

Part 1 – assignment 5

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Q1

1.a

Two lazy lists $lzl1$ $lzl2$ are equivalent if:

- $lzl1$ and $lzl2$ are empty

Or

- $Lzl1(T) = \text{Pair} (T1 , (f : [\text{empty} \rightarrow lzl (T)]))$
 $Lzl2(D) = \text{Pair} (D1 , (g : [\text{empty} \rightarrow lzl (D)]))$

Such that $D = T$ and $T1 = D1$ and in the i 'th call to f the result is equivalent to the result of the i 'th call to g for every i .

1.b

Even-square-1 applies $lzl\text{-filter}$ (with the predicate that returns only the even numbers on the list) on the lazy list that is returned from $lzl\text{-map}$ which applies the square function on all the numbers of the list that is returned from 'integers-from' which returns a list of all natural numbers starting from 0.

Even-square-2 applies $lzl\text{-map}$ which applies the square function on all the numbers of the list that is returned from $lzl\text{-filter}$ (with the predicate that returns only the even numbers on the list) on the lazy list that is returned from 'integers-from' which returns a list of all natural numbers starting from 0.

Therefore, the lists returned are a list of the even squared natural numbers and the squared even natural numbers.

By our definition:

- $T=D$ - Both lists always return a number
- $T1 = D1$ - both start with 0.
- in the i 'th call to even-square-1 we will get the i 'th even squared number and in the i 'th call to even-square-2 we will get the i 'th even number squared which are the same.

Q2

a

f is equivalent to f\$ (its success – fail – continuations version)

\Leftrightarrow

1) if $f : [D1 * D2 * \dots * Dn \rightarrow T]$ then $f\$: [D1 * D2 * \dots * Dn * g * h \rightarrow S]$ when g and f are functions
and

2) if $f\$ (x1 \dots xn, succ, fail)$ is successful then $f\$ (x1 \dots xn, succ, fail) = succ(f(x1 \dots xk))$
else $f\$ (x1 \dots xn, succ, fail) = fail(f(x1 \dots xk))$

d

Get-value is of type $[List<Pair<Symbol, T>> * Symbol \rightarrow T \mid 'fail]$ and get-value\$ is of type

$[List<Pair<Symbol, T>> * Symbol * [T \rightarrow T1] * [Empty \rightarrow T2]] \rightarrow T1 \mid T2]$ so (1) is satisfied.

Now we'll show that (2) is satisfied too.

When get-value\$ is successful, it means that a value v is found, and the result of the function $succ$ on that value is returned. This means, that the function get-value is successful too and it will return v , which in turn will mean that $get_value\$ (list, key, succ, fail) = succ(v) = succ(get_value(list, key))$.

When get-value\$ has failed, it means that no value has been found, and so the result of fail is returned. This means, that the function get-value is unsuccessful as well and returns a failure too.

Q3

1.a

$Equations = [t(s(s), G, H, p, t(E), s) = t(s(H), G, p, p, t(E), K)], sub = \{\}$

if the predicate symbols and the number of arguments are the same: $eq_1 = (p(t_1, \dots, t_n) = p(s_1, \dots, s_n))$:
split eq_1 into equations: $equations = equation \cup (t_i = s_i)$ for $i=1..n$.

$Equations = [s(s) = s(H), G = G, H = p, p = p, t(E) = t(E), s = K], sub = \{\}$

$Equations = Equations \circ \{s = H\}$

$Equations = [s = H, G = G, H = p, p = p, t(E) = t(E), s = K], sub = \{\}$

$sub = sub \circ \{s = H\}$

$Equations = [G = G, H = p, p = p, t(E) = t(E), s = K], sub = \{s = H\}$

if the equation is the same variable on both sides, continue.

$Equations = [H = p, p = p, t(E) = t(E), s = K], sub = \{s = H\}$

$\{H = p\} \circ \{s = H\} \rightarrow \{s = p\}$

FAIL - both sides are different constant symbols

1.b

Equations = [$g(c, v(U), g, G, U, E, v(M)) = g(c, M, g, v(M), v(G), g, v(M))$]

Sub = {}

if the predicate symbols and the number of arguments are the same: $eq'_1 = (p(t_1, \dots, t_n) = p(s_1, \dots, s_n))$:
 split eq'_1 into equations: equations = equation $U(t_i = s_i)$ for $i=1..n$

