# GROUP 3 CAPSTONE PROJECT

AI GARDENING PROBLEM

#### Members



Đỗ Tuấn Minh 20200390



Nguyễn Kim Tuyến 20205196



Hoàng Huy Chiến 20200084



Tô Thái Dương 20205180



Đinh Ngọc Hạnh Trang 20204928

# **CONTENTS**

- Problem description
- Detailed description of the problem
- A\* search algorithm for searching
- BFS for searching
- DFS for traversing
- Summary of algorithms



There is a m $\times$ n square garden. We denote the location of a square by (x, y). Each square contains exactly one of the following:

Scarecrow (Denoted by		by 0)
-----------------------	--	-------

Warehouse (Only at (0, 0), denoted by 1)

Newly planted seed (Denoted by 2)

Live tree (Denoted by 3)

Dead tree (Denoted by 4)

Χ	0	1	2

1	2	3
4	0	2
4	0	3

.



An AI machine, starts at the warehouse (0, 0), has following functionalities:

- If the AI machine gets to the newly planted seed, it will water it
- If the AI machine gets to the live tree, it will pick all the fruit and bring it back to the warehouse.
- If the AI machine gets to the dead tree, it will remove the dead tree, bring it back to the warehouse, then take a new seed to the previous spot, plant and water it (plant and water at the same time).
- The AI machine can not go pass the scarecrow.

	2	3
4	0	2
4	0	3

1

0



#### The goal is to:

- Water all the newly planted seed
- Bring all the fruit back to the warehouse
- Bring all the dead trees back to the warehouse.
- Plant and water all the new tree in the dead tree spot.

With a minimum number of actions.

3

Χ

0



#### The goal is to:

- Water all the newly planted seed
- Bring all the fruit back to the warehouse
- Bring all the dead trees back to the warehouse.
- Plant and water all the new tree in the dead tree spot.

With a minimum number of actions.

If it is not possible, print out "Can not find the path".

( 0 1

0

	2	3
4	2	0
4	0	3



#### DETAILED DESCRIPTION OF THE PROBLEM

The environment is known, partially observable, deterministic, dynamic and discrete. Thus, the solution is always a fixed sequence of actions.

There are two main parts in this problem to deal with:

- Searching: to find the shortest path in the matrix (m×n square garden).
- Traversing: to visit and do the task in every cell (except the scarecrow).

So, we have to combine at least two algorithms.

Already get the fruit, now let's find the shortest path back to the warehouse

\_

0

1	2	
4	0	2
4	0	3

```
function heuristic(a, b)
     return abs(a[0] - b[0]) + abs(a[1] - b[1])
function init()
     OPEN = \{\}
     CLOSED = {}
     add start to OPEN
     set g_cost of start to 0
     set h_cost of start to heuristic(start, target)
     set f_cost of start to g_cost(start) + h_cost(start)
```

We choose heuristic function is the Manhattan distance (the sum of the absolute differences between the two vectors) between two nodes because:

- It is admissible
- Since we can only move the blocks 1 at a time and in only one of 4 directions (right, left, top, bottom), the optimal scenario for each block is that it has a clear, unobstructed path to its goal state

```
function heuristic(a, b)
     return abs(a[0] - b[0]) + abs(a[1] - b[1])
function init()
     OPEN = \{\}
     CLOSED = {}
     add start to OPEN
     set g cost of start to 0
     set h cost of start to heuristic(start, target)
     set f cost of start to g cost(start) + h cost(start)
```

We choose heuristic function is the Manhattan distance (the sum of the absolute differences between the two vectors) between two nodes because:

- It is admissible
- Since we can only move the blocks 1 at a time and in only one of 4 directions (right, left, top, bottom), the optimal scenario for each block is that it has a clear, unobstructed path to its goal state

Initialize a set to store nodes to be traversed (OPEN) and a set to store already traverse nodes (CLOSED).

```
function heuristic(a, b)
     return abs(a[0] - b[0]) + abs(a[1] - b[1])
function init()
     OPEN = \{\}
     CLOSED = {}
     add start to OPEN
     set g cost of start to 0
     set h cost of start to heuristic(start, target)
     set f cost of start to g cost(start) + h cost(start)
```

We choose heuristic function is the Manhattan distance (the sum of the absolute differences between the two vectors) between two nodes because:

- It is admissible
- Since we can only move the blocks 1 at a time and in only one of 4 directions (right, left, top, bottom), the optimal scenario for each block is that it has a clear, unobstructed path to its goal state

Initialize a set to store nodes to be traversed (OPEN) and a set to store already traverse nodes (CLOSED).

