

M6809PM (AD)

MC6809-MC6809E
8-BIT MICROPROCESSOR
PROGRAMMING MANUAL

Original Issue: March 1, 1981

TABLE OF CONTENTS

<i>Paragraph No.</i>	<i>Title</i>	<i>Page No.</i>
SECTION 1 GENERAL DESCRIPTION		
1.1	Introduction.....	1-1
1.2	Features.....	1-1
1.3	Software Features.....	1-2
1.4	Programming Model.....	1-3
1.5	Index Registers (X, Y)	1-3
1.6	Stack Pointer Registers (U, S)	1-3
1.7	Program Counter (PC)	1-4
1.8	Accumulator Registers (A, B, D)	1-4
1.9	Direct Page Register (DP)	1-4
1.10	Condition Code Register (CC)	1-4
1.10.1	Condition Code Bits	1-5
1.10.1.1	Half Carry (H), Bit 5.....	1-5
1.10.1.2	Negative (N), Bit 3	1-5
1.10.1.3	Zero (Z), Bit 2.....	1-5
1.10.1.4	Overflow (V), Bit 1	1-5
1.10.1.5	Carry (C), Bit 0	1-5
1.10.2	Interrupt Mask Bits and Stacking Indicator	1-5
1.10.2.1	Fast Interrupt Request Mask (F), Bit 6	1-5
1.10.2.2	Interrupt Request Mask (I), Bit 4	1-5
1.10.2.3	Entire Flag (E), Bit 7	1-6
1.11	Pin Assignments and Signal Description.....	1-6
1.11.1	MC6809 Clocks	1-6
1.11.1.1	Oscillator (EXTAL, XTAL).....	1-6
1.11.1.2	Enable (E).....	1-7
1.11.1.3	Quadrature (Q)	1-7
1.11.2	MC6809E Clocks (E and Q)	1-7
1.11.3	Three State Control (TSC) (MC6809E)	1-7
1.11.4	Last Instruction Cycle (LIC) (MC6809E)	1-7
1.11.5	Address Bus (A0-A15).....	1-7
1.11.6	Data Bus (D0-D7)	1-7
1.11.7	Read/Write (R/W).....	1-8
1.11.8	Processor State Indicators (BA, BS)	1-8
1.11.8.1	Normal	1-8
1.11.8.2	Interrupt or Reset Acknowledge.....	1-8
1.11.8.3	Sync Acknowledge.....	1-8

TABLE OF CONTENTS (CONTINUED)

<i>Paragraph No.</i>	<i>Title</i>	<i>Page No.</i>
1.11.8.4	Halt/Bus Grant	1-8
1.11.9	Reset (RESET)	1-9
1.11.10	Interrupts	1-9
1.11.10.1	Non-Maskable Interrupt (NMI)	1-9
1.11.10.2	Fast Interrupt Request (FIRQ)	1-9
1.11.10.3	Interrupt Request (IRQ)	1-9
1.11.11	Memory Ready (MRDY) (MC6809)	1-9
1.11.12	Advanced Valid Memory Address (AVMA) (MC6809E)	1-10
1.11.13	Halt (HALT)	1-10
1.11.14	Direct Memory Access/Bus Request (DMA/BREQ) (MC6809)	1-10
1.11.15	Busy (MC6809E)	1-10
1.11.16	Power	1-11

SECTION 2 ADDRESSING MODES

2.1	Introduction	2-1
2.2	Addressing Modes	2-1
2.2.1	Inherent	2-1
2.2.2	Immediate	2-1
2.2.3	Extended	2-2
2.2.4	Direct	2-2
2.2.5	Indexed	2-2
2.2.5.1	Constant Offset from Register	2-2
2.2.5.2	Accumulator Offset from Register	2-3
2.2.5.3	Autoincrement/Decrement from Register	2-3
2.2.5.4	Indirection	2-4
2.2.5.5	Extended Indirect	2-4
2.2.5.6	Program Counter Relative	2-4
2.2.6	Branch Relative	2-4

SECTION 3 INTERRUPT CAPABILITIES

3.1	Introduction	3-1
3.2	Non-Maskable Interrupt (NMI)	3-1
3.3	Fast Maskable Interrupt Request (FIRQ)	3-2
3.4	Normal Maskable Interrupt Request (IRQ)	3-2
3.5	Software Interrupts (SWI, SWI2, SWI3)	3-2

TABLE OF CONTENTS (CONCLUDED)

<i>Paragraph No.</i>	<i>Title</i>	<i>Page No.</i>
----------------------	--------------	-----------------

SECTION 4 PROGRAMMING

4.1	Introduction.....	4-1
4.1.1	Position-Independence.....	4-1
4.1.2	Modular Programming.....	4-1
4.1.2.1	Local Storage.....	4-1
4.1.2.2	Global Storage.....	4-2
4.1.3	Reentrancy/Recursion.....	4-2
4.2	M6809 Capabilities.....	4-2
4.2.1	Module Construction	4-2
4.2.1.1	Parameters.....	4-3
4.2.1.2	Local Storage.....	4-3
4.2.1.3	Global Storage.....	4-3
4.2.2	Position-Independent Code	4-4
4.2.3	Reentrant Programs.....	4-5
4.2.4	Recursive Programs.....	4-5
4.2.5	Loops	4-5
4.2.6	Stack Programming	4-6
4.2.6.1	M6809 Stacking Operations.....	4-6
4.2.6.2	Subroutine Linkage.....	4-7
4.2.6.3	Software Stacks	4-8
4.2.7	Real Time Programming.....	4-8
4.3	Program Documentation	4-8
4.4	Instruction Set	4-9

APPENDIX A INSTRUCTION SET DETAILS

A.1	Introduction	A-1
A.2	Notation	A-1
	Instructions (listed in alphabetical order).....	A-3

APPENDIX B ASSIST09 MONITOR PROGRAM

B.1	General Description.....	B-1
B.2	Implementation Requirements	B-1
B.3	Interrupt Control	B-2
B.4	Initialization.....	B-3

TABLE OF CONTENTS (CONTINUED)

<i>Paragraph No.</i>	<i>Title</i>	<i>Page No.</i>
B.5	Input/Output Control	B-4
B.6	Command Format	B-4
B.7	Command List	B-5
B.8	Commands..... Breakpoint	B-5 B-6
	Call.....	B-6
	Display.....	B-7
	Encode	B-7
	Go.....	B-8
	Load	B-8
	Memory	B-9
	Null	B-10
	Offset.....	B-10
	Punch.....	B-11
	Register.....	B-11
	Stlevel.....	B-12
	Trace.....	B-12
	Verify	B-13
	Window	B-13
B.9	Services..... BKPT	B-14 B-15
	INCHP.....	B-15
	MONITR.....	B-16
	OUTCH	B-17
	OUT2HS.....	B-17
	OUT4HS.....	B-18
	PAUSE	B-18
	PCRLF	B-19
	PDATA	B-19
	PDATA1.....	B-20
	SPACE	B-21
	VTRSW	B-21
B.10	Vector Swap Service..... .ACIA.....	B-22 B-23
	.AVTBL.....	B-23
	.BSDTA	B-24
	.BSOFF	B-24
	.BSON	B-25
	.CIDTA	B-25
	.CIOFF	B-26
	.CION	B-26
	.CMDL1.....	B-27
	.CMDL2.....	B-28

TABLE OF CONTENTS (CONTINUED)

<i>Paragraph No.</i>	<i>Title</i>	<i>Page No.</i>
	.CODTA.....	B-28
	.COOFF.....	B-29
	.COON.....	B-29
	.ECHO.....	B-30
	.FIRQ.....	B-30
	.HSDATA.....	B-31
	.IRQ.....	B-31
	.NMI.....	B-32
	.PAD.....	B-32
	.PAUSE.....	B-33
	.PTM.....	B-33
	.RESET.....	B-34
	.RSVD.....	B-34
	.SWI.....	B-35
	.SWI2.....	B-35
	.SWI3.....	B-36
B.11	Monitor Listing	B-37

APPENDIX C MACHINE CODE TO INSTRUCTION CROSS REFERENCE

C.1	Introduction	C-1
-----	--------------------	-----

APPENDIX D PROGRAMMING AID

D.1	Introduction	D-1
-----	--------------------	-----

APPENDIX E ASCII CHARACTER SET

E.1	Introduction	E-1
E.2	Character Representation and Code Identification	E-1
E.3	Control Characters	E-2
E.4	Graphic Characters	E-2

TABLE OF CONTENTS (CONTINUED)

<i>Paragraph No.</i>	<i>Title</i>	<i>Page No.</i>
	APPENDIX F OPCODE MAP	
F.1	Introduction	F-1
F.2	Opcode Map.....	F-1

APPENDIX G PIN ASSIGNMENTS

G.1	Introduction.....	G-1
-----	-------------------	-----

APPENDIX H CONVERSION TABLES

H.1	Introduction.....	H-1
H.2	Powers of 2; Powers of 16.....	H-1
H.3	Hexadecimal and Decimal Conversion.....	H-2
H.3.1	Converting Hexadecimal to Decimal	H-2
H.3.2	Converting Decimal to Hexadecimal	H-2

LIST OF ILLUSTRATIONS

<i>Figure No.</i>	<i>Title</i>	<i>Page No.</i>
1-1	Programming Model.....	1-3
1-2	Condition Code Register	1-4
1-3	Processor Pin Assignments	1-6
2-1	Postbyte Usage for EXG/TFR, PSH/PUL Instructions.....	2-2
3-1	Interrupt Processing Flowchart.....	3-5
4-1	Stacking Order	4-7
B-1	Memory Map	B-2
E-1	ASCII Character Set.....	E-1
G-1	Pin Assignments.....	G-1

LIST OF TABLES

<i>Table No.</i>	<i>Title</i>	<i>Page No.</i>
1-1	BA/BS Signal Encoding.....	1-8
2-1	Postbyte Usage for Indexed Addressing Modes	2-3
3-1	Interrupt Vector Locations	3-1
4-1	Instruction Set	4-9
4-2	8-Bit Accumulator and Memory Instructions.....	4-11
4-3	16-Bit Accumulator and Memory Instructions.....	4-12
4-4	Index/Stack Pointer Instructions.....	4-12
4-5	Branch Instructions	4-13
4-6	Miscellaneous Instructions	4-13
A-1	Operation Notation.....	A-1
A-2	Register Notation.....	A-2
B-1	Command List	B-5
B-2	Services.....	B-14
B-3	Vector Table Entries	B-22
C-1	Machine Code to Instruction Cross Reference.....	C-2
D-1	Programming Aid	D-1
E-1	Control Characters	E-2
E-2	Graphic Characters	E-3
F-1	Opcode Map.....	F-2
F-2	Indexed Addressing Mode Data	F-3
H-1	Powers of 2; Powers of 16.....	H-1
H-2	Hexadecimal and Decimal Conversion Chart.....	H-2

SECTION 1 GENERAL DESCRIPTION

1.1 INTRODUCTION

This section contains a general description of the Motorola MC6809 and MC6809E Microprocessor Units (MPU). Pin assignments and a brief description of each input/output signal are also given. The term MPU, processor, or M6809 will be used throughout this manual to refer to both the MC6809 and MC6809E processors. When a topic relates to only one of the processors, that specific designator (MC6809 or MC6809E) will be used.

1.2 FEATURES

The MC6809 and MC6809E microprocessors are greatly enhanced, upward compatible, computationally faster extensions of the MC6800 microprocessor.

Enhancements such as additional registers (a Y index register, a U stack pointer, and a direct page register) and instructions (such as MUL) simplify software design. Improved addressing modes have also been implemented.

Upward compatibility is guaranteed as MC6800 assembly language programs may be assembled using the Motorola MC6809 Macro Assembler. This code, while not as compact as native M6809 code, is, in most cases, 100% functional.

Both address and data are available from the processor earlier in an instruction cycle than from the MC6800 which simplifies hardware design. Two clock signals, E (the MC6800 ϕ_2) and a new quadrature clock Q (which leads E by one-quarter cycle) also simplify hardware design.

A memory ready (MRDY) input is provided on the MC6809 for working with slow memories. This input stretches both the processor internal cycle and direct memory access bus cycle times but allows internal operations to continue at full speed. A direct memory access request (DMA/BREQ) input is provided for immediate memory access or dynamic memory refresh operations; this input halts the internal MC6809 clocks. Because the processor's registers are dynamic, an internal counter periodically recovers the bus from direct memory access operations and performs a true processor refresh cycle to allow unlimited length direct memory access operation. An interrupt acknowledge signal is available to allow development of vectoring by interrupt device hardware or detection of operating system calls.

Three prioritized, vectored, hardware interrupt levels are available: non-maskable, fast, and normal. The highest and lowest priority interrupts, non-maskable and interrupt request respectively, are the normal interrupts used in the M6800 family. A new interrupt on this processor is the fast interrupt request which provides faster service to its interrupt input by only stacking the program counter and condition code register and then servicing the interrupt.

Modern programming techniques such as position-independent, system independent, and reentrant programming are readily supported by these processors.

A Memory Management Unit (MMU), the MC6829, allows a M6809 based system to address a two megabyte memory space. Note: An arbitrary number of tasks may be supported — slower — with software.

This advanced family of processors is compatible with all M6800 peripheral parts.

1.3 SOFTWARE FEATURES

Some of the software features of these processors are itemized in the following paragraphs. Programs developed for the MC6800 can be easily converted for use with the MC6809 or MC6809E by running the source code through a M6809 Macro Assembler or any one of the many cross assemblers that are available.

The addressing modes of any microprocessor provide it with the capability to efficiently address memory to obtain data and instructions. The MC6809 and MC6809E have a versatile set of addressing modes which allow them to function using modern programming techniques.

The addressing modes and instructions of the MC6809 and MC6809E are upward compatible with the MC6800. The old addressing modes have been retained and many new ones have been added.

A direct page register has been added which allows a 256 byte "direct" page anywhere in the 64K logical address space. The direct page register is used to hold the most-significant byte of the address used in direct addressing and decrease the time required for address calculation.

Branch relative addressing to anywhere in the memory map (-32768 to +32767) is available.

Program counter relative addressing is also available for data access as well as branch instructions.

The indexed addressing modes have been expanded to include:

- 0-, 5-, 8-, 16-bit constant offsets,
- 8- or 16-bit accumulator offsets,
- autoincrement/decrement (stack operation).

In addition, most indexed addressing modes may have an additional level of indirection added.

Any or all registers may be pushed on to or pulled from either stack with a single instruction.

A multiply instruction is included which multiplies unsigned binary numbers in accumulators A and B and places the unsigned result in the 16-bit accumulator D. This unsigned multiply instruction also allows signed or unsigned multiple precision multiplication.

1.4 PROGRAMMING MODEL

The programming model (Figure 1-1) for these processors contains five 16-bit and four 8-bit registers that are available to the programmer.

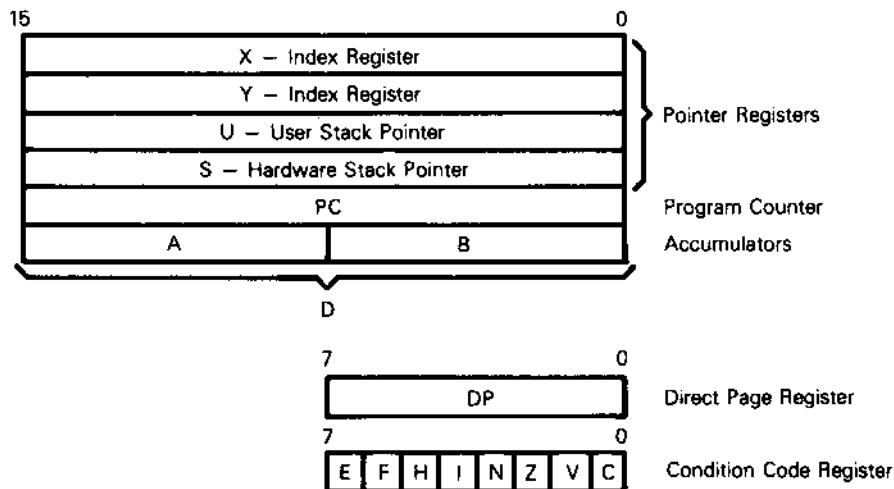


Figure 1-1. Programming Model

1.5 INDEX REGISTERS (X, Y)

The index registers are used during the indexed addressing modes. The address information in an index register is used in the calculation of an effective address. This address may be used to point directly to data or may be modified by an optional constant or register offset to produce the effective address.

1.6 STACK POINTER REGISTERS (U, S)

Two stack pointer registers are available in these processors. They are: a user stack pointer register (U) controlled exclusively by the programmer, and a hardware stack pointer register (S) which is used automatically by the processor during subroutine calls

and interrupts, but may also be used by the programmer. Both stack pointers always point to the top of the stack.

These registers have the same indexed addressing mode capabilities as the index registers, and also support push and pull instructions. All four indexable registers (X, Y, U, S) are referred to as pointer registers.

1.7 PROGRAM COUNTER (PC)

The program counter register is used by these processors to store the address of the next instruction to be executed. It may also be used as an index register in certain addressing modes.

1.8 ACCUMULATOR REGISTERS (A, B, D)

The accumulator registers (A, B) are general-purpose 8-bit registers used for arithmetic calculations and data manipulation.

Certain instructions concatenate these registers into one 16-bit accumulator with register A positioned as the most-significant byte. When concatenated, this register is referred to as accumulator D.

1.9 DIRECT PAGE REGISTER (DP)

This 8-bit register contains the most-significant byte of the address to be used in the direct addressing mode. The contents of this register are concatenated with the byte following the direct addressing mode operation code to form the 16-bit effective address. The direct page register contents appear as bits A15 through A8 of the address. This register is automatically cleared by a hardware reset to ensure M6800 compatibility.

1.10 CONDITION CODE REGISTER (CC)

The condition code register contains the condition codes and the interrupt masks as shown in Figure 1-2.

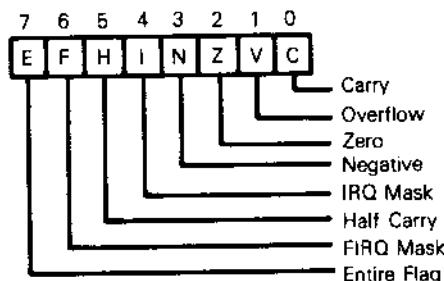


Figure 1-2. Condition Code Register

1.10.1 CONDITION CODE BITS. Five bits in the condition code register are used to indicate the results of instructions that manipulate data. They are: half carry (H), negative (N), zero (Z), overflow (V), and carry (C). The effect each instruction has on these bits is given in the detail information for each instruction (see Appendix A).

1.10.1.1 Half Carry (H), Bit 5. This bit is used to indicate that a carry was generated from bit three in the arithmetic logic unit as a result of an 8-bit addition. This bit is undefined in all subtract-like instructions. The decimal addition adjust (DAA) instruction uses the state of this bit to perform the adjust operation.

1.10.1.2 Negative (N), Bit 3. This bit contains the value of the most-significant bit of the result of the previous data operation.

1.10.1.3 Zero (Z), Bit 2. This bit is used to indicate that the result of the previous operation was zero.

1.10.1.4 Overflow (V), Bit 1. This bit is used to indicate that the previous operation caused a signed arithmetic overflow.

1.10.1.5 Carry (C), Bit 0. This bit is used to indicate that a carry or a borrow was generated from bit seven in the arithmetic logic unit as a result of an 8-bit mathematical operation.

1.10.2 INTERRUPT MASK BITS AND STACKING INDICATOR. Two bits (I and F) are used as mask bits for the interrupt request and the fast interrupt request inputs. When either or both of these bits are set, their associated input will not be recognized.

One bit (E) is used to indicate how many registers (all, or only the program counter and condition code) were stacked during the last interrupt.

1.10.2.1 Fast Interrupt Request Mask (F), Bit 6. This bit is used to mask (disable) any fast interrupt request line ($\overline{\text{FIRQ}}$). This bit is set automatically by a hardware reset or after recognition of another interrupt. Execution of certain instructions such as SWI will also inhibit recognition of a $\overline{\text{FIRQ}}$ input.

1.10.2.2 Interrupt Request Mask (I), Bit 4. This bit is used to mask (disable) any interrupt request input ($\overline{\text{IRQ}}$). This bit is set automatically by a hardware reset or after recognition of another interrupt. Execution of certain instructions such as SWI will also inhibit recognition of an $\overline{\text{IRQ}}$ input.

1.10.2.3 Entire Flag (E), Bit 7. This bit is used to indicate how many registers were stacked. When set, all the registers were stacked during the last interrupt stacking operation. When clear, only the program counter and condition code registers were stacked during the last interrupt.

The state of the E bit in the stacked condition code register is used by the return from interrupt (RTI) instruction to determine the number of registers to be unstacked.

1.11 PIN ASSIGNMENTS AND SIGNAL DESCRIPTION

Figure 1-3 shows the pin assignments for the processors. The following paragraphs provide a short description of each of the input and output signals.

MC6809		MC6809E	
VSS	1	40	HALT
NMI	2	39	XTAL
IRQ	3	38	EXTAL
FIRQ	4	37	RESET
BS	5	36	MRDY
BA	6	35	Q
VCC	7	34	E
A0	8	33	DMA/BREQ
A1	9	32	R/W
A2	10	31	D0
A3	11	30	D1
A4	12	29	D2
A5	13	28	D3
A6	14	27	D4
A7	15	26	D5
A8	16	25	D6
A9	17	24	D7
A10	18	23	A15
A11	19	22	A14
A12	20	21	A13

Figure 1-3. Processor Pin Assignments

1.11.1 MC6809 CLOCKS. The MC6809 has four pins committed to developing the clock signals needed for internal and system operation. They are: the oscillator pins EXTAL and XTAL; the standard M6800 enable (E) clock; and a new, quadrature (Q) clock.

1.11.1.1 Oscillator (EXTAL, XTAL). These pins are used to connect the processor's internal oscillator to an external, parallel-resonant crystal. These pins can also be used for input of an external TTL timing signal by grounding the XTAL pin and applying the input to the EXTAL pin. The crystal or the external timing source is four times the resulting bus frequency.

1.11.1.2 Enable (E). The E clock is similar to the phase 2 (ϕ_2) MC6800 bus timing clock. The leading edge indicates to memory and peripherals that the data is stable and to begin write operations. Data movement occurs after the Q clock is high and is latched on the trailing edge of E. Data is valid from the processor (during a write operation) by the rising edge of E.

1.11.1.3 Quadrature (Q). The Q clock leads the E clock by approximately one half of the E clock time. Address information from the processor is valid with the leading edge of the Q clock. The Q clock is a new signal in these processors and does not have an equivalent clock within the MC6800 bus timing.

1.11.2 MC6809E CLOCKS (E and Q). The MC6809E has two pins provided for the TTL clock signal inputs required for internal operation. They are the standard M6800 enable (E) clock and the quadrature (Q) clock. The Q Input must lead the E input.

Addresses will be valid from the processor (on address delay time after the falling edge of E) and data will be latched from the bus by the falling edge of E. The Q input is fully TTL compatible. The E input is used to drive the internal MOS circuitry directly and therefore requires input levels above the normal TTL levels.

1.11.3 THREE STATE CONTROLS (TSC) (MC6809E). This input is used to place the address and data lines and the R/W line in the high-impedance state and allows the address bus to be shared with other bus masters.

1.11.4 LAST INSTRUCTION CYCLE (LIC) (MC6809E). This output goes high during the last cycle of every instruction and its high-to-low transition indicates that the first byte of an opcode will be latched at the end of the present bus cycle.

1.11.5 ADDRESS BUS (A0-A15). This 16-bit, unidirectional, three-state bus is used by the processor to provide address information to the address bus. Address information is valid on the rising edge of the Q clock. All 16 outputs are in the high-impedance state when the bus available (BA) signal is high, and for one bus cycle thereafter.

When the processor does not require the address bus for a data transfer, it outputs address FFFF₁₆, and read/write (R/W) high. This is a "dummy access" of the least-significant byte of the reset vector which replaces the valid memory address (VMA) functions of the MC6800. For the MC6809, the memory read signal internal circuitry inhibits stretching of the clocks during non-access cycles.

1.11.6 DATA BUS (D0-D7). This 8-bit, bidirectional, three-state bus is the general purpose data path. All eight outputs are in the high-impedance state when the bus available (BA) output is high.

1.11.7 READ/WRITE (R/W). This output indicates the direction of data transfer on the data bus. A low indicates that the processor is writing onto the data bus; a high indicates that the processor is reading data from the data bus. The signal at the R/W output is valid at the leading edge of the Q clock. The R/W output is in the high-impedance state when the bus available (BA) output is high.

1.11.8 PROCESSOR STATE INDICATORS (BA, BS). The processor uses these two output lines to indicate the present processor state. These pins are valid with the leading edge of the Q clock.

The bus available (BA) output is used to indicate that the buses (address and data) and the read/write output are in the high-impedance state. This signal can be used to indicate to bus-sharing or direct memory access systems that the buses are available. When BA goes low, an additional dead cycle will elapse before the processor regains control of the buses.

The bus status (BS) output is used in conjunction with the BA output to indicate the present state of the processor. Table 1-1 is a listing of the BA and BS outputs and the processor states that they indicate. The following paragraphs briefly explain each processor state.

Table 1-1. BA/BS Signal Encoding

<u>BA</u>	<u>BS</u>	<u>Processor State</u>
0	0	Normal (Running)
0	1	Interrupt or Reset Acknowledge
1	0	Sync Acknowledge
1	1	Halt/Bus Grant Acknowledged

1.11.8.1 Normal. The processor is running and executing instructions.

1.11.8.2 Interrupt or Reset Acknowledge. This processor state is indicated during both cycles of a hardware vector fetch which occurs when any of the following interrupts have occurred: RESET, NMI, FIRQ, IRQ, SWI, SWI2, and SWI3.

This output, plus decoding of address lines A3 through A1 provides the user with an indication of which interrupt is being serviced.

1.11.8.3 Sync Acknowledge. The processor is waiting for an external synchronization input on an interrupt line. See SYNC instruction in Appendix A.

1.11.8.4 Halt/Bus Grant. The processor is halted or bus control has been granted to some other device.

1.11.9 RESET (RESET). This input is used to reset the processor. A low input lasting longer than one bus cycle will reset the processor.

The reset vector is fetched from locations \$FFFE and \$FFFF when the processor enters the reset acknowledge state as indicated by the BA output being low and the BS output being high.

During initial power-on, the reset input should be held low until the clock oscillator is fully operational.

1.11.10 INTERRUPTS. The processor has three separate interrupt input pins: non-maskable interrupt (NMI), fast interrupt request (FIRQ), and interrupt request (IRQ). These interrupt inputs are latched by the falling edge of every Q clock except during cycle stealing operations where only the NMI input is latched. Using this point as a reference, a delay of at least one bus cycle will occur before the interrupt is recognized by the processor.

1.11.10.1 Non-Maskable Interrupt (NMI). A negative edge on this input requests that a non-maskable interrupt sequence be generated. This input, as the name indicates, cannot be masked by software and has the highest priority of the three interrupt inputs. After a reset has occurred, a NMI input will not be recognized by the processor until the first program load of the hardware stack pointer. The entire machine state is saved on the hardware stack during the processing of a non-maskable interrupt. This interrupt is internally blocked after a hardware reset until the stack pointer is initialized.

1.11.10.2 Fast Interrupt Request (FIRQ). This input is used to initiate a fast interrupt request sequence. Initiation depends on the F (fast interrupt request mask) bit in the condition code register being clear. This bit is set during reset. During the interrupt, only the contents of the condition code register and the program counter are stacked resulting in a short amount of time required to service this interrupt. This interrupt has a higher priority than the normal interrupt request (IRQ).

1.11.10.3 Interrupt Request (IRQ). This input is used to initiate what might be considered the "normal" interrupt request sequence. Initiation depends on the I (interrupt mask) bit in the condition code register being clear. This bit is set during reset. The entire machine state is saved on the hardware stack during processing of an IRQ input. This input has the lowest priority of the three hardware interrupts.

1.11.11 MEMORY READ (MRDY) (MC6809). This input allows extension of the E and Q clocks to allow a longer data access time. A low on this input allows extension of the E and Q clocks (E high and Q low) in integral multiples of quarter bus cycles (up to 10 cycles) to allow interface with slow memory devices.

Memory ready does not extend the E and Q clocks during non-valid memory access cycles and therefore the processor does not slow down for "don't care" bus accesses. Memory ready may also be used to extend the E and Q clocks when an external device is using the halt and direct memory access/bus request inputs.

1.11.12 ADVANCED VALID MEMORY ADDRESS (AVMA) (MC6809E). This output signal indicates that the MC6809E will use the bus in the following bus cycle. This output is low when the MC6809E is in either a halt or sync state.

1.11.13 HALT. This input is used to halt the processor. A low input halts the processor at the end of the present instruction execution cycle and the processor remains halted indefinitely without loss of data.

When the processor is halted, the BA output is high to indicate that the buses are in the high-impedance state and the BS output is also high to indicate that the processor is in the halt/bus grant state.

During the halt/bus grant state, the processor will not respond to external real-time requests such as FIRQ or IRQ. However, a direct memory access/bus request input will be accepted. A non-maskable interrupt or a reset input will be latched for processing later. The E and Q clocks continue to run during the halt/bus grant state.

1.11.14 DIRECT MEMORY ACCESS/BUS REQUEST (DMA/BREQ) (MC6809). This input is used to suspend program execution and make the buses available for another use such as a direct memory access or a dynamic memory refresh.

A low level on this input occurring during the Q clock high time suspends instruction execution at the end of the current cycle. The processor acknowledges acceptance of this input by setting the BA and BS outputs high to signify the bus grant state. The requesting device now has up to 15 bus cycles before the processor retrieves the bus for self-refresh.

Typically, a direct memory access controller will request to use the bus by setting the DMA/BREQ input low when E goes high. When the processor acknowledges this input by setting the BA and BS outputs high, that cycle will be a dead cycle used to transfer bus mastership to the direct memory access controller. False memory access during any dead cycle should be prevented by externally developing a system DMAVMA signal which is low in any cycle when the BA output changes.

When the BA output goes low, either as a result of a direct memory access/bus request or a processor self-refresh, the direct memory access device should be removed from the bus. Another dead cycle will elapse before the processor accesses memory, to allow transfer of bus mastership without contention.

1.11.15 BUSY (MC6809E). This output indicates that bus re-arbitration should be deferred and provides the indivisible memory operation required for a "test-and-set" primitive.

This output will be high for the first two cycles of any Read-Modify-Write instruction, high during the first byte of a double-byte access, and high during the first byte of any indirect access or vector-fetch operation.

1.11.16 POWER. Two inputs are used to supply power to the processor: V_{CC} is +5.0 ± 5%, while V_{SS} is ground or 0 volts.

SECTION 2 ADDRESSING MODES

2.1 INTRODUCTION

This section contains a description of each of the addressing modes available on these processors.

2.2 ADDRESSING MODES

The addressing modes available on the MC6809 and MC6809E are: Inherent, Immediate, Extended, Direct, Indexed (with various offsets and autoincrementing/decrementing), and Branch Relative. Some of these addressing modes require an additional byte after the opcode to provide additional addressing interpretation. This byte is called a postbyte.

The following paragraphs provide a description of each addressing mode. In these descriptions the term effective address is used to indicate the address in memory from which the argument for an instruction is fetched or stored, or from which instruction processing is to proceed.

2.2.1 INHERENT. The information necessary to execute the instruction is contained in the opcode. Some operations specifying only the index registers or the accumulators, and no other arguments, are also included in this addressing mode.

Example: MUL

2.2.2 IMMEDIATE. The operand is contained in one or two bytes immediately following the opcode. This addressing mode is used to provide constant data values that do not change during program execution. Both 8-bit and 16-bit operands are used depending on the size of the argument specified in the opcode.

Example: LDA #CR
 LDB #7
 LDA #\$F0
 LDB #%1110000
 LDX #\$8004

Another form of immediate addressing uses a postbyte to determine the registers to be manipulated. The exchange (EXG) and transfer (TFR) instructions use the postbyte as shown in Figure 2-1(A). The push and pull instructions use the postbyte to designate the registers to be pushed or pulled as shown in Figure 2-1(B).

b7	b6	b5	b4	b3	b2	b1	b0
SOURCE (R1)				DESTINATION (R2)			
Code*	Register	Code*	Register				
0000	D (A:B)	0101	Program Counter				
0001	X Index	1000	A Accumulator				
0010	Y Index	1001	B Accumulator				
0011	U Stack Pointer	1010	Condition Code				
0100	S Stack Pointer	1011	Direct Page				

*All other combinations of bits produce undefined results.

(A) Exchange (EXG) or Transfer (TFR) Instruction Postbyte

b7	b6	b5	b4	b3	b2	b1	b0
PC	S/U	Y	X	DP	B	A	CC

PC = Program Counter
 S/U = Hardware/User Stack Pointer
 Y = Y Index Register
 X = U Index Register
 DP = Direct Page Register
 B = B Accumulator
 A = A Accumulator
 CC = Condition Code Register

(B) Push (PSH) or Pull (PUL) Instruction Postbyte

Figure 2-1. Postbyte Usage for EXG/TFR, PSH/PUL Instructions

2.2.3 EXTENDED. The effective address of the argument is contained in the two bytes following the opcode. Instructions using the extended addressing mode can reference arguments anywhere in the 64K addressing space. Extended addressing is generally not used in position independent programs because it supplies an absolute address.

Example: LDA < CAT

2.2.4 DIRECT. The effective address is developed by concatenation of the contents of the direct page register with the byte immediately following the opcode. The direct page register contents are the most-significant byte of the address. This allows accessing 256 locations within any one of 256 pages. Therefore, the entire addressing range is available for access using a single two-byte instruction.

Example: LDA > CAT

2.2.5 INDEXED. In these addressing modes, one of the pointer registers (X, Y, U, or S), and sometimes the program counter (PC) is used in the calculation of the effective address of the instruction operand. The basic types (and their variations) of indexed addressing available are shown in Table 2-1 along with the postbyte configuration used.

2.2.5.1 Constant Offset from Register. The contents of the register designated in the postbyte are added to a two's complement offset value to form the effective address of

the instruction operand. The contents of the designated register are not affected by this addition. The offset sizes available are:

- No offset — designated register contains the effective address
- 5-bit — 16 to + 15
- 8-bit — 128 to + 127
- 16-bit — 32768 to + 32767

Table 2-1. Postbyte Usage for Indexed Addressing Modes

Mode Type	Variation	Direct	Indirect
Constant Offset from Register (twos Complement Offset)	No Offset 5-Bit Offset 8-Bit Offset 16-Bit Offset	1RR00100 0RRnnnnn 1RR01000 1RR01001	1RR10100 Defaults to 8-bit 1RR11000 1RR11001
Accumulator Offset from Register (twos Complement Offset)	A Accumulator Offset B Accumulator Offset D Accumulator Offset	1RR00110 1RR00101 1RR01011	1RR10110 1RR10101 1RR11011
Auto Increment/Decrement from Register	Increment by 1 Increment by 2 Decrement by 1 Decrement by 2	1RR00000 1RR00001 1RR00010 1RR00011	Not Allowed 1RR10001 Not Allowed 1RR10011
Constant Offset from Program Counter	8-Bit Offset 16-Bit Offset	1XX01100 1XX01101	1XX11100 1XX11101
Extended Indirect	16-Bit Address	-----	10011111

The 5-bit offset value is contained in the postbyte. The 8- and 16-bit offset values are contained in the byte or bytes immediately following the postbyte. If the Motorola assembler is used, it will automatically determine the most efficient offset; thus, the programmer need not be concerned about the offset size.

Examples: LDA ,X LDY - 64000,U
 LDB 0,Y LDA 17,PC
 LDX 64,000,S LDA There,PCR

2.2.5.2 Accumulator Offset from Register. The contents of the index or pointer register designed in the postbyte are temporarily added to the twos complement offset value contained in an accumulator (A, B, or D) also designated in the postbyte. Neither the designated register nor the accumulator contents are affected by this addition.

Example: LDA A,X LDA D,U
 LDA B,Y

2.2.5.3 Autoincrement/Decrement from Register. This addressing mode works in a postincrementing or predecrementing manner. The amount of increment or decrement, one or two positions, is designated in the postbyte.

In the autoincrement mode, the contents of the effective address contained in the pointer register, designated in the postbyte, and then the pointer register is automatically incremented; thus, the pointer register is postincremented.

In the autodecrement mode, the pointer register, designated in the postbyte, is automatically decremented first and then the contents of the new address are used; thus, the pointer register is predecremented.

Examples:

	Autoincrement	Autodecrement
LDA ,X+	LDY ,X+ +	LDA , - X LDY , - - X
LDA ,Y+	LDX ,Y+ +	LDA , - Y LDX , - - Y
LDA ,S+	LDX ,U+ +	LDA , - S LDX , - - U
LDA ,U+	LDX ,S+ +	LDA , - U LDX , - - S

2.2.5.4 Indirection. When using indirection, the effective address of the base indexed addressing mode is used to fetch two bytes which contain the final effective address of the operand. It can be used with all the indexed addressing modes and the program counter relative addressing mode.

2.2.5.5 Extended Indirect. The effective address of the argument is located at the address specified by the two bytes following the postbyte. The postbyte is used to indicate indirection.

Example: LDA [\$F000]

2.2.5.6 Program Counter Relative. The program counter can also be used as a pointer with either an 8- or 16-bit signed constant offset. The offset value is added to the program counter to develop an effective address. Part of the postbyte is used to indicate whether the offset is 8 or 16 bits.

2.2.6 BRANCH RELATIVE. This addressing mode is used when branches from the current instruction location to some other location relative to the current program counter are desired. If the test condition of the branch instruction is true, then the effective address is calculated (program counter plus two's complement offset) and the branch is taken. If the test condition is false, the processor proceeds to the next in-line instruction. Note that the program counter is always pointing to the next instruction when the offset is added. Branch relative addressing is always used in position independent programs for all control transfers.

For short branches, the byte following the branch instruction opcode is treated as an 8-bit signed offset to be used to calculate the effective address of the next instruction if the branch is taken. This is called a short relative branch and the range is limited to plus 127 or minus 128 bytes from the following opcode.

For long branches, the two bytes after the opcode are used to calculate the effective address. This is called a long relative branch and the range is plus 32,767 or minus 32,768

SECTION 3 INTERRUPT CAPABILITIES

3.1 INTRODUCTION

The MC6809 and MC6809E microprocessors have six vectored interrupts (three hardware and three software). The hardware interrupts are the non-maskable interrupt (NMI), the fast maskable interrupt request (FIRQ), and the normal maskable interrupt request (IRQ). The software interrupts consist of SWI, SWI2, and SWI3. When an interrupt request is acknowledged, all the processor registers are pushed onto the hardware stack, except in the case of FIRQ where only the program counter and the condition code register is saved, and control is transferred to the address in the interrupt vector. The priority of these interrupts is, highest to lowest, NMI, SWI, FIRQ, IRQ, SWI2, and SWI3. Figure 3-1 is a detailed flowchart of interrupt processing in these processors. The interrupt vector locations are given in Table 3-1. The vector locations contain the address for the interrupt routine.

Additional information on the SWI, SWI2, and SWI3 interrupts is given in Appendix A. The hardware interrupts, NMI, FIRQ, and IRQ are listed alphabetically at the end of Appendix A.

Table 3-1. Interrupt Vector Locations

Interrupt Description	Vector Location	
	MS Byte	LS Byte
Reset (RESET)	FFFE	FFFF
Non-Maskable Interrupt (NMI)	FFFC	FFFD
Software Interrupt (SWI)	FFFA	FFFB
Interrupt Request (IRQ)	FFF8	FFF9
Fast Interrupt Request (FIRQ)	FFF6	FFF7
Software Interrupt 2 (SWI2)	FFF4	FFF5
Software Interrupt 3 (SWI3)	FFF2	FFF3
Reserved	FFF0	FFF1

3.2 NON-MASKABLE INTERRUPT (NMI)

The non-maskable interrupt is edge-sensitive in the sense that if it is sampled low one cycle after it has been sampled high, a non-maskable interrupt will be triggered. Because the non-maskable interrupt cannot be masked by execution of the non-maskable interrupt handler routine, it is possible to accept another non-maskable interrupt before executing the first instruction of the interrupt routine. A fatal error will exist if a non-maskable interrupt is repeatedly allowed to occur before completing the return from interrupt (RTI) instruction of the previous non-maskable interrupt request, since the stack

will eventually overflow. This interrupt is especially applicable to gaining immediate processor response for powerfall, software dynamic memory refresh, or other non-delayable events.

3.3 FAST MASKABLE INTERRUPT REQUEST (FIRQ)

A low level on the FIRQ input with the F (fast interrupt request mask) bit in the condition code register clear triggers this interrupt sequence. The fast interrupt request provides fast interrupt response by stacking only the program counter and condition code register. This allows fast context switching with minimal overhead. If any registers are used by the interrupt routine then they can be saved by a single push instruction.

After accepting a fast interrupt request, the processor clears the E flag, saves the program counter and condition code register, and then sets both the I and F bits to mask any further IRQ and FIRQ interrupts. After servicing the original interrupt, the user may selectively clear the I and F bits to allow multiple-level interrupts if so desired.

3.4 NORMAL MASKABLE INTERRUPT REQUEST (IRQ)

A low level on the IRQ input with the I (interrupt request mask) bit in the condition code register clear triggers this interrupt sequence. The normal maskable interrupt request provides a slower hardware response to interrupts because it causes the entire machine state to be stacked. However, this means that interrupting software routines can use all processor resources without fear of damaging the interrupted routine. A normal interrupt request, having lower priority than the fast interrupt request, is prevented from interrupting the fast interrupt handler by the automatic setting of the I bit by the fast interrupt request handler.

After accepting a normal interrupt request, the processor sets the E flag, saves the entire machine state, and then sets the I bit to mask any further interrupt request inputs. After servicing the original interrupt, the user may clear the I bit to allow multiple-level normal interrupts.

All interrupt handling routines should return to the formerly executing tasks using a return from interrupt (RTI) instruction. This instruction recovers the saved machine state from the hardware stack and control is returned to the interrupted program. If the recovered E bit is clear, it indicates that a fast interrupt request occurred and only the program counter address and condition code register are to be recovered.

3.5 SOFTWARE INTERRUPTS (SWI, SWI2, SWI3)

The software interrupts cause the processor to go through the normal interrupt request sequence of stacking the complete machine state even though the interrupting source is the processor itself. These interrupts are commonly used for program debugging and for calls to an operating system.

Normal processing of the SWI input sets the I and F bits to prevent either of these interrupt requests from affecting the completion of a software interrupt request. The remaining software interrupt request inputs (SWI2 and SWI3) do not have the priority of the SWI input and therefore do not mask the two hardware interrupt request inputs (FIRQ and IRQ).

