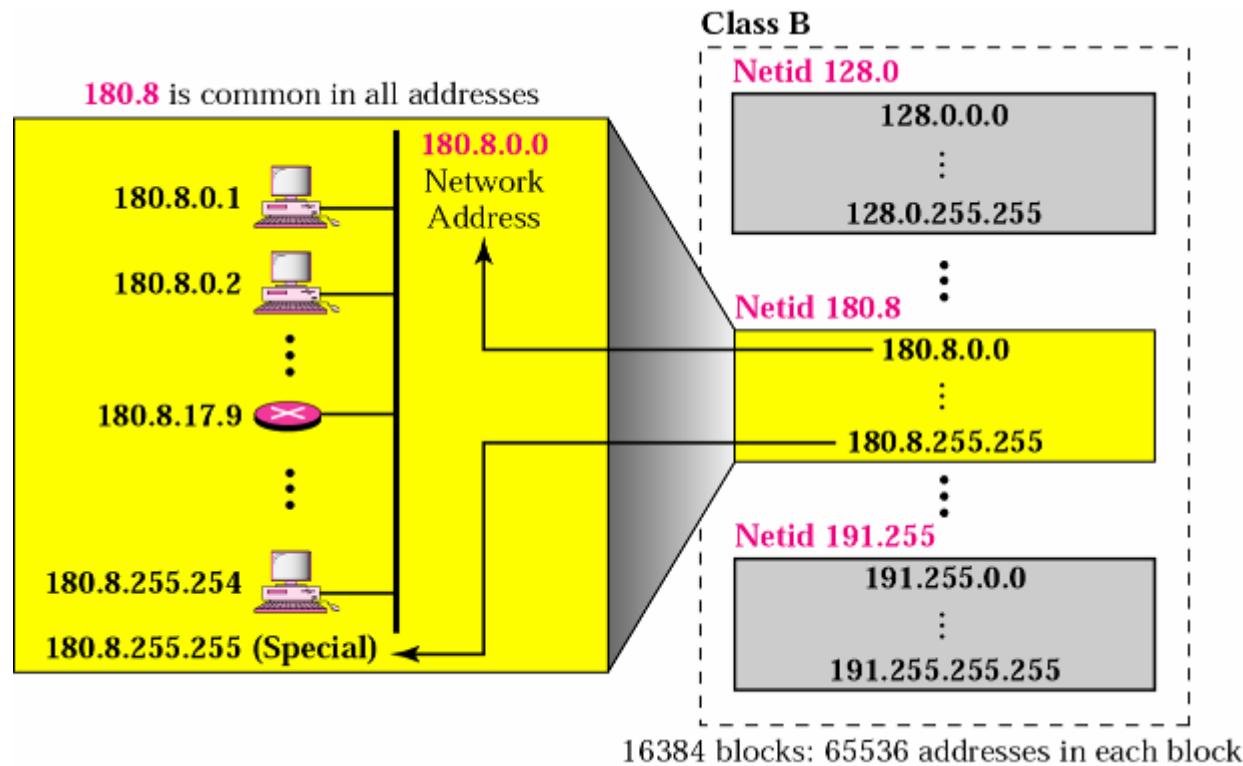
- Class A (international organisations)
  - 126 networks with 16,277,214 hosts each
- Class B (large companies)
  - 16,384 networks with 65,354 hosts each
- Class C (smaller companies)
  - 2,097,152 networks with 254 hosts each

# Class A



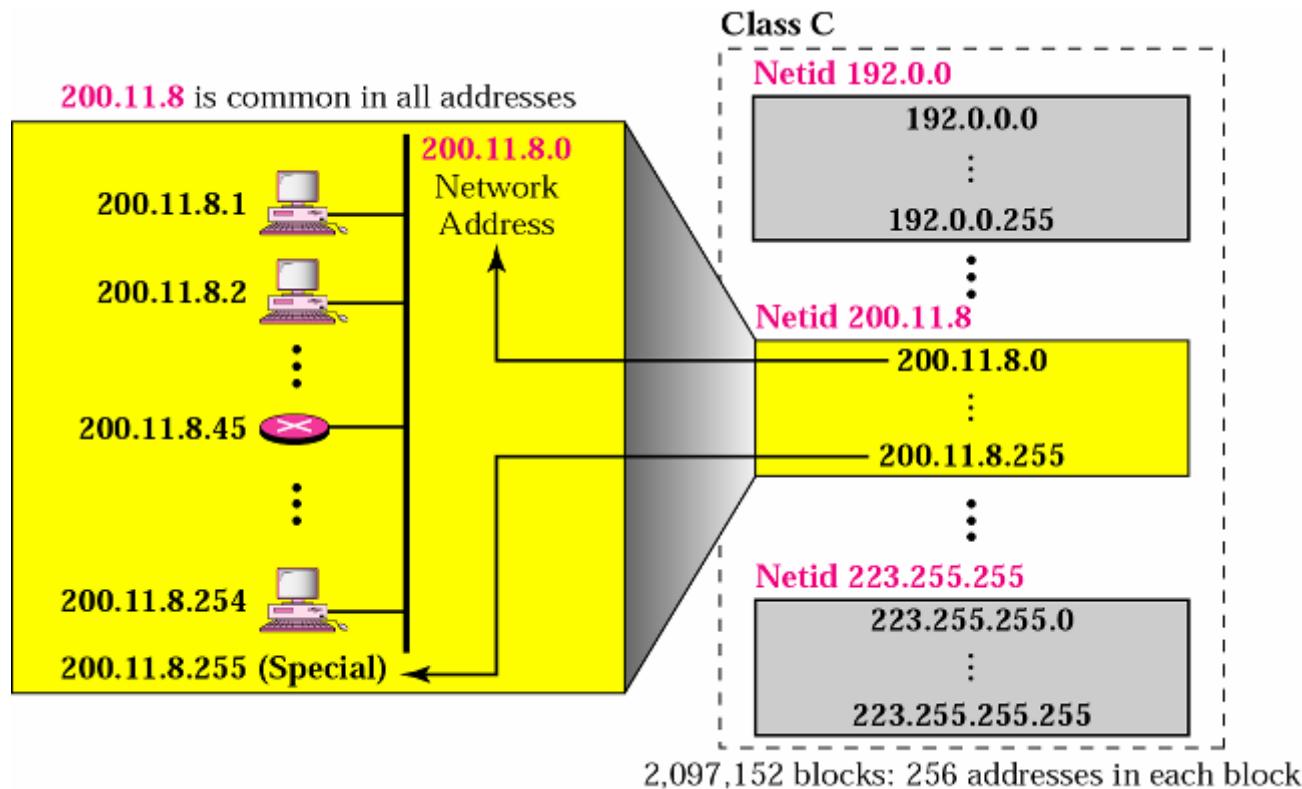
- Limited number of very large networks
- 126 networks with 16,277,214 hosts each

# Class B



- Limited number of relatively large address ranges
- 16,384 networks with 65,354 hosts each

# Class C



- Large number of small address ranges
- 2,097,152 networks with 254 hosts each



# Classes & Private Addresses

	First byte	Second byte	Third byte	Fourth byte
Class A	0			
Class B	10			
Class C	110			
Class D	1110			
Class E	1111			

a. Binary notation

	First byte	Second byte	Third byte	Fourth byte
Class A	0–127			
Class B	128–191			
Class C	192–223			
Class D	224–239			
Class E	240–255			

b. Dotted-decimal notation

Private Address Ranges:

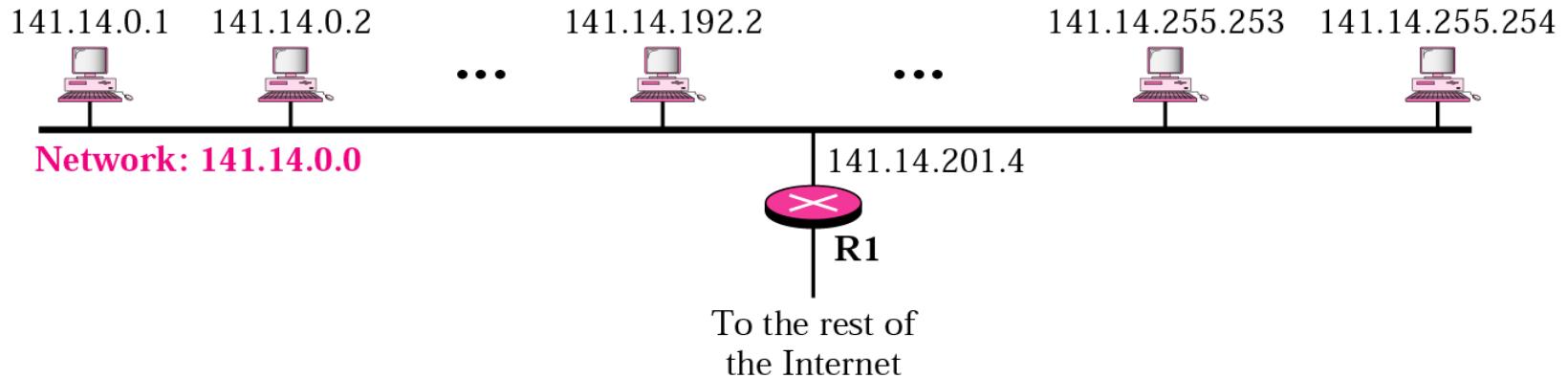
Range	Total
10.0.0.0 to 10.255.255.255	$2^{24}$
172.16.0.0 to 172.31.255.255	$2^{20}$
192.168.0.0 to 192.168.255.255	$2^{16}$

# Special Addresses

- Loopback device: 127.x.x.x
  - e.g. 127.0.0.1 = localhost
- This network: 0.0.0.x (all zero's)
  - e.g. 0.0.0.54 = host 54 on this network
- Broadcast: x.x.255.255 (all one's)
  - e.g. 134.226.36.255 = all nodes in this network



# 2-Level Hierarchy with Classful Addresses



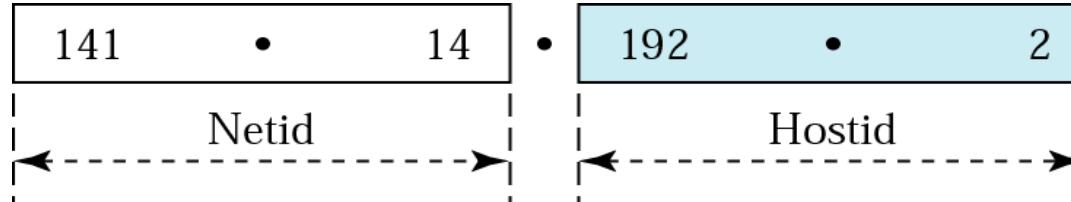
- 1<sup>st</sup> level: NetworkID within the Internet
- 2<sup>nd</sup> level: HostID within the network