Then, set all cost to node start...

function heuristic(a, b)

return abs(a[0] - b[0]) + abs(a[1] - b[1])

function init()

**OPEN** = {}

CLOSED = {}

add start to OPEN

set g\_cost of start to 0

set h cost of start to heuristic(start, target)

set f\_cost of start to g\_cost(start) + h\_cost(start)

```
function astar(start, target)
        init()
        while OPEN is not empty
                current = node in OPEN with lowest f_cost
                remove current from OPEN
                if current is target
                         return
                add current to CLOSED
                foreach neighbor of current
                temp_cost = g_cost(current) + 1
                if neighbor is not Obstacle
                         if neighbor in OPEN or CLOSED and g_cost(neighbor) <= temp_cost
                                 continue
                         if neighbor in CLOSED
                                 remove neighbor from CLOSED
                                 add neighbor to OPEN
                         if neighbor not in OPEN or CLOSED
                                 add neighbor to OPEN
                                 set f_cost of neighbor
                                 set parent of neighbor to current
```

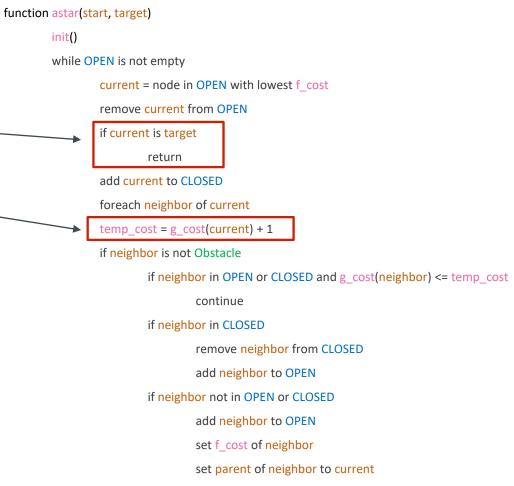
#### For each loop:

Check if node current is the target or not

```
function astar(start, target)
        init()
        while OPEN is not empty
                current = node in OPEN with lowest f_cost
                remove current from OPEN
                if current is target
                         return
                add current to CLOSED
                foreach neighbor of current
                temp_cost = g_cost(current) + 1
                if neighbor is not Obstacle
                         if neighbor in OPEN or CLOSED and g_cost(neighbor) <= temp_cost
                                 continue
                         if neighbor in CLOSED
                                 remove neighbor from CLOSED
                                 add neighbor to OPEN
                         if neighbor not in OPEN or CLOSED
                                 add neighbor to OPEN
                                 set f_cost of neighbor
                                 set parent of neighbor to current
        return FALSE
```

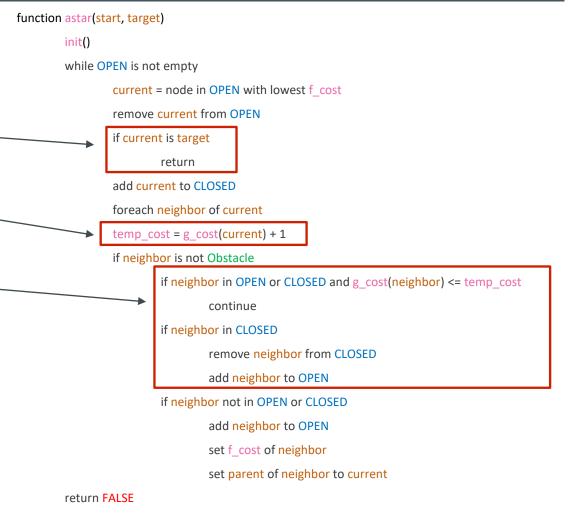
#### For each loop:

- Check if node current is the target or not
- Temporarily store the new g\_cost of neighbor one larger than that of current (the distance of two adjacent nodes is one).



