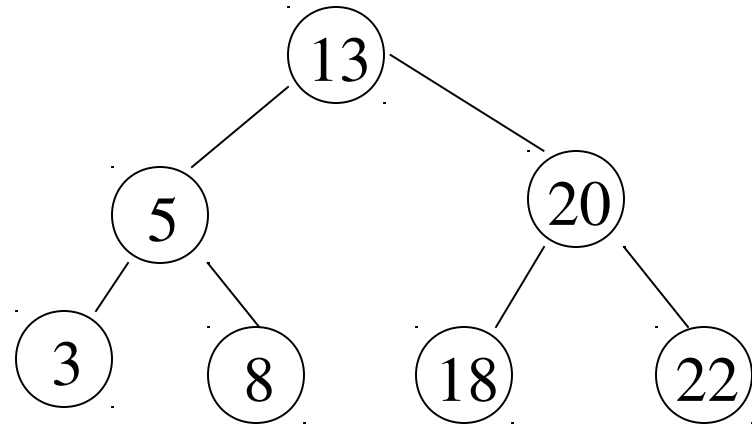
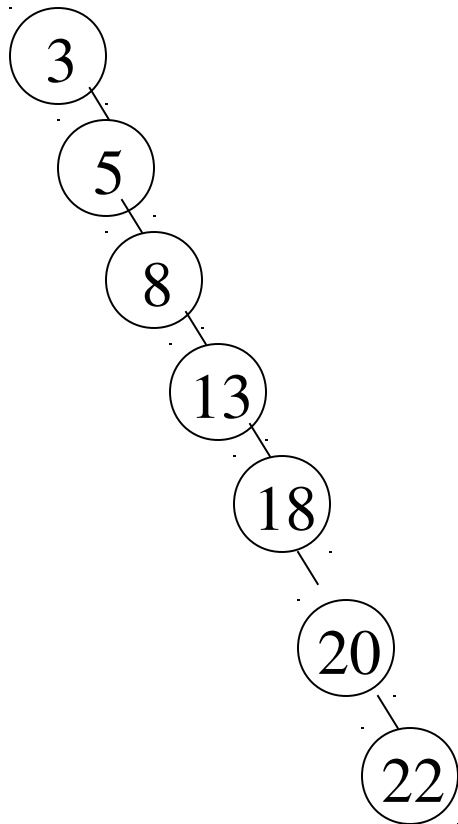


Trees 4: AVL Trees

- Section 4.4

Motivation

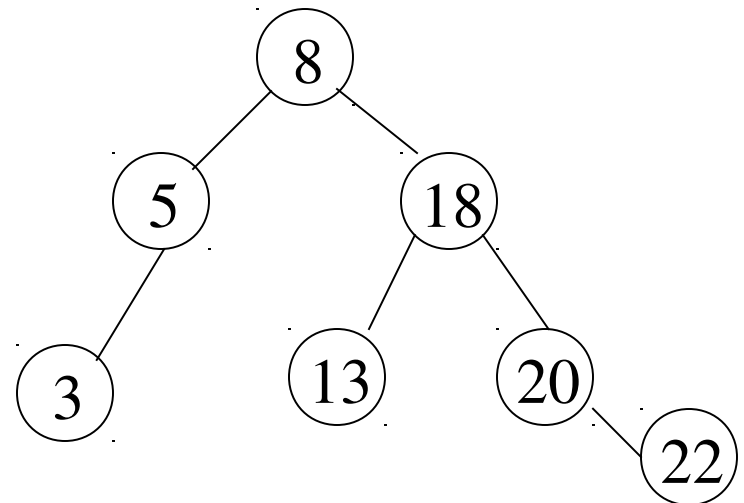
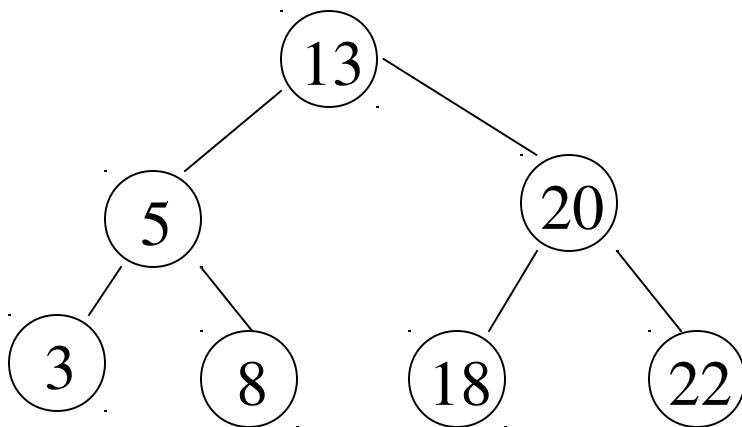
- When building a binary search tree, what type of trees would we like? Example: 3, 5, 8, 20, 18, 13, 22



Motivation



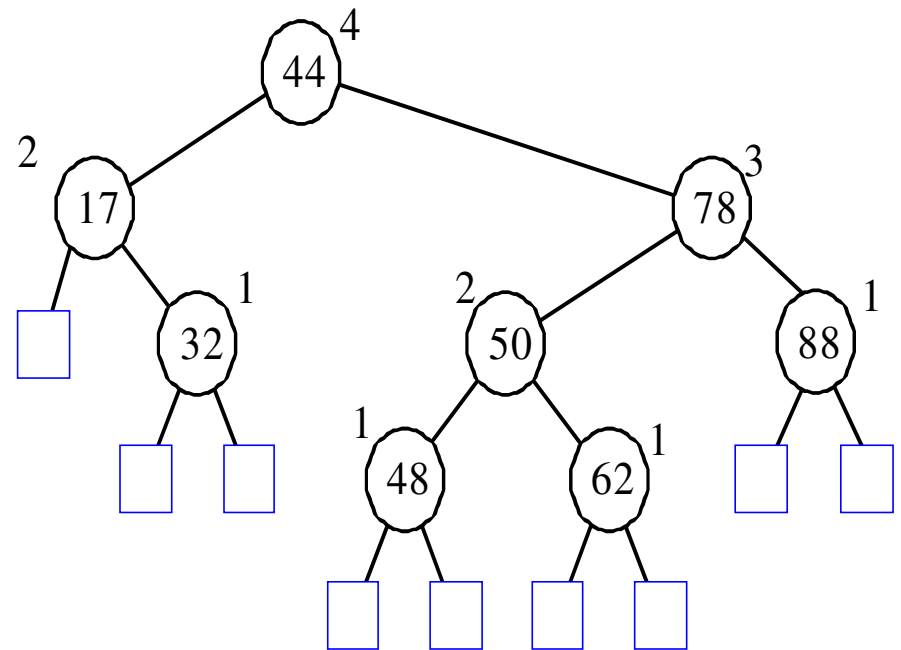
- Complete binary tree is hard to build when we allow dynamic insert and remove.
 - We want a tree that has the following properties
 - Tree height = $O(\log(N))$
 - allows dynamic insert and remove with $O(\log(N))$ time complexity.
 - The AVL tree is one of this kind of trees.



AVL (Adelson-Velskii and Landis) Trees



- An AVL Tree is a **binary search tree** such that for every internal node v of T , the *heights of the children of v can differ by at most 1*.



