

알고리즘

개요

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학습목표

- 알고리즘 문제 해결 전략 익히기
 - 복잡도
 - 샘플 문제들
 - 완전탐색(Exhaustive search or brute-force)
 - 대안
 - 정렬
 - 자료구조
 - (재귀)

코딩 테스트 문제 해결의 이해

- 정확한 입출력
- 시간제한
 - 예) 복잡도 이해: $O(n)$, $O(n \log n)$, $O(n^2)$, $O(n^3)$, $O(2^n)$, $O(n!)$
 - 컴퓨터의 CPU 클럭 1 GHz
 - 대략 1초에 1,000,000,000번 명령 수행 가능
 - $O(n)$ 이면 입력 크기 최대 1,000,000,000
 - $O(n^2)$ 이면 입력 크기 최대 30,000
 - $O(n^3)$ 이면 입력 크기 최대 1,000

| Notation | Name | Example |
|--|---|---|
| $O(1)$ | constant | Determining if a binary number is even or odd; Calculating $(-1)^n$; Using a constant-size lookup table |
| $O(\log \log n)$ | double logarithmic | Number of comparisons spent finding an item using interpolation search in a sorted array of uniformly distributed values |
| $O(\log n)$ | logarithmic | Finding an item in a sorted array with a binary search or a balanced search tree as well as all operations in a Binomial heap |
| $O((\log n)^c)$ $c > 1$ | polylogarithmic | Matrix chain ordering can be solved in polylogarithmic time on a parallel random-access machine . |
| $O(n^c)$ $0 < c < 1$ | fractional power | Searching in a k-d tree |
| $O(n)$ | linear | Finding an item in an unsorted list or in an unsorted array; adding two n -bit integers by ripple carry |
| $O(n \log^* n)$ | $n \log$ -star n | Performing triangulation of a simple polygon using Seidel's algorithm , or the union-find algorithm . Note that $\log^*(n) = \begin{cases} 0, & \text{if } n \leq 1 \\ 1 + \log^*(\log n), & \text{if } n > 1 \end{cases}$ |
| $O(n \log n) = O(\log n!)$ | linearithmic, loglinear, quasilinear, or " $n \log n$ " | Performing a fast Fourier transform ; Fastest possible comparison sort ; heapsort and merge sort |
| $O(n^2)$ | quadratic | Multiplying two n -digit numbers by a simple algorithm; simple sorting algorithms, such as bubble sort , selection sort and insertion sort ; (worst case) bound on some usually faster sorting algorithms such as quicksort , Shellsort , and tree sort |
| $O(n^c)$ | polynomial or algebraic | Tree-adjointing grammar parsing; maximum matching for bipartite graphs ; finding the determinant with LU decomposition |
| $L_n[\alpha, c] = e^{(c+o(1))(\ln n)^\alpha (\ln \ln n)^{1-\alpha}}$ $0 < \alpha < 1$ | L-notation or sub-exponential | Factoring a number using the quadratic sieve or number field sieve |
| $O(c^n)$ $c > 1$ | exponential | Finding the (exact) solution to the travelling salesman problem using dynamic programming ; determining if two logical statements are equivalent using brute-force search |
| $O(n!)$ | factorial | Solving the travelling salesman problem via brute-force search; generating all unrestricted permutations of a poset ; finding the determinant with Laplace expansion ; enumerating all partitions of a set |

