

Solutions for Exercise Sheet 2

Handout: September 16th — Deadline: September 23rd, 4pm

Question 2.1 (0.3 marks)

Express the following running times in Θ -notation. Justify your answer by referring to the definition of Θ (i.e. work out suitable c_1, c_2, n_0).

a) $3n^2 + 5n - 2 = \Theta(n^2)$.

To see this, we need to show that $c_1 n^2 \leq 3n^2 + 5n - 2 \leq c_2 n^2$ for all $n \geq n_0$ and suitable positive constants c_1, c_2 and a suitable n_0 . Divide by n^2 and we get $c_1 \leq 3 + \frac{5}{n} - \frac{2}{n^2} \leq c_2$ which is true, for example, for $c_1 := 3, c_2 := 4$, and $n_0 := 5$.

b) $42 = \Theta(1)$.

To see this, we need to show that $c_1 \cdot 1 \leq 42 \leq c_2 \cdot 1$ for all $n \geq n_0$ and suitable positive constants c_1, c_2 and a suitable n_0 . We can choose, for example, $c_1 := 42, c_2 := 42$, and $n_0 := 1$.

c) $4n^2 \cdot (1 + \log n) - 2n^2 = \Theta(n^2 \log n)$.

To see this, we need to show that $c_1 n^2 \log n \leq 4n^2 \cdot (1 + \log n) - 2n^2 \leq c_2 n^2 \log n$ for all $n \geq n_0$ and suitable positive constants c_1, c_2 and a suitable n_0 . Divide by $n^2 \log n$ and we get $c_1 \leq 4 \cdot \left(\frac{1}{\log n} + 1 \right) - \frac{2}{\log n} \leq c_2$, which we can rewrite as $c_1 \leq 4 + \frac{2}{\log n} \leq c_2$. This is true, for example, for $n_0 := 2$ (thus $\log n \geq 1$ for $n \geq n_0$) and $c_1 := 2, c_2 := 8$.

Question 2.2 (0.7 marks)

(a) Indicate for each pair of functions $f(n), g(n)$ in the following table whether $f(n)$ is O, o, Ω, ω , or Θ of $g(n)$ by writing “yes” or “no” in each box.

Solutions: (a) To fill in the table, it helps to remember that Θ (“equal asymptotic growth”) holds if and only if both O (“grows at most as fast as”) and Ω (“grows at least as fast as”) hold. Also o (“grows slower than”) implies O (“grows at most as fast as”) and ω (“grows faster than”) implies Ω (“grows at least as fast as”). So the answer in all cases below is either “ o and O ” or “ ω and Ω ” or “ Θ and O and Ω ”.

The hint explains that $\log n$ grows slower than \sqrt{n} , hence $\log n = o(\sqrt{n})$ and, consequently, $\log n = O(\sqrt{n})$ (if it grows slower than \sqrt{n} , it also grows at most as fast as \sqrt{n}). The other symbols Θ, Ω, ω do not apply.

Obviously, n grows faster than \sqrt{n} , so we need to put $n = \omega(\sqrt{n})$ and, consequently, $n = \Omega(\sqrt{n})$ (if it grows faster than \sqrt{n} , it also grows at least as fast as \sqrt{n}). None of the other symbols Θ, O, o apply.

For the third line, n grows slower than $n \log n$, thus $n = o(n \log n)$ and $n = O(n \log n)$.

In row 4, the two expressions are asymptotically the same since by the hint, $(\log n)^3 = o(n^2)$ and hence $n^2 + (\log n)^3 = O(n^2)$. Alternatively, we can argue that $n^2 \leq n^2 + (\log n)^3 \leq c_2 n^2$

for some constant c_2 , some n_0 and all $n \geq n_0$ (for example, $c_2 := 2$ and $n_0 := 100$). Hence $n^2 = \Theta(n^2 + (\log n)^3)$ and this implies that O and Ω also hold. (Remember that every time you tick Θ , you also need to tick O and Ω , but don't tick o and ω .)