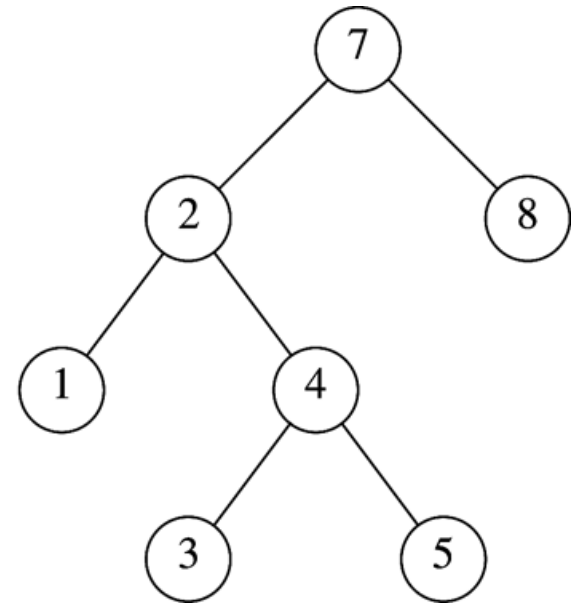
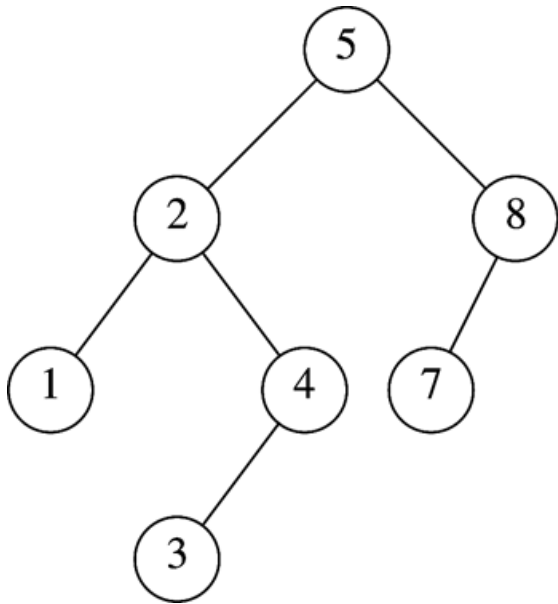
An example of an AVL tree where the heights are shown next to the nodes:

AVL (Adelson-Velskii and Landis) Trees

- AVL tree is a binary search tree with balance condition
 - To ensure depth of the tree is $O(\log(N))$
 - And consequently, search/insert/remove complexity bound $O(\log(N))$
- Balance condition
 - For every node in the tree, height of left and right subtree can differ by at most 1



Which is an AVL Tree?





Height of an AVL tree

- Theorem: The **height** of an AVL tree storing n keys is $O(\log n)$.
- **Proof:**
 - Let us bound $n(h)$, the minimum number of internal nodes of an AVL tree of height h .
 - We easily see that $n(0) = 1$ and $n(1) = 2$
 - For $h > 2$, an AVL tree of height h contains the root node, one AVL subtree of height $h-1$ and another of height $h-2$ (at worst).
 - That is, $n(h) \geq 1 + n(h-1) + n(h-2)$
 - Knowing $n(h-1) > n(h-2)$, we get $n(h) > 2n(h-2)$. So
$$n(h) > 2n(h-2), n(h) > 4n(h-4), n(h) > 8n(h-6), \dots \text{ (by induction),}$$
$$n(h) > 2^i n(h-2i)$$
 - Solving the base case we get: $n(h) > 2^{h/2-1}$
 - Taking logarithms: $h < 2\log n(h) + 2$
 - Since $n \geq n(h)$, $h < 2\log(n) + 2$ and the height of an AVL tree is $O(\log n)$

AVL Tree Insert and Remove

- Do binary search tree insert and remove
- The balance condition can be violated sometimes
 - Do something to fix it : **rotations**
 - After rotations, the balance of the whole tree is maintained





Balance Condition Violation

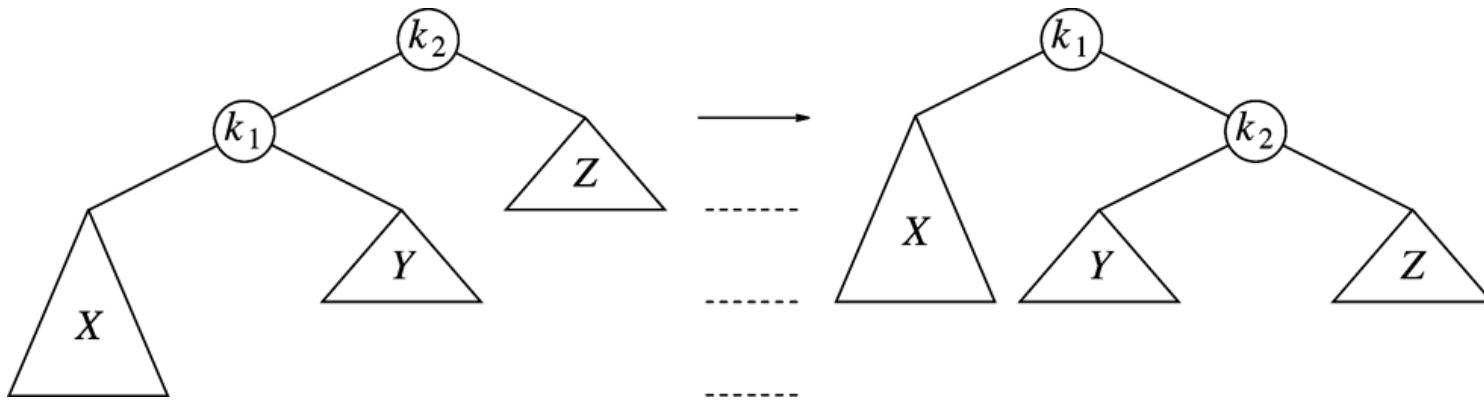
- If condition violated after a node insertion
 - Which nodes do we need to rotate?
 - Only nodes on path from insertion point to root may have their balance altered
- Rebalance the tree through rotation at the deepest node with balance violated
 - The entire tree will be rebalanced
- Violation cases at node k (deepest node)
 1. An insertion into left subtree of left child of k
 2. An insertion into right subtree of left child of k
 3. An insertion into left subtree of right child of k
 4. An insertion into right subtree of right child of k
 - Cases 1 and 4 equivalent
 - Single rotation to rebalance
 - Cases 2 and 3 equivalent
 - Double rotation to rebalance



AVL Trees Complexity

- **Overhead**
 - Extra space for maintaining height information at each node
 - Insertion and deletion become more complicated, but still $O(\log N)$
- **Advantage**
 - Worst case $O(\log(N))$ for insert, delete, and search