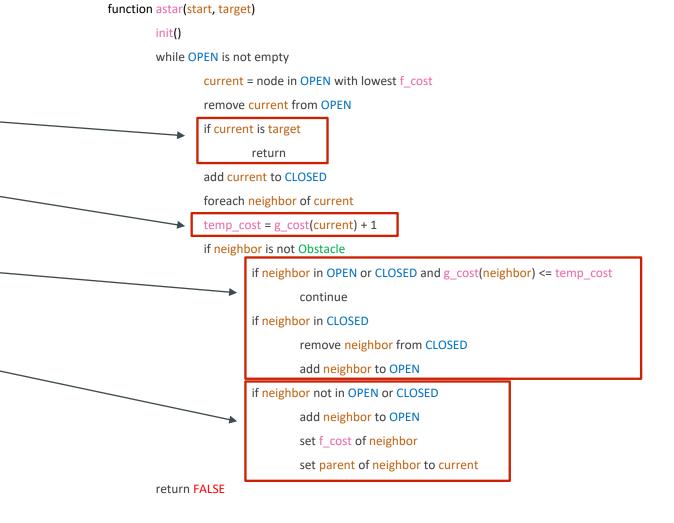
#### For each loop:

- Check if node current is the target or not
- Temporarily store the new g\_cost of neighbor one larger than that of current (the distance of two adjacent nodes is one).
- Compare the old and new g\_cost to make \_\_ decision about neighbor.



#### For each loop:

- Check if node current is the target or not
- Temporarily store the new g\_cost of neighbor one larger than that of current (the distance of two adjacent nodes is one).
- Compare the old and new g\_cost to make \_ decision about neighbor.
- If neighbor is neither in two sets (so haven't has g\_cost yet), add it to OPEN, set all cost to it and set its parent node in the path.



#### function astar(start, target) init() while **OPEN** is not empty For each loop: current = node in OPEN with lowest f cost remove current from OPEN Check if node current is the target or not if current is target return Temporarily store the new g cost of neighbor add current to CLOSED one larger than that of current (the distance of foreach neighbor of current two adjacent nodes is one). temp\_cost = g\_cost(current) + 1 if neighbor is not Obstacle Compare the old and new g\_cost to make \_ if neighbor in OPEN or CLOSED and g\_cost(neighbor) <= temp\_cost decision about neighbor. continue if neighbor in CLOSED If neighbor is neither in two sets (so haven't remove neighbor from CLOSED has g cost yet), add it to OPEN, set all cost to add neighbor to OPEN it and set its parent node in the path. if neighbor not in OPEN or CLOSED add neighbor to OPEN Return FALSE when can not find a path set f cost of neighbor set parent of neighbor to current return FALSE

```
function BFS(start, target)
     VISITED = []
     QUEUE
     add start to QUEUE and VISITED
     set parent of start to None
     while QUEUE is not empty
           remove front from QUEUE store it to current
           if current is target
                 return path
     foreach neighbor of current
           if neighbor not in VISITED
                 add neighbor to QUEUE
                 set parent of neighbor to current
                 add neighbor to VISITED
```

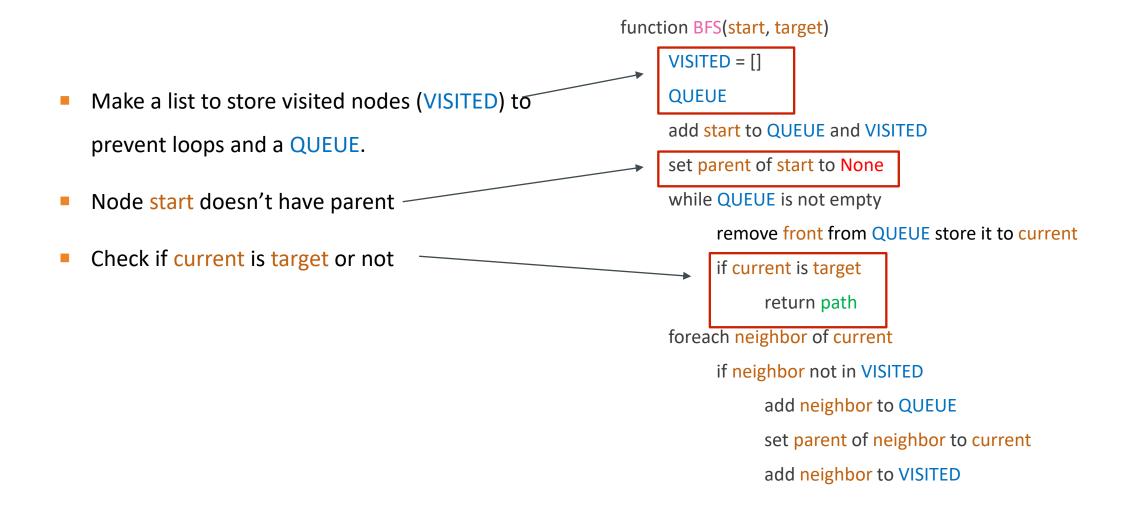
 Make a list to store visited nodes (VISITED) to prevent loops and a QUEUE.

