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# https://towardsdatascience.com/fast-fourier-transform-937926e591cb  
# Hsu H.P. - Schaum's outline of signals and systems-McGraw Hill (2011).pdf  
# page 324
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Discrete Fourier Transform

The Discrete Fourier Transform (DTF) can be written as follows.

$$x[k] = \sum_{n=0}^{N-1} x[n] e^{\frac{-j2\pi kn}{N}}$$

Discrete Fourier transform $O(N^2)$

Fast Fourier transform $O(N \log_2 N)$

Fast Fourier Transform Algorithm

Suppose, we separated the Fourier Transform into even and odd indexed sub-sequences.

$$\begin{cases} n = 2r & \text{if } \textit{even} \\ n = 2r + 1 & \text{if } \textit{odd} \end{cases}$$

$$\text{where } r = 1, 2, \dots, \frac{N}{2} - 1$$

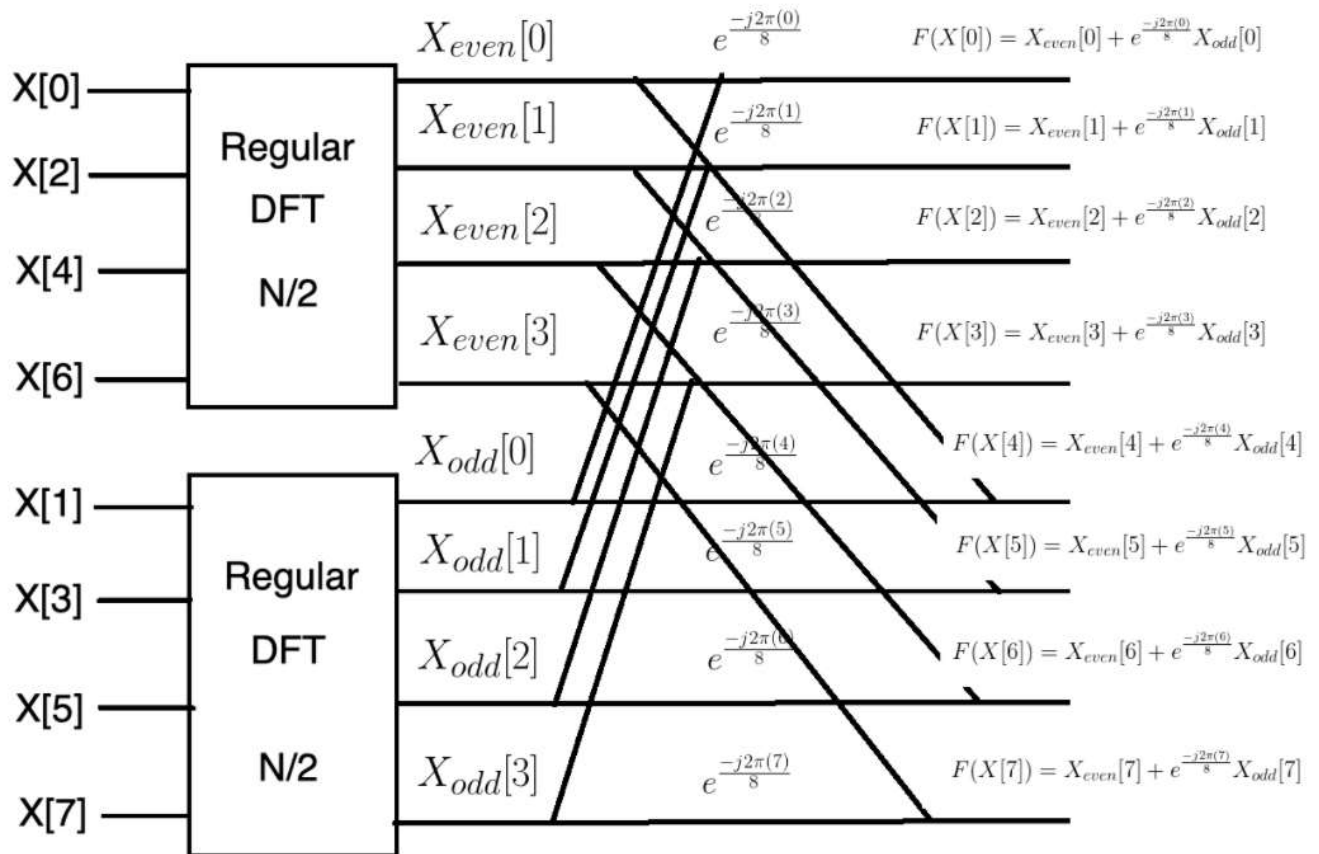
$$x[k] = \sum_{n=0}^{N-1} x[n] e^{\frac{-j2\pi kn}{N}}$$

$$x[k] = \sum_{r=0}^{\frac{N}{2}-1} x[2r] e^{\frac{-j2\pi k(2r)}{N}} + \sum_{r=0}^{\frac{N}{2}-1} x[2r+1] e^{\frac{-j2\pi k(2r+1)}{N}}$$

$$x[k] = \sum_{r=0}^{\frac{N}{2}-1} x[2r] e^{\frac{-j2\pi k(2r)}{N}} + e^{\frac{-j2\pi k}{N}} \sum_{r=0}^{\frac{N}{2}-1} x[2r+1] e^{\frac{-j2\pi k(2r)}{N}}$$

$$x[k] = \sum_{r=0}^{\frac{N}{2}-1} x[2r] e^{\frac{-j2\pi k(r)}{N/2}} + e^{\frac{-j2\pi k}{N}} \sum_{r=0}^{\frac{N}{2}-1} x[2r+1] e^{\frac{-j2\pi k(r)}{N/2}}$$

$$x[k] = x_{\text{even}}[k] + e^{\frac{-j2\pi k}{N}} x_{\text{odd}}[k]$$



DFT on $N/2$ elements

2 DFT

$$x[k] = x_{even}[k] + e^{-j\frac{2\pi k}{N}} x_{odd}[k]$$

$$2 \times \left(\frac{N}{2}\right)^2 + N$$

$$= 2 \times \frac{N^2}{4} + N$$

$$= \frac{N^2}{2} + N$$

$$O\left(\frac{N^2}{2} + N\right) \sim O(N^2)$$

What if we keep splitting?