# College Network

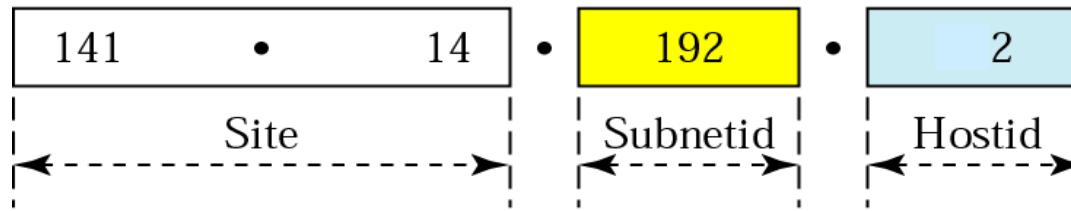
- Flat Network with 1000s of Nodes



# Subnetting



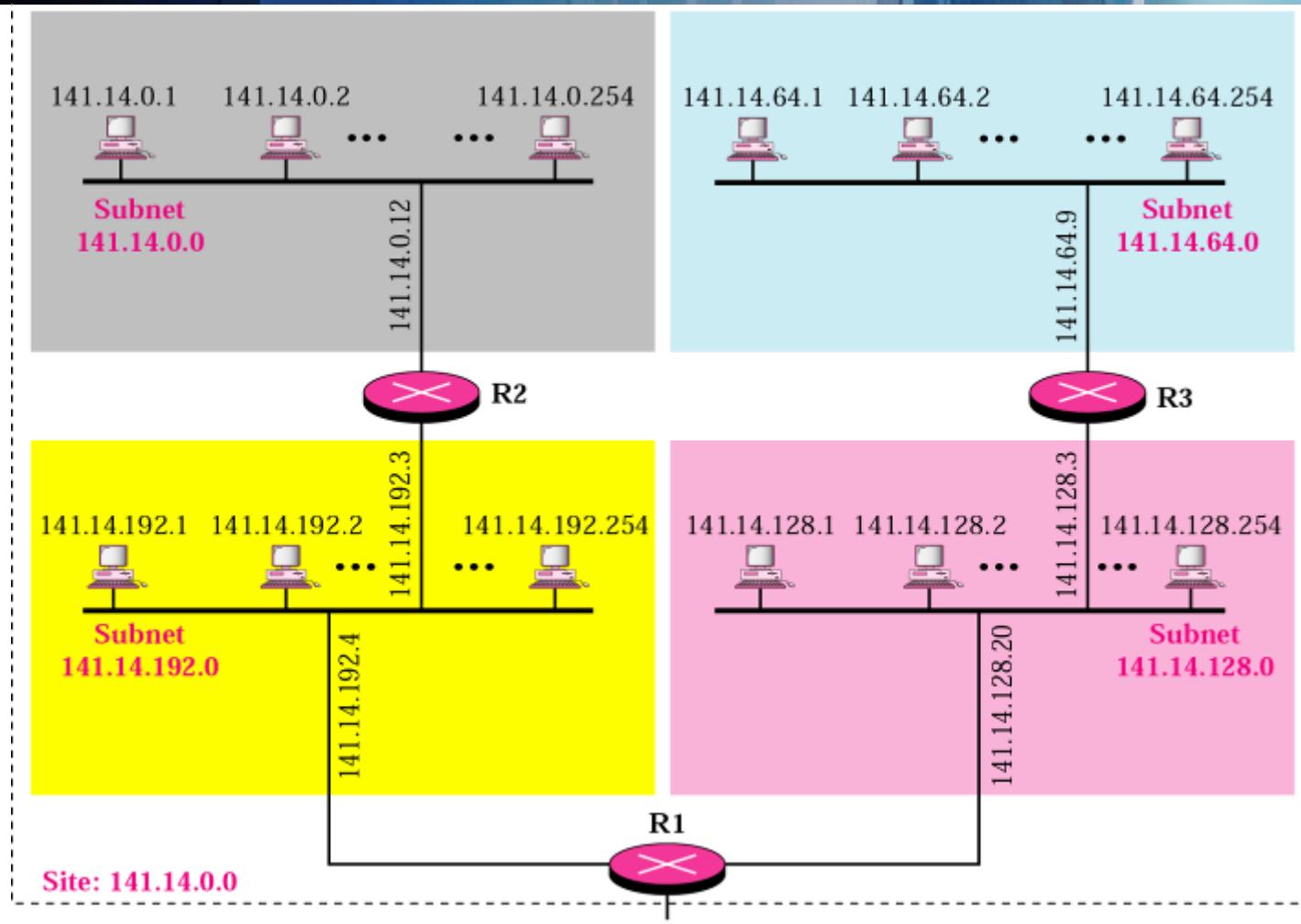
a. Without subnetting



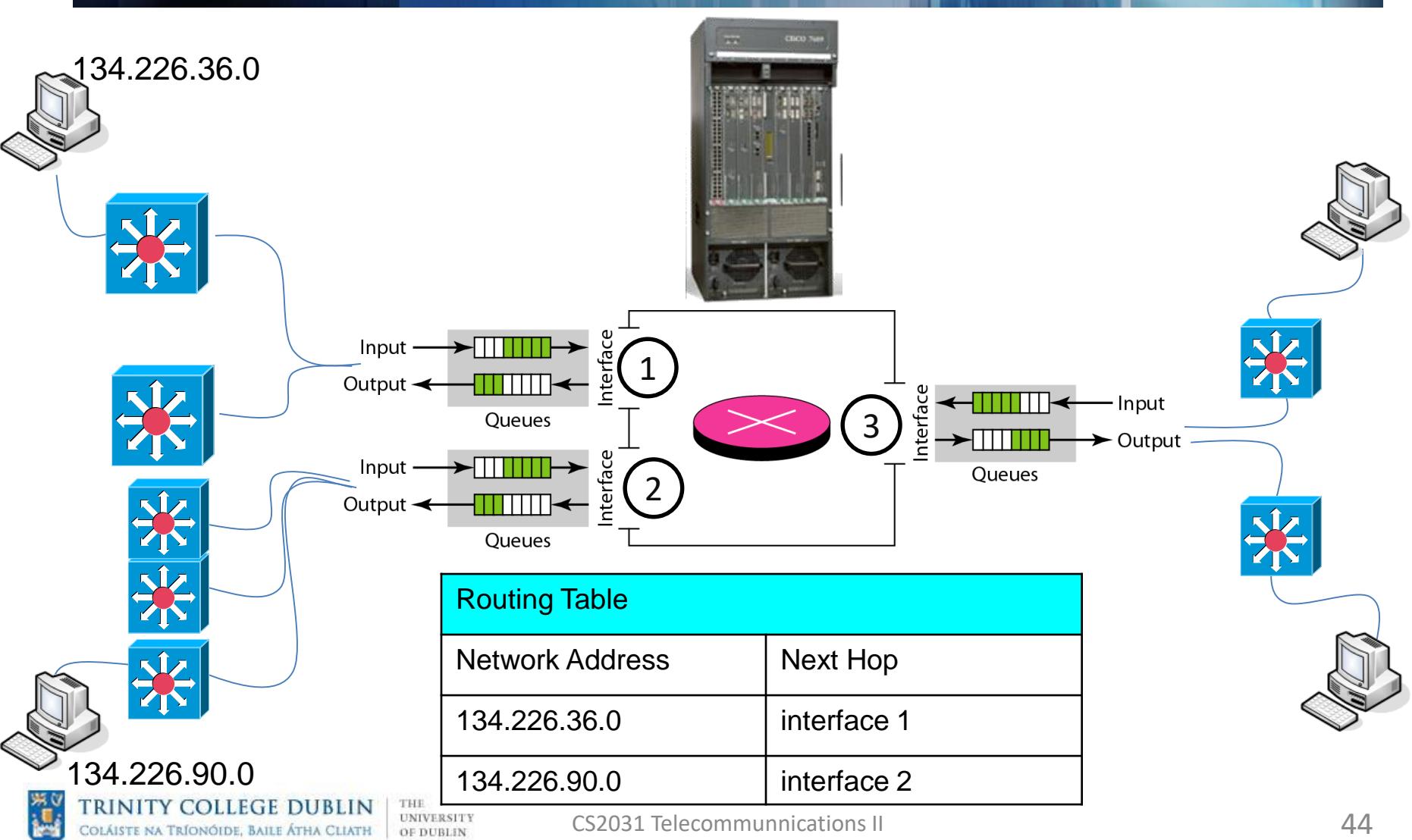
b. With subnetting

- Add another level to address hierarchy: *subnet*
- Splitting a class B network into a number of class C subnets

# 3-Level Hierarchy through Subnetting

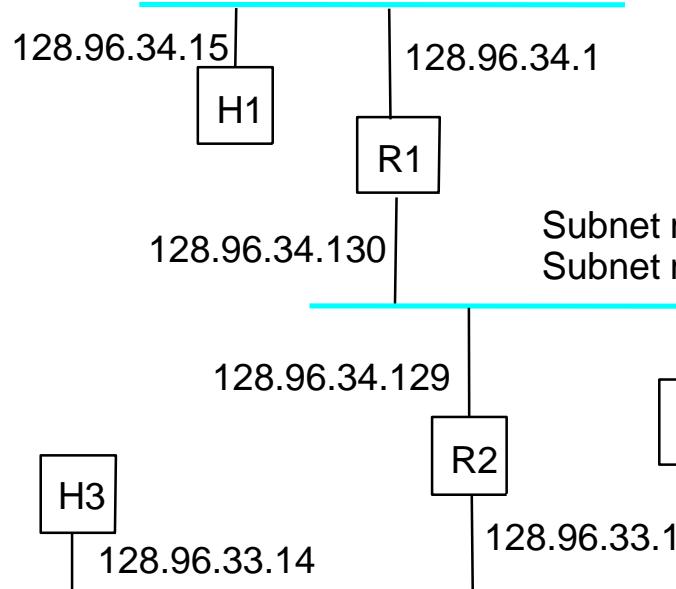


# Task of Routers



# Subnet Example

Subnet mask: 255.255.255.128  
 Subnet number: 128.96.34.0



Subnet mask: 255.255.255.128  
 Subnet number: 128.96.34.128

Forwarding table at router R1

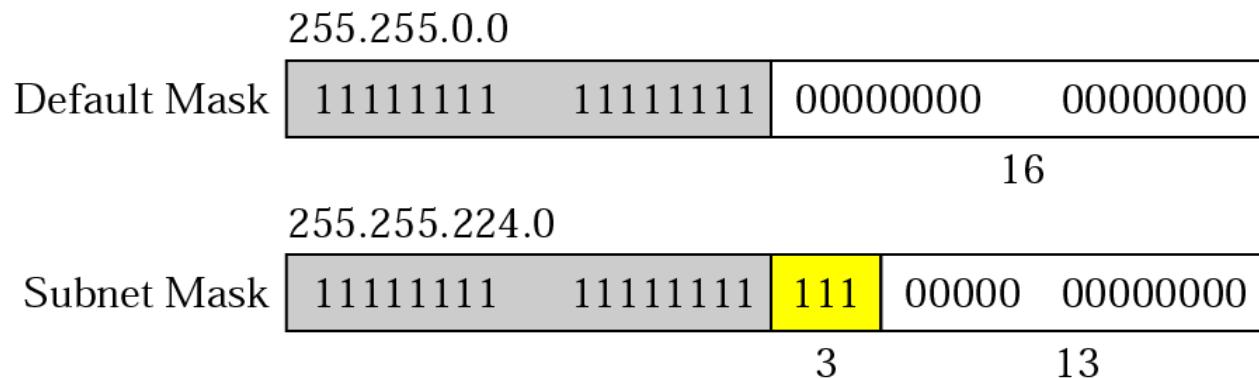
Subnet mask: 255.255.255.0  
 Subnet number: 128.96.33.0

Subnet Number	Subnet Mask	Next Hop
128.96.34.0	255.255.255.128	interface 0
128.96.34.128	255.255.255.128	interface 1
128.96.33.0	255.255.255.0	R2

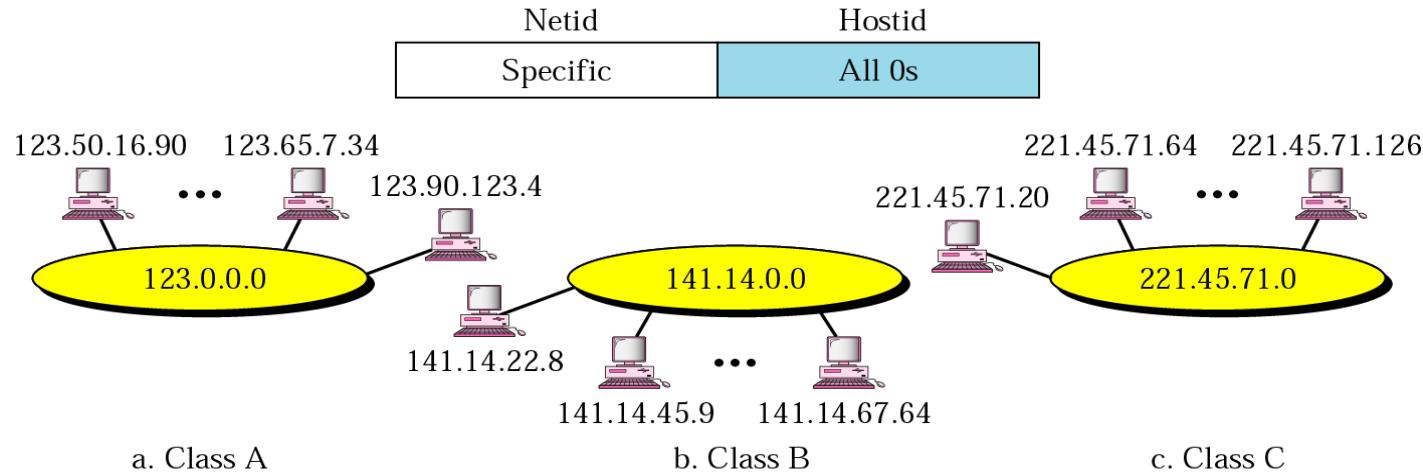
# Default Masks

Class	In Binary	Dotted-Decimal	Using Slash
A	11111111 00000000 00000000 00000000	255.0.0.0	/8
B	11111111 11111111 00000000 00000000	255.255.0.0	/16
C	11111111 11111111 11111111 00000000	255.255.255.0	/24

- Subnet Masks



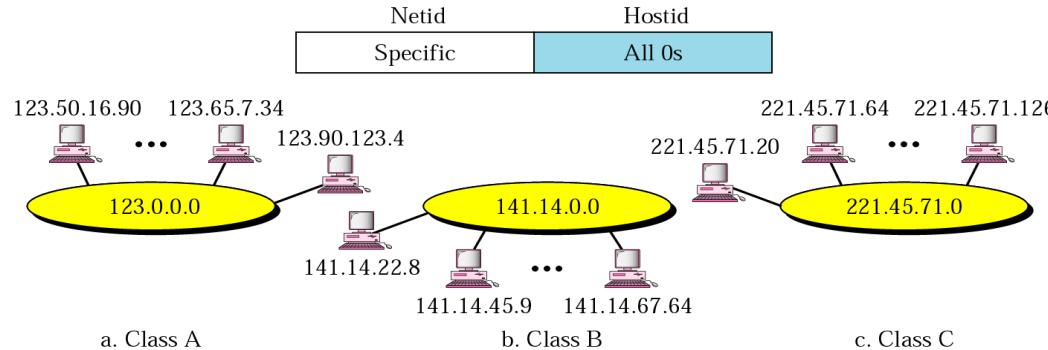
# Classful Addresses & NetworkID



- Classful Addresses:

Class	Networks	Addresses
A	126	16,777,214
B	16,382	65,534
C	2,097,152	254

# Inefficiency of Classful Addresses



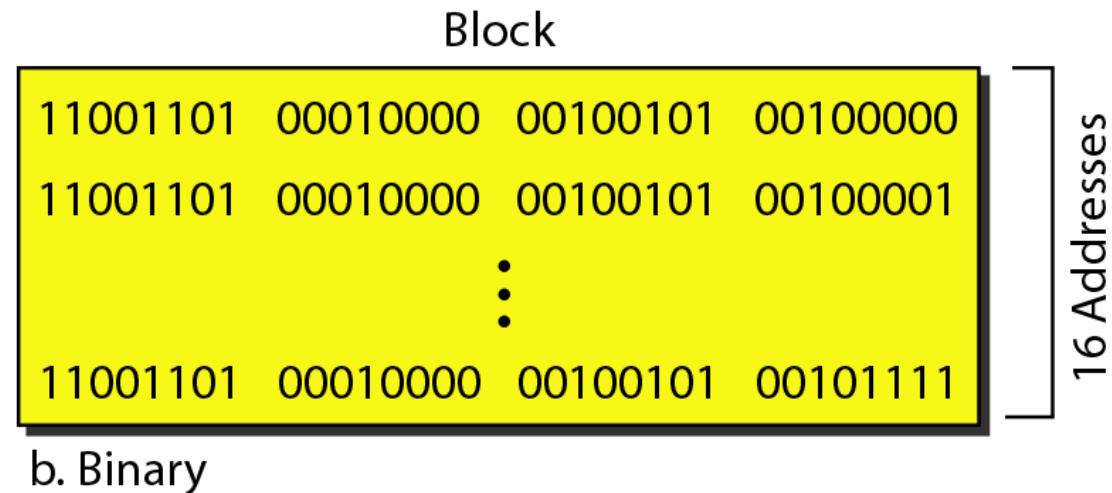
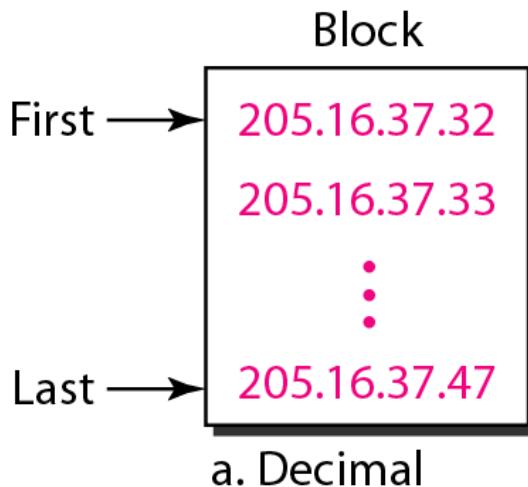
- Classful Addresses:

Class	Networks	Addresses
A	126	16,777,214
B	16,382	65,534
C	2,097,152	254

- Inefficient use of Hierarchical Address Space
  - Class C with 2 hosts ( $2/254 = 0.78\%$  efficient)
  - Class B with 256 hosts ( $256/65534 = 0.39\%$  efficient)

# Classless Inter-Domain Routing(CIDR)

- Allow address space to be divided into blocks of addresses
  - only limited to the power of 2
- Notation as decimal number of the significant bits e.g.  
134.226.36.0 /29
- 205.16.37.32/28
  - 32 bits – 28 bits are static - 4 bits are varied



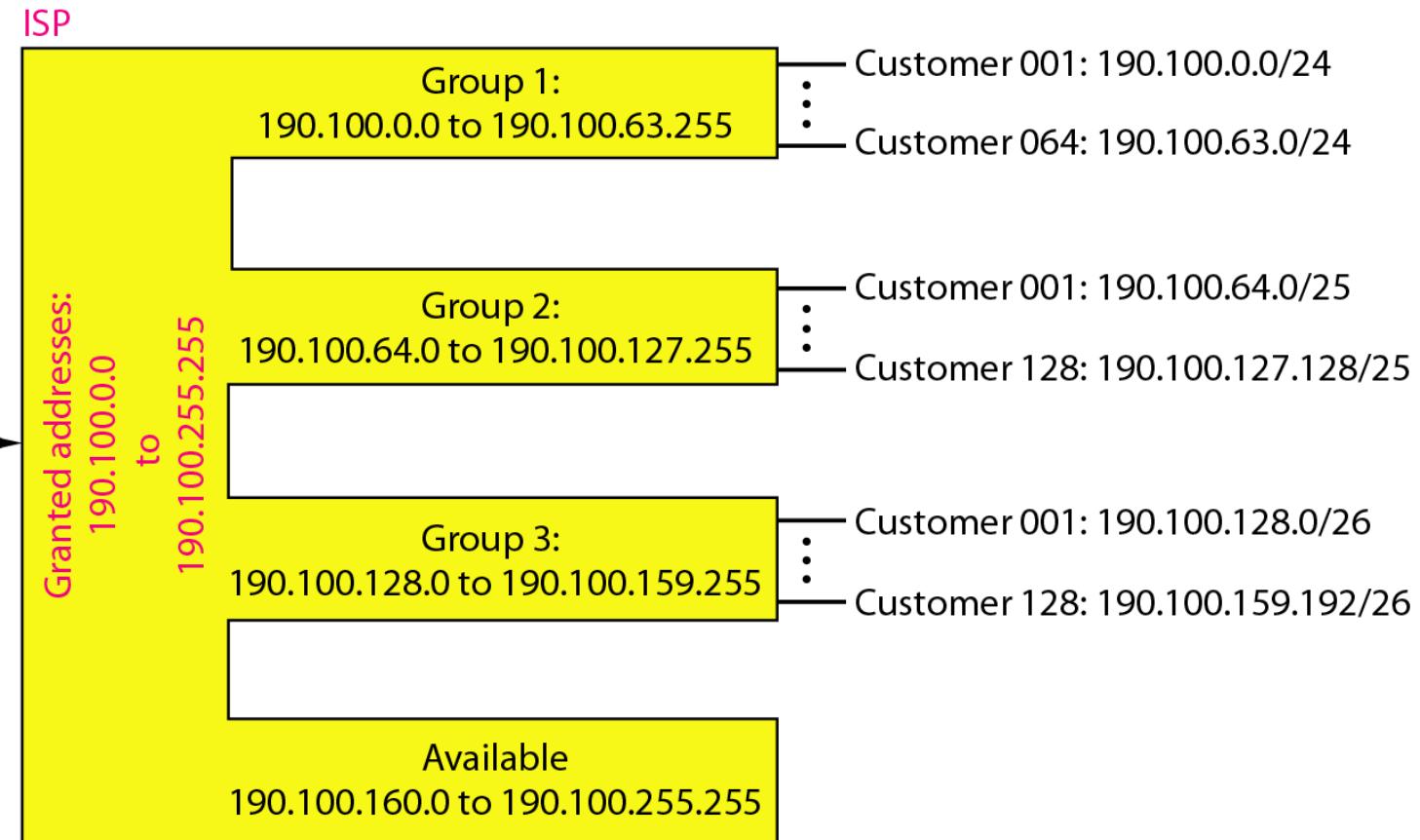
# Classless Inter-Domain Routing(CIDR)

	Dotted Decimal	32-bit binary equivalent
Lowest	128.211.168.0	10000000 11010011 10101 <b>000 00000000</b>
Highest	128.211.175.255	10000000 11010011 10101 <b>111 11111111</b>

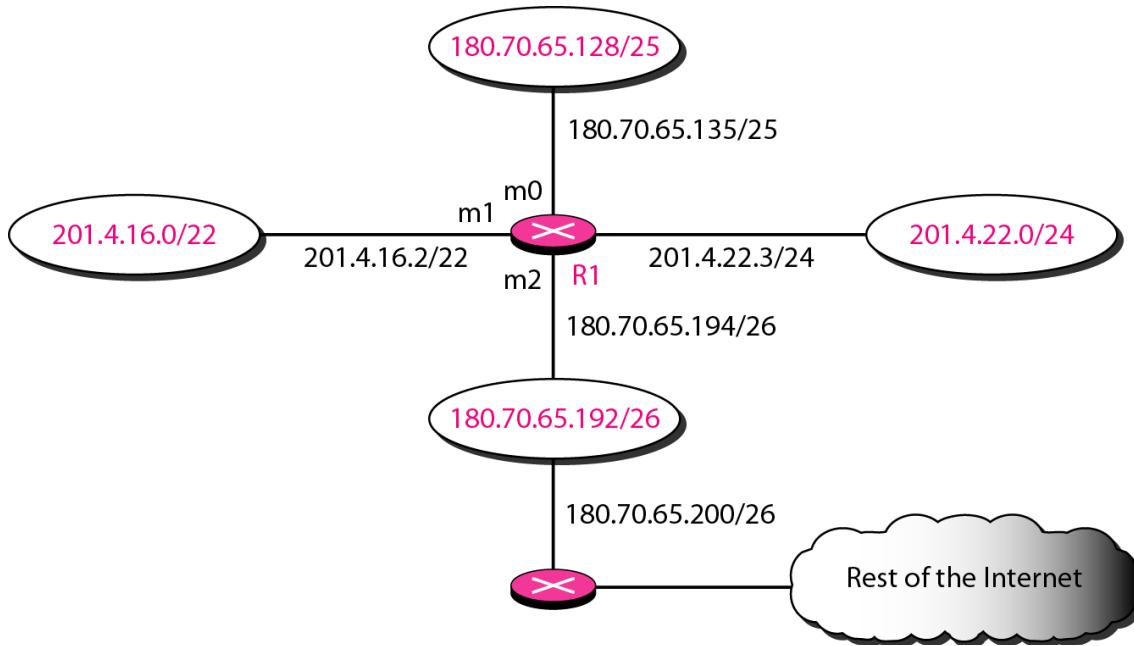
= 128.211.168.0/21

- Aggregation: For example, class C networks can be combined to larger networks
- /21 = 8 class C = 8 \* 256 addresses = 2048 addresses

# ISPs & Classless Addresses

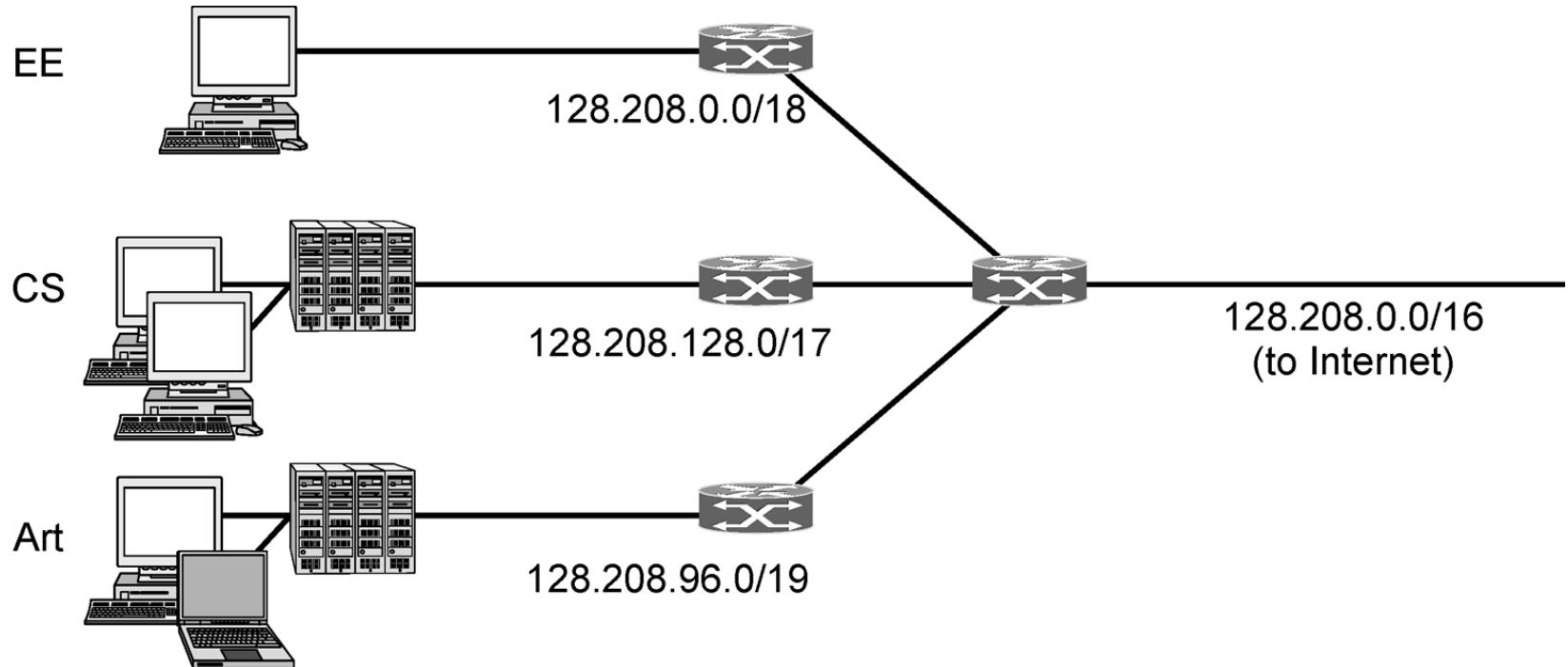


# CIDR & Router

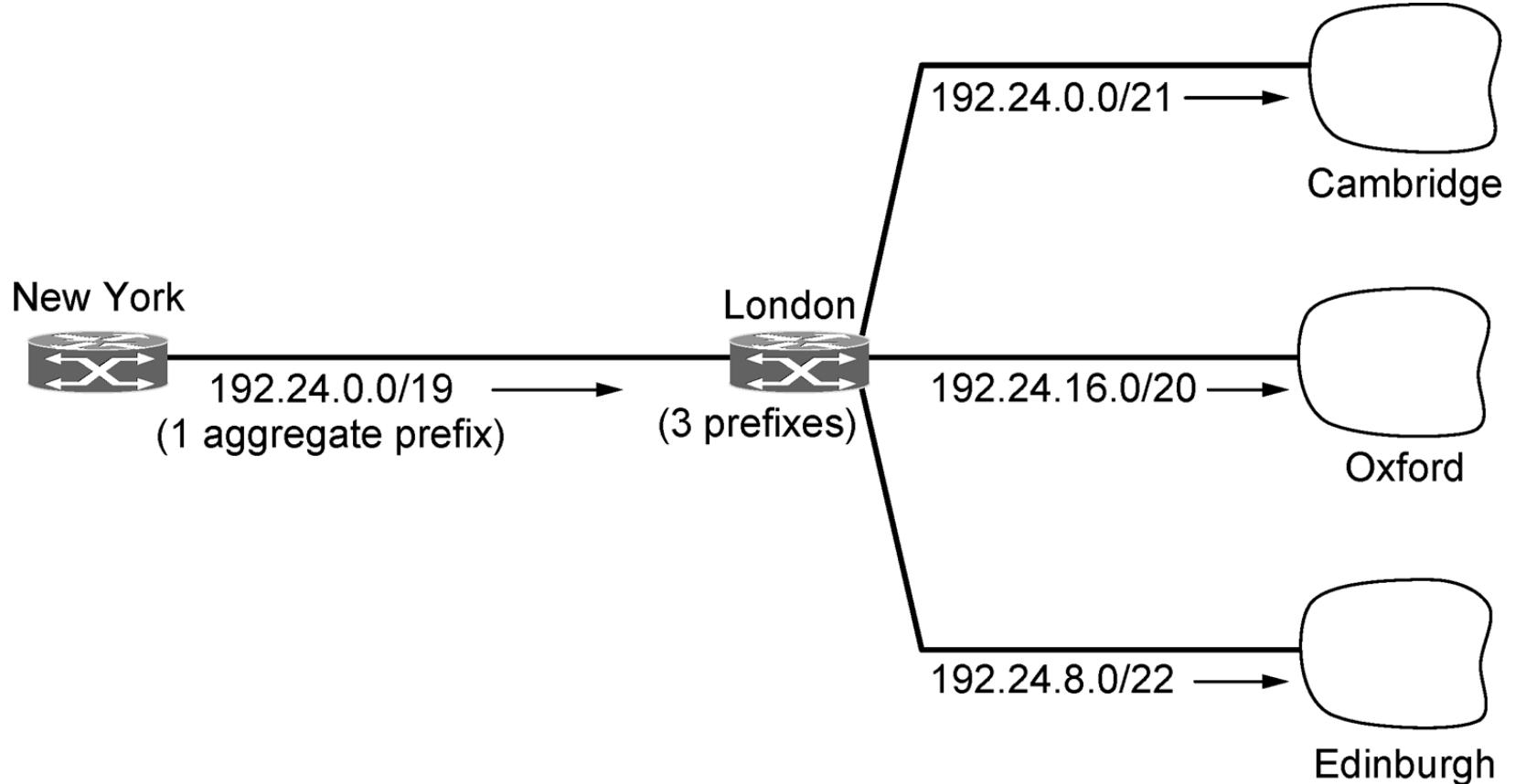


Mask	Network Address	Next Hop	Interface
/26	180.70.65.192	—	m2
/25	180.70.65.128	—	m0
/24	201.4.22.0	—	m3
/22	201.4.16.0	....	m1
Any	Any	180.70.65.200	m2