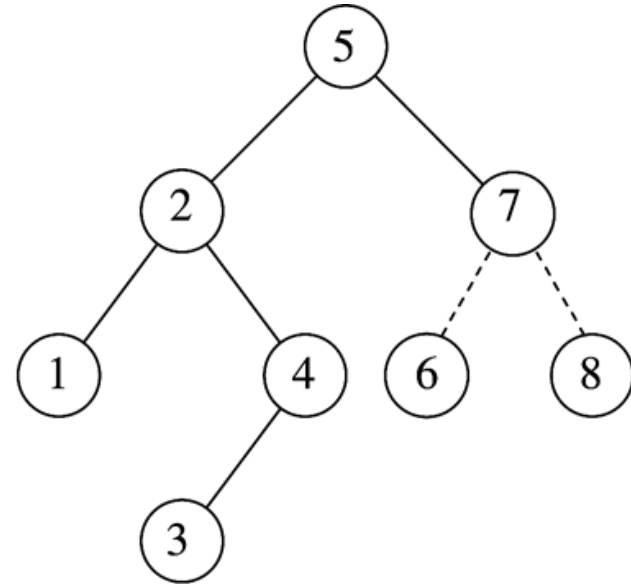
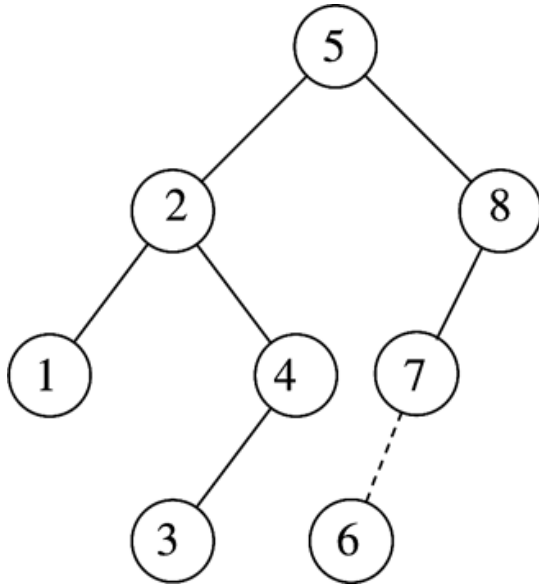
Single Rotation (Case 1)



- Replace node k_2 by node k_1
- Set node k_2 to be right child of node k_1
- Set subtree Y to be left child of node k_2
- Case 4 is similar



Example

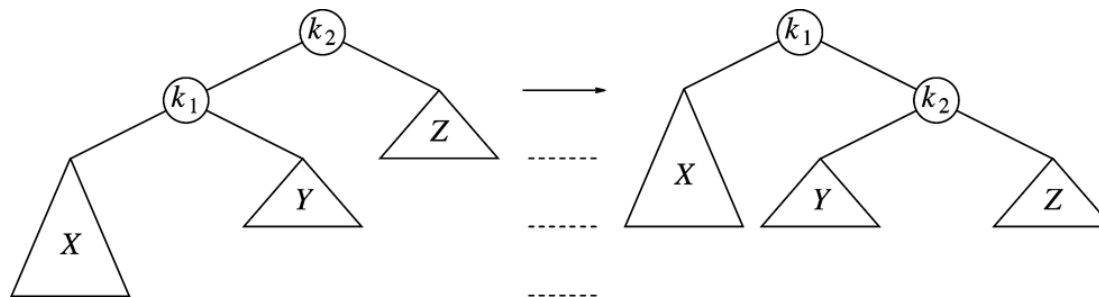


- After inserting 6
 - Balance condition at node 8 is violated



Single Rotation (Case 1)

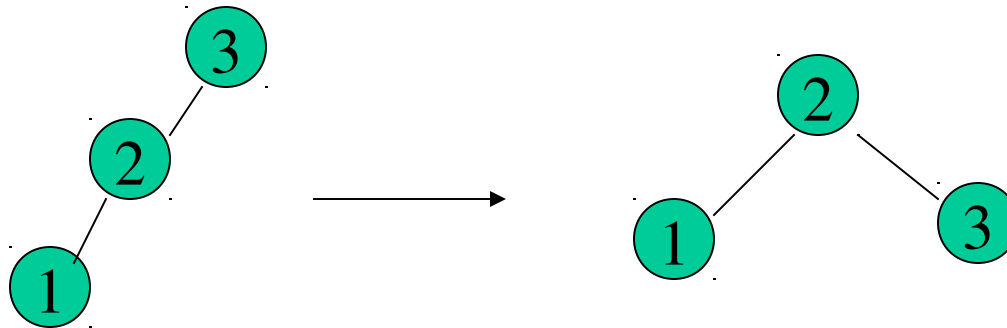
```
1  /**
2   * Rotate binary tree node with left child.
3   * For AVL trees, this is a single rotation for case 1.
4   * Update heights, then set new root.
5   */
6  void rotateWithLeftChild( AvlNode * & k2 )
7  {
8      AvlNode *k1 = k2->left;
9      k2->left = k1->right;
10     k1->right = k2;
11     k2->height = max( height( k2->left ), height( k2->right ) ) + 1;
12     k1->height = max( height( k1->left ), k2->height ) + 1;
13     k2 = k1;
14 }
```





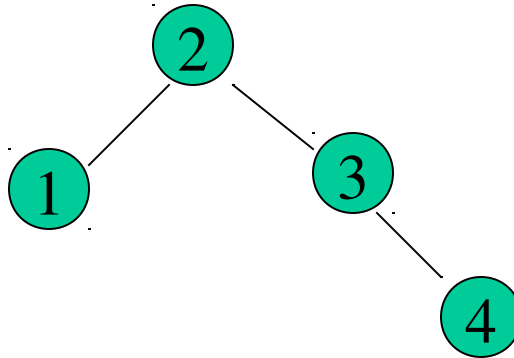
Example

- Inserting 3, 2, 1, and then 4 to 7 sequentially into empty AVL tree

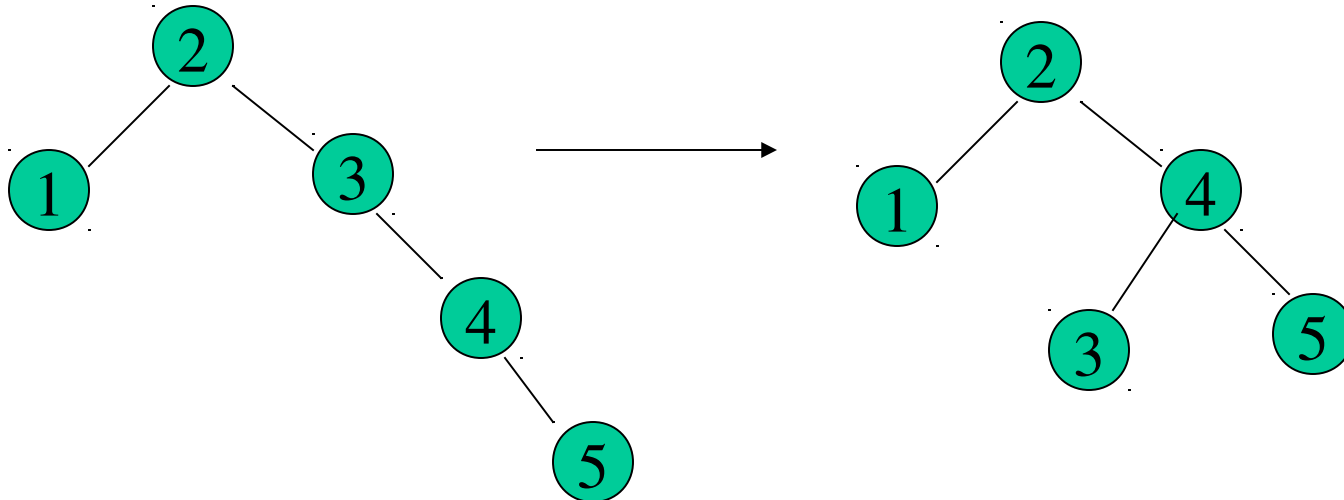


Example (Cont'd)

