## CSC3150 Operating System

# Assignment Report #4

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1. Design (Definition of variables)

fs->volume space allocation:

- (1) SUPERBLOCK: first 4KB (superblock size)
- (2) FCB: 32KB after SUPERBLOCK

A maximum of 1k files can be collected, each file information has 32B

- 1) 0-19 B stored name
- 2) 20-21B stored create time
- 3) 22-23B stored modified time
- 4) 24-25B stored address (of super block number)
- 5) 26-27B stored size (The units are B)

Specially, in bonus, for directory it stored all of the file name length under the directory.

- 6) 28B stored valid bit (ff means invalid, 00 means valid, 01 means in sort) Specially, in bonus, f0 means valid as directory, x1means under the now directory (when sorting), x2 means has sort (when sorting).
- 7) 29-30B stored parent point in bonus, when its parent is root or no parent, these bits should be 0xff
- (3) Contains of file: 1024K for total. Stored files with block is 32B.

gtime: gtime is increased every time a file (or a directory) is updated gsize: gsize marks the number of blocks used.

file number: Records the number of existing files

now\_directory: Mark the current directory, which is 0xffff when in the root directory

- 2. Design (Definition of function)
  - (1) FCB\_init: Set the valid bit of all FCBs to 0xff

    Specially, in bonus, also set the parent bit of FCB to 0xff
  - (2) SUPERBLOCK init: set all of superblock to zero.
  - (3) string print: print string with pointer s.
  - (4) check name: check strings s1 and s2 same or not
  - (5) update\_parent\_size: Change the size of the directory when the file name changes

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(6) update\_file\_name: Try writing s as the name of the fp file to check if the name is too long

- (7) Modified FCB: Update FCB of file based on size and gtime.
- (8) get len: Returns the length of the string s
- (9) set\_superblock\_bit: Changes the superblock information for the addr position based on bit
- (10) Remove\_file: Clear the storage space of the file fp, and move the storage location of the following file forward
- (11)Cmp: The preference of file b is true according to the judgment of file a and file b which has a higher priority in the arrangement according to op.
- (12) Output parent: Rely on recursive output file paths
- (13) Others function: Definition as required
- 3. Environment

Use Cluster.

System Type x86\_64

Opearing System CentOS Linux release 7.5.1804

CPU Intel(R) Xeon(R) Silver 4210R CPU @ 2.40GHz

20 Cores, 40 Threads

Memory 100GB RAM

GPU Nvidia Quadro RTX 4000 GPU x 1

CUDA 11.7

GCC Red Hat 7.3.1-5

CMake 3.14.1

- 4. Screenshot
  - (1) Case1:

```
===sort by modified time===
t.txt
b.txt
===sort by file size===
t.txt 32
b.txt 32
===sort by file size===
t.txt 32
b.txt 12
===sort by modified time===
b.txt
t.txt
===sort by file size===
b.txt
```

### (2) Case 2

```
41 ===sort by modified time===
42 t.txt
   b.txt
43
44 ===sort by file size===
45 t.txt 32
   b.txt 32
46
47
    ===sort by file size===
   t.txt 32
48
49
    b.txt 12
50
    ===sort by modified time===
51
    t.txt
53
    ===sort by file size===
    b.txt 12
   ===sort by file size===
55
    *ABCDEFGHIJKLMNOPQR 33
56
57 )ABCDEFGHIJKLMNOPQR 32
58 (ABCDEFGHIJKLMNOPQR 31
59
   'ABCDEFGHIJKLMNOPQR 30
60 &ABCDEFGHIJKLMNOPQR 29
61 %ABCDEFGHIJKLMNOPQR 28
62 $ABCDEFGHIJKLMNOPQR 27
63 #ABCDEFGHIJKLMNOPQR 26
64 "ABCDEFGHIJKLMNOPQR 25
65 !ABCDEFGHIJKLMNOPQR 24
66 b.txt 12
67 ===sort by modified time===
68 *ABCDEFGHIJKLMNOPQR
69 ) ABCDEFGHIJKLMNOPQR
70 (ABCDEFGHIJKLMNOPQR
71 'ABCDEFGHIJKLMNOPQR
72 &ABCDEFGHIJKLMNOPQR
73
   b.txt
74
```

