CS 457, Data Structures and Algorithms I First Problem Set

September 26, 2018

Due on October 7. Collaboration is allowed; please mention your collaborators.

- 1. (16 pts) Find two functions f(n) and g(n) that satisfy the following relationship. If no such f and g exist, then try to explain why this is the case.
 - a) f(n) = o(g(n)) and $f(n) \notin \Theta(g(n))$ b) $f(n) = \Theta(g(n))$ and f(n) = o(g(n))c) $f(n) = \Theta(g(n))$ and $f(n) \notin O(g(n))$ d) $f(n) = \Omega(g(n))$ and $f(n) \notin O(g(n))$
- 2. (16 pts) For each of the following questions, briefly explain your answer.
 - a) If I prove that an algorithm takes $O(n^2)$ worst-case time, is it possible that it takes O(n) time on some inputs?
 - b) If I prove that an algorithm takes $O(n^2)$ worst-case time, is it possible that it takes O(n) time on all inputs?
 - c) If I prove that an algorithm takes $\Theta(n^2)$ worst-case time, is it possible that it takes O(n) time on some inputs?
 - d) If I prove that an algorithm takes $\Theta(n^2)$ worst-case time, is it possible that it takes O(n) time on all inputs?
- 3. (10 pts) Let f(n) and g(n) be any two functions with f(n) > 1 and g(n) > 1 for all n. Is it true that $15f(n)^2 + 23g(n)^2 = \Theta((f(n) + g(n))^2)$? Give a formal justification for your answer.
- 4. (15 pts) Is the following statement true or false, and why? (provide a formal argument either way)

If
$$f(n) = O(g(n))$$
 then $2^{f(n)} = O(2^{g(n)})$.

5. (20 pts) What value is returned by the following function? Express your answer as a function of n. From this, deduce the worst-case running time using Big Oh notation. For full credit, provide a lower bound as well, thus leading to a bound with $\Theta(\cdot)$ notation. Then, also deduce the worst-case running time if we change the while loop clause to $j \leq i$ instead of $j \leq n$.

```
1 function fool(n)

2 r := 0

3 for i := 1 to n do

4 j := 1

5 while j \le n do

6 r := r + 1

7 j := 2 \cdot j

8 return r
```

6. (14 pts) What value is returned by the following function? Express your answer as a function of n. From this, deduce the worst-case running time using Big Oh notation. For full credit, provide a lower bound as well, thus leading to a bound with $\Theta(\cdot)$ notation.

```
1 function foo2(n)

2 r := 0

3 j := n

4 while j \ge 1 do

5 for i := 1 to j do

6 r := r + 1

7 j := j/2

8 return r
```

- 7. (9 pts) Show that, if c is a positive real number, then $g(n) = 1 + c + c^2 + \cdots + c^n$ is
 - a) $\Theta(1)$ if c < 1
 - b) $\Theta(n)$ if c=1
 - c) $\Theta(c^n)$ if c > 1