



Arrays

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Outline

- ❑ Array ADT
- ❑ Matrix
- ❑ Symmetric Matrix
- ❑ Triangular Matrix
- ❑ Symmetric Band Matrix
- ❑ Sparse Matrix

Representation, Transposing

Arrays

- Array:
a set of pairs (**index** and **value**)
- data structure:
For each index, there is a value associated with that index.
- representation (possible):
implemented by using consecutive memory.

The Array ADT

- ▣ **Objects:** A set of pairs $\langle \text{index}, \text{value} \rangle$ where for each value of index there is a value from the set item. **Index** is a finite ordered set of one or more dimensions, for example, $\{0, \dots, n-1\}$ for one dimension, $\{(0,0),(0,1),(0,2),(1,0),(1,1),(1,2),(2,0),(2,1),(2,2)\}$ for two dimensions, etc.

Methods:

for all A Array, i index, x item, j, size integer

Array Create(j, list)

// **return** an array of **j dimensions** where **list** is a **j-tuple** whose **kth element** is the

//**size** of the **kth** dimension. Items are undefined.

Item Retrieve(A, i)

// **if** (i index) **return** the item associated with index value i in array A

// **else return** error

Array Store(A, i, x)

// **if** (i in index) **return** an array that is identical to array A except the new pair

// $\langle i, x \rangle$ has been inserted **else return** error



Matrices

- Two-dimensional arrays are a particularly common representation for matrices.
- A matrix, also referred to as a general matrix, is an m by n ordered collection of numbers. It is represented symbolically as:

$$A = \begin{bmatrix} a_{11} & \cdot & \cdot & \cdot & a_{1n} \\ \cdot & & & & \cdot \\ \cdot & & & & \cdot \\ \cdot & & & & \cdot \\ a_{m1} & \cdot & \cdot & \cdot & a_{mn} \end{bmatrix}$$

- where the matrix is named **A** and has m rows and n columns. And a_{ij} is the element in i th row and j th column of matrix A.

Matrices

- A matrix appears as two-dimensional, but physically it is stored in a linear fashion. How to represent this two-dimensional array?

$$A = \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{bmatrix}$$

Matrices

- Common ways to index into multi-dimensional arrays include:
- Row-major order:

The elements of each row are stored in order.

1	2	3	4	5	6	7	8	9
---	---	---	---	---	---	---	---	---

- Column-major order:

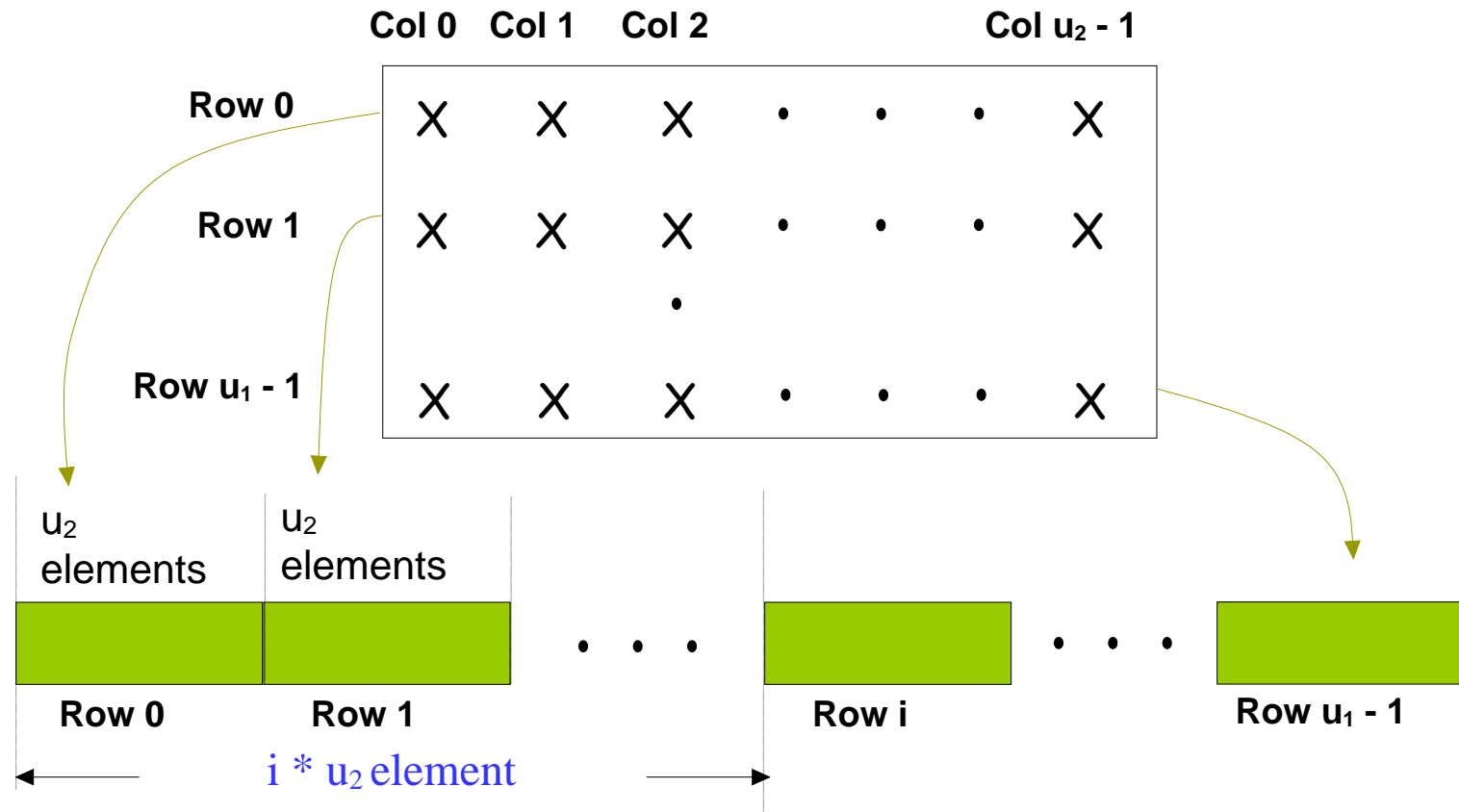
The elements of each column are stored in order.

