

Asymmetric Cryptanalysis – Submission Guidelines

- Timeline: Release **12. 05. 2016**; Question time **02. 06. 2016**; Submission **23. 06. 2016**
- Upload your team's submission on <https://stics.iaik.tugraz.at/>. Each task is uploaded separately. Don't forget to tick the selected tasks!
- Submit your code as a `{zip,tar.gz}` archive. Add a file `README.{md,txt,pdf}` on the top level that documents interesting properties and limitations of your submission.
- You can use your favourite programming language and libraries for the necessary big-integer and modular arithmetic. Consider using a computer algebra system, such as SAGE (<https://sage.tugraz.at> with TUGRAZonline login). Please document your choices, as well as how to compile and use your implementations in the `README`. If you intend to use non-free software, please clarify with us beforehand.

2–A Wiener's Attack on RSA (4 Points)

Implement the attack by Wiener [Wie90] to recover small RSA private exponents d .

- (a) **Continued fractions for \mathbb{Q} (2 Points)**: Implement the computation of n -th convergents of the continued fraction expansion to approximate rational numbers.
- (b) **Recover RSA private key (2 Points)**: Use your implementation to recover Bob's private key d, p, q from the following 1024-bit RSA public key with modulus N and exponent e :

```
N = 0x1cb33aad9d96e572b13204d77700515cf489028e35811ba23a2aa7763fb443a2
    2d6a7bac40a279e98faa91f82898efa3cfc57e11417f47f604093cd97371ff4e
    6d6e572c9c476ac6155719a4dbba3d93ce8891ea5116fe7a0502612052879f1c
    3c82c699dfb69518c6ed43871b88edb40da0b3cec6eabec9988adf8a4294547

e = 0x07f5e658edd83082b49740286814b8c63a2b0c4957d909af3a9b574e2f9d8d72
    bbe332244148a313350dac5657287e1e383e4b50f0a7d5078d4f6d48c19144e0
    4be6cb5bfef0e7fe100d03d51a5c6df8911ad49cf61eeaae849866ae9bbfa17a
    820374ad62ac1fec5692dc88cc85006975e87836cdda7888ddd822c0b47d5311
```

[Wie90] M. J. Wiener. "Cryptanalysis of short RSA secret exponents". In: IEEE Transactions on Information Theory 36.3 (1990), pp. 553–558. DOI: 10.1109/18.54902.

2–B Discrete Logarithms with Pollard- ρ (8 Points)

Implement the algorithm by Pollard [Pol78] to compute the discrete logarithm x in $y = g^x \pmod{p}$. Use it to recover Alice's and Bob's private keys a, b from a Diffie-Hellman key exchange.

- (a) **Find equation for 32-bit problem (4 Points):** Implement the Pollard- ρ algorithm to find an equation $x \cdot (a_j - a_k) \equiv (b_k - b_j) \pmod{p-1}$ for x . Apply your implementation to get equations for α, β from the following Diffie-Hellman key exchange:

Domain parameters:	generator $g = 0x00000002 \in \mathbb{Z}_p^*$
	prime $p = 0xffffffffb$
Alice \rightarrow Bob:	$g^\alpha = 0x4ebb660a$
Bob \rightarrow Alice:	$g^\beta = 0xe9467263$

- (b) **Determine correct solution for 32-bit problem (2 Points):** Recover α, β from the key exchange in (a) by testing all candidate solutions of the equations.
- (c) **Find solution for 64-bit problem (2 Points):** Find α or β to recover the key $g^{\alpha\beta}$:

Domain parameters:	generator $g = 0x0000000000000005 \in \mathbb{Z}_p^*$
	prime $p = 0x8ac25c704f8e5947$
Alice \rightarrow Bob:	$g^\alpha = 0x82c2c001a3ccff71$
Bob \rightarrow Alice:	$g^\beta = 0x07a3fb4211394fd1$

[Pol78] J. M. Pollard. “Monte Carlo Methods for Index Computation \pmod{p} ”. In: Mathematics of Computation 32.143 (1978), pp. 918–924. DOI: 10.2307/2006496.

2–C Factoring with Continued Fractions (12 Points)

Implement CFRAC factoring [LP31; MB75], and use it to factor the RSA modulus N .

- (a) **Factoring with random squares (6 Points):** Implement Dixon's basic factor-base factoring algorithm: Pick random x_i , factor $x_i^2 \pmod{n}$ with respect to a factor base \mathcal{B} , and use linear algebra to combine these factorizations to find solutions $x^2 \equiv y^2 \pmod{n}$.
- (b) **Continued fractions for \mathbb{R} (2 Points):** Implement the computation of n -th convergents of the continued fraction expansion to approximate irrational numbers like \sqrt{N} .
- (c) **Factor 32-bit number (2 Points):** Combine your ingredients and factor $N = 0x347b702f$.
- (d) **Factor 64-bit number (2 Points):** Factor $N = 0x8de50360f22507bf$.

[LP31] D. H. Lehmer and R. E. Powers. “On Factoring Large Numbers”. In: Bulletin of the American Mathematical Society 37.10 (1931), pp. 770–776. DOI: 10.1090/S0002-9904-1931-05271-X.

[MB75] M. A. Morrison and J. Brillhart. “A Method of Factoring and the Factorization of F_7 ”. In: Mathematics of Computation 29.129 (1975), pp. 183–205. DOI: 10.2307/2005475.