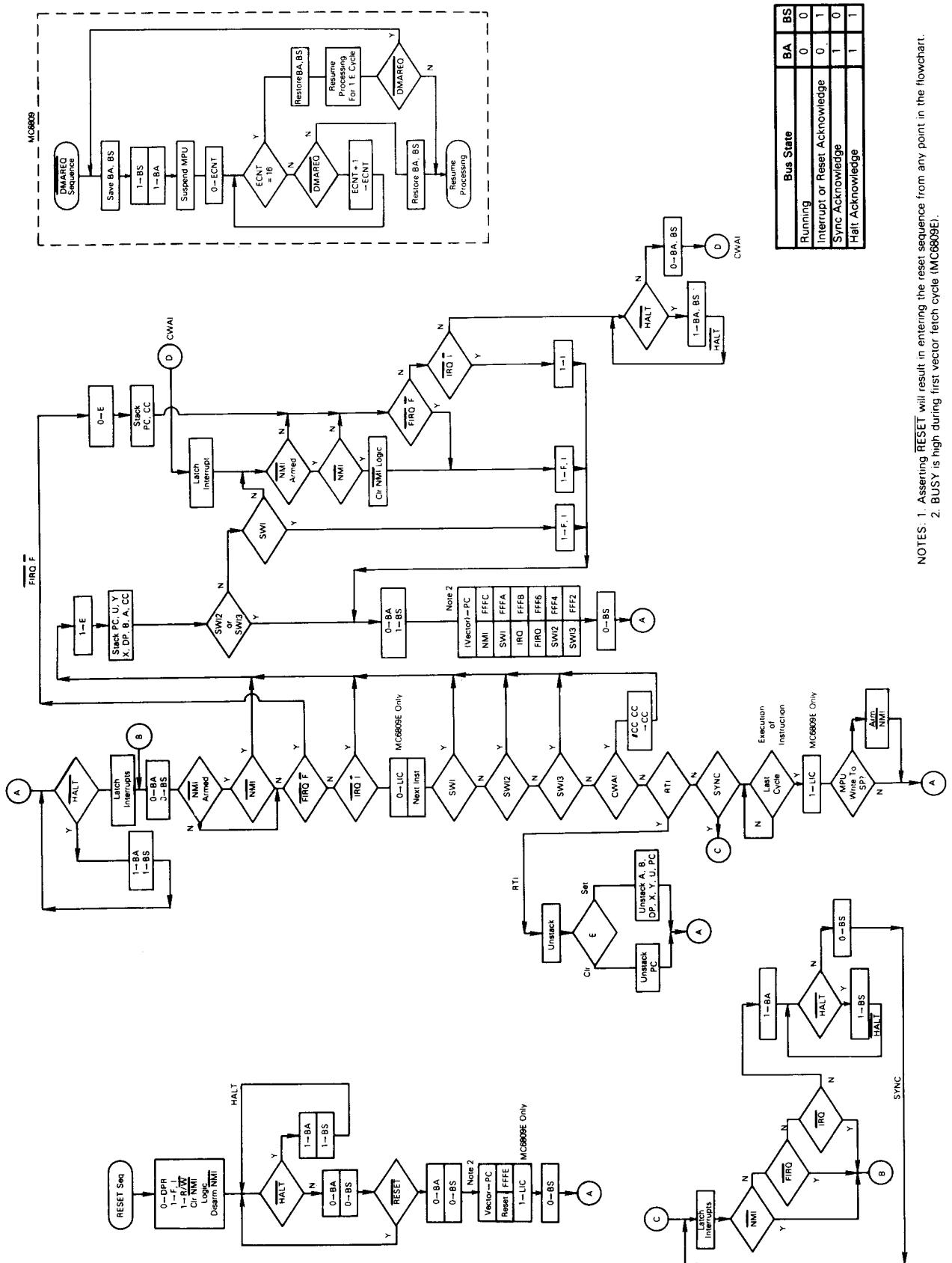


Figure 3-1. Interrupt Processing Flowchart

SECTION 4 PROGRAMMING

4.1 INTRODUCTION

These processors are designed to be source-code compatible with the M6800 to make use of the substantial existing base of M6800 software and training. However, this asset should not overshadow the capabilities built into these processors that allow more modern programming techniques such as position-independence, modular programming, and reentrancy/recursion to be used on a microprocessor-based system. A brief review of these methods is given in the following paragraphs.

4.1.1 POSITION INDEPENDENCE. A program is said to be "position-independent" if it will run correctly when the same machine code is positioned arbitrarily in memory. Such a program is useful in many different hardware configurations, and might be copied from a disk into RAM when the operating system first sees a request to use a system utility. Position-independent programs never use absolute (extended or direct) addressing; instead, inherent immediate, register, indexed and relative modes are used. In particular, there should be no jump (absolute) or jump to subroutine instructions nor should absolute addresses be used. A position-independent program is almost always preferable to a position-dependent program (although position-independent code is usually 5 to 10% slower than normal code).

4.1.2 MODULAR PROGRAMMING. Modular programming is another indication of quality code. A module is a program element which can be easily disconnected from the rest of the program either for re-use in a new environment or for replacement. A module is usually a subroutine (although a subroutine is not necessarily a module); frequently, the programmer isolates register changes internal to the module by pushing these registers onto the stack upon entry, and pulling them off the stack before the return. Isolating register changes in the called module, to that module alone, allows the code in the calling program to be more easily analyzed since it can be assumed that all registers (except those specifically used for parameter transfer) are unchanged by each called module. This leaves the processor's registers free at each level for loop counts, address comparisons, etc.

4.1.2.1 Local Storage. A clean method for allocating "local" storage is required both by position-independent programs as well as modular programs. Local or temporary storage is used to hold values only during execution of a module (or called modules) and is released upon return. One way to allocate local storage is to decrement the hardware stack

pointer(s) by the number of bytes needed. Interrupts will then leave this area intact and it can be de-allocated on exiting the module. A module will almost always need more temporary storage than just the MPU registers.

4.1.2.2 Global Storage. Even in a modular environment there may be a need for "global" values which are accessible by many modules within a given system. These provide a convenient means for storing values from one invocation to another invocation of the same routine. Global storage may be created as local storage at some level, and a pointer register (usually U) used to point at this area. This register is passed unchanged in all subroutines, and may be used to index into the global area.

4.1.3 REENTRANCY/RECURSION. Many programs will eventually involve execution in an interrupt-driven environment. If the interrupt handlers are complex, they might well call the same routine which has just been interrupted. Therefore, to protect present programs against certain obsolescence, all programs should be written to be reentrant. A reentrant routine allocates different local variable storage upon each entry. Thus, a later entry does not destroy the processing associated with an earlier entry.

The same technique which was implemented to allow reentrancy also allows recursion. A recursive routine is defined as a routine that calls itself. A recursive routine might be written to simplify the solution of certain types of problems, especially those which have a data structure whose elements may themselves be a structure. For example, a parenthetical equation represents a case where the expression in parenthesis may be considered to be a value which is operated on by the rest of the equation. A programmer might choose to write an expression evaluator passing the parenthetical expression (which might also contain parenthetical expressions) in the call, and receive back the returned value of the expression within the parenthesis.

4.2 M6809 CAPABILITIES

The following paragraphs briefly explain how the MC6809 is used with the programming techniques mentioned earlier.

4.2.1 MODULE CONSTRUCTION. A module can be defined as a logically self-contained and discrete part of a larger program. A properly constructed module accepts well defined inputs, carries out a set of processing actions, and produces a specified output. The use of parameters, local storage, and global storage by a program module is given in the following paragraphs. Since registers will be used inside the module (essentially a form of local storage), the first thing that is usually done at entry to a module is to push (save) them on to the stack. This can be done with one instruction (e.g., PSHS Y, X, B, A). After the body of the module is executed, the saved registers are collected, and a subroutine return is performed, at one time, by pulling the program counter from the stack (e.g., PULS A,B,X,Y,PC).

4.2.1.1 Parameters. Parameters may be passed to or from modules either in registers, if they will provide sufficient storage for parameter passage, or on the stack. If parameters are passed on the stack, they are placed there before calling the lower level module. The called module is then written to use local storage inside the stack as needed (e.g., ADDA offset,S). Notice that the required offset consists of the number of bytes pushed (upon entry), plus two from the stacked return address, plus the data offset at the time of the call. This value may be calculated, by hand, by drawing a "stack picture" diagram representing module entry, and assigning convenient mnemonics to these offsets with the assembler. Returned parameters replace those sent to the routine. If more parameters are to be returned on the stack than would normally be sent, space for their return is allocated by the calling routine before the actual call (if four additional bytes are to be returned, the caller would execute LEAS - 4,S to acquire the additional storage).

4.2.1.2 Local Storage. Local storage space is acquired from the stack while the present routine is executing and then returned to the stack prior to exit. The act of pushing registers which will be used in later calculations essentially saves those registers in temporary local storage. Additional local storage can easily be acquired from the stack e.g., executing LEAS - 2048,S acquires a buffer area running from the 0,S to 2047,S inclusive. Any byte in this area may be accessed directly by any instruction which has an indexed addressing mode. At the end of the routine, the area acquired for local storage is released (e.g., LEAS 2048,S) prior to the final pull. For cleaner programs, local storage should be allocated at entry to the module and released at the exit of the module.

4.2.1.3 Global Storage. The area required for global storage is also most effectively acquired from the stack, probably by the highest level routine in the standard package. Although this is local storage to the highest level routine, it becomes "global" by positioning a register to point at this storage, (sometimes referred to as a stack mark) then establishing the convention that all modules pass that same pointer value when calling lower level modules. In practice, it is convenient to leave this stack mark register unchanged in all modules, especially if global accesses are common. The highest level routine in the standard package would execute the following sequence upon entry (to initialize the global area):

PSHS	U	higher level mark, if any
TFR	S,U	new stack mark
LEAS	- 17,U	allocate global storage

Note that the U register now defines 17-bytes of locally allocated (permanent) globals (which are - 1,U through - 17,U) as well as other external globals (2,U and above) which have been passed on the stack by the routine which called the standard package. Any global may be accessed by any module using exactly the same offset value at any level (e.g., ROL, RAT,U; where RAT EQU - 11 has been defined). Furthermore, the values stacked prior to invoking the standard package may include pointers to data or I/O peripherals. Any indexed operation may be performed indexed indirect through those pointers, which means, for example, that the module need know nothing about the actual hardware configuration, except that (upon entry) the pointer to an I/O register has been placed at a given location on the stack.

4.2.2 POSITION-INDEPENDENT CODE. Position-independent code means that the same machine language code can be placed anywhere in memory and still function correctly. The M6809 has a long relative (16-bit offset) branch mode along with the common MC6800 branches, plus program-counter relative addressing. Program-counter relative addressing uses the program counter like an indexable register, which allows all instructions that reference memory to also reference data relative to the program counter. The M6809 also has load effective address (LEA) instructions which allow the user to point to data in a ROM in a position-independent manner.

An important rule for generating position-independent code is: NEVER USE ABSOLUTE ADDRESSING.

Program-counter relative addressing on the M6809 is a form of indexed addressing that uses the program counter as the base register for a constant-offset indexing operation. However, the M6809 assembler treats the PCR address field differently from that used in other indexed instructions. In PCR addressing, the assembly time location value is subtracted from the (constant) value of the PCR offset. The resulting distance to the desired symbol is the value placed into the machine language object code. During execution, the processor adds the value of the run time PC to the distance to get a position-independent absolute address.

The PCR indexed addressing form can be used to point at any location relative to the program regardless of position in memory. The PCR form of indexed addressing allows access to tables within the program space in a position-independent manner via use of the load effective address instruction.

In a program which is completely position-independent, some absolute locations are usually required, particularly for I/O. If the locations of I/O devices are placed on the stack (as globals) by a small setup routine before the standard package is invoked, all internal modules can do their I/O through that pointer (e.g., STA [ACIAD, U]), allowing the hardware to be easily changed, if desired. Only the single, small, and obvious setup routine need be rewritten for each different hardware configuration.

Global, permanent, and temporary values need to be easily available in a position-independent manner. Use the stack for this data since the stacked data is directly accessible. Stack the absolute address of I/O devices before calling any standard software package since the package can use the stacked addresses for I/O in any system.

The LEA Instructions allow access to tables, data, or immediate values in the text of the program in a position-independent manner as shown in the following example:

MSG1 FCC /PRINT THIS!/

Here we wish to point at a message to be printed from the body of the program. By writing "MSG1, PCR" we signal the assembler to compute the distance between the present address (the address of the LBSR) and MSG1. This result is inserted as a constant into the LEA instruction which will be indexed from the program counter value at the time of execution. Now, no matter where the code is located, when it is executed the computer offset from the program counter will point at MSG1. This code is position-independent.

It is common to use space in the hardware stack for temporary storage. Space is made for temporary variables from 0,S through TEMP-1, S by decrementing the stack pointer equal to the length of required storage. We could use:

LEAS - TEMP,S.

Not only does this facilitate position-independent code but it is structured and helps reentrancy and recursion.

4.2.3 REENTRANT PROGRAMS. A program that can be executed by several different users sharing the same copy of it in memory is called reentrant. This is important for interrupt driven systems. This method saves considerable memory space, especially with large interrupt routines. Stacks are required for reentrant programs, and the M6809 can support up to four stacks by using the X and Y index registers as stack pointers.

Stacks are simple and convenient mechanisms for generating reentrant programs. Subroutines which use stacks for passing parameters and results can be easily made to be reentrant. Stack accesses use the indexed addressing mode for fast, efficient execution. Stack addressing is quick.

Pure code, or code that is not self-modifying, is mandatory to produce reentrant code. No internal information within the code is subject to modification. Reentrant code never has internal temporary storage, is simpler to debug, can be placed in ROM, and must be interruptable.

4.2.4 RECURSIVE PROGRAMS. A recursive program is one that can call itself. They are quite convenient for parsing mechanisms and certain arithmetic functions such as computing factorials. As with reentrant programming, stacks are very useful for this technique.

4.2.5 LOOPS. The usual structured loops (i.e., REPEAT...UNTIL, WHILE...DO, FOR..., etc.) are available in assembly language in exactly the same way a high-level language compiler could translate the construct for execution on the target machine. Using a FOR...NEXT loop as an example, it is possible to push the loop count, increment value, and termination value on the stack as variables local to that loop. On each pass through the loop, the working register is saved, the loop count picked up, the increment added in, and the result compared to the termination value. Based on this comparison, the loop counter might be updated, the working register recovered and the loop resumed, or the working register recovered and the loop variables de-allocated. Reasonable macros

could make the source form for loop trivial, even in assembly language. Such macros might reduce errors resulting from the use of multiple instructions simply to implement a standard control structure.

4.2.6 STACK PROGRAMMING. Many microprocessor applications require data stored as contiguous pieces of information in memory. The data may be temporary, that is, subject to change or it may be permanent. Temporary data will most likely be stored in RAM. Permanent data will most likely be stored in ROM.

It is important to allow the main program as well as subroutines access to this block of data, especially if arguments are to be passed from the main program to the subroutines and vice versa.

4.2.6.1 M6809 Stacking Operations. Stack pointers are markers which point to the stack and its internal contents. Although all four index registers may be used as stack registers, the S (hardware stack pointer) and the U (user stack pointer) are generally preferred because the push and pull instructions apply to these registers. Both are 16-bit indexable registers. The processor uses the S register automatically during interrupts and subroutine calls. The U register is free for any purpose needed. It is not affected by interrupts or subroutine calls implemented by the hardware.

Either stack pointer can be specified as the base address in indexed addressing. One use of the indirect addressing mode uses stack pointers to allow addresses of data to be passed to a subroutine on a stack as arguments to a subroutine. The subroutine can now reference the data with one instruction. High-level language calls that pass arguments by reference are now more efficiently coded. Also, each stack push or pull operation in a program uses a postbyte which specifies any register or set of registers to be pushed or pulled from either stack. With this option, the overhead associated with subroutine calls in both assembly and high-level language programs is greatly decreased. In fact, with the large number of instructions that use autoincrement and autodecrement, the M6809 can emulate a true stack computer architecture.

Using the S or U stack pointer, the order in which the registers are pushed or pulled is shown in Figure 4-1. Notice that we push "onto" the stack towards decreasing memory locations. The program counter is pushed first. Then the stack pointer is decremented and the "other" stack pointer is pushed onto the stack. Decrementing and storing continues until all the registers requested by the postbyte are pushed onto the stack. The stack pointer points to the top of the stack after the push operation.

The stacking order is specified by the processor. The stacking order is identical to the order used for all hardware and software interrupts. The same order is used even if a subset of the registers is pushed.

Without stacks, most modern block-structured high-level languages would be cumbersome to implement. Subroutine linkage is very important in high-level language generation. Paragraph 4.2.6.2 describes how to use a stack mark pointer for this important task.

Good programming practice dictates the use of the hardware stack for temporary storage. To reserve space, decrement the stack pointer by the amount of storage required with the instruction LEAS – TEMPS, S. This instruction makes space for temporary variables from 0,S through TEMPS – 1,S.

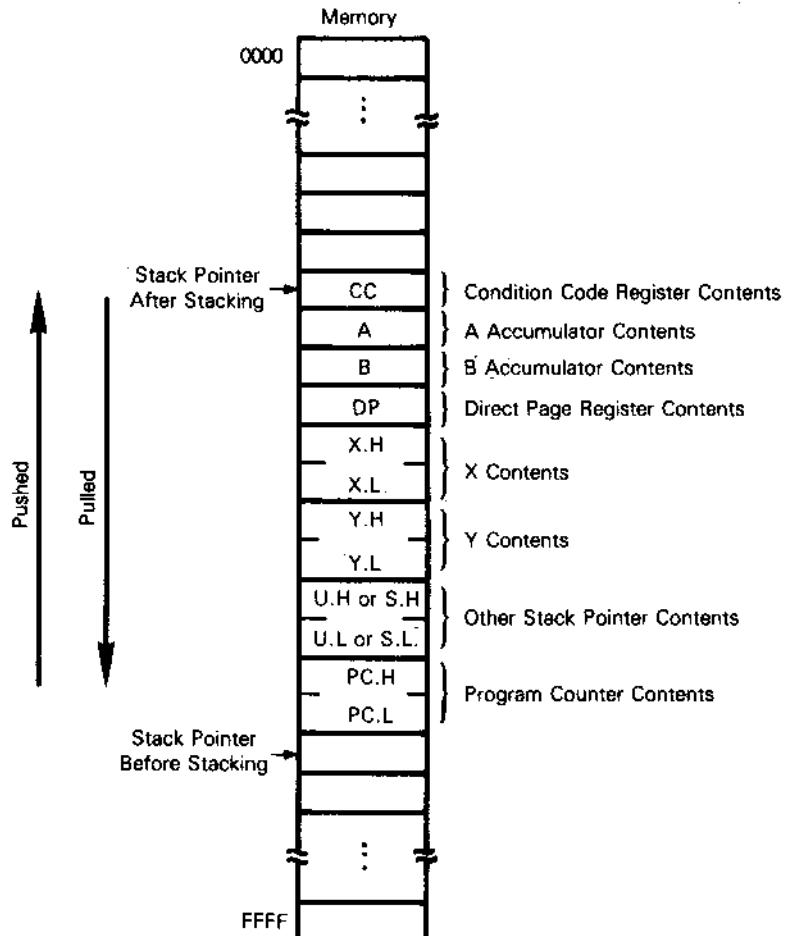


Figure 4-1. Stacking Order

4.2.6.2 Subroutine Linkage. In the highest level routine, global variables are sometimes considered to be local. Therefore, global storage is allocated at this point, but access to these same variables requires different offset values depending on subroutine depth. Because subroutine depth changes dynamically, the length may not be known beforehand. This problem is solved by assigning one pointer (U will be used in the following description, but X or Y could also be used) to "mark" a location on the hardware stack by using the instruction TFR S,U. If the programmer does this immediately prior to allocating global storage, then all variables will then be available at a constant negative offset location from this stack mark. If the stack is marked after the global variables are

allocated, then the global variables are available at a constant positive offset from U. Register U is then called the stack mark pointer. Recall that the hardware stack pointer may be modified by hardware interrupts. For this reason, it is fatal to use data referred to by a negative offset with respect to the hardware stack pointer, S.

4.2.6.3 Software Stacks. If more than two stacks are needed, autoincrement and autodecrement mode of addressing can be used to generate additional software stack pointers.

The X, Y, and U Index registers are quite useful in loops for incrementing and decrementing purposes. The pointer is used for searching tables and also to move data from one area of memory to another (block moves). This autoincrement and autodecrement feature is available in the Indexed addressing mode of the M6809 to facilitate such operations.

In autoincrement, the value contained in the index register (X or Y, U or S) is used as the effective address and then the register is incremented (postincremented). In autodecrement, the index register is first decremented and then used to obtain the effective address (predecremented). Postincrement or predecrement is always performed in this addressing mode. This is equivalent in operation to the push and pull from a stack. This equivalence allows the X and Y index registers to be used as software stack pointers. The indexed addressing mode can also implement an extra level of post indirection. This feature supports parameter and pointer operations.

4.2.7 REAL TIME PROGRAMMING. Real time programming requires special care. Sometimes a peripheral or task demands an immediate response from the processor, other times it can wait. Most real time applications are demanding in terms of processor response.

A common solution is to use the interrupt capability of the processor in solving real time problems. Interrupts mean just that; they request a break in the current sequence of events to solve an asynchronous service request. The system designer must consider all variations of the conditions to be encountered by the system including software interaction with interrupts. As a result, problems due to software design are more common in interrupt implementation code for real time programming than most other situations. Software timeouts, hardware interrupts, and program control interrupts are typically used in solving real time programming problems.

4.3 PROGRAM DOCUMENTATION

Common sense dictates that a well documented program is mandatory. Comments are needed to explain each group of instructions since their use is not always obvious from looking at the code. Program boundaries and branch instructions need full clarification. Consider the following points when writing comments: up-to-date, accuracy, completeness, conciseness, and understandability.

Accurate documentation enables you and others to maintain and adapt programs for updating and/or additional use with other programs.

The following program documentation standards are suggested.

- A) Each subroutine should have an associated header block containing at least the following elements:
 - 1) A full specification for this subroutine — including associated data structures — such that replacement code could be generated from this description alone.
 - 2) All usage of memory resources must be defined, including:
 - a) All RAM needed from temporary (local) storage used during execution of this subroutine or called subroutines.
 - b) All RAM needed for permanent storage (used to transfer values from one execution of the subroutine to future executions).
 - c) All RAM accessed as global storage (used to transfer values from or to higher-level subroutines).
 - d) All possible exit-state conditions, if these are to be used by calling routines to test occurrences internal to the subroutine.
- B) Code internal to each subroutine should have sufficient associated line comments to help in understanding the code.
- C) All code must be non-self-modifying and position-independent.
- D) Each subroutine which includes a loop must be separately documented by a flowchart or pseudo high-level language algorithm.
- E) Any module or subroutine should be executable starting at the first location and exit at the last location.

4.4 INSTRUCTION SET

The complete instruction set for the M6809 is given in Table 4-1.

Table 4-1. Instruction Set

Instruction	Description
ABX	Add Accumulator B into Index Register X
ADC	Add with Carry into Register
ADD	Add Memory into Register
AND	Logical AND Memory into Register
ASL	Arithmetic Shift Left
ASR	Arithmetic Shift Right
BCC	Branch on Carry Clear
BCS	Branch on Carry Set
BEQ	Branch on Equal
BGE	Branch on Greater Than or Equal to Zero
BGT	Branch on Greater
BHI	Branch if Higher
BHS	Branch if Higher or Same
BIT	Bit Test
BLE	Branch if Less than or Equal to Zero

Table 4-1. Instruction Set (Continued)

Instruction	Description
BLO	Branch on Lower
BLS	Branch on Lower or Same
BLT	Branch on Less than Zero
BMI	Branch on Minus
BNE	Branch Not Equal
BPL	Branch on Plus
BRA	Branch Always
BRN	Branch Never
BSR	Branch to Subroutine
BVC	Branch on Overflow Clear
BVS	Branch on Overflow Set
CLR	Clear
CMP	Compare Memory from a Register
COM	Complement
CWAI	Clear CC bits and Wait for Interrupt
DAA	Decimal Addition Adjust
DEC	Decrement
EOR	Exclusive OR
EXG	Exchange Registers
INC	Increment
JMP	Jump
JSR	Jump to Subroutine
LD	Load Register from Memory
LEA	Load Effective Address
LSL	Logical Shift Left
LSR	Logical Shift Right
MUL	Multiply
NEG	Negate
NOP	No Operation
OR	Inclusive OR Memory into Register
PSH	Push Registers
PUL	Pull Registers
ROL	Rotate Left
ROR	Rotate Right
RTI	Return from Interrupt
RTS	Return from Subroutine
SBC	Subtract with Borrow
SEX	Sign Extend
ST	Store Register into Memory
SUB	Subtract Memory from Register
SWI	Software Interrupt
SYNC	Synchronize to External Event
TFR	Transfer Register to Register
TST	Test

The instruction set can be functionally divided into five categories. They are:

8-Bit Accumulator and Memory Instructions

16-Bit Accumulator and Memory Instructions

Index Register/Stack Pointer Instructions

Branch Instructions

Miscellaneous Instructions

Tables 4-2 through 4-6 are listings of the M6809 instructions and their variations grouped into the five categories listed.

Table 4-2. 8-Bit Accumulator and Memory Instructions

Instruction	Description
ADCA, ADCB	Add memory to accumulator with carry
ADDA, ADDB	Add memory to accumulator
ANDA, ANDB	And memory with accumulator
ASL, ASLA, ASLB	Arithmetic shift of accumulator or memory left
ASR, ASRA, ASRB	Arithmetic shift of accumulator or memory right
BITA, BITB	Bit test memory with accumulator
CLR, CLRA, CLRB	Clear accumulator or memory location
CMPA, CMPB	Compare memory from accumulator
COM, COMA, COMB	Complement accumulator or memory location
DAA	Decimal adjust A accumulator
DEC, DECA, DECB	Decrement accumulator or memory location
EORA, EORB	Exclusive or memory with accumulator
EXG R1, R2	Exchange R1 with R2 (R1, R2=A, B, CC, DP)
INC, INCA, INCB	Increment accumulator or memory location
LDA, LDB	Load accumulator from memory
LSL, LSLA, LSLB	Logical shift left accumulator or memory location
LSR, LSRA, LSRB	Logical shift right accumulator or memory location
MUL	Unsigned multiply (A × B → D)
NEG, NEGA, NEG B	Negate accumulator or memory
ORA, ORB	Or memory with accumulator
ROL, ROLA, ROLB	Rotate accumulator or memory left
ROR, RORA, RORB	Rotate accumulator or memory right
SBCA, SBCB	Subtract memory from accumulator with borrow
STA, STB	Store accumulator to memory
SUBA, SUBB	Subtract memory from accumulator
TST, TSTA, TSTB	Test accumulator or memory location
TFR R1, R2	Transfer R1 to R2 (R1, R2=A, B, CC, DP)

NOTE: A, B, CC, or DP may be pushed to (pulled from) either stack with PSHS, PSHU (PULS, PULU) instructions.

Table 4-3. 16-Bit Accumulator and Memory Instructions

Instruction	Description
ADDD	Add memory to D accumulator
CMPD	Compare memory from D accumulator
EXG D, R	Exchange D with X, Y, S, U, or PC
LDD	Load D accumulator from memory
SEX	Sign Extend B accumulator into A accumulator
STD	Store D accumulator to memory
SUBD	Subtract memory from D accumulator
TFR D, R	Transfer D to X, Y, S, U, or PC
TFR R, D	Transfer X, Y, S, U, or PC to D

NOTE: D may be pushed (pulled) to either stack with PSHS, PSHU (PULS, PULU) instructions.

Table 4-4. Index/Stack Pointer Instructions

Instruction	Description
CMPS, CMPU	Compare memory from stack pointer
CMPX, CMPY	Compare memory from index register
EXG R1, R2	Exchange D, X, Y, S, U or PC with D, X, Y, S, U or PC
LEAS, LEAU	Load effective address into stack pointer
LEAX, LEAY	Load effective address into index register
LDS, LDU	Load stack pointer from memory
LDX, LDY	Load index register from memory
PSHS	Push A, B, CC, DP, D, X, Y, U, or PC onto hardware stack
PSHU	Push A, B, CC, DP, D, X, Y, X, or PC onto user stack
PULS	Pull A, B, CC, DP, D, X, Y, U, or PC from hardware stack
PULU	Pull A, B, CC, DP, D, X, Y, S, or PC from hardware stack
STS, STU	Store stack pointer to memory
STX, STY	Store index register to memory
TFR R1, R2	Transfer D, X, Y, S, U, or PC to D, X, Y, S, U, or PC
ABX	Add B accumulator to X (unsigned)

Table 4-5. Branch Instructions

Instruction	Description
SIMPLE BRANCHES	
BEQ, LB EQ	Branch if equal
BNE, LB NE	Branch if not equal
BMI, LB MI	Branch if minus
BPL, LB PL	Branch if plus
BCS, LB CS	Branch if carry set
BCC, LB CC	Branch if carry clear
BVS, LB VS	Branch if overflow set
BVC, LB VC	Branch if overflow clear
SIGNED BRANCHES	
BGT, LB GT	Branch if greater (signed)
BVS, LB VS	Branch if invalid twos complement result
BGE, LB GE	Branch if greater than or equal (signed)
BEQ, LB EQ	Branch if equal
BNE, LB NE	Branch if not equal
BLE, LB LE	Branch if less than or equal (signed)
BVC, LB VC	Branch if valid twos complement result
BLT, LB LT	Branch if less than (signed)
UNSIGNED BRANCHES	
BHI, LB HI	Branch if higher (unsigned)
BCC, LB CC	Branch if higher or same (unsigned)
BHS, LB HS	Branch if higher or same (unsigned)
BEQ, LB EQ	Branch if equal
BNE, LB NE	Branch if not equal
BLS, LB LS	Branch if lower or same (unsigned)
BCS, LB CS	Branch if lower (unsigned)
BLO, LB LO	Branch if lower (unsigned)
OTHER BRANCHES	
BSR, LB SR	Branch to subroutine
BRA, LB RA	Branch always
BRN, LB RN	Branch never

Table 4-6. Miscellaneous Instructions

Instruction	Description
ANDCC	AND condition code register
CWAI	AND condition code register, then wait for interrupt
NOP	No operation
ORCC	OR condition code register
JMP	Jump
JSR	Jump to subroutine
RTI	Return from interrupt
RTS	Return from subroutine
SWI, SWI2, SWI3	Software interrupt (absolute indirect)
SYNC	Synchronize with interrupt line

APPENDIX A INSTRUCTION SET DETAILS

A.1 INTRODUCTION

This appendix contains detailed information about each instruction in the MC6809 instruction set. They are arranged in an alphabetical order with the mnemonic heading set in larger type for easy reference.

A.2 NOTATION

In the operation description for each instruction, symbols are used to indicate the operation. Table A-1 lists these symbols and their meanings. Abbreviations for the various registers, bits, and bytes are also used. Table A-2 lists these abbreviations and their meanings.

Table A-1. Operation Notation

<u>Symbol</u>	<u>Meaning</u>
\leftarrow	Is transferred to
\wedge	Boolean AND
\vee	Boolean OR
\oplus	Boolean exclusive OR
$\overline{-}$ (Overline)	Boolean NOT
:	Concatenation
+	Arithmetic plus
-	Arithmetic minus
\times	Arithmetic multiply

Table A-2. Register Notation

<u>Abbreviation</u>	<u>Meaning</u>
ACCA or A	Accumulator A
ACCB or B	Accumulator B
ACCA:ACCB or D	Double accumulator D
ACCX	Either accumulator A or B
CCR or CC	Condition code register
DPR or DP	Direct page register
EA	Effective address
IFF	If and only if
IX or X	Index register X
IY or Y	Index register Y
LSN	Least significant nibble
M	Memory location
MI	Memory immediate
MSN	Most significant nibble
PC	Program counter
R	A register before the operation
R'	A register after the operation
TEMP	Temporary storage location
xxH	Most significant byte of any 16-bit register
xxL	Least significant byte of any 16-bit register
Sp or S	Hardware Stack pointer
Us or U	User Stack pointer
P	A memory argument with Immediate, Direct, Extended, and Indexed addressing modes
Q	A read-modify-write argument with Direct, Indexed, and Extended addressing modes
()	The data pointed to by the enclosed (16-bit address)
dd	8-bit branch offset
DDDD	16-bit branch offset
#	Immediate value follows
\$	Hexadecimal value follows
{ }	Indirection
'	Indicates indexed addressing

ABX

Add Accumulator B into Index Register X

ABX

Source Form: ABX

Operation: $IX' \leftarrow IX + ACCB$

Condition Codes: Not affected.

Description: Add the 8-bit unsigned value in accumulator B into index register X.

Addressing Mode: Inherent

ADC

Add with Carry into Register

ADC

Source Forms: ADCA P; ADCB P

Operation: $R' \leftarrow R + M + C$

Condition Codes:

- H — Set if a half-carry is generated; cleared otherwise.
- N — Set if the result is negative; cleared otherwise.
- Z — Set if the result is zero; cleared otherwise.
- V — Set if an overflow is generated; cleared otherwise.
- C — Set if a carry is generated; cleared otherwise.

Description: Adds the contents of the C (carry) bit and the memory byte into an 8-bit accumulator.

Addressing Modes: Immediate
Extended
Direct
Indexed

ADD (8-Bit)

Add Memory into Register

ADD (8-Bit)

Source Forms: ADDA P; ADDB P

Operation: $R' \leftarrow R + M$

Condition Codes: H — Set if a half-carry is generated; cleared otherwise.
N — Set if the result is negative; cleared otherwise.
Z — Set if the result is zero; cleared otherwise.
V — Set if an overflow is generated; cleared otherwise.
C — Set if a carry is generated; cleared otherwise.

Description: Adds the memory byte into an 8-bit accumulator.

Addressing Modes: Immediate
Extended
Direct
Indexed

ADD (16-Bit)

Add Memory Into Register

ADD (16-Bit)

Source Forms: ADDD P

Operation: $R' \leftarrow R + M:M + 1$

Condition Codes: H — Not affected.

N — Set if the result is negative; cleared otherwise.

Z — Set if the result is zero; cleared otherwise.

V — Set if an overflow is generated; cleared otherwise.

C — Set if a carry is generated; cleared otherwise.

Description: Adds the 16-bit memory value into the 16-bit accumulator

Addressing Modes: Immediate

Extended

Direct

Indexed

AND

Logical AND Memory into Register

AND

Source Forms: ANDA P; ANDB P

Operation: $R' \leftarrow R \wedge M$

Condition Codes:

- H — Not affected.
- N — Set if the result is negative; cleared otherwise.
- Z — Set if the result is zero; cleared otherwise.
- V — Always cleared.
- C — Not affected.

Description: Performs the logical AND operation between the contents of an accumulator and the contents of memory location M and the result is stored in the accumulator.

Addressing Modes: Immediate
Extended
Direct
Indexed

AND Logical AND Immediate Memory Into Condition Code Register **AND**

Source Form: ANDCC #xx

Operation: $R' \leftarrow R \wedge M$

Condition Codes: Affected according to the operation.

Description: Performs a logical AND between the condition code register and the immediate byte specified in the instruction and places the result in the condition code register.

Addressing Mode: Immediate

ASL

Arithmetic Shift Left

ASL

Source Forms: ASL Q; ASLA; ASLB

Operation: $C \leftarrow \boxed{\quad \quad \quad \quad \quad \quad \quad} \leftarrow 0$

Condition Codes: H — Undefined

N — Set if the result is negative; cleared otherwise.

Z — Set if the result is zero; cleared otherwise.

V — Loaded with the result of the exclusive OR of bits six and seven of the original operand.

C — Loaded with bit seven of the original operand.

Description: Shifts all bits of the operand one place to the left. Bit zero is loaded with a zero. Bit seven is shifted into the C (carry) bit.

Addressing Modes: Inherent
Extended
Direct
Indexed

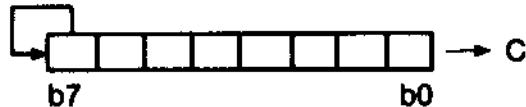
ASR

Arithmetic Shift Right

ASR

Source Forms: ASR Q; ASRA; ASRB

Operation:



Condition Codes: H — Undefined.

N — Set if the result is negative; cleared otherwise.

Z — Set if the result is zero; cleared otherwise.

V — Not affected.

C — Loaded with bit zero of the original operand.

Description: Shifts all bits of the operand one place to the right. Bit seven is held constant. Bit zero is shifted into the C (carry) bit.

Addressing Modes: Inherent
Extended
Direct
Indexed

BCC

Branch on Carry Clear

BCC

Source Forms: BCC dd; LBCC DDDD

Operation: TEMP ← MI
IFF C = 0 then PC' ← PC + TEMP

Condition Codes: Not affected.

Description: Tests the state of the C (carry) bit and causes a branch if it is clear.

Addressing Mode: Relative

Comments: Equivalent to BHS dd; LBHS DDDD

BCS**Branch on Carry Set****BCS****Source Forms:** BCS dd; LBCS DDDD**Operation:** TEMP \leftarrow MI
IFF C = 1 then PC' \leftarrow PC + TEMP**Condition Codes:** Not affected.**Description:** Tests the state of the C (carry) bit and causes a branch if it is set.**Addressing Mode:** Relative**Comments:** Equivalent to BLO dd; LBLO DDDD

BEQ

Branch on Equal

BEQ

Source Forms: BEQ dd; LBEQ DDDD

Operation: TEMP \leftarrow MI
IFF Z = 1 then PC' \leftarrow PC + TEMP

Condition Codes: Not affected.

Description: Tests the state of the Z (zero) bit and causes a branch if it is set. When used after a subtract or compare operation, this instruction will branch if the compared values, signed or unsigned, were exactly the same.

Addressing Mode: Relative

BGE

Branch on Greater than or Equal to Zero

BGE**Source Forms:** BGE dd; LBGE DDDD**Operation:** TEMP \leftarrow MI
IFF [N \oplus V] = 0 then PC' \leftarrow PC + TEMP**Condition Codes:** Not affected.**Description:** Causes a branch if the N (negative) bit and the V (overflow) bit are either both set or both clear. That is, branch if the sign of a valid twos complement result is, or would be, positive. When used after a subtract or compare operation on twos complement values, this instruction will branch if the register was greater than or equal to the memory operand.**Addressing Mode:** Relative

BGT

Branch on Greater

BGT

- Source Forms:** BGT dd; LBGT DDDD
- Operation:** TEMP ← MI
IFF Z A [N ⊕ V] = 0 then PC' ← PC + TEMP
- Condition Codes:** Not affected.
- Description:** Causes a branch if the N (negative) bit and V (overflow) bit are either both set or both clear and the Z (zero) bit is clear. In other words, branch if the sign of a valid twos complement result is, or would be, positive and not zero. When used after a subtract or compare operation on twos complement values, this instruction will branch if the register was greater than the memory operand.
- Addressing Mode:** Relative

BHI**Branch If Higher****BHI**

Source Forms: BHI dd; LBHI DDDD

Operation: TEMP ← M₁
IFF [C v Z]=0 then PC' ← PC + TEMP

Condition Codes: Not affected.

Description: Causes a branch if the previous operation caused neither a carry nor a zero result. When used after a subtract or compare operation on unsigned binary values, this instruction will branch if the register was higher than the memory operand.

Addressing Mode: Relative

Comments: Generally not useful after INC/DEC, LD/TST, and TST/CLR/COM instructions.

BHS

Branch if Higher or Same

BHS

- Source Forms:** BHS dd; LBHS DDDD
- Operation:** TEMP ← MI
IFF C = 0 then PC' ← PC + MI
- Condition Codes:** Not affected.
- Description:** Tests the state of the C (carry) bit and causes a branch if it is clear. When used after a subtract or compare on unsigned binary values, this instruction will branch if the register was higher than or the same as the memory operand.
- Addressing Mode:** Relative
- Comments:** This is a duplicate assembly-language mnemonic for the single machine instruction BCC. Generally not useful after INC/DEC, LD/ST, and TST/CLR/COM instructions.

BIT

Bit Test

BIT

Source Form: Bit P

Operation: TEMP \leftarrow R \wedge M

Condition Codes:

- H — Not affected.
- N — Set if the result is negative; cleared otherwise.
- Z — Set if the result is zero; cleared otherwise.
- V — Always cleared.
- C — Not affected.

Description: Performs the logical AND of the contents of accumulator A or B and the contents of memory location M and modifies the condition codes accordingly. The contents of accumulator A or B and memory location M are not affected.

Addressing Modes: Immediate
Extended
Direct
Indexed

BLE**Branch on Less than or Equal to Zero****BLE**

- Source Forms:** BLE dd; LBLE DDDD
- Operation:** TEMP ← MI
IFF $Z \vee [N \oplus V] = 1$ then $PC' \leftarrow PC + TEMP$
- Condition Codes:** Not affected.
- Description:** Causes a branch if the exclusive OR of the N (negative) and V (overflow) bits is 1 or if the Z (zero) bit is set. That is, branch if the sign of a valid two's complement result is, or would be, negative. When used after a subtract or compare operation on two's complement values, this instruction will branch if the register was less than or equal to the memory operand.
- Addressing Mode:** Relative

BLO

Branch on Lower

BLO**Source Forms:** BLO dd; LBLO DDDD**Operation:** TEMP ← MI
IFF C = 1 then PC' ← PC + TEMP**Condition Codes:** Not affected.**Description:** Tests the state of the C (carry) bit and causes a branch if it is set. When used after a subtract or compare on unsigned binary values, this instruction will branch if the register was lower than the memory operand.**Addressing Mode:** Relative**Comments:** This is a duplicate assembly-language mnemonic for the single machine instruction BCS. Generally not useful after INC/DEC, LD/ST, and TST/CLR/COM instructions.

BLS

Branch on Lower or Same

BLS

Source Forms: BLS dd; LBLS DDDD

Operation: TEMP \leftarrow MI
IFF (C v Z) = 1 then PC' \leftarrow PC + TEMP

Condition Codes: Not affected.

Description: Causes a branch if the previous operation caused either a carry or a zero result. When used after a subtract or compare operation on unsigned binary values, this instruction will branch if the register was lower than or the same as the memory operand.

Addressing Mode: Relative

Comments: Generally not useful after INC/DEC, LD/ST, and TST/CLR/COM instructions.

BLT**Branch on Less than Zero****BLT****Source Forms:** BLT dd; LBLT DDDD**Operation:** TEMP ← MI
IFF [N ⊕ V] = 1 then PC' ← PC + TEMP**Condition Codes:** Not affected.**Description:** Causes a branch if either, but not both, of the N (negative) or V (overflow) bits is set. That is, branch if the sign of a valid two's complement result is, or would be, negative. When used after a subtract or compare operation on two's complement binary values, this instruction will branch if the register was less than the memory operand.**Addressing Mode:** Relative

BMI

Branch on Minus

BMI

Source Forms: BMI dd; LBMI DDDD

Operation: TEMP \leftarrow MI
IFF N = 1 then PC' \leftarrow PC + TEMP

Condition Codes: Not affected.

Description: Tests the state of the N (negative) bit and causes a branch if set.
That is, branch if the sign of the two's complement result is negative.

Addressing Mode: Relative

Comments: When used after an operation on signed binary values, this instruction will branch if the result is minus. It is generally preferred to use the LBLT instruction after signed operations.

BNE

Branch Not Equal

BNE

Source Forms: BNE dd; LBNE DDDD

Operation: TEMP \leftarrow M1
IFF Z = 0 then PC' \leftarrow PC + TEMP

Condition Codes: Not affected.

Description: Tests the state of the Z (zero) bit and causes a branch if it is clear. When used after a subtract or compare operation on any binary values, this instruction will branch if the register is, or would be, not equal to the memory operand.

Addressing Mode: Relative

BPL

Branch on Plus

BPL

Source Forms: BPL dd; LBPL DDDD

Operation: TEMP ← MI
IFF N = 0 then PC' ← PC + TEMP

Condition Codes: Not affected.

Description: Tests the state of the N (negative) bit and causes a branch if it is clear. That is, branch if the sign of the two's complement result is positive.

Addressing Mode: Relative

Comments: When used after an operation on signed binary values, this instruction will branch if the result (possibly invalid) is positive. It is generally preferred to use the BGE instruction after signed operations.

BRA

Branch Always

BRA

Source Forms: BRA dd; LBRA DDDD

Operation: TEMP ← MI
PC' ← PC + TEMP

Condition Codes: Not affected.

Description: Causes an unconditional branch.

Addressing Mode: Relative

BRN

Branch Never

BRN

Source Forms: BRN dd; LBRN DDDD

Operation: TEMP ← MI

Condition Codes: Not affected.

Description: Does not cause a branch. This instruction is essentially a no operation, but has a bit pattern logically related to branch always.