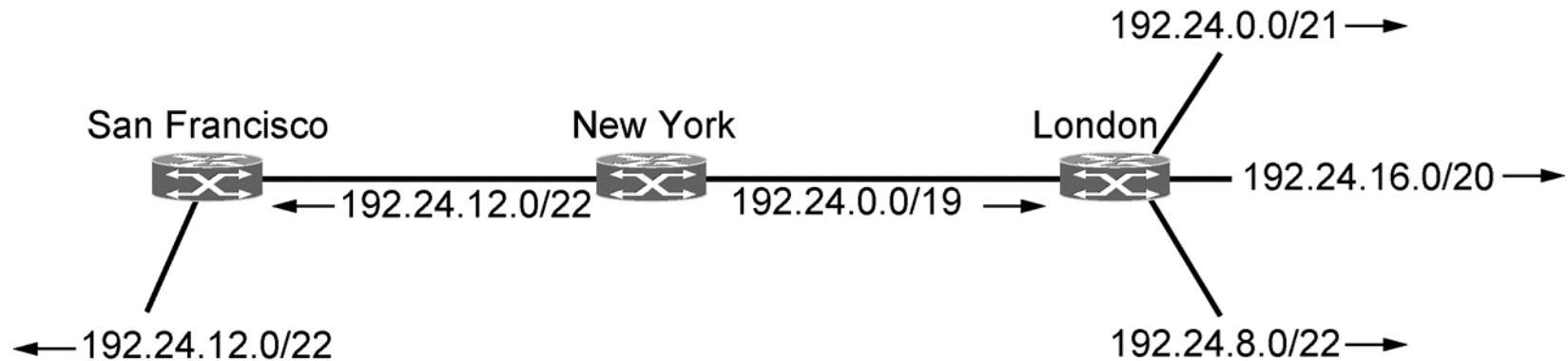
# IP Prefix & Subnets



# Aggregate Prefix

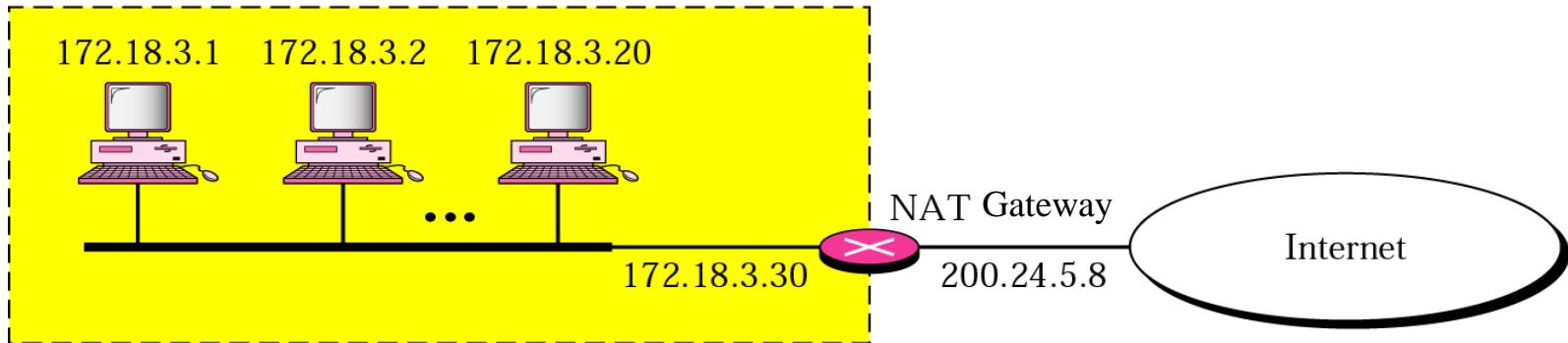


# Longest matching Prefix



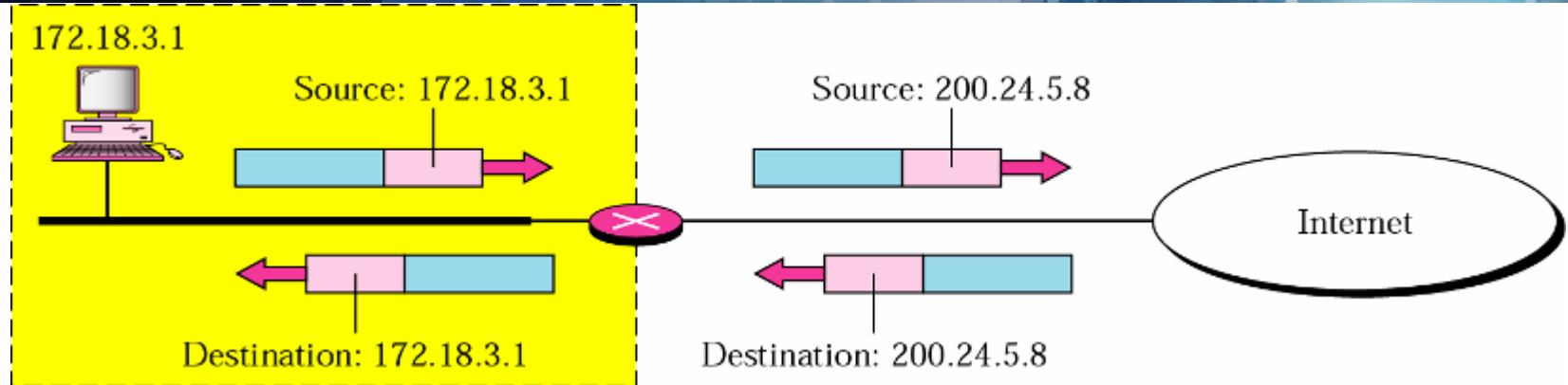
# Network Address Translation (NAT)

Site using private addresses



- NAT gateway translates traffic from the local network to the IP address of the gateway
- Involves processing of outgoing & incoming packet e.g. translation between addresses, recalculation of checksums, etc

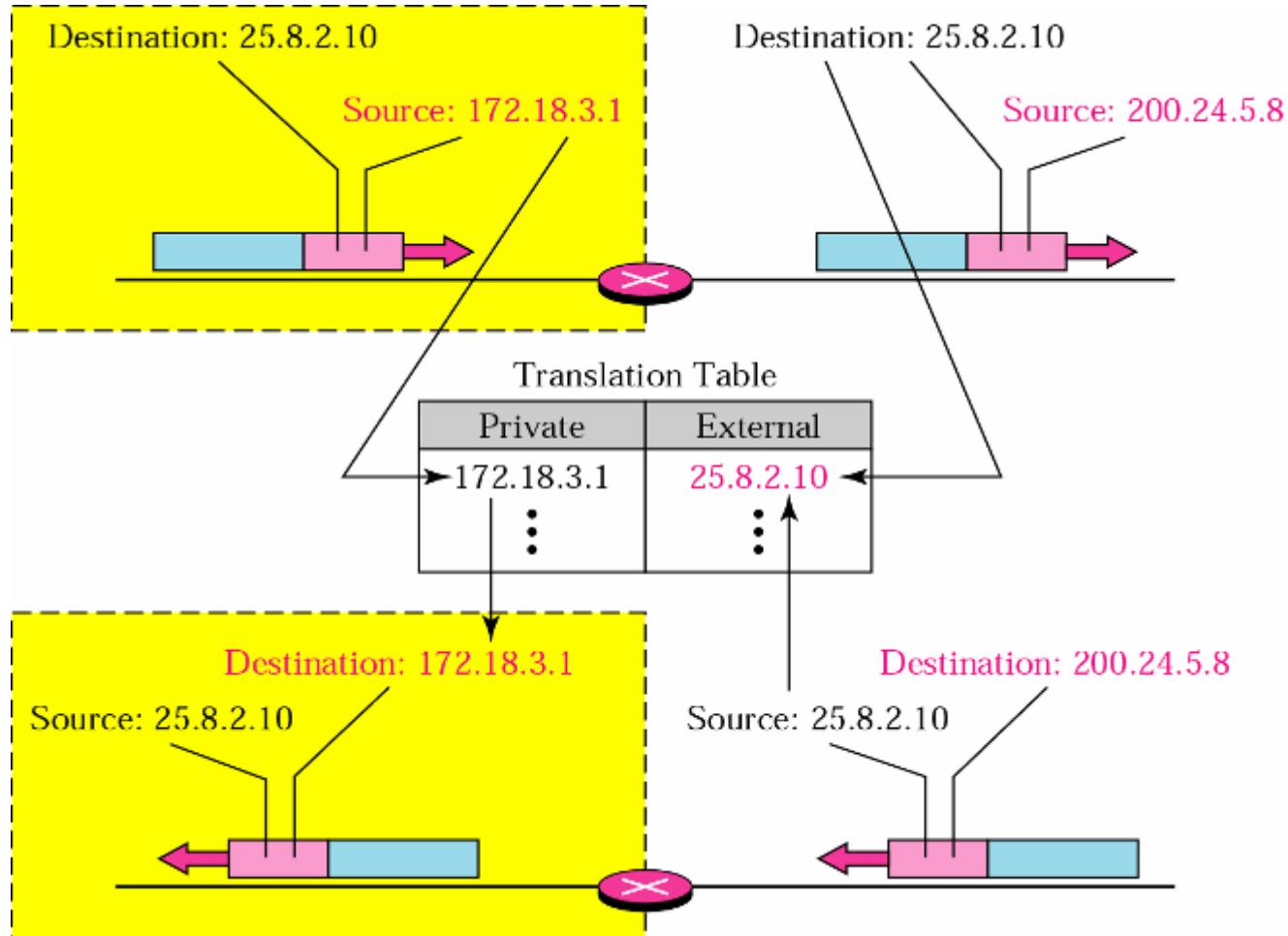
# Address Translation



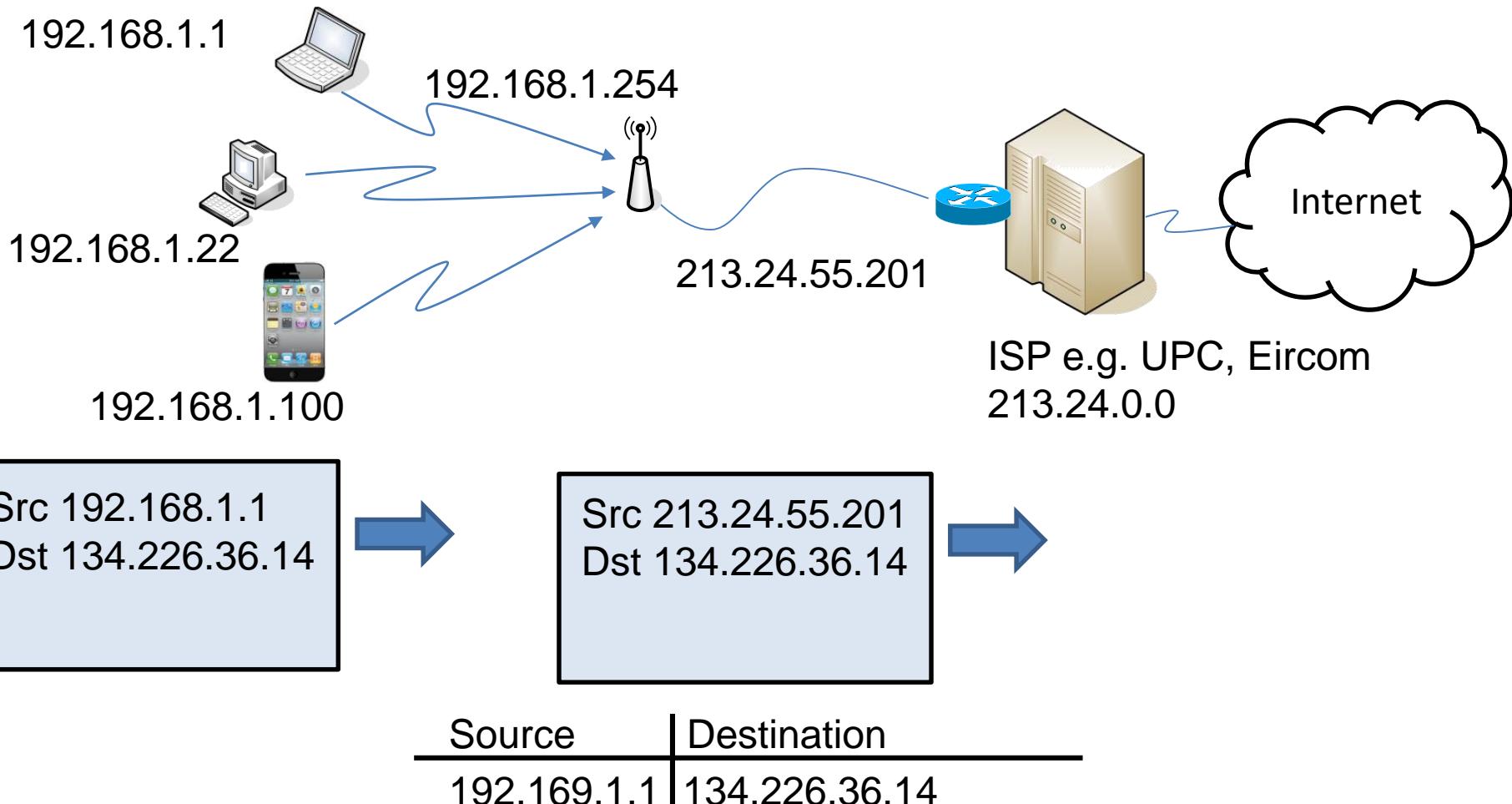
Private Address	Private Port	External Address	External Port	Transport Protocol
172.18.3.1	1400	25.8.3.2	80	TCP
172.18.3.2	1401	25.8.3.2	80	TCP
...	...	...	...	...

