Motivation
Definitions
Algorithms
Implementation
Applications

cis112 Binary Search Trees

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- Motivation
- 2 Definitions
- 3 Algorithms
- 4 Implementation
- 6 Applications



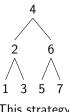
Motivation

Game of secret number

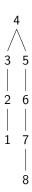
Game.

- Player-A selects a secret number in the range of {1,2,...,8}.
- Game repeats till the secret is found.
 - Player-B tries a number.
 - Player-A should answer as "found", or "smaller" or "larger".

Q. What is the best strategy if the range becomes large such as $\{1, 2, ..., 1000, 000\}$.



This strategy finds the secret faster.



This strategy is slower.

Definitions

Binary Search Tree

Definition

Binary search tree property:

Let x be a node in a binary tree.

- 1 If y is a node in the left subtree of x, then $y.key \le x.key$.
- 2 If y is a node in the right subtree of x, then $y.key \ge x.key$.

A binary tree with the binary-search-tree property is called a binary search tree.

Example



- 2 is in the left subtree of 5 since $2 \le 5$. 2 is in the left subtree of 6 since $2 \le 6$.
- 2 is in the left subtree of 6 since $2 \le 6$. 8 is in the right subtree of 7 since 8 > 7.
- 8 is in the right subtree of 6 since $8 \ge 6$.

[[1], [2], [3], [4], [5], [6], [7]]

Tree Traversal Searching Minimum, Maximum, Successor, Predecessor Insertion, Deletion

Algorithms

Node

Algorithm 1: Node for Binary Search Tree

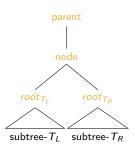
```
| begin | 2 | data | // reference to the parent | 4 | left | // reference to the left child | 5 | right | // reference to the right child |
```

[[2]

Warning. key is comparable:

Comparison of a.data and b.data for Nodes, a and b, is defined as

$$\label{eq:a.key.comparedTo(b.key)} \text{a.key.comparedTo(b.key)} = \begin{cases} <0, & \text{a.data < b.data,} \\ 0, & \text{a.data == b.data,} \\ >0, & \text{a.data > b.data.} \end{cases}$$



Tree Traversal

earching inimum, Maximum, Successor, Predecessor sertion, Deletion

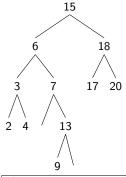
Algorithms

Tree Traversal

Inorder Tree Traversal

Function InorderTreeWalk(x) [[2]]

```
1 if x ≠ N/L then
2 | InorderTreeWalk(x.left)
3 | print x.data
4 | InorderTreeWalk(x.right)
```



Output of InorderTreeWalk (15)
2, 3, 4, 6, 7, 9, 13, 15, 17, 18, 20.

Algorithms

Searching

Tree Search

```
Function TreeSearch(x, k) [2]

if x == NIL or key == x.key then

treturn x

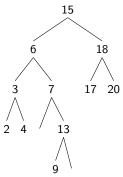
if k < x.key then

return TreeSearch(x.left, k)

else

return TreeSearch(x.right, k)
```

Trace of TreeSearch (15, 14)
TreeSearch (15,14)
TreeSearch (6,14)
TreeSearch (7,14)
TreeSearch (13,14)
TreeSearch (NIL,14)

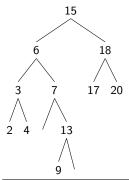


Trace of TreeSearch (15, 13)
TreeSearch (15,13)
TreeSearch (6,13)
TreeSearch (7,13)
TreeSearch (13.13)

Iterative Tree Search

Function IterativeTreeSearch(x, k) [[2]]

6 return x



Trace of IterativeTreeSearch (15, 13)
x: 15 // lines: 1-5
x: 6 // lines: 1-5

x: 7 // lines: 1-5 x: 13 // lines: 1-5

Algorithms

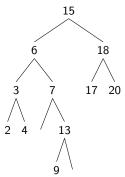
Minimum, Maximum, Successor, Predecessor

Tree Minimum

Function TreeMinimum(x) [[2]]

- while x.left \neq NIL do
- x = x.left
- 3 return x

- Q. TreeMinimum (6)?
- Q. TreeMinimum (7)?
- Q. TreeMinimum (18)?