$$\frac{N}{2} \longrightarrow 2 \left(\frac{N}{2} \right)^2 + N = \frac{N^2}{2} + N$$

$$\frac{N}{4} \longrightarrow 2 \left(2 \left(\frac{N}{4} \right)^2 + \frac{N}{2} \right) + N = \frac{N^2}{4} + 2N$$

$$\frac{N}{8} \longrightarrow 2 \left(2 \left(2 \left(\frac{N}{8} \right)^2 + \frac{N}{4} \right) + \frac{N}{2} \right) + N = \frac{N^2}{8} + 3N$$

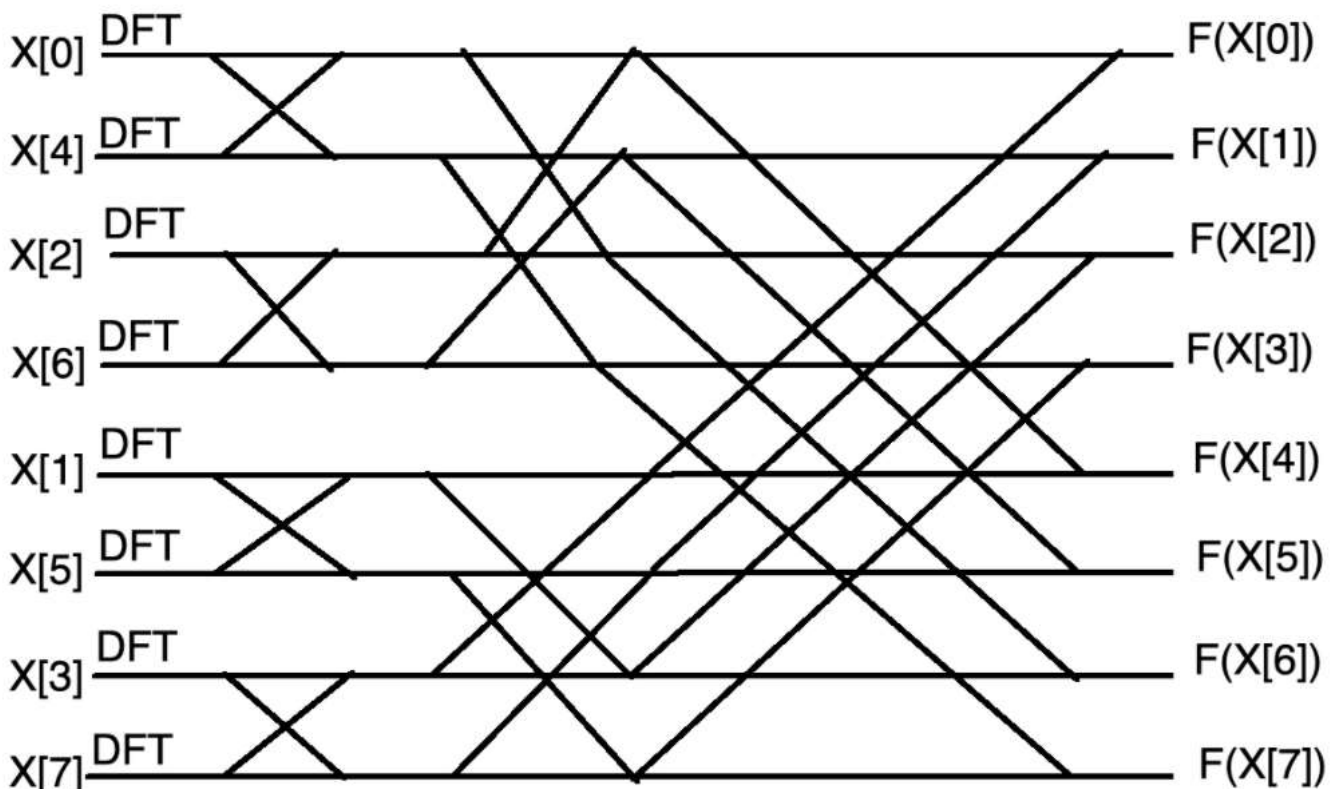
⋮

$$\frac{N}{2^p} \longrightarrow \frac{N^2}{2^p} + pN = \frac{N^2}{N} + (\log_2 N)N = N + (\log_2 N)N$$

$$\sim O(N + N \log_2 N) \sim O(N \log_2 N)$$

Split a maximum of $p = \log_2 N$ times

$$N = 8, p = 3$$



- (b) Let $W_{n+1, k+1}$ denote the entry in the $(n+1)$ st row and $(k+1)$ st column of the \mathbf{W}_4 matrix. Then, from Eq. (6.207)

$$W_{n+1, k+1} = W_4^{nk} = e^{-j(2\pi/4)nk} = e^{-j(\pi/2)nk} = (-j)^{nk} \quad (6.211)$$

and we have

$$\mathbf{W}_4 = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & -j & -1 & j \\ 1 & -1 & 1 & -1 \\ 1 & j & -1 & -j \end{bmatrix} \quad \mathbf{W}_4^{-1} = \frac{1}{4} \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & j & -1 & -j \\ 1 & -1 & 1 & -1 \\ 1 & -j & -1 & j \end{bmatrix} \quad (6.212)$$