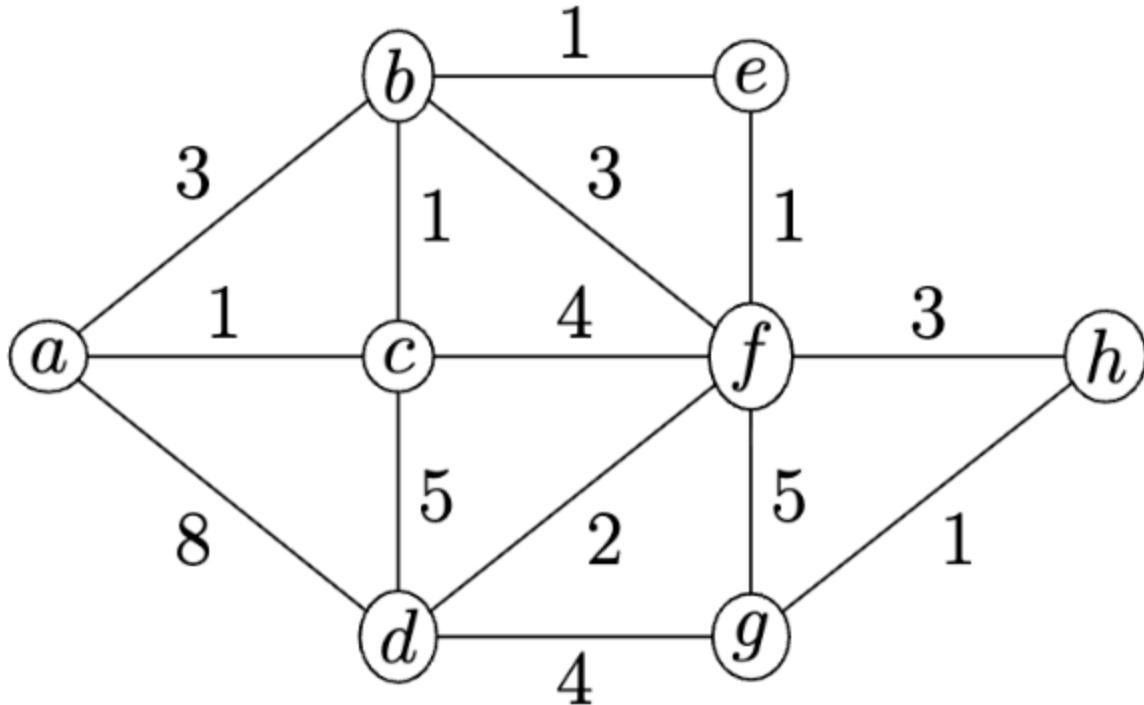


**Exercise Sheet 15****Question 15.1 (0.25 marks)**

Compute the minimum spanning tree of the following weighted graph both with Prim's and Kruskal's algorithm. List the edges in the order considered, draw the tree and calculate its weight.

**Kruskal's algorithm**

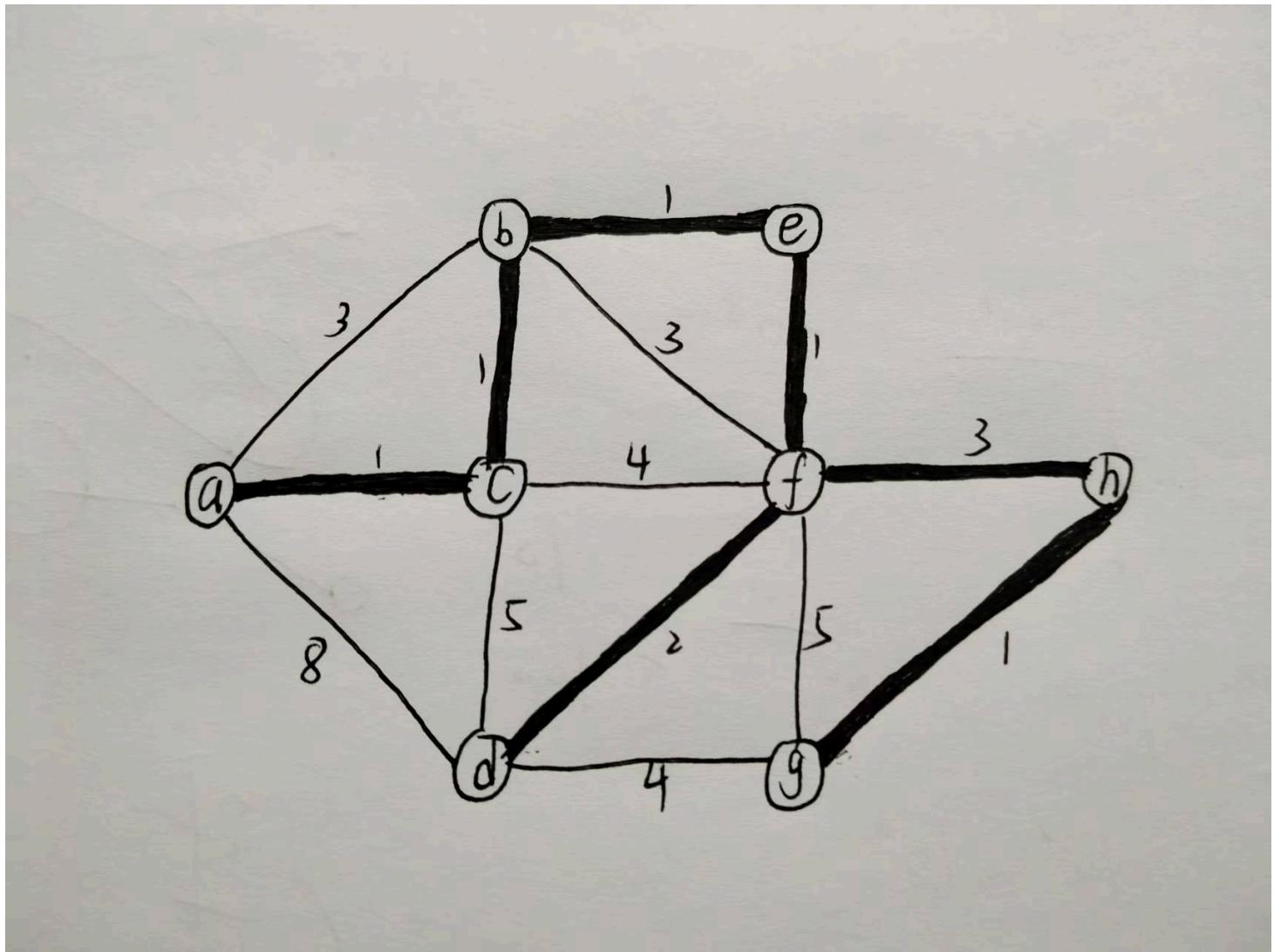
The idea of Kruskal's algorithm is to sort all edges in non-decreasing order of their weight, choose the smallest edge whose two vertices belong to different trees and merge two trees.