Trace of TreeMinimum (15)
15 // lines: 1-2
6 // lines: 1-2
3 // lines: 1-2
2 // lines: 1-2

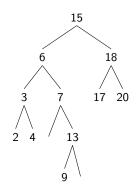
Tree Maximum

Function TreeMaximum(x) [[2]]

- while x.right ≠ NIL do
- 2 | x = x.right
- 3 return x

- Q. TreeMaximum (3)?
- Q. TreeMaximum (4)?

Trace of TreeMaximum (6)
6 // lines: 1-2
7 // lines: 1-2
13 // lines: 1-2



Trace of TreeMaximum (15)	
15 // lines: 1-2	
18 // lines: 1-2	
20 // lines: 1-2	
400400455455	-

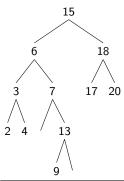
Tree Successor

```
Trace of TreeSuccessor(15)

15 // line: 2

18 // line: 2

17 // line: 2
```



```
Trace of TreeSuccessor (13)
x:13, y:7 // lines: 4-7
x:7, y:6 // lines: 4-7
x:6, y:15 // lines: 4-7
```

Algorithms

Insertion, Deletion

Tree Insert

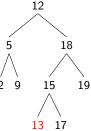
y.left = z

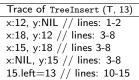
y.right = z

7

8

```
Function TreeInsert(T, z) [[2]]
  x = T root
                                // node being compared with z
                                       // will be parent of z
v = NIL
3 while x \neq NIL do
       // descend until reaching a leaf
       y = x
       if z.key < x.key then
             x = x.left
       else
            x = x.right
9 z.parent = y // found the location- insert z with parent y
10 if y == NIL then
       T.root = z
                                           // tree was empty
  else if z.key < y.key then
```





Deleting node z

Deleting a node z, in blue, from a binary search tree. Node z may be the root, a left child of node q, or a right child of q. The node that will replace node z in its position in the tree is colored orange.

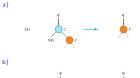
a) Node z has no left child.



Deleting node z

Deleting a node z, in blue, from a binary search tree. Node z may be the root, a left child of node q, or a right child of q. The node that will replace node z in its position in the tree is colored orange.

- a) Node z has no left child.
- b) Node z has a left child ℓ but no right child.

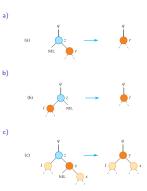




Deleting node z

Deleting a node z, in blue, from a binary search tree. Node z may be the root, a left child of node q, or a right child of q. The node that will replace node z in its position in the tree is colored orange.

- a) Node z has no left child.
- b) Node z has a left child ℓ but no right child.
- c) Node z has two children. Its left child is node ℓ, its right child is its successor y (which has no left child), and y's right child is node x.

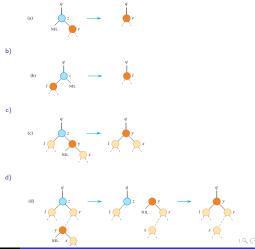


a)

Deleting node z

Deleting a node z, in blue, from a binary search tree. Node z may be the root, a left child of node q, or a right child of q. The node that will replace node z in its position in the tree is colored orange.

- a) Node z has no left child.
- b) Node z has a left child ℓ but no right child.
- c) Node z has two children. Its left child is node ℓ, its right child is its successor y (which has no left child), and y's right child is node x.
- d) Node has two children (left child ℓ and right child r), and its successor y ≠ r lies within the subtree rooted at r.