1	4	7	2	5	8	3	6	9
---	---	---	---	---	---	---	---	---



Matrices

Row-major order:



Matrices

- So, in order to map logical view to physical structure, we need indexing formula.
 - Row-major order: Assume that the base address is at M, the address of a_{ij} will be obtained as
$$\text{Address}(a_{ij}) = M + (i-1)*n + j - 1$$
 - Column-major order: Considering the base address at M, the formula will stand as
$$\text{Address}(a_{ij}) = M + (j-1)*n + i - 1$$

Symmetric Matrix

- The matrix **A** is symmetric if it has the property **A** equal to **A**^T, which means:
 - It has the same number of rows as it has columns; that is, it has n rows and n columns.
 - The value of every element a_{ij} on one side of the main diagonal equals its mirror image a_{ji} on the other side: a_{ij} equal to a_{ji} .

Symmetric Matrix

- The following matrix illustrates a symmetric matrix of order n ; that is, it has n rows and n columns. The subscripts on each side of the diagonal appear the same to show which elements are equal:

$$A = \begin{bmatrix} a_{11} & a_{21} & a_{31} & \dots & \dots & a_{n1} \\ a_{21} & a_{22} & a_{32} & & & \vdots \\ a_{31} & a_{32} & a_{33} & & & \vdots \\ \vdots & & & \ddots & & \vdots \\ \vdots & & & & \ddots & \vdots \\ a_{n1} & \vdots & \vdots & \vdots & \vdots & a_{nn} \end{bmatrix}$$

Symmetric Matrix

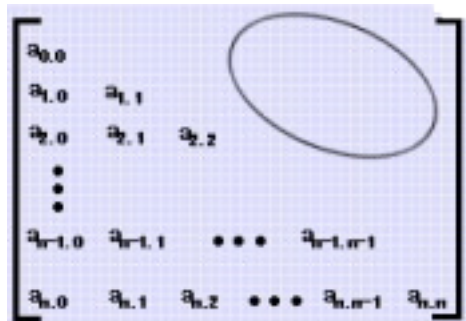
- When a symmetric matrix is stored in lower-packed storage mode, the lower triangular part of the symmetric matrix is stored, including the diagonal, in a one-dimensional array.
- The lower triangle can be packed by row or columns. The matrix is packed sequentially row by row (column by column) in $n(n+1)/2$ elements of a one-dimensional array.
- When the matrix is packed sequentially row by row, to calculate the location k of element a_{ij} of matrix **A** in an array, use the following formula:

$$k = i*(i-1)/2 + j - 1 \quad i \geq j, \text{ lower triangular part}$$

$$k = j*(j-1)/2 + i - 1 \quad i < j, \text{ upper triangular part}$$

Triangular Matrix

A matrix of the form



The diagram shows a square matrix of size $n \times n$ enclosed in large square brackets. The elements are arranged as follows:

- Row 0: $a_{0,0}$
- Row 1: $a_{1,0}$ followed by $a_{1,1}$
- Row 2: $a_{2,0}$ followed by $a_{2,1}$ followed by $a_{2,2}$
- Row 3: Three vertical dots (\vdots)
- Row $r-1$: $a_{r-1,0}$ followed by $a_{r-1,1}$ followed by three dots (\dots) followed by $a_{r-1,r-1}$
- Row n : $a_{n,0}$ followed by $a_{n,1}$ followed by $a_{n,2}$ followed by three dots (\dots) followed by $a_{n,n-1}$ followed by $a_{n,n}$

An empty oval is drawn in the upper right portion of the matrix, indicating that all elements above the main diagonal are zero.

is called a **triangular matrix**.

Triangular Matrix

- There are two types of triangular matrices: upper triangular matrix and lower triangular matrix. Triangular matrices have the same number of rows as they have columns; that is, they have n rows and n columns.
- A matrix **U** is an upper triangular matrix if its nonzero elements are found only in the upper triangle of the matrix, including the main diagonal; that is: u_{ij} equal to 0 (or constant C) if i greater than j
- A matrix **L** is an lower triangular matrix if its nonzero elements are found only in the lower triangle of the matrix, including the main diagonal; that is: l_{ij} equal to 0 (or constant C) if i less than j



Triangular Matrix

- The following matrices, **U** and **L**, illustrate upper and lower triangular matrices of order n, respectively:

$$\mathbf{U} = \begin{bmatrix} u_{11} & u_{12} & u_{13} & . & . & . & u_{1n} \\ 0 & u_{22} & u_{23} & & & & . \\ 0 & 0 & u_{33} & & & & . \\ . & & & . & & & . \\ . & & & & . & & . \\ . & & & & & . & . \\ 0 & . & . & . & . & 0 & u_{nn} \end{bmatrix} \quad \mathbf{L} = \begin{bmatrix} l_{11} & 0 & 0 & . & . & . & 0 \\ l_{21} & l_{22} & 0 & & & & . \\ l_{31} & l_{32} & l_{33} & & & & . \\ . & & & . & & & . \\ . & & & & . & & . \\ . & & & & & . & 0 \\ l_{n2} & . & . & . & . & . & l_{nn} \end{bmatrix}$$

Triangular Matrix

- ❑ When a lower-triangular matrix is stored in lower-triangular-packed storage mode, the lower triangle of the matrix is stored, including the diagonal, in a one-dimensional array. The lower triangle is packed by row or by columns. The elements are packed sequentially, row by row (column by column), in $n(n+1)/2$ elements of a one-dimensional array. To calculate the location of each element of the triangular matrix in the array, use the technique described in Symmetric Matrix.
- ❑ When an upper-triangular matrix is stored in upper-triangular-packed storage mode, the upper triangle of the matrix is stored, including the diagonal, in a one-dimensional array.

Symmetric Band Matrix

- A symmetric band matrix is a symmetric matrix whose nonzero elements are arranged uniformly near the diagonal, such that: a_{ij} equal to 0 if $|i-j|$ greater than k , where k is the half band width.

Symmetric Band Matrix

- The following matrix illustrates a symmetric band matrix of order n , where the half band width k equal to $q-1$:

$$A = \begin{array}{c} \begin{array}{c} | \quad \leftarrow k \rightarrow \quad | \\ a_{11} \ a_{21} \ a_{31} \ \dots \ a_{q1} \ 0 \ \dots \ 0 \\ a_{21} \ a_{22} \ a_{32} \ \dots \ \dots \ 0 \ \dots \ . \\ a_{31} \ a_{32} \ a_{33} \ \dots \ \dots \ \dots \ 0 \ \dots \ . \\ . \ \dots \ . \ \dots \ . \ \dots \ . \ \dots \ . \ \dots \ . \ \dots \ . \\ . \ \dots \ . \ \dots \ . \ \dots \ . \ \dots \ . \ \dots \ . \ \dots \ . \\ a_{q1} \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ . \\ 0 \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ . \\ . \ 0 \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ . \\ . \ \dots \ 0 \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ . \\ 0 \ \dots \ \dots \ 0 \ \dots \ \dots \ \dots \ \dots \ \dots \ a_{nn} \end{array} \end{array}$$

- Only the band elements of the symmetric band matrix are stored.