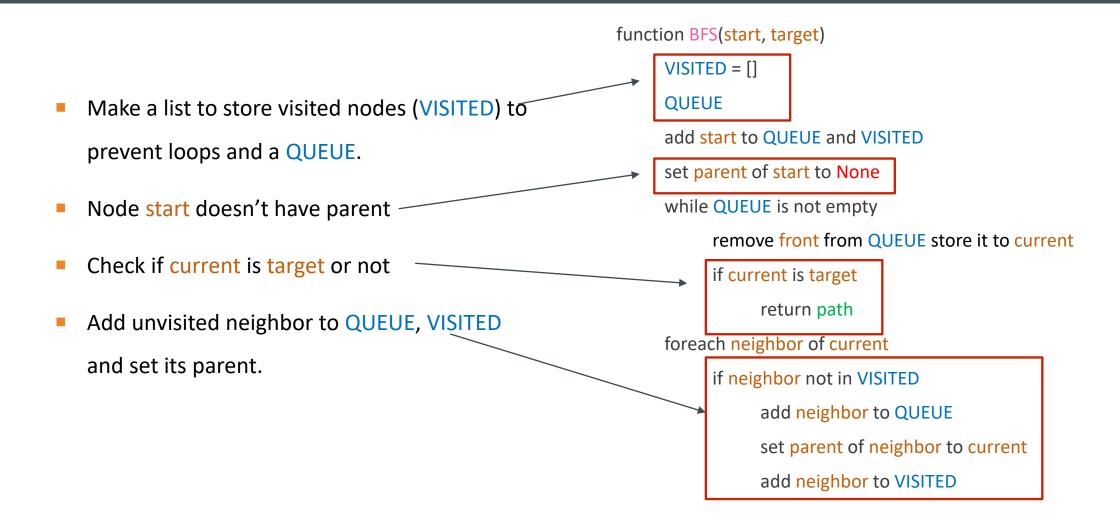
```
function BFS(start, target)
     VISITED = []
     QUEUE
     add start to QUEUE and VISITED
     set parent of start to None
     while QUEUE is not empty
           remove front from QUEUE store it to current
           if current is target
                 return path
     foreach neighbor of current
           if neighbor not in VISITED
                 add neighbor to QUEUE
                 set parent of neighbor to current
                 add neighbor to VISITED
```

 Make a list to store visited nodes (VISITED) to prevent loops and a QUEUE.

Node start doesn't have parent

function BFS(start, target) VISITED = [] QUEUE add start to QUEUE and VISITED set parent of start to None while QUEUE is not empty remove front from QUEUE store it to current if current is target return path foreach neighbor of current if neighbor not in VISITED add neighbor to QUEUE set parent of neighbor to current add neighbor to VISITED





```
DFS(s)
        update PERCEPT SEQUENCE
        if visited(s) == TRUE:
                 return
        visited(s) = TRUE
        if s is Obstacle:
                 return
        if distance(last_visited_node, s) != 1
                 add shortest_path(last_visited_node, s) to PATH
        else
                 add s to PATH
        add move(PATH) to ACTION
        add do_task(s) to ACTION
        print PERCEPT SEQUENCE and ACTION
        foreach neighbor of s
                 if visited(neighbor) == FALSE
                         DFS(neighbor)
forall cell in MATRIX
        visited(cell) = FALSE
forall cell in MATRIX
        DFS(cell)
```

