

1 **CS370 Programming Languages (Zaring)**  
 2 **Spring 2020**  
 3 **Assignment 3**  
 4 **Due by the beginning of lecture on Thursday, March 5**

5  
 6 **Description:**

7 Write the Scheme procedures described below. Unless told otherwise, you may use only the  
 8 primitive procedures and expressions used in lecture, the example procedures defined in lecture,  
 9 and the primitive procedures and expressions used and defined in chapters 1-8 of *TLIS*. If you're  
 10 uncertain about using some procedure (e.g., one you discovered via the DrRacket help system),  
 11 ask if it's okay to use it: **contrary to the usual belief, in this course, it's better to ask for**  
 12 **permission before the fact than to ask for forgiveness after the fact.** Unless told otherwise in a  
 13 particular problem, you may define and use auxiliary/helper procedures in your solutions.

14  
 15 When writing operators (i.e., a procedure that returns a funval),

- 16  
 17 (1) Avoid having the funval returned by the operator depend on any top-level user-defined  
 18 procedures  
 19 (2) Think hard before using `letrec` (although you may find it's necessary in some cases)

20  
 21 For each procedure you write, include a header comment that (at the very least)

- 22  
 23 • States the purpose of the procedure  
 24 • Gives any/all pre-conditions for the procedure (i.e., a description of any special properties  
 25 the actual parameters must have in order for the procedure to work correctly)  
 26 • Gives a big-O statement of the procedure's worst-case asymptotic runtime

27  
 28 • (curried-binary *binaryProc*)

29 Assuming *binaryProc* is a binary procedure, returns a unary procedure that when applied to a  
 30 value *x* returns a unary procedure that when applied to a value *y* returns the value of  
 31 *binaryProc* applied to *x* and *y*. For example,

32  
 33 `((curried-binary cons) 'a ' (b))` returns (a b)  
 34 `((curried-binary +) 10) 2)` returns 12  
 35 `((curried-binary >) 1) 2)` returns #f

36  
 37 • (uncurried-binary *curriedBinaryProc*)

38 Assuming *curriedBinaryProc* is a binary procedure that has been curried (perhaps one that was  
 39 “curried” by `curried-binary`), returns a binary procedure that when applied to a value *x*  
 40 and a value *y* returns the value of the original uncurried version of *curriedBinaryProc* applied  
 41 to *x* and *y*. For example,

42  
 43 `((uncurried-binary (curry-binary cons)) 'a ' (b))` returns (a b)  
 44 `((uncurried-binary (curry-binary +)) 10 2)` returns 12  
 45 `((uncurried-binary (curry-binary >)) 1 2)` returns #f

46

47 • (adjacent-related-grouped *related?* *lis*)  
 48 (Reminiscent of the procedure adjacent-equals-grouped from Assignment 1.)  
 49 Assume *lis* is a list and *related?* is a binary predicate (which should be assumed to be O(1)).  
 50 Returns the list in which all runs of two or more adjacent elements of *lis* for which *related?*  
 51 returns a non-#f value have been grouped into lists. For example,

```

52
53 (adjacent-related-grouped eq? '()) returns ()
54 (adjacent-related-grouped eq? '(a)) returns ((a))
55 (adjacent-related-grouped eq? '(a b)) returns ((a) (b))
56 (adjacent-related-grouped eq? '(a b c)) returns ((a) (b) (c))
57 (adjacent-related-grouped eq? '(a a b c)) returns ((a a) (b) (c))
58 (adjacent-related-grouped eq? '(a b a a b b c a a a b b)) returns
59 ((a) (b) (a a) (b b) (c) (a a a) (b b))
60 (adjacent-related-grouped eq? '(a b a a b b c a a a b b b c c c d)) returns
61 ((a) (b) (a a) (b b) (c) (a a a) (b b b) (c c c) (d))
62 (adjacent-related-grouped (lambda (x y) (not (eq? x y)))
63 '(a b a a b b c a a a b b b c c c d)) returns
64 ((a b a) (a b) (b c a) (a) (a b) (b) (b c) (c) (c d))
65 (adjacent-related-grouped (lambda (x y) #f) '(1 1 2 1 2 3 1 2 3 4)) returns
66 ((1) (1) (2) (1) (2) (3) (1) (2) (3) (4))
67 (adjacent-related-grouped (lambda (x y) #t) '(1 1 2 1 2 3 1 2 3 4)) returns
68 ((1 1 2 1 2 3 1 2 3 4))
69 (adjacent-related-grouped < '(1 1 2 1 2 3 1 2 3 4)) returns
70 ((1) (1 2) (1 2 3) (1 2 3 4))
71 (adjacent-related-grouped <= '(1 1 2 1 2 3 1 2 3 4)) returns
72 ((1 1 2) (1 2 3) (1 2 3 4))
73 (adjacent-related-grouped
74 (lambda (x y) (<= (car x) (car y))))
75 '( (1 a) (1 b) (2 c) (1 d) (2 e) (3 f) (1 g) (2 h) (3 i) (4 j)))
76 returns (((1 a) (1 b) (2 c)) ((1 d) (2 e) (3 f))
77 ((1 g) (2 h) (3 i) (4 j)))
78