- Inserting 4

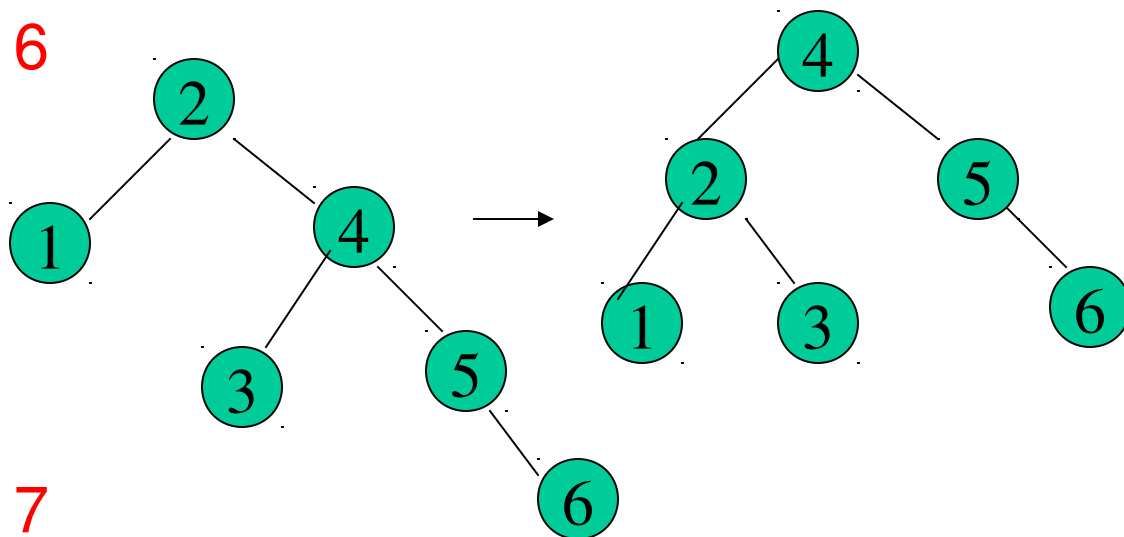


- Inserting 5

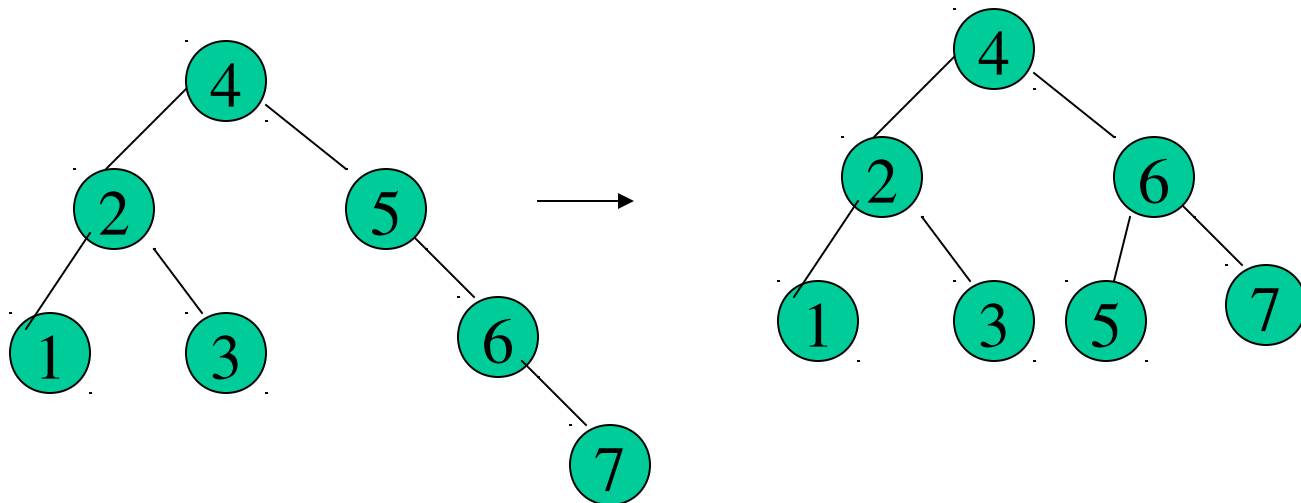


Example (Cont'd)

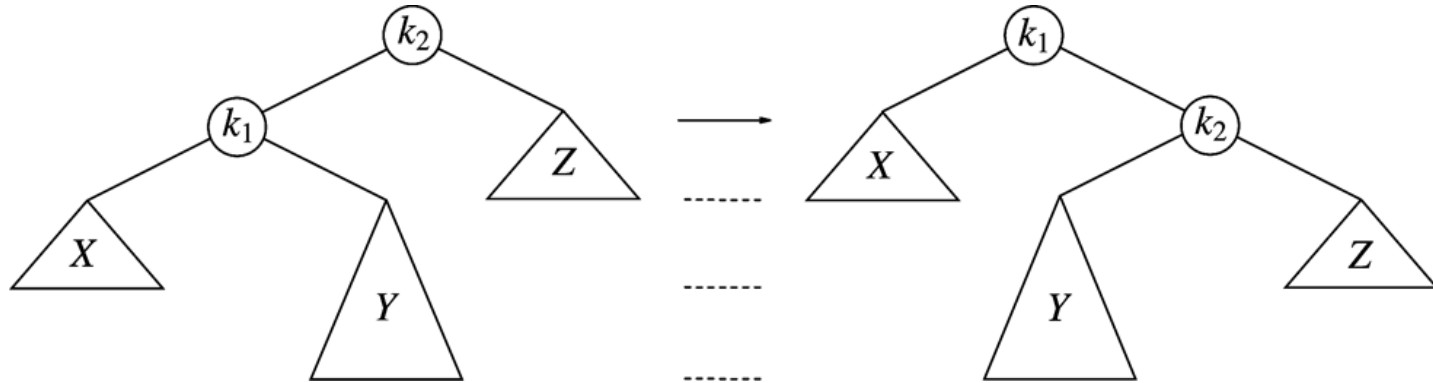
- Inserting 6



- Inserting 7

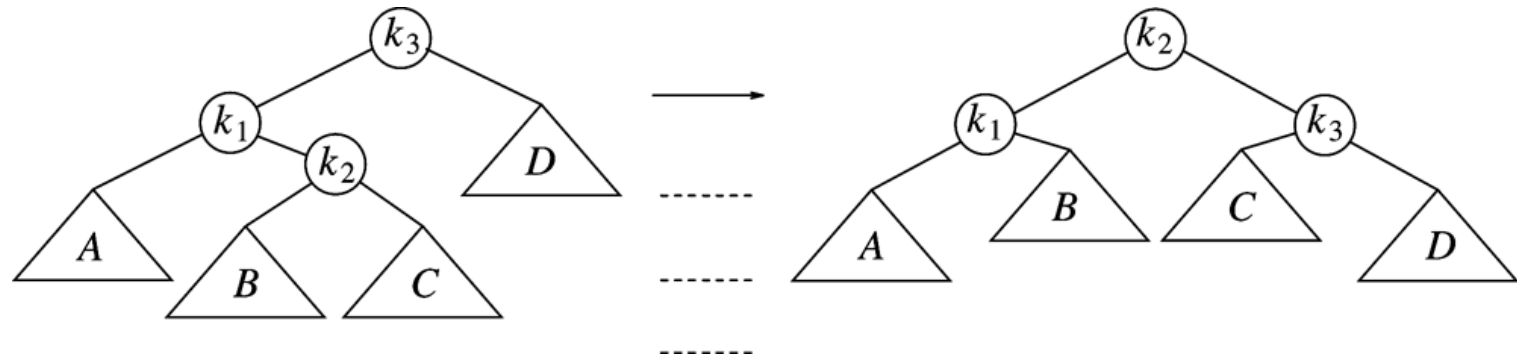


Single Rotation Will Not Work for the Other Case



- For case 2
- After single rotation, k_1 still not balanced
- Double rotations needed for case 2 and case 3