1. Sort edges by weight:  $(b, e)=1, (a, c)=1, (b, c)=1, (e, f)=1, (g, h)=1, (d, f)=2, (a, b)=3, (b, f)=3, (f, h)=3, (c, f)=4, (d, g)=4, (c, d)=5, (f, g)=5, (a, d)=8$
2. Initialize forest with each vertex as a separate tree.
3. choose edges in sorted order, adding them to the forest if they connect two different trees:
  - Add  $(b, e)=1$
  - Add  $(a, c)=1$
  - Add  $(b, c)=1$
  - Add  $(e, f)=1$

- Add (g, h)=1
- Add (d, f)=2
- Skip (a, b)=3 (creates a cycle)
- Skip (b, f)=3 (creates a cycle)
- Add (f, h)=3
- Skip (c, f)=4 (creates a cycle)
- Skip (d, g)=4 (creates a cycle)
- Skip (c, d)=5 (creates a cycle)
- Skip (f, g)=5 (creates a cycle)
- Skip (a, d)=8 (creates a cycle)

total weight =  $1 + 1 + 1 + 1 + 1 + 2 + 3 = 10$

The result is:



## Prim's algorithm

The idea of Prim's algorithm is to start from an arbitrary vertex and grow the minimum spanning tree by adding the smallest edge that connects a vertex in the tree to a vertex outside the tree.

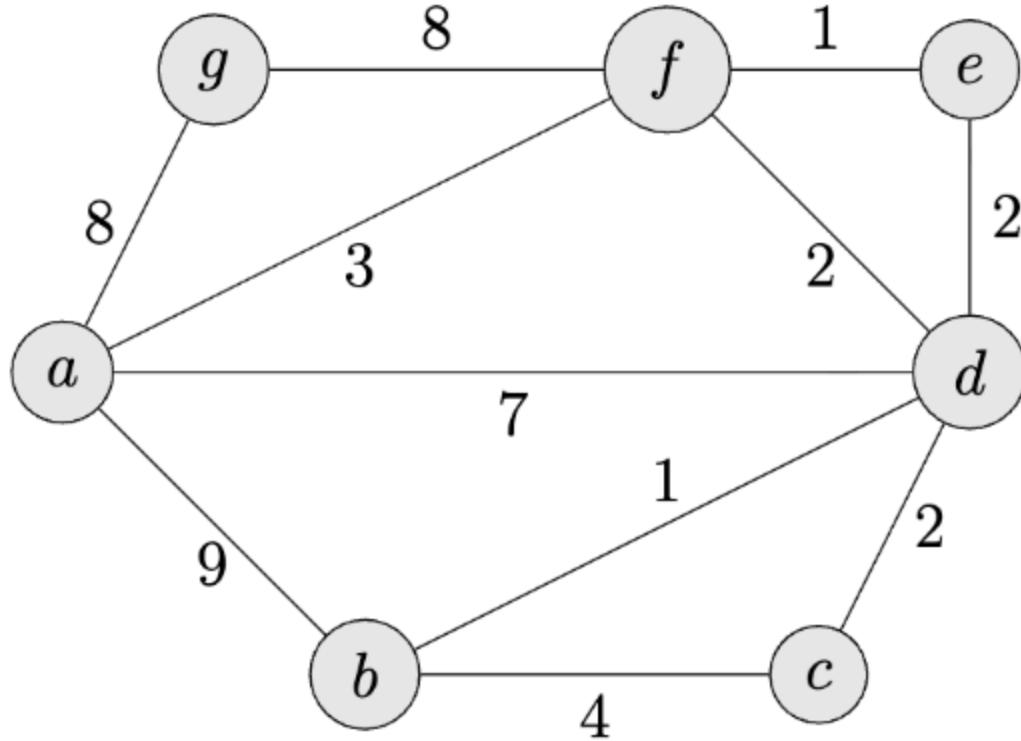
Edges: (b, e)=1, (a, c)=1, (b, c)=1, (e, f)=1, (g, h)=1, (d, f)=2, (a, b)=3, (b, f)=3, (f, h)=3, (c, f)=4, (d, g)=4, (c, d)=5, (f, g)=5, (a, d)=8