For row 5, we can use the hint that tells us that every polynomial function grows slower than any exponential function. Hence, $2^n = \omega(n^3)$, thus $2^n = \Omega(n^3)$ and none of the other symbols apply.

For row 6, we have $2^{n/2} = o(2^n)$ as the former is the square root of the latter. Another justification for o is that dividing $2^{n/2}$ by 2^n gives $2^{n/2}/2^n = 2^{-n/2}$, which goes to 0 as n grows.

In the final row, we use that two logarithms of n have the same order of growth as $\log_x(n) = \log_y(n)/\log_y(x)$ and $\log_y(x)$ is constant (for $x, y > 1$). Hence $\log_2 n$ and $\log_{10} n$ have the same order of growth. We can put Θ and this also implies that we need to tick O and Ω .

$f(n)$	$g(n)$	O	o	Ω	ω	Θ
$\log n$	\sqrt{n}	yes	yes	no	no	no
n	\sqrt{n}	no	no	yes	yes	no
n	$n \log n$	yes	yes	no	no	no
n^2	$n^2 + (\log n)^3$	yes	no	yes	no	yes
2^n	n^3	no	no	yes	yes	no
$2^{n/2}$	2^n	yes	yes	no	no	no
$\log_2 n$	$\log_{10} n$	yes	no	yes	no	yes

Hints: the book states that every polynomial of $\log n$ grows strictly slower than every polynomial n^ε , for constant $\varepsilon > 0$. For example, $(\log n)^{100} = o(n^{0.01})$. Likewise, every polynomial grows slower than every exponential function 2^{n^ε} , for example $n^{100} = o(2^{n^{0.01}})$.

To convert the base of a logarithm, use $\log_x(n) = \log_y(n)/\log_y(x)$.

Solution: Question 2.3 (0.3 marks)

State the number of “foo” operations for each of the following algorithms in Θ -notation. Pay attention to indentation and how long loops are run for. Justify your answer by stating constants $c_1, c_2, n_0 > 0$ from the definition of $\Theta(g(n))$ in your answer.

Example: Line 1 is executed once and line 3 is executed $n - 4$ times. So the number of foos is $1 + n - 4 = n - 3 = \Theta(n)$ as $c_1 n \leq n - 3 \leq c_2 n$ for all $n \geq n_0$ when choosing, say, $n_0 = 6, c_1 = 1/2, c_2 = 1$.

EXAMPLE ALGORITHM

```

1: foo
2: for i = 1 to n - 4 do
3:     foo

```

ALGORITHM A

```

1: foo
2: for i = 1 to n do
3:     for j = 1 to n - 2 do
4:         foo
5:         foo
6:         foo

```

ALGORITHM B

```

1: foo
2: for i = 1 to n do
3:     foo
4: for i = 1 to n/2 do
5:     foo
6:     foo

```

ALGORITHM C

```

1: foo
2: for i = 1 to n do
3:     for j = 1 to i do
4:         foo
5:         foo
6: foo

```

Solutions: For Algorithm A, line 1 is executed once and lines 4–6 are executed $n \cdot (n - 2)$ times each. The total is $1 + 3n(n - 2) = 3n^2 - 6n + 1 = \Theta(n^2)$ as $c_1 n^2 \leq 3n^2 - 6n + 1 \leq c_2 n^2$ for all $n \geq 3$ (thus $n_0 = 3$) when choosing, say, $c_1 = 1/2$ and $c_2 = 3$.

(If you have chosen other constants $0 < c_1 < 3$ and $c_2 \geq 3$ then your answer is correct so long as your n_0 is chosen large enough.)

For Algorithm B, line 1 is executed once, line 3 is executed n times, and lines 5 and 6 are executed $n/2$ times each. Hence the number of foos is $1 + n + 2n/2 = 2n + 1 = \Theta(n)$ as $c_1 n \leq 2n + 1 \leq c_2 n$ for all $n \geq 1$ (thus $n_0 = 1$) when choosing, say, $c_1 = 2$ and $c_2 = 3$.