- Gateway maintains table to match incoming and outgoing packets; including IDs for applications

# Example for NAT



# NAT Example

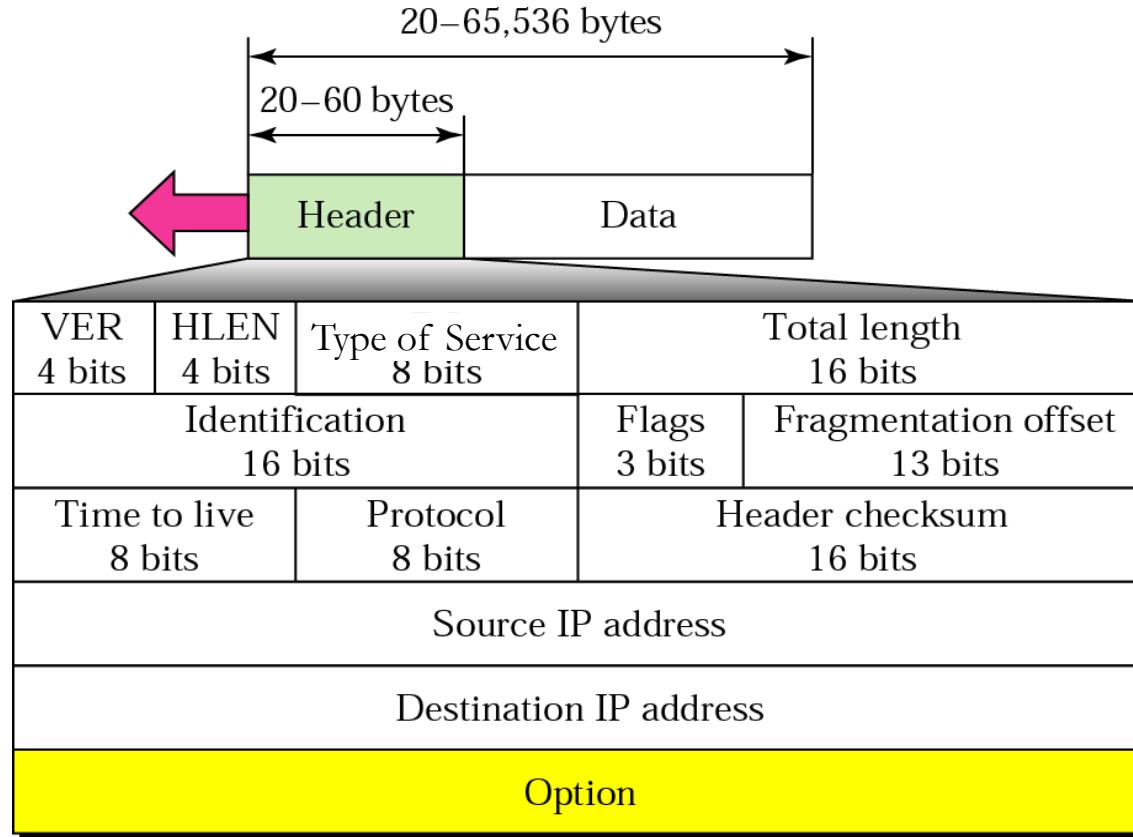


# Summary: Addresses

- 32-bit number / Dotted decimal notation
- IP addresses are unique and universal
  - Exception: Private Addresses
- Classful addresses
  - Classes A, B, and C for networks, D for multicast
  - Routing on Network IDs
- Subnetting + Netmasks
  - Dealing with scale in local networks
- Classless Inter-Domain Routing (CIDR)
  - / notation – significant bits of address
- Network address translation (NAT)

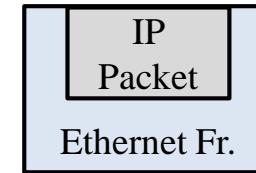
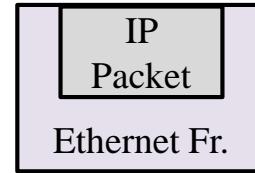
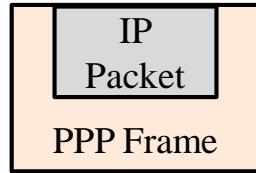
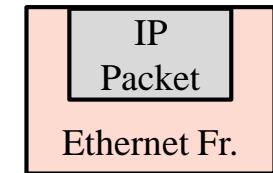
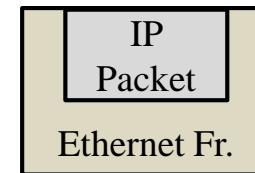
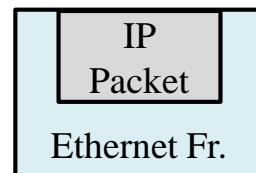
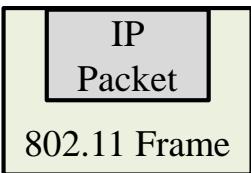
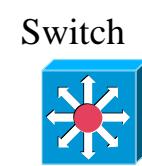
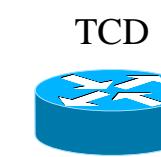


# IP Datagram

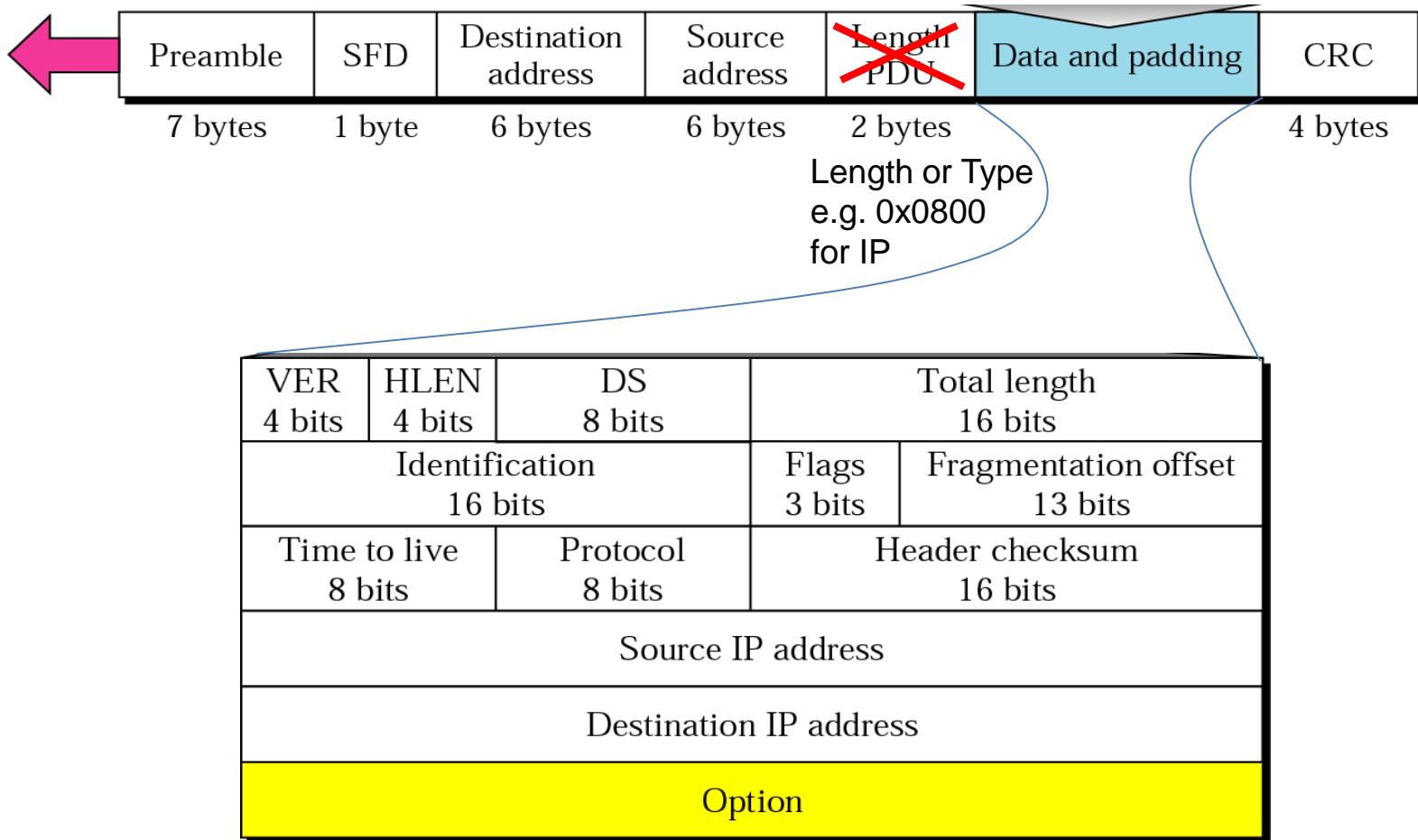


- The total length field defines the total length of the datagram including the header.

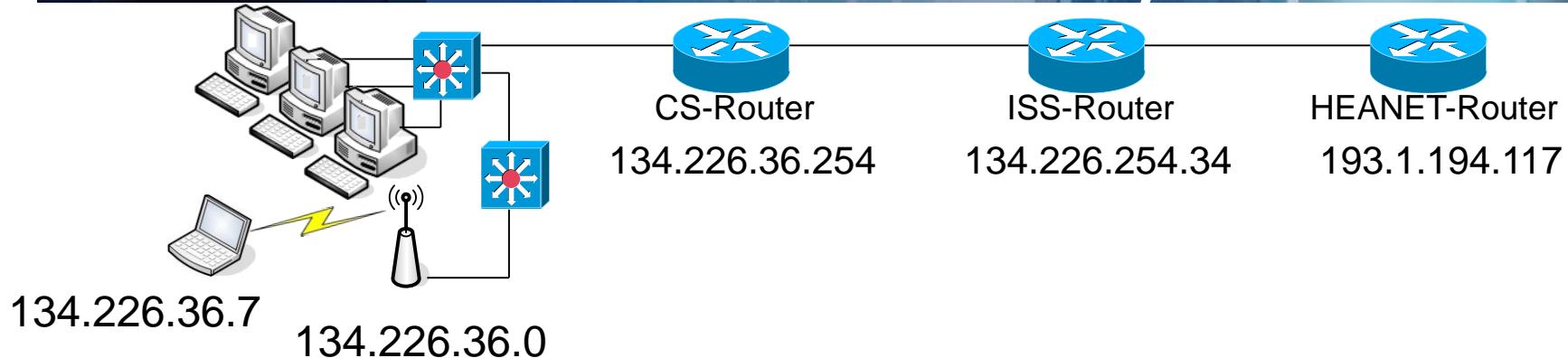
# Encapsulation



# Ethernet & IP



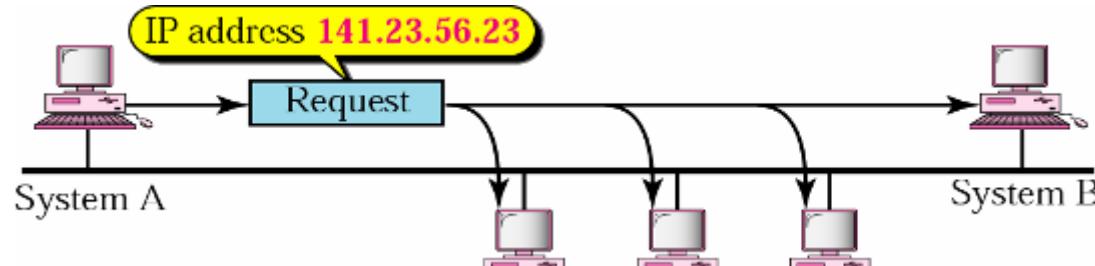
# Default Gateway



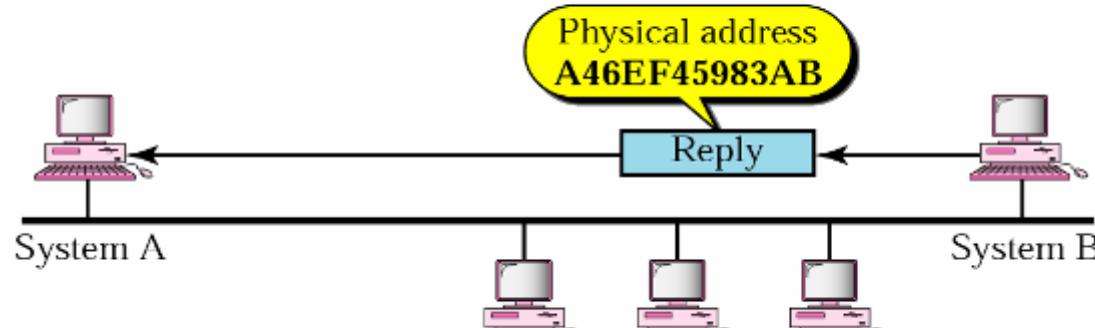
Subnet	Gateway	Netmask	Interface
134.226.36.0	0.0.0.0	255.255.0.0	eth1
0.0.0.0	134.226.36.254	0.0.0.0	eth1

- All nodes within the subnet can communicate directly with each other
- All communication with nodes in other networks passes through the default gateway e.g. router 134.226.36.254

# Address Resolution Protocol (ARP)



a. ARP request is broadcast



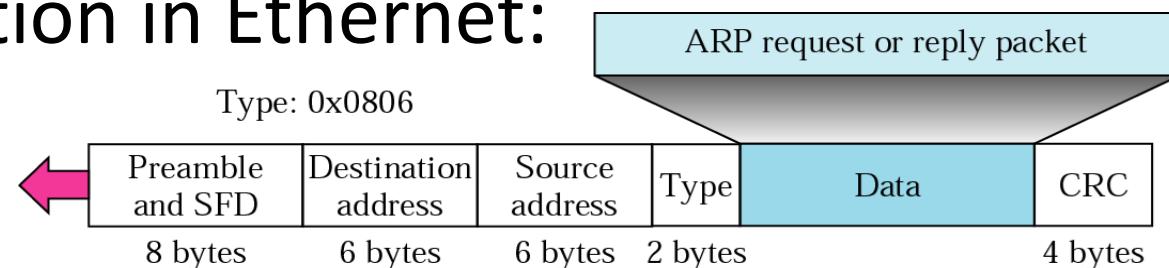
b. ARP reply is unicast

- Association between hardware address and IP address

# ARP Packet

Hardware Type		Protocol Type
Hardware length	Protocol length	Operation Request 1, Reply 2
Sender hardware address (For example, 6 bytes for Ethernet)		
Sender protocol address (For example, 4 bytes for IP)		
Target hardware address (For example, 6 bytes for Ethernet) (It is not filled in a request)		
Target protocol address (For example, 4 bytes for IP)		