Addressing Mode: Relative

BSR

Branch to Subroutine

BSR

Source Forms: BSR dd; LBSR DDDD

Operation:

$$\begin{aligned} \text{TEMP} &\leftarrow \text{MI} \\ \text{SP}' &\leftarrow \text{SP} - 1, (\text{SP}) \leftarrow \text{PCL} \\ \text{SP}' &\leftarrow \text{SP} - 1, (\text{SP}) \leftarrow \text{PCH} \\ \text{PC}' &\leftarrow \text{PC} + \text{TEMP} \end{aligned}$$

Condition Codes: Not affected.

Description: The program counter is pushed onto the stack. The program counter is then loaded with the sum of the program counter and the offset.

Addressing Mode: Relative

Comments: A return from subroutine (RTS) instruction is used to reverse this process and must be the last instruction executed in a subroutine.

BVC**Branch on Overflow Clear****BVC****Source Forms:** BVC dd; LBVC DDDD**Operation:** TEMP \leftarrow MI
IFF V = 0 then PC' \leftarrow PC + TEMP**Condition Codes:** Not affected.**Description:** Tests the state of the V (overflow) bit and causes a branch if it is clear. That is, branch if the twos complement result was valid. When used after an operation on twos complement binary values, this instruction will branch if there was no overflow.**Addressing Mode:** Relative

BVS**Branch on Overflow Set****BVS**

Source Forms: BVS dd; LBVS DDDD

Operation: TEMP ← MI
IFF V = 1 then PC' ← PC + TEMP

Condition Codes: Not affected.

Description: Tests the state of the V (overflow) bit and causes a branch if it is set. That is, branch if the twos complement result was invalid. When used after an operation on twos complement binary values, this instruction will branch if there was an overflow.

Addressing Mode: Relative

CLR

Clear

CLR

Source Form:	CLR Q
Operation:	TEMP ← M M ← 0016
Condition Codes:	H — Not affected. N — Always cleared. Z — Always set. V — Always cleared. C — Always cleared.
Description:	Accumulator A or B or memory location M is loaded with 00000000. Note that the EA is read during this operation.
Addressing Modes:	Inherent Extended Direct Indexed

CMP (8-Bit) Compare Memory from Register

CMP (8-Bit)

Source Forms: CMPA P; CMPB P

Operation: TEMP \leftarrow R - M

Condition Codes: H — Undefined.

N — Set if the result is negative; cleared otherwise.

Z — Set if the result is zero; cleared otherwise.

V — Set if an overflow is generated; cleared otherwise.

C — Set if a borrow is generated; cleared otherwise.

Description: Compares the contents of memory location to the contents of the specified register and sets the appropriate condition codes. Neither memory location M nor the specified register is modified. The carry flag represents a borrow and is set to the inverse of the resulting binary carry.

Addressing Modes: Immediate

Extended

Direct

Indexed

CMP (16-Bit) Compare Memory from Register CMP (16-Bit)

Source Forms: CMPD P; CMPX P; CMPY P; CMPU P; CMPS P

Operation: TEMP \leftarrow R - M:M + 1

Condition Codes: H — Not affected.

N — Set if the result is negative; cleared otherwise.

Z — Set if the result is zero; cleared otherwise.

V — Set if an overflow is generated; cleared otherwise.

C — Set if a borrow is generated; cleared otherwise.

Description: Compares the 16-bit contents of the concatenated memory locations M:M + 1 to the contents of the specified register and sets the appropriate condition codes. Neither the memory locations nor the specified register is modified unless autoincrement or autodecrement are used. The carry flag represents a borrow and is set to the inverse of the resulting binary carry.

Addressing Modes: Immediate
Extended
Direct
Indexed

COM

Complement

COM

Source Forms: COM Q; COMA; COMB

Operation: $M' \leftarrow O + \bar{M}$

Condition Codes:
H — Not affected.
N — Set if the result is negative; cleared otherwise.
Z — Set if the result is zero; cleared otherwise.
V — Always cleared.
C — Always set.

Description: Replaces the contents of memory location M or accumulator A or B with its logical complement. When operating on unsigned values, only BEQ and BNE branches can be expected to behave properly following a COM instruction. When operating on two's complement values, all signed branches are available.

Addressing Modes: Inherent
Extended
Direct
Indexed

CWAI

Clear CC bits and Wait for Interrupt

CWAI

Source Form:

CWAI #\$XX

E	F	H	I	N	Z	V	C
---	---	---	---	---	---	---	---

Operation:

CCR \leftarrow CCR \wedge MI (Possibly clear masks)
Set E (entire state saved)
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow DPR$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCB$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCA$
 $SP' \leftarrow SP - 1, (SP) \leftarrow CCR$

Condition Codes: Affected according to the operation.

Description:

This instruction ANDs an immediate byte with the condition code register which may clear the interrupt mask bits I and F, stacks the entire machine state on the hardware stack and then looks for an interrupt. When a non-masked interrupt occurs, no further machine state information need be saved before vectoring to the interrupt handling routine. This instruction replaced the MC6800 CLI WAI sequence, but does not place the buses in a high-impedance state. A FIRQ (fast interrupt request) may enter its interrupt handler with its entire machine state saved. The RTI (return from interrupt) instruction will automatically return the entire machine state after testing the E (entire) bit of the recovered condition code register.

Addressing Mode: Immediate

Comments: The following immediate values will have the following results:

FF = enable neither

EF = enable IRQ

BF = enable FIRQ

AF = enable both

DAA

Decimal Addition Adjust

DAA

Source Form: DAA

Operation: $ACCA' \leftarrow ACCA + CF(MSN):CF(LSN)$
where CF is a Correction Factor, as follows: the CF for each nibble (BCD) digit is determined separately, and is either 6 or 0.

Least Significant Nibble

$CF(LSN) = 6$ IFF 1) $C = 1$
or 2) $LSN > 9$

Most Significant Nibble

$CF(MSN) = 6$ IFF 1) $C = 1$
or 2) $MSN > 9$
or 3) $MSN > 8$ and $LSN > 9$

Condition Codes: H — Not affected.
N — Set if the result is negative; cleared otherwise.
Z — Set if the result is zero; cleared otherwise.
V — Undefined.
C — Set if a carry is generated or if the carry bit was set before the operation; cleared otherwise.

Description: The sequence of a single-byte add instruction on accumulator A (either ADDA or ADCA) and a following decimal addition adjust instruction results in a BCD addition with an appropriate carry bit. Both values to be added must be in proper BCD form (each nibble such that: $0 \leq \text{nibble} \leq 9$). Multiple-precision addition must add the carry generated by this decimal addition adjust into the next higher digit during the add operation (ADCA) immediately prior to the next decimal addition adjust.

Addressing Mode: Inherent

DEC

Decrement

DEC

Source Forms: DEC Q; DECA; DECB

Operation: $M' \leftarrow M - 1$

Condition Codes: H — Not affected.
N — Set if the result is negative; cleared otherwise.
Z — Set if the result is zero; cleared otherwise.
V — Set if the original operand was 10000000; cleared otherwise.
C — Not affected.

Description: Subtract one from the operand. The carry bit is not affected, thus allowing this instruction to be used as a loop counter in multiple-precision computations. When operating on unsigned values, only BEQ and BNE branches can be expected to behave consistently. When operating on two's complement values, all signed branches are available.

Addressing Modes: Inherent
Extended
Direct
Indexed

EOR

Exclusive OR

EOR

Source Forms: EORA P; EORB P

Operation: $R' \leftarrow R \oplus M$

Condition Codes: H — Not affected.
N — Set if the result is negative; cleared otherwise.
Z — Set if the result is zero; cleared otherwise.
V — Always cleared.
C — Not affected.

Description: The contents of memory location M is exclusive ORed into an 8-bit register.

Addressing Modes: Immediate
Extended
Direct
Indexed

EXG**Exchange Registers****EXG****Source Form:** EXG R1,R2**Operation:** R1 \leftarrow R2**Condition Codes:** Not affected (unless one of the registers is the condition code register).**Description:** Exchanges data between two designated registers. Bits 3-0 of the postbyte define one register, while bits 7-4 define the other, as follows:

0000 = A:B	1000 = A
0001 = X	1001 = B
0010 = Y	1010 = CCR
0011 = US	1011 = DPR
0100 = SP	1100 = Undefined
0101 = PC	1101 = Undefined
0110 = Undefined	1110 = Undefined
0111 = Undefined	1111 = Undefined

Only like size registers may be exchanged. (8-bit with 8-bit or 16-bit with 16-bit.)

Addressing Mode: Immediate

INC

Increment

INC**Source Forms:** INC Q; INCA; INCB**Operation:** $M' \leftarrow M + 1$ **Condition Codes:** H — Not affected.

N — Set if the result is negative; cleared otherwise.

Z — Set if the result is zero; cleared otherwise.

V — Set if the original operand was 01111111; cleared otherwise.

C — Not affected.

Description: Adds to the operand. The carry bit is not affected, thus allowing this instruction to be used as a loop counter in multiple-precision computations. When operating on unsigned values, only the BEQ and BNE branches can be expected to behave consistently. When operating on two's complement values, all signed branches are correctly available.**Addressing Modes:** Inherent
Extended
Direct
Indexed

JMP

Jump

JMP

Source Form: JMP EA

Operation: $PC' \leftarrow EA$

Condition Codes: Not affected.

Description: Program control is transferred to the effective address.

Addressing Modes: Extended
Direct
Indexed

JSR

Jump to Subroutine

JSR

Source Form: JSR EA

Operation:
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCH$
 $PC' \leftarrow EA$

Condition Codes: Not affected.

Description: Program control is transferred to the effective address after storing the return address on the hardware stack. A RTS instruction should be the last executed instruction of the subroutine.

Addressing Modes: Extended
Direct
Indexed

LD (8-Bit)

Load Register from Memory

LD (8-Bit)

Source Forms: LDA P; LDB P

Operation: $R' \leftarrow M$

Condition Codes: H — Not affected.

N — Set if the loaded data is negative; cleared otherwise.

Z — Set if the loaded data is zero; cleared otherwise.

V — Always cleared.

C — Not affected.

Description: Loads the contents of memory location M into the designated register.

Addressing Modes: Immediate

Extended

Direct

Indexed

LD (16-Bit)

Load Register from Memory

LD (16-Bit)

Source Forms: LDD P; LDX P; LDY P; LDS P; LDU P

Operation: $R' \leftarrow M:M + 1$

Condition Codes: H — Not affected.

N — Set if the loaded data is negative; cleared otherwise.

Z — Set if the loaded data is zero; cleared otherwise.

V — Always cleared.

C — Not affected.

Description: Load the contents of the memory location M:M + 1 into the designated 16-bit register.

Addressing Modes: Immediate

Extended

Direct

Indexed

LEA

Load Effective Address

LEA

Source Forms: LEAX, LEAY, LEAS, LEAU

Operation: $R' \leftarrow EA$

Condition Codes: H — Not affected.
N — Not affected.
Z — LEAX, LEAY: Set if the result is zero; cleared otherwise.
LEAS, LEAU: Not affected.
V — Not affected.
C — Not affected.

Description: Calculates the effective address from the indexed addressing mode and places the address in an indexable register.

LEAX and LEAY affect the Z (zero) bit to allow use of these registers as counters and for MC6800 INX/DEX compatibility.

LEAU and LEAS do not affect the Z bit to allow cleaning up the stack while returning the Z bit as a parameter to a calling routine, and also for MC6800 INS/DES compatibility.

Addressing Mode: Indexed

Comments: Due to the order in which effective addresses are calculated internally, the LEAX, X + + and LEAX, X + do not add 2 and 1 (respectively) to the X register; but instead leave the X register unchanged. This also applies to the Y, U, and S registers. For the expected results, use the faster instruction LEAX 2, X and LEAX 1, X.

Some examples of LEA instruction uses are given in the following table.

Instruction	Operation	Comment
LEAX 10, X	$X + 10 - X$	Adds 5-bit constant 10 to X
LEAX 500, X	$X + 500 - X$	Adds 16-bit constant 500 to X
LEAY A, Y	$Y + A - Y$	Adds 8-bit accumulator to Y
LEAY D, Y	$Y + D - Y$	Adds 16-bit D accumulator to Y
LEAU - 10, U	$U - 10 - U$	Subtracts 10 from U
LEAS - 10, S	$S - 10 - S$	Used to reserve area on stack
LEAS 10, S	$S + 10 - S$	Used to 'clean up' stack
LEAX 5, S	$S + 5 - X$	Transfers as well as adds

LSL

Logical Shift Left

LSL

Source Forms: LSL Q; LSLA; LSLB

Operation: C ←  ← 0
b7 b0

Condition Codes:
H — Undefined.
N — Set if the result is negative; cleared otherwise.
Z — Set if the result is zero; cleared otherwise.
V — Loaded with the result of the exclusive OR of bits six and seven of the original operand.
C — Loaded with bit seven of the original operand.

Description: Shifts all bits of accumulator A or B or memory location M one place to the left. Bit zero is loaded with a zero. Bit seven of accumulator A or B or memory location M is shifted into the C (carry) bit.

Addressing Modes: Inherent
Extended
Direct
Indexed

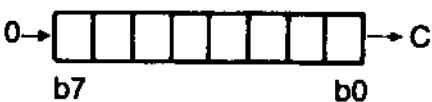
Comments: This is a duplicate assembly-language mnemonic for the single machine instruction ASL.

LSR

Logical Shift Right

LSR

Source Forms: LSR Q; LSRA; LSRB

Operation: 

Condition Codes: H — Not affected.
N — Always cleared.
Z — Set if the result is zero; cleared otherwise.
V — Not affected.
C — Loaded with bit zero of the original operand.

Description: Performs a logical shift right on the operand. Shifts a zero into bit seven and bit zero into the C (carry) bit.

Addressing Modes: Inherent
Extended
Direct
Indexed

MUL

~~multiply~~

MUL

Source Form: MUL

Operation: ACCA':ACCB' ← ACCA × ACCB

Condition Codes: H — Not affected.
N — Not affected.
Z — Set if the result is zero; cleared otherwise.
V — Not affected.
C — Set if ACCB bit 7 of result is set; cleared otherwise.

Description: Multiply the unsigned binary numbers in the accumulators and place the result in both accumulators (ACCA contains the most-significant byte of the result). Unsigned multiply allows multiple-precision operations.

Addressing Mode: Inherent

Comments: The C (carry) bit allows rounding the most-significant byte through the sequence: MUL, ADCA #0.

NEG

Negate

NEG

Source Forms: NEG Q; NEGA; NEGB

Operation: $M' \leftarrow 0 - M$

Condition Codes: H — Undefined.

N — Set if the result is negative; cleared otherwise.

Z — Set if the result is zero; cleared otherwise.

V — Set if the original operand was 10000000.

C — Set if a borrow is generated; cleared otherwise.

Description: Replaces the operand with its twos complement. The C (carry) bit represents a borrow and is set to the inverse of the resulting binary carry. Note that 80₁₆ is replaced by itself and only in this case is the V (overflow) bit set. The value 00₁₆ is also replaced by itself, and only in this case is the C (carry) bit cleared.

Addressing Modes: Inherent
Extended
Direct

NOP

No Operation

NOP

Source Form: NOP

Operation: Not affected.

Condition Codes: This instruction causes only the program counter to be incremented.
No other registers or memory locations are affected.

Addressing Mode: Inherent

OR

Inclusive OR Memory into Register

OR

Source Forms: ORA P; ORB P

Operation: $R' \leftarrow R \vee M$

Condition Codes:

- H — Not affected.
- N — Set if the result is negative; cleared otherwise.
- Z — Set if the result is zero; cleared otherwise.
- V — Always cleared.
- C — Not affected.

Description: Performs an inclusive OR operation between the contents of accumulator A or B and the contents of memory location M and the result is stored in accumulator A or B.

Addressing Modes: Immediate
Extended
Direct
Indexed

OR**Inclusive OR Memory Immediate Into Condition Code Register****OR****Source Form:** ORCC #XX**Operation:** $R' \leftarrow R \vee MI$ **Condition Codes:** Affected according to the operation.**Description:** Performs an Inclusive OR operation between the contents of the condition code registers and the immediate value, and the result is placed in the condition code register. This instruction may be used to set interrupt masks (disable interrupts) or any other bit(s).**Addressing Mode:** Immediate

PSHS

Push Registers on the Hardware Stack

PSHS

Source Form:

PSHS register list

PSHS #LABEL

Postbyte:

b7	b6	b5	b4	b3	b2	b1	b0
PC	U	Y	X	DP	B	A	CC

push order----->

Operation:

IFF b7 of postbyte set, then: $SP' \leftarrow SP - 1$, $(SP) \leftarrow PCL$
 $SP' \leftarrow SP - 1$, $(SP) \leftarrow PCH$

IFF b6 of postbyte set, then: $SP' \leftarrow SP - 1$, $(SP) \leftarrow USL$
 $SP' \leftarrow SP - 1$, $(SP) \leftarrow USH$

IFF b5 of postbyte set, then: $SP' \leftarrow SP - 1$, $(SP) \leftarrow IYL$
 $SP' \leftarrow SP - 1$, $(SP) \leftarrow IYH$

IFF b4 of postbyte set, then: $SP' \leftarrow SP - 1$, $(SP) \leftarrow IXL$
 $SP' \leftarrow SP - 1$, $(SP) \leftarrow IXH$

IFF b3 of postbyte set, then: $SP' \leftarrow SP - 1$, $(SP) \leftarrow DPR$

IFF b2 of postbyte set, then: $SP' \leftarrow SP - 1$, $(SP) \leftarrow ACCB$

IFF b1 of postbyte set, then: $SP' \leftarrow SP - 1$, $(SP) \leftarrow ACCA$

IFF b0 of postbyte set, then: $SP' \leftarrow SP - 1$, $(SP) \leftarrow CCR$

Condition Codes: Not affected.

Description: All, some, or none of the processor registers are pushed onto the hardware stack (with the exception of the hardware stack pointer itself).

Addressing Mode: Immediate

Comments: A single register may be placed on the stack with the condition codes set by doing an autodecrement store onto the stack (example: STX, -- S).

PSHU

Push Registers on the User Stack

PSHU

Source Form: PSHU register list

PSHU #LABEL

Postbyte:

b7	b6	b5	b4	b3	b2	b1	b0
PC	U	Y	X	DP	B	A	CC

push order----->

Operation:

IFF b7 of postbyte set, then: $US' \leftarrow US - 1$, $(US) \leftarrow PCL$
 $US' \leftarrow US - 1$, $(US) \leftarrow PCH$

IFF b6 of postbyte set, then: $US' \leftarrow US - 1$, $(US) \leftarrow SPL$
 $US' \leftarrow US - 1$, $(US) \leftarrow SPH$

IFF b5 of postbyte set, then: $US' \leftarrow US - 1$, $(US) \leftarrow IYL$
 $US' \leftarrow US - 1$, $(US) \leftarrow IYH$

IFF b4 of postbyte set, then: $US' \leftarrow US - 1$, $(US) \leftarrow IXL$
 $US' \leftarrow US - 1$, $(US) \leftarrow IXH$

IFF b3 of postbyte set, then: $US' \leftarrow US - 1$, $(US) \leftarrow DPR$

IFF b2 of postbyte set, then: $US' \leftarrow US - 1$, $(US) \leftarrow ACCB$

IFF b1 of postbyte set, then: $US' \leftarrow US - 1$, $(US) \leftarrow ACCA$

IFF b0 of postbyte set, then: $US' \leftarrow US - 1$, $(US) \leftarrow CCR$

Condition Codes: Not affected.

Description: All, some, or none of the processor registers are pushed onto the user stack (with the exception of the user stack pointer itself).

Addressing Mode: Immediate

Comments: A single register may be placed on the stack with the condition codes set by doing an autodecrement store onto the stack (example: STX, -- U).

PULS

Pull Registers from the Hardware Stack

PULS

Source Form:

PULS register list

PULS #LABEL

Postbyte:

b7	b6	b5	b4	b3	b2	b1	b0
PC	U	Y	X	DP	B	A	CC

←-----pull order

Operation:

IFF b0 of postbyte set, then: CCR' ←(SP), SP'← SP + 1
IFF b1 of postbyte set, then: ACCA' ←(SP), SP'← SP + 1
IFF b2 of postbyte set, then: ACCB' ←(SP), SP'← SP + 1
IFF b3 of postbyte set, then: DPR' ←(SP), SP'← SP + 1
IFF b4 of postbyte set, then: IXH' ←(SP), SP'← SP + 1
IFF b5 of postbyte set, then: IXL' ←(SP), SP'← SP + 1
IFF b6 of postbyte set, then: IYH' ←(SP), SP'← SP + 1
IFF b7 of postbyte set, then: IYL' ←(SP), SP'← SP + 1
IFF b6 of postbyte set, then: USH' ←(SP), SP'← SP + 1
IFF b7 of postbyte set, then: USL' ←(SP), SP'← SP + 1
IFF b0 of postbyte set, then: PCH' ←(SP), SP'← SP + 1
IFF b1 of postbyte set, then: PCL' ←(SP), SP'← SP + 1

Condition Codes: May be pulled from stack; not affected otherwise.

Description: All, some, or none of the processor registers are pulled from the hardware stack (with the exception of the hardware stack pointer itself).

Addressing Mode: Immediate

Comments: A single register may be pulled from the stack with condition codes set by doing an autoincrement load from the stack (example: LDX ,S++).

PULU

Pull Registers from the User Stack

PULU

Source Form: PULU register list

PULU #LABEL

Postbyte:

b7 b6 b5 b4 b3 b2 b1 b0

PC	U	Y	X	DP	B	A	CC
----	---	---	---	----	---	---	----

←----- pull order

Operation: IFF b0 of postbyte set, then: CCR' ←(US), US'← US + 1
IFF b1 of postbyte set, then: ACCA' ←(US), US'← US + 1
IFF b2 of postbyte set, then: ACCB' ←(US), US'← US + 1
IFF b3 of postbyte set, then: DPR' ←(US), US'← US + 1
IFF b4 of postbyte set, then: IXH' ←(US), US'← US + 1
 IXL' ←(US), US'← US + 1
IFF b5 of postbyte set, then: IYH' ←(US), US'← US + 1
 IYL' ←(US), US'← US + 1
IFF b6 of postbyte set, then: SPH' ←(US), US'← US + 1
 SPL' ←(US), US'← US + 1
IFF b7 of postbyte set, then: PCH' ←(US), US'← US + 1
 PCL' ←(US), US'← US + 1

Condition Codes: May be pulled from stack; not affected otherwise.

Description: All, some, or none of the processor registers are pulled from the user stack (with the exception of the user stack pointer itself).

Addressing Mode: Immediate

Comments: A single register may be pulled from the stack with condition codes set by doing an autoincrement load from the stack (example: LDX ,U ++).

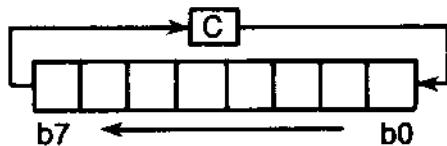
ROL

Rotate Left

ROL

Source Forms: ROL Q; ROLA; ROLB

Operation:



Condition Codes:

- H — Not affected.
- N — Set if the result is negative; cleared otherwise.
- Z — Set if the result is zero; cleared otherwise.
- V — Loaded with the result of the exclusive OR of bits six and seven of the original operand.
- C — Loaded with bit seven of the original operand.

Description: Rotates all bits of the operand one place left through the C (carry) bit. This is a 9-bit rotation.

Addressing Mode:

- Inherent
- Extended
- Direct
- Indexed

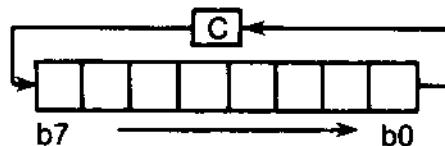
ROR

Rotate Right

ROR

Source Forms: ROR Q; RORA; RORB

Operation:



Condition Codes: H — Not affected.

N — Set if the result is negative; cleared otherwise.

Z — Set if the result is zero; cleared otherwise.

V — Not affected.

C — Loaded with bit zero of the previous operand.

Description: Rotates all bits of the operand one place right through the C (carry) bit. This is a 9-bit rotation.

Addressing Modes: Inherent
Extended
Direct
Indexed

RTI**Return from Interrupt****RTI****Source Form:** RTI**Operation:** $CCR' \leftarrow (SP)$, $SP' \leftarrow SP + 1$, thenIFF CCR bit E is set, then: $ACCA' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ $ACCB' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ $DPR' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ $IXH' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ $IXL' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ $IYH' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ $IYL' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ $USH' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ $USL' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ $PCH' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ $PCL' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ IFF CCR bit E is clear, then: $PCH' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ $PCL' \leftarrow (SP)$, $SP' \leftarrow SP + 1$ **Condition Codes:** Recovered from the stack.**Description:** The saved machine state is recovered from the hardware stack and control is returned to the interrupted program. If the recovered E (entire) bit is clear, it indicates that only a subset of the machine state was saved (return address and condition codes) and only that subset is recovered.**Addressing Mode:** Inherent

RTS

Return from Subroutine

RTS

Source Form: RTS

Operation: $PCH' \leftarrow (SP)$, $SP' \leftarrow SP + 1$
 $PCL' \leftarrow (SP)$, $SP' \leftarrow SP + 1$

Condition Codes: Not affected.

Description: Program control is returned from the subroutine to the calling program. The return address is pulled from the stack.

Addressing Mode: Inherent

SBC

Subtract with Borrow

SBC

Source Forms: SBCA P; SBCB P

Operation: $R' \leftarrow R - M - C$

Condition Codes:

- H — Undefined.
- N — Set if the result is negative; cleared otherwise.
- Z — Set if the result is zero; cleared otherwise.
- V — Set if an overflow is generated; cleared otherwise.
- C — Set if a borrow is generated; cleared otherwise.

Description: Subtracts the contents of memory location M and the borrow (in the C (carry) bit) from the contents of the designated 8-bit register, and places the result in that register. The C bit represents a borrow and is set to the inverse of the resulting binary carry.

Addressing Modes: Immediate
Extended
Direct
Indexed

SEX

Sign Extended

SEX

Source Form: SEX

Operation: If bit seven of ACCB is set then $ACCA' \leftarrow FF16$
else $ACCA' \leftarrow 0016$

Condition Codes: H — Not affected.
N — Set if the result is negative; cleared otherwise.
Z — Set if the result is zero; cleared otherwise.
V — Not affected.
C — Not affected.

Description: This instruction transforms a two's complement 8-bit value in accumulator B into a two's complement 16-bit value in the D accumulator.

Addressing Mode: Inherent

ST (8-Bit)

Store Register into Memory

ST (8-Bit)

Source Forms: STA P; STB P

Operation: $M' \leftarrow R$

Condition Codes: H — Not affected.
N — Set if the result is negative; cleared otherwise.
Z — Set if the result is zero; cleared otherwise.
V — Always cleared.
C — Not affected.

Description: Writes the contents of an 8-bit register into a memory location.

Addressing Modes: Extended
Direct
Indexed

ST (16-Bit)

Store Register Into Memory

ST (16-Bit)

Source Forms: STD P; STX P; STY P; STS P; STU P

Operation: $M':M + 1' \leftarrow R$

Condition Codes: H — Not affected.
N — Set if the result is negative; cleared otherwise.
Z — Set if the result is zero; cleared otherwise.
V — Always cleared.
C — Not affected.

Description: Writes the contents of a 16-bit register into two consecutive memory locations.

Addressing Modes: Extended

Direct

Indexed

SUB (8-Bit)

Subtract Memory from Register

SUB (8-Bit)

Source Forms: SUBA P; SUBB P

Operation: $R' \leftarrow R - M$

Condition Codes: H — Undefined.

N — Set if the result is negative; cleared otherwise.

Z — Set if the result is zero; cleared otherwise.

V — Set if the overflow is generated; cleared otherwise.

C — Set if a borrow is generated; cleared otherwise.

Description: Subtracts the value in memory location M from the contents of a designated 8-bit register. The C (carry) bit represents a borrow and is set to the inverse of the resulting binary carry.

Addressing Modes: Immediate
Extended
Direct
Indexed

SUB (16-Bit) Subtract Memory from Register **SUB (16-Bit)**

Source Forms: SUBD P

Operation: $R' \leftarrow R - M:M + 1$

Condition Codes: H — Not affected.

N — Set if the result is negative; cleared otherwise.

Z — Set if the result is zero; cleared otherwise.

V — Set if the overflow is generated; cleared otherwise.

C — Set if a borrow is generated; cleared otherwise.

Description: Subtracts the value in memory location M:M + 1 from the contents of a designated 16-bit register. The C (carry) bit represents a borrow and is set to the inverse of the resulting binary carry.

Addressing Modes: Immediate
Extended
Direct
Indexed

SWI**Software Interrupt****SWI****Source Form:** SWI**Operation:** Set E (entire state will be saved)
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow DPR$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCB$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCA$
 $SP' \leftarrow SP - 1, (SP) \leftarrow CCR$
Set I, F (mask interrupts)
 $PC' \leftarrow (FFFA):(FFFF)$ **Condition Codes:** Not affected.**Description:** All of the processor registers are pushed onto the hardware stack (with the exception of the hardware stack pointer itself), and control is transferred through the software interrupt vector. Both the normal and fast interrupts are masked (disabled).**Addressing Mode:** Inherent

SWI2

Software Interrupt 2

SWI2

Source Form: SWI2

Operation: Set E (entire state saved)
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow DPR$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCB$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCA$
 $SP' \leftarrow SP - 1, (SP) \leftarrow CCR$
 $PC' \leftarrow (FFF4):(FFF5)$

Condition Codes: Not affected.

Description: All of the processor registers are pushed onto the hardware stack (with the exception of the hardware stack pointer itself), and control is transferred through the software interrupt 2 vector. This interrupt is available to the end user and must not be used in packaged software. This interrupt does not mask (disable) the normal and fast interrupts.

Addressing Mode: Inherent

SWI3

Software Interrupt 3

SWI3

Source Form: SWI 3

Operation: Set E (entire state will be saved)
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow DPR$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCB$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCA$
 $SP' \leftarrow SP - 1, (SP) \leftarrow CCR$
 $PC' \leftarrow (FFF2):(FFF3)$

Condition Codes: Not affected.

Description: All of the processor registers are pushed onto the hardware stack (with the exception of the hardware stack pointer itself), and control is transferred through the software interrupt 3 vector. This interrupt does not mask (disable) the normal and fast interrupts.

Addressing Mode: Inherent

SYNC

Synchronize to External Event

SYNC

Source Form: SYNC

Operation: Stop processing instructions

Condition Codes: Not affected.

Description: When a SYNC instruction is executed, the processor enters a synchronizing state, stops processing instructions, and waits for an interrupt. When an interrupt occurs, the synchronizing state is cleared and processing continues. If the interrupt is enabled, and it lasts three cycles or more, the processor will perform the interrupt routine. If the interrupt is masked or is shorter than three cycles, the processor simply continues to the next instruction. While in the synchronizing state, the address and data buses are in the high-impedance state.

This instruction provides software synchronization with a hardware process. Consider the following example for high-speed acquisition of data:

FAST	SYNC Interrupt!	WAIT FOR DATA
LDA	DISC	DATA FROM DISC AND CLEAR INTERRUPT
STA	,X+	PUT IN BUFFER
DECB		COUNT IT, DONE?
BNE	FAST	GO AGAIN IF NOT.

The synchronizing state is cleared by any interrupt. Of course, enabled interrupts at this point may destroy the data transfer and, as such, should represent only emergency conditions.

The same connection used for interrupt-driven I/O service may also be used for high-speed data transfers by setting the interrupt mask and using the SYNC instruction as the above example demonstrates.

Addressing Mode: Inherent

TFR

Transfer Register to Register

TFR

Source Form: TFR R1, R2

Operation: R1 → R2

Condition Code: Not affected unless R2 is the condition code register.

Description: Transfers data between two designated registers. Bits 7-4 of the postbyte define the source register, while bits 3-0 define the destination register, as follows:

0000 = A:B	1000 = A
0001 = X	1001 = B
0010 = Y	1010 = CCR
0011 = US	1011 = DPR
0100 = SP	1100 = Undefined
0101 = PC	1101 = Undefined
0110 = Undefined	1110 = Undefined
0111 = Undefined	1111 = Undefined

Only like size registers may be transferred. (8-bit to 8-bit, or 16-bit to 16-bit.)

Addressing Mode: Immediate

TST

Test

TST

Source Forms: TST Q; TSTA; TSTB

Operation: TEMP ← M – 0

Condition Codes: H — Not affected.

N — Set if the result is negative; cleared otherwise.

Z — Set if the result is zero; cleared otherwise.

V — Always cleared.

C — Not affected.

Description: Set the N (negative) and Z (zero) bits according to the contents of memory location M, and clear the V (overflow) bit. The TST instruction provides only minimum information when testing unsigned values; since no unsigned value is less than zero, BLO and BLS have no utility. While BHI could be used after TST, it provides exactly the same control as BNE, which is preferred. The signed branches are available.

Addressing Modes: Inherent
Extended
Direct
Indexed

Comments: The MC6800 processor clears the C (carry) bit.

FIRQ

Fast Interrupt Request (Hardware Interrupt)

FIRQ

Operation: IFF F bit clear, then: $SP' \leftarrow SP - 1$, $(SP) \leftarrow PCL$
 $SP' \leftarrow SP - 1$, $(SP) \leftarrow PCH$
Clear E (subset state is saved)
 $SP' \leftarrow SP - 1$, $(SP) \leftarrow CCR$
Set F, I (mask further interrupts)
 $PC' \leftarrow (FFF6):(FFF7)$

Condition Codes: Not affected.

Description: A FIRQ (fast interrupt request) with the F (fast interrupt request mask) bit clear causes this interrupt sequence to occur at the end of the current instruction. The program counter and condition code register are pushed onto the hardware stack. Program control is transferred through the fast interrupt request vector. An RTI (return from interrupt) instruction returns the processor to the original task. It is possible to enter the fast interrupt request routine with the entire machine state saved if the fast interrupt request occurs after a clear and wait for interrupt instruction. A normal interrupt request has lower priority than the fast interrupt request and is prevented from interrupting the fast interrupt request routine by automatic setting of the I (interrupt request mask) bit. This mask bit could then be reset during the interrupt routine if priority was not desired. The fast interrupt request allows operations on memory, TST, INC, DEC, etc. instructions without the overhead of saving the entire machine state on the stack.

Addressing Mode: Inherent

IRQ**Interrupt Request (Hardware Interrupt)****IRQ**

Operation: IFF I bit clear, then:

$SP' \leftarrow SP - 1, (SP) \leftarrow PCL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow DPR$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCB$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCA$
Set E (entire state saved)
 $SP' \leftarrow SP - 1, (SP) \leftarrow CCR$
Set I (mask further IRQ interrupts)
 $PC' \leftarrow (FFF8):(FFF9)$

Condition Codes: Not affected.

Description: If the I (interrupt request mask) bit is clear, a low level on the IRQ input causes this interrupt sequence to occur at the end of the current instruction. Control is returned to the interrupted program using a RTI (return from interrupt) instruction. A FIRQ (fast interrupt request) may interrupt a normal IRQ (interrupt request) routine and be recognized anytime after the interrupt vector is taken.

Addressing Mode: Inherent

NMI

Non-Maskable Interrupt (Hardware Interrupt)

NMI

Operation:

$SP' \leftarrow SP - 1, (SP) \leftarrow PCL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow PCH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow USH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IYH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXL$
 $SP' \leftarrow SP - 1, (SP) \leftarrow IXH$
 $SP' \leftarrow SP - 1, (SP) \leftarrow DPR$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCB$
 $SP' \leftarrow SP - 1, (SP) \leftarrow ACCA$
Set E (entire state save)
 $SP' \leftarrow SP - 1, (SP) \leftarrow CCR$
Set I, F (mask interrupts)
 $PC' \leftarrow (FFFC):(FFFD)$

Condition Codes: Not affected.

Description: A negative edge on the NMI (non-maskable interrupt) input causes all of the processor's registers (except the hardware stack pointer) to be pushed onto the hardware stack, starting at the end of the current instruction. Program control is transferred through the NMI vector. Successive negative edges on the NMI input will cause successive NMI operations. Non-maskable interrupt operation can be internally blocked by a RESET operation and any non-maskable interrupt that occurs will be latched. If this happens, the non-maskable interrupt operation will occur after the first load into the stack pointer (LDS; TFR r,s; EXG r,s; etc.) after RESET.

Addressing Mode: Inherent

RESTART

Restart (Hardware Interrupt)

RESTART

Operation: $CCR' \leftarrow X1X1XXXX$
 $DPR' \leftarrow 0016$
 $PC' \leftarrow (FFFE):(FFFF)$

Condition Codes: Not affected.

Description: The processor is initialized (required after power-on) to start program execution. The starting address is fetched from the restart vector.

Addressing Mode: Extended Indirect

APPENDIX B

ASSIST09 MONITOR PROGRAM

B.1 GENERAL DESCRIPTION

The M6809 is a high-performance microprocessor which supports modern programming techniques such as position-independent, reentrancy, and modular programming. For a software monitor to take advantage of such capabilities demands a more refined and sophisticated user interface than that provided by previous monitors. ASSIST09 is a monitor which supports the advanced features that the M6809 makes possible. ASSIST09 features include the following:

- Coded in a position (address) independent manner. Will execute anywhere in the 64K address space.
- Multiple means available for installing user modifications and extensions.
- Full complement of commands for program development including breakpoint and trace.
- Sophisticated monitor calls for completely address-independent user program services.
- RAM work area is located relative to the ASSIST09 ROM, not at a fixed address as with other monitors.
- Easily adapted to real-time environments.
- Hooks for user command tables, I/O handlers, and default specifications.
- A complete user interface with services normally only seen in full disk operating systems.

The concise instruction set of the M6809 allows all of these functions and more to be contained in only 2048 bytes.

The ASSIST09 monitor is easily adapted to run under control of a real-time operating system. A special function is available which allows voluntary time-slicing, as well as forced time-slicing upon the use of several service routines by a user program.

B.2 IMPLEMENTATION REQUIREMENTS

Since ASSIST09 was coded in an address-independent manner, it will properly execute anywhere in the 64K address space of the M6809. However, an assumption must be made regarding the location of a work area needed to hold miscellaneous variables and the default stack location. This work area is called the page work area and it is addressed within ASSIST09 by use of the direct page register. It is located relative to the start of the

ASSIST09 ROM by an offset of -1900 hexadecimal. Assuming ASSIST09 resides at the top of the memory address space for direct control of the hardware interrupt vectors, the memory map would appear as shown in Figure B-1.

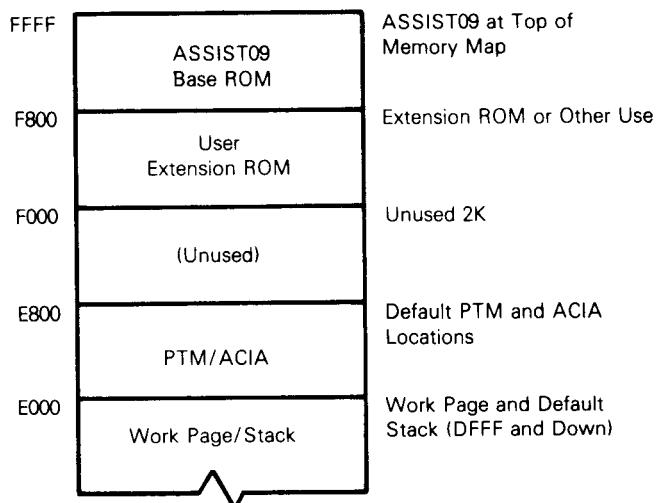


Figure B-1. Memory Map

If F800 is not the start of the monitor ROM the addresses would change, but the relative locations would remain the same except for the programmable timer module (PTM) and asynchronous communications interface adapter (ACIA) default addresses which are fixed.

The default console input/output handlers access an ACIA located at E008. For trace commands, a PTM with default address E000 is used to force an NMI so that single instructions may be executed. These default addresses may easily be changed using one of several methods. The console I/O handlers may also be replaced by user routines. The PTM is initialized during the MONITR service call (see Paragraph B.9 SERVICES) to fireup the monitor unless its default address has been changed to zero, in which case no PTM references will occur.

B.3 INTERRUPT CONTROL

Upon reset, a vector table is created which contains, among other things, default interrupt vector handler appendage addresses. These routines may easily be replaced by user appendages with the vector swap service described later. The default actions taken by the appendages are as follows:

RESET — Build the ASSIST09 vector table and setup monitor defaults, then invoke the monitor startup routine.

SWI — Request a service from ASSIST09.

FIRQ — An immediate RTI is done.

SWI2, SWI3, IRQ, Reserved, NMI — Force a breakpoint and enter the command processor.

The use of IRQ is recommended as an abort function during program debugging sessions, as breakpoints and other ASSIST09 defaults are reinitialized upon RESET. Only the primary software interrupt instruction (SWI) is used, not the SWI2 or SWI3. This avoids page fault problems which would otherwise occur with a memory management unit as the SWI2 and SWI3 instructions do not disable interrupts.

Counter number one of the PTM is used to cause an NMI interrupt for the trace and breakpoint commands. At RESET the control register for timer one is initialized for tracing purposes. If no tracing or breakpointing is done then the entire PTM is available to the user. Otherwise, only counters two and three are available. Although control register two must be used to initialize control register one, ASSIST09 returns control register two to the same value it has after a RESET occurs. Therefore, the only condition imposed on a user program is that if the "operate/preset" bit in control register one must be turned on, \$A7 should be stored, \$A6 should be stored if it must be turned off.

B.4 INITIALIZATION

During ASSIST09 execution, a vector table is used to address certain service routines and default values. This table is generated to provide easily changed control information for user modifications. The first byte of the ASSIST09 ROM contains the start of a subroutine which initializes the vector table along with setting up certain default values before returning to the caller.

If the ASSIST09 RESET vector receives control, it does three things:

1. Assigns a default stack in the work space,
2. Calls the aforementioned subroutine to initialize the vector table, and
3. Fires up the ASSIST09 monitor proper with a MONITR SWI service request.

However, a user routine can perform the same functions with a bonus. After calling the vector initialization subroutine, it may examine or alter any of the vector table values before starting normal ASSIST09 processing. Thus, a user routine may "bootstrap" ASSIST09 and alter the default standard values.

Another method of inserting user modifications is to have a user routine reside at an extension ROM location 2K below the start of the ASSIST09 ROM. The vector table initialization routine mentioned above, looks for a "BRA*" flag (\$20FE) at this address, and if found calls the location following the flag as a subroutine with the U register pointing to the vector table. Since this is done after vector table initialization, any or all defaults may be altered at this time. A big advantage to using this method is that the modifications are "automatic" in that upon a RESET condition the changes are made without overt action required such as the execution of a memory change command.

No special stack is used during ASSIST09 processing. This means that the stack pointer must be valid at all interruptable times and should contain enough room for the stacking of at least 21 bytes of information. The stack in use during the initial MONITR service call to start up ASSIST09 processing becomes the "official" stack. If any later stack validity checks occur, this same stack will be re-based before entering the command handler.

ASSIST09 uses a work area which is addressed at an offset from the start of the ASSIST09 ROM. The offset value is -1900 hexadecimal. This points to the base page used during monitor execution and contains the vector table as well as the start of the default stack. If the default stack is used and it exceeds 81 bytes in size, then contiguous RAM must exist below this base work page for proper extension of the stack.

B5. INPUT/OUTPUT CONTROL

Output generated by use of the ASSIST09 services may be halted by pressing any key, causing a 'FREEZE' mode to be entered. The next keyboard entry will release this condition allowing normal output to continue. Commands which generate large amounts of output may be aborted by entering CANCEL (CONTROL-X). User programs may also monitor for CANCEL along with the 'FREEZE' condition even when not performing console I/O (PAUSE service).

B.6 COMMAND FORMAT

There are three possible formats for a command:

<Command> CR

<Command> <Expression1> CR

<Command> <Expression1> <Expression2> CR

The space character is used as the delimiter between the command and all arguments. Two special quick commands need no carriage return, “.” and “/”. To re-enter a command once a mistake is made, type the CANCEL (CONTROL-X) key.