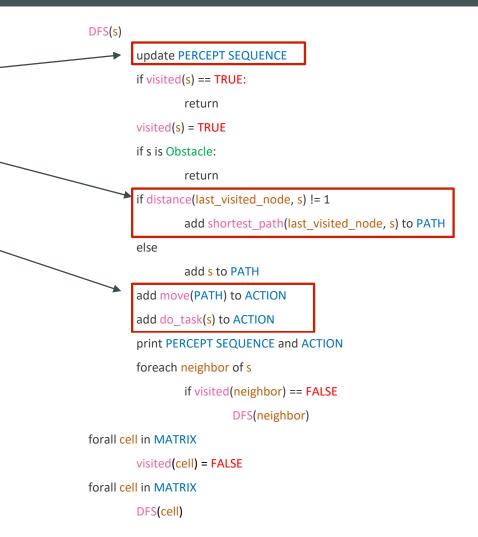
Update the matrix that the AI machine sees

```
DFS(s)
        update PERCEPT SEQUENCE
        if visited(s) == TRUE:
                 return
        visited(s) = TRUE
        if s is Obstacle:
                 return
        if distance(last_visited_node, s) != 1
                 add shortest_path(last_visited_node, s) to PATH
        else
                 add s to PATH
        add move(PATH) to ACTION
        add do_task(s) to ACTION
        print PERCEPT SEQUENCE and ACTION
        foreach neighbor of s
                 if visited(neighbor) == FALSE
                         DFS(neighbor)
forall cell in MATRIX
        visited(cell) = FALSE
forall cell in MATRIX
        DFS(cell)
```

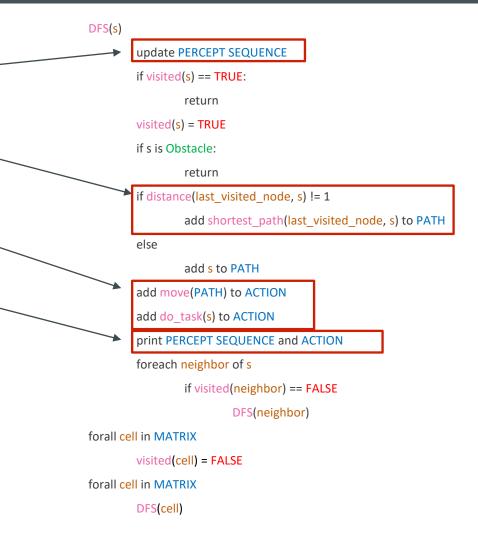
- Update the matrix that the AI machine sees
- If DFS a cell that is not adjacent to the latest one, then find the shortest path between them.

```
DFS(s)
        update PERCEPT SEQUENCE
         if visited(s) == TRUE:
                 return
        visited(s) = TRUE
        if s is Obstacle:
                 return
        if distance(last_visited_node, s) != 1
                 add shortest path(last visited node, s) to PATH
        else
                 add s to PATH
        add move(PATH) to ACTION
        add do_task(s) to ACTION
        print PERCEPT SEQUENCE and ACTION
        foreach neighbor of s
                 if visited(neighbor) == FALSE
                         DFS(neighbor)
forall cell in MATRIX
        visited(cell) = FALSE
forall cell in MATRIX
        DFS(cell)
```

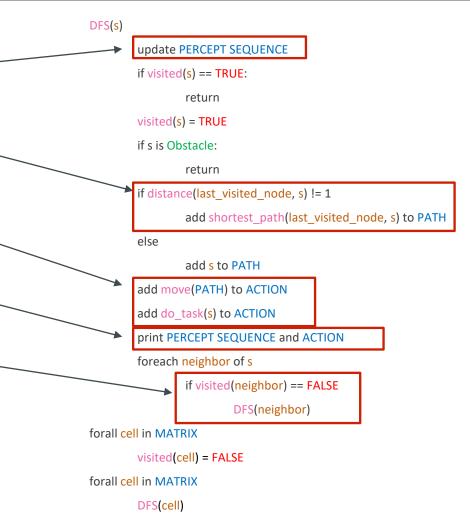
- Update the matrix that the AI machine sees
- If DFS a cell that is not adjacent to the latest one, then find the shortest path between them.
- Update ACTION needed to get to the cell through PATH and to finish tasks in the cell.



- Update the matrix that the AI machine sees
- If DFS a cell that is not adjacent to the latest one, then find the shortest path between them.
- Update ACTION needed to get to the cell through PATH and to finish tasks in the cell.
- Print out the PERCEPT SEQUENCE and list of ACTION for each run.