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### (3) Case 3(Fore and aft)

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2111

b.txt 12

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### (4) Case 4 (Fore and aft)

58 b.txt 12

===sort by file size===

29	triggering gc		
30	===sort by modified time===		
31	1024-block-1023		
32	1024-block-1022	1046	1024-D10CK-0008
32	1024-D10CK-1022	1047	1024-block-0007
33	1024-block-1021	1047	1024 DIOCK 0007
2.4	1004 13 1 1000	1048	1024-block-0006
34	1024-block-1020	1010	1024   1 - 1 0005
35	1024-block-1019	1049	1024-block-0005
22	1024-D10CK-1019	1050	1024-block-0004
36	1024-block-1018		
27	1004   1   1017	1051	1024-block-0003
37	1024-block-1017	1052	1024 blask 0002
38	1024-block-1016	1052	1024-block-0002
50	1024 DIOCK 1010	1053	1024-block-0001
39	1024-block-1015		
	4004 13 1 4044	1054	1024-block-0000

### (5) Case bonus (aft)

```
109 /app/soft
110 ===sort by file size===
111
    B.txt 1024
112 C.txt 1024
113 D.txt 1024
114
    A.txt 64
     ===sort by file size===
115
116
     a.txt 64
117 b.txt 32
118
    soft 24 d
119
     /app
120
     ===sort by file size===
121
    t.txt 32
122 b.txt 32
     app 17 d
123
     ===sort by file size===
124
125
     a.txt 64
126 b.txt 32
127
    ===sort by file size===
    t.txt 32
128
129
     b.txt 32
130 app 12 d
```

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#### 5. Flow of some function

- (1) fs\_open
  - Check whether there is a valid file named s in the FCB.
  - 2 If yes, return the serial number
  - 3 Else, Look for an invalid FCB
  - 4 Use update file name to write the file name.
  - © Use modified\_FCB to update the update time and file size in the FCB information.
  - @ gtime, file\_number is increased, and the FCB is set to valid
  - In bonus, set the father bit as now directory
- (2) fs\_read
  - Terrify that the file is readable. If not, output ERROR
  - ② Use the FCB address bit to find the corresponding file location in the storage
  - 3 Copy the size character to output
- (3) fs\_write
  - Verify that the file is readable. If not, output ERROR
  - 2) Check whether the input file size is too large. If yes, output ERROR.
  - Check whether the original location needs to be cleared (the original size is not 0 and the new input occupies a different block size from the original occupied block size)
    - 1) The previous size was 0: stored after the occupied block, after the address bit is occupied
    - 2) The previous size and the current size occupy the same number of blocks: data is directly rewritten to the original address
    - 3) Use remove\_file to erase the previous storage, rewrite the data at the end, and update the address
  - 4 After writing, update the data in FCB (size, address, modified time)
- (4) fs\_gsys
  - PWD: Use output parent to write out the path
  - 2 CP\_D: Update now\_directory with the parent bit in FCB of now directory

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- LS\_D/LS\_S: Sort by marking files that need to be sorted and files that have already been sorted, relying on the cmp function to compare sizes. The valid bit is restored after the sorting is complete
  - In bonus, you first do a markup to filter out the files and directories in the now directory
- RM: Locate the file in FCB, use remove\_file to clear the address, and set FCB to invalid
- MKDIR: Find the invaild page in FCB, write the information, and set the valid bit to 0xf0
- © CD: Search for a directory in FCB where the valid bit is 0xf0(is directory) and the parent node is now\_directory. Update now\_directory to the directory.
- RM\_RF: Locate the folder in the subfolder of now\_directory and search for all the subfiles and subfolders of that folder. Subfiles are deleted with RM, and subfolders are recursively deleted with RM\_RF instruction and subfolder name. Finally, set the folder to invalid.

The bonus content is highlighted