# x86 Intel Assembly - Registers and The Stack

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  - Segment Registers
  - Special Registers

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- 3 The Stack
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### Objective

This tutorial/instruction slide-set is intended to provide a beginners look at registers and the stack as operated on by x86 instructions.

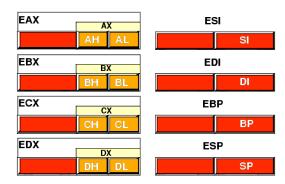
```
FunctionPrologue:

push ebp ; Comment

mov ebp, esp; Comment Two

sub esp, 8
```

# 32-bit Register Layout



# 32-bit Registers By Types

- General Registers
  - EAX
  - EBX
  - ECX
  - EDX
- 2 Segment Registers
  - CS
  - DS
  - ES
  - FS
  - GS
  - SS

- Index and Pointers
  - ESI
  - EDI
  - EBP
  - EIP
  - ESP
- Indicator
  - EFLAGS

### **General Registers**

Overview

Registers

General registers are used by most instructions in x86. These will be extremely common and can be broken down to 16 and 8 bit seaments.

```
ebx
              edx
       ecx
bx
              dx
 bh bl
       ch cl
              dh dl
```

The 'h' and 'l' suffixes on the 8 bit registers stands for the higher (h) and lower (I) bytes of the overall register.

### Register Value Example

In the example below, assume eax contained the value 0 (zero) before the execution of the code below - for now, sign is ignored

```
add eax, 0xFF00h; Add 65280 to eax; eax = 00000000 00000000 11111111 00000000 32; ax = 11111111 00000000 16; ah = 11111111 8; al = 000000000 8
```

For common reverse engineering the segment registers will not be of much interest. They are 16-bits in size and contain a 'segment selector.' A 'segment selector' is defined as a pointer to a place in memory where a segment exists.

Segment Registers

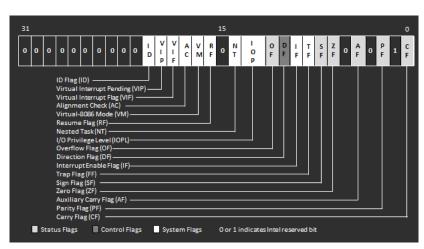
- CS Code segment points to where the instructions are stored and executed
- 2 DS, ES, FS, GS Points to the four data segments
- 3 SS Points to the stack segment, where the procedure stack is stored

### Pointers and Indexes

These registers are slightly misleading, the only register in this group that is classified separate from the General Registers is EIP. The other are used in convention for special purposes. EIP stores an offset to the next instruction to be executed. This register cannot be set by an application directly, it is set through implicit processes and functions.

### **EFLAGS** Register

The EFLAGS register is different in operation than the others. It is 32-bits in size and stores a variety of 1-bit values, commonly called Flags.



## Status Flags

The flags at bits 0, 2, 4, 6, 7, and 11 are called status flags. These are set after arithmetic to provide information about the 'status' of the computation. All flags are cleared if not set to 1 after a computation

- Bit 0 Carry Flag Set to 1 is the computation generates a carry of the most significant bit(MSB). Indicates an overflow in unsigned arithmetic.
- 2 Bit 2 Parity Flag Set if the results least significant byte contains an even number of one bits.
- 3 Bit 4 Aux. Carry Flag- Set if arithmetic generates a carry of bit 3 of the result Used in Binary Coded Decimal arithmetic.
- Bit 6 Zero Flag Set if the result is zero
- 5 Bit 7 Sign Flag Set equal to the MSB of the result (0 is positive, 1 is negative)
- Bit 11 Overflow Flag Set if the result is too large of a positive number or too small of a negative number. Indicates overflow in signed arithmetic.



These flags generally control operating-system and executive operations. They are not usually modified by applications. The DF is the direction flag and is used in string operations. Unset, the DF string instructions will auto-increment (low-address to high-address) and with DF set string functions will auto-decrement (high-address to low-address).

# Registers Overview

- 1 EAX Commonly used for I/O Port Access, arithmetic, interrupts
- EBX Commonly used as memory access base pointer
- ECX Commonly used as a loop counter or a shift counter
- EDX Very similar to EAX in common usage
- EDI Used for string, memory array copying. Far pointer addressing
- 6 ESI Used for string and memory array copying
- EBP Stores the stack base pointer
- 8 ESP Stores a pointer to the top of the stack
- EIP Holds offset to next instruction



Stack Overview

### Stack overview

#### What is a stack?

- Data Structure
- FILO First In, Last Out

### Why?

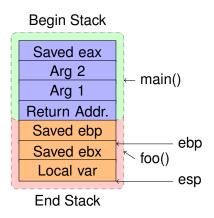
Computers use the stack to keep track of relevant data and addresses during execution. Understanding the stack will be essential in reverse engineering programs.