Equations = [$c = c, v(U) = M, g = g, G = v(M), U = v(G), E = g, v(M) = v(M)$], sub={}

if both sides are the same constant symbol then continue

Equations = [$v(U) = M, g = g, G = v(M), U = v(G), E = g, v(M) = v(M)$], sub={}

sub = sub $\circ \{ v(U) = M \}$

Equations = [$g = g, G = v(M), U = v(G), E = g, v(M) = v(M)$], sub={ $v(U) = M$ }

if both sides are the same constant symbol then continue

Equations = [$G = v(M), U = v(G), E = g, v(M) = v(M)$], sub={ $v(U) = M$ }

$\{G = v(M)\} \circ \text{sub} = \{G = v(v(U))\}$

sub = sub $\circ \{G = v(v(U))\}$

Equations = [$U = v(G), E = g, v(M) = v(M)$], sub={ $v(U) = M, G = v(v(U))$ }

$\{U = v(G)\} \circ \text{sub} = \{U = v(v(v(U)))\}$

FAIL - occurs check

1.c

Equations = [$s([v | ([v | V] | A)]) = s([v | [v | A]])$], Sub = {}

if the predicate symbols and the number of arguments are the same: $eq'_1 = (p(t_1, \dots, t_n) = p(s_1, \dots, s_n))$:
 split eq'_1 into equations: equations = equation $U(t_i = s_i)$ for $i=1..n$

Equations = [$[v | ([v | V] | A)] = [v | [v | A]]$], sub={}

the predicate symbols and the number of arguments are the same: $eq'_1 = (p(t_1, \dots, t_n) = p(s_1, \dots, s_n))$:
 split eq'_1 into equations: equations = equation $U(t_i = s_i)$ for $i=1..n$,
 continue.

Equations = [$v = v, [v | V] | A = [v | A]$], sub={}

if both sides are the same constant symbol then continue

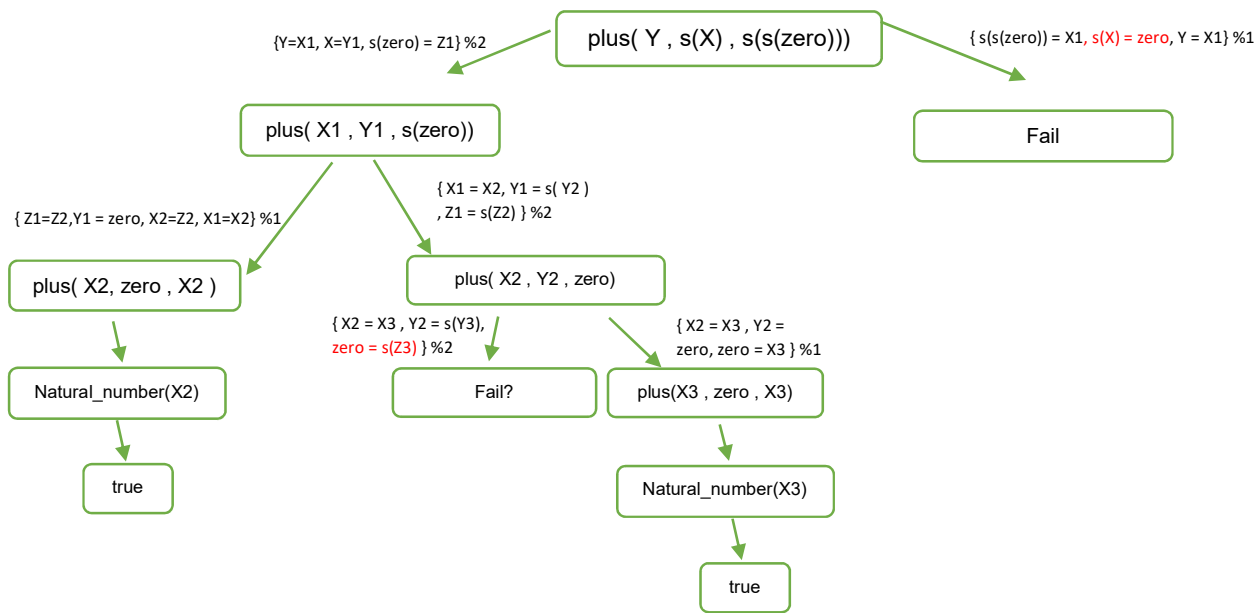
Equations = [$[v | V] | A = [v | A]$], sub={}

the predicate symbols and the number of arguments are the same: $eq'_1 = (p(t_1, \dots, t_n) = p(s_1, \dots, s_n))$:
 split eq'_1 into equations: equations = equation $U(t_i = s_i)$ for $i=1..n$,
 continue.

Equations = [$[v | V] = v, A = A$], sub={}

FAIL - Constant symbol value v can't be equal to list predicate

2.a



2.b

$Y = \text{zero}, X = s(\text{zero})$

$Y = s(\text{zero}), X = \text{zero}$

2.c

This is a success proof tree because it has at least one successful path.

2.d

This tree is finite because it has no infinite path .