- Encapsulation in Ethernet:



# “Who has”-Packet

```
ARP      Who has 192.168.1.2? Tell 192.168.1.5

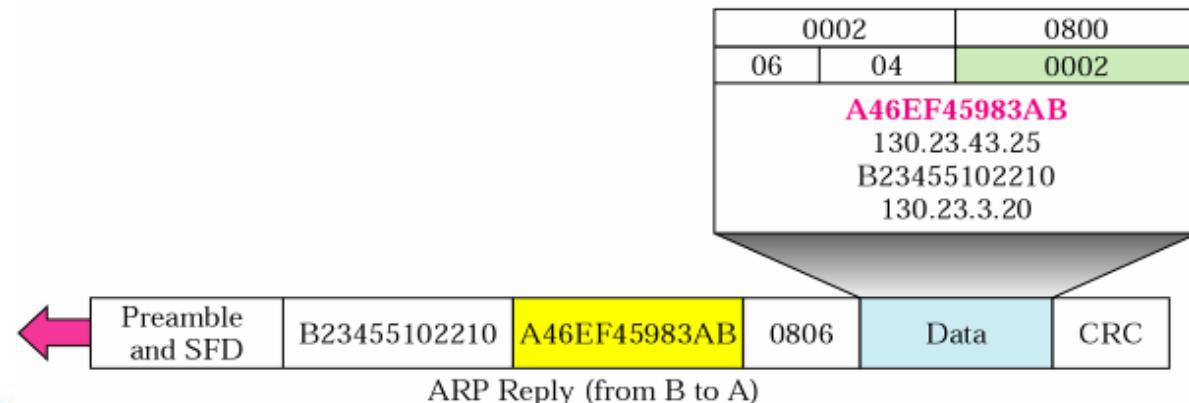
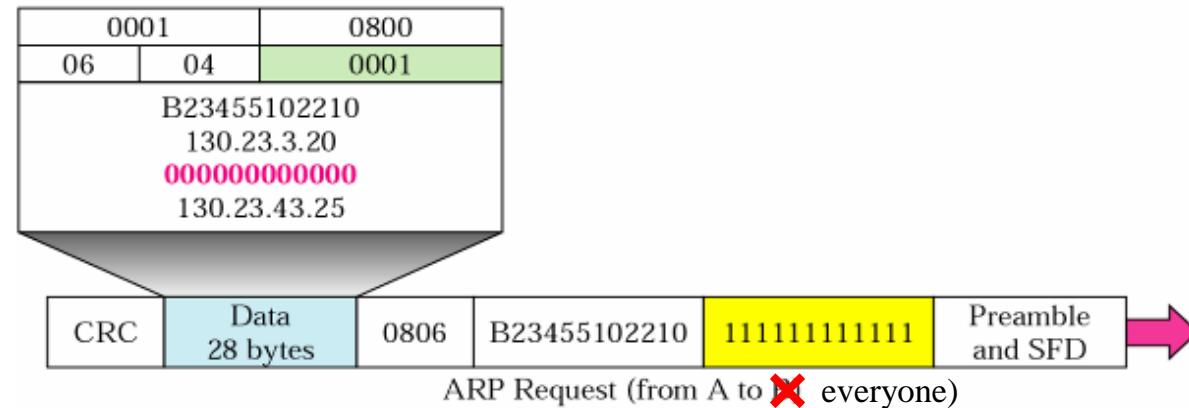
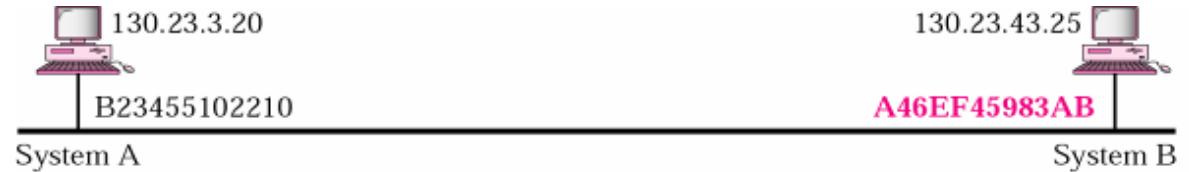
Frame 34 (42 bytes on wire, 42 bytes captured)
Arrival Time: Nov 27, 2005 02:30:26.608147000
[Time delta from previous packet: 59.766200000 seconds]
[Time since reference or first frame: 176.316976000 seconds]
Frame Number: 34
Packet Length: 42 bytes
Capture Length: 42 bytes
[Protocols in frame: eth:arp]
Ethernet II, src: 00:0f:b5:96:19:e5 (00:0f:b5:96:19:e5), dst: ff:ff:ff:ff:ff:ff
Destination: ff:ff:ff:ff:ff:ff (ff:ff:ff:ff:ff:ff)
Source: 00:0f:b5:96:19:e5 (00:0f:b5:96:19:e5)
Type: ARP (0x0806)
Address Resolution Protocol (request)
Hardware type: Ethernet (0x0001)
Protocol type: IP (0x0800)
Hardware size: 6
Protocol size: 4
Opcode: request (0x0001)
Sender MAC address: 00:0f:b5:96:19:e5 (00:0f:b5:96:19:e5)
Sender IP address: 192.168.1.5 (192.168.1.5)
Target MAC address: 00:00:00:00:00:00 (00:00:00:00:00:00)
Target IP address: 192.168.1.2 (192.168.1.2)

0000  ff ff ff ff ff 00 0f b5 96 19 e5 08 06 00 01
0010  08 00 06 04 00 01 00 0f b5 96 19 e5 c0 a8 01 05
0020  00 00 00 00 00 00 c0 a8 01 02  ....:....:....:
```

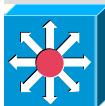
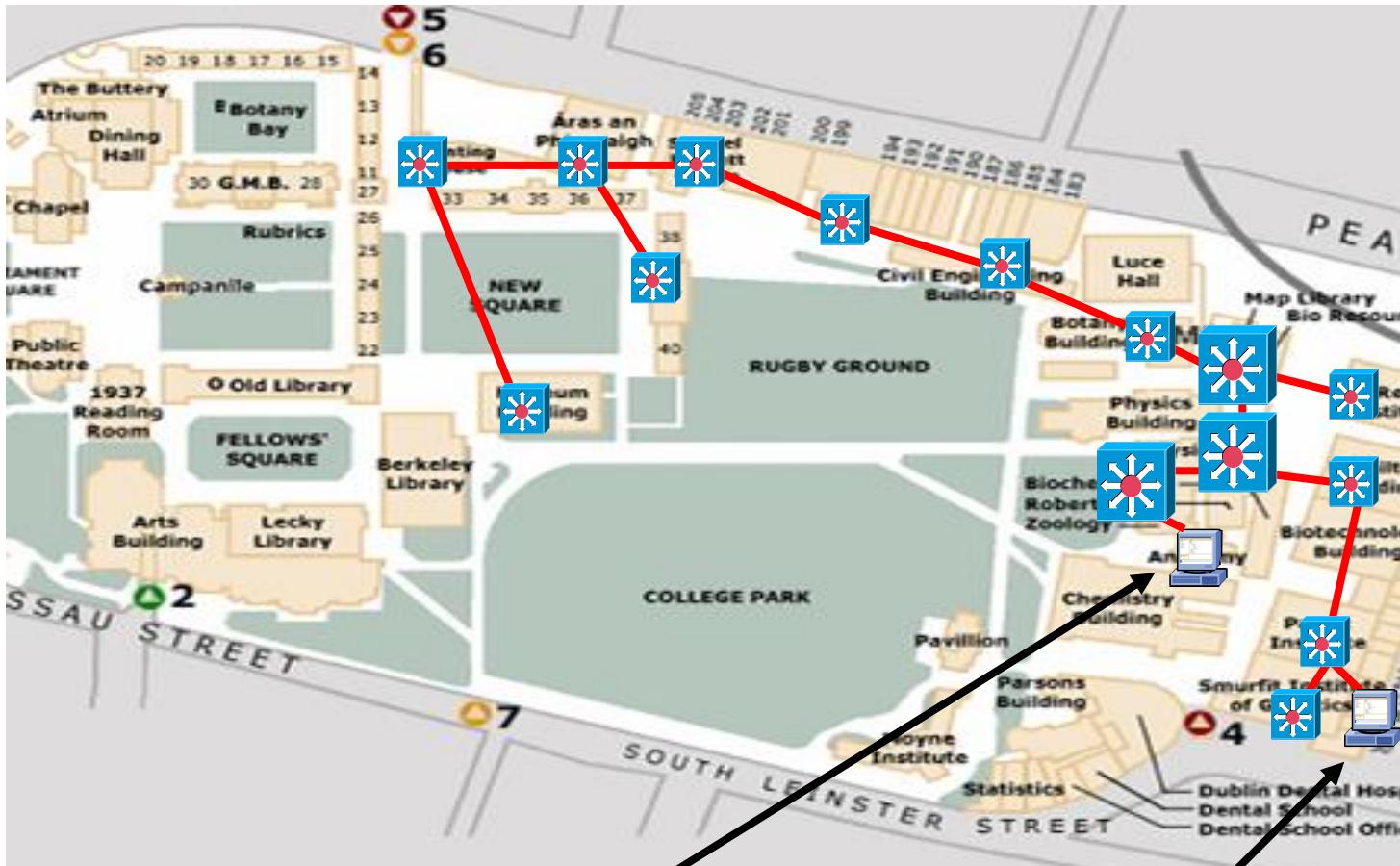
- ARP Request: Who has 192.168.1.2? Tell 192.168.1.5



# ARP Example



# Trinity Network with Switches



Switch

Address: 134.226.38.55  
Subnet: 134.226.0.0

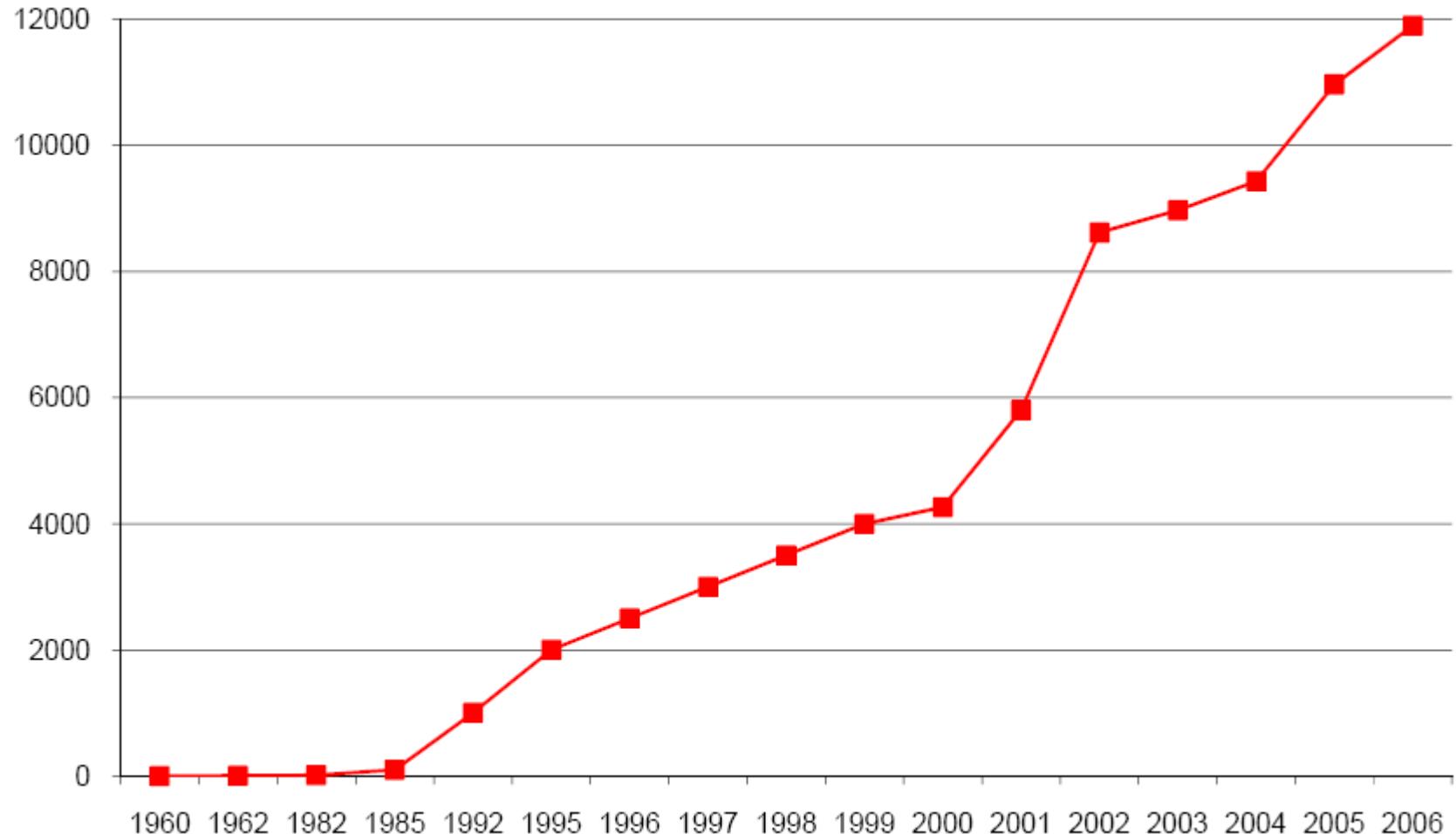
Address: 134.226.32.55  
Subnet: 134.226.0.0