1. Start from vertex a.
2. Add edge (a, c)=1
3. Add edge (b, c)=1
4. Add edge (b, e)=1
5. Add edge (e, f)=1
6. Add edge (d, f)=2
7. Add edge (f, h)=3
8. Add edge (g, h)=1

The result is the same as Kruskal's algorithm, with total weight = 10.

## Question 15.2 (0.25 marks)

Execute Dijkstra's algorithm on the following weighted graph to find a shortest path from vertex a to c. Show for each iteration of the while loop which vertex is added to the set S and how the distance estimates of adjacent vertices are being refined.



Edges:  $(a, b)=9$ ,  $(a, f)=3$ ,  $(a, g)=8$ ,  $(b, c)=4$ ,  $(b, d)=1$ ,  $(a, d)=7$ ,  $(c, d)=2$ ,  $(g, f)=8$ ,  $(f, d)=2$ ,  $(f, e)=1$ ,  $(e, d)=2$

Initialization:

- $S = \{\}$
- $Q = \{a=0, b=\infty, c=\infty, d=\infty, e=\infty, f=\infty, g=\infty\}$
- Distance estimates:  $d(a)=0$ ,  $d(b)=\infty$ ,  $d(c)=\infty$ ,  $d(d)=\infty$ ,  $d(e)=\infty$ ,  $d(f)=\infty$ ,  $d(g)=\infty$

1. Add a to S:

- Update distances:  $d(b)=9$ ,  $d(f)=3$ ,  $d(g)=8$ ,  $d(d)=7$
- $S = \{a\}$
- $Q = \{f=3, d=7, g=8, b=9, e=\infty, c=\infty\}$

2. Add f to S

- Update distances:  $d(d)=5$ ,  $d(e)=4$
- $S = \{a, f\}$
- $Q = \{e=4, d=5, g=8, b=9, c=\infty\}$

3. Add e to S

- Update distances: none
- $S = \{a, f, e\}$
- $Q = \{d=5, g=8, b=9, c=\infty\}$

4. Add d to S

- Update distances:  $d(c)=7$ ,  $d(b)=6$
- $S = \{a, f, e, d\}$

- $Q = \{b=6, c=7, g=8\}$
5. Add b to S
- Update distances: none
  - $S = \{a, f, e, d, b\}$
  - $Q = \{c=7, g=8\}$

6. Add c to S
- Update distances: none
  - $S = \{a, f, e, d, b, c\}$
  - $Q = \{g=8\}$

The shortest path from a to c is a  $\rightarrow$  f  $\rightarrow$  d  $\rightarrow$  c with a total weight of 7.

## Question 15.3 (0.5 marks)

A precondition for Dijkstra's algorithm is that all edges of the directed graph under consideration have non-negative weight. Somebody on the internet claims that the algorithm works for graphs with negative edge weights as well: just add an appropriate constant  $c$  to each edge weight to make all weights positive, then run Dijkstra's algorithm, and finally remove the constants from the shortest paths computed. Give a directed acyclic graph as a counterexample to falsify this claim. Explain in your own words what goes wrong.

### Answer:

This claim is wrong because adding a constant to each edge weight can change the relative weights of different paths with different numbers of edges, leading to incorrect shortest path calculations.

Consider the following directed acyclic graph:

Edges:  $(s, b) = 2, (s, a) = 1, (a, b) = -2$

If we run Dijkstra's algorithm directly on this graph starting from vertex s, we find the shortest path to b is  $s \rightarrow a \rightarrow b$  with a total weight of -1.

Now, let's add a constant  $c = 5$  to each edge weight to make all weights positive:

Edges after adding  $c$ :  $(s, b) = 7, (s, a) = 6, (a, b) = 3$

Running Dijkstra's algorithm on this modified graph starting from vertex s, we find the shortest path to b is  $s \rightarrow b$  with a total weight of 7.

This is because the originally correct path include two edges, leading to more punishment when we add the constant, while the direct edge from s to b is less affected.