

# CS 330: Network Applications & Protocols

## Introduction to Computer Networks & the Internet

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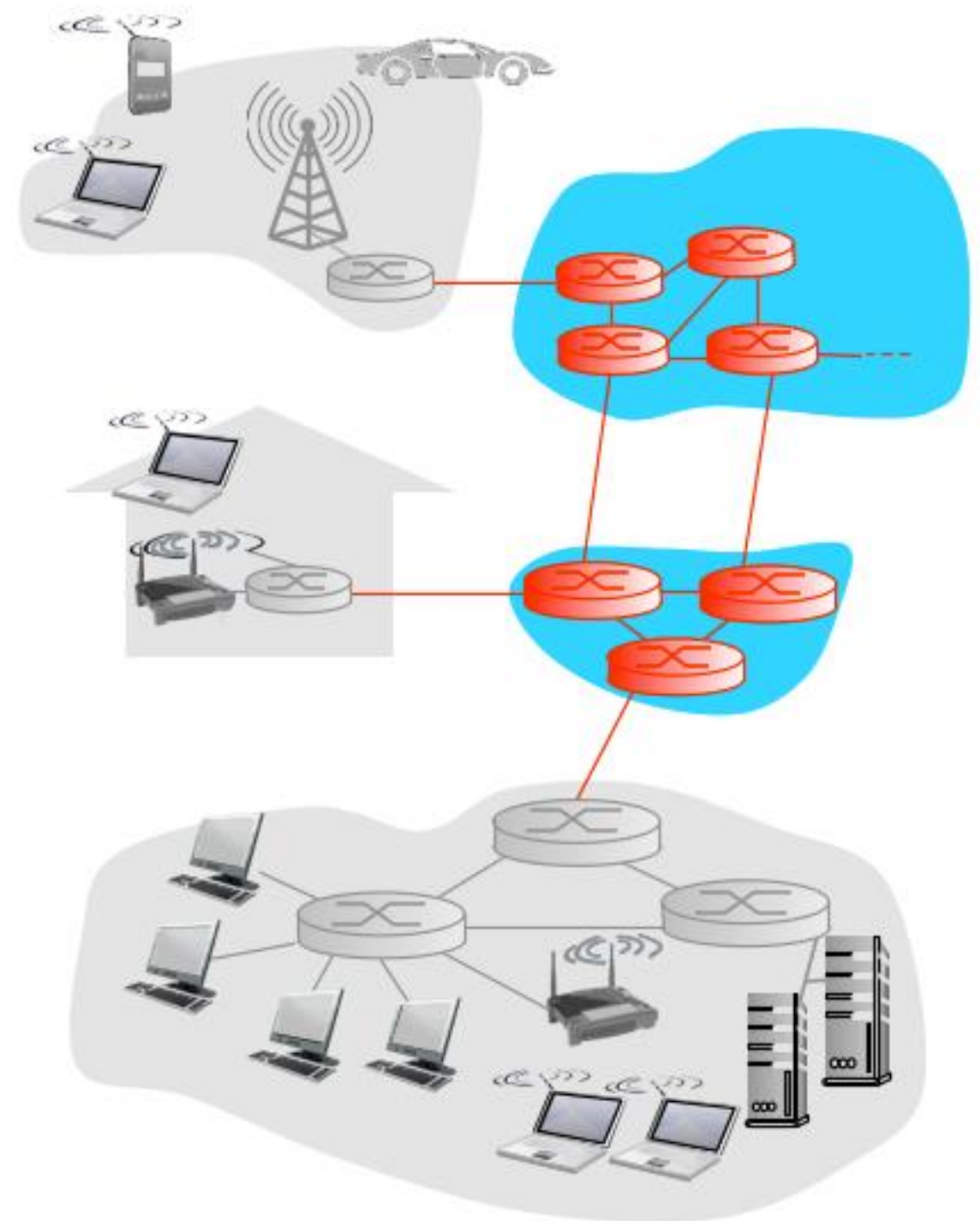
# Introduction