Each “expression” above consists of one or more values separated by an operator. Values can be hex strings, the letters “P”, “M”, and “W”, or the result of a function. Each hexadecimal string is converted internally to a 16-bit binary number. The letter “P” stands for the current program counter, “M” for the last memory examine/change address, and “W” for the window value. The window value is set by using the WINDOW command.

One function exists and it is the INDIRECT function. The character “@” following a value replaces that value with the 16-bit number obtained by using that value as an address.

Two operators are allowed, “+” and “-” which cause addition and subtraction. Values are operated on in a left-to-right order.

Examples:

480 — hexadecimal 480

W + 3 — value of window plus three

P-200 — current program counter minus 200 hexadecimal

M - W — current memory pointer minus window value

100@ — value of word addressed by the two bytes at 100 hexadecimal

P + 1@ — value addressed by the word located one byte up from the current program counter

B.7 COMMAND LIST

Table B-1 lists the commands available in the ASSIST09 monitor.

Table B-1. Command List

Command Name	Description	Command Entry
Breakpoint	Set, clear, display, or delete breakpoints	B
Call	Call program as subroutine	C
Display	Display memory block in hex and ASCII	D
Encode	Return indexed postbyte value	E
Go	Start or resume program execution	G
Load	Load memory from tape	L
Memory	Examine or alter memory	M
	Memory change or examine last referenced	/
	Memory change or examine	hex/
Null	Set new character and new line padding	N
Offset	Compute branch offsets	O
Punch	Punch memory on tape	P
Registers	Display or alter registers	R
Stlevel	Alter stack trace level value	S
Trace	Trace number of instructions	T
	Trace one instruction	.
Verify	Verify tape to memory load	V
Window	Set a window value	W

B.8 COMMANDS

Each of the commands are explained on the following pages. They are arranged in alphabetical order by the command name used in the command list. The command name appears at each margin and in slightly larger type for easy reference.

BREAKPOINT

BREAKPOINT

Format: Breakpoint
Breakpoint –
Breakpoint <Address>
Breakpoint – <Address>

Operation: Set or change the breakpoint table. The first format displays all breakpoints. The second clears the breakpoint table. The third enters an address into the table. The fourth deletes an address from the table. At reset, all breakpoints are deleted. Only instructions in RAM may be breakpointed.

CALL

CALL

Format: Call
Call <Address>

Operation: Call and execute a user routine as a subroutine. The current program counter will be used unless the address is specified. The user routine should eventually terminate with a “RTS” instruction. When this occurs, a breakpoint will ensue and the program counter will point into the monitor.

DISPLAY

DISPLAY

Format: Display <From>
 Display <From> <Length>
 Display <From> <To>

Operation: Display contents of memory in hexadecimal and ASCII characters. The second argument, when entered, is taken to be a length if it is less than the first, otherwise it is the ending address. A default length of 16 decimal is assumed for the first format. The addresses are adjusted to include all bytes within the surrounding modulo 16 address byte boundary. The CANCEL (CONTROL-X) key may be entered to abort the display. Care must be exercised when the last 15 bytes of memory are to be displayed. The <Length> option should always be used in this case to assure proper termination: D FFE0 40

Examples:

D M 10 — Display 16 bytes surrounding the last memory location examined.
D E000 F000 — Display memory from E000 to F000 hex.

ENCODE

ENCODE

Format: Encode <Indexed operand>

Operation: The encode command will return the indexing instruction mode postbyte value from the entered assembler-like syntax operand. This is useful when hand coding instructions. The letter "H" is used to indicate the number of hex digits needed in the expression as shown in the following examples:

E ,Y — Return zero offset to Y register postbyte.
E [HHHH,PCR] — Return two byte PCR offset using indirection.
E [S + +] — Return autoincrement S by two indirect.
E H,X — Return 5-bit offset from X.

Note that one "H" specifies a 5-bit offset, and that the result given will have zeros in the offset value position. This command does not detect all incorrectly specified syntax or illegal indexing modes.

GO

Format: Go
Go <Address>

Operation: Execute starting from the address given. The first format will continue from the current program counter setting. If it is a breakpoint no break will be taken. This allows continuation from a breakpoint. The second format will breakpoint if the address specified is in the breakpoint list.

LOAD

Format: Load
Load <Offset>

Operation: Load a tape file created using the S1-S9 format. The offset option, if used, is added to the address on the tape to specify the actual load address. All offsets are positive, but wrap around memory modulo 64K. Depending on the equipment involved, after the load is complete a few spurious characters may still be sent by the input device and interpreted as command characters. If this happens, a CANCEL (CONTROL-X) should be entered to cause such characters to be ignored. If the load was not successful a "?" is displayed.

GO

LOAD

MEMORY

MEMORY

Format: MEMORY <Address>/
 <Address>/
 /

Operation: Initiate the memory examine/change function. The second format will not accept an expression for the address, only a hex string. The third format defaults to the address displayed during the last memory change/examine function. (The same value is obtained in expressions by use of the letter "M".) After activation, the following actions may be taken until a carriage return is entered:

<Expr>	Replaces the byte with the specified value. The value may be an expression.
SPACE	Go to next address and print the byte value.
,	(Comma) Go to next address without printing the byte value.
LF	(Line feed) Go to next address and print it along with the byte value on the next line.
^	(Circumflex or Up arrow) Go the previous address and print it along with the byte value on the next line.
/	Print the current address with the byte value on the next line.
CR	(Carriage return) Terminate the command.
'<Text>'	Replace succeeding bytes with ASCII characters until the second apostrophe is entered.

If a change attempt fails (i.e., the location is not valid RAM) then a question mark will appear and the next location displayed.

NULL

NULL

Format: Null <Specification>

Operation: Set the new line and character padding count values. The expression value is treated as two values. The upper two hex represent the character pad count, and the lower two the new line pad count (triggered by a carriage return). An expression of less than three hex digits will set the character pad count to zero. The values must range from zero to 7F hexadecimal (127 decimal).

Example:

N 3 — Set the character count to zero and new line count to three.

N 207 — Set character padding count to two and new line count to seven.

Settings for TI Silent 700 terminals are:

Baud	Setting
100	0
300	4
1200	317
2400	72F

OFFSET

OFFSET

Format: Offset <Offset addr> <To instruction>

Operation: Print the one and two byte offsets needed to perform a branch from the first expression to the instruction. Thus, offsets for branches as well as indexed mode instructions which use offsets may be obtained. If only a four byte value is printed, then a short branch count cannot be done between the two addresses.

Example:

0 P+2 A000 — Compute offsets needed from the current program counter plus two to A000.

PUNCH

PUNCH

Format: Punch <From> <To>

Operation: Punch or record formatted binary object tape in S1-S9 (MIKBUG) format.

REGISTER

REGISTER

Format: Register

Operation: Print the register set and prompt for a change. At each prompt the following may be entered.

- | | |
|--------------|---|
| SPACE | Skip to the next register prompt |
| <Expr> SPACE | Replace with the specified value and prompt for the next register. |
| <Expr> CR | (carriage return) Replace with the specified value and terminate the command. |
| CR | Terminate the command. |

STLEVEL

STLEVEL

Format: Stlevel
Stlevel <Address>

Operation: Set the stack trace level for inhibiting tracing information. As long as the stack is at or above the stack level address, the trace display will continue. However, when lower than the address it is inhibited. This allows tracing of a routine without including all subroutine and lower level calls in the trace information. Note that tracing through a ASSIST09 "SWI" service request may also temporarily suppress trace output as explained in the description of the trace command. The first format sets the stack trace level to the current program stack value.

TRACE

TRACE

Format: Trace <Count>
. (period)

Operation: Trace the specified number of instructions. At each trace, the opcode just executed will be shown along with the register set. The program counter in the register display points to the NEXT instruction to be executed. A CANCEL (CONTROL-X) will prematurely halt tracing. The second format (period) will cause a single trace to occur. Breakpoints have no effect during the trace. Selected portions of a trace may be disabled using the STLEVEL command. Instructions in ROM and RAM may be traced, whereas breakpoints may be done only in RAM. When tracing through a ASSIST09 service request, the trace display will be suppressed starting two instructions into the monitor until shortly before control is returned to the user program. This is done to avoid an inordinate amount of displaying because ASSIST09, at times, performs a sizeable amount of processing to provide the requested services.

VERIFY

VERIFY

Format: Verify
Verify <Offset>

Operation: Verify or compare the contents of memory to the tape file. This command has the same format and operation as a LOAD command except the file is compared to memory. If the verify fails for any reason a "?" is displayed.

WINDOW

WINDOW

Format: Window <Value>

Operation: Set the window to a value. This value may be referred to when entering expressions by use of the letter "W". The window may be set to any 16-bit value.

B.9 SERVICES

The following describes services provided by the ASSIST09 monitor. These services are invoked by using the "SWI" instruction followed by a one byte function code. All services are designed to allow complete address independence both in invocation and operation. Unless specified otherwise, all registers are transparent over the "SWI" call. In the following descriptions, the terms "input handler" and "output handler" are used to refer to appendage routines which may be replaced by the user. The default routines perform standard I/O through an ACIA for console operations to a terminal. The ASCII CANCEL code can be entered on most terminals by depressing the CONTROL and X keys simultaneously. A list of services is given in Table B-2.

Table B-2. Services

Service	Entry	Code	Description
Obtain input character	INCHP	0	Obtain the input character in register A from the input handler
Output a character	OUTCH	1	Send the character in the register A to the output handler
Send string	PDATA1	2	Send a string of characters to the output handler
Send new line and string	PDATA	3	Send a carriage return, line feed, and string of characters to the output handler
Convert byte to hex	OUT2HS	4	Display the byte pointed to by the X register in hex
Convert word to hex	OUT4HS	5	Display the word pointed to by the X register in hex
Output to next line	PCRLF	6	Send a carriage return and line feed to the output handler
Send space	SPACE	7	Send a blank to the output handler
Fireup ASSIST09	MONITR	8	Enter the ASSIST09 monitor
Vector swap	VCTRSW	9	Examine or exchange a vector table entry
User breakpoint	BRKPT	10	Display registers and enter the command handler
Program break and check	PAUSE	11	Stop processing and check for a freeze or cancel condition

BRKPT

User Breakpoint

BRKPT

Code: 10

Arguments: None

Result: A disabled breakpoint is taken. The registers are displayed and the command handler of ASSIST09 is entered.

Description: Establishes user breakpoints. Both SWI2 and SWI3 default appendages cause a breakpoint as well, but do not set the I and F mask bits. However, since they may both be replaced by user routines the breakpoint service always ensures breakpoint availability. These user breakpoints have nothing to do with system breakpoints which are handled differently by the ASSIST09 monitor.

Example:	BRKPT	EQU 10	INPUT CODE FOR BRKPT
		SWI	REQUEST SERVICE
		FCB BRKPT	FUNCTION CODE BYTE

INCHP

Obtain Input Character

INCHP

Code: 0

Arguments: None

Result: Register A contains a character obtained from the input handler.

Description: Control is not returned until a valid input character is received from the input handler. The input character will have its parity bit (bit 7) stripped and forced to a zero. All NULL (\$00) and RUBOUT (\$7F) characters are ignored and not returned to the caller. The ECHO flag, which may be changed by the vector SWAP service, determines whether or not the input character is echoed to the output handler (full duplex operation). The default at reset is to echo input. When a carriage return (\$0D) is received, line feed (\$A0) is automatically sent back to the output handler.

Example:	INCHNP	EQU 0	INPUT CODE FOR INCHP
	SWI		PERFORM SERVICE CALL
	FCB	INCHNP	FUNCTION FOR INCHNP

A REGISTER NOW CONTAINS NEXT CHARACTER

MONITR

Startup ASSIST09

MONITR

Code: 8

Arguments: S→ Stack to become the “official” stack
DP→ Direct page default for executed user programs
A=0 Call input and output console initialization handlers and give the
“ASSIST09” startup message
A#0 Go directly to the command handler

Result: ASSIST09 is entered and the command handler given control

Description: The purpose for this function is to enter ASSIST09, either after a system reset, or when a user program desires to terminate. Control is not returned unless a “GO” or “CALL” command is done without altering the program counter. ASSIST09 runs on the passed stack, and if a stack error is detected during user program execution this is the stack that is rebased. The direct page register value in use remains the default for user program execution.

The ASSIST09 restart vector routine uses this function to startup monitor processing after calling the vector build subroutine as explained in INITIALIZATION.

If indicated by the A register, the input and output initialization handlers are called followed by the sending of the string “ASSIST09” to the output handler. The programmable timer (PTM) is initialized, if its address is not zero, such that register 1 can be used for causing an NMI during trace commands. The command handler is then entered to perform the command request prompt.

Example:	MONITR EQU 8	INPUT CODE FOR MONITR
	LOOP CLRA	PREPARE ZERO PAGE REGISTER AND INITIALIZATION PARAMETER
	* TFR A,DP	SET DEFAULT PAGE VALUE
	LEAS STACK, PCR	SETUP DEFAULT STACK VALUE
	SWI	REQUEST SERVICE
	FCB MONITR	FUNCTION CODE BYTE
	BRA LOOP	REENTER IF FALLOUT OCCURS

OUTCH

Output a Character

OUTCH

Code: 1

Arguments: Register A contains the byte to transmit.

Result: The character is sent to the output handler

The character is set as follows ONLY if a LINEFEED was the character to transmit:

CC=0 if normal output occurred.

CC=1 if CANCEL was entered during output.

Description: If a FREEZE Occurs (any input character is received) then control is not returned to the user routine until the condition is released. The FREEZE condition is checked for only when a linefeed is being sent. Padding null characters (\$00) may be sent following the outputted character depending on the current setting of the NULLS command. For DLE (Data Link Escape), character nulls are never sent. Otherwise, carriage returns (\$0D) receive the new line count of nulls, all other characters the character count of nulls.

Example: OUTCH EQU 1 INPUT CODE FOR OUTCH

LDA #'0	LOAD CHARACTER "0"
SWI	SEND OUT WITH MONITOR CODE
FCB OUTCH	SERVICE CODE BYTE

OUT2HS

Convert Byte to Hex

OUT2HS

Code: 4

Arguments: Register X points to a byte to display in hex.

Result: The byte is converted to two hex digits and sent to the output handler followed by a blank.

Example: OUT2HS EQU 4 INPUT CODE FOR OUT2HS

LEAX DATA, PCR	POINT TO 'DATA' TO DECODE
SWI	REQUEST SERVICE
FCB OUT2HS	SERVICE CODE BYTE

OUT4HS

Convert Word to Hex

OUT4HS

Code: 5

Arguments: Register X points to a word (two bytes) to display in hex.

Result: The word is converted to four hex digits and sent to the output handler followed by a blank.

Example:	OUT4HS EQU 5	INPUT CODE FOR OUT4HS
	LEAX DATA, PCR	LOAD 'DATA' ADDRESS TO DECODE
	SWI	REQUEST ASSIST09 SERVICE
	FCB OUT4HS	SERVICE CODE BYTE

PAUSE

Program Break and Check

PAUSE

Code: 11

Arguments: None

Result: CC=0 For a normal return.
CC=1 If a CANCEL was entered during the interim.

Description: The PAUSE service should be used whenever a significant amount of processing is done by a program without any external interaction (such as console I/O). Another use of the PAUSE service is for the monitoring of FREEZE or CANCEL requests from the input handler. This allows multi-tasking operating systems to receive control and possibly re-dispatch other programs in a timeslice-like fashion. Testing for FREEZE and CANCEL conditions is performed before return. Return may be after other tasks have had a chance to execute, or after a FREEZE condition is lifted. In a one task system, return is always immediate unless a FREEZE occurs.

PCRLF

Output to Next Line

PCRLF

Code: 6

Arguments: None

Result: A carriage return and line feed are sent to the output handler.
C = 1 if normal output occurred.
C = 1 if CONTROL-X was entered during output.

Description: If a FREEZE occurs (any input character is received), then control is not returned to the user routine until the condition is released. The string is completely sent regardless of any FREEZE or CANCEL events occurring. Padding characters may be sent as described under the OUTCH service.

Example:	PCRLF	EQU 6	INPUT CODE PCRLF
		SWI	REQUEST SERVICE
		FCB PCRLF	SERVICE CODE BYTE

PDATA

Send New Line and String

PDATA

Code: 3

Arguments: Register X points to an output string terminated with an ASCII EOT (\$04).

Result: The string is sent to the output handler following a carriage return and line feed.
CC = 0 if normal output occurred.
CC = 1 if CONTROL-X was entered during output.

Description: The output string may contain embedded carriage returns and line feeds thus allowing several lines of data to be sent with one function call. If a FREEZE occurs (any input character is received), then control is not returned to the user routine until the condition is released. The string is completely sent regardless of any FREEZE or CANCEL events occurring. Padding characters may be sent as described by the OUTCH function.

PDATA

Send New Line and String (Continued)

PDATA

Example: PDATA EQU 3 INPUT CODE FOR PDATA

MSGOUT FCC 'THIS IS A MULTIPLE LINE MESSAGE.'
FCB \$0A, \$0D LINE FEED, CARRIAGE RETURN
FCC 'THIS IS THE SECOND LINE.'
FCB \$04 STRING TERMINATOR

LEAX MSGOUT, PCR LOAD MESSAGE ADDRESS
SWI REQUEST A SERVICE
FCB PDATA SERVICE CODE BYTE

PDATA1

Send String

PDATA1

Code: 2

Arguments: Register X points to an output string terminated with an ASCII EOT (\$04).

Result: The string is sent to the output handler.

CC=0 if normal output occurred.

CC=1 if CONTROL-X was entered during output.

Description: The output string may contain embedded carriage returns and line feeds thus allowing several lines of data to be sent with one function call. If a FREEZE occurs (any input character is received), then control is not returned to the user routine until the condition is released. The string is completely sent regardless of any FREEZE or CANCEL events occurring. Padding characters may be sent as described by the OUTCH function.

Example: PDATA EQU 2 INPUT CODE FOR PDATA1

MSG FCC 'THIS IS AN OUTPUT STRING'
FCB \$04 STRING TERMINATOR

LEAX MSG, PCR LOAD 'MSG' STRING ADDRESS
SWI REQUEST A SERVICE
FCB PDATA1 SERVICE CODE BYTE

SPACE

Single Space Output

SPACE

Code: 7

Arguments: None

Result: A space is sent to the output handler.

Description: Padding characters may be sent as described under the OUTCH service.

Example: SPACE EQU 7 INPUT CODE SPACE
SWI REQUEST ASSIST09 SERVICE
FCB SPACE SERVICE CODE BYTE

VCTRSW

Vector Swap

VCTRSW

Code: 9

Arguments: Register A contains the vector swap input code.
Register X contains zero or a replacement value.

Result: Register X contains the previous value for the vector.

Description: The vector swap service examines/alters a word entry in the ASSIST09 vector table. This table contains pointers and default values used during monitor processing. The entry is replaced with the value contained in the X register unless it is zero. The codes available are listed in Table B-3.

Example: VCTRSW EQU 9 INPUT CODE VCTRSW
.IRQ EQU 12 IRQ APPENDAGE SWAP FUNCTION
CODE

LEAX MYIRQH,PCR LOAD NEW IRQ HANDLER ADDRESS
LDA #.IRQ LOAD SUBCODE FOR VECTOR SWAP
SWI REQUEST SERVICE
FCB VCTRSW SERVICE CODE BYTE
X NOW HAS THE PREVIOUS APPENDAGE ADDRESS

B.10 VECTOR SWAP SERVICE

The vector swap service allows user modifications of the vector table to be easily installed. Each vector handler, including the one for SWI, performs a validity check on the stack before any other processing. If the stack is not pointing to valid RAM, it is reset to the initial value passed to the MONITR request which fired-up ASSIST09 after RESET. Also, the current register set is printed following a “?” (question mark) and then the command handler is entered. A list of each entry in the vector table is given in Table B-3.

Table B-3. Vector Table Entries

Entry	Code	Description
.AVTBL	0	Returns address of vector table
.CMDL1	2	Primary command list
.RSVD	4	Reserved MC6809 interrupt vector appendage
.SWI3	6	Software interrupt 3 interrupt vector appendage
.SWI2	8	Software interrupt 2 interrupt vector appendage
.FIRQ	10	Fast interrupt request vector appendage
.IRQ	12	Interrupt request vector appendage
.SWI	14	Software interrupt vector appendage
.NMI	16	Non-maskable interrupt vector appendage
.RESET	18	Reset interrupt vector appendage
.CION	20	Input console initialization routine
.CIDTA	22	Input data byte from console routine
.CIOFF	24	Input console shutdown routine
.COON	26	Output console initialization routine
.CODTA	28	Output/data byte to console routine
.COOFF	30	Output console shutdown routine
.HSDTA	32	High speed display handler routine
.BSON	34	Punch/load initialization routine
.BSDTA	36	Punch/load handler routine
.BSOFF	38	Punch/load shutdown routine
.PAUSE	40	Processing pause routine
.CMDL2	44	Secondary command list
.ACIA	46	Address of ACIA
.PAD	48	Character and new line pad counts
.ECHO	50	Echo flag
.PTM	52	Programmable timer module address

The following pages describe the purpose of each entry and the requirements which must be met for a user replaceable value or routine to be successfully substituted.

.ACIA

ACIA Address

.ACIA

Code: 46

Description: This entry contains the address of the ACIA used by the default console input and output device handlers. Standard ASSIST09 initialization sets this value to hexadecimal E008. If this must be altered, then it must be done before the MONITR startup service is invoked, since that service calls the .COON and .COIN input and output device initialization routines which initialize the ACIA pointed to by this vector slot.

.AVTBL

Return Address of Vector Table

.AVTBL

Code: 0

Description: The address of the vector table is returned with this code. This allows mass changes to the table without individual calls to the vector swap service. The code values are identical to the offsets in the vector table. This entry should never be changed, only examined.

.BSDTA

Punch/Load Handler Routine

.BSDTA

Code: 36

Description: This entry contains the address of a routine which performs punch, load, and verify operations. The .BSON routine is always executed before the routine is given control. This routine is given the same parameter list documented for .BSON. The default handler uses the .CODTA routine to punch or the .CIDTA routine to read data in S1/S9 (MIKBUG) format. The function code byte must be examined to determine the type request being handled.

A return code must be given which reflects the final processing disposition:

Z = 1 Successful completion

or

Z = 0 Unsuccessful completion.

The .BSOFF routine will be called after this routine is completed.

.BSOFF

Punch/Load Shutdown Routine

.BSOFF

Code: 38

Description: This entry points to a subroutine which is designated to terminate device processing for the punch, load, and verify handler .BSDTA. The stack contains a parameter list as documented for the .BSON entry. The default ASSIST09 routine issues DC4 (\$14 or stop) and DC3 (\$13 or x-off) followed by a one second delay to give the reader/punch time to stop. Also, an internally used flag by the INCHP service routine is cleared to reverse the effect caused by its setting in the .BSON handler. See that description for an explanation of the proper use of this flag.

.BSON

Punch/Load Initialization Routine

.BSON

Code: 34

Description: This entry points to a subroutine with the assigned task of turning on the device used for punch, load, and verify processing. The stack contains a parameter list describing which function is requested. The default routine sends an ASCII "reader on" or "punch on" code of DC1 (\$11) or DC2 (\$12) respectively to the output handler (.CODTA). A flag is also set which disables test for FREEZE conditions during INCHNP processing. This is done so characters are not lost by being interpreted as FREEZE mode indicators. If a user replacement routine also uses the INCHNP service, then it also should set this same byte non-zero and clear it in the .BSOFF routine. The ASSIST09 source listing should be consulted for the location of this byte.

The stack is setup as follows:

S + 6 = Code byte, VERIFY (-1), PUNCH (0), LOAD (1)

S + 4 = Start address for punch only

S + 2 = End address for punch, or offset for READ/LOAD

S + 0 = Return address

.CIDTA

Input Data Byte from Console Routine

.CIDTA

Code: 22

Description: This entry determines the console input handler appendage. The responsibility of this routine is to furnish the requested next input character in the A register, if available, and return with a condition code. The INCHP service routine calls this appendage to supply the next character. Also, a "FREEZE" mode routine calls at various times to test for a FREEZE condition or determine if the CANCEL key has been entered. Processing for this appendage must abide by the following conventions:

Input: PC → ASSIST09 work page

S → Return address

Output: C = 0, A = input character

C = 1 if no input character is yet available

Volatile Registers: U, B

The handler should always pass control back immediately even if no character is yet available. This enables other tasks to do productive work while input is unavailable. The default routine reads an ACIA as explained in Paragraph B.2 Implementation Requirements.

.CIOFF

Input Console Shutdown Routine

.CIOFF

Code: 24

Description: This entry points to a routine which is called to terminate input processing. It is not called by ASSIST09 at any time, but is included for consistency. The default routine merely does an "RTS". The environment is as follows:

Input: None
Output: Input device terminated
Volatile Registers: None

.CION

Input Console Initialization Routine

.CION

Code: 20

Description: This entry is called to initiate the input device. It is called once during the MONITR service which initializes the monitor so the command processor may obtain commands to process. The default handler resets the ACIA used for standard input and output and sets up the following default conditions: 8-bit word length, no parity checking, 2 stop bits, divide-by-16 counter ratio. The effect of an 8-bit word with no parity checking is to accept 7-bit ASCII and ignore the parity bit.

Input: .ACIA Memory address of the ACIA
Output: The output device is initialized
Volatile Registers: A, X

.CMDL1

Primary Command List

.CMDL1

Code: 2

Description: User supplied command tables may either substitute or replace the ASSIST09 standard tables. The command handler scans two lists, the primary table first followed by the secondary table. The primary table is pointed to by this entry and contains, as a default, the ASSIST09 command table. The secondary table defaults to a null list. A user may insert their own table into either position. If a user list is installed in the secondary table position, then the ASSIST09 list will be searched first. The default ASSIST09 list contains all one character command names. Thus, a user command "PRINT" would be matched if the letters "PR" are typed, but not just a "P" since the system command list would match first. A user may replace the primary system list if desired. A command is chosen on a first match basis comparing only the character(s) entered. This means that two or more commands may have the same initial characters and that if only that much is entered then the first one in the list(s) is chosen.

Each entry in the users command list must have the following format:

+0	FCB	L	Where "L" is the size of the entry including this byte
+1	FCC	'<string>'	Where "<string>" is the command name
+N	FDB	EP - *	Where "EP" represents the symbol defining the start of the command routine

The first byte is an entry length byte and is always three more than the length of the command string (one for the length itself plus two for the routine offset). The command string must contain only ASCII alphanumeric characters, no special characters. An offset to the start of the command routine is used instead of an absolute address so that position-independent programs may contain command tables. The end of the command table is a one byte flag. A -1 (\$FF) specifies that the secondary table is to be searched, or a -2 (\$FE) that command list searching is to be terminated. The table represented as the secondary command list must end with -2. The first list must end with a -1 if both lists are to be searched, or a -2 if only one list is to be used.

A command routine is entered with the following registers set:

DPR→	ASSIST09 page work area.
S→	A return address to the command processor.
Z=1	A carriage return terminated the command name.
Z=0	A space delimiter followed the command name.

.CMDL1

Primary Command List (Continued)

.CMDL1

A command routine is entered after the delimiter following the command name is typed in. This means that a carriage return may be the delimiter entered with the input device resting on the next line. For this reason the Z bit in the condition code is set so the command routine may determine the current position of the input device. The command routine should ensure that the console device is left on a new line before returning to the command handler.

.CMDL2

Secondary Command List

.CMDL2

Code: 44

Description: This entry points to the second list table. The default is a null list followed by a byte of -2. A complete explanation of the use for this entry is provided under the description of the .CMDL1 entry.

.CODTA

Output Data Byte to Console Routine

.CODTA

Code: 28

Description: The responsibility of this handler is to send the character in the A register to the output device. The default routine also follows with padding characters as explained in the description of the OUTCH service. If the output device is not ready to accept a character, then the "pause" subroutine should be called repeatedly while this condition lasts. The address of the pause routine is obtained from the .PAUSE entry in the vector table. The character counts for padding are obtained from the .PAD entry in the table. All ASSIST09 output is done with a call to this appendage. This includes punch processing as well. The default routine sends the character to an ACIA as explained in Paragraph B.2 Implementation Requirements. The operating environment is as follows:

Input: A = Character to send
 DP = ASSIST09 work page
 .PAD = Character and new line padding counts
 (in vector table)
 .PAUSE = Pause routine (in vector table)

Output: Character sent to the output device

Volatile Registers: None. All work registers must be restored

.COOFF

Output Console Shutdown Routine

.COOFF

Code: 30

Description: This entry addresses the routine to terminate output device processing. ASSIST09 does not call this routine. It is included for completeness. The default routine is an "RTS".

Input: DP → ASSIST09 work page
Output: The output device is terminated
Volatile Registers: None

.COON

Output Console Initialization Routine

.COON

Code: 26

Description: This entry points to a routine to initialize the standard output device. The default routine initializes an ACIA and is the very same one described under the .CION vector swap definition.

Input: .ACIA vector entry for the ACIA address
Output: The output device is initialized
Volatile Registers: A, X

.ECHO

Echo Flag

.ECHO

Code: 50

Description: The first byte of this word is used as a flag for the INCHP service routine to determine the requirement of echoing input received from the input handler. A non-zero value means to echo the input; zero not to echo. The echoing will take place even if user handlers are substituted for the default .CIDTA handler as the INCHP service routine performs the echo.

.FIRQ

Fast Interrupt Request Vector Appendage

.FIRQ

Code: 10

Description: The fast interrupt request routine is located via this pointer. The MC6809 addresses hexadecimal FFF6 to locate the handler when processing a FIRQ. The stack and machine status is as defined for the FIRQ interrupt upon entry to this appendage. It should be noted that this routine is "jumped" to with an indirect jump instruction which adds eleven cycles to the interrupt time before the handler actually receives control. The default handler does an immediate "RTI" which, in essence, ignores the interrupt.

.HSDTA

High Speed Display Handler Routine

.HSDTA

Code: 32

Description: This entry is invoked as a subroutine by the DISPLAY command and passed a parameter list containing the "TO" and "FROM" addresses. The from value is rounded down to a 16 byte address boundary. The default routine displays memory in both hexadecimal and ASCII representations, with a title produced on every 128 byte boundary. The purpose for this vector table entry is for easy implementation of a user routine for special purpose handling of a block of data. (The data could, for example, be sent to a high speed printer for later analysis.) The parameters are all passed on the stack. The environment is as follows:

Input: S + 4 = Start address
 S + 2 = Stop address
 S + 0 = Return Address
 DP → ASSIST09 work page
Output: Any purpose desired
Volatile Registers: X, D

.IRQ

Interrupt Request Vector Appendage

.IRQ

Code: 12

Description: All interrupt requests are passed to the routine pointed to by this vector. Hexadecimal FFF8 is the MC6809 location where this interrupt vector is fetched. The stack and processor status is that defined for the IRQ interrupt upon entry to the handler. Since the routine's address is in the vector table, an indirect jump must be done to invoke it. This adds eleven cycles to the interrupt time before the IRQ handler receives control. The default IRQ handler prints the registers and enters the ASSIST09 command handler.

.NMI

Non-Maskable Interrupt Vector Appendage

.NMI

Code: 16

Description: This entry points to the non-maskable interrupt handler to receive control whenever the processor branches to the address at hexadecimal FFFC. Since ASSIST09 uses the NMI interrupt during trace and breakpoint processing, such commands should not be used if a user handler is in control. This is true unless the user handler has the intelligence to forward control to the default handler if the NMI interrupt has not been generated due to user facilities. The NMI handler given control will have an eleven cycle overhead as its address must be fetched from the vector table.

.PAD

Character and New Line Pad Count

.PAD

Code: 48

Description: This entry contains the pad count for characters and new lines. The first of the two bytes is the count of nulls for other characters, and the second is the number of nulls (\$00) to send out after any line feed is transmitted. The ASCII Escape character (\$10) never has nulls sent following it. The default .CODTA handler is responsible for transmitting these nulls. A user handler may or may not use these counts as required.

The “NULLS” command also sets these two bytes with user specified values.

.PAUSE

Processing Pause Routine

.PAUSE

Code: 40

Description: In order to support real-time (also known as multi-tasking) environments ASSIST09 calls a dead-time routine whenever processing must wait for some external change of state. An example would be when the OUTCH service routine attempts the sending of a character to the ACIA through the default .CODTA handler and the ACIA status registers shows that it cannot yet be accepted. The default dead-time routine resides in a reserved four byte area which contains the single instruction, "RTS". The .PAUSE vector entry points to this routine after standard initialization. This pointer may be changed to point to a user routine which dispatches other programs so that the MC6809 may be utilized more efficiently. Another example of use would be to increment a counter so that dead-time cycle counts may be accumulated for statistical or debugging purposes. The reason for the four byte reserved area (which exists in the ASSIST09 work page) is so other code may be overlayed without the need for another space in the address map to be assigned. For example, a master monitor may be using a memory management unit to assign a complete 64K block of memory to ASSIST09 and the programs being executed/tested under ASSIST09 control. The master monitor wishes, or course, to be reentered when any "dead time" occurs, so it overlays the default routine ("RTS") with its own "SWI". Since the master monitor would be "front ending" all "SWI's" anyway, it knows when a "pause" call is being performed and can redispach other systems on a time-slice basis.

All registers must be transparent across the pause handler. Along with selected points in ASSIST09 user service processing, there is a special service call specifically for user programs to invoke the pause routine. It may be suggested that if no services are being requested for a given time period (say 10 ms) user programs should call the .PAUSE service routine so that fair-task dispatching can be guaranteed.

.PTM

Programmable Timer Module Address

.PTM

Code: 53

Description: This entry contains the address of the MC6840 programmable timer module (PTM). Alteration of this slot should occur before the MONITR startup service is called as explained in Paragraph B.4 Initialization. If no PTM is available, then the address should be changed to a zero so that no initialization attempt will take place. Note that if a zero is supplied, ASSIST09 Breakpoint and Trace commands should not be issued.

.RESET

Reset Interrupt Vector Appendage

.RESET

Code: 18

Description: This entry returns the address of the RESET routine which initializes ASSIST09. Changing it has no effect, but it is included in the vector table in case a user program wishes to determine where the ASSIST09 restart code resides. For example, if ASSIST09 resides in the memory map such that it does not control the MC6809 hardware vectors, a user routine may wish to start it up and thus need to obtain the standard RESET vector code address. The ASSIST09 reset code assigns the default in the work page, calls the vector build subroutine, and then starts ASSIST09 proper with the MONITR service call.

.RSVD

Reserved MC6809 Interrupt Vector Appendage

.RSVD

Code: 4

Description: This is a pointer to the reserved interrupt vector routine addressed at hexadeciml FFF0. This MC6809 hardware vector is not defined as yet. The default routine setup by ASSIST09 will cause a register display and entrance to the command handler.

.SWI

.SWI

Software Interrupt Vector Appendage

Code: 14

Description: This vector entry contains the address of the Software Interrupt routine. Normally, ASSIST09 handles these interrupts to provide services for user programs. If a user handler is in place, however, these facilities cannot be used unless the user routine "passes on" such requests to the ASSIST09 default handler. This is easy to do, since the vector swap function passes back the address of the default handler when the switch is made by the user. This "front ending" allows a user routine to examine all service calls, or alter/replace/extend them to his requirements. Of course, the registers must be transparent across the transfer of control from the user to the standard handler. A "JMP" instruction branches directly to the routine pointed to by this vector entry when a SWI occurs. Therefore, the environment is that as defined for the "SWI" interrupt.

.SWI2

Software Interrupt 2 Vector Appendage

.SWI2

Code: 8

Description: This entry contains a pointer to the SWI2 handler entered whenever that instruction is executed. The status of the stack and machine are those defined for the SWI2 interrupt which has its interrupt vector address at FFF4 hexadecimal. The default handler prints the registers and enters the ASSIST09 command handler.

.SWI3

Software Interrupt 3 Vector Appendage

.SWI3

Code: 6

Description: This entry contains a pointer to the SWI3 handler entered whenever that instruction is executed. The status of the stack and machine are those defined for the SWI3 interrupt which has its interrupt vector address located at hexadecimal FFF2. The default handler prints the registers and enters the ASSIST09 command handler.

PLEASE NOTE:

I did not scan this ASSIST09 listing from the Motorola book. The listing was very large, and I was scanning a bound book which required each page to be scanned individually. Instead, I assembled the ASSIST09 source code using my own ASM09 assembler and placed the resultant listing into these pages. This has the added advantage that the text is searchable and can be extracted if you wish to use code snippets in other places. It does mean however that this listing is not precisely identical to the one originally printed in the Motorola MC6809-MC6809E Programming Reference manual.

Also note: I found the source code on which this listing is based via an internet search. It appears to be the original Motorola source code, however it had been modified both in function (code changes) and source format (different assembler). I have attempted to restore it as closely as possible to the original - please notify me if you find errors. For the most part, the source was compatible with my assembler. I did have to change the single-quote character constants to double-quote format ('a' instead of 'a'), and some of the directives are slightly different. (TITLE instead of TTL for example).

Dave Dunfield

```

0000      1 ****
0000      2 * COPYRIGHT (C) MOTOROLA, INC. 1979 *
0000      3 ****
0000      4 ****
0000      5 ****
0000      6 * THIS IS THE BASE ASSIST09 ROM.
0000      7 * IT MAY RUN WITH OR WITHOUT THE
0000      8 * EXTENSION ROM WHICH
0000      9 * WHEN PRESENT WILL BE AUTOMATICALLY
0000     10 * INCORPORATED BY THE BLDVTR
0000     11 * SUBROUTINE.
0000     12 ****
0000     13 ****
0000     14 ****
0000     15 *      GLOBAL MODULE EQUATES
0000     16 ****
F800     17 ROMBEG EQU    $F800      ROM START ASSEMBLY ADDRESS
E700     18 RAMOFS EQU   -$1900      ROM OFFSET TO RAM WORK PAGE
0800     19 ROMSIZ EQU    2048      ROM SIZE
F000     20 ROM2OF EQU    ROMBEG-ROMSIZ  START OF EXTENSION ROM
E008     21 ACIA   EQU    $E008      DEFAULT ACIA ADDRESS
E000     22 PTM    EQU    $E000      DEFAULT PTM ADDRESS
0000     23 DFTCHP EQU    0          DEFAULT CHARACTER PAD COUNT
0005     24 DFTNLP EQU    5          DEFAULT NEW LINE PAD COUNT
003E     25 PROMPT EQU    '>'        PROMPT CHARACTER
0008     26 NUMBKP EQU    8          NUMBER OF BREAKPOINTS
0000     27 ****
0000     28 ****
0000     29 ****
0000     30 *      MISCELLANEOUS EQUATES
0000     31 ****
0004     32 EOT    EQU    $04        END OF TRANSMISSION
0007     33 BELL   EQU    $07        BELL CHARACTER
000A     34 LF     EQU    $0A        LINE FEED
000D     35 CR     EQU    $0D        CARRIAGE RETURN
0010     36 DLE    EQU    $10        DATA LINK ESCAPE
0018     37 CAN    EQU    $18        CANCEL (CTL-X)
0000     38 *      PTM ACCESS DEFINITIONS
E001     39 PTMSTA EQU    PTM+1      READ STATUS REGISTER
E000     40 PTMC13 EQU    PTM        CONTROL REGISTERS 1 AND 3
E001     41 PTMC2  EQU    PTM+1      CONTROL REGISTER 2
E002     42 PTMTM1 EQU    PTM+2      LATCH 1
E004     43 PTMTM2 EQU    PTM+4      LATCH 2
E006     44 PTMTM3 EQU    PTM+6      LATCH 3
0000     45 ****
0000     46 SKIP2  EQU    $8C        "CMPX #" OPCODE - SKIPS TWO BYTES
0000     47 ****
0000     48 ****
0000     49 *      ASSIST09 MONITOR SWI FUNCTIONS
0000     50 *      THE FOLLOWING EQUATES DEFINE FUNCTIONS PROVIDED
0000     51 *      BY THE ASSIST09 MONITOR VIA THE SWI INSTRUCTION.
0000     52 ****
0000     53 INCHNP EQU    0          INPUT CHAR IN A REG - NO PARITY
0001     54 OUTCH  EQU    1          OUTPUT CHAR FROM A REG
0002     55 PDATA1 EQU    2          OUTPUT STRING
0003     56 PDATA  EQU    3          OUTPUT CR/LF THEN STRING
0004     57 OUT2HS EQU    4          OUTPUT TWO HEX AND SPACE
0005     58 OUT4HS EQU    5          OUTPUT FOUR HEX AND SPACE
0006     59 PCRLF  EQU    6          OUTPUT CR/LF
0007     60 SPACE   EQU    7          OUTPUT A SPACE
0008     61 MONITR EQU    8          ENTER ASSIST09 MONITOR
0009     62 VCTRSW EQU    9          VECTOR EXAMINE/SWITCH
000A     63 BRKPT  EQU    10         USER PROGRAM BREAKPOINT
000B     64 PAUSE   EQU    11         TASK PAUSE FUNCTION
000B     65 NUMFUN  EQU    11         NUMBER OF AVAILABLE FUNCTIONS
0000     66 *      NEXT SUB-CODES FOR ACCESSING THE VECTOR TABLE.
0000     67 *      THEY ARE EQUIVALENT TO OFFSETS IN THE TABLE.
0000     68 *      RELATIVE POSITIONING MUST BE MAINTAINED.