TRINITY COLLEGE DUBLIN  
COLÁISTE NA TRÍONÓIDE, BAILE ÁTHA CLIATH

THE  
UNIVERSITY  
OF DUBLIN

# Number of Computers in College

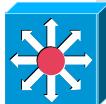
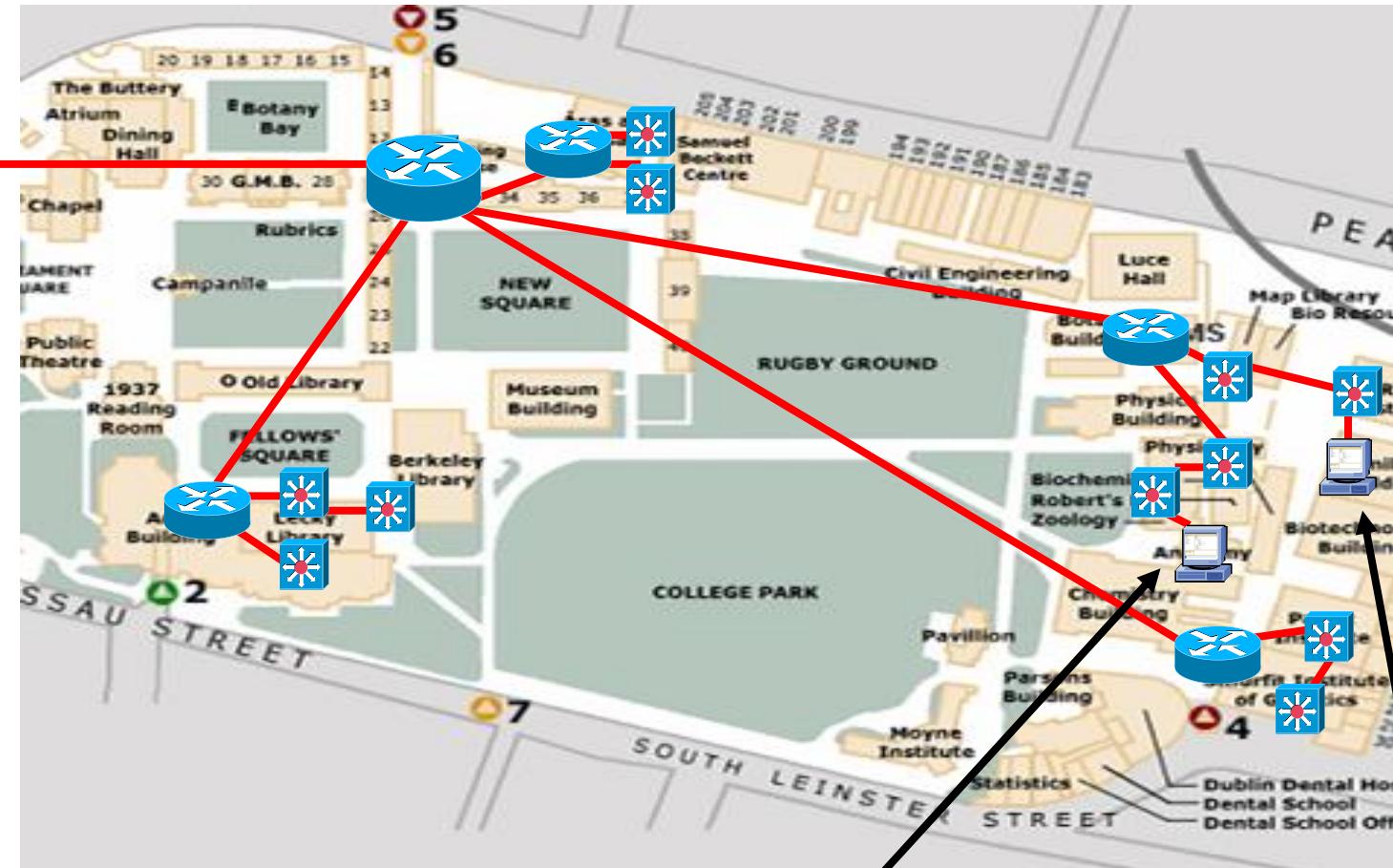


# Broadcast-Storms

- Imagine 1000s of machines sending ARP requests



# Trinity Network with Subnets



Switch



Router



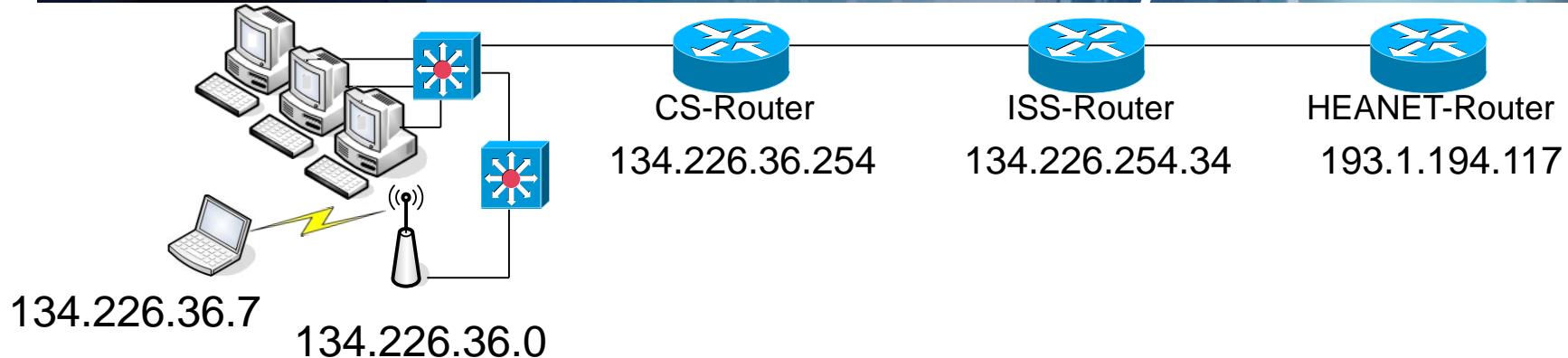
TRINITY COLLEGE DUBLIN  
COLÁISTE NA TRÍONÓIDE, BAILE ÁTHA CLIATH

THE  
UNIVERSITY  
OF DUBLIN

Address: 134.226.38.55  
Subnet: 134.226.38.0

Address: 134.226.32.55  
Subnet: 134.226.32.0

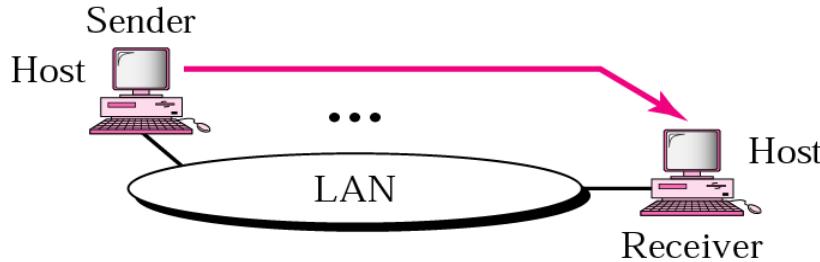
# Default Gateway



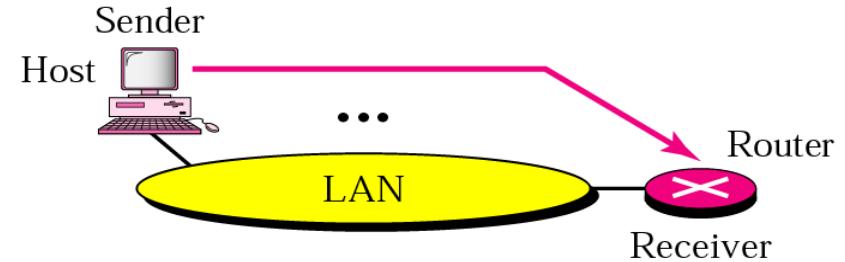
<b>Subnet</b>	<b>Gateway</b>	<b>Netmask</b>	<b>Interface</b>
134.226.36.0	0.0.0.0	255.255.0.0	eth1
0.0.0.0	134.226.36.254	0.0.0.0	eth1

- All nodes within the subnet can communicate directly with each other
- All communication with nodes in other networks passes through the default gateway e.g. router 134.226.36.254

# 4 Cases for ARP



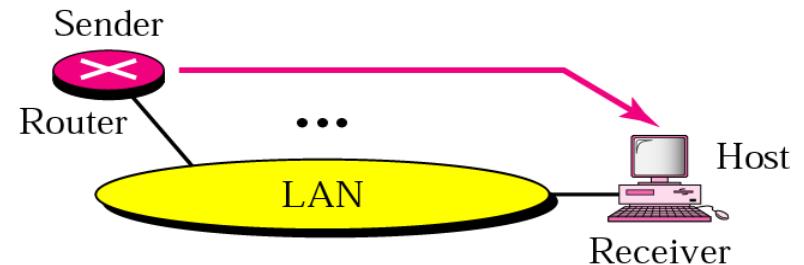
Case 1. A host has a packet to send to another host on the same network.



Case 2. A host wants to send a packet to another host on another network.  
It must first be delivered to the appropriate router.



Case 3. A router receives a packet to be sent to a host on another network.  
It must first be delivered to the appropriate router.



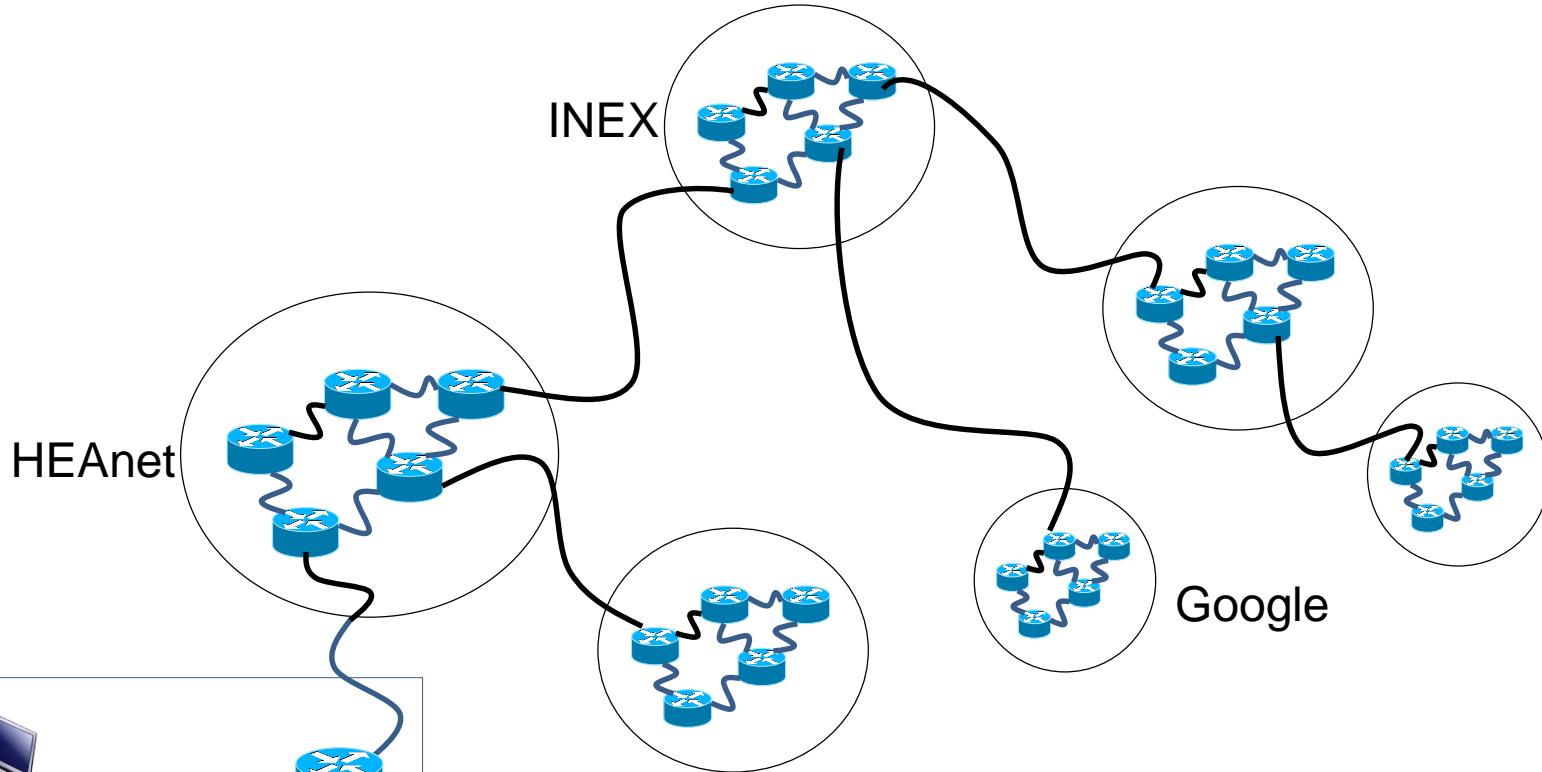
Case 4. A router receives a packet to be sent to a host on the same network.