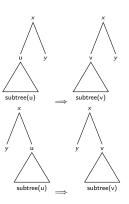


Transplant

Function Transplant(T, u, v) [[2]]

- 1 if u.p == NIL then
- 2 | T.root = v
- 3 else if u == u.p.left then
- 4 | u.p.left = v
- 5 else
- 6 | u.p.right = v
- 7 if $v \neq NIL$ then
- 8 | v.p = u.p

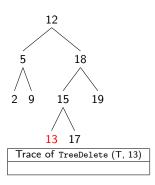
TRANSPLANT replaces one subtree as a child of its parent with another subtree. When TRANSPLANT replaces the subtree rooted at node u with the subtree rooted at node v, node u's parent becomes node v's parent, and u's parent ends up having v as its appropriate child.



Tree Delete

Function TreeDelete(T, z) [[2]]

```
if z left == NII. then
        Transplant (T, z, z.right)
                                         // replace z by its right child
  else if z.right == NIL then
        Transplant (T, z, z.left)
                                          // replace z by its left child
5 else
        y = TreeMinimum (z.right)
                                                  // v is z's successor
        if y \neq z.right then
                                            // is farther down the tree?
             Transplant (T, y, y.right) // replace y by its right child
 8
             y.right = z.right
                                          // z's right child becomes
             v.right.parent = v
10
                                                     // y's right child
        Transplant (T, z, y)
                                      // replace z by its successor y
11
        v.left = z.left
                                         // and give z's left child to v.
12
        v.left.parent = v
13
                                              // which had no left child
```



Implementation

Node for Binary Search Tree

Algorithm 2: Node for Binary

Search Tree

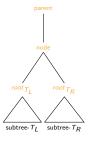
```
1 begin
2 data
3 parent // reference to the parent
4 left // reference to the left child
5 right // reference to the right child
```

[[2]

In binary search tree, each node has 0, 1 or 2 subtrees.

Use node with three pointers:

- parent: points the parent
- left: points the left child
- right: points the right child



Binary Search Tree property:

- Every node in subtree T_L is smaller than node
- Every node in subtree T_R is larger than node
- Q. Is node.data smaller than parent.data?
- Q. Is node.data larger than parent.data?

Assume that data at each node is different.

MyNodeBST

Algorithm 3: Node for Binary

Search Tree

[[2]

MyNodeBST implements interface NodeBSTInterface.

T implements Comparable

```
y.data.compareTo(x.data) = \begin{cases} < 0, & y.data < x.data, \frac{17}{18} \\ 0, & y.data = x.data, \\ > 0, & y.data > x.data. \end{cases}
```

```
1 MyNodeBST<T> x = ...;
2 MyNodeBST<T> y = ...;
3 4 // compare x and y
5 if (y.data.compareTo(x.data) < 0) {
6     // y is smaller than x
7 } else {
8     // y is larger than x
9 }</pre>
```

```
public class MyNodeBST<T extends Comparable<T>>>
       implements MyNodeBSTInterface<T> {
     T data:
     MyNodeBST<T> parent;
     MyNodeBST<T> left;
     MyNodeBST<T> right;
8
     public MyNodeBST() {
9
       this (null);
     public MyNodeBST(T data) {
       this . data = data;
       parent = null;
       left = null:
14
       right = null:
```

```
public interface NodeBSTInterface<T> {

    T data();

    MyNodeBSTInterface<T> parent();

    MyNodeBSTInterface<T> left();

    MyNodeBSTInterface<T> right();

    T canonical();

    T canonical();
```