Sparse Matrix

A sparse matrix is a matrix having a relatively small number of nonzero elements.

	col 1	col 2	col 3
row 1	-27	3	4
row 2	6	82	-2
row 3	109	-64	11
row 4	12	8	9
row 5	48	27	47

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	col1	col2	col3	col4	col5	col6
row0	15	0	0	22	0	15
row1	0	11	3	0	0	0
row2	0	0	0	6	0	0
row3	0	0	0	0	0	0
row4	91	0	0	0	0	0
row5	0	0	28	0	0	0

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sparse matrix
data structure?

Sparse Matrix Representation

- The standard representation of a matrix is a two dimensional array defined as

$a[\text{MAX_ROWS}][\text{MAX_COLS}]$

- We can locate quickly any element by writing **$a[i][j]$**

- Sparse matrix wastes space

- We must consider alternate forms of representation.
- Our representation of sparse matrices should store only nonzero elements.
- Each element is characterized by <row, col, value>.



Transposing A Matrix

- Transpose a Matrix
 - For each row i
 - take element $\langle i, j, \text{value} \rangle$ and store it in element $\langle j, i, \text{value} \rangle$ of the transpose.
 - difficulty: where to put $\langle j, i, \text{value} \rangle$
 - $(0, 0, 15) \implies (0, 0, 15)$
 - $(0, 3, 22) \implies (3, 0, 22)$
 - $(0, 5, -15) \implies (5, 0, -15)$
 - $(1, 1, 11) \implies (1, 1, 11)$
 - Move elements down very often.
 - For all elements in column j ,
 - place element $\langle i, j, \text{value} \rangle$ in element $\langle j, i, \text{value} \rangle$



Transposing A Matrix

Assign

$A[i][j]$ to $B[j][i]$

place element $\langle i, j, \text{value} \rangle$
in element $\langle j, i, \text{value} \rangle$

For all columns i

For all elements in column j

Scan the array “columns”
times. The array has
“elements” elements.

```
void transpose(term a[], term b[])
/* b is set to the transpose of a */
{
    int n,i,j, currentb;
    n = a[0].value;          /* total number of elements */
    b[0].row = a[0].col; /* rows in b = columns in a */
    b[0].col = a[0].row; /* columns in b = rows in a */
    b[0].value = n;
    if (n > 0 ) { /* non zero matrix */
        currentb = 1;
        for (i = 0; i < a[0].col; i++)
            /* transpose by the columns in a */
            for (j = 1; j <= n; j++)
                /* find elements from the current column */
                if (a[j].col == i) {
                    /* element is in current column, add it to b */
                    b[currentb].row = a[j].col;
                    b[currentb].col = a[j].row;
                    b[currentb].value = a[j].value;
                    currentb++;
                }
            }
    }
}
```

$\Rightarrow O(\text{columns} * \text{elements})$



EX: A[6][6] transpose to
B[6][6]

Matrix A

Row Col Value

a[0]	6	6	8
[1]	0	0	15
[2]	0	3	22
[3]	0	5	-15
[4]	1	1	11
[5]	1	2	3
[6]	2	3	-6
[7]	4	0	91
[8]	5	2	28

Row Col Value

0	6	6	8
1	0	0	15
2	0	4	91
3	1	1	11

i=1 j=8
a[i].col = 2 != i

```
void transpose(term a[], term b[])
/* b is set to the transpose of a */
{
    int n,i,j, currentb;
    n = a[0].value; /* total number of elements */
    b[0].row = a[0].col; /* rows in b = columns in a */
    b[0].col = a[0].row; /* columns in b = rows in a */
    b[0].value = n;
    if (n > 0 ) { /* non zero matrix */
        currentb = 1;
        for (i = 0; i < a[0].col; i++)
            /* transpose by the columns in a */
            for (j = 1; j <= n; j++)
                /* find elements from the current column */
                if (a[j].col == i) {
                    /* element is in current column, add it to b */
                    b[currentb].row = a[j].col;
                    b[currentb].col = a[j].row;
                    b[currentb].value = a[j].value;
                    currentb++;
                }
    }
}
```

Set Up row & column
in B[6][6]

And So on...



Reference

- ▣ 《数据结构（C语言版）》，严蔚敏，吴伟民编著，清华大学出版社，1997年第1版, P91-99



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