```

0000	69	.AVTBL	EQU	0	ADDRESS OF VECTOR TABLE
0002	70	.CMDL1	EQU	2	FIRST COMMAND LIST
0004	71	.RSVD	EQU	4	RESERVED HARDWARE VECTOR
0006	72	.SWI3	EQU	6	SWI3 ROUTINE
0008	73	.SWI2	EQU	8	SWI2 ROUTINE
000A	74	.FIRQ	EQU	10	FIRQ ROUTINE
000C	75	.IRQ	EQU	12	IRQ ROUTINE
000E	76	.SWI	EQU	14	SWI ROUTINE
0010	77	.NMI	EQU	16	NMI ROUTINE
0012	78	.RESET	EQU	18	RESET ROUTINE
0014	79	.CION	EQU	20	CONSOLE ON
0016	80	.CIDTA	EQU	22	CONSOLE INPUT DATA
0018	81	.CIOFF	EQU	24	CONSOLE INPUT OFF
001A	82	.COON	EQU	26	CONSOLE OUTPUT ON
001C	83	.CODTA	EQU	28	CONSOLE OUTPUT DATA
001E	84	.COOFF	EQU	30	CONSOLE OUTPUT OFF
0020	85	.HSDTA	EQU	32	HIGH SPEED PRINTDATA
0022	86	.BSON	EQU	34	PUNCH/LOAD ON
0024	87	.BSDTA	EQU	36	PUNCH/LOAD DATA
0026	88	.BSOFF	EQU	38	PUNCH/LOAD OFF
0028	89	.PAUSE	EQU	40	TASK PAUSE ROUTINE
002A	90	.EXPAN	EQU	42	EXPRESSION ANALYZER
002C	91	.CMDL2	EQU	44	SECOND COMMAND LIST
002E	92	.ACIA	EQU	46	ACIA ADDRESS
0030	93	.PAD	EQU	48	CHARACTER PAD AND NEW LINE PAD
0032	94	.ECHO	EQU	50	ECHO/LOAD AND NULL BKPT FLAG
0034	95	.PTM	EQU	52	PTM ADDRESS
001B	96	NUMVTR	EQU	52/2+1	NUMBER OF VECTORS
0034	97	HIVTR	EQU	52	HIGHEST VECTOR OFFSET

```

0000          99  ****
0000         100  *      WORK AREA
0000         101  * THIS WORK AREA IS ASSIGNED TO THE PAGE ADDRESSED BY
0000         102  * -$1800,PCR FROM THE BASE ADDRESS OF THE ASSIST09
0000         103  * ROM. THE DIRECT PAGE REGISTER DURING MOST ROUTINE
0000         104  * OPERATIONS WILL POINT TO THIS WORK AREA. THE STACK
0000         105  * INITIALLY STARTS UNDER THE RESERVED WORK AREAS AS
0000         106  * DEFINED HEREIN.
0000         107  ****
DF00        108  WORKPG EQU     ROMBEG+RAMOFS   SETUP DIRECT PAGE ADDRESS
0000        109  SETDP  =WORKPG    NOTIFY ASSEMBLER
E000        110  ORG     WORKPG+256  READY PAGE DEFINITIONS
E000        111  * THE FOLLOWING THRU BKPTOP MUST RESIDE IN THIS ORDER
E000        112  * FOR PROPER INITIALIZATION
DFFC        113  ORG     *-4
DFFC        114  PAUSER EQU     *      PAUSE ROUTINE
DFFB        115  ORG     *-1
DFFB        116  SWIBFL EQU     *      BYPASS SWI AS BREAKPOINT FLAG
DFFA        117  ORG     *-1
DFFA        118  BKPTCT EQU     *      BREAKPOINT COUNT
DFF8        119  ORG     *-2
DFF8        120  SLEVEL EQU     *      STACK TRACE LEVEL
DFC2        121  ORG     *-(NUMVTR*2)  VECTOR TABLE
DFC2        122  VECTAB EQU     *      VECTOR TABLE
DFB2        123  ORG     *-(2*NUMBKP)  BREAKPOINT TABLE
DFB2        124  BKPTBL EQU     *      BREAKPOINT TABLE
DFA2        125  ORG     *-(2*NUMBKP)  BREAKPOINT OPCODE TABLE
DFA2        126  BKPTOP EQU     *      BREAKPOINT OPCODE TABLE
DFA0        127  ORG     *-2
DFA0        128  WINDOW EQU     *      WINDOW
DF9E        129  ORG     *-2
DF9E        130  ADDR    EQU     *      ADDRESS POINTER VALUE
DF9D        131  ORG     *-1
DF9D        132  BASEPG EQU     *      BASE PAGE VALUE
DF9B        133  ORG     *-2
DF9B        134  NUMBER EQU     *      BINARY BUILD AREA
DF99        135  ORG     *-2
DF99        136  LASTOP EQU     *      LAST OPCODE TRACED
DF97        137  ORG     *-2
DF97        138  RSTACK EQU     *      RESET STACK POINTER
DF95        139  ORG     *-2
DF95        140  PSTACK EQU     *      COMMAND RECOVERY STACK
DF93        141  ORG     *-2
DF93        142  PCNTER EQU     *      LAST PROGRAM COUNTER
DF91        143  ORG     *-2
DF91        144  TRACEC EQU     *      TRACE COUNT
DF90        145  ORG     *-1
DF90        146  SWICNT EQU     *      TRACE "SWI" NEST LEVEL COUNT
DF8F        147  ORG     *-1  (MISFLG MUST FOLLOW SWICNT)
DF8F        148  MISFLG EQU     *      LOAD CMD/THRU BREAKPOINT FLAG
DF8E        149  ORG     *-1
DF8E        150  DELIM  EQU     *      EXPRESSION DELIMITER/WORK BYTE
DF66        151  ORG     *-40
DF66        152  ROM2WK EQU     *      EXTENSION ROM RESERVED AREA
DF51        153  ORG     *-21
DF51        154  TSTACK EQU     *      TEMPORARY STACK HOLD
DF51        155  STACK  EQU     *      START OF INITIAL STACK

```

```

DF51      157 ****
DF51      158 * DEFAULT THE ROM BEGINNING ADDRESS TO 'ROMBEG'
DF51      159 * ASSIST09 IS POSITION ADDRESS INDEPENDENT, HOWEVER
DF51      160 * WE ASSEMBLE ASSUMING CONTROL OF THE HARDWARE VECTORS.
DF51      161 * NOTE THAT THE WORK RAM PAGE MUST BE 'RAMOFS'
DF51      162 * FROM THE ROM BEGINNING ADDRESS.
DF51      163 ****
F800      164     ORG     ROMBEG           ROM ASSEMBLY/DEFAULT ADDRESS
F800      165
F800      166 ****
F800      167 *
F800      168 * BLDVTR - BUILD ASSIST09 VECTOR TABLE
F800      169 * HARDWARE RESET CALLS THIS SUBROUTINE TO BUILD THE
F800      170 * ASSIST09 VECTOR TABLE. THIS SUBROUTINE RESIDES AT
F800      171 * THE FIRST BYTE OF THE ASSIST09 ROM, AND CAN BE
F800      172 * CALLED VIA EXTERNAL CONTROL CODE FOR REMOTE
F800      173 * ASSIST09 EXECUTION.
F800      174 * INPUT: S->VALID STACK RAM
F800      175 * OUTPUT: U->VECTOR TABLE ADDRESS
F800      176 * DPR->ASSIST09 WORK AREA PAGE
F800      177 * THE VECTOR TABLE AND DEFAULTS ARE INITIALIZED
F800      178 * ALL REGISTERS VOLATILE
F800      179 ****
F800 30 8D E7 BE 180 BLDVTR LEAX   VECTAB,PCR   ADDRESS VECTOR TABLE
F804 1F 10 181 TFR    X,D    OBTAIN BASE PAGE ADDRESS
F806 1F 8B 182 TFR    A,DP   SETUP DPR
F808 97 9D 183 STA    BASEPG STORE FOR QUICK REFERENCE
F80A 33 84 184 LEAU   ,X    RETURN TABLE TO CALLER
F80C 31 8C 35 185 LEAY   <INITVT,PCR LOAD FROM ADDR
F80F EF 81 186 STU    ,X++  INIT VECTOR TABLE ADDRESS
F811 C6 16 187 LDB    #NUMVTR-5 NUMBER RELOCATABLE VECTORS
F813 34 04 188 PSHS   B     STORE INDEX ON STACK
F815 1F 20 189 BLD2   TFR   Y,D   PREPARE ADDRESS RESOLVE
F817 E3 A1 190 ADDD   ,Y++  TO ABSOLUTE ADDRESS
F819 ED 81 191 STD    ,X++  INTO VECTOR TABLE
F81B 6A E4 192 DEC    ,S    COUNT DOWN
F81D 26 F6 193 BNE    BLD2  BRANCH IF MORE TO INSERT
F81F C6 0D 194 LDB    #INTVE-INTVS STATIC VALUE INIT LENGTH
F821 A6 A0 195 BLD3   LDA   ,Y+  LOAD NEXT BYTE
F823 A7 80 196 STA    ,X+  STORE INTO POSITION
F825 5A 197 DECB
F826 26 F9 198 BNE    BLD3  LOOP UNTIL DONE
F828 31 8D F7 D4 199 LEAY   ROM2OF,PCR TEST POSSIBLE EXTENSION ROM
F82C 8E 20 FE 200 LDX   #$20FE LOAD "BRA *" FLAG PATTERN
F82F AC A1 201 CMPX   ,Y++ ? EXTENDED ROM HERE
F831 26 02 202 BNE    BLDRTN BRANCH NOT OUR ROM TO RETURN
F833 AD A4 203 JSR    ,Y    CALL EXTENDED ROM INITIALIZE
F835 35 84 204 BLDRTN PULS  PC,B  RETURN TO INITIALIZER
F837 205
F837 206 ****
F837 207 *
F837 208 * RESET ENTRY POINT
F837 209 * HARDWARE RESET ENTERS HERE IF ASSIST09 IS ENABLED
F837 210 * TO RECEIVE THE MC6809 HARDWARE VECTORS. WE CALL
F837 211 * THE BLDVTR SUBROUTINE TO INITIALIZE THE VECTOR
F837 212 * TABLE, STACK, AND THEN FIREUP THE MONITOR VIA SWI
F837 213 ****
F837 32 8D E7 16 214 RESET  LEAS   STACK,PCR  SETUP INITIAL STACK
F83B 8D C3 215 BSR    BLDVTR BUILD VECTOR TABLE
F83D 4F 216 RESET2 CLRA
F83E 1F 8B 217 TFR    A,DP   ISSUE STARTUP MESSAGE
F840 3F 218 SWI
F841 08 219 FCB    MONITR TO ENTER COMMAND PROCESSING
F842 20 F9 220 BRA    RESET2 REENTER MONITOR IF 'CONTINUE'
F844 221
F844 222 ****
F844 223 * INITVT - INITIAL VECTOR TABLE
F844 224 * THIS TABLE IS RELOCATED TO RAM AND REPRESENTS THE

```

```

F844          225 * INITIAL STATE OF THE VECTOR TABLE. ALL ADDRESSES
F844          226 * ARE CONVERTED TO ABSOLUTE FORM. THIS TABLE STARTS
F844          227 * WITH THE SECOND ENTRY, ENDS WITH STATIC CONSTANT
F844          228 * INITIALIZATION DATA WHICH CARRIES BEYOND THE TABLE.
F844          229 ****
F844 01 58    230 INITVT FDB  CMDTBL-*      DEFAULT FIRST COMMAND TABLE
F846 02 92    231 FDB  RSRVDR-*      DEFAULT UNDEFINED HARDWARE VECTOR
F848 02 90    232 FDB  SWI3R-*       DEFAULT SWI3
F84A 02 8E    233 FDB  SWI2R-*       DEFAULT SWI2
F84C 02 70    234 FDB  FIRQR-*       DEFAULT IRQ
F84E 02 8A    235 FDB  IRQR-*        DEFAULT IRQ ROUTINE
F850 00 45    236 FDB  SWIR-*        DEFAULT SWI ROUTINE
F852 02 2B    237 FDB  NMIR-*        DEFAULT NMI ROUTINE
F854 FF E3    238 FDB  RESET-*       RESTART VECTOR
F856 02 90    239 FDB  CION-*        DEFAULT CION
F858 02 84    240 FDB  CIDTA-*       DEFAULT CIDTA
F85A 02 96    241 FDB  CIOFF-*       DEFAULT CIOFF
F85C 02 8A    242 FDB  COON-*        DEFAULT COON
F85E 02 93    243 FDB  CODTA-*       DEFAULT CODTA
F860 02 90    244 FDB  COOFF-*       DEFAULT COOFF
F862 03 9A    245 FDB  HSDTA-*       DEFAULT HSDTA
F864 02 B7    246 FDB  BSON-*        DEFAULT BSON
F866 02 D2    247 FDB  BSDTA-*       DEFAULT BSDTA
F868 02 BF    248 FDB  BSOFF-*       DEFAULT BSOFF
F86A E7 92    249 FDB  PAUSER-*      DEFAULT PAUSE ROUTINE
F86C 04 7D    250 FDB  EXP1-*        DEFAULT EXPRESSION ANALYZER
F86E 01 2D    251 FDB  CMDTB2-*      DEFAULT SECOND COMMAND TABLE
F870          252 * CONSTANTS
F870 E0 08    253 INTVS FDB  ACIA      DEFAULT ACIA
F872 00 05    254 FCB  DFTCHP,DFTNLP  DEFAULT NULL PADDs
F874 00 00    255 FDB  0           DEFAULT ECHO
F876 E0 00    256 FDB  PTM         DEFAULT PTM
F878 00 00    257 FDB  0           INITIAL STACK TRACE LEVEL
F87A 00      258 FCB  0           INITIAL BREAKPOINT COUNT
F87B 00      259 FCB  0           SWI BREAKPOINT LEVEL
F87C 39      260 FCB  $39         DEFAULT PAUSE ROUTINE (RTS)
F87D          261 INTVE EQU  *
F87D          262 *B
F87D          263
F87D          264 ****
F87D          265 * ASSIST09 SWI HANDLER
F87D          266 * THE SWI HANDLER PROVIDES ALL INTERFACING NECESSARY
F87D          267 * FOR A USER PROGRAM. A FUNCTION BYTE IS ASSUMED TO
F87D          268 * FOLLOW THE SWI INSTRUCTION. IT IS BOUND CHECKED
F87D          269 * AND THE PROPER ROUTINE IS GIVEN CONTROL. THIS
F87D          270 * INVOCATION MAY ALSO BE A BREAKPOINT INTERRUPT.
F87D          271 * IF SO, THE BREAKPOINT HANDLER IS ENTERED.
F87D          272 * INPUT: MACHINE STATE DEFINED FOR SWI
F87D          273 * OUTPUT: VARIES ACCORDING TO FUNCTION CALLED. PC ON
F87D          274 * CALLERS STACK INCREMENTED BY ONE IF VALID CALL.
F87D          275 * VOLATILE REGISTERS: SEE FUNCTIONS CALLED
F87D          276 * STATE: RUNS DISABLED UNLESS FUNCTION Clears I FLAG.
F87D          277 ****
F87D          278
F87D          279 * SWI FUNCTION VECTOR TABLE
F87D 01 94    280 SWIVTB FDB  ZINCH-SWIVTB  INCHNP
F87F 01 B1    281 FDB  ZOTCH1-SWIVTB  OUTCH
F881 01 CB    282 FDB  ZPDTA1-SWIVTB  PDATA1
F883 01 C3    283 FDB  ZPDATA-SWIVTB  PDATA
F885 01 75    284 FDB  ZOT2HS-SWIVTB  OUT2HS
F887 01 73    285 FDB  ZOT4HS-SWIVTB  OUT4HS
F889 01 C0    286 FDB  ZPCRLF-SWIVTB  PCRLF
F88B 01 79    287 FDB  ZSPACE-SWIVTB  SPACE
F88D 00 55    288 FDB  ZMONTR-SWIVTB  MONITR
F88F 01 7D    289 FDB  ZVSWTH-SWIVTB  VCTRSW
F891 02 56    290 FDB  ZBKPNTR-SWIVTB  BREAKPOINT
F893 01 D1    291 FDB  ZPAUSE-SWIVTB  TASK PAUSE
F895          292

```

```

F895  6A 8D E6 F7      293  SWIR   DEC     SWICNT,PCR    UP "SWI" LEVEL FOR TRACE
F899  17 02 25          294  LBSR   LDDP    SETUP PAGE AND VERIFY STACK
F89C
F89C  EE 6A             295  * CHECK FOR BREAKPOINT TRAP
F89E  33 5F             296  LDU    10,S    LOAD PROGRAM COUNTER
F8A0  0D FB             297  LEAU   -1,U    BACK TO SWI ADDRESS
F8A2  26 11             298  TST    SWIBFL  ? THIS "SWI" BREAKPOINT
F8A4  17 06 9B           299  BNE    SWIDNE BRANCH IF SO TO LET THROUGH
F8A7  50                 300  LBSR   CBKLDR OBTAIN BREAKPOINT POINTERS
F8A8  5A                 301  NEGB
F8A9  2B 0A             302  SWILP  DECB   OBTAIN POSITIVE COUNT
F8AB  11 A3 A1           303  BMI    SWIDNE COUNT DOWN
F8AE  26 F8             304  CMPU   ,Y++   BRANCH WHEN DONE
F8B0  EF 6A             305  BNE    SWILP   ? WAS THIS A BREAKPOINT
F8B2  16 02 1E           306  STU    10,S    BRANCH IF NOT
F8B5  0F FB             307  LBRA   ZBKPNTR SET PROGRAM COUNTER BACK
F8B7  37 06             308  SWIDNE CLR    GO DO BREAKPOINT
F8B9  C1 0B             309  PULU   D      CLEAR IN CASE SET
F8BB  10 22 02 0F         310  CMPB   #NUMFUN OBTAIN FUNCTION BYTE, UP PC
F8BF  EF 6A             311  LBHI   ERROR  ? TOO HIGH
F8C1  58                 312  STU    10,S    YES, DO BREAKPOINT
F8C2  33 8C B8           313  ASLB    FUNCTION BUMP PROGRAM COUNTER PAST SWI
F8C5  EC C5             314  LEAU   SWIVTB,PCR OBTAIN VECTOR BRANCH ADDRESS
F8C7  6E CB             315  LDD    B,U     LOAD OFFSET
F8C9
F8C9  316  JMP    D,U     JUMP TO ROUTINE
F8C9
F8C9  317
F8C9  318  ****
F8C9  319  * REGISTERS TO FUNCTION ROUTINES:
F8C9  320  * DP-> WORK AREA PAGE
F8C9  321  * D,Y,U=UNRELIABLE          X=AS CALLED FROM USER
F8C9  322  * S=AS FROM SWI INTERRUPT
F8C9  323  ****
F8C9
F8C9  324
F8C9  325  ****
F8C9  326  * [SWI FUNCTION 8]
F8C9  327  * MONITOR ENTRY
F8C9  328  * FIREUP THE ASSIST09 MONITOR.
F8C9  329  * THE STACK WITH ITS VALUES FOR THE DIRECT PAGE
F8C9  330  * REGISTER AND CONDITION CODE FLAGS ARE USED AS IS.
F8C9  331  * 1) INITIALIZE CONSOLE I/O
F8C9  332  * 2) OPTIONALLY PRINT SIGNON
F8C9  333  * 3) INITIALIZE PTM FOR SINGLE STEPPING
F8C9  334  * 4) ENTER COMMAND PROCESSOR
F8C9  335  * INPUT: A=0 INIT CONSOLE AND PRINT STARTUP MESSAGE
F8C9  336  * A#0 OMIT CONSOLE INIT AND STARTUP MESSAGE
F8C9  337  ****
F8C9  338
F8C9  41 53 53 49 53 54 + 339  SIGNON FCC    /ASSIST09/   SIGNON EYE-CATCHER
F8D1  04                 340  FCB    EOT
F8D2
F8D2  10 DF 97           341
F8D5  6D 61             342  ZMONTR STS    RSTACK  SAVE FOR BAD STACK RECOVERY
F8D7  26 0D             343  TST    1,S     ? INIT CONSOLE AND SEND MSG
F8D9  AD 9D E6 F9         344  BNE    ZMONT2 BRANCH IF NOT
F8DD  AD 9D E6 FB         345  JSR    [VECTAB+.CION,PCR] READY CONSOLE INPUT
F8E1  30 8C E5           346  JSR    [VECTAB+.COON,PCR] READY CONSOLE OUTPUT
F8E4  3F                 347  LEAX   SIGNON,PCR  READY SIGNON EYE-CATCHER
F8E5  03                 348  SWI
F8E6  9E F6             349  FCB    PDATA  PRINT STRING
F8E8  27 0D             350  ZMONT2 LDX    VECTAB+.PTM  LOAD PTM ADDRESS
F8EA  6F 02             351  BEQ    CMD    BRANCH IF NOT TO USE A PTM
F8EC  6F 03             352  CLR    PTMTM1-PTM,X  SET LATCH TO CLEAR RESET
F8EE  CC 01 A6           353  CLR    PTMTM1+1-PTM,X AND SET GATE HIGH
F8F1  A7 01             354  LDD    #$01A6  SETUP TIMER 1 MODE
F8F3  E7 00             355  STA    PTMC2-PTM,X  SETUP FOR CONTROL REGISTER1
F8F5
F8F5  6F 01             356  STB    PTMC13-PTM,X SET OUTPUT ENABLED/
F8F7
F8F7  357  * SINGLE SHOT/ DUAL 8 BIT/INTERNAL MODE/OPERATE
F8F7  358  CLR    PTMC2-PTM,X  SET CR2 BACK TO RESET FORM
F8F7  359  * FALL INTO COMMAND PROCESSOR
F8F7  360

```

```

F8F7          361 ****
F8F7          362 *      COMMAND HANDLER
F8F7          363 *      BREAKPOINTS ARE REMOVED AT THIS TIME.
F8F7          364 *      PROMPT FOR A COMMAND, AND STORE ALL CHARACTERS
F8F7          365 *      UNTIL A SEPARATOR ON THE STACK.
F8F7          366 *      SEARCH FOR FIRST MATCHING COMMAND SUBSET,
F8F7          367 *      CALL IT OR GIVE '?' RESPONSE.
F8F7          368 *      DURING COMMAND SEARCH:
F8F7          369 *          B=OFFSET TO NEXT ENTRY ON X
F8F7          370 *          U=SAVED S
F8F7          371 *          U-1=ENTRY SIZE+2
F8F7          372 *          U-2=VALID NUMBER FLAG (>=0 VALID)/COMPARE CNT
F8F7          373 *          U-3=CARRIAGE RETURN FLAG (0=CR HAS BEEN DONE)
F8F7          374 *          U-4=START OF COMMAND STORE
F8F7          375 *          S+0=END OF COMMAND STORE
F8F7          376 ****
F8F7    3F    377 CMD     SWI           TO NEW LINE
F8F8    06    378   FCB     PCRLF         FUNCTION
F8F9          379 *      DISARM THE BREAKPOINTS
F8F9    17 06 46  380 CMDNEP  LBSR  CBKLDR    OBTAIN BREAKPOINT POINTERS
F8FC    2A 0C  381   BPL     CMDNOL       BRANCH IF NOT ARMED OR NONE
F8FE    50    382   NEGB    STB      BKPTCT    MAKE POSITIVE
F8FF    D7 FA  383   STB     CMDDDL      FLAG AS DISARMED
F901    5A    384   CMDDDL  DECB    ? FINISHED
F902    2B 06  385   BMI     CMDNOL       BRANCH IF SO
F904    A6 30  386   LDA     -NUMBK*2,Y  LOAD OPCODE STORED
F906    A7 B1  387   STA     [,Y+]        STORE BACK OVER "SWI"
F908    20 F7  388   BRA     CMDDDL      LOOP UNTIL DONE
F90A    AE 6A  389   CMDNOL  LDX    10,S        LOAD USERS PROGRAM COUNTER
F90C    9F 93  390   STX     PCNTER      SAVE FOR EXPRESSION ANALYZER
F90E    86 3E  391   LDA     #PROMPT    LOAD PROMPT CHARACTER
F910    3F    392   SWI           SEND TO OUTPUT HANDLER
F911    01    393   FCB     OUTCH       FUNCTION
F912    33 E4  394   LEAU    ,S          REMEMBER STACK RESTORE ADDRESS
F914    DF 95  395   STU     PSTACK      REMEMBER STACK FOR ERROR USE
F916    4F    396   CLRA    CLRB    PREPARE ZERO
F917    5F    397   CLRB    CLRB    PREPARE ZERO
F918    DD 9B  398   STD     NUMBER      CLEAR NUMBER BUILD AREA
F91A    DD 8F  399   STD     MISFLG     CLEAR MISCEL. AND SWICNT FLAGS
F91C    DD 91  400   STD     TRACEC    CLEAR TRACE COUNT
F91E    C6 02  401   LDB     #2        SET D TO TWO
F920    34 07  402   PSHS    D,CC      PLACE DEFAULTS ONTO STACK
F922          403 *      CHECK FOR "QUICK" COMMANDS.
F922    17 04 54  404   LBSR    READ       OBTAIN FIRST CHARACTER
F925    30 8D 05 81 405   LEAX    CDOT+2,PCR  PRESET FOR SINGLE TRACE
F929    81 2E  406   CMPA    #'.'      ? QUICK TRACE
F92B    27 5A  407   BEQ     CMDXQT    BRANCH EQUAL FOR TRACE ONE
F92D    30 8D 04 E9 408   LEAX    CMPADP+2,PCR  READY MEMORY ENTRY POINT
F931    81 2F  409   CMPA    #'/'      ? OPEN LAST USED MEMORY
F933    27 52  410   BEQ     CMDXQT    BRANCH TO DO IT IF SO
F935          411 *      PROCESS NEXT CHARACTER
F935    81 20  412   CMD2    CMPA    #' '      ? BLANK OR DELIMITER
F937    23 14  413   BLS     CMDGOT    BRANCH YES, WE HAVE IT
F939    34 02  414   PSHS    A          BUILD ONTO STACK
F93B    6C 5F  415   INC     -1,U      COUNT THIS CHARACTER
F93D    81 2F  416   CMPA    #'/'      ? MEMORY COMMAND
F93F    27 4F  417   BEQ     CMDMEM    BRANCH IF SO
F941    17 04 0B  418   LBSR    BLDHXC   TREAT AS HEX VALUE
F944    27 02  419   BEQ     CMD3      BRANCH IF STILL VALID NUMBER
F946    6A 5E  420   DEC     -2,U      FLAG AS INVALID NUMBER
F948    17 04 2E  421   CMD3    LBSR    READ      OBTAIN NEXT CHARACTER
F94B    20 E8  422   BRA     CMD2      TEST NEXT CHARACTER
F94D          423 *      GOT COMMAND, NOW SEARCH TABLES
F94D    80 0D  424   CMDGOT  SUBA    #CR      SET ZERO IF CARRIAGE RETURN
F94F    A7 5D  425   STA     -3,U      SETUP FLAG
F951    9E C4  426   LDX     VECTAB+.CMDL1  START WITH FIRST CMD LIST
F953    E6 80  427   CMDSCH  LDB     ,X+      LOAD ENTRY LENGTH
F955    2A 10  428   BPL     CMDSME    BRANCH IF NOT LIST END

```

```

F957 9E EE          429      LDX      VECTAB+.CMDL2    NOW TO SECOND CMD LIST
F959 5C             430      INCB     ? TO CONTINUE TO DEFAULT LIST
F95A 27 F7          431      BEQ      CMDSCH   BRANCH IF SO
F95C 10 DE 95       432      CMDBAD  LDS      PSTACK   RESTORE STACK
F95F 30 8D 01 5A   433      LEAX    ERRMSG,PCR  POINT TO ERROR STRING
F963 3F             434      SWI      SEND OUT
F964 02             435      FCB     PDATA1  TO CONSOLE
F965 20 90          436      BRA     CMD     AND TRY AGAIN
F967               437      * SEARCH NEXT ENTRY
F967 5A             438      CMDSME DECB    TAKE ACCOUNT OF LENGTH BYTE
F968 E1 5F          439      CMPB    -1,U    ? ENTERED LONGER THAN ENTRY
F96A 24 03          440      BHS     CMDSIZ  BRANCH IF NOT TOO LONG
F96C 3A             441      CMDFLS ABX    TO NEXT ENTRY
F96D 20 E4          442      BRA     CMDSCH  AND TRY NEXT
F96F 31 5D          443      CMDSIZ LEAY    PREPARE TO COMPARE
F971 A6 5F          444      LDA     -1,U    LOAD SIZE+2
F973 80 02          445      SUBA   #2     TO ACTUAL SIZE ENTERED
F975 A7 5E          446      STA     -2,U    SAVE SIZE FOR COUNTDOWN
F977 5A             447      CMDCMP DECB    DOWN ONE BYTE
F978 A6 80          448      LDA     ,X+    NEXT COMMAND CHARACTER
F97A A1 A2          449      CMPA    ,-Y    ? SAME AS THAT ENTERED
F97C 26 EE          450      BNE     CMDFLS  BRANCH TO FLUSH IF NOT
F97E 6A 5E          451      DEC     -2,U    COUNT DOWN LENGTH OF ENTRY
F980 26 F5          452      BNE     CMDCMP  BRANCH IF MORE TO TEST
F982 3A             453      ABX     TO     NEXT ENTRY
F983 EC 1E          454      LDD     -2,X    LOAD OFFSET
F985 30 8B          455      LEAX   D,X    COMPUTE ROUTINE ADDRESS+2
F987 6D 5D          456      CMIDXQT TST    SET CC FOR CARRIAGE RETURN TEST
F989 32 C4          457      LEAS   ,U    DELETE STACK WORK AREA
F98B AD 1E          458      JSR     -2,X    CALL COMMAND
F98D 16 FF 7A       459      LBRA   CMDNOL  GO GET NEXT COMMAND
F990 6D 5E          460      CMDMEM TST    ? VALID HEX NUMBER ENTERED
F992 2B C8          461      BMI    CMDBAD  BRANCH ERROR IF NOT
F994 30 88 AE       462      LEAX   <CMEMN-CMPADP,X TO DIFFERENT ENTRY
F997 DC 9B          463      LDD    NUMBER  LOAD NUMBER ENTERED
F999 20 EC          464      BRA    CMIDXQT AND ENTER MEMORY COMMAND
F99B
F99B               465      ** COMMANDS ARE ENTERED AS A SUBROUTINE WITH:
F99B               466      ** DPR->ASSIST09 DIRECT PAGE WORK AREA
F99B               467      ** Z=1 CARRIAGE RETURN ENTERED
F99B               468      ** Z=0 NON CARRIAGE RETURN DELIMITER
F99B               469      ** S=NORMAL RETURN ADDRESS
F99B               470      ** THE LABEL "CMDBAD" MAY BE ENTERED TO ISSUE AN
F99B               471      ** AN ERROR FLAG (*).
F99B               472      *
F99B               473      *
F99B               474      ****
F99B               475      * ASSIST09 COMMAND TABLES
F99B               476      * THESE ARE THE DEFAULT COMMAND TABLES. EXTERNAL
F99B               477      * TABLES OF THE SAME FORMAT MAY EXTEND/REPLACE
F99B               478      * THESE BY USING THE VECTOR SWAP FUNCTION.
F99B               479      *
F99B               480      * ENTRY FORMAT:
F99B               481      * +0...TOTAL SIZE OF ENTRY (INCLUDING THIS BYTE)
F99B               482      * +1...COMMAND STRING
F99B               483      * +N...TWO BYTE OFFSET TO COMMAND (ENTRYADDR-*)
F99B               484      *
F99B               485      * THE TABLES TERMINATE WITH A ONE BYTE -1 OR -2.
F99B               486      * THE -1 CONTINUES THE COMMAND SEARCH WITH THE
F99B               487      *      SECOND COMMAND TABLE.
F99B               488      * THE -2 TERMINATES COMMAND SEARCHES.
F99B               489      ****
F99B               490      *
F99B               491      * THIS IS THE DEFAULT LIST FOR THE SECOND COMMAND
F99B               492      * LIST ENTRY.
F99B   FE           493      CMDTB2 FCB     -2      STOP COMMAND SEARCHES
F99C
F99C               494      *
F99C               495      * THIS IS THE DEFAULT LIST FOR THE FIRST COMMAND
F99C               496      * LIST ENTRY.

```

F99C		497	CMDTBL	EQU	*	MONITOR COMMAND TABLE
F99C	04	498	FCB		4	
F99D	42	499	FCC		/B/	'BREAKPOINT' COMMAND
F99E	05 4D	500	FDB		CBKPT-*	
F9A0	04	501	FCB		4	
F9A1	43	502	FCC		/C/	'CALL' COMMAND
F9A2	04 17	503	FDB		CCALL-*	
F9A4	04	504	FCB		4	
F9A5	44	505	FCC		/D/	'DISPLAY' COMMAND
F9A6	04 9D	506	FDB		CDISP-*	
F9A8	04	507	FCB		4	
F9A9	45	508	FCC		/E/	'ENCODE' COMMAND
F9AA	05 9F	509	FDB		CENCDE-*	
F9AC	04	510	FCB		4	
F9AD	47	511	FCC		/G/	'GO' COMMAND
F9AE	03 D2	512	FDB		CGO-*	
F9B0	04	513	FCB		4	
F9B1	4C	514	FCC		/L/	'LOAD' COMMAND
F9B2	04 DD	515	FDB		CLOAD-*	
F9B4	04	516	FCB		4	
F9B5	4D	517	FCC		/M/	'MEMORY' COMMAND
F9B6	04 0D	518	FDB		CMEM-*	
F9B8	04	519	FCB		4	
F9B9	4E	520	FCC		/N/	'NULLS' COMMAND
F9BA	04 FD	521	FDB		CNULLS-*	
F9BC	04	522	FCB		4	
F9BD	4F	523	FCC		/O/	'OFFSET' COMMAND
F9BE	05 0A	524	FDB		COFFS-*	
F9C0	04	525	FCB		4	
F9C1	50	526	FCC		/P/	'PUNCH' COMMAND
F9C2	04 AF	527	FDB		CPUNCH-*	
F9C4	04	528	FCB		4	
F9C5	52	529	FCC		/R/	'REGISTERS' COMMAND
F9C6	02 84	530	FDB		CREG-*	
F9C8	04	531	FCB		4	
F9C9	53	532	FCC		/S/	'STLEVEL' COMMAND
F9CA	04 F2	533	FDB		CSTLEV-*	
F9CC	04	534	FCB		4	
F9CD	54	535	FCC		/T/	'TRACE' COMMAND
F9CE	04 D6	536	FDB		CTRACE-*	
F9D0	04	537	FCB		4	
F9D1	56	538	FCC		/V/	'VERIFY' COMMAND
F9D2	04 CF	539	FDB		CVER-*	
F9D4	04	540	FCB		4	
F9D5	57	541	FCC		/W/	'WINDOW' COMMAND
F9D6	04 68	542	FDB		CWINDO-*	
F9D8	FF	543	FCB		-1	END, CONTINUE WITH THE SECOND
F9D9		544				
F9D9		545				*****
F9D9		546	*			[SWI FUNCTIONS 4 AND 5]
F9D9		547	*	4 - OUT2HS	- DECODE BYTE TO HEX AND ADD SPACE	
F9D9		548	*	5 - OUT4HS	- DECODE WORD TO HEX AND ADD SPACE	
F9D9		549	*	INPUT: X->BYTE OR WORD TO DECODE		
F9D9		550	*	OUTPUT: CHARACTERS SENT TO OUTPUT HANDLER		
F9D9		551	*	X->NEXT BYTE OR WORD		
F9D9		552				*****
F9D9		553				
F9D9	A6 80	554	ZOUT2H	LDA	,X+	LOAD NEXT BYTE
F9DB	34 06	555	PSHS	D		SAVE - DO NOT REREAD
F9DD	C6 10	556	LDB	#16		SHIFT BY 4 BITS
F9DF	3D	557	MUL			WITH MULTIPLY
F9E0	8D 04	558	BSR	ZOUTHX		SEND OUT AS HEX
F9E2	35 06	559	PULS	D		RESTORE BYTES
F9E4	84 0F	560	ANDA	#\$0F		ISOLATE RIGHT HEX
F9E6	8B 90	561	ZOUTHX	ADDA	#\$90	PREPARE A-F ADJUST
F9E8	19	562	DAA		ADJUST	
F9E9	89 40	563	ADCA	#\$40		PREPARE CHARACTER BITS
F9EB	19	564	DAA		ADJUST	

```

F9EC 6E 9D E5 EE      565 SEND    JMP     [VECTAB+.CODTA,PCR]    SEND TO OUT HANDLER
F9F0
F9F0 8D E7      566
F9F2 8D E5      567 ZOT4HS  BSR     ZOUT2H      CONVERT FIRST BYTE
F9F4 AF 64      568 ZOT2HS  BSR     ZOUT2H      CONVERT BYTE TO HEX
F9F6
F9F6 569 STX    4,S      UPDATE USERS X REGISTER
F9F6 570 * FALL INTO SPACE ROUTINE
F9F6 571
F9F6 572 ****
F9F6 573 *          [SWI FUNCTION 7]
F9F6 574 *          SPACE - SEND BLANK TO OUTPUT HANDLER
F9F6 575 * INPUT: NONE
F9F6 576 * OUTPUT: BLANK SEND TO CONSOLE HANDLER
F9F6 577 ****
F9F6 86 20      578 ZSPACE LDA    #' '      LOAD BLANK
F9F8 20 3D      579 BRA     ZOTCH2    SEND AND RETURN
F9FA
F9FA 580
F9FA 581 ****
F9FA 582 *          [SWI FUNCTION 9]
F9FA 583 *          SWAP VECTOR TABLE ENTRY
F9FA 584 * INPUT: A=VECTOR TABLE CODE (OFFSET)
F9FA 585 *          X=0 OR REPLACEMENT VALUE
F9FA 586 * OUTPUT: X=PREVIOUS VALUE
F9FA 587 ****
F9FA A6 61      588 ZVSWTH LDA    1,S      LOAD REQUESTERS A
F9FC 81 34      589 CMPA   #HIVTR    ? SUB-CODE TOO HIGH
F9FE 22 39      590 BHI    ZOTCH3    IGNORE CALL IF SO
FA00 10 9E C2      591 LDY     VECTAB+.AVTBL  LOAD VECTOR TABLE ADDRESS
FA03 EE A6      592 LDU    A,Y      U=OLD ENTRY
FA05 EF 64      593 STU    4,S      RETURN OLD VALUE TO CALLERS X
FA07 AF 7E      594 STX    -2,S    ? X=0
FA09 27 2E      595 BEQ    ZOTCH3    YES, DO NOT CHANGE ENTRY
FA0B AF A6      596 STX    A,Y      REPLACE ENTRY
FA0D 20 2A      597 BRA     ZOTCH3    RETURN FROM SWI
FA0F
FA0F 598 *D
FA0F 599
FA0F 600 ****
FA0F 601 *          [SWI FUNCTION 0]
FA0F 602 *          INCHNP - OBTAIN INPUT CHAR IN A (NO PARITY)
FA0F 603 *          NULLS AND RUBOUTS ARE IGNORED.
FA0F 604 *          AUTOMATIC LINE FEED IS SENT UPON RECEIVING A
FA0F 605 *          CARRIAGE RETURN.
FA0F 606 *          UNLESS WE ARE LOADING FROM TAPE.
FA0F 607 ****
FA0F 8D 5D      608 ZINCHP BSR     XQPAUS    RELEASE PROCESSOR
FA11 8D 5F      609 ZINCH  BSR     XQCIDT    CALL INPUT DATA APPENDAGE
FA13 24 FA      610 BCC    ZINCHP    LOOP IF NONE AVAILABLE
FA15 4D
FA16 27 F9      611 TSTA   ?       TEST FOR NULL
FA18 81 7F      612 BEQ    ZINCH    IGNORE NULL
FA1A 27 F5      613 CMPA   #$7F    ? RUBOUT
FA1A 27 F5      614 BEQ    ZINCH    BRANCH YES TO IGNORE
FA1C A7 61      615 STA    1,S      STORE INTO CALLERS A
FA1E 0D 8F      616 TST     MISFLG   ? LOAD IN PROGRESS
FA20 26 17      617 BNE    ZOTCH3   BRANCH IF SO TO NOT ECHO
FA22 81 0D      618 CMPA   #CR     ? CARRIAGE RETURN
FA24 26 04      619 BNE    ZIN2     NO, TEST ECHO BYTE
FA26 86 0A      620 LDA    #LF     LOAD LINE FEED
FA28 8D C2      621 BSR    SEND    ALWAYS ECHO LINE FEED
FA2A 0D F4      622 ZIN2    TST     VECTAB+.ECHO ? ECHO DESIRED
FA2C 26 0B      623 BNE    ZOTCH3   NO, RETURN
FA2E
FA2E 624 * FALL THROUGH TO OUTCH
FA2E 625
FA2E 626 ****
FA2E 627 *          [SWI FUNCTION 1]
FA2E 628 *          OUTCH - OUTPUT CHARACTER FROM A
FA2E 629 * INPUT: NONE
FA2E 630 * OUTPUT: IF LINEFEED IS THE OUTPUT CHARACTER THEN
FA2E 631 *          C=0 NO CTL-X RECIEVED, C=1 CTL-X RECIEVED
FA2E 632 ****

```

```

FA2E A6 61          633 ZOTCH1 LDA    1,S      LOAD CHARACTER TO SEND
FA30 30 8C 09        634 LEAX   <ZPCRLS,PCR  DEFAULT FOR LINE FEED
FA33 81 0A          635 CMPA   #LF     ? LINE FEED
FA35 27 0F          636 BEQ    ZPDTLP  BRANCH TO CHECK PAUSE IF SO
FA37 8D B3          637 ZOTCH2 BSR    SEND    SEND TO OUTPUT ROUTINE
FA39 0C 90          638 ZOTCH3 INC    SWICNT BUMP UP "SWI" TRACE NEST LEVEL
FA3B 3B             639 RTI    RETURN FROM "SWI" FUNCTION
FA3C 640
FA3C 641 **** [SWI FUNCTION 6]
FA3C 642 * PCRLF - SEND CR/LF TO CONSOLE HANDLER
FA3C 643 * INPUT: NONE
FA3C 644 * OUTPUT: CR AND LF SENT TO HANDLER
FA3C 645 * C=0 NO CTL-X, C=1 CTL-X RECIEVED
FA3C 646 ****
FA3C 647 ****
FA3C 648
FA3C 04            649 ZPCRLS FCB    EOT    NULL STRING
FA3D 650
FA3D 30 8C FC        651 ZPCRLF LEAX   ZPCRLS,PCR  READY CR,LF STRING
FA40 652 * FALL INTO CR/LF CODE
FA40 653
FA40 654 **** [SWI FUNCTION 3]
FA40 655 * PDATA - OUTPUT CR/LF AND STRING
FA40 656 * INPUT: X->STRING
FA40 657 * OUTPUT: CR/LF AND STRING SENT TO OUTPUT CONSOLE
FA40 658 * HANDLER.
FA40 659 * C=0 NO CTL-X, C=1 CTL-X RECIEVED
FA40 660 * NOTE: LINE FEED MUST FOLLOW CARRIAGE RETURN FOR
FA40 661 * PROPER PUNCH DATA.
FA40 662 *
FA40 663 ****
FA40 86 0D          664 ZPDATA LDA    #CR    LOAD CARRIAGE RETURN
FA42 8D A8          665 BSR    SEND   SEND IT
FA44 86 0A          666 LDA    #LF    LOAD LINE FEED
FA46 667 * FALL INTO PDATA1
FA46 668
FA46 669 **** [SWI FUNCTION 2]
FA46 670 * PDATA1 - OUTPUT STRING TILL EOT ($04)
FA46 671 * THIS ROUTINE PAUSES IF AN INPUT BYTE BECOMES
FA46 672 * AVAILABLE DURING OUTPUT TRANSMISSION UNTIL A
FA46 673 * SECOND IS RECIEVED.
FA46 674 * INPUT: X->STRING
FA46 675 * OUTPUT: STRING SENT TO OUTPUT CONSOLE DRIVER
FA46 676 * C=0 NO CTL-X, C=1 CTL-X RECIEVED
FA46 677 *
FA46 678 ****
FA46 8D A4          679 ZPDTLP BSR    SEND   SEND CHARACTER TO DRIVER
FA48 A6 80          680 ZPDTA1 LDA    ,X+    LOAD NEXT CHARACTER
FA4A 81 04          681 CMPA   #EOT   ? EOT
FA4C 26 F8          682 BNE    ZPDTLP LOOP IF NOT
FA4E 683 * FALL INTO PAUSE CHECK FUNCTION
FA4E 684
FA4E 685 **** [SWI FUNCTION 12]
FA4E 686 * PAUSE - RETURN TO TASK DISPATCHING AND CHECK
FA4E 687 * FOR FREEZE CONDITION OR CTL-X BREAK
FA4E 688 * THIS FUNCTION ENTERS THE TASK PAUSE HANDLER SO
FA4E 689 * OPTIONALLY OTHER 6809 PROCESSES MAY GAIN CONTROL.
FA4E 690 * UPON RETURN, CHECK FOR A 'FREEZE' CONDITION
FA4E 691 * WITH A RESULTING WAIT LOOP, OR CONDITION CODE
FA4E 692 * RETURN IF A CONTROL-X IS ENTERED FROM THE INPUT
FA4E 693 * HANDLER.
FA4E 694 * C=1 IF CTL-X HAS ENTERED, C=0 OTHERWISE
FA4E 695 *
FA4E 696 ****
FA4E 8D 1E          697 ZPAUSE BSR    XQPAUS RELEASE CONTROL AT EVERY LINE
FA50 8D 06          698 BSR    CHKABT CHECK FOR FREEZE OR ABORT
FA52 1F A9          699 TFR    CC,B   PREPARE TO REPLACE CC
FA54 E7 E4          700 STB    ,S    OVERLAY OLD ONE ON STACK