MyBST, constructor

```
public class MyBST<T extends Comparable<T>>> {
     MyNodeBST<T> root;
 4
     public MyBST() {
       root = null;
 6
     public void insert (T data) {
8
        if (data = null) {
9
          return:
       MyNodeBST<T> x = root;
       MyNodeBST<T> y = null;
       MyNodeBST<T> z = new MyNodeBST<>(data);
14
       while (x != null) {
         v = x:
16
          if (z, data, compareTo(x, data) < 0) {
           x = x. left:
          } else {
            x = x.right:
20
       // location is found.
       z.parent = v:
        if (v = null) {
24
          root = z:
26
       } else if (z.data.compareTo(y.data) < 0) {</pre>
         y.left = z;
28
         else {
29
         y.right = z;
30
31
32
33
```

```
public class MyNodeBST<T extends Comparable<T>>>
        implements MyNodeBSTInterface<T> {
     T data:
 4
     MyNodeBST<T> parent;
     MyNodeBST<T> left;
     MyNodeBST<T> right;
8
     public MyNodeBST() {
9
        this (null);
     public MyNodeBST(T data) {
        this . data = data;
        parent = null;
        left = null:
14
        right = null:
16
18
```

```
public class MyBSTConstructor {

public static MyBST<String>
    constructBST.S.412.nodeByNode() {

MyBST<String> tree = new MyBST<>);

tree.insert("4");

tree.insert("2");

return tree;

}

...

}
```

MyBST, insert

```
public class MyBST<T extends Comparable<T>>> {
     MvNodeBST<T> root;
 4
     public MyBST() {
       root = null;
 6
     public void insert (T data) {
8
        if (data = null) {
9
          return:
       MyNodeBST<T> x = root;
       MyNodeBST<T> y = null;
       MyNodeBST<T> z = new MyNodeBST<>(data);
       while (x != null) {
14
         v = x:
16
          if (z, data, compareTo(x, data) < 0) {
            x = x.left:
          } else {
19
            x = x, right:
20
       // location is found.
       z.parent = v:
        if (v = null) {
24
25
          root = z:
26
       } else if (z.data.compareTo(y.data) < 0) {</pre>
         y.left = z;
28
         else {
29
         y.right = z;
30
31
32
33
```

```
public class MyBSTConstructor_Test {
     private static final boolean DEBUG = true:
4
     public static void main(String[] args) {
       MyBST<String> tree;
8
       tree = MyBSTConstructor.
          constructBST_S_412_nodeByNode();
9
10
       if (DEBUG) {
         tree.plot();
         System.out.println(tree.canonical());
14
15
16
   public class MyBSTConstructor {
3
     public static MyBST<String>
          constructBST_S_412_nodeBvNode() {
       MyBST < String > tree = new MyBST <> ();
4
       tree . insert ("4");
6
       tree . insert ("1");
       tree.insert("2"):
8
       return tree:
q
11
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```

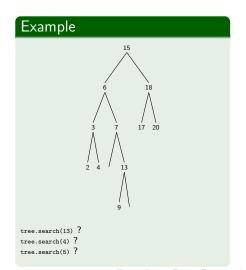
/[1]/[2]\[4]\

MyBST, insert

```
public class MyBST<T extends Comparable<T>>> {
     MyNodeBST<T> root;
 4
     public MyBST() {
       root = null;
 6
     public void insert (T data) {
8
        if (data = null) {
9
          return:
       MyNodeBST<T> x = root;
       MyNodeBST<T> y = null;
       MyNodeBST<T> z = new MyNodeBST<>(data);
       while (x != null) {
14
         v = x:
16
          if (z, data, compareTo(x, data) < 0) {
            x = x.left:
18
          } else {
19
            x = x, right:
20
       // location is found.
       z.parent = v:
        if (v = null) {
24
          root = z:
26
       } else if (z.data.compareTo(y.data) < 0) {</pre>
         y.left = z;
28
         else {
29
         y.right = z;
30
31
32
33
```

```
public class MyBSTConstructor_Test {
     private static final boolean DEBUG = true:
4
     public static void main(String[] args) {
       MyBST<String> tree;
8
       tree = MyBSTConstructor.
          constructBST_S_412_nodeByNode();
9
10
       if (DEBUG) {
         tree.plot();
         System.out.println(tree.canonical());
14
15
16
   public class MyBSTConstructor {
3
     public static MyBST<String>
          constructBST_S_412_nodeBvNode() {
       MyBST < String > tree = new MyBST <> ();
4
       tree . insert ("4");
6
       tree . insert ("1");
       tree.insert("2"):
8
       return tree:
q
11
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```