Double Rotation (Case 2)

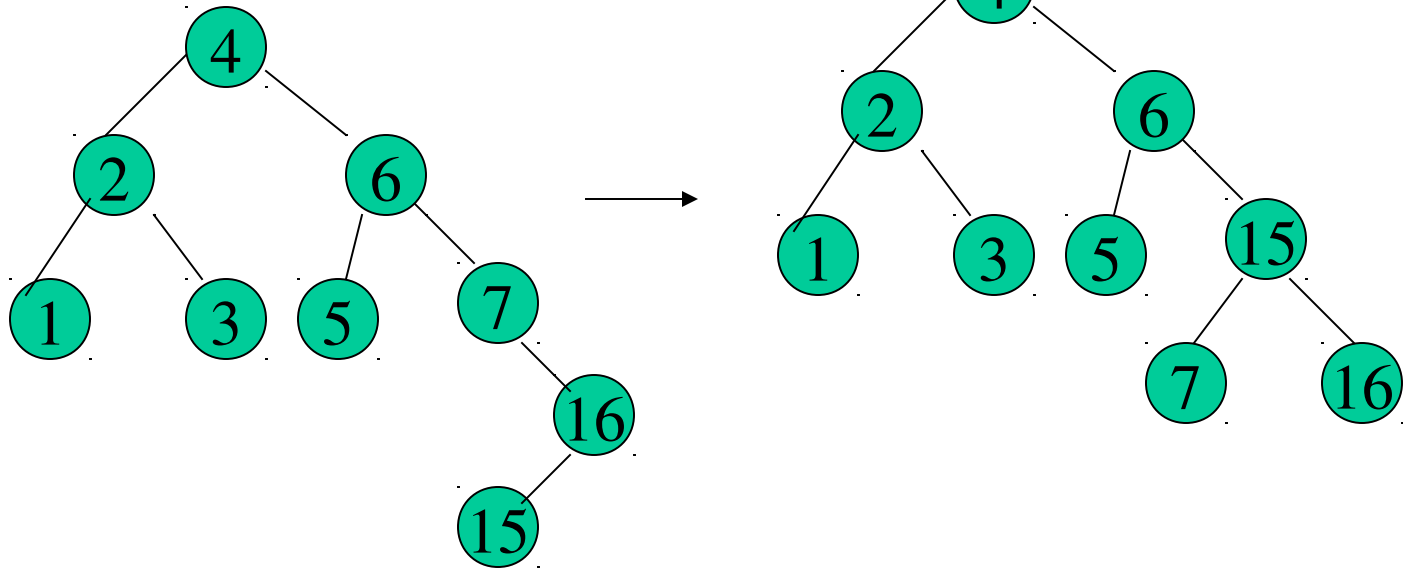


- Left-right double rotation to fix case 2
- First rotate between k_1 and k_2
- Then rotate between k_2 and k_3
- Case 3 is similar



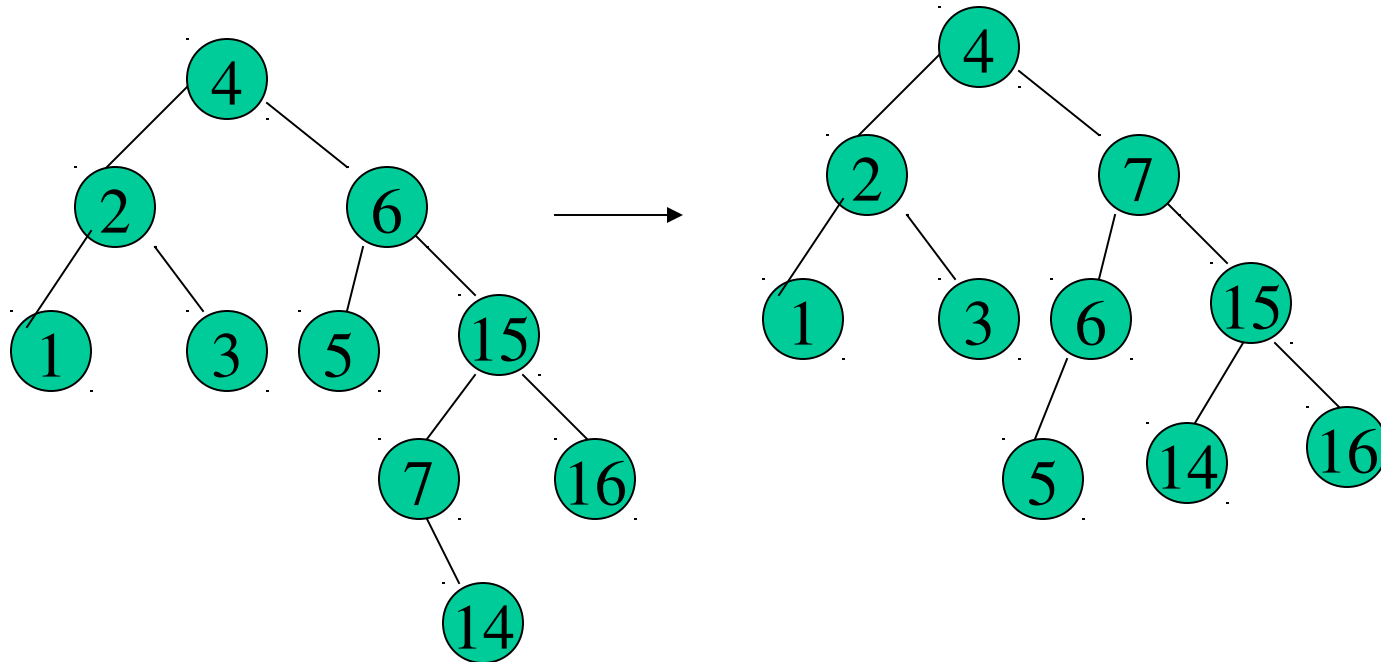
Example

- Continuing the previous example by inserting
 - 16 down to 10, and then 8 and 9
- Inserting 16 and 15



Example (Cont'd)

- Inserting 14

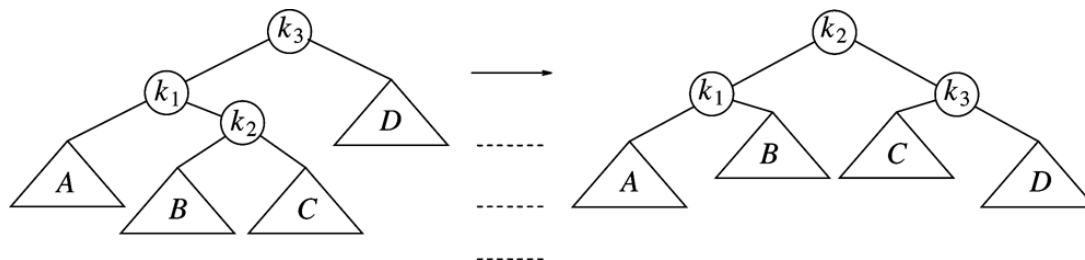


- Other cases as exercises



Double Rotation (Case 2)

