

T2 Asymmetric Analysis Question Time

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A Wiener's Attack on RSA

Wiener's Attack on RSA

4 Points

Implement Wiener's attack on RSA to recover small $d (\rightarrow L6)$:

- a Compute n-th convergents of continued fractions for \mathbb{Q}
- b Recover private key from given 1024-bit RSA public key

A Questions

Example:

$$N = 9449868410449, \ e = 6792605526025, \ d < \frac{1}{3} \sqrt[4]{N} \approx 584.$$

1 Perform Wiener's attack by computing the convergents of $\frac{\theta}{N}$:

$$\frac{p_i}{q_i} = \quad \frac{1}{1}, \quad \frac{2}{3}, \quad \frac{3}{4}, \quad \frac{5}{7}, \quad \frac{18}{25} \quad \frac{23}{32}, \quad \frac{409}{569}, \quad \dots$$

- **2** Test: d = 569 if $M^{e \cdot 569} = M$
- 3 Pick x = 2:

$$x^{(ed-1)/2^1} = x^{(ed-1)/2^2} = \dots = x^{(ed-1)/2^5} = 1$$

 $x^{(ed-1)/2^6} = x^{60390508504816} \equiv 8233548335126 = y \neq \pm 1$

4 $p = \gcd(N, y - 1) = 1234577 \Rightarrow N = 1234577 \cdot 7654337$

B Discrete Logarithms with Pollard-p

Discrete Logarithms with Pollard- ρ

8 Points

Implement Pollard- ρ algorithm to find discrete logarithms (\rightarrow L5)

- a Implement Pollard- ρ and apply to 32-bit challenge
- **b** Solve resulting equation to complete the solution
- Solve 64-bit challenge

B Questions

To solve $y = g^x \pmod{p}$, find $r_j = y^{a_j}g^{b_j} = y^{a_k}g^{b_k} = r_k \pmod{p}$. Then $x \cdot (a_j - a_k) \equiv (b_k - b_j) \pmod{p-1}$.



$$\begin{aligned} & (\textbf{r}_0, \textbf{a}_0, \textbf{b}_0) = (\textbf{1}, \textbf{0}, \textbf{0}) \\ & (\textbf{r}_{i+1}, \textbf{a}_{i+1}, \textbf{b}_{i+1}) = \begin{cases} ([\textbf{y} \cdot \textbf{r}_i]_p, [\textbf{a}_i + \textbf{1}]_{p-1}, [\textbf{b}_i]_{p-1} &) & 0 < r_i < \frac{p}{3} \\ ([\textbf{r}_i^2]_p, [2\textbf{a}_i]_{p-1}, [2\textbf{b}_i]_{p-1}) & \frac{p}{3} < r_i < \frac{2p}{3} \\ ([\textbf{g} \cdot \textbf{r}_i]_p, [\textbf{a}_i]_{p-1}, [\textbf{b}_i + \textbf{1}]_{p-1}) & \frac{2p}{3} < r_i < p \end{cases}$$

C Factoring with Continued Fractions

Factoring with Continued Fractions

12 Points

Implement Cfrac factorization (\rightarrow L6):

- a Implement factoring with factor bases and random squares
- **b** Compute *n*-th convergents of continued fractions for \mathbb{R}
- Piece this together to get CFRAC and factor 32-bit N
- d Factor 64-bit N

C Questions

Example: Factor n = 9073 with CFRAC with $\mathcal{B} = \{-1, 2, 3, 5, 7\}$

1 Compute convergents for $\sqrt{9073} = 95.2523...$:

$$\frac{\rho_0}{q_0} = \frac{95}{1}, \quad \frac{\rho_1}{q_1} = \frac{286}{3}, \quad \frac{\rho_2}{q_2} = \frac{381}{4}, \quad \frac{\rho_3}{q_3} = \frac{10192}{107}, \quad \frac{\rho_4}{q_4} = \frac{20765}{218}$$

Smallest absolute residue b_i of p_i^2 mod 9073:

$$b_0 = -48, \ b_1 = 139, \ b_2 = -7, \ b_3 = 87, \ b_4 = -27$$

3 Check smoothness of the b_i and factorize:

$$b_0=(1,4,1,0,0),\quad b_2=(1,0,0,0,1),\quad b_4=(1,0,3,0,0).$$

4 Combine to get $(-36)^2 \equiv 3834^2 \pmod{9073}$:

$$x = b_0 \cdot b_4 = -1 \cdot 2^2 \cdot 3^2 = -36$$

 $y = p_0 \cdot p_4 = 95 \cdot 20765 \equiv 3834 \pmod{9073}$

5 Factor *n*: $gcd(3834 + 36,9073) = 43 \Rightarrow 9073 = 43 \cdot 211$