MyBST, search

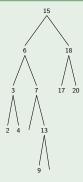


MyBST, search

```
1 public class MyBST<T extends Comparable<T> {
            MyNodeBST<T root;
            ...
4 public MyNodeBST<T> successor() {
            return (MyNodeBST<T>) LibTreeBS.successor(
            root);
6 }
7
8 public MyNodeBST<T> successor(MyNodeBST<T> x)
            {
            return (MyNodeBST<T> bubTreeBS.successor(x)
            ;
            ;
10 }
11 ...
12 }
```

```
1 search (13): 15 \rightarrow 6 \rightarrow 7 \rightarrow 13
2 search (4): 15 \rightarrow 6 \rightarrow 3 \rightarrow 4
3 search (5): 15 \rightarrow 6 \rightarrow 3 \rightarrow 4 \rightarrow \text{not found}
```

Example



```
tree.search(13) ?
tree.search(4) ?
tree.search(5) ?
```

MyBST, minimum

```
public class MyBST<T extends Comparable<T>>> {
     MyNodeBST<T> root;
     public MyNodeBST<T> minimum() {
       return (MyNodeBST<T>) LibTreeBS.minimum(root
          ):
8
9
     public MyNodeBST<T> minimum (MyNodeBST<T> x) {
       return (MyNodeBST<T>) LibTreeBS.minimum(x);
   public class LibTreeBS<T extends Comparable<T>>>
     public static <T> NodeBSTInterface<T> minimum(
         NodeBSTInterface<T> x
       while (x.left() != null) {
         x = x.left();
Q
       return x:
12
```

Example 15 17 20 tree.minimum() ? tree.minimum(tree.root.right) ? tree.minimum(tree.root.right.right) ?

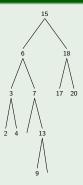
MyBST, minimum

12

```
public class MyBST<T extends Comparable<T>>> {
    MyNodeBST<T> root;
    public MyNodeBST<T> minimum() {
      return (MyNodeBST<T>) LibTreeBS.minimum(root
         ):
8
9
    public MyNodeBST<T> minimum (MyNodeBST<T> x) {
      return (MyNodeBST<T>) LibTreeBS.minimum(x);
  public class LibTreeBS<T extends Comparable<T>>
    public static <T> NodeBSTInterface<T> minimum(
         NodeBSTInterface<T> x
      while (x.left() != null) {
        x = x.left();
8
Q
      return x:
```

```
minimum (15): 2
minimum (18): 17
minimum (20): 20
```

Example



```
tree.minimum() ?
tree.minimum(tree.root.right) ?
tree.minimum(tree.root.right.right) ?
```

MyBST, maximum

```
public class MyBST<T extends Comparable<T>>> {
     MyNodeBST<T> root;
     public MyNodeBST<T> maximum() {
       return (MyNodeBST<T>) LibTreeBS.maximum(root
         ):
6
8
     public MyNodeBST<T> maximum(MyNodeBST<T> x) {
9
       return (MyNodeBST<T>) LibTreeBS.maximum(x);
   public class LibTreeBS<T extends Comparable<T>>>
     public static <T> NodeBSTInterface<T> maximum(
         NodeBSTInterface<T> x
       while (x.right() != null) {
         x = x.right();
9
       return x:
10
```

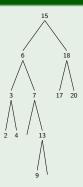
Example 15 17 20 tree.maximum() ? tree.maximum(tree.root.left) ? tree.maximum(tree.root.left.left) ?