```

FA56 20 E1	701	BRA	ZOTCH3	RETURN FROM "SWI"
FA58	702			
FA58	703	*	CHKABT - SCAN FOR INPUT PAUSE/ABORT DURING OUTPUT	
FA58	704	*	OUTPUT: C=0 OK, C=1 ABORT (CTL-X ISSUED)	
FA58	705	*	VOLATILE: U,X,D	
FA58 8D 18	706	CHKABT	BSR XQCIDT	ATTEMPT INPUT
FA5A 24 05	707		BCC CHKRPN	BRANCH NO TO RETURN
FA5C 81 18	708		CMPA #CAN	? CTL-X FOR ABORT
FA5E 26 02	709		BNE CHKWWT	BRANCH NO TO PAUSE
FA60 53	710	CHKSEC	COMB SET	CARRY
FA61 39	711	CHKRTN	RTS	RETURN TO CALLER WITH CC SET
FA62 8D 0A	712	CHKWT	BSR XQPAUS	PAUSE FOR A MOMENT
FA64 8D 0C	713		BSR XQCIDT	? KEY FOR START
FA66 24 FA	714		BCC CHKWWT	LOOP UNTIL RECEIVED
FA68 81 18	715		CMPA #CAN	? ABORT SIGNALLED FROM WAIT
FA6A 27 F4	716		BEQ CHKSEC	BRANCH YES
FA6C 4F	717		CLRA	SET C=0 FOR NO ABORT
FA6D 39	718		RTS	AND RETURN
FA6E	719			
FA6E	720	*	SAVE MEMORY WITH JUMPS	
FA6E 6E 9D E5 78	721	XQPAUS	JMP [VECTAB+.PAUSE,PCR]	TO PAUSE ROUTINE
FA72 AD 9D E5 62	722	XQCIDT	JSR [VECTAB+.CIDTA,PCR]	TO INPUT ROUTINE
FA76 84 7F	723		ANDA #\$7F	STRIP PARITY
FA78 39	724		RTS	RETURN TO CALLER
FA79	725			
FA79	726	*****	*****	*****
FA79	727	*	NMI DEFAULT INTERRUPT HANDLER	
FA79	728	*	THE NMI HANDLER IS USED FOR TRACING INSTRUCTIONS.	
FA79	729	*	TRACE PRINTOUTS OCCUR ONLY AS LONG AS THE STACK	
FA79	730	*	TRACE LEVEL IS NOT BREACHED BY FALLING BELOW IT.	
FA79	731	*	TRACING CONTINUES UNTIL THE COUNT TURNS ZERO OR	
FA79	732	*	A CTL-X IS ENTERED FROM THE INPUT CONSOLE DEVICE.	
FA79	733	*****	*****	*****
FA79	734			
FA79 4F 50 2D 04	735	MSHOWP	FCB 'O','P','-',EOT OPCODE PREP	
FA7D	736			
FA7D 8D 42	737	NMIR	BSR LDDP	LOAD PAGE AND VERIFY STACK
FA7F 0D 8F	738		TST MISFLG	? THRU A BREAKPOINT
FA81 26 34	739		BNE NMICON	BRANCH IF SO TO CONTINUE
FA83 0D 90	740		TST SWICNT	? INHIBIT "SWI" DURING TRACE
FA85 2B 29	741		BMI NMITRC	BRANCH YES
FA87 30 6C	742		LEAX 12,S	OBTAIN USERS STACK POINTER
FA89 9C F8	743		CMPX SLEVEL	? TO TRACE HERE
FA8B 25 23	744		BLO NMITRC	BRANCH IF TOO LOW TO DISPLAY
FA8D 30 8C E9	745		LEAX MSHOWP,PCR	LOAD OP PREP
FA90 3F	746		SWI	SEND TO CONSOLE
FA91 02	747		FCB PDATA1	FUNCTION
FA92 09 8E	748		ROL DELIM	SAVE CARRY BIT
FA94 30 8D E5 01	749		LEAX LASTOP,PCR	POINT TO LAST OP
FA98 3F	750		SWI	SEND OUT AS HEX
FA99 05	751		FCB OUT4HS	FUNCTION
FA9A 8D 17	752		BSR REGPRS	FOLLOW MEMORY WITH REGISTERS
FA9C 25 37	753		BCS ZBKCMD	BRANCH IF "CANCEL"
FA9E 06 8E	754		ROR DELIM	RESTORE CARRY BIT
FAA0 25 33	755		BCS ZBKCMD	BRANCH IF "CANCEL"
FAA2 9E 91	756		LDX TRACEC	LOAD TRACE COUNT
FAA4 27 2F	757		BEQ ZBKCMD	IF ZERO TO COMMAND HANDLER
FAA6 30 1F	758		LEAX -1,X	MINUS ONE
FAA8 9F 91	759		STX TRACEC	REFRESH
FAAA 27 29	760		BEQ ZBKCMD	STOP TRACE WHEN ZERO
FAAC 8D AA	761		BSR CHKABT	? ABORT THE TRACE
FAAE 25 25	762		BCS ZBKCMD	BRANCH YES TO COMMAND HANDLER
FAB0 16 03 F7	763	NMITRC	LBRA CTRCE3	NO, TRACE ANOTHER INSTRUCTION
FAB3	764			
FAB3 17 01 B9	765	REGPRS	LBSR REGPRT	PRINT REGISTERS AS FROM COMMAND
FAB6 39	766		RTS	RETURN TO CALLER
FAB7	767			
FAB7	768	*	JUST EXECUTED THRU A BRKPNT.	NOW CONTINUE NORMALLY

```

FAB7 0F 8F          769 NMICON CLR    MISFLG      CLEAR THRU FLAG
FAB9 17 02 EB       770 LBSR   ARMBK2     ARM BREAKPOINTS
FABC 3B            771 RTI    RTI    AND      CONTINUE USERS PROGRAM
FABD
FABD             772
FABD             773 * LDDP - SETUP DIRECT PAGE REGISTER, VERIFY STACK.
FABD             774 * AN INVALID STACK CAUSES A RETURN TO THE COMMAND
FABD             775 * HANDLER.
FABD             776 * INPUT: FULLY STACKED REGISTERS FROM AN INTERRUPT
FABD             777 * OUTPUT: DPR LOADED TO WORK PAGE
FABD             778
FABD 3F 07 20 04   779 ERRMSG FCB    '?',BELL,$20,EOT ERROR RESPONSE
FAC1
FAC1 E6 8D E4 D8  780
FAC5 1F 9B          781 LDDP   LDB    BASEPG,PCR  LOAD DIRECT PAGE HIGH BYTE
FAC7 A1 63          782 TFR    B,DP   ? IS STACK VALID
FAC9 27 25          783 CMPA   3,S    YES, RETURN
FACB 10 DE 97       784 BEQ    RTS    RESET TO INITIAL STACK POINTER
FACE 30 8C EC       785 LDS    RSTACK LOAD ERROR REPORT
FAD1 3F             786 ERROR  LEAX   ERRMSG,PCR SEND OUT BEFORE REGISTERS
FAD2 03             787 SWI    FCB    PDATA ON NEXT LINE
FAD3
FAD3             788
FAD3             789 * FALL INTO BREAKPOINT HANDLER
FAD3             790
FAD3             791 ****
FAD3             792 * [SWI FUNCTION 10]
FAD3             793 * BREAKPOINT PROGRAM FUNCTION
FAD3             794 * PRINT REGISTERS AND GO TO COMMAND HANLER
FAD3             795 ****
FAD3 8D DE          796 ZBKPNTR BSR    REGPRS  PRINT OUT REGISTERS
FAD5 16 FE 21       797 ZBKCMD  LBRA   CMDNEP NOW ENTER COMMAND HANDLER
FAD8
FAD8             798 ****
FAD8             799 *
FAD8             800 * IRQ, RESERVED, SWI2 AND SWI3 INTERRUPT HANDLERS
FAD8             801 * THE DEFAULT HANDLING IS TO CAUSE A BREAKPOINT.
FAD8             802 ****
FAD8             803 SWI2R  EQU    *      SWI2 ENTRY
FAD8             804 SWI3R  EQU    *      SWI3 ENTRY
FAD8             805 IRQR   EQU    *      IRQ ENTRY
FAD8 8D E7          806 RSRVDR BSR    LDDP   SET BASE PAGE, VALIDATE STACK
FADA 20 F7          807 BRA    ZBKPNTR FORCE A BREAKPOINT
FADC
FADC             808
FADC             809 ****
FADC             810 * FIRQ HANDLER
FADC             811 * JUST RETURN FOR THE FIRQ INTERRUPT
FADC             812 ****
FABC             813 FIRQR  EQU    RTI    IMMEDIATE RETURN
FADC             814
FADC             815 ****
FADC             816 * DEFAULT I/O DRIVERS
FADC             817 ****
FADC             818
FADC             819 * CIDTA - RETURN CONSOLE INPUT CHARACTER
FADC             820 * OUTPUT: C=0 IF NO DATA READY, C=1 A=CHARACTER
FADC             821 * U VOLATILE
FADC DE F0          822 CIDTA  LDU    VECTAB+.ACIA LOAD ACIA ADDRESS
FADE A6 C4          823 LDA    ,U      LOAD STATUS REGISTER
FAE0 44            824 LSRA   CIRTN TEST RECEIVER REGISTER FLAG
FAE1 24 02          825 BCC    CIRTN RETURN IF NOTHING
FAE3 A6 41          826 LDA    1,U    LOAD DATA BYTE
FAE5 39            827 CIRTN RTS    RETURN TO CALLER
FAE6
FAE6             828
FAE6             829 * CION - INPUT CONSOLE INITIALIZATION
FAE6             830 * COON - OUTPUT CONSOLE INITIALIZATION
FAE6             831 * A,X VOLATILE
FAE6             832 CION   EQU    *
FAE6             833 COON  LDA    #3    RESET ACIA CODE
FAE8 9E F0          834 LDX    VECTAB+.ACIA LOAD ACIA ADDRESS
FAEA A7 84          835 STA    ,X    STORE INTO STATUS REGISTER
FAEC 86 51          836 LDA    #$51 SET CONTROL

```

FAEE	A7	84	837	STA	,X	REGISTER UP
FAF0	39		838	RTS	RTS	RETURN TO CALLER
FAF1			839			
FAF1			840	* THE FOLLOWING HAVE NO DUTIES TO PERFORM		
FAF0			841	CIOFF EQU	RTS	CONSOLE INPUT OFF
FAF0			842	COOFF EQU	RTS	CONSOLE OUTPUT OFF
FAF1			843			
FAF1			844	* CODTA - OUTPUT CHARACTER TO CONSOLE DEVICE		
FAF1			845	* INPUT: A=CHARACTER TO SEND		
FAF1			846	* OUTPUT: CHAR SENT TO TERMINAL WITH PROPER PADDING		
FAF1			847	* ALL REGISTERS TRANSPARENT		
FAF1			848			
FAF1	34	47	849	CODTA PSHS	U,D,CC	SAVE REGISTERS, WORK BYTE
FAF3	DE	F0	850	LDU	VECTAB+.ACIA	ADDRESS ACIA
FAF5	8D	1B	851	BSR	CODTAO	CALL OUTPUT CHAR SUBROUTINE
FAF7	81	10	852	CMPA	#DLE	? DATA LINE ESCAPE
FAF9	27	12	853	BEQ	CODTRT	YES, RETURN
FAFB	D6	F2	854	LDB	VECTAB+.PAD	DEFAULT TO CHAR PAD COUNT
FAFD	81	0D	855	CMPA	#CR	? CR
FAFF	26	02	856	BNE	CODTPD	BRANCH NO
FB01	D6	F3	857	LDB	VECTAB+.PAD+1	LOAD NEW LINE PAD COUNT
FB03	4F		858	CODTPD CLRA		CREATE NULL
FB04	E7	E4	859	STB	,S	SAVE COUNT
FB06	8C		860	FCB	SKIP2	ENTER LOOP
FB07	8D	09	861	CODTLR BSR	CODTAO	SEND NULL
FB09	6A	E4	862	DEC	,S	? FINISHED
FB0B	2A	FA	863	BPL	CODTLR	NO, CONTINUE WITH MORE
FB0D	35	C7	864	CODTRT PULS	PC,U,D,CC	RESTORE REGISTERS AND RETURN
FB0F			865			
FB0F	17	FF	866	CODTAD LBSR	XQPAUS	TEMPORARY GIVE UP CONTROL
FB12	E6	C4	867	CODTAO LDB	,U	LOAD ACIA CONTROL REGISTER
FB14	C5	02	868	BITB	#\$02	? TX REGISTER CLEAR >LSAB FIXME
FB16	26	F7	869	BNE	CODTAD	RELEASE CONTROL IF NOT
FB18	A7	41	870	STA	1,U	STORE INTO DATA REGISTER
FB1A	39		871	RTS		RETURN TO CALLER
FB1B			872	*E		
FB1B			873			
FB1B			874	* BSON - TURN ON READ/VERIFY/PUNCH MECHANISM		
FB1B			875	* A IS VOLATILE		
FB1B			876			
FB1B	86	11	877	BSON LDA	#\$11	SET READ CODE
FB1D	6D	66	878	TST	6,S	? READ OR VERIFY
FB1F	26	01	879	BNE	BSON2	BRANCH YES
FB21	4C		880	INCA		SET TO WRITE
FB22	3F		881	BSON2 SWI		PERFORM OUTPUT
FB23	01		882	FCB	OUTCH	FUNCTION
FB24	0C	8F	883	INC	MISFLG	SET LOAD IN PROGRESS FLAG
FB26	39		884	RTS		RETURN TO CALLER
FB27			885			
FB27			886	* BSOFF - TURN OFF READ/VERIFY/PUNCH MECHANISM		
FB27			887	* A,X VOLATILE		
FB27	86	14	888	BSOFF LDA	#\$14	TO DC4 - STOP
FB29	3F		889	SWI		SEND OUT
FB2A	01		890	FCB	OUTCH	FUNCTION
FB2B	4A		891	DECA		CHANGE TO DC3 (X-OFF)
FB2C	3F		892	SWI		SEND OUT
FB2D	01		893	FCB	OUTCH	FUNCTION
FB2E	0A	8F	894	DEC	MISFLG	CLEAR LOAD IN PROGRESS FLAG
FB30	8E	61	A8	895	LDX #25000	DELAY 1 SECOND (2MHZ CLOCK)
FB33	30	1F	896	BSOFLP LEAX	-1,X	COUNT DOWN
FB35	26	FC	897	BNE	BSOFLP	LOOP TILL DONE
FB37	39		898	RTS		RETURN TO CALLER
FB38			899			
FB38			900	* BSDTA - READ/VERIFY/PUNCH HANDLER		
FB38			901	* INPUT: S+6=CODE BYTE, VERIFY(-1), PUNCH(0), LOAD(1)		
FB38			902	*	S+4=START ADDRESS	
FB38			903	*	S+2=STOP ADDRESS	
FB38			904	*	S+0=RETURN ADDRESS	

FB38		905	* OUTPUT: Z=1 NORMAL COMPLETION, Z=0 INVALID LOAD/VER			
FB38		906	* REGISTERS ARE VOLATILE			
FB38		907				
FB38 EE 62		908	BSDTA	LDU	2,S	U=TO ADDRESS OR OFFSET
FB3A 6D 66		909	TST	6,S		? PUNCH
FB3C 27 54		910	BEQ	BSDPUN		BRANCH YES
FB3E		911	* DURING READ/VERIFY: S+2=MSB ADDRESS SAVE BYTE			
FB3E		912	*		S+1=BYTE COUNTER	
FB3E		913	*		S+0=CHECKSUM	
FB3E		914	*		U HOLDS OFFSET	
FB3E 32 7D		915	LEAS	-3,S	ROOM FOR WORK/COUNTER/CHECKSUM	
FB40 3F		916	BSDLD1	SWI	GET NEXT CHARACTER	
FB41 00		917	FCB	INCHNP	FUNCTION	
FB42 81 53		918	BSDLD2	CMPA	#'S'	? START OF S1/S9
FB44 26 FA		919	BNE	BSDLD1	BRANCH NOT	
FB46 3F		920	SWI		GET NEXT CHARACTER	
FB47 00		921	FCB	INCHNP	FUNCTION	
FB48 81 39		922	CMPA	#'9'	? HAVE S9	
FB4A 27 22		923	BEQ	BSDSRT	YES, RETURN GOOD CODE	
FB4C 81 31		924	CMPA	#'1'	? HAVE NEW RECORD	
FB4E 26 F2		925	BNE	BSDLD2	BRANCH IF NOT	
FB50 6F E4		926	CLR	,S	CLEAR CHECKSUM	
FB52 8D 21		927	BSR	BYTE	OBTAIN BYTE COUNT	
FB54 E7 61		928	STB	1,S	SAVE FOR DECREMENT	
FB56		929	* READ ADDRESS			
FB56 8D 1D		930	BSR	BYTE	OBTAIN HIGH VALUE	
FB58 E7 62		931	STB	2,S	SAVE IT	
FB5A 8D 19		932	BSR	BYTE	OBTAIN LOW VALUE	
FB5C A6 62		933	LDA	2,S	MAKE D=VALUE	
FB5E 31 CB		934	LEAY	D,U	Y=ADDRESS+OFFSET	
FB60		935	* STORE TEXT			
FB60 8D 13		936	BSDNXT	BSR	BYTE	NEXT BYTE
FB62 27 0C		937	BEQ	BSDEOL	BRANCH IF CHECKSUM	
FB64 6D 69		938	TST	9,S	? VERIFY ONLY	
FB66 2B 02		939	BMI	BSDCMP	YES, ONLY COMPARE	
FB68 E7 A4		940	STB	,Y	STORE INTO MEMORY	
FB6A E1 A0		941	BSDCMP	CMPB	,Y+	? VALID RAM
FB6C 27 F2		942	BEQ	BSDNXT	YES, CONTINUE READING	
FB6E 35 92		943	BSDSRT	PULS	PC,X,A	RETURN WITH Z SET PROPER
FB70		944				
FB70 4C		945	BSDEOL	INCA		? VALID CHECKSUM
FB71 27 CD		946	BEQ	BSDLD1	BRANCH YES	
FB73 20 F9		947	BRA	BSDSRT	RETURN Z=0 INVALID	
FB75		948				
FB75 8D 12		949	* BYTE BUILDS 8 BIT VALUE FROM TWO HEX DIGITS IN			
FB75		950	BYTE	BSR	BYTHEX	OBTAINT FIRST HEX
FB77 C6 10		951	LDB	#16	PREPARE SHIFT	
FB79 3D		952	MUL		OVER TO A	
FB7A 8D 0D		953	BSR	BYTHEX	OBTAINT SECOND HEX	
FB7C 34 04		954	PSHS	B	SAVE HIGH HEX	
FB7E AB E0		955	ADDA	,S+	COMBINE BOTH SIDES	
FB80 1F 89		956	TFR	A,B	SEND BACK IN B	
FB82 AB 62		957	ADDA	2,S	COMPUTE NEW CHECKSUM	
FB84 A7 62		958	STA	2,S	STORE BACK	
FB86 6A 63		959	DEC	3,S	DECREMENT BYTE COUNT	
FB88 39		960	BYTRTS	RTS	RETURN TO CALLER	
FB89		961				
FB89 3F		962	BYTHEX	SWI	GET NEXT HEX	
FB8A 00		963	FCB	INCHNP	CHARACTER	
FB8B 17 01 D4		964	LBSR	CNVHEX	CONVERT TO HEX	
FB8E 27 F8		965	BEQ	BYTRTS	RETURN IF VALID HEX	
FB90 35 F2		966	PULS	PC,U,Y,X,A	RETURN TO CALLER WITH Z=0	
FB92		967				
FB92		968	* PUNCH STACK USE: S+8=TO ADDRESS			
FB92		969	*		S+6=RETURN ADDRESS	
FB92		970	*		S+4=SAVED PADDING VALUES	
FB92		971	*		S+2 FROM ADDRESS	
FB92		972	*		S+1=FRAME COUNT/CHECKSUM	

FB92		973	*	S+0=BYTE COUNT
FB92	DE F2	974	BSDPUN	LDU VECTAB+.PAD LOAD PADDING VALUES
FB94	AE 64	975	LDX 4,S X=FROM ADDRESS	
FB96	34 56	976	PSHS U,X,D CREATE STACK WORK AREA	
FB98	CC 00 18	977	LDD #24 SET A=0, B=24	
FB9B	D7 F2	978	STB VECTAB+.PAD SETUP 24 CHARACTER PADS	
FB9D	3F	979	SWI SEND NULLS OUT	
FB9E	01	980	FCB OUTCH FUNCTION	
FB9F	C6 04	981	LDB #4 SETUP NEW LINE PAD TO 4	
FBA1	DD F2	982	STD VECTAB+.PAD SETUP PUNCH PADDING	
FBA3		983	* CALCULATE SIZE	
FBA3	EC 68	984	BSPGO LDD 8,S LOAD TO	
FBA5	A3 62	985	SUBD 2,S MINUS FROM=LENGTH	
FBA7	10 83 00 18	986	CMPD #24 ? MORE THAN 23	
FBAB	25 02	987	BLO BSPOK NO, OK	
FBAD	C6 17	988	LDB #23 FORCE TO 23 MAX	
FBAF	5C	989	BSPOK INCB PREPARE COUNTER	
FBB0	E7 E4	990	STB ,S STORE BYTE COUNT	
FBB2	CB 03	991	ADD8 #3 ADJUST TO FRAME COUNT	
FBB4	E7 61	992	STB 1,S SAVE	
FBB6		993	*PUNCH CR,LF,NULS,S,1	
FBB6	30 8C 33	994	LEAX <BSPSTR,PCR LOAD START RECORD HEADER	
FBB9	3F	995	SWI SEND OUT	
FBBB	03	996	FCB PDATA FUNCTION	
FBBB		997	* SEND FRAME COUNT	
FBBB	5F	998	CLRB INITIALIZE CHECKSUM	
FBBC	30 61	999	LEAX 1,S POINT TO FRAME COUNT AND ADDR	
FBBE	8D 27	1000	BSR BSPUN2 SEND FRAME COUNT	
FBC0		1001	*DATA ADDRESS	
FBC0	8D 25	1002	BSR BSPUN2 SEND ADDRESS HI	
FBC2	8D 23	1003	BSR BSPUN2 SEND ADDRESS LOW	
FBC4		1004	*PUNCH DATA	
FBC4	AE 62	1005	LDX 2,S LOAD START DATA ADDRESS	
FBC6	8D 1F	1006	BSPMRE BSR BSPUN2 SEND OUT NEXT BYTE	
FBC8	6A E4	1007	DEC ,S ? FINAL BYTE	
FBCA	26 FA	1008	BNE BSPMRE LOOP IF NOT DONE	
FBCC	AF 62	1009	STX 2,S UPDATE FROM ADDRESS VALUE	
F BCE		1010	*PUNCH CHECKSUM	
F BCE	53	1011	COMB COMPLEMENT	
FBCF	E7 61	1012	STB 1,S STORE FOR SENDOUT	
FBD1	30 61	1013	LEAX 1,S POINT TO IT	
FBD3	8D 14	1014	BSR BSPUNC SEND OUT AS HEX	
FBD5	AE 68	1015	LDX 8,S LOAD TOP ADDRESS	
FBD7	AC 62	1016	CMPX 2,S ? DONE	
FBD9	24 C8	1017	BHS BSPGO BRANCH NOT	
FBDB	30 8C 11	1018	LEAX <BSPEOF,PCR PREPARE END OF FILE	
FBDE	3F	1019	SWI SEND OUT STRING	
FBDF	03	1020	FCB PDATA FUNCTION	
FBE0	EC 64	1021	LDD 4,S RECOVER PAD COUNTS	
FBE2	DD F2	1022	STD VECTAB+.PAD RESTORE	
FBE4	4F	1023	CLRA SET Z=1 FOR OK RETURN	
FBE5	35 D6	1024	PULS PC,U,X,D RETURN WITH OK CODE	
FBE7		1025		
FBE7	EB 84	1026	BSPUN2 ADDB ,X ADD TO CHECKSUM	
FBE9	16 FD ED	1027	BSPUNC LBRA ZOUT2H SEND OUT AS HEX AND RETURN	
F BEC		1028		
F BEC	53 31 04	1029	BSPSTR FCB 'S','1',EOT CR,LF,NULS,S,1	
F BEF	53 39 30 33 30 30 +	1030	BSPEOF FCC /S9030000FC/ EOF STRING	
F BF9	0D 0A 04	1031	FCB CR,LF,EOT	
FBFC		1032		
FBFC		1033	* HSDTA - HIGH SPEED PRINT MEMORY	
FBFC		1034	* INPUT: S+4=START ADDRESS	
FBFC		1035	* S+2=STOP ADDRESS	
FBFC		1036	* S+0=RETURN ADDRESS	
FBFC		1037	* X,D VOLATILE	
FBFC		1038		
FBFC		1039	* SEND TITLE	
FBFC	3F	1040	HSDTA SWI SEND NEW LINE	

FBFD 06	1041	FCB	PCRLF	FUNCTION
FBFE C6 06	1042	LDB	#6	PREPARE 6 SPACES
FC00 3F	1043	HSBLNK	SWI	SEND BLANK
FC01 07	1044	FCB	SPACE	FUNCTION
FC02 5A	1045	DEC B		COUNT DOWN
FC03 26 FB	1046	BNE	HSBLNK	LOOP IF MORE
FC05 5F	1047	CLRB		SETUP BYTE COUNT
FC06 1F 98	1048	HSHTTL	TFR B,A	PREPARE FOR CONVERT
FC08 17 FD DB	1049	LBSR	ZOUTHX	CONVERT TO A HEX DIGIT
FC0B 3F	1050	SWI		SEND BLANK
FC0C 07	1051	FCB	SPACE	FUNCTION
FC0D 3F	1052	SWI		SEND ANOTHER
FC0E 07	1053	FCB	SPACE	BLANK
FC0F 5C	1054	INC B		UP ANOTHER
FC10 C1 10	1055	CMPB	#\$10	? PAST 'F'
FC12 25 F2	1056	BLO	HSHTTL	LOOP UNTIL SO
FC14 3F	1057	HSHLNE	SWI	TO NEXT LINE
FC15 06	1058	FCB	PCRLF	FUNCTION
FC16 25 2F	1059	BCS	HSDRTN	RETURN IF USER ENTERED CTL-X
FC18 30 64	1060	LEAX	4,S	POINT AT ADDRESS TO CONVERT
FC1A 3F	1061	SWI		PRINT OUT ADDRESS
FC1B 05	1062	FCB	OUT4HS	FUNCTION
FC1C AE 64	1063	LDX	4,S	LOAD ADDRESS PROPER
FC1E C6 10	1064	LDB	#16	NEXT SIXTEEN
FC20 3F	1065	HSHNXT	SWI	CONVERT BYTE TO HEX AND SEND
FC21 04	1066	FCB	OUT2HS	FUNCTION
FC22 5A	1067	DEC B		COUNT DOWN
FC23 26 FB	1068	BNE	HSHNXT	LOOP IF NOT SIXTEENTH
FC25 3F	1069	SWI		SEND BLANK
FC26 07	1070	FCB	SPACE	FUNCTION
FC27 AE 64	1071	LDX	4,S	RELOAD FROM ADDRESS
FC29 C6 10	1072	LDB	#16	COUNT
FC2B A6 80	1073	HSHCHR	LDA ,X+	NEXT BYTE
FC2D 2B 04	1074	BMI	HSHDOT	TOO LARGE, TO A DOT
FC2F 81 20	1075	CMPA	' '	? LOWER THAN A BLANK
FC31 24 02	1076	BHS	HSHCOK	NO, BRANCH OK
FC33 86 2E	1077	HSHDOT	LDA #'. '	CONVERT INVALID TO A BLANK
FC35 3F	1078	HSHCOK	SWI	SEND CHARACTER
FC36 01	1079	FCB	OUTCH	FUNCTION
FC37 5A	1080	DEC B		? DONE
FC38 26 F1	1081	BNE	HSHCHR	BRANCH NO
FC3A AC 62	1082	CMPX	2,S	? PAST LAST ADDRESS
FC3C 24 09	1083	BHS	HSDRTN	QUIT IF SO
FC3E AF 64	1084	STX	4,S	UPDATE FROM ADDRESS
FC40 A6 65	1085	LDA	5,S	LOAD LOW BYTE ADDRESS
FC42 48	1086	ASLA	?	TO SECTION BOUNDARY
FC43 26 CF	1087	BNE	HSHLNE	BRANCH IF NOT
FC45 20 B5	1088	BRA	HSDTA	BRANCH IF SO
FC47 3F	1089	HSDRTN	SWI	SEND NEW LINE
FC48 06	1090	FCB	PCRLF	FUNCTION
FC49 39	1091		RTS	RETURN TO CALLER
FC4A	1092	*		
FC4A	1093			
FC4A	1094	*****ASSIST09 COMMANDS*****		
FC4A	1095	*	A S S I S T 0 9	C O M M A N D S
FC4A	1096	*	*****	
FC4A	1097			
FC4A	1098	*****REGISTERS - DISPLAY AND CHANGE REGISTERS		
FC4A 8D 23	1099	CREG	BSR	PRINT REGISTERS
FC4C 4C	1100	INCA		SET FOR CHANGE FUNCTION
FC4D 8D 21	1101	BSR	REGCHG	GO CHANGE, DISPLAY REGISTERS
FC4F 39	1102		RTS	RETURN TO COMMAND PROCESSOR
FC50	1103			
FC50	1104	*****		
FC50	1105	*	REGPRT - PRINT/CHANGE REGISTERS SUBROUTINE	
FC50	1106	*	* WILL ABORT TO 'CMDBAD' IF OVERFLOW DETECTED DURING	
FC50	1107	*	* A CHANGE OPERATION. CHANGE DISPLAYS REGISTERS WHEN	
FC50	1108	*	* DONE.	

```

FC50          1109 * REGISTER MASK LIST CONSISTS OF:
FC50          1110 * A) CHARACTERS DENOTING REGISTER
FC50          1111 * B) ZERO FOR ONE BYTE, -1 FOR TWO
FC50          1112 * C) OFFSET ON STACK TO REGISTER POSITION
FC50          1113 * INPUT: SP+4=STACKED REGISTERS
FC50          1114 *           A=0 PRINT, A#0 PRINT AND CHANGE
FC50          1115 * OUTPUT: (ONLY FOR REGISTER DISPLAY)
FC50          1116 *           C=1 CONTROL-X ENTERED, C=0 OTHERWISE
FC50          1117 * VOLATILE: D,X (CHANGE)
FC50          1118 *           B,X (DISPLAY)
FC50          1119 ****
FC50 50 43 FF 13 1120 REGMSK FCB 'P','C',-1,19 PC REG
FC54 41 00 0A 1121 FCB 'A',0,10 A REG
FC57 42 00 0B 1122 FCB 'B',0,11 B REG
FC5A 58 FF 0D 1123 FCB 'X',-1,13 X REG
FC5D 59 FF 0F 1124 FCB 'Y',-1,15 Y REG
FC60 55 FF 11 1125 FCB 'U',-1,17 U REG
FC63 53 FF 01 1126 FCB 'S',-1,1 S REG
FC66 43 43 00 09 1127 FCB 'C','C',0,9 CC REG
FC6A 44 50 00 0C 1128 FCB 'D','P',0,12 DP REG
FC6E 00 1129 FCB 0 END OF LIST
FC6F 1130
FC6F 4F 1131 REGPRT CLRA
FC70 30 E8 10 1132 REGCHG LEAX 4+12,S SETUP PRINT ONLY FLAG
FC73 34 32 1133 PSHS Y,X,A READY STACK VALUE
FC75 31 8C D8 1134 LEAY REGMSK,PCR SAVE ON STACK WITH OPTION
FC78 EC A0 1135 REGP1 LDD ,Y+ LOAD REGISTER MASK
FC7A 4D 1136 TSTA ? LOAD NEXT CHAR OR <=0
FC7B 2F 04 1137 BLE REGP2 END OF CHARACTERS
FC7D 3F 1138 SWI BRANCH NOT CHARACTER
FC7E 01 1139 FCB OUTCH SEND TO CONSOLE
FC7F 20 F7 1140 BRA REGP1 FUNCTION BYTE
FC81 86 2D 1141 REGP2 LDA #'-' CHECK NEXT
FC83 3F 1142 SWI READY '-'
FC84 01 1143 FCB OUTCH SEND OUT
FC85 30 E5 1144 LEAX B,S WITH OUTCH
FC87 6D E4 1145 TST ,S X->REGISTER TO PRINT
FC89 26 12 1146 BNE REGCNG ? CHANGE OPTION
FC8B 6D 3F 1147 TST -1,Y BRANCH YES
FC8D 27 03 1148 BEQ REGP3 ? ONE OR TWO BYTES
FC8F 3F 1149 SWI BRANCH ZERO MEANS ONE
FC90 05 1150 FCB OUT4HS PERFORM WORD HEX
FC91 8C 1151 FCB SKIP2 FUNCTION
FC92 3F 1152 REGP3 SWI SKIP BYTE PRINT
FC93 04 1153 FCB OUT2HS PERFORM BYTE HEX
FC94 EC A0 1154 REG4 LDD ,Y+ FUNCTION
FC96 5D 1155 TSTB ? TO FRONT OF NEXT ENTRY
FC97 26 DF 1156 BNE REGP1 END OF ENTRIES
FC99 3F 1157 SWI LOOP IF MORE
FC9A 06 1158 FCB PCRLF FORCE NEW LINE
FC9B 35 B2 1159 REGRTN PULS PC,Y,X,A FUNCTION
FC9D 1160 RESTORE STACK AND RETURN
FC9D 8D 40 1161 REGCNG BSR BLDNNB INPUT BINARY NUMBER
FC9F 27 10 1162 BEQ REGNXC IF CHANGE THEN JUMP
FCA1 81 0D 1163 CMPA #CR ? NO MORE DESIRED
FCA3 27 1E 1164 BEQ REGAGN BRANCH NOPE
FCA5 E6 3F 1165 LDB -1,Y LOAD SIZE FLAG
FCA7 5A 1166 DECB MINUS ONE
FCA8 50 1167 NEGB MAKE POSITIVE
FCA9 58 1168 ASLB TIMES TWO (=2 OR =4)
FCAA 3F 1169 REGSKP SWI PERFORM SPACES
FCAB 07 1170 FCB SPACE FUNCTION
FCAC 5A 1171 DECB
FCAD 26 FB 1172 BNE REGSKP LOOP IF MORE
FCAF 20 E3 1173 BRA REG4 CONTINUE WITH NEXT REGISTER
FCB1 A7 E4 1174 REGNXC STA ,S SAVE DELIMITER IN OPTION
FCB3 1175 * (ALWAYS > 0)
FCB3 DC 9B 1176 LDD NUMBER OBTAIN BINARY RESULT

```

```

FCB5  6D 3F          1177    TST    -1,Y      ? TWO BYTES WORTH
FCB7  26 02          1178    BNE    REGTWO   BRANCH YES
FCB9  A6 82          1179    LDA    ,-X      SETUP FOR TWO
FCBB  ED 84          1180    REGTWO  STD    ,X      STORE IN NEW VALUE
FCBD  A6 E4          1181    LDA    ,S       RECOVER DELIMITER
FCBF  81 0D          1182    CMPA   #CR     ? END OF CHANGES
FCC1  26 D1          1183    BNE    REG4     NO, KEEP ON TRUCK'N
FCC3
FCC3  30 8D E2 8A    1184    * MOVE STACKED DATA TO NEW STACK IN CASE STACK
FCC3  30 8D E2 8A    1185    * POINTER HAS CHANGED
FCC3  30 8D E2 8A    1186    REGAGN  LEAX   TSTACK,PCR  LOAD TEMP AREA
FCC7  C6 15          1187    LDB    #21     LOAD COUNT
FCC9  35 02          1188    REGTF1  PULS   A       NEXT BYTE
FCCB  A7 80          1189    STA    ,X+     STORE INTO TEMP
FCCD  5A             1190    DECB   -
FCCE  26 F9          1191    BNE    REGTF1  COUNT DOWN
FCD0  10 EE 88 EC    1192    LDS    -20,X   LOOP IF MORE
FCD4  C6 15          1193    LDB    #21     LOAD NEW STACK POINTER
FCD6  A6 82          1194    REGTF2  LDA    ,-X   LOAD COUNT AGAIN
FCD8  34 02          1195    PSHS   A       NEXT TO STORE
FCDA  5A             1196    DECB   -
FCDB  26 F9          1197    BNE    REGTF2  BACK ONTO NEW STACK
FCDD  20 BC          1198    BRA    REGRTN  COUNT DOWN
FCDF
FCDF  1199          1200    ****
FCDF  1200    * BLDNUM - BUILDS BINARY VALUE FROM INPUT HEX
FCDF  1201    * THE ACTIVE EXPRESSION HANDLER IS USED.
FCDF  1202    * INPUT: S=RETURN ADDRESS
FCDF  1203    * OUTPUT: A=DELIMITER WHICH TERMINATED VALUE
FCDF  1204    * (IF DELM NOT ZERO)
FCDF  1205    * "NUMBER"=WORD BINARY RESULT
FCDF  1206    * Z=1 IF INPUT RECIEVED, Z=0 IF NO HEX RECIEVED
FCDF  1207    * REGISTERS ARE TRANSPARENT
FCDF  1208    ****
FCDF  1209
FCDF  1210
FCDF  1211    * EXECUTE SINGLE OR EXTENDED ROM EXPRESSION HANDLER
FCDF  1212    *
FCDF  1213    * THE FLAG "DELIM" IS USED AS FOLLOWS:
FCDF  1214    * DELIM=0 NO LEADING BLANKS, NO FORCED TERMINATOR
FCDF  1215    * DELIM=CHR ACCEPT LEADING 'CHR'S, FORCED TERMINATOR
FCDF  4F             1216    BLDNNB  CLRA   NO DYNAMIC DELIMITER
FCE0  8C             1217    FCB    SKIP2  SKIP NEXT INSTRUCTION
FCE1
FCE1  86 20          1218    * BUILD WITH LEADING BLANKS
FCE3  97 8E          1219    BLDNUM LDA   #' ' ALLOW LEADING BLANKS
FCE5  6E 9D E3 03    1220    STA    DELIM  STORE AS DELIMITER
FCE9
FCE9  1221    JMP    [VECTAB+.EXPAN,PCR] TO EXP ANALYZER
FCE9  1222
FCE9  1223    * THIS IS THE DEFAULT SINGLE ROM ANALYZER. WE ACCEPT:
FCE9  1224    * 1) HEX INPUT
FCE9  1225    * 2) 'M' FOR LAST MEMORY EXAMINE ADDRESS
FCE9  1226    * 3) 'P' FOR PROGRAM COUNTER ADDRESS
FCE9  1227    * 4) 'W' FOR WINDOW VALUE
FCE9  1228    * 5) '@' FOR INDIRECT VALUE
FCE9  34 14          1229    EXP1    PSHS   X,B   SAVE REGISTERS
FCEB  8D 5C          1230    EXPDLM  BSR    BLDHXI  CLEAR NUMBER, CHECK FIRST CHAR
FCED  27 18          1231    BEQ    EXP2   IF HEX DIGIT CONTINUE BUILDING
FCEF
FCEF  91 8E          1232    * SKIP BLANKS IF DESIRED
FCEF  1233    CMPA   DELIM  ? CORRECT DELIMITER
FCF1  27 F8          1234    BEQ    EXPDLM  YES, IGNORE IT
FCF3
FCF3  9E 9E          1235    * TEST FOR M OR P
FCF3  1236    LDX    ADDR   DEFAULT FOR 'M'
FCF5  81 4D          1237    CMPA   #'M'  ? MEMORY EXAMINE ADDR WANTED
FCF7  27 16          1238    BEQ    EXPTDL BRANCH IF SO
FCF9  9E 93          1239    LDX    PCNTER DEFAULT FOR 'P'
FCFB  81 50          1240    CMPA   #'P'  ? LAST PROGRAM COUNTER WANTED
FCFD  27 10          1241    BEQ    EXPTDL BRANCH IF SO
FCFF  9E A0          1242    LDX    WINDOW DEFAULT TO WINDOW
FD01  81 57          1243    CMPA   #'W'  ? WINDOW WANTED
FD03  27 0A          1244    BEQ    EXPTDL