(If you have chosen other constants $0 < c_1 \leq 2$ and $c_2 > 2$ then your answer is correct so long as n_0 is chosen large enough.)

For Algorithm C, lines 1 and 6 are executed once. The inner for loop goes from 1 to i , hence line 4 is executed $\sum_{i=1}^n i = n(n+1)/2$ times. Line 5 is executed n times. The total is $2 + n(n+1)/2 + n = n^2/2 + 3n/2 + 2 = \Theta(n^2)$ as $c_1 n^2 \leq n^2/2 + 3n/2 + 2 \leq c_2 n^2$ for $n \geq 1$ when choosing, say, $c_1 = 1/2$ and $c_2 = 4$.

(If you have chosen other constants $0 < c_1 \leq 1/2$ and $c_2 > 1/2$ then your answer is correct so long as n_0 is chosen large enough.)

Question 2.4 (0.3 marks)

Recall from Lecture 2 that a statement like $2n^2 + \Theta(n) = \Theta(n^2)$ is true if *no matter how the anonymous functions are chosen on the left of the equal sign, there is a way to choose the anonymous functions on the right of the equal sign to make the equation valid*. You might want to think of the $\Theta(n)$ on the left-hand side being a placeholder for some (anonymous) function that grows as fast as n .

For each of the following statements, state whether it is true or false. Justify your answers.

1. $O(\sqrt{n}) = O(n)$

Solution: true.

This equation can be read as “some (anonymous) function that grows at most as fast as \sqrt{n} grows at most as fast as n ”. Whatever this anonymous function is, if it grows at most as fast as \sqrt{n} , it also grows at most as fast as n . Hence the statement is true. We can also express this in set notation as $O(\sqrt{n}) \subseteq O(n)$.

2. $n + o(n^2) = \omega(n)$

Solution: false.

The statement can be read as “the sum of n plus some (anonymous) function that grows slower than n^2 grows faster than n ”. The statement needs to hold for *all* anonymous functions $o(n^2)$, that is, all functions that grow slower than n^2 . This includes, say, n , which obviously grows slower than n^2 and hence is included in the set $o(n^2)$. Then the left-hand side would be $n + n = 2n$, which is not in $\omega(n)$. So the statement is false as it does not hold for all functions in $o(n^2)$. In set notation, $n + o(n^2) \not\subseteq \omega(n)$.

3. $3n \log n + O(n) = \Theta(n \log n)$

Solution: true.

The left-hand side can be read as “ $3n \log n$ plus some (anonymous) function that grows at most as fast as n ” and then the statement asserts that this sum grows as fast as $n \log n$. The term $3n \log n$ dominates the left-hand side as every function in $O(n)$ grows more slowly than $3n \log n$, hence $O(n)$ is just a small-order term compared to $3n \log n$. The left-hand side grows asymptotically like $n \log n$, hence it is in $\Theta(n \log n)$. In set notation, we have $3n \log n + O(n) \subseteq \Theta(n \log n)$.

Also, explain why the statement “The running time of Algorithm A is at least $O(n^2)$ ” is meaningless.

Solution: “ $O(n^2)$ ” is read as “at most cn^2 ” for a constant $c > 0$ and all $n \geq n_0$. So the statement spells out as “The running time of Algorithm A is at least at most cn^2 ”, which is obviously pointless.

Question 2.5 (0.3 marks)

The following algorithm computes the product C of two $n \times n$ matrices A and B , where $A[i, j]$ corresponds to the element in the i -th row and the j -th column.