MyBST, maximum

```
public class MyBST<T extends Comparable<T>>> {
     MyNodeBST<T> root;
     public MyNodeBST<T> maximum() {
       return (MyNodeBST<T>) LibTreeBS.maximum(root
         ):
6
8
     public MyNodeBST<T> maximum(MyNodeBST<T> x) {
9
       return (MyNodeBST<T>) LibTreeBS.maximum(x);
   public class LibTreeBS<T extends Comparable<T>>>
3
     public static <T> NodeBSTInterface<T> maximum(
         NodeBSTInterface<T> x
       while (x.right() != null) {
         x = x.right();
9
       return x:
10
```

maximum(15): 20 maximum(6): 13 maximum(3): 4

Example



```
tree.maximum() ?
tree.maximum(tree.root.left) ?
tree.maximum(tree.root.left.left) ?
```

MyBST, successor

```
public class MyBST<T extends Comparable<T>>> {
     MyNodeBST<T> root;
      public MyNodeBST<T> maximum() {
        return (MyNodeBST<T>) LibTreeBS.maximum(root
          ):
8
      public MyNodeBST<T> maximum(MyNodeBST<T> x) {
9
        return (MyNodeBST<T>) LibTreeBS.maximum(x);
   public class LibTreeBS<T extends Comparable<T>>>
      public static <T> NodeBSTInterface<T>
          successor(NodeBSTInterface<T> x) {
        if (x.right() != null) {
          // leftmost node in the right subtree
          return minimum(x.right());
        } else {
          // find the lowest ancestor of x whose
          left child is an ancestor of x
          NodeBSTInterface\langle T \rangle y = null;
9
          y = x.parent();
10
          while (y \mid = \text{null } \&\& x \Longrightarrow y. \text{right}()) {
            x = y;
            y = y.parent();
14
          return y;
16
```

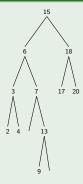
Example 15 17 20 tree.successor() ? tree.successor(tree.root.left) ? tree.successor(tree.root.left.right.right) ?

MyBST, successor

```
public class MyBST<T extends Comparable<T>>> {
     MyNodeBST<T> root;
     public MyNodeBST<T> maximum() {
       return (MyNodeBST<T>) LibTreeBS.maximum(root
          ):
8
     public MyNodeBST<T> maximum(MyNodeBST<T> x) {
9
       return (MyNodeBST<T>) LibTreeBS.maximum(x);
   public class LibTreeBS<T extends Comparable<T>>>
     public static <T> NodeBSTInterface<T>
          successor(NodeBSTInterface<T> x) {
        if (x.right() != null) {
         // leftmost node in the right subtree
          return minimum(x.right());
       } else {
         // find the lowest ancestor of x whose
          left child is an ancestor of x
          NodeBSTInterface\langle T \rangle y = null;
9
         y = x.parent();
10
          while (y \mid = \text{null } \&\& x == y.right()) {
            x = y;
            y = y.parent();
14
          return y;
16
```

```
successor (15): 17
successor (6): 7
successor (13): 15
```

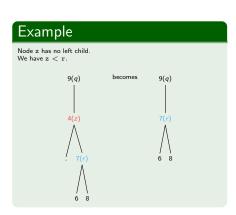
Example



```
tree.successor() ?
tree.successor(tree.root.left) ?
tree.successor(tree.root.left.right.right) ?
```

MyBST, delete (i)

```
public class MyBST<T extends Comparable<T>>> {
     MyNodeBST<T> root;
     public void delete(MyNodeBST<T> z) {
       MyNodeBST < T > y = null;
       if (z.left = null) {
         // replace z by its right child
8
         transplant(z, z.right);
9
       } else if (z.right == null) {
         // replace z by its left child
         transplant(z, z.left);
       } else {
         // y is z's successor
14
         y = minimum(z.right);
         if (y != z.right) {
           // v is farther down in the tree
16
           // replace y by its right child
           transplant(v. v.right):
           // z's right child becomes v's right
          child
           v.right = z.right;
20
           y.right.parent = y;
         // replace z by its successor y
24
         transplant(z. v):
25
         // z's right child, which has no let child
           becomes v's right child
         v.left = z.left:
26
         y.left.parent = y;
28
29
30
```