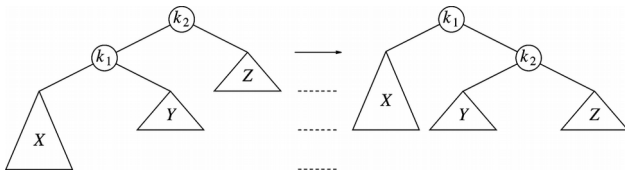
```
1  /**
2   * Double rotate binary tree node: first left child
3   * with its right child; then node k3 with new left child.
4   * For AVL trees, this is a double rotation for case 2.
5   * Update heights, then set new root.
6   */
7  void doubleWithLeftChild( AvlNode * & k3 )
8  {
9      rotateWithRightChild( k3->left );
10     rotateWithLeftChild( k3 );
11 }
```



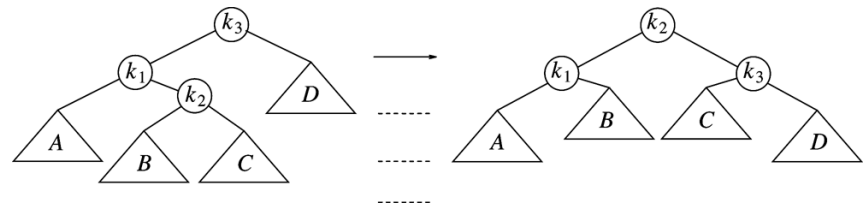
Summary

Violation cases at node k (deepest node)

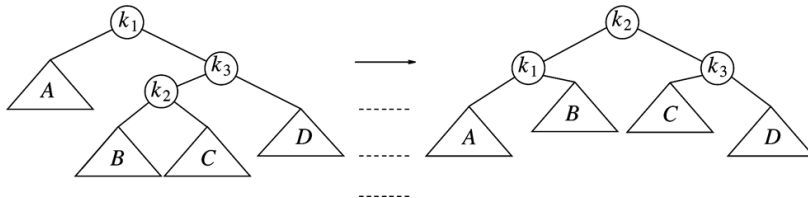
1. An insertion into left subtree of left child of k
2. An insertion into right subtree of left child of k
3. An insertion into left subtree of right child of k
4. An insertion into right subtree of right child of k



Case 1



Case 2



Case 3

Case 4?



Implementation of AVL Tree

```
1 struct AvlNode
2 {
3     Comparable element;
4     AvlNode *left;
5     AvlNode *right;
6     int height;
7
8     AvlNode( const Comparable & theElement, AvlNode *lt,
9             AvlNode *rt, int h = 0 )
10         : element( theElement ), left( lt ), right( rt ), height( h )
11 };
1
2 /**
3  * Return the height of node t or -1 if NULL.
4  */
5 int height( AvlNode *t ) const
6 {
7     return t == NULL ? -1 : t->height;
8 }
```

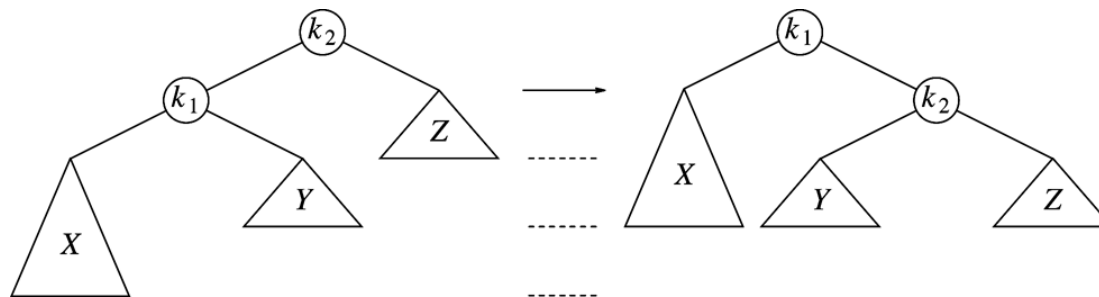


```
1      /**
2      * Internal method to insert into a subtree.
3      * x is the item to insert.
4      * t is the node that roots the subtree.
5      * Set the new root of the subtree.
6      */
7      void insert( const Comparable & x, AvlNode * & t )
8      {
9          if( t == NULL )
10             t = new AvlNode( x, NULL, NULL );
11         else if( x < t->element )
12         {
13             insert( x, t->left );
14             if( height( t->left ) - height( t->right ) == 2 )
15                 if( x < t->left->element )
16                     rotateWithLeftChild( t ); ← Case 1
17                 else
18                     doubleWithLeftChild( t ); ← Case 2
19         }
20         else if( t->element < x )
21         {
22             insert( x, t->right );
23             if( height( t->right ) - height( t->left ) == 2 )
24                 if( t->right->element < x )
25                     rotateWithRightChild( t ); ← Case 4
26                 else
27                     doubleWithRightChild( t ); ← Case 3
28         }
29         else
30             ; // Duplicate; do nothing
31         t->height = max( height( t->left ), height( t->right ) ) + 1;
32     }
```




Single Rotation (Case 1)

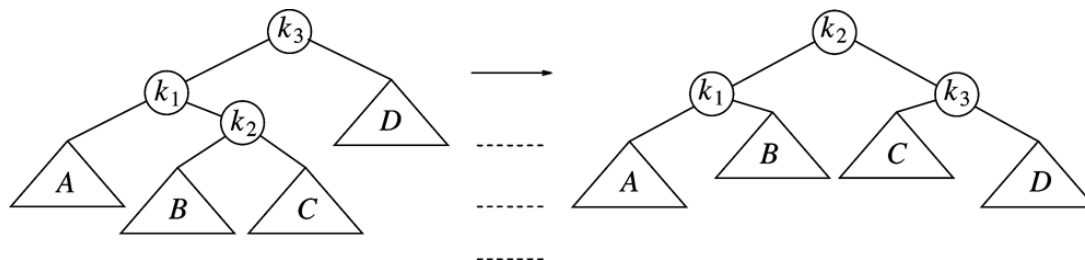
```
1  /**
2   * Rotate binary tree node with left child.
3   * For AVL trees, this is a single rotation for case 1.
4   * Update heights, then set new root.
5   */
6  void rotateWithLeftChild( AvlNode * & k2 )
7  {
8      AvlNode *k1 = k2->left;
9      k2->left = k1->right;
10     k1->right = k2;
11     k2->height = max( height( k2->left ), height( k2->right ) ) + 1;
12     k1->height = max( height( k1->left ), k2->height ) + 1;
13     k2 = k1;
14 }
```