```

79 NOTE: For the following procedure, you may not define any top-level or local auxiliary  
 80 procedures as part of your answer nor may the returned procedure *p* (described below) call any  
 81 top-level or local user-defined auxiliary procedures (other than, if applicable, *binaryProc*).

82  
 83 • (rreducer *binaryProc* *unaryProc* *zeroaryProc*)  
 84 Assume *binaryProc* is a binary procedure, *unaryProc* is a unary procedure, and *zeroaryProc* is  
 85 a *thunk* (a procedure of zero parameters). Returns a unary procedure *p* that “reduces” a list  
 86 using *binaryProc*, treating *binaryProc* as if it’s right-associative. That is, applying *p* to a list  
 87 (*x*<sub>1</sub> *x*<sub>2</sub> *x*<sub>3</sub> *x*<sub>4</sub> *x*<sub>5</sub>) would return a value equivalent to the application

```

88
89 (binaryProc x1 (binaryProc x2 (binaryProc x3 (binaryProc x4 x5))))
90

```

91 applying *p* to the list (*x*<sub>1</sub>) would return a value equivalent to the application

```

92
93 (unaryProc x1)
94

```

95 and applying *p* to the list () would return a value equivalent to the application

```

96
97 (zeroaryProc)

```

For example,

```
((rreducer - (lambda (x) x) (lambda () 0)) '()) returns 0
((rreducer - (lambda (x) x) (lambda () 0)) '(10)) returns 10
((rreducer - (lambda (x) x) (lambda () 0)) '(10 20 30 40)) returns -20
((rreducer cons (lambda (x) x) (lambda () '())) '()) returns ()
((rreducer cons (lambda (x) x) (lambda () '())) '(a)) returns a
((rreducer cons (lambda (x) x) (lambda () '())) '(a b c d)) returns
  (a b c . d)
```

NOTE: For the following procedure, you may not define any top-level or local auxiliary procedures as part of your answer nor may the returned procedure *p* (described below) call any top-level or local user-defined auxiliary procedures (other than, if applicable, *binaryProc*).

• (lreducer *binaryProc unaryProc zeroaryProc*)

Like *rreducer*, but instead returns a unary procedure *p* that “reduces” a list using *binaryProc*, treating *binaryProc* as if it’s left-associative. For example,

```
((lreducer - (lambda (x) x) (lambda () 0)) '()) returns 0
((lreducer - (lambda (x) x) (lambda () 0)) '(10)) returns 10
((lreducer - (lambda (x) x) (lambda () 0)) '(10 20 30 40)) returns -80
((lreducer cons (lambda (x) x) (lambda () '())) '()) returns ()
((lreducer cons (lambda (x) x) (lambda () '())) '(a)) returns a
((lreducer cons (lambda (x) x) (lambda () '())) '(a b c d)) returns
  (((a . b) . c) . d)
```

• (subst-every-other-sf *old new los+ succeed*)