```

FD05 35 94	1245 EXPRTN	PULS	PC,X,B	RETURN AND RESTORE REGISTERS
FD07	1246 *	GOT HEX, NOW	CONTINUE BUILDING	
FD07 8D 44	1247 EXP2	BSR	BLDHEX	COMPUTE NEXT DIGIT
FD09 27 FC	1248	BEQ	EXP2	CONTINUE IF MORE
FD0B 20 0A	1249	BRA	EXPCDL	SEARCH FOR +/-
FD0D	1250	*	STORE VALUE AND CHECK IF NEED	DELIMITER
FD0D AE 84	1251 EXPTDI	LDX	,X	INDIRECTION DESIRED
FD0F 9F 9B	1252 EXPTDL	STX	NUMBER	STORE RESULT
FD11 0D 8E	1253	TST	DELIM	? TO FORCE A DELIMITER
FD13 27 F0	1254 BEQ	EXPRTN		RETURN IF NOT WITH VALUE
FD15 8D 62	1255 BSR	READ		OBTAIN NEXT CHARACTER
FD17	1256 *	TEST FOR + OR -		
FD17 9E 9B	1257 EXPCDL	LDX	NUMBER	LOAD LAST VALUE
FD19 81 2B	1258 CMPA	#'+'		? ADD OPERATOR
FD1B 26 0E	1259 BNE	EXPCHM		BRANCH NOT
FD1D 8D 23	1260 BSR	EXPTRM		COMPUTE NEXT TERM
FD1F 34 02	1261 PSHS	A		SAVE DELIMITER
FD21 DC 9B	1262 LDD	NUMBER		LOAD NEW TERM
FD23 30 8B	1263 EXPADD	LEAX	,X	ADD TO X
FD25 9F 9B	1264 STX	NUMBER		STORE AS NEW RESULT
FD27 35 02	1265 PULS	A		RESTORE DELIMITER
FD29 20 EC	1266 BRA	EXPCDL		NOW TEST IT
FD2B 81 2D	1267 EXPCHM	CMPA	#'-'	? SUBTRACT OPERATOR
FD2D 27 07	1268 BEQ	EXPSUB		BRANCH IF SO
FD2F 81 40	1269 CMPA	#'@'		? INDIRECT DESIRED
FD31 27 DA	1270 BEQ	EXPTDI		BRANCH IF SO
FD33 5F	1271 CLR	B		SET DELIMITER RETURN
FD34 20 CF	1272 BRA	EXPRTN		AND RETURN TO CALLER
FD36 8D 0A	1273 EXPSUB	BSR	EXPTRM	OBTAIN NEXT TERM
FD38 34 02	1274 PSHS	A		SAVE DELIMITER
FD3A DC 9B	1275 LDD	NUMBER		LOAD UP NEXT TERM
FD3C 40	1276 NEGA			NEGATE A
FD3D 50	1277 NEGB			NEGATE B
FD3E 82 00	1278 SBCA	#0		CORRECT FOR A
FD40 20 E1	1279 BRA	EXPADD		GO ADD TO EXPRESION
FD42	1280 *	COMPUTE NEXT EXPRESSION TERM		
FD42	1281 *	OUTPUT: X=OLD VALUE		
FD42	1282 *	'NUMBER' =NEXT TERM		
FD42 8D 9D	1283 EXPTRM	BSR	BLDNUM	OBTAIN NEXT VALUE
FD44 27 32	1284 BEQ	CNVRTS		RETURN IF VALID NUMBER
FD46 16 FC 13	1285 BLDBAD	LBRA	CMDBAD	ABORT COMMAND IF INVALID
FD49	1286			
FD49	1287 *****			
FD49	1288 *	BUILD BINARY VALUE USING INPUT CHARACTERS.		
FD49	1289 *	INPUT: A=ASCII HEX VALUE OR DELIMITER		
FD49	1290 *	SP+0=RETURN ADDRESS		
FD49	1291 *	SP+2=16 BIT RESULT AREA		
FD49	1292 *	OUTPUT: Z=1 A=BINARY VALUE		
FD49	1293 *	Z=0 IF INVALID HEX CHARACTER (A UNCHANGED)		
FD49	1294 *	VOLATILE: D		
FD49	1295 *****			
FD49 0F 9B	1296 BLDHXI	CLR	NUMBER	CLEAR NUMBER
FD4B 0F 9C	1297	CLR	NUMBER+1	CLEAR NUMBER
FD4D 8D 2A	1298 BLDHEX	BSR	READ	GET INPUT CHARACTER
FD4F 8D 11	1299 BLDHXC	BSR	CNVHEX	CONVERT AND TEST CHARACTER
FD51 26 25	1300 BNE	CNVRTS		RETURN IF NOT A NUMBER
FD53 C6 10	1301 LDB	#16		PREPARE SHIFT
FD55 3D	1302 MUL			BY FOUR PLACES
FD56 86 04	1303 LDA	#4		ROTATE BINARY INTO VALUE
FD58 58	1304 BLDSHF	ASLB	OBTAIN	NEXT BIT
FD59 09 9C	1305 ROL	NUMBER+1		INTO LOW BYTE
FD5B 09 9B	1306 ROL	NUMBER		INTO HI BYTE
FD5D 4A	1307 DECA			COUNT DOWN
FD5E 26 F8	1308 BNE	BLDSHF		BRANCH IF MORE TO DO
FD60 20 14	1309 BRA	CNVOK		SET GOOD RETURN CODE
FD62	1310			
FD62	1311 *****			
FD62	1312 *	CONVERT ASCII CHARACTER TO BINARY BYTE		

```

FD62          1313 * INPUT: A=ASCII
FD62          1314 * OUTPUT: Z=1 A=BINARY VALUE
FD62          1315 *           Z=0 IF INVALID
FD62          1316 * ALL REGISTERS TRANSPARENT
FD62          1317 * (A UNALTERED IF INVALID HEX)
FD62          1318 ****
FD62 81 30   1319 CNVHEX CMPA #'0'      ? LOWER THAN A ZERO
FD64 25 12   1320 BLO CNVRTS      BRANCH NOT VALUE
FD66 81 39   1321 CMPA #'9'       ? POSSIBLE A-F
FD68 2F 0A   1322 BLE CNVGOT      BRANCH NO TO ACCEPT
FD6A 81 41   1323 CMPA #'A'       ? LESS THEN TEN
FD6C 25 0A   1324 BLO CNVRTS      RETURN IF MINUS (INVALID)
FD6E 81 46   1325 CMPA #'F'       ? NOT TOO LARGE
FD70 22 06   1326 BHI CNVRTS      NO, RETURN TOO LARGE
FD72 80 07   1327 SUBA #7        DOWN TO BINARY
FD74 84 0F   1328 CNVGOT ANDA #$0F    CLEAR HIGH HEX
FD76 1A 04   1329 CNVOK ORCC #4      FORCE ZERO ON FOR VALID HEX
FD78 39     1330 CNVRTS RTS       RETURN TO CALLER
FD79          1331
FD79          1332 * GET INPUT CHAR, ABORT COMMAND IF CONTROL-X (CANCEL)
FD79 3F      1333 READ SWI        GET NEXT CHARACTER
FD7A 00      1334 FCB INCHNP      FUNCTION
FD7B 81 18   1335 CMPA #CAN      ? ABORT COMMAND
FD7D 27 C7   1336 BEQ BLDBAD     BRANCH TO ABORT IF SO
FD7F 39     1337 RTS          RETURN TO CALLER
FD80          1338 *G
FD80          1339
FD80          1340 ****GO - START PROGRAM EXECUTION
FD80 8D 01   1341 CGO BSR GOADDR     BUILD ADDRESS IF NEEDED
FD82 3B      1342 RTI START       EXECUTING
FD83          1343
FD83          1344 * FIND OPTIONAL NEW PROGRAM COUNTER. ALSO ARM THE
FD83          1345 * BREAKPOINTS.
FD83 35 30   1346 GOADDR PULS Y,X      RECOVER RETURN ADDRESS
FD85 34 10   1347 PSHS X         STORE RETURN BACK
FD87 26 19   1348 BNE GONDFT     IF NO CARRIAGE RETURN THEN NEW PC
FD89          1349 * DEFAULT PROGRAM COUNTER, SO FALL THROUGH IF
FD89          1350 * IMMEDIATE BREAKPOINT.
FD89 17 01 B6 1351 LBSR CBKldr      SEARCH BREAKPOINTS
FD8C AE 6C   1352 LDX 12,S       LOAD PROGRAM COUNTER
FD8E 5A      1353 ARMBLP DECB      COUNT DOWN
FD8F 2B 16   1354 BMI ARMBk2      DONE, NONE TO SINGLE TRACE
FD91 A6 30   1355 LDA -NUMBKp*2,Y PRE-FETCH OPCODE
FD93 AC A1   1356 CMPX ,Y++      ? IS THIS A BREAKPOINT
FD95 26 F7   1357 BNE ARMBLP     LOOP IF NOT
FD97 81 3F   1358 CMPA #$3F      ? SWI BREAKPOINTED
FD99 26 02   1359 BNE ARMNSW     NO, SKIP SETTING OF PASS FLAG
FD9B 97 FB   1360 STA SWIBFL      SHOW UPCOMMING SWI NOT BRKPNT
FD9D 0C 8F   1361 ARMNSW INC MISFLG    FLAG THRU A BREAKPOINT
FD9F 16 01 06 1362 LBRA CDOT      DO SINGLE TRACE W/O BREAKPOINTS
FDA2          1363 * OBTAIN NEW PROGRAM COUNTER
FDA2 17 00 BB 1364 GONDFT LBSR CDNUM     OBTAIN NEW PROGRAM COUNTER
FDA5 ED 6C   1365 STD 12,S       STORE INTO STACK
FDA7 17 01 98 1366 ARMBk2 LBSR CBKldr     OBTAIN TABLE
FDAA 00 FA   1367 NEG BKPTCT      COMPLEMENT TO SHOW ARMED
FDAC 5A      1368 ARMLOP DECB      ? DONE
FDAD 2B C9   1369 BMI CNVRTS      RETURN WHEN DONE
FDAF A6 B4   1370 LDA [,Y]       LOAD OPCODE
FDB1 A7 30   1371 STA -NUMBKp*2,Y STORE INTO OPCODE TABLE
FDB3 86 3F   1372 LDA #$3F      READY "SWI" OPCODE
FDB5 A7 B1   1373 STA [,Y++]    STORE AND MOVE UP TABLE
FDB7 20 F3   1374 BRA ARMLOP      AND CONTINUE
FDB9          1375
FDB9          1376 ****CALL - CALL ADDRESS AS SUBROUTINE
FDB9 8D C8   1377 CCALL BSR GOADDR     FETCH ADDRESS IF NEEDED
FDBB 35 7F   1378 PULS U,Y,X,DP,D,CC RESTORE USERS REGISTERS
FDBD AD F1   1379 JSR [,S++]     CALL USER SUBROUTINE
FDBF 3F     1380 CGOBRK SWI      PERFORM BREAKPOINT

```

FDC0 0A	1381	FCB	BRKPT	FUNCTION	
FDC1 20 FC	1382	BRA	CGOBRK	LOOP UNTIL USER CHANGES PC	
FDC3	1383				
FDC3	1384	*****MEMORY - DISPLAY/CHANGE MEMORY			
FDC3	1385	* CMEMN AND CMPADP ARE DIRECT ENTRY POINTS FROM			
FDC3	1386	* THE COMMAND HANDLER FOR QUICK COMMANDS			
FDC3 17 00 9A	1387	CMEM	LBSR	CNUM	OBTAIN ADDRESS
FDC6 DD 9E	1388	CMEMN	STD	ADDR	STORE DEFAULT
FDC8 9E 9E	1389	CMEM2	LDX	ADDR	LOAD POINTER
FDCA 17 FC 0C	1390	LBSR	ZOUT2H		SEND OUT HEX VALUE OF BYTE
FDCC 86 2D	1391	LDA	#'-'		LOAD DELIMITER
FDCE 3F	1392	SWI			SEND OUT
FDD0 01	1393	FCB	OUTCH		FUNCTION
FDD1 17 FF 0B	1394	CMEM4	LBSR	BLDNNB	OBTAIN NEW BYTE VALUE
FDD4 27 0A	1395	BEQ	CMENUM		BRANCH IF NUMBER
FDD6	1396	* COMA	- SKIP	BYTE	
FDD6 81 2C	1397	CMPA	#', '		? COMMA
FDD8 26 0E	1398	BNE	CMNOTC		BRANCH NOT
FDDA 9F 9E	1399	STX	ADDR		UPDATE POINTER
FDDC 30 01	1400	LEAX	,X		TO NEXT BYTE
FDDE 20 F1	1401	BRA	CMEM4		AND INPUT IT
FDE0 D6 9C	1402	CMENUM	LDB	NUMBER+1	LOAD LOW BYTE VALUE
FDE2 8D 47	1403	BSR	MUPDAT		GO OVERLAY MEMORY BYTE
FDE4 81 2C	1404	CMPA	#', '		? CONTINUE WITH NO DISPLAY
FDE6 27 E9	1405	BEQ	CMEM4		BRANCH YES
FDE8	1406	* QUOTED STRING			
FDE8 81 27	1407	CMNOTC	CMPA	#\$27	? QUOTED STRING
FDEA 26 0C	1408	BNE	CMNOTQ		BRANCH NO
FDEC 8D 8B	1409	CMESTR	BSR	READ	OBTAIN NEXT CHARACTER
FDEE 81 27	1410	CMPA	#\$27		? END OF QUOTED STRING
FDF0 27 0C	1411	BEQ	CMSPCE		YES, QUIT STRING MODE
FDF2 1F 89	1412	TFR	A,B		TO B FOR SUBROUTINE
FDF4 8D 35	1413	BSR	MUPDAT		GO UPDATE BYTE
FDF6 20 F4	1414	BRA	CMESTR		GET NEXT CHARACTER
FDF8	1415	* BLANK	- NEXT	BYTE	
FDF8 81 20	1416	CMNOTQ	CMPA	#\$20	? BLANK FOR NEXT BYTE
FDFA 26 06	1417	BNE	CMNOTB		BRANCH NOT
FDFFC 9F 9E	1418	STX	ADDR		UPDATE POINTER
FDFF 3F	1419	CMSPCE	SWI		GIVE SPACE
FDFD 07	1420	FCB	SPACE		FUNCTION
FE00 20 C6	1421	BRA	CMEM2		NOW PROMPT FOR NEXT
FE02	1422	* LINE FEED	- NEXT	BYTE WITH ADDRESS	
FE02 81 0A	1423	CMNOTB	CMPA	#LF	? LINE FEED FOR NEXT BYTE
FE04 26 08	1424	BNE	CMNOTL		BRANCH NO
FE06 86 0D	1425	LDA	#CR		GIVE CARRIAGE RETURN
FE08 3F	1426	SWI			TO CONSOLE
FE09 01	1427	FCB	OUTCH		HANDLER
FE0A 9F 9E	1428	STX	ADDR		STORE NEXT ADDRESS
FE0C 20 0A	1429	BRA	CMPADP		BRANCH TO SHOW
FE0E	1430	* UP ARROW	- PREVIOUS	BYTE AND ADDRESS	
FE0E 81 5E	1431	CMNOTL	CMPA	#'^'	? UP ARROW FOR PREVIOUS BYTE
FE10 26 0A	1432	BNE	CMNOTU		BRANCH NOT
FE12 30 1E	1433	LEAX	-2,X		DOWN TO PREVIOUS BYTE
FE14 9F 9E	1434	STX	ADDR		STORE NEW POINTER
FE16 3F	1435	CMPADS	SWI		FORCE NEW LINE
FE17 06	1436	FCB	PCRLF		FUNCTION
FE18 8D 07	1437	CMPADP	BSR	PRTADR	GO PRINT ITS VALUE
FE1A 20 AC	1438	BRA	CMEM2		THEN PROMPT FOR INPUT
FE1C	1439	* SLASH	- NEXT	BYTE WITH ADDRESS	
FE1C 81 2F	1440	CMNOTU	CMPA	#'/'	? SLASH FOR CURRENT DISPLAY
FE1E 27 F6	1441	BEQ	CMPADS		YES, SEND ADDRESS
FE20 39	1442	RTS			RETURN FROM COMMAND
FE21	1443				
FE21	1444	* PRINT	CURRENT	ADDRESS	
FE21 9E 9E	1445	PRTADR	LDX	ADDR	LOAD POINTER VALUE
FE23 34 10	1446	PSHS	X		SAVE X ON STACK
FE25 30 E4	1447	LEAX	,S		POINT TO IT FOR DISPLAY
FE27 3F	1448	SWI			DISPLAY POINTER IN HEX

FE28 05	1449	FCB	OUT4HS	FUNCTION
FE29 35 90	1450	PULS	PC,X	RECOVER POINTER AND RETURN
FE2B	1451			
FE2B	1452 * UPDATE BYTE			
FE2B 9E 9E	1453 MUPDAT	LDX	ADDR	LOAD NEXT BYTE POINTER
FE2D E7 80	1454	STB	,X+	STORE AND INCREMENT X
FE2F E1 1F	1455	CMPB	-1,X	? SUCCESFULL STORE
FE31 26 03	1456	BNE	MUPBAD	BRANCH FOR '?' IF NOT
FE33 9F 9E	1457	STX	ADDR	STORE NEW POINTER VALUE
FE35 39	1458	RTS		BACK TO CALLER
FE36 34 02	1459 MUPBAD	PSHS	A	SAVE A REGISTER
FE38 86 3F	1460	LDA	#'?'	SHOW INVALID
FE3A 3F	1461	SWI		SEND OUT
FE3B 01	1462	FCB	OUTCH	FUNCTION
FE3C 35 82	1463	PULS	PC,A	RETURN TO CALLER
FE3E	1464			
FE3E	1465 *****WINDOW -			SET WINDOW VALUE
FE3E 8D 20	1466 CWINDO	BSR	CDNUM	OBTAİN WINDOW VALUE
FE40 DD A0	1467	STD	WINDOW	STORE IT IN
FE42 39	1468	RTS		END COMMAND
FE43	1469			
FE43	1470 *****DISPLAY - HIGH SPEED DISPLAY MEMORY			
FE43 8D 1B	1471 CDISP	BSR	CDNUM	FETCH ADDRESS
FE45 C4 F0	1472	ANDB	##\$F0	FORCE TO 16 BOUNDARY
FE47 1F 02	1473	TFR	D,Y	SAVE IN Y
FE49 30 2F	1474	LEAX	15,Y	DEFAULT LENGTH
FE4B 25 04	1475	BCS	CDISPS	BRANCH IF END OF INPUT
FE4D 8D 11	1476	BSR	CDNUM	OBTAİN COUNT
FE4F 30 AB	1477	LEAX	D,Y	ASSUME COUNT, COMPUTE END ADDR
FE51 34 30	1478 CDISPS	PSHS	Y,X	SETUP PARAMETERS FOR HSDATA
FE53 10 A3 62	1479	CMPD	2,S	? WAS IT COUNT
FE56 23 02	1480	BLS	CDCNT	BRANCH YES
FE58 ED E4	1481	STD	,S	STORE HIGH ADDRESS
FE5A AD 9D E1 84	1482 CDCNT	JSR	[VECTAB+.HSDTA,PCR]	CALL PRINT ROUTINE
FE5E 35 E0	1483	PULS	PC,U,Y	CLEAN STACK AND END COMMAND
FE60	1484			
FE60	1485 * OBTAIN NUMBER - ABORT IF NONE			
FE60	1486 * ONLY DELIMITERS OF CR, BLANK, OR '/' ARE ACCEPTED			
FE60	1487 * OUTPUT: D=VALUE, C=1 IF CARRIAGE RETURN DELMITER,			
FE60	1488 *			ELSE C=0
FE60 17 FE 7E	1489 CDNUM	LBSR	BLDNUM	OBTAİN NUMBER
FE63 26 09	1490	BNE	CDBADN	BRANCH IF INVALID
FE65 81 2F	1491	CMPA	#'/'	? VALID DELIMITER
FE67 22 05	1492	BHI	CDBADN	BRANCH IF NOT FOR ERROR
FE69 81 0E	1493	CMPA	#CR+1	LEAVE COMPARE FOR CARRIAGE RET
FE6B DC 9B	1494	LDD	NUMBER	LOAD NUMBER
FE6D 39	1495	RTS		RETURN WITH COMPARE
FE6E 16 FA EB	1496 CDBADN	LBRA	CMDBAD	RETURN TO ERROR MECHANISM
FE71	1497			
FE71	1498 *****PUNCH - PUNCH MEMORY IN S1-S9 FORMAT			
FE71 8D ED	1499 CPUNCH	BSR	CDNUM	OBTAİN START ADDRESS
FE73 1F 02	1500	TFR	D,Y	SAVE IN Y
FE75 8D E9	1501	BSR	CDNUM	OBTAİN END ADDRESS
FE77 6F E2	1502	CLR	,-S	SETUP PUNCH FUNCTION CODE
FE79 34 26	1503	PSHS	Y,D	STORE VALUES ON STACK
FE7B AD 9D E1 65	1504 CCALBS	JSR	[VECTAB+.BSON,PCR]	INITIALIZE HANDLER
FE7F AD 9D E1 63	1505	JSR	[VECTAB+.BSDTA,PCR]	PERFORM FUNCTION
FE83 34 01	1506	PSHS	CC	SAVE RETURN CODE
FE85 AD 9D E1 5F	1507	JSR	[VECTAB+.BSOFF,PCR]	TURN OFF HANDLER
FE89 35 01	1508	PULS	CC	OBTAİN CONDITION CODE SAVED
FE8B 26 E1	1509	BNE	CDBADN	BRANCH IF ERROR
FE8D 35 B2	1510	PULS	PC,Y,X,A	RETURN FROM COMMAND
FE8F	1511			
FE8F	1512 *****LOAD - LOAD MEMORY FROM S1-S9 FORMAT			
FE8F 8D 01	1513 CLOAD	BSR	CLVOFS	CALL SETUP AND PASS CODE
FE91 01	1514	FCB	1	LOAD FUNCTION CODE FOR PACKET
FE92	1515			
FE92 33 F1	1516 CLVOFS	LEAU	[,S++]	LOAD CODE IN HIGH BYTE OF U

FE94 33 D4	1517	LEAU	[,U]	NOT CHANGING CC AND RESTORE S
FE96 27 03	1518	BEQ	CLVDFT	BRANCH IF CARRIAGE RETURN NEXT
FE98 8D C6	1519	BSR	CNUM	OBTAIN OFFSET
FE9A 8C	1520	FCB	SKIP2	SKIP DEFAULT OFFSET
FE9B 4F	1521	CLVDFT	CLRA	CREATE ZERO OFFSET
FE9C 5F	1522	CLRB		AS DEFAULT
FE9D 34 4E	1523	PSHS	U,DP,D	SETUP CODE, NULL WORD, OFFSET
FE9F 20 DA	1524	BRA	CCALBS	ENTER CALL TO BS ROUTINES
FEA1	1525			
FEA1	1526	*****VERIFY - COMPARE MEMORY WITH FILES		
FEA1 8D EF	1527	CVER	BSR CLVOFS	COMPUTE OFFSET IF ANY
FEA3 FF	1528	FCB	-1	VERIFY FNCTN CODE FOR PACKET
FEA4	1529			
FEA4	1530	*****TRACE - TRACE INSTRUCTIONS		
FEA4	1531	***** . - SINGLE STEP TRACE		
FEA4 8D BA	1532	CTRACE	BSR CNUM	OBTAIN TRACE COUNT
FEA6 DD 91	1533	STD	TRACEC	STORE COUNT
FEA8 32 62	1534	CDOT	LEAS 2,S	RID COMMAND RETURN FROM STACK
FEAA EE F8 0A	1535	CTRCE3	LDU [10,S]	LOAD OPCODE TO EXECUTE
FEAD DF 99	1536	STU	LASTOP	STORE FOR TRACE INTERRUPT
FEAF DE F6	1537	LDU	VECTAB+.PTM	LOAD PTM ADDRESS
FEB1 CC 07 01	1538	LDD	#\$0701	7,1 CYCLES DOWN+CYCLES UP
FEB4 ED 42	1539	STD	PTMTM1-PTM,U	START NMI TIMEOUT
FEB6 3B	1540	RTI	RETURN	FOR ONE INSTRUCTION
FEB7	1541			
FEB7	1542	*****NULLS - SET NEW LINE AND CHAR PADDING		
FEB7 8D A7	1543	CNULLS	BSR CNUM	OBTAIN NEW LINE PAD
FEB9 DD F2	1544	STD	VECTAB+.PAD	RESET VALUES
FEBB 39	1545	RTS		END COMMAND
FEBC	1546			
FEBC	1547	*****STLEVEL - SET STACK TRACE LEVEL		
FEBC 27 05	1548	CSTLEV	BEQ STLDFT	TAKE DEFAULT
FEBE 8D A0	1549	BSR	CNUM	OBTAIN NEW STACK LEVEL
FEC0 DD F8	1550	STD	SLEVEL	STORE NEW ENTRY
FEC2 39	1551	RTS		TO COMMAND HANDLER
FEC3 30 6E	1552	STLDFT	LEAX 14,S	COMPUTE NMI COMPARE
FEC5 9F F8	1553	STX	SLEVEL	AND STORE IT
FEC7 39	1554	RTS		END COMMAND
FEC8	1555			
FEC8	1556	*****OFFSET - COMPUTE SHORT AND LONG		
FEC8	1557	***** BRANCH OFFSETS		
FEC8 8D 96	1558	COFFS	BSR CNUM	OBTAIN INSTRUCTION ADDRESS
FECA 1F 01	1559	TFR	D,X	USE AS FROM ADDRESS
FECC 8D 92	1560	BSR	CNUM	OBTAIN TO ADDRESS
FECE	1561	* D=TO INSTRUCTION, X=FROM INSTRUCTION OFFSET BYTE(S)		
FECE 30 01	1562	LEAX	1,X	ADJUST FOR *+2 SHORT BRANCH
FED0 34 30	1563	PSHS	Y,X	STORE WORK WORD AND VALUE ON S
FED2 A3 E4	1564	SUBD	,S	FIND OFFSET
FED4 ED E4	1565	STD	,S	SAVE OVER STACK
FED6 30 61	1566	LEAX	1,S	POINT FOR ONE BYTE DISPLAY
FED8 1D	1567	SEX	SIGN	EXTEND LOW BYTE
FED9 A1 E4	1568	CMPA	,S	? VALID ONE BYTE OFFSET
FEDB 26 02	1569	BNE	COFNO1	BRANCH IF NOT
FEDD 3F	1570	SWI		SHOW ONE BYTE OFFSET
FEDE 04	1571	FCB	OUT2HS	FUNCTION
FEDF EE E4	1572	COFNO1	LDU ,S	RELOAD OFFSET
FEE1 33 5F	1573	LEAU	-1,U	CONVERT TO LONG BRANCH OFFSET
FEE3 EF 84	1574	STU	,X	STORE BACK WHERE X POINTS NOW
FEE5 3F	1575	SWI		SHOW TWO BYTE OFFSET
FEE6 05	1576	FCB	OUT4HS	FUNCTION
FEE7 3F	1577	SWI		FORCE NEW LINE
FEE8 06	1578	FCB	PCRLF	FUNCTION
FEE9 35 96	1579	PULS	PC,X,D	RESTORE STACK AND END COMMAND
FEEB	1580	*H		
FEEB	1581			
FEEB	1582	*****BREAKPOINT - DISPLAY/ENTER/DELETE/CLEAR		
FEEB	1583	***** BREAKPOINTS		
FEEB 27 23	1584	CBKPT	BEQ CBKDSP	BRANCH DISPLAY OF JUST 'B'

FEED	17 FD F1	1585	LBSR	BLDNUM	ATTEMPT VALUE ENTRY	
FEF0	27 2C	1586	BEQ	CBKADD	BRANCH TO ADD IF SO	
FEF2	81 2D	1587	CMPA	#' - '	? CORRECT DELIMITER	
FEF4	26 3F	1588	BNE	CBKERR	NO, BRANCH FOR ERROR	
FEF6	17 FD E8	1589	LBSR	BLDNUM	ATTEMPT DELETE VALUE	
FEF9	27 03	1590	BEQ	CBKDLE	GOT ONE, GO DELETE IT	
FEFB	0F FA	1591	CLR	BKPTCT	WAS 'B - ', SO ZERO COUNT	
FEFD	39	1592	CBKRTS	RTS	END COMMAND	
FEFE		1593	* DELETE THE ENTRY			
FEFE	8D 40	1594	CBKDLE	BSR	CBKSET	SETUP REGISTERS AND VALUE
FF00	5A	1595	CBKDLPL	DEC B	? ANY ENTRIES IN TABLE	
FF01	2B 32	1596	BMI	CBKERR	BRANCH NO, ERROR	
FF03	AC A1	1597	CMPX	, Y++	? IS THIS THE ENTRY	
FF05	26 F9	1598	BNE	CBKDLPL	NO, TRY NEXT	
FF07		1599	* FOUND, NOW MOVE OTHERS UP IN ITS PLACE			
FF07	AE A1	1600	CBKDLM	LD X	, Y++	LOAD NEXT ONE UP
FF09	AF 3C	1601	ST X	-4, Y	MOVE DOWN BY ONE	
FF0B	5A	1602	DEC B		? DONE	
FF0C	2A F9	1603	BPL	CBKDLM	NO, CONTINUE MOVE	
FF0E	0A FA	1604	DEC	BKPTCT	DECREMENT BREAKPOINT COUNT	
FF10	8D 2E	1605	CBKDSP	BSR	CBKSET	SETUP REGISTERS AND LOAD VALUE
FF12	27 E9	1606	BEQ	CBKRTS	RETURN IF NONE TO DISPLAY	
FF14	30 A1	1607	CBKDSL	LEAX	, Y++	POINT TO NEXT ENTRY
FF16	3F	1608	SWI		DISPLAY IN HEX	
FF17	05	1609	FCB	OUT4HS	FUNCTION	
FF18	5A	1610	DEC B		COUNT DOWN	
FF19	26 F9	1611	BNE	CBKDSL	LOOP IF NGABLE RAM	
FF1B	3F	1612	SWI		SKIP TO NEW LINK	
FF1C	06	1613	FCB	PCRLF	FUNCTIONRTS	
FF1D	39	1614	RTS			
FF1E		1615	* ADD NEW ENTRY			
FF1E	8D 20	1616	CBKADD	BSR	CBKSET	SETUP REGISTERS
FF20	C1 08	1617	CMPB	#NUMBKPK	? ALREADY FULL	
FF22	27 11	1618	BEQ	CBKERR	BRANCH ERROR IF SO	
FF24	A6 84	1619	LDA	, X	LOAD BYTE TO TRAP	
FF26	E7 84	1620	STB	, X	TRY TO CHANGE	
FF28	E1 84	1621	CMPB	, X	? CHANGEABLE RAM	
FF2A	26 09	1622	BNE	CBKERR	BRANCH ERROR IF NOT	
FF2C	A7 84	1623	STA	, X	RESTORE BYTE	
FF2E	5A	1624	CBKADL	DEC B	COUNT DOWN	
FF2F	2B 07	1625	BMI	CBKADT	BRANCH IF DONE TO ADD IT	
FF31	AC A1	1626	CMPX	, Y++	? ENTRY ALREADY HERE	
FF33	26 F9	1627	BNE	CBKADL	LOOP IF NOT	
FF35	16 FA 24	1628	CBKERR	LBRA	CMDBAD	
FF38	AF A4	1629	CBKADT	ST X	, Y	
FF3A	6F 31	1630	CLR	-NUMBKPK*2+1, Y	RETURN TO ERROR PRODUCE	
FF3C	0C FA	1631	INC	BKPTCT	ADD THIS ENTRY	
FF3E	20 D0	1632	BRA	CBKDSP	CLEAR OPTIONAL BYTE	
FF40		1633	* SETUP REGISTERS FOR SCAN			
FF40	9E 9B	1634	CBKSET	LD X	ADD ONE TO COUNT	
FF42	31 8D E0 6C	1635	CBKLDR	LEAY	NUMBER	
FF46	D6 FA	1636	LDB	BKPTCT	AND NOW DISPLAY ALL OF 'EM	
FF48	39	1637	RTS		LOAD VALUE DESIRED	
FF49		1638			LOAD START OF TABLE	
FF49		1639	*****ENCODE - ENCODE A POSTBYTE			
FF49	6F E2	1640	CENCODE	CLR	, -S	DEFAULT TO NOT INDIRECT
FF4B	5F	1641	CLRB			ZERO POSTBYTE VALUE
FF4C	30 8C 3F	1642	LEAX	<CONV1, PCR	START TABLE SEARCH	
FF4F	3F	1643	SWI		OBTAIN FIRST CHARACTER	
FF50	00	1644	FCB	INCHNP	FUNCTION	
FF51	81 5B	1645	CMPA	#'[]'	? INDIRECT HERE	
FF53	26 06	1646	BNE	CEN2	BRANCH IF NOT	
FF55	86 10	1647	LDA	#\$10	SET INDIRECT BIT ON	
FF57	A7 E4	1648	STA	, S	SAVE FOR LATER	
FF59	3F	1649	CENGET	SWI	OBTAIN NEXT CHARACTER	
FF5A	00	1650	FCB	INCHNP	FUNCTION	
FF5B	81 0D	1651	CEN2	CMPA	? END OF ENTRY	
FF5D	27 OC	1652	BEQ	CEND1	BRANCH YES	

FF5F	6D	84	1653	CENLP1	TST	,X	? END OF TABLE
FF61	2B	D2	1654	BMI	CBKERR		BRANCH ERROR IF SO
FF63	A1	81	1655	CMPA	,X++		? THIS THE CHARACTER
FF65	26	F8	1656	BNE	CENLP1		BRANCH IF NOT
FF67	EB	1F	1657	ADDB	-1,X		ADD THIS VALUE
FF69	20	EE	1658	BRA	CENGET		GET NEXT INPUT
FF6B	30	8C 49	1659	CEND1	LEAX	<CONV2, PCR	POINT AT TABLE 2
FF6E	1F	98	1660	TFR	B,A		SAVE COPY IN A
FF70	84	60	1661	ANDA	#\$60		ISOLATE REGISTER MASK
FF72	AA	E4	1662	ORA	,S		ADD IN INDIRECTION BIT
FF74	A7	E4	1663	STA	,S		SAVE BACK AS POSTBYTE SKELETON
FF76	C4	9F	1664	ANDB	#\$9F		CLEAR REGISTER BITS
FF78	6D	84	1665	CENLP2	TST	,X	? END OF TABLE
FF7A	27	B9	1666	BEQ	CBKERR		BRANCH ERROR IF SO
FF7C	E1	81	1667	CMPB	,X++		? SAME VALUE
FF7E	26	F8	1668	BNE	CENLP2		LOOP IF NOT
FF80	E6	1F	1669	LDB	-1,X		LOAD RESULT VALUE
FF82	EA	E4	1670	ORB	,S		ADD TO BASE SKELETON
FF84	E7	E4	1671	STB	,S		SAVE POSTBYTE ON STACK
FF86	30	E4	1672	LEAX	,S		POINT TO IT
FF88	3F		1673	SWI			SEND OUT AS HEX
FF89	04		1674	FCB	OUT2HS		FUNCTION
FF8A	3F		1675	SWI			TO NEXT LINE
FF8B	06		1676	FCB	PCRLF		FUNCTION
FF8C	35	84	1677	PULS	PC,B		END OF COMMAND
FF8E			1678				
FF8E			1679	*	TABLE ONE DEFINES VALID INPUT IN SEQUENCE		
FF8E	41	04	42	05	44	06	+ 1680 CONV1 FCB 'A', \$04, 'B', \$05, 'D', \$06, 'H', \$01
FF96	48	01	48	01	48	00	+ 1681 FCB 'H', \$01, 'H', \$01, 'H', \$00, ' ', '\$00
FF9E	2D	09	2D	01	53	70	+ 1682 FCB '-' , \$09, '-' , \$01, 'S' , \$70, 'Y' , \$30
FFA6	55	50	58	10	2B	07	+ 1683 FCB 'U' , \$50, 'X' , \$10, '+' , \$07, '+' , \$01
FFAE	50	80	43	00	52	00	+ 1684 FCB 'P' , \$80, 'C' , \$00, 'R' , \$00, ']' , \$00
FFB6	FF		1685	FCB	\$FF		END OF TABLE
FFB7			1686	*	CONV2 USES ABOVE CONVERSION TO SET POSTBYTE		
FFB7			1687	*	BIT SKELETON.		
FFB7	10	84	11	00	CONV2 FDB \$1084, \$1100	R, H,R	
FFBB	12	88	13	89	FDB \$1288, \$1389	HH,R HHHH,R	
FFBF	14	86	15	85	FDB \$1486, \$1585	A,R B,R	
FFC3	16	88	17	80	FDB \$168B, \$1780	D,R ,R+	
FFC7	18	81	19	82	FDB \$1881, \$1982	,R++ , -R	
FFCB	1A	83	82	8C	FDB \$1A83, \$828C	,--R HH,PCR	
FFCF	83	8D	03	9F	FDB \$838D, \$039F	HHHH,PCR [HHHH]	
FFD3	00		1695	FCB	0		END OF TABLE
FFD4			1696				
FFD4			1697	*****	*****	*****	*****
FFD4			1698	*	DEFAULT INTERRUPT TRANSFERS	*	*
FFD4			1699	*****	*****	*****	*****
FFD4	6E	9D	DF	EE	1700 RSRVD JMP [VECTAB+.RSVD, PCR]	RESERVED VECTOR	
FFD8	6E	9D	DF	EC	1701 SWI3 JMP [VECTAB+.SWI3, PCR]	SWI3 VECTOR	
FFDC	6E	9D	DF	EA	1702 SWI2 JMP [VECTAB+.SWI2, PCR]	SWI2 VECTOR	
FFE0	6E	9D	DF	E8	1703 FIRQ JMP [VECTAB+.FIRQ, PCR]	FIRQ VECTOR	
FFE4	6E	9D	DF	E6	1704 IRQ JMP [VECTAB+.IRQ, PCR]	IRQ VECTOR	
FFE8	6E	9D	DF	E4	1705 SWI JMP [VECTAB+.SWI, PCR]	SWI VECTOR	
FFEC	6E	9D	DF	E2	1706 NMI JMP [VECTAB+.NMI, PCR]	NMI VECTOR	
FFF0			1707				
FFF0			1708	*****	*****	*****	*****
FFF0			1709	*	ASSIST09 HARDWARE VECTOR TABLE		
FFF0			1710	*	THIS TABLE IS USED IF THE ASSIST09 ROM ADDRESSES		
FFF0			1711	*	THE MC6809 HARDWARE VECTORS.		
FFF0			1712	*****	*****	*****	*****
FFF0			1713	ORG ROMBEG+ROMSIZ-16		SETUP HARDWARE VECTORS	
FFF0	FF	D4	1714	FDB RSRVD		RESERVED SLOT	
FFF2	FF	D8	1715	FDB SWI3		SOFTWARE INTERRUPT 3	
FFF4	FF	DC	1716	FDB SWI2		SOFTWARE INTERRUPT 2	
FFF6	FF	E0	1717	FDB FIRQ		FAST INTERRUPT REQUEST	
FFF8	FF	E4	1718	FDB IRQ		INTERRUPT REQUEST	
FFFA	FF	E8	1719	FDB SWI		SOFTWARE INTERRUPT	
FFFC	FF	EC	1720	FDB NMI		NON-MASKABLE INTERRUPT	

DUNFIELD 6809 ASSEMBLER: ASSIST09

PAGE: 27

FFFE F8 37
0000

1721
1722

FDB

RESET

RESTART

SYMBOL TABLE:

.ACIA	-002E	.AVTBL	-0000	.BSDTA	-0024	.BSOFF	-0026	.BSON	-0022
.CIDTA	-0016	.CIOFF	-0018	.CION	-0014	.CMDL1	-0002	.CMDL2	-002C
.CODTA	-001C	.COOFF	-001E	.COON	-001A	.ECHO	-0032	.EXPAN	-002A
.FIRQ	-000A	.HSDTA	-0020	.IRQ	-000C	.NMI	-0010	.PAD	-0030
.PAUSE	-0028	.PTM	-0034	.RESET	-0012	.RSVD	-0004	.SWI	-000E
.SWI2	-0008	.SWI3	-0006	.ACIA	-E008	ADDR	-DF9E	ARMBK2	-FDA7
ARMBLP	-FD8E	ARLOP	-FDAC	ARMNSW	-FD9D	BASEPG	-DF9D	BELL	-0007
BKPTBL	-DFB2	BKPTCT	-DFFA	BKPTOP	-DFA2	BLD2	-F815	BLD3	-F821
BLDBAD	-FD46	BLDHEX	-FD4D	BLDHXC	-FD4F	BLDHXI	-FD49	BLDNNB	-FCDF
BLDNUM	-FCE1	BLDRTN	-F835	BLDSHF	-FD58	BLDVTR	-F800	BRKPT	-000A
BSDCMP	-FB6A	BSDEOL	-FB70	BSDLD1	-FB40	BSDLD2	-FB42	BSDNXT	-FB60
BSDPUN	-FB92	BSDSRT	-FB6E	BSDTA	-FB38	BSOFF	-FB27	BSOFLP	-FB33
BSON	-FB1B	BSON2	-FB22	BSPEOF	-FBEF	BSPGO	-FBA3	BSPMRE	-FBC6
BSPOK	-FBAF	BSPSTR	-FBEC	BSPUN2	-FBE7	BSPUNC	-FBE9	BYTE	-FB75
BYTHEX	-FB89	BYTRTS	-FB88	CAN	-0018	CBKADD	-FF1E	CBKADL	-FF2E
CBKADT	-FF38	CBKDLE	-FEFE	CBKDLM	-FF07	CBKDLP	-FF00	CBKDSL	-FF14
CBKDSP	-FF10	CBKERR	-FF35	CBKLDR	-FF42	CBKPT	-FEEB	CBKRITS	-FEFD
CBKSET	-FF40	CCALBS	-FE7B	CCALL	-FDB9	CDBADN	-FE6E	CDCNT	-FE5A
CDISP	-FE43	CDISPS	-FE51	CDNUM	-FE60	CDOT	-FEA8	CEN2	-FF5B
CENCDE	-FF49	CEND1	-FF6B	CENGET	-FF59	CENLP1	-FF5F	CENLP2	-FF78
CGO	-FD80	CGOBRK	-FDBF	CHKABT	-FA58	CHKRTN	-FA61	CHKSEC	-FA60
CHKWT	-FA62	CIDTA	-FADC	CIOFF	-FAF0	CION	-FAE6	CIRTN	-FAE5
CLOAD	-FE8F	CLVDFT	-FE9B	CLVOFS	-FE92	CMD	-F8F7	CMD2	-F935
CMD3	-F948	CMDBAD	-F95C	CMDCMP	-F977	CMDDDL	-F901	CMDFLS	-F96C
CMDGOT	-F94D	CMDMEM	-F990	CMDNEP	-F8F9	CMDNOL	-F90A	CMDSCH	-F953
CMDSIZ	-F96F	CMDSMS	-F967	CMDTB2	-F99B	CMDTBL	-F99C	CMDXQT	-F987
CMEM	-FDC3	CMEM2	-FDC8	CMEM4	-FDD1	CMEMN	-FDC6	CMEMUM	-FDE0
CMESTR	-FDEC	CMNOTB	-FE02	CMNOTC	-FDE8	CMNOTL	-FE0E	CMNOTQ	-FDF8
CMNOTU	-FE1C	CMPADP	-FE18	CMPADS	-FE16	CMPCCE	-FDFE	CNULLS	-FEB7
CNVGOT	-FD74	CNVHEX	-FD62	CNVOK	-FD76	CNRVTS	-FD78	CODTA	-FAF1
CODTAD	-FB0F	CODTAO	-FB12	CODTLT	-FB07	CODTPD	-FB03	CODTRT	-FB0D
COFFS	-FEC8	COFNO1	-FEDF	CONV1	-FF8E	CONV2	-FFB7	COOFF	-FAF0
COON	-FAE6	CUPNCH	-FE71	CR	-000D	CREG	-FC4A	CSTLEV	-FEBC
CTRACE	-FEA4	CTRCE3	-FEAA	CVER	-FEA1	CWINDO	-FE3E	DELIM	-DF8E
DFTCHP	-0000	DFTNLP	-0005	DLE	-0010	EOT	-0004	ERRMSG	-FABD
ERROR	-FACE	EXP1	-FCE9	EXP2	-FD07	EXPADD	-FD23	EXPCDL	-FD17
EXPCHM	-FD2B	EXPDL	-FCEB	EXPRTN	-FD05	EXPSSB	-FD36	EXPTDI	-FD0D
EXPTDL	-FD0F	EXPTRM	-FD42	FIRQ	-FFE0	FIRQR	-FABC	GOADDR	-FD83
GONDFT	-FDA2	HIVTR	-0034	HSBLNK	-FC00	HSDRTN	-FC47	HSDTA	-FBFC
HSHCHR	-FC2B	HSHCOK	-FC35	HSHDOT	-FC33	HSHLNE	-FC14	HSHNXT	-FC20
HSHTTL	-FC06	INCHNP	-0000	INITVT	-F844	INTVE	-F87D	INTVS	-F870
IRQ	-FEF4	IRQR	-FAD8	LASTOP	-DF99	LDDP	-FAC1	LF	-000A
MISFLG	-DF8F	MONITR	-0008	MSHOWP	-FA79	MUPBAD	-FE36	MUPDAT	-FE2B
NMI	-FFEC	NMICON	-FAB7	NMIR	-FA7D	NMITRC	-FAB0	NUMBER	-DF9B
NUMBK	-0008	NUMFUN	-000B	NUMVTR	-001B	OUT2HS	-0004	OUT4HS	-0005
OUTCH	-0001	PAUSE	-000B	PAUSER	-DFFC	PCNTER	-DF93	PCRLF	-0006
PDATA	-0003	PDATA1	-0002	PROMPT	-003E	PRTADR	-FE21	PSTACK	-DF95
PTM	-E000	PTMC13	-E000	PTMC2	-E001	PTMSTA	-E001	PTMTM1	-E002
PTMTM2	-E004	PTMTM3	-E006	RAMOFS	-E700	READ	-FD79	REG4	-FC94
REGAGN	-FCC3	REGCHG	-FC70	REGCNG	-FC9D	REGMSK	-FC50	REGNXC	-FCB1
REGP1	-FC78	REGP2	-FC81	REGP3	-FC92	REGPRS	-FAB3	REGRT	-FC6F
REGRTN	-FC9B	REGSKP	-FCAA	REGTF1	-FCC9	REGTF2	-FCD6	REGTWO	-FCBB
RESET	-F837	RESET2	-F83D	ROM2OF	-F000	ROM2WK	-DF66	ROMBEG	-F800
ROMSIZ	-0800	RSRVD	-FFD4	RSRVRD	-FAD8	RSTACK	-DF97	RTI	-FABC
RTS	-FAF0	SEND	-F9EC	SIGNON	-F8C9	SKIP2	-008C	SLEVEL	-DFF8
SPACE	-0007	STACK	-DF51	STLDFT	-FEC3	SWI	-FFE8	SWI2	-FFDC
SWI2R	-FAD8	SWI3	-FFD8	SWI3R	-FAD8	SWIBFL	-DFFB	SWICNT	-DF90
SWIDNE	-F8B5	SWILP	-F8A8	SWIR	-F895	SWIVTB	-F87D	TRACEC	-DF91
TSTACK	-DF51	VCTRSW	-0009	VECTAB	-DFC2	WINDOW	-DFA0	WORKPG	-DF00
XQCIDT	-FA72	XQPAUS	-FA6E	ZBKCMD	-FAD5	ZBKPN	-FAD3	ZIN2	-FA2A
ZINCH	-FA11	ZINCHP	-FA0F	ZMONT2	-F8E6	ZMONTR	-F8D2	ZOT2HS	-F9F2
ZOT4HS	-F9F0	ZOTCH1	-FA2E	ZOTCH2	-FA37	ZOTCH3	-FA39	ZOUT2H	-F9D9
ZOUTHX	-F9E6	ZPAUSE	-FA4E	ZPCRLF	-FA3D	ZPCRLS	-FA3C	ZPDATA	-FA40
ZPDATA1	-FA48	ZPDTLP	-FA46	ZSPACE	-F9F6	ZVSWTH	-F9FA		

Symbol	Define	References									
.ACIA	92										
.AVTBL	69										
.BSDTA	87										
.BSOFF	88										
.BSON	86										
.CIDTA	80										
.CIOFF	81										
.CION	79										
.CMDL1	70										
.CMDL2	91										
.CODTA	83										
.COOFF	84										
.COON	82										
.ECHO	94										
.EXPAN	90										
.FIRQ	74										
.HSDTA	85										
.IRQ	75										
.NMI	77										
.PAD	93										
.PAUSE	89										
.PTM	95										
.RESET	78										
.RSVD	71										
.SWI	76										
.SWI2	73										
.SWI3	72										
ACIA	21	253	822	834	850						
ADDR	130	1236	1388	1389	1399	1418	1428	1434	1445	1453	1457
ARMBK2	1366	770	1354								
ARMBLP	1353	1357									
ARMLOP	1368	1374									
ARMNSW	1361	1359									
BASEPG	132	183	781								
BELL	33	779									
BKPTBL	124	1635									
BKPTCT	118	383	1367	1591	1604	1631	1636				
BKPTOP	126										
BLD2	189	193									
BLD3	195	198									
BLDBAD	1285	1336									
BLDHEX	1298	1247									
BLDHXC	1299	418									
BLDHXI	1296	1230									
BLDNNB	1216	1161	1394								
BLDNUM	1219	1283	1489	1585	1589						
BLDRTN	204	202									
BLDSHF	1304	1308									
BLDVTR	180	215									
BRKPT	63	1381									
BSDCMP	941	939									
BSDEOL	945	937									
BSDLD1	916	919	946								
BSDLD2	918	925									
BSDNXT	936	942									
BSDPUN	974	910									
BSDSRT	943	923	947								
BSDTA	908	247	1505								
BSOFF	888	248	1507								
BSOFLP	896	897									
BSON	877	246	1504								
BSON2	881	879									
BSPEOF	1030	1018									
BSPGO	984	1017									
BSPMRE	1006	1008									
BSPOK	989	987									
BSPSTR	1029	994									
BSPUN2	1026	1000	1002	1003	1006						
BSPUNC	1027	1014									
BYTE	950	927	930	932	936						
BYTHEX	962	950	953								
BYTRTS	960	965									
CAN	37	708	715	1335							
CBKADD	1616	1586									
CBKADL	1624	1627									
CBKADT	1629	1625									
CBKDLE	1594	1590									
CBKDLM	1600	1603									
CBKDLR	1595	1598									
CBKDSL	1607	1611									
CBKDSP	1605	1584	1632								
CBKERR	1628	1588	1596	1618	1622	1654	1666				

APPENDIX C

MACHINE CODE TO INSTRUCTION CROSS REFERENCE

C.1 INTRODUCTION

This appendix contains a cross reference between the machine code, represented in hexadecimal and the instruction and addressing mode that it represents. The number of MPU cycles and the number of program bytes is also given. Refer to Table C-1.