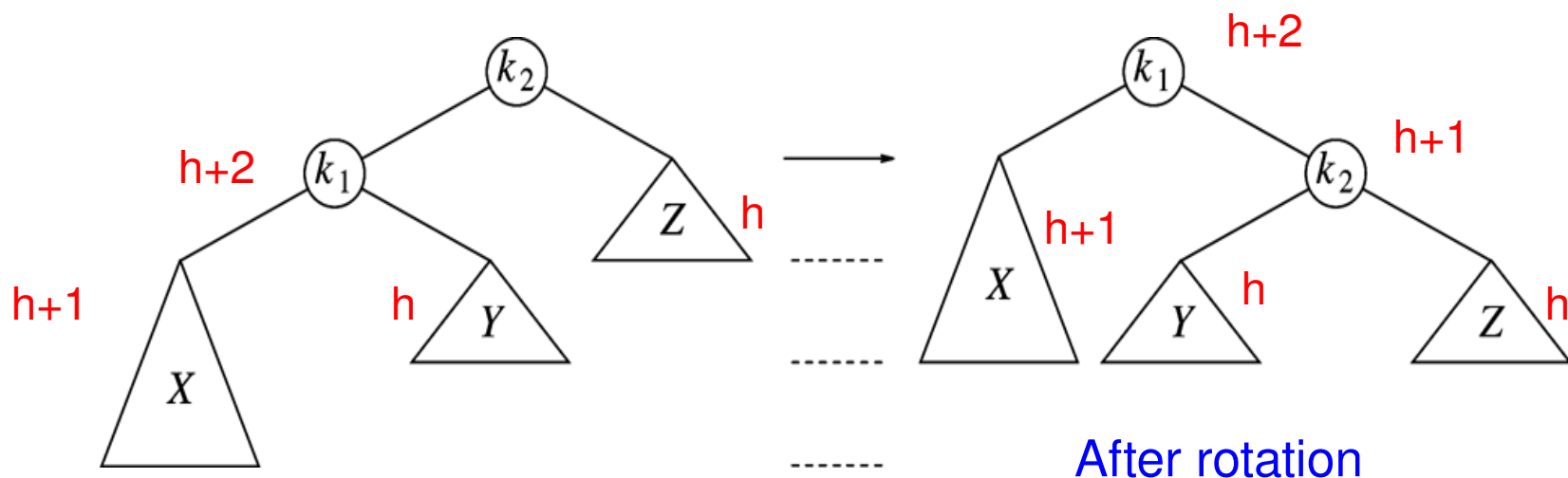
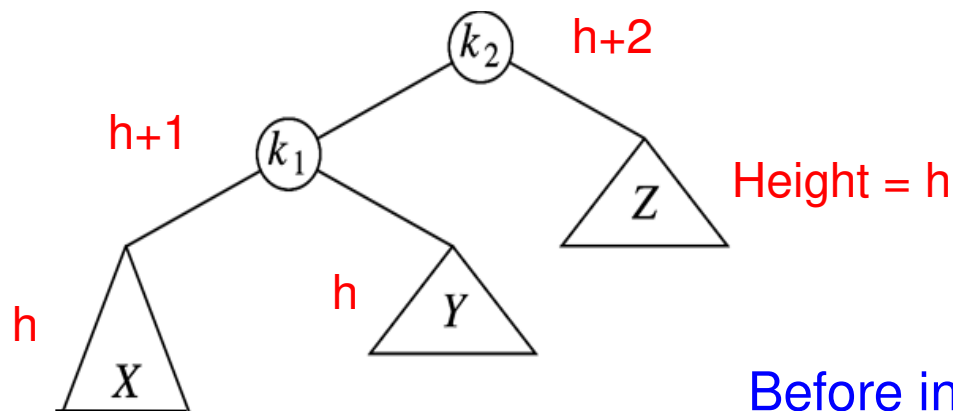


Double Rotation (Case 2)

```
1  /**
2   * Double rotate binary tree node: first left child
3   * with its right child; then node k3 with new left child.
4   * For AVL trees, this is a double rotation for case 2.
5   * Update heights, then set new root.
6   */
7  void doubleWithLeftChild( AvlNode * & k3 )
8  {
9      rotateWithRightChild( k3->left );
10     rotateWithLeftChild( k3 );
11 }
```