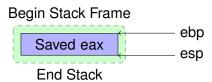
In other words, items are pushed onto the stack and then popped off of the top of the stack. The most recently added item will be the item that is popped.

Standards

### **CDECL Stack Arrangement**



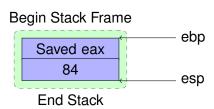
CDECL is called a caller clean-up convention. The calling function is responsible for pushing arguments onto the stack resizing the stack after the called function returns. This slide-set uses CDECL as it is the most common.



```
push eax
       push ebp
       mov ebp, esp
       push ebx
       sub esp,
       mov eax,
       add esp,
16
           ebp
```

Overview

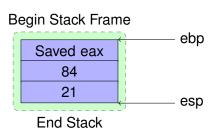
Registers



```
push eax
       add esp, 12
       push ebp
       mov ebp, esp
       push ebx
       sub esp,
       mov eax,
       add esp,
16
           ebp
```

Overview

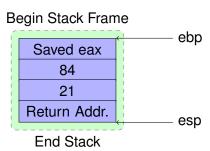
Registers



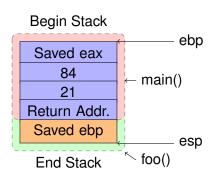
```
push eax
       add esp, 12
       push ebp
       mov ebp, esp
       push ebx
       sub esp,
       mov eax,
       add esp,
       pop ebx
16
       pop ebp
```

Overview

Registers



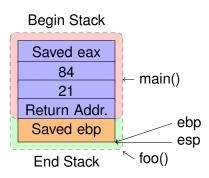
```
push eax
       push ebp
       mov ebp, esp
       push ebx
10
       sub esp,
       mov eax,
       add esp,
       pop ebx
16
       pop ebp
```



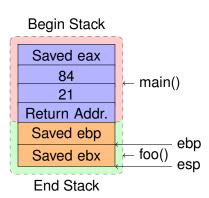
```
push eax
       add esp,
       push ebp
10
       push ebx
       sub esp,
       mov eax,
       add esp,
            ebx
16
            ebp
```

Overview

Registers



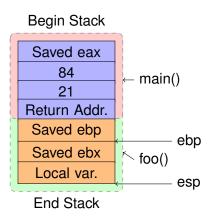
```
push eax
       add esp. 12
       push ebp
       mov ebp, esp ; <--Here
10
       push ebx
       sub esp,
       add esp,
       pop ebx
16
           ebp
```



```
push eax
       add esp. 12
       push ebp
       mov ebp, esp
10
       push ebx
       sub esp,
       mov eax,
       add esp,
           ebx
16
           ebp
```

Overview The Stack Stack Overview Example Standards

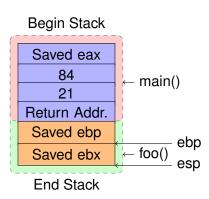
### Stack Example - Part 8



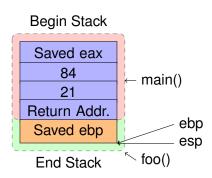
```
push eax
       add esp. 12
       push ebp
       mov ebp, esp
       push ebx
       sub esp,
       add esp,
       pop ebx
16
       pop ebp
```

Overview

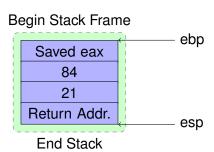
Registers



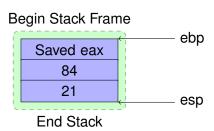
```
push eax
       add esp. 12
       push ebp
       mov ebp, esp
10
       push ebx
       sub esp,
       add esp,
       pop ebx
16
       pop ebp
```



```
push eax
       add esp. 12
       push ebp
       mov ebp, esp
10
       push ebx
       sub esp,
       mov eax,
       add esp,
16
           ebp
```



```
push eax
       add esp, 12
       push ebp
       mov ebp, esp
       push ebx
       sub esp,
       mov eax,
       add esp,
       pop ebx
16
       pop ebp
```



```
push eax
       add esp,
       push ebp
       mov ebp, esp
       push ebx
10
       sub esp,
       mov eax,
            ebx
            ebp
16
```

Overview

Registers

The program is now finished. If there was a ret, the program would have continued backward into the next stack frame. Check out slide-set 2 to learn more about the x86 instructions and what they are used for.

```
<u>pu</u>sh eax
        add esp. 12
        push ebp
        mov ebp, esp
        push ebx
10
        sub esp,
        mov eax,
        add esp,
             ebx
16
             ebp
```