MATRIX-MULTIPLY(A, B)

```
1: for  $i = 1$  to  $n$  do
2:   for  $j = 1$  to  $n$  do
3:      $C[i, j] := 0$ 
4:     for  $k = 1$  to  $n$  do
5:        $C[i, j] := C[i, j] + A[i, k] \cdot B[k, j]$ 
6: return  $C$ 
```

Give the running time of the algorithm (number of operations in a RAM machine) in Θ -notation. Justify your answer. Feel free to use the rules on calculating with Θ -notation from the lecture.

Solution: The first line is executed $\Theta(n)$ times, lines 2 to 3 are each executed $\Theta(n^2)$ times and cost $\Theta(1)$. The inner for loop is executed $\Theta(n^3)$ times, and the time for one execution of lines 4 to 5 is $\Theta(1)$. The time for the return statement is $\Theta(1)$. So the total time is

$$\Theta(n) + \Theta(n^2) \cdot \Theta(1) + \Theta(n^3) \cdot \Theta(1) + \Theta(1) = \Theta(n^3).$$

Question 2.6 (marks 0.75)

BUBBLESORT is a popular, but inefficient, sorting algorithm. It works by repeatedly swapping adjacent elements that are out of order. The effect is that small elements “bubble” to the left-hand side of the array, accumulating to form a growing sorted subarray. (You might want to work out your own example to understand this better.)

BUBBLE-SORT(A)

```
1: for  $i = 1$  to  $A.\text{length} - 1$  do
2:   for  $j = A.\text{length}$  downto  $i + 1$  do
3:     if  $A[j] < A[j - 1]$  then
4:       exchange  $A[j]$  with  $A[j - 1]$ 
```

Prove the correctness of BUBBLESORT and analyse its running time as follows. Try to keep your answers brief.

1. The inner loop “bubbles” a small element to the left-hand side of the array. State a loop invariant for the inner loop that captures this effect and prove that this loop invariant holds, addressing the three properties initialisation, maintenance, and termination.

Solution: the loop invariant is: *at the start of the inner loop, $A[j]$ contains a smallest element out of $A[j], \dots, A[n]$.*

Initialisation: $A[n]$ is the smallest element out of $A[n]$.

Maintenance: if $A[j - 1] \leq A[j]$, then $A[j - 1]$ is a smallest element out of $A[j - 1], \dots, A[n]$. Decreasing j establishes the loop invariant. Otherwise, after swapping $A[j]$ and $A[j - 1]$ we have $A[j - 1] < A[j]$ and continue as in the previous case.

Termination: at the end of the inner for loop, $j = i$ and $A[i]$ is the smallest element out of $A[i], \dots, A[n]$.

2. Using the termination condition of the loop invariant for the inner loop, state and prove a loop invariant for the outer loop in the same way as in part 1. that allows you to conclude that at the end of the algorithm the array is sorted.

Solution: the loop invariant is: *at the start of the outer loop, the subarray $A[1 \dots i - 1]$ contains the $i - 1$ smallest elements in sorted order.*

Initialisation: for $i = 1$, $A[1 \dots 0]$ is empty and contains the 0 smallest elements in sorted order.

Maintenance: after the end of the inner for loop, $A[i]$ contains the smallest element out of $A[i], \dots, A[n]$. Then $A[1 \dots i]$ contains the i smallest elements in sorted order. Incrementing i establishes the loop invariant.

Termination: at the end of the inner for loop, $i = n$ and $A[1 \dots n - 1]$ contains the $n - 1$ smallest elements in sorted order. This implies that $A[n]$ is no smaller and the whole array is sorted.

3. State the runtime of BUBBLESORT in asymptotic notation. Justify your answer. One iteration of the inner loop takes time $\Theta(1)$. The inner loop is executed $\sum_{i=1}^{n-1} (n - i)$ times. This is $\sum_{i=1}^{n-1} i = \frac{n(n-1)}{2} = \Theta(n^2)$. Hence the total time is $\Theta(1) \cdot \Theta(n^2) = \Theta(n^2)$.

Programming Question 2.7 (0.1 marks)

Implement MATRIX-MULTIPLY(A,B) and BUBBLESORT on the new Judge system.