MyBST, delete (ii)

```
public class MyBST<T extends Comparable<T>>> {
     MyNodeBST<T> root;
     public void delete(MyNodeBST<T> z) {
       MyNodeBST < T > y = null;
       if (z.left = null) {
         // replace z by its right child
8
         transplant(z, z.right);
9
       } else if (z.right = null) {
         // replace z by its left child
         transplant(z, z.left);
       } else {
         // y is z's successor
14
         y = minimum(z.right);
         if (y != z.right) {
           // v is farther down in the tree
16
           // replace y by its right child
           transplant(v. v.right):
           // z's right child becomes v's right
          child
           v.right = z.right;
20
           y.right.parent = y;
         // replace z by its successor y
24
         transplant(z. v):
25
         // z's right child, which has no let child
          . becomes v's right child
         v.left = z.left:
26
         y.left.parent = y;
28
29
30
```

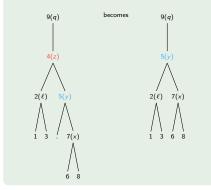
Example Node z has a left child ℓ but no right child. We have $\ell < z$. becomes 9(q)9(q)4(z)

MyBST, delete (iii)

```
public class MyBST<T extends Comparable<T>>> {
     MyNodeBST<T> root;
     public void delete(MyNodeBST<T> z) {
       MyNodeBST<T> y = null;
       if (z.left = null) {
         // replace z by its right child
         transplant(z, z.right);
9
       } else if (z.right = null) {
         // replace z by its left child
         transplant(z, z.left);
       } else {
         // y is z's successor
14
         y = minimum(z.right);
         if (y != z.right) {
           // v is farther down in the tree
16
           // replace y by its right child
           transplant(v. v.right):
           // z's right child becomes v's right
          child
20
           y.right = z.right;
           y.right.parent = y;
         // replace z by its successor y
24
         transplant(z. v):
25
         // z's right child, which has no let child
           becomes v's right child
         v.left = z.left:
26
         y.left.parent = y;
28
29
30
```

Example

Node z has two children (ℓ and y). Node y is the successor of z, that is, y has no left child. Node x is the right child of y. Hence, y = successor(z) and $\ell < z < y < x$.

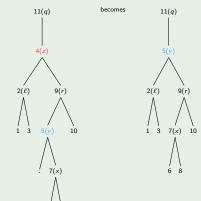


MyBST, delete (iv)

```
public class MyBST<T extends Comparable<T>>> {
     MyNodeBST<T> root;
     public void delete(MyNodeBST<T> z) {
       MyNodeBST < T > y = null;
       if (z.left = null) {
         // replace z by its right child
8
         transplant(z, z.right);
9
       } else if (z.right == null) {
         // replace z by its left child
         transplant(z, z.left);
       } else {
         // y is z's successor
14
         y = minimum(z.right);
         if (y != z.right) {
           // v is farther down in the tree
16
           // replace v by its right child
           transplant(v. v.right):
           // z's right child becomes v's right
          child
20
           v.right = z.right:
           y.right.parent = y;
         // replace z by its successor y
24
         transplant(z, y);
25
         // z's right child, which has no let child
           becomes v's right child
         v.left = z.left:
26
         y.left.parent = y;
28
29
30
```

Example

Node z has two children (ℓ and r) and its successor y \neq r lies within the subtree rooted at r. Hence, y = successor(z) and y \neq r.



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Motivation
Definitions
Algorithms
Implementation
Applications

Applications

Some Applications of BST

Implement a set in math as a BST.

- Modify insert so that no duplication is allowed.
- If the BST is balanced, it would be fast to search, to insert and to delete in O(log n).

Database Management Systems (DBMS)

 DBMS keeps data in a system similar to BSTs.

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