## Summary: ARP

- “Who-has?” packet
- Associating IP addresses with hardware addresses
- Based on broadcast



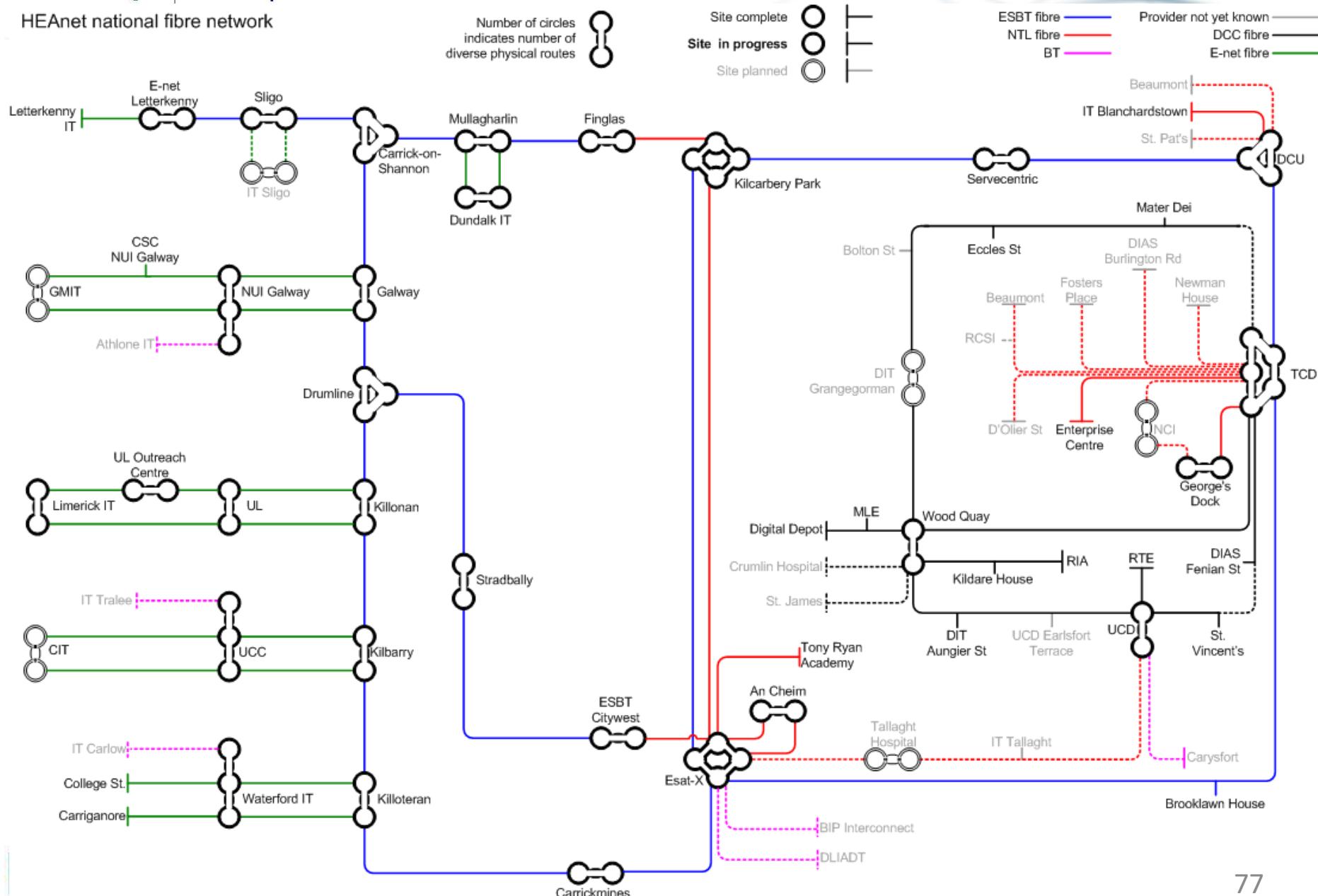
# Internet = Network of Networks



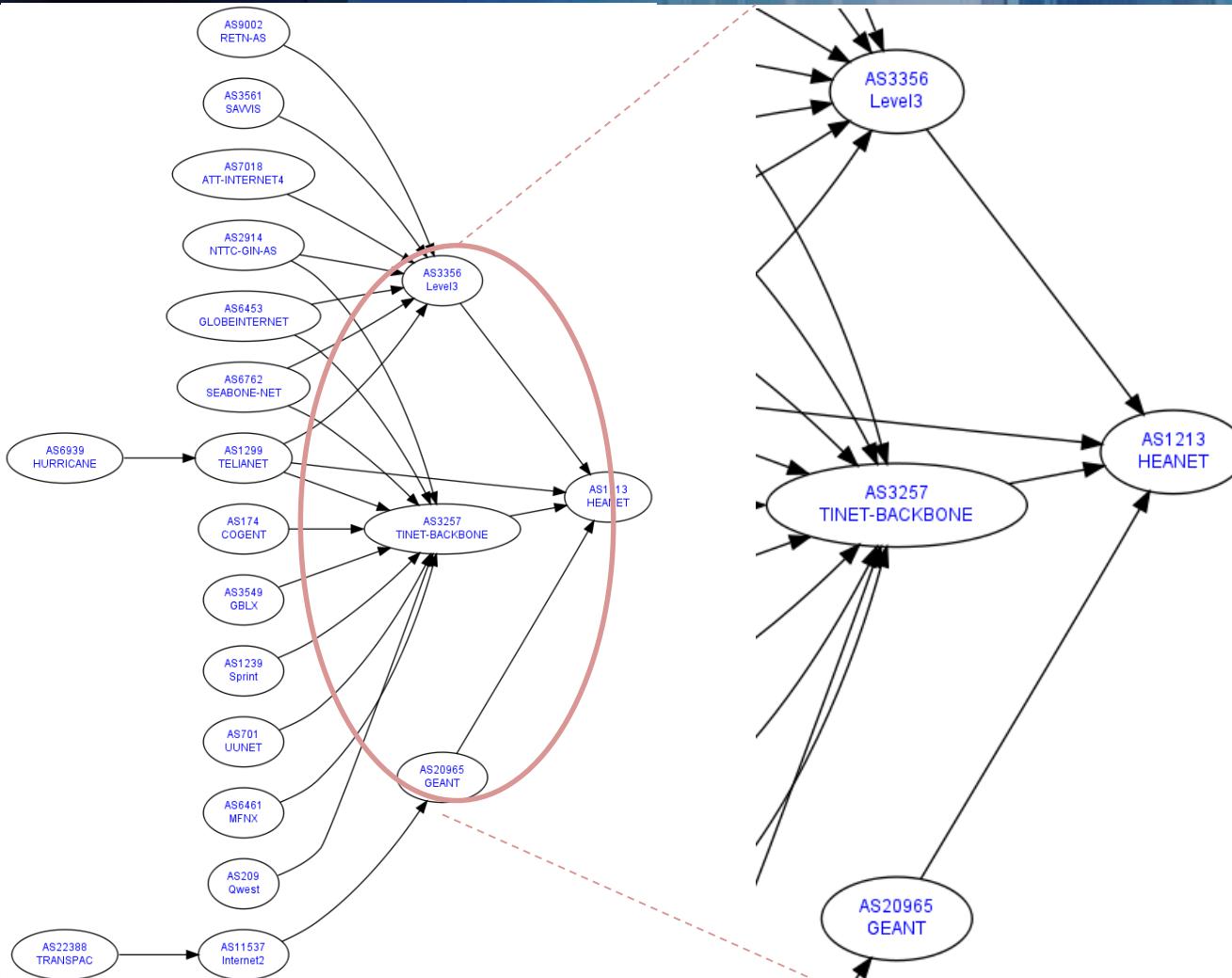
TCD/Edge Network

*school of* Computer Science & Statistics

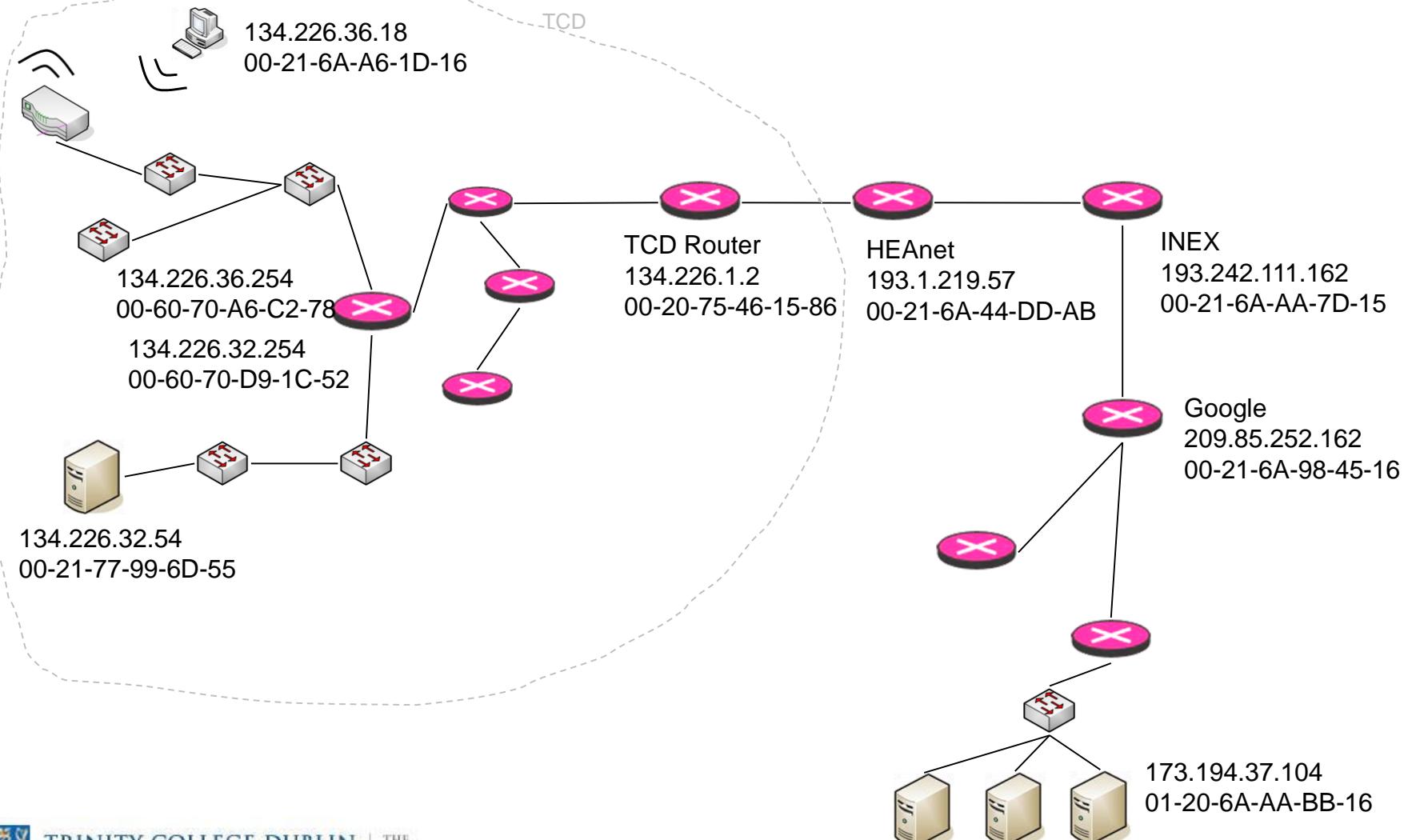
HEAnet national fibre network



# AS1213 - HEANET

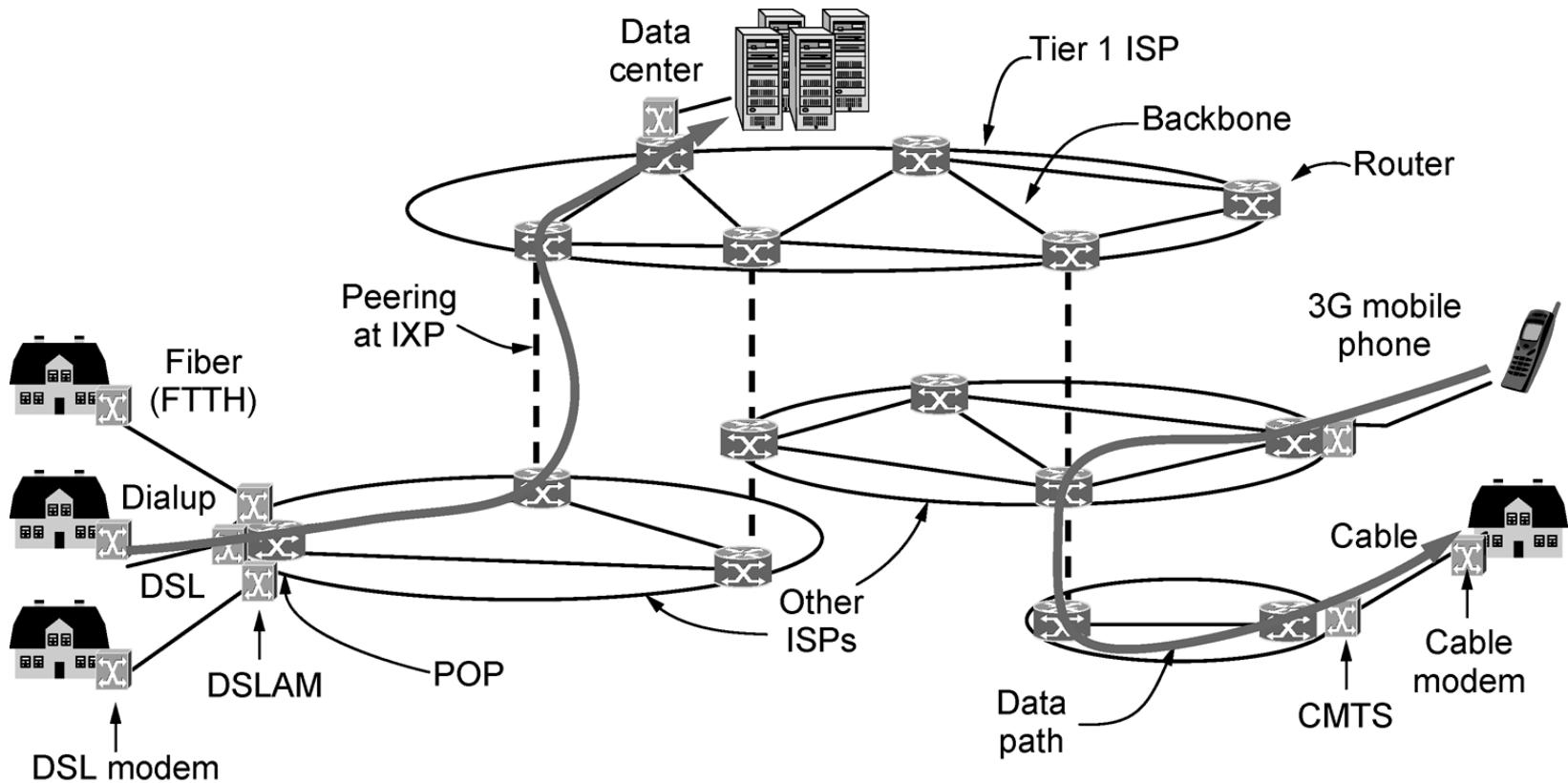


# Network Layer Scenario





# Overview of the Internet Architecture



# Collection of Networks