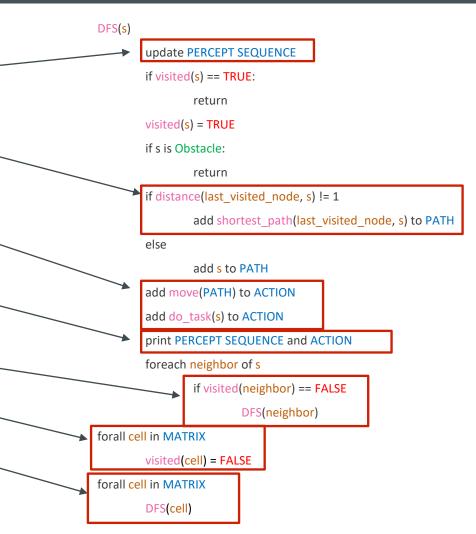
- Update the matrix that the AI machine sees
- If DFS a cell that is not adjacent to the latest one, then find the shortest path between them.
- Update ACTION needed to get to the cell through PATH and to finish tasks in the cell.
- Print out the PERCEPT SEQUENCE and list of ACTION for each run.
- Continue run DFS on unvisited neighbors



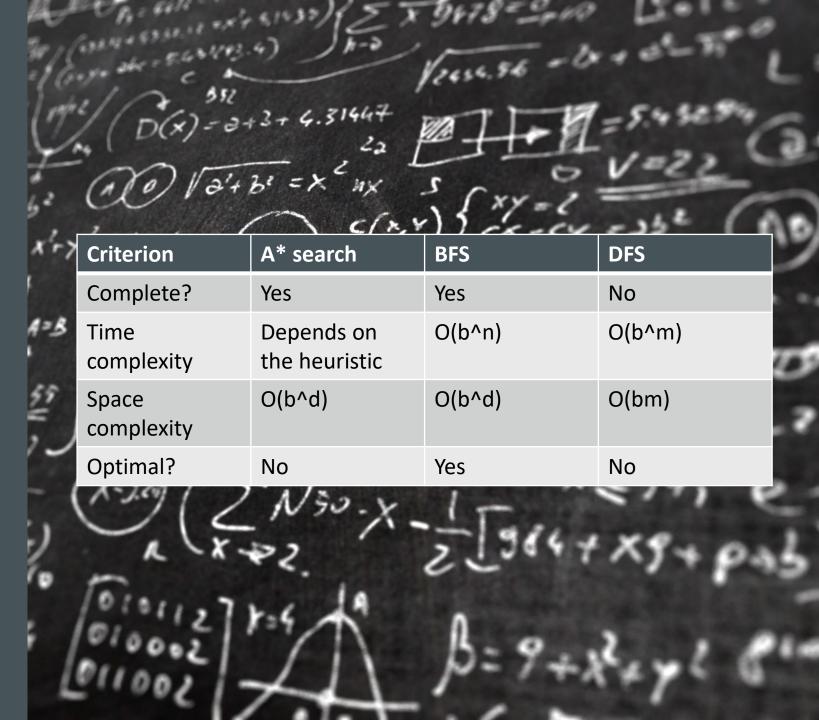
DFS(s) update PERCEPT SEQUENCE Update the matrix that the AI machine sees if visited(s) == TRUE: return If DFS a cell that is not adjacent to the latest visited(s) = TRUE one, then find the shortest path between if s is Obstacle: them. return if distance(last visited node, s) != 1 Update ACTION needed to get to the cell add shortest path(last visited node, s) to PATH through PATH and to finish tasks in the cell. else add s to PATH Print out the PERCEPT SEQUENCE and list of add move(PATH) to ACTION ACTION for each run. add do task(s) to ACTION print PERCEPT SEQUENCE and ACTION Continue run DFS on unvisited neighbors foreach neighbor of s if visited(neighbor) == FALSE Visited check is initial set FALSE to all cells DFS(neighbor) forall cell in MATRIX visited(cell) = FALSE forall cell in MATRIX

DFS(cell)

- Update the matrix that the AI machine sees
- If DFS a cell that is not adjacent to the latest one, then find the shortest path between them.
- Update ACTION needed to get to the cell through PATH and to finish tasks in the cell.
- Print out the PERCEPT SEQUENCE and list of ACTION for each run.
- Continue run DFS on unvisited neighbors
- Visited check is initial set FALSE to all cells
- Run DFS on all cell to make sure no cell is unvisited.



# SUMMARY OF ALGORITHMS



Q&A?

# THANK YOU FOR LISTENING