Works like *subst-every-other* from Assignment 2, but is written using the success-fail style; however, since failure can’t really occur in this case, no failure thunk is needed. *subst-every-other-sf* succeeds with two values: a list like the arbitrarily-complex list of symbols *los+*, but with every other occurrence (as read from left to right, starting with the leftmost occurrence) of the symbol *old* replaced by an occurrence of the symbol *new* and a Boolean value that’s #f if the last occurrence of *old* wasn’t replaced by *new* but #t if the last occurrence of *old* was replaced by *new*. Since *succeed* succeeds with two values, all the success procedures will be binary procedures. For example,

```
(subst-every-other-sf 'a 'b '() (lambda (result replaced) result)) returns ()
(subst-every-other-sf 'a 'b '() (lambda (result replaced) replaced)) returns
  #f
(subst-every-other-sf 'a 'b '(a) (lambda (result replaced) result)) returns
  (b)
(subst-every-other-sf 'a 'b '(a) (lambda (result replaced) replaced)) returns
  #t
(subst-every-other-sf 'a 'b '(a a) (lambda (result replaced) result)) returns
  (b a)
(subst-every-other-sf 'a 'b '(a a) (lambda (result replaced) replaced))
  returns #f
(subst-every-other-sf 'a 'b '(a a a a a a) (lambda (result replaced)
  result)) returns (b a b a b a)
(subst-every-other-sf 'a 'b '(a a a a a a) (lambda (result replaced)
  replaced)) returns #f
```

```

150 (subst-every-other-sf 'a 'b '(a a a a a a) (lambda (result replaced)
151   result)) returns (b a b a b a b)
152 (subst-every-other-sf 'a 'b '(a a a a a a a) (lambda (result replaced)
153   replaced)) returns #t
154 (subst-every-other-sf 'a 'b '(a (a (a a) a) a) (lambda (result replaced)
155   result)) returns (b (a (b a) b) a)
156 (subst-every-other-sf 'a 'b '(a (a (a a) a) a) (lambda (result replaced)
157   replaced)) returns #f
158 (subst-every-other-sf 'a 'b '(a (a (a (a (a (a ())))))) (lambda (result
159   replaced) result)) returns (b (a (b (a (b (a ()))))))
160 (subst-every-other-sf 'a 'b '(a (a (a (a (a (a ())))))) (lambda (result
161   replaced) replaced)) returns #f
162 (subst-every-other-sf 'a 'x '(a (b (a (c (a (d ())))))) (lambda (result
163   replaced) result)) returns (x (b (a (c (x (d ()))))))
164 (subst-every-other-sf 'a 'x '(a (b (a (c (a (d ())))))) (lambda (result
165   replaced) replaced)) returns #t

```

**167 Strategy:**

168 This is a first exercise in writing Scheme higher-order procedures. Use the techniques for  
169 designing higher-order procedures exemplified in *TLS* and in lecture.

171 You'll be graded on program correctness, style (including choice of identifiers/symbols), and  
172 documentation (including the required header comment described at the beginning of this  
173 handout), just as you have been in your earlier computer science courses.

175 Avoid major big-O inefficiencies where it's possible to do so without seriously obfuscating your  
176 code. Some of the big-O reckoning might take a bit of thought. Big-O bounds should be stated  
177 in terms of the properties of the parameters (e.g., the length of a list) rather than in terms of  
178 undefined variables (e.g., *n*). Assume that `car`, `cdr`, `cons`, `eq?`, `atom?`, and `null?` work in  
179 time  $O(1)$ .

181 Bundle up all your procedure definitions into a single file named `assign03.rkt`. If you use  
182 any of the example procedures from lecture (e.g., `rac`), include those definitions at the very end  
183 of `assign03.rkt`. If you don't finish a problem, and the associated procedure definitions  
184 aren't syntactically correct, please comment out those incomplete definitions and place a note at  
185 the top of your file indicating which definitions have been commented out.

## 187 What to Hand in:

- 188 • A printed listing of `assign03.rkt`. Format your listing (using landscape orientation,  
189 smaller fonts, etc.) to avoid illegible line-wrapping in your listing.
- 190 • Your file `assign03.rkt` submitted using the Assignment 3 item on the  
191 Assignments page of the CS370 Katie course