Table C-1. Machine Code to Instruction Cross Reference

OP	Mnem	Mode	~	#	OP	Mnem	Mode	~	#	OP	Mnem	Mode	~	#
00	NEG	Direct	6	2	30	LEAX	Indexed	4+	2+	60	NEG	Indexed	6+	2+
01	•				31	LEAY		4+	2+	61	•			
02	•				32	LEAS		4+	2+	62	•			
03	COM		6	2	33	LEAU	Indexed	4+	2+	63	COM		6+	2+
04	LSR		6	2	34	PSHS	Immed	5+	2	64	LSR		6+	2+
05	•				35	PULS		5+	2	65	•			
06	ROR		6	2	36	PSHU		5+	2	66	ROR		6+	2+
07	ASR		6	2	37	PULU	Immed	5+	2	67	ASR		6+	2+
08	ASL, LSL		6	2	38	•	Inherent	5	1	68	ASL, LSL		6+	2+
09	ROL		6	2	39	RTS		3	1	69	ROL		6+	2+
0A	DEC		6	2	3A	ABX		6/15	1	6A	DEC		6+	2+
0B	•				3B	RTI		20	2	6C	INC		6+	2+
0C	INC		6	2	3C	CWAI		11	1	6D	TST		6+	2+
0D	TST		6	2	3D	MUL				6E	JMP		3+	2+
0E	JMP		3	2	3E	•				6F	CLR		6+	2+
0F	CLR	Direct	6	2	3F	SWI	Inherent	19	1			Indexed		
10	Page 2		—	—	40	NEGA	Inherent	2	1	70	NEG	Extended	7	3
11	Page 3		—	—	41	•				71	•			
12	NOP		Inherent	2	42	•				72	•			
13	SYNC		Inherent	4	43	COMA		2	1	73	COM		7	3
14	•				44	LSRA		2	1	74	LSR		7	3
15	•				45	•				75				
16	LBRA	Relative	5	3	46	RORA		2	1	76	ROR		7	3
17	LBSR	Relative	9	3	47	ASRA		2	1	77	ASR		7	3
18	•				48	ASLA, LSLA		2	1	78	ASL, LSL		7	3
19	DAA	Inherent	2	1	49	ROLA		2	1	79	ROL		7	3
1A	ORCC	Immed	3	2	4A	DECA		2	1	7A	DEC		7	3
1B	•	—			4B	•				7B	•			
1C	ANDCC	Immed	3	2	4C	INCA		2	1	7C	INC		7	3
1D	SEX	Inherent	2	1	4D	TSTA		2	1	7D	TST		7	3
1E	EXG	Immed	8	2	4E	•				7E	JMP		4	3
1F	TFR	Immed	6	2	4F	CLRA	Inherent	2	1	7F	CLR	Extended	7	3
20	BRA	Relative	3	2	50	NEGB	Inherent	2	1	80	SUBA	Immed	2	2
21	BRN		3	2	51	•				81	CMPA		2	2
22	BHI		3	2	52	•				82	SBCA		2	2
23	BLS		3	2	53	COMB		2	1	83	SUBD		4	3
24	BHS, BCC		3	2	54	LSRB		2	1	84	ANDA		2	2
25	BLO, BCS		3	2	55	•				85	BITA		2	2
26	BNE		3	2	56	RORB		2	1	86	LDA		2	2
27	BEQ		3	2	57	ASRB		2	1	87	•			
28	BVC		3	2	58	ASLB, LSLB		2	1	88	EORA		2	2
29	BVS		3	2	59	ROLB		2	1	89	ADCA		2	2
2A	BPL		3	2	5A	DEC B		2	1	8A	ORA		2	2
2B	BMI		3	2	5B	•				8B	ADDA		2	2
2C	BGE		3	2	5C	INCB		2	1	8C	CMPX		4	3
2D	BLT		3	2	5D	TSTB		2	1	8D	BSR	Relative	7	2
2E	BGT		3	2	5E	•				8E	LDX	Immed	3	3
2F	BLE	Relative	3	2	5F	CLRB	Inherent	2	1	8F	•			

LEGEND:

~ Number of MPU cycles (less possible push pull or indexed-mode cycles)

Number of program bytes

• Denotes unused opcode

Table C-1. Machine Code to Instruction Cross Reference (Continued)

OP	Mnem	Mode	-	#	OP	Mnem	Mode	-	#	OP	Mnem	Mode	-	#
90	SUBA	Direct	4	2	C0	SUBB	Immed	2	2	1021	LBRN	Relative	5	4
91	CMPA		4	2	C1	CMPB		2	2	1022	LBHI		5(6)	4
92	SBCA		4	2	C2	SBCB		2	2	1023	LBLS		5(6)	4
93	SUBD		6	2	C3	ADDD		4	3					
94	ANDA		4	2	C4	ANDB		2	2	1024	LBHS, LBCC		5(6)	4
95	BITA		4	2	C5	BITB	Immed	2	2	1025	LBCS, LBLO		5(6)	4
96	LDA		4	2	C6	LDB	Immed	2	2	1026	LBNE		5(6)	4
97	STA		4	2	C7	*				1027	LBEQ		5(6)	4
98	EORA		4	2	C8	EORB		2	2	1028	LBVC		5(6)	4
99	ADCA		4	2	C9	ADCB		2	2	1029	LBVS		5(6)	4
9A	ORA		4	2	CA	ORB		2	2	102A	LBPL		5(6)	4
9B	ADDA		4	2	CB	ADT 3		2	2	102B	LBMI		5(6)	4
9C	CMPX		6	2	CC	LDD		3	3	102C	LBGE		5(6)	4
9D	JSR		7	2	CD	*				102D	LBLT		5(6)	4
9E	LDX		5	2	CE	LDU	Immed	3	3	102E	LBGT		5(6)	4
9F	STX	Direct	5	2	CF	*				102F	LBLE	Relative	5(6)	4
A0	SUBA	Indexed	4+	2+	D0	SUBB	Direct	4	2	103F	SWI2	Inherent	20	2
A1	CMPA		4+	2+	D1	CMPB		4	2	1083	CMPD	Immed	5	4
A2	SBCA		4+	2+	D2	SBCB		6	2	108C	CMPY		5	4
A3	SUBD		6+	2+	D3	ADDD		4	2	108E	LDY	Immed	4	4
A4	ANDA		4+	2+	D4	ANDB		4	2	1093	CMPD	Direct	7	3
A5	BITA		4+	2+	D5	BITB		4	2	109C	CMPY		7	3
A6	LDA		4+	2+	D6	LDB		4	2	109E	LDY		6	3
A7	STA		4+	2+	D7	STB		4	2	109F	STY		6	3
A8	EORA		4+	2+	D8	EORB		4	2	10A3	CMPD	Indexed	7+	3+
A9	ADCA		4+	2+	D9	ADCB		4	2	10AC	CMPY		7+	3+
AA	ORA		4+	2+	DA	ORB		4	2	10AE	LDY		6+	3+
AB	ADDA		4+	2+	DB	ADDB		4	2	10AF	STY		6+	3+
AC	CMPX		6+	2+	DC	LDD		5	2	10B3	CMPD	Extended	8	4
AD	JSR		7+	2+	DD	STD		5	2	10BC	CMPY		8	4
AE	LDX		5+	2+	DE	LDU		5	2	10BE	LDY		7	4
AF	STX	Indexed	5+	2+	DF	STU	Direct	5	2	10BF	STY	Extended	7	4
B0	SUBA	Extended	5	3	E0	SUBB	Indexed	4+	2+	10CE	LDS	Immed	4	4
B1	CMPA		5	3	E1	CMPB		4+	2+	10DF	STS	Direct	6	3
B2	SBCA		5	3	E2	SBCB		6+	2+	10EE	LDS	Indexed	6+	3+
B3	SUBD		7	3	E3	ADDD		4+	2+	10EF	STS	Indexed	6+	3+
B4	ANDA		5	3	E4	ANDB		4+	2+	10FE	LDS	Extended	7	4
B5	BITA		5	3	E5	BITB		4+	2+	10FF	STS	Extended	7	4
B6	LDA		5	3	E6	LDB		4+	2+	113F	SWI3	Inherent	20	2
B7	STA		5	3	E7	STB		4+	2+	1183	CMPU	Immed	5	4
B8	EORA		5	3	E8	EORB		4+	2+	118C	CMPS	Immed	5	4
B9	ADCA		5	3	E9	ADCB		5+	2+	1193	CMPU	Direct	7	3
BA	ORA		5	3	EA	ORB		4+	2+	119C	CMPS	Direct	7	3
BB	ADDA		5	3	EB	ADDB		5+	2+	11A3	CMPU	Indexed	7+	3+
BC	CMPX		7	3	EC	LDD		5+	2+	11AC	CMPS	Indexed	7+	3+
BD	JSR		8	3	ED	STD		5+	2+	11B3	CMPU	Extended	8	4
BE	LDX		6	3	EE	LDU		5+	2+	11BC	CMPS	Extended	8	4
BF	STX	Extended	6	3	EF	STU	Extended	5	3					
					F0	SUBB								
					F1	CMPB								
					F2	SBCB								
					F3	ADDD								
					F4	ANDB								
					F5	BITB								
					F6	LDB								
					F7	STB								
					F8	EORB								
					F9	ADCB								
					FA	ORB								
					FB	ADDB								
					FC	LDD	Extended	6	3					
					FD	STD	Extended	6	3					
					FE	LDU								
					FF	STU	Extended	6	3					

NOTE: All unused opcodes are both undefined and illegal

APPENDIX D PROGRAMMING AID

D.1 INTRODUCTION

This appendix contains a compilation of data that will assist you in programming the M6809 processor. Refer to Table D-1.

Table D-1. Programming Aid

Branch Instructions

Instruction	Forms	Addressing Mode			Description	5 H	3 N	2 Z	1 V	0 C						
		Relative														
		OP	-	#												
BCC	BCC LBCC	24 10 24	3 5(6) 4	2	Branch C=0 Long Branch C=0	• •	• •	• •	• •	• •						
BCS	BCS LBCS	25 10 25	3 5(6) 4	2	Branch C=1 Long Branch C=1	• •	• •	• •	• •	• •						
BEQ	BEQ LBEQ	27 10 27	3 5(6) 4	2	Branch Z=0 Long Branch Z=0	• •	• •	• •	• •	• •						
BGE	BGE LBGE	2C 10 2C	3 5(6) 4	2	Branch \geq Zero Long Branch \geq Zero	• •	• •	• •	• •	• •						
BGT	BGT LBGT	2E 10 2E	3 5(6) 4	2	Branch > Zero Long Branch > Zero	• •	• •	• •	• •	• •						
BHI	BHI LBHI	22 10 22	3 5(6) 4	2	Branch Higher Long Branch Higher	• •	• •	• •	• •	• •						
BHS	BHS LBHS	24 10 24	3 5(6) 4	2	Branch Higher or Same Long Branch Higher or Same	• •	• •	• •	• •	• •						
BLE	BLE LBLE	2F 10 2F	3 5(6) 4	2	Branch \leq Zero Long Branch \leq Zero	• •	• •	• •	• •	• •						
BLO	BLO LBLO	25 10 25	3 5(6) 4	2	Branch lower Long Branch Lower	• •	• •	• •	• •	• •						

Instruction	Forms	Addressing Mode			Description	5 H	3 N	2 Z	1 V	0 C						
		Relative														
		OP	-	#												
BLS	BLS LBLS	23 10 23	3 5(6) 4	2	Branch Lower or Same Long Branch Lower or Same	• •	• •	• •	• •	• •						
BLT	BLT LBLT	2D 10 2D	3 5(6) 4	2	Branch < Zero Long Branch < Zero	• •	• •	• •	• •	• •						
BMI	BMI LBMI	2B 10 2B	3 5(6) 4	2	Branch Minus Long Branch Minus	• •	• •	• •	• •	• •						
BNE	BNE LBNE	26 10 26	3 5(6) 4	2	Branch Z \neq 0 Long Branch Z \neq 0	• •	• •	• •	• •	• •						
BPL	BPL LBPL	2A 10 2A	?	2	Branch Plus Long Branch Plus	• •	• •	• •	• •	• •						
BRA	BRA LBRA	20 16	3 5	2	Branch Always Long Branch Always	• •	• •	• •	• •	• •						
BRN	BRN LBRN	21 10 21	3 5 4	2	Branch Never Long Branch Never	• •	• •	• •	• •	• •						
BSR	BSR LBSR	8D 17	7 9	2	Branch to Subroutine Long Branch to Subroutine	• •	• •	• •	• •	• •						
BVC	BVC LBVC	28 10 28	3 5(6) 4	2	Branch V=0 Long Branch V=0	• •	• •	• •	• •	• •						
BVS	BVS LBVS	29 10 29	3 5(6) 4	2	Branch V=1 Long Branch V=1	• •	• •	• •	• •	• •						

Table D-1. Programming Aid (Continued)

SIMPLE BRANCHES

	OP	~	#
BRA	20	3	2
LBRA	16	5	3
BRN	21	3	2
LBRN	1021	5	4
BSR	8D	7	2
LBSR	17	9	3

SIMPLE CONDITIONAL BRANCHES (Notes 1-4)

Test	True	OP	False	OP
N = 1	BMI	2B	BPL	2A
Z = 1	BEQ	27	BNE	26
V = 1	BVS	29	BVC	28
C = 1	BCS	25	BCC	24

SIGNED CONDITIONAL BRANCHES (Notes 1-4)

Test	True	OP	False	OP
r > m	BGT	2E	BLE	2F
r ≥ m	BGE	2C	BLT	2D
r = m	BEQ	27	BNE	26
r ≤ m	BLE	2F	BGT	2E
r < m	BLT	2D	BGE	2C

UNSIGNED CONDITIONAL BRANCHES (Notes 1-4)

Test	True	OP	False	OP
r > m	BHI	22	BLS	23
r ≥ m	BHS	24	BLO	25
r = m	BEQ	27	BNE	26
r ≤ m	BLS	23	BHI	22
r < m	BLO	25	BHS	24

Notes:

1. All conditional branches have both short and long variations.
2. All short branches are 2 bytes and require 3 cycles.
3. All conditional long branches are formed by prefixing the short branch opcode with \$10 and using a 16-bit destination offset.
4. All conditional long branches require 4 bytes and 6 cycles if the branch is taken or 5 cycles if the branch is not taken.

Table D-1. Programming Aid (Continued)

Instruction	Forms	Addressing Modes															Description	5	3	2	1	0	
		Immediate			Direct			Indexed			Extended			Inherent				H	N	Z	V	C	
		Op	-	#	Op	-	#	Op	-	#	Op	-	#	Op	-	#							
ABX																	3A	3	1				
ADC	ADCA ADCB	89 2 2 99 4 2 C9 2 2 D9 4 2	A9 4+ 2+ E9 4+ 2+	B9 5 3 F9 5 3																			
ADD	ADDA ADDB ADD	8B 2 2 9B 4 2 CB 2 2 DB 4 2 C3 4 3 D3 6 2	AB 4+ 2+ EB 4+ 2+ E3 6+ 2+	BB 5 3 FB 5 3 F3 7 3																			
AND	ANDA ANDB ANDCC	84 2 2 94 4 2 C4 2 2 D4 4 2 1C 3 2	A4 4+ 2+ E4 4+ 2+	B4 5 3 F4 5 3																			
ASL	ASLA ASLB ASL		08 6 2	68 6+ 2+							78 7 3			48 58 2 1					8 1 1 1 1 1	8 1 1 1 1 1	8 1 1 1 1 1	8 1 1 1 1 1	
ASR	ASRB ASR ASR		07 6 2	67 6+ 2+							77 7 3			47 57 2 1					8 1 1 1 1 1	8 1 1 1 1 1	8 1 1 1 1 1	8 1 1 1 1 1	
BIT	BITA BITB	85 2 2 95 4 2 C5 2 2 D5 4 2	A5 4+ 2+ E5 4+ 2+	B5 5 3 F5 5 3																			
CLR	CLRA CLRB CLR		OF 6 2	6F 6+ 2+							7F 7 3			4F 5F 2 1					0 1 0 0 0	0 1 0 0 0	0 1 0 0 0	0 1 0 0 0	
CMP	CMPA CMPB CMPD 83 CMPS 8C CMPU 83 CMPX 8C CMPY 8C	81 2 2 91 4 2 C1 2 2 D1 4 2 10 5 4 10 7 3 93 11 5 4 11 7 3 8C 11 5 4 11 7 3 93 8C 4 3 9C 6 2 10 5 4 10 7 3 AC	A1 4+ 2+ E1 4+ 2+ 10 7+ 3+ A3 11 7+ 3+ AC 11 7+ 3+ A3 6+ 2+ 10 7+ 3+ BC	B1 5 3 F1 5 3 10 8 4 B3 11 8 4 B3 11 8 4 BC 7 3 BC															8 1 1 1 1 1	8 1 1 1 1 1	8 1 1 1 1 1	8 1 1 1 1 1	
COM	COMA COMB COM		03 6 2	63 6+ 2+							73 7 3			43 53 2 1					1 A→A B→B M→M	1 A→A B→B M→M	1 A→A B→B M→M	1 A→A B→B M→M	
CWAI		3C ≥20	2																			7	
DAA														19 2 1									
DEC	DECA DEC B C		0A 6 2	6A 6+ 2+							7A 7 3			4A 5A 2 1					1 A-1→A B-1→B M-1→M	1 A-1→A B-1→B M-1→M	1 A-1→A B-1→B M-1→M	1 A-1→A B-1→B M-1→M	
EOR	EORA EORB	88 2 2 98 4 2 C8 2 2 D8 4 2	A8 4+ 2+ E8 4+ 2+	B8 5 3 F8 5 3															1 A¬M→A B¬M→B	1 A¬M→A B¬M→B	1 A¬M→A B¬M→B	1 A¬M→A B¬M→B	
EXG	R1, R2	1E 8 2																					
INC	INCA INCB INC		0C 6 2	6C 6+ 2+							7C 7 3			4C 5C 2 1					1 A+1→A B+1→B M+1→M	1 A+1→A B+1→B M+1→M	1 A+1→A B+1→B M+1→M	1 A+1→A B+1→B M+1→M	
JMP			0E 3 2	6E 3+ 2+							7E 4 3								EA ³ →PC	EA ³ →PC	EA ³ →PC	EA ³ →PC	
JSR			9D 7 2	AD 7+ 2+							BD 8 3								Jump to Subroutine	Jump to Subroutine	Jump to Subroutine	Jump to Subroutine	
LD	LDA LDB LDD LDS LDU LDX LDY	86 2 2 96 4 2 C6 2 2 D6 4 2 CC 3 3 DC 5 2 10 4 4 10 6 3 CE 3 3 DE 5 2 8E 3 3 9E 5 2 10 4 4 10 6 3 8E	A6 4+ 2+ E6 4+ 2+ EC 5+ 2+ 10 6+ 3+ EE 5+ 2+ AE 5+ 2+ 10 6+ 3+ AE	B6 5 3 F6 5 3 FC 6 3 10 7 4 FE 6 3 BE 6 3 10 7 4 BE														M→A M→B M:M+1→D M:M+1→S M:M+1→U M:M+1→X M:M+1→Y	M→A M→B M:M+1→D M:M+1→S M:M+1→U M:M+1→X M:M+1→Y	M→A M→B M:M+1→D M:M+1→S M:M+1→U M:M+1→X M:M+1→Y	M→A M→B M:M+1→D M:M+1→S M:M+1→U M:M+1→X M:M+1→Y	M→A M→B M:M+1→D M:M+1→S M:M+1→U M:M+1→X M:M+1→Y	
LEA	LEAS LEAU LEAX LEAY										32 4+ 2+ 33 4+ 2+ 30 4+ 2+ 31 4+ 2+								EA ³ →S EA ³ →U EA ³ →X EA ³ →Y	EA ³ →S EA ³ →U EA ³ →X EA ³ →Y	EA ³ →S EA ³ →U EA ³ →X EA ³ →Y	EA ³ →S EA ³ →U EA ³ →X EA ³ →Y	EA ³ →S EA ³ →U EA ³ →X EA ³ →Y

Legend:

OP Operation Code (Hexadecimal)

~ Number of MPU Cycles

Number of Program Bytes

+ Arithmetic Plus

- Arithmetic Minus

• Multiply

M Complement of M

→ Transfer Into

H Half-carry (from bit 3)

N Negative (sign bit)

Z Zero (Reset)

V Overflow, 2's complement

C Carry from ALU

† Test and set if true, cleared otherwise

• Not Affected

CC Condition Code Register

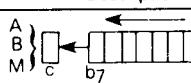
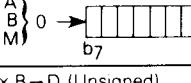
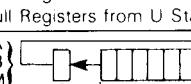
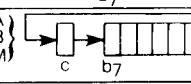
: Concatenation

V Logical or

Λ Logical and

✗ Logical Exclusive or

Table D-1. Programming Aid (Continued)

Instruction	Forms	Addressing Modes												Description	5	3	2	1	0					
		Immediate			Direct			Indexed ¹			Extended				H	N	Z	V	C					
		Op	-	#	Op	-	#	Op	-	#	Op	-	#	Op	-	#								
LSL	LSLA LSLB LSL				08	6	2	68	6+	2+	78	7	3	48 58	2	1		• 1 • 1 • 1	1 1 1	1 1 1	1 1 1	1 1 1		
LSR	LSRA LSRB LSR				04	6	2	64	6+	2+	74		3	44 54	2	1		0 1 0 1 0 1	1 1 1	1 1 1	1 1 1	1 1 1		
MUL														3D	11	1	A × B → D (Unsigned)							
NEG	NEGA NEGB NEG				00	6	2	60	6+	2+	70	7	3	40 50	2	1	A + 1 → A B + 1 → B M + 1 → M	8 1 8 1 8 1	1 1 1	1 1 1	1 1 1	1 1 1		
NOP														12	2	1	No Operation	• •	• •	• •	• •	• •		
OR	ORA ORB ORCC	8A CA 1A	2 2 3	2 2 2	9A DA EA	4 4 2	2 + 2+	AA BA FA	4+ 5 5	2+ 3 3					A V M → A B V M → B CC V IMM → CC				• 1 • 1 • 1	1 0 0				
PSH	PSHS PSHU	34 36	5+ 5+	4 4											Push Registers on S Stack Push Registers on U Stack				• •	• •	• •	• •	• •	
PUL	PULS PULU	35 37	5+ 5+	4 4											Pull Registers from S Stack Pull Registers from U Stack				• •	• •	• •	• •	• •	
ROL	ROLA ROLB ROL				09	6	2	69	6+	2+	79	7	3	49 59	2	1		• 1 • 1 • 1	1 1 1	1 1 1	1 1 1	1 1 1		
ROR	RORA RORB ROR				06	6	2	66	6+	2+	76	7	3	46 56	2	1		• 1 • 1 • 1	1 1 1	1 1 1	1 1 1	1 1 1		
RTI														3B	6/15	1	Return From Interrupt					7		
RTS														39	5	1	Return from Subroutine	• •	• •	• •	• •	• •		
SBC	SBCA SBCB	82 C2	2 2	2 2	92 D2	4 4	2 2	A2 E2	4+ 4+	2+ 2+	B2 F2	5 5	3 3			A - M → C → A B - M → C → B				8 1 8 1	1 1 1	1 1 1	1 1 1	1 1 1
SEX														1D	2	1	Sign Extend B into A	• 1	1	1	0	•		
ST	STA STB STD STS STU STX STY				97 D7 DD 10 DF DF 9F 10	4 4 5 6 EF 5 5 6	2 2 2 3 + 2 + 2 3	A7 E7 ED 10 EF FF AF 10	4+ 4+ 5+ 6+ 3+ 2+ 5+ 3+	2+ 2+ 2+ 3+ FF FF 2+ 3+	B7 F7 FD 10 FF 6 BF 10	5 5 6 7 4 3 6 4			A → M B → M D → M M + 1 S → M M + 1 U → M M + 1 X → M M + 1 Y → M M + 1				• 1 • 1 • 1 • 1 • 1 • 1 • 1 • 1	1 0 0 0 0 0 0 0	0 • • • 0 0 0 0	• • • • • • • •	• • • • • • • •	
SUB	SUBA SUBB SUBD	80 C0 83	2 2 4	2 2 3	90 D0 93	4 4 6	2 2 2	A0 E0 A3	4+ 4+ 6+	2+ 2+ 2+	B0 F0 B3	5 5 7	3 3 3			A - M → A B - M → B D - M: M + 1 → D				8 1 8 1 • 1	1 1 1	1 1 1	1 1 1	1 1 1
SWI	SWI ⁶ SWI2 ⁶ SWI3 ⁶													3F 10 3F 11 3F	19 20 20	1 1	Software Interrupt 1 Software Interrupt 2 Software Interrupt 3	• •	• •	• •	• •	• •		
SYNC														13	≥ 4	1	Synchronize to Interrupt	• •	• •	• •	• •	• •		
TFR	R1, R2	1F	6	2												R1 → R2 ²				• •	• •	• •	• •	• •
TST	TSTA TSTB TST				0D	6	2	6D	6+	2+	7D	7	3	4D 5D	2	1	Test A	• 1	1	0	•	•		
															2	1	Test B	• 1	1	0	•	•		
																3	Test M	• 1	1	0	•	•		

Notes:

- This column gives a base cycle and byte count. To obtain total count, add the values obtained from the INDEXED ADDRESSING MODE table, in Appendix F.
- R1 and R2 may be any pair of 8 bit or any pair of 16 bit registers.
The 8 bit registers are: A, B, CC, DP
The 16 bit registers are: X, Y, U, S, D, PC
- EA is the effective address.
- The PSH and PUL instructions require 5 cycles plus 1 cycle for each byte pushed or pulled.
- 5(6) means: 5 cycles if branch not taken, 6 cycles if taken (Branch instructions).
- SWI sets I and F bits. SWI2 and SWI3 do not affect I and F.
- Conditions Codes set as a direct result of the instruction.
- Value of half-carry flag is undefined.
- Special Case – Carry set if b7 is SET.

APPENDIX E

ASCII CHARACTER SET

E.1 INTRODUCTION

This appendix contains the standard 112 character ASCII character set (7-bit code).

E.2 CHARACTER REPRESENTATION AND CODE IDENTIFICATION

The ASCII character set is given in Figure E-1.

							0	0	0	0	1	0	1	0	1	0	1
							0	0	0	1	0	1	1	0	0	1	1
b7	b6	b5	b4	b3	b2	b1	Row	Column	0	1	2	3	4	5	6	7	
1	1	1	1	1	1	1	0	0	0	1	2	3	4	5	6	7	
0	0	0	0	0	0	0	0	NUL	DLE	SP	0	@	P	'	p		
0	0	0	1	1	1	1	1	SOH	DC1	!	1	A	Q	a	q		
0	0	1	0	2	2	2	2	STX	DC2	"	2	B	R	b	r		
0	0	1	1	3	3	3	3	ETX	DC3	#	3	C	S	c	s		
0	1	0	0	4	4	4	4	EOT	DC4	\$	4	D	T	d	t		
0	1	0	1	5	5	5	5	ENQ	NAK	%	5	E	U	e	u		
0	1	1	0	6	6	6	6	ACK	SYN	&	6	F	V	f	v		
0	1	1	1	7	7	7	7	BEL	ETB	'	7	G	W	g	w		
1	0	0	0	8	8	8	8	BS	CAN	(8	H	X	h	x		
1	0	0	1	9	9	9	9	HT	EM)	9	I	Y	i	y		
1	0	1	0	10	A	LF	10	SUB	*	:	J	Z	j	z			
1	0	1	1	11	B	VT	11	ESC	+	;	K	[k	{			
1	1	0	0	12	C	FF	12	FS	,	<	L	\	l				
1	1	0	1	13	D	CR	13	GS	-	=	M]	m	}			
1	1	1	0	14	E	SO	14	RS	.	>	N	^	n	~			
1	1	1	1	15	F	SI	15	US	/	?	O	_	o	DEL			

Figure E-1. ASCII Character Set

Each 7-bit character is represented with bit seven as the high-order bit and bit one as the low-order bit as shown in the following example:

b7	b6	b5	b4	b3	b2	b1	b0
1	0	0	0	0	0	0	1

The bit representation for the character "A" is developed from the bit pattern for bits seven through five found above the column designated 4 and the bit pattern for bits four through one found to the left of the row designated 1.

A hexadecimal notation is commonly used to indicate the code for each character. This is easily developed by assuming a logic zero in the non-existent bit eight position for the column numbers and using the hexadecimal number for the row numbers.

E.3 CONTROL CHARACTERS

The characters located in columns zero and one of Figure E-1 are considered control characters. By definition, these are characters whose occurrence in a particular context initiates, modifies, or stops an action that affects the recording, processing, transmission, or interpretation of data. Table E-1 provides the meanings of the control characters.

Table E-1. Control Characters

Mnemonic	Meaning	Mnemonic	Meaning
NUL	Null	DLE	Data Link Escape
SOH	Start of Heading	DC1	Device Control 1
STX	Start of Text	DC2	Device Control 2
ETX	End of Text	DC3	Device Control 3
EOT	End of Transmission	DC4	Device Control 4
ENQ	Enquiry	NAK	Negative Acknowledge
ACK	Acknowledge	SYN	Synchronous Idle
BEL	Bell	ETB	End of Transmission Block
BS	Backspace	CAN	Cancel
HT	Horizontal Tabulation	EM	End of Medium
LF	Line Feed	SUB	Substitute
VT	Vertical Tabulation	ESC	Escape
FF	Form Feed	FS	File Separator
CR	Carriage Return	GS	Group Separator
SO	Shift Out	RS	Record Separator
SI	Shift In	US	Unit Separator
		DEL	Delete

E.4 GRAPHIC CHARACTERS

The characters in columns two through seven are considered graphic characters. These characters have a visual representation which is normally displayed or printed. These characters and their names are given in Table E-2.

Table E-2. Graphic Characters

Symbol	Name
SP	Space (Normally Nonprinting)
!	Exclamation Point
"	Quotation Marks (Diaeresis)
#	Number Sign
\$	Dollar Sign
%	Percent Sign
&	Ampersand
'	Apostrophe (Closing Single Quotation Mark; Acute Accent)
(Opening Parenthesis
)	Closing Parenthesis
*	Asterisk
+	Plus
,	Comma (Cedilla)
-	Hyphen (Minus)
.	Period (Decimal Point)
/	Slant
0...9	Digits 0 Through 9
:	Colon
;	Semicolon
<	Less Than
=	Equals
>	Greater Than
?	Question Mark
@	Commercial At
A...Z	Uppercase Latin Letters A Through Z
[Opening Bracket
\	Reverse Slant
]	Closing Bracket
^	Circumflex
—	Underline
'	Opening Single Quotation Mark (Grave Accent)
a...z	Lowercase Latin Letters a Through z
{	Opening Brace
	Vertical Line
}	Closing Brace
~	Tilde

APPENDIX F OPCODE MAP

F.1 INTRODUCTION

This appendix contains the opcode map and additional information for calculating required machine cycles.

F.2 OPCODE MAP

Table F-1 is the opcode map for M6809 processors. The number(s) by each instruction indicates the number of machine cycles required to execute that instruction. When the number contains an “I” (e.g., 4 + I), it indicates that the indexed addressing mode is being used and that an additional number of machine cycles may be required. Refer to Table F-2 to determine the additional machine cycles to be added.

Some instructions in the opcode map have two numbers, the second one in parenthesis. This indicates that the instruction involves a branch. The parenthetical number applies if the branch is taken.

The “page 2, page 3” notation in column one means that all page 2 instructions are preceded by a hexadecimal 10 opcode and all page 3 instructions are preceded by a hexadecimal 11 opcode.

Table F-1. Opcode Map

		Most-Significant Four Bits															
		DIR	REL	ACCA	ACCB	IND	EXT	IMM	DIR	IND	EXT	IMM	DIR	IND	EXT		
0000	0001	0010	0011	0100	0101	0111	1000	1001	1010	1011	1100	1101	1110	1111			
0	1	2	3	4	5	6	7	8	A	B	C	D	E	F			
0000 0	NEG	PAGE2	3 BRA	4+1	2	2	6+1	7	2	4	4+1	5	2	4	4+1	5	
0001 1	—	PAGE3	3 BRN/5 LBRN	4+1	—	—	—	—	2	4	4+1	5	2	4	4+1	5	0
0010 2	—	NOP	3 BH/5(6) LBHI	4+1	—	—	—	—	2	4	4+1	5	2	4	4+1	5	
0011 3	COM	SYNC	3 BL/5(6) LBLS	4+1	2	2	6+1	7	4,6,6+1,7 / SUBD	5,7,7+1,8 / CMPD	5,7,7+1,8 / CMPU	4	6	6+1	7		1
0100 4	LSR	—	3 BHS	5+1/by PSHS	2	2	LSR	6+1	2	4	ANDA	5	2	4	4+1	5	2
0101 5	—	—	3 BLO	5+1/by PULS	—	—	—	—	2	4	BITA	5	2	4	4+1	5	3
0110 6	ROR	LBRA	3 BNE/5(6) LBNE	5+1/by PSHU	2	2	ROR	6+1	2	4	LDA	5	2	4	4+1	5	4
0111 7	ASR	LBSR	3 BEQ/5(6) LBEQ	5+1/by PULL	2	2	ASR	6+1	—	4	4+1	5	—	4	4+1	5	5
1000 8	ASL	(LSL)	3 BVC/5(6) LBVC	—	2	2	6+1	7	2	4	4+1	5	2	4	4+1	5	6
1001 9	ROL	DAA	3 BVS/5(6) LBVS	5	2	2	6+1	7	2	4	4+1	5	2	4	4+1	5	9
1010 A	DEC	ORCC	3 BMI/5(6) LBPL	6/15 RTI	RTS	ROL	—	—	ADCA	ADCA	ADCA	ADCA	ADCA	ADCA	ADCA	ADCA	A
1011 B	—	—	3 BGE/5(6) LBMI	RTI	DEC	—	—	—	ORA	ORA	ORA	ORA	ORA	ORA	ORA	ORA	B
1100 C	INC	ANDCC	3 BG/E/5(6) LBGE	CWAI	20	2	2	6+1	7	2	4	4+1	5	2	4	4+1	5
1101 D	TST	SEX	3 BLT/5(6) LBLT	MUL	11	2	2	6+1	7	BSR	7+1 JSR	7+1 JSR	—	5	5+1 STD	6	C
1110 E	JMP	EXG	3 BGT/5(6) LBGT	—	—	3+1	JMP	4	3,5,5+1,6 / LDY	4,6,6+1,7 / LDY	3,5,5+1,6 / LDY	4,6,6+1,7 / LDY	—	5,5+1,6 STU	5,5+1,6 STU	E	
1111 F	CLR	TFR	3 BLE/5(6) LBLE	SW/2/3	19/20/20	2	2	6+1	7	—	5,5+1,6 STX	6,6+1,7 STY	—	6,6+1,7 STU	6,6+1,7 STU	F	

Table F-2. Indexed Addressing Mode Data

Type	Forms	Non Indirect				Indirect			
		Assembler Form	Postbyte OP Code	x ~	+ #	Assembler Form	Postbyte OP Code	+ ~	+ #
Constant Offset From R (twos complement offset)	No Offset	.R	1RR00100	0	0	[.R]	1RR10100	3	0
	5 Bit Offset	n, R	0RRnnnnn	1	0	defaults to 8-bit			
	8 Bit Offset	n, R	1RR01000	1	1	[n, R]	1RR11000	4	1
	16 Bit Offset	n, R	1RR01001	4	2	[n, R]	1RR11001	7	2
Accumulator Offset From R (twos complement offset)	A — Register Offset	A, R	1RR00110	1	0	[A, R]	1RR10110	4	0
	B — Register Offset	B, R	1RR00101	1	0	[B, R]	1RR10101	4	0
	D — Register Offset	D, R	1RR01011	4	0	[D, R]	1RR11011	7	0
Auto Increment/Decrement R	Increment By 1	.R+	1RR00000	2	0	not allowed			
	Increment By 2	.R++	1RR00001	3	0	[.R++]	1RR10001	6	0
	Decrement By 1	.-R	1RR00010	2	0	not allowed			
	Decrement By 2	.--R	1RR00011	3	0	[--R]	1RR10011	6	0
Constant Offset From PC (twos complement offset)	8 Bit Offset	n, PCR	1XX01100	1	1	[n, PCR]	1XX11100	4	1
	16 Bit Offset	n, PCR	1XX01101	5	2	[n, PCR]	1XX11101	8	2
Extended Indirect	16 Bit Address	—	—	—	—	[n]	10011111	5	2

R = X, Y, U or S X = 00 Y = 01
X = Don't Care U = 10 S = 11

⁺ and ⁺ indicate the number of additional cycles and bytes for the particular variation.
#

APPENDIX G PIN ASSIGNMENTS

G.1 INTRODUCTION

This appendix is provided for a quick reference of the pin assignments for the MC6809 and MC6809E processors. Refer to Figure G-1. Descriptions of these pin assignments are given in Section 1.

MC6809		MC6809E	
VSS	1	40	HALT
NMI	2	39	XTAL
IRQ	3	38	EXTAL
FIRQ	4	37	RESET
BS	5	36	MRDY
BA	6	35	Q
VCC	7	34	E
A0	8	33	DMA/BREQ
A1	9	32	R/W
A2	10	31	D0
A3	11	30	D1
A4	12	29	D2
A5	13	28	D3
A6	14	27	D4
A7	15	26	D5
A8	16	25	D6
A9	17	24	D7
A10	18	23	A15
A11	19	22	A14
A12	20	21	A13

Figure G-1. Pin Assignments

APPENDIX H

CONVERSION TABLES

H.1 INTRODUCTION

This appendix provides some conversion tables for your convenience.

H.2 POWERS OF 2, POWERS OF 16

Refer to Table H-1.

Table H-1. Powers of 2; Powers of 16

16^m m =	2^n n =	Value	16^m m =	2^n n =	Value
0	0	1	4	16	65,536
—	1	2	—	17	131,072
—	2	4	—	18	262,144
—	3	8	—	19	524,288
1	4	16	5	20	1,048,576
—	5	32	—	21	2,097,152
—	6	64	—	22	4,194,304
—	7	128	—	23	8,388,608
2	8	256	6	24	16,777,216
—	9	512	—	25	33,554,432
—	10	1,024	—	26	67,108,864
—	11	2,048	—	27	134,217,728
3	12	4,096	7	28	268,435,456
—	13	8,192	—	29	536,870,912
—	14	16,384	—	30	1,073,741,824
—	15	32,768	—	31	2,147,483,648

H.3 HEXADECIMAL AND DECIMAL CONVERSION

Table H-2 is a chart that can be used for converting numbers from either hexadecimal to decimal or decimal to hexadecimal.

H.3.1 CONVERTING HEXADECIMAL TO DECIMAL. Find the decimal weights for corresponding hexadecimal characters beginning with the least-significant character. The sum of the decimal weights is the decimal value of the hexadecimal number.

H.3.2 CONVERTING DECIMAL TO HEXADECIMAL. Find the highest decimal value in the table which is lower than or equal to the decimal number to be converted. The corresponding hexadecimal character is the most-significant digit of the final number. Subtract the decimal value found from the decimal number to be converted. Repeat the above step to determine the hexadecimal character. Repeat this process to find the subsequent hexadecimal numbers.

Table H-2. Hexadecimal and Decimal Conversion Chart

15		Byte		8		7		Byte		0	
15	Char	12	11	Char	8	7	Char	4	3	Char	0
Hex	Dec	Hex	Dec	Hex	Dec	Hex	Dec	Hex	Dec	Hex	Dec
0	0	0	0	0	0	0	0	0	0	0	0
1	4,096	1	256	1	16	1	1	1	1	1	1
2	8,192	2	512	3	32	2	2	2	2	2	2
3	12,288	3	768	3	48	3	3	3	3	3	3
4	16,384	4	1,024	4	64	4	4	4	4	4	4
5	20,480	5	1,280	5	80	5	5	5	5	5	5
6	24,576	6	1,536	6	96	6	6	6	6	6	6
7	28,672	7	1,792	7	112	7	7	7	7	7	7
8	32,768	8	2,048	8	128	8	8	8	8	8	8
9	36,864	9	2,304	9	144	9	9	9	9	9	9
A	40,960	A	2,560	A	160	A	10				
B	45,056	B	2,816	B	176	B	11				
C	49,152	C	3,072	C	192	C	12				
D	53,248	D	3,328	D	208	D	13				
E	57,344	E	3,584	E	224	E	14				
F	61,440	F	3,840	F	240	F	15				