Common Data Structure Operations

| Data Structure | Time Complexity | | | | | | | | Space Complexity |
|---------------------------|-------------------|-------------------|-------------------|-------------------|-------------------|-------------------|-------------------|-------------------|---------------------|
| | Average | | | | Worst | | | | Worst |
| | Access | Search | Insertion | Deletion | Access | Search | Insertion | Deletion | |
| <u>Array</u> | $\theta(1)$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ | $\theta(1)$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ |
| <u>Stack</u> | $\theta(n)$ | $\theta(n)$ | $\theta(1)$ | $\theta(1)$ | $\theta(n)$ | $\theta(n)$ | $\theta(1)$ | $\theta(1)$ | $\theta(n)$ |
| <u>Queue</u> | $\theta(n)$ | $\theta(n)$ | $\theta(1)$ | $\theta(1)$ | $\theta(n)$ | $\theta(n)$ | $\theta(1)$ | $\theta(1)$ | $\theta(n)$ |
| <u>Singly-Linked List</u> | $\theta(n)$ | $\theta(n)$ | $\theta(1)$ | $\theta(1)$ | $\theta(n)$ | $\theta(n)$ | $\theta(1)$ | $\theta(1)$ | $\theta(n)$ |
| <u>Doubly-Linked List</u> | $\theta(n)$ | $\theta(n)$ | $\theta(1)$ | $\theta(1)$ | $\theta(n)$ | $\theta(n)$ | $\theta(1)$ | $\theta(1)$ | $\theta(n)$ |
| <u>Skip List</u> | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ | $\theta(n \log(n))$ |
| <u>Hash Table</u> | N/A | $\theta(1)$ | $\theta(1)$ | $\theta(1)$ | N/A | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ |
| <u>Binary Search Tree</u> | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ |
| <u>Cartesian Tree</u> | N/A | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | N/A | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ |
| <u>B-Tree</u> | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(n)$ |
| <u>Red-Black Tree</u> | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(n)$ |
| <u>Splay Tree</u> | N/A | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | N/A | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(n)$ |
| <u>AVL Tree</u> | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(n)$ |
| <u>KD Tree</u> | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(\log(n))$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ | $\theta(n)$ |

Array Sorting Algorithms

| Algorithm | Time Complexity | | | Space Complexity |
|-----------------------|---------------------|------------------------|-------------------|------------------|
| | Best | Average | Worst | Worst |
| <u>Quicksort</u> | $\Omega(n \log(n))$ | $\Theta(n \log(n))$ | $O(n^2)$ | $O(\log(n))$ |
| <u>Mergesort</u> | $\Omega(n \log(n))$ | $\Theta(n \log(n))$ | $O(n \log(n))$ | $O(n)$ |
| <u>Timsort</u> | $\Omega(n)$ | $\Theta(n \log(n))$ | $O(n \log(n))$ | $O(n)$ |
| <u>Heapsort</u> | $\Omega(n \log(n))$ | $\Theta(n \log(n))$ | $O(n \log(n))$ | $O(1)$ |
| <u>Bubble Sort</u> | $\Omega(n)$ | $\Theta(n^2)$ | $O(n^2)$ | $O(1)$ |
| <u>Insertion Sort</u> | $\Omega(n)$ | $\Theta(n^2)$ | $O(n^2)$ | $O(1)$ |
| <u>Selection Sort</u> | $\Omega(n^2)$ | $\Theta(n^2)$ | $O(n^2)$ | $O(1)$ |
| <u>Tree Sort</u> | $\Omega(n \log(n))$ | $\Theta(n \log(n))$ | $O(n^2)$ | $O(n)$ |
| <u>Shell Sort</u> | $\Omega(n \log(n))$ | $\Theta(n(\log(n))^2)$ | $O(n(\log(n))^2)$ | $O(1)$ |
| <u>Bucket Sort</u> | $\Omega(n+k)$ | $\Theta(n+k)$ | $O(n^2)$ | $O(n)$ |
| <u>Radix Sort</u> | $\Omega(nk)$ | $\Theta(nk)$ | $O(nk)$ | $O(n+k)$ |
| <u>Counting Sort</u> | $\Omega(n+k)$ | $\Theta(n+k)$ | $O(n+k)$ | $O(k)$ |
| <u>Cubesort</u> | $\Omega(n)$ | $\Theta(n \log(n))$ | $O(n \log(n))$ | $O(n)$ |

입력 크기와 실전 복잡도

Table 3.1 Estimating time complexity from input size

| Input size | Expected time complexity |
|---------------|--------------------------|
| $n \leq 10$ | $O(n!)$ |
| $n \leq 20$ | $O(2^n)$ |
| $n \leq 500$ | $O(n^3)$ |
| $n \leq 5000$ | $O(n^2)$ |
| $n \leq 10^6$ | $O(n \log n)$ or $O(n)$ |
| n is large | $O(1)$ or $O(\log n)$ |

정리

- 기본적인 복잡도
- 문제해결방법의 이해
 - 완전탐색