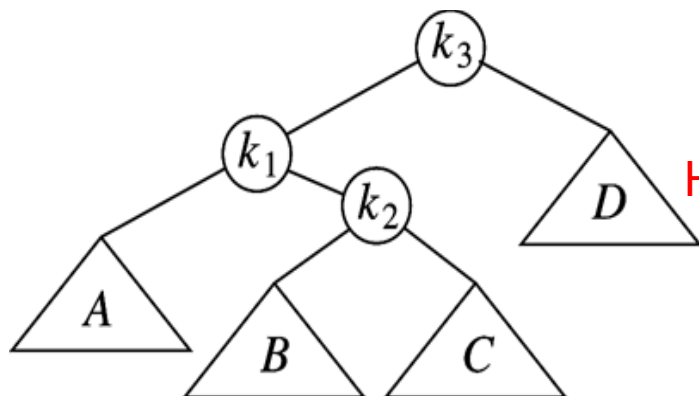
Review Insertion -- Case 1



After insert

After rotation

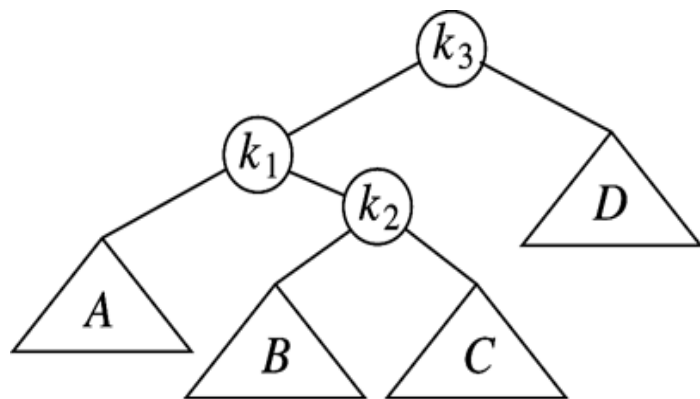
Review Insertion -- Case 2



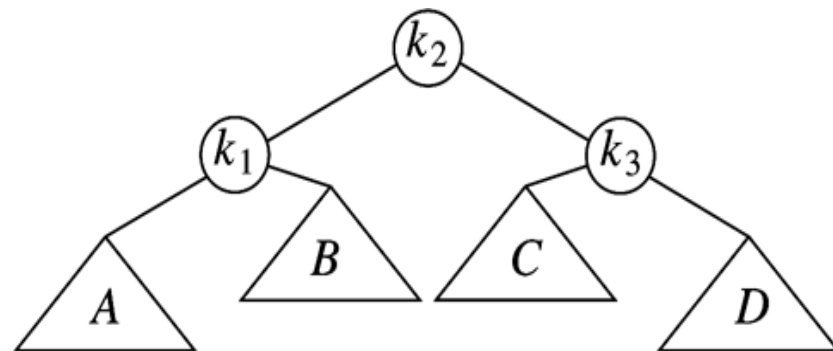
Height = h

Determine all heights

Before insert



After insert

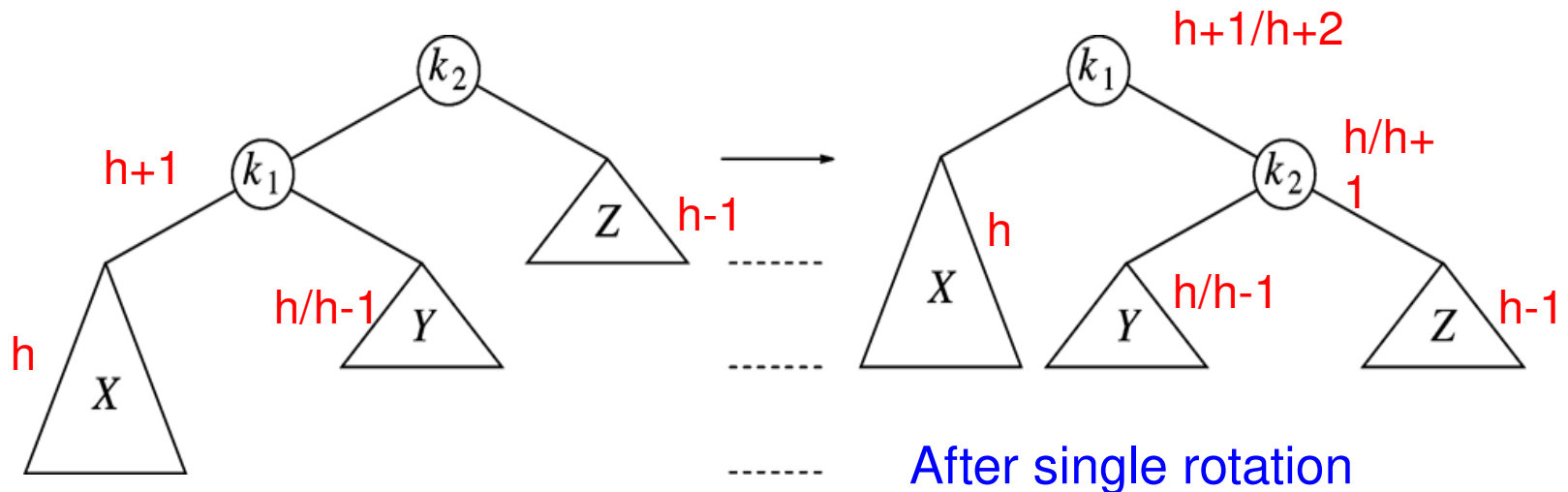
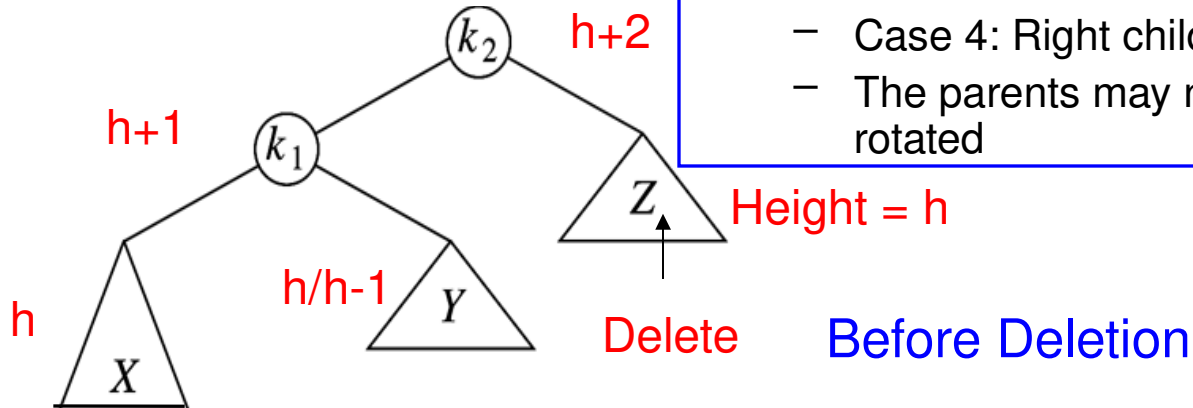


After double rotation



Delete -- Case 1

- Consider deepest unbalanced node
 - Case 1: Left child's left side is too high
 - Case 4: Right child's right side is too high
 - The parents may need to be recursively rotated

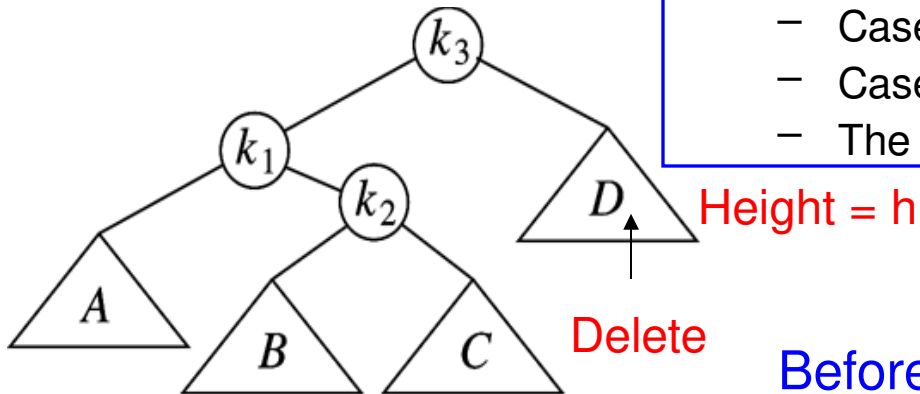


After delete

Delete -- Case 2

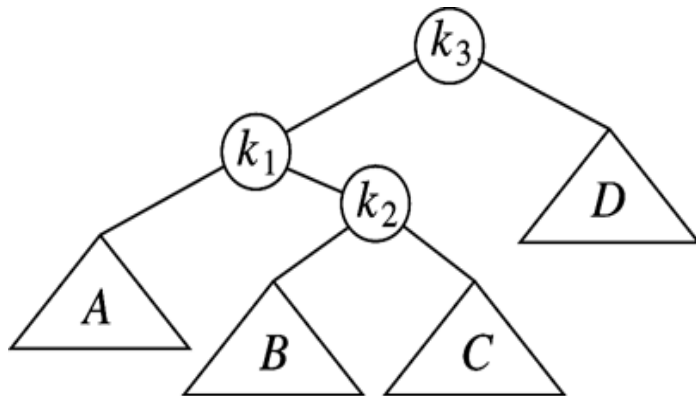


- Consider deepest unbalanced node
 - Case 2: Left child's right side is too high
 - Case 3: Right child's left side is too high
 - The parents may need to be recursively rotated

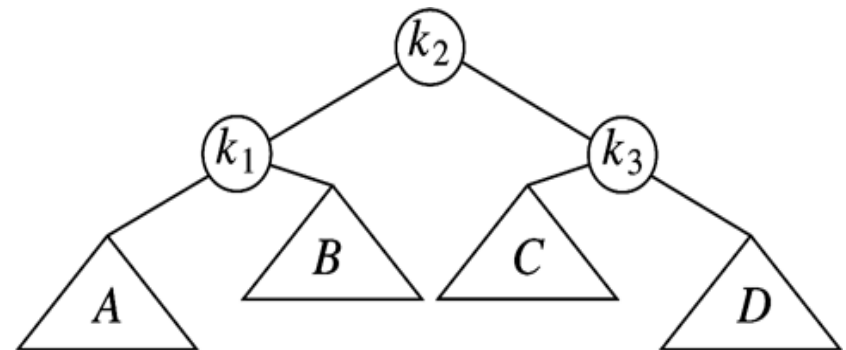
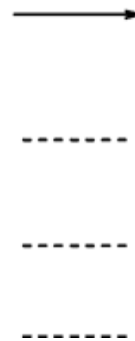


Determine all heights

Before Deletion



After Delete



After double rotation