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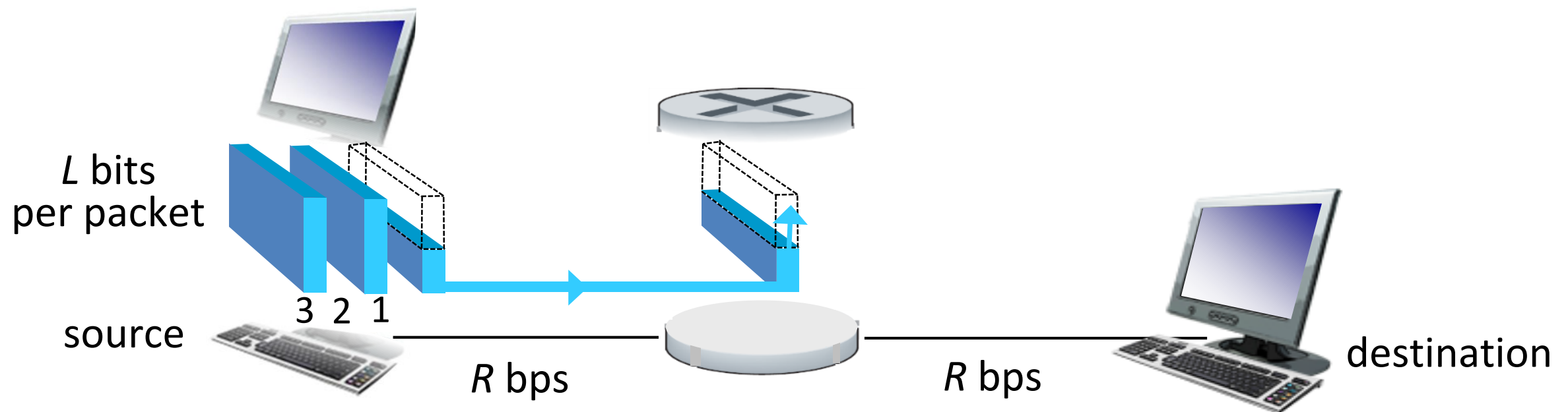
- What is the Internet?
- Network edge
  - End systems, access networks, links
- Network core
  - Packet switching, circuit switching, network structure
- Delay, loss, throughput in networks
- Protocol layers, service models
- Networks under attack: security
- History

# The Network Core

- **Mesh of interconnected routers**
- **Packet-switching:** hosts break application-layer messages into packets
  - Forward packets from one router to the next, across links on path from source to destination
  - Each packet transmitted at full link capacity



# Packet-switching: store-and-forward



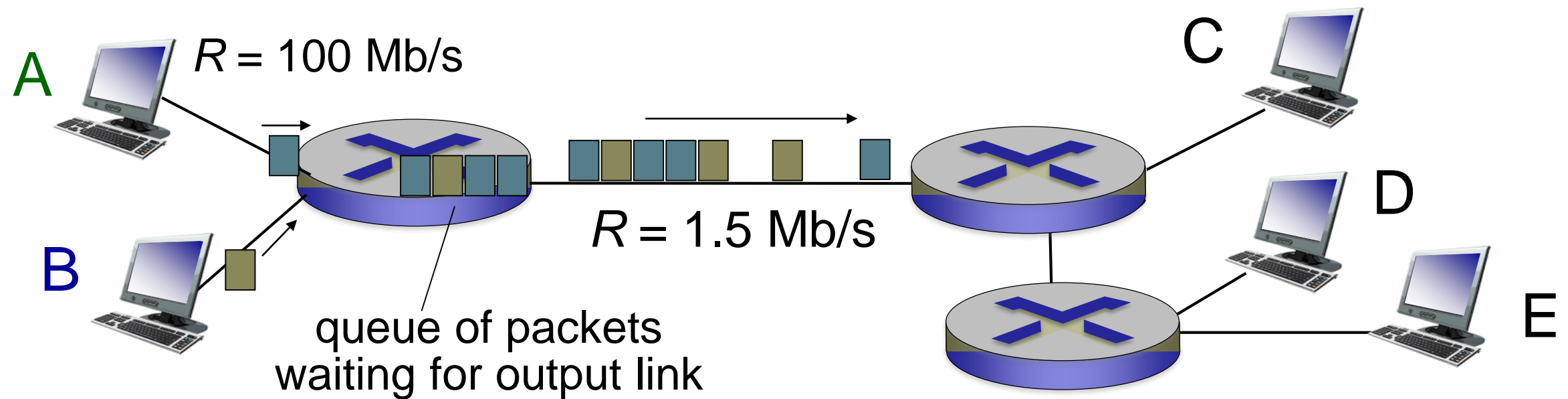
- takes  $L/R$  seconds to transmit (push out)  $L$ -bit packet into link at  $R$  bps
- **store and forward:** entire packet must arrive at router before it can be transmitted on next link
  - end-end delay =  $2L/R$  (assuming zero propagation delay)

## **one-hop numerical example:**

- $L = 7.5$  Mbits
- $R = 1.5$  Mbps
- one-hop transmission delay = 5 sec

} more on delay shortly ...

# Packet Switching: queueing delay, loss



## queuing and loss:

- if arrival rate (in bits) to link exceeds transmission rate of link for a period of time:
  - packets will queue, wait to be transmitted on link
  - packets can be dropped (lost) if memory (buffer) fills up

# Two Key Functions of the Network-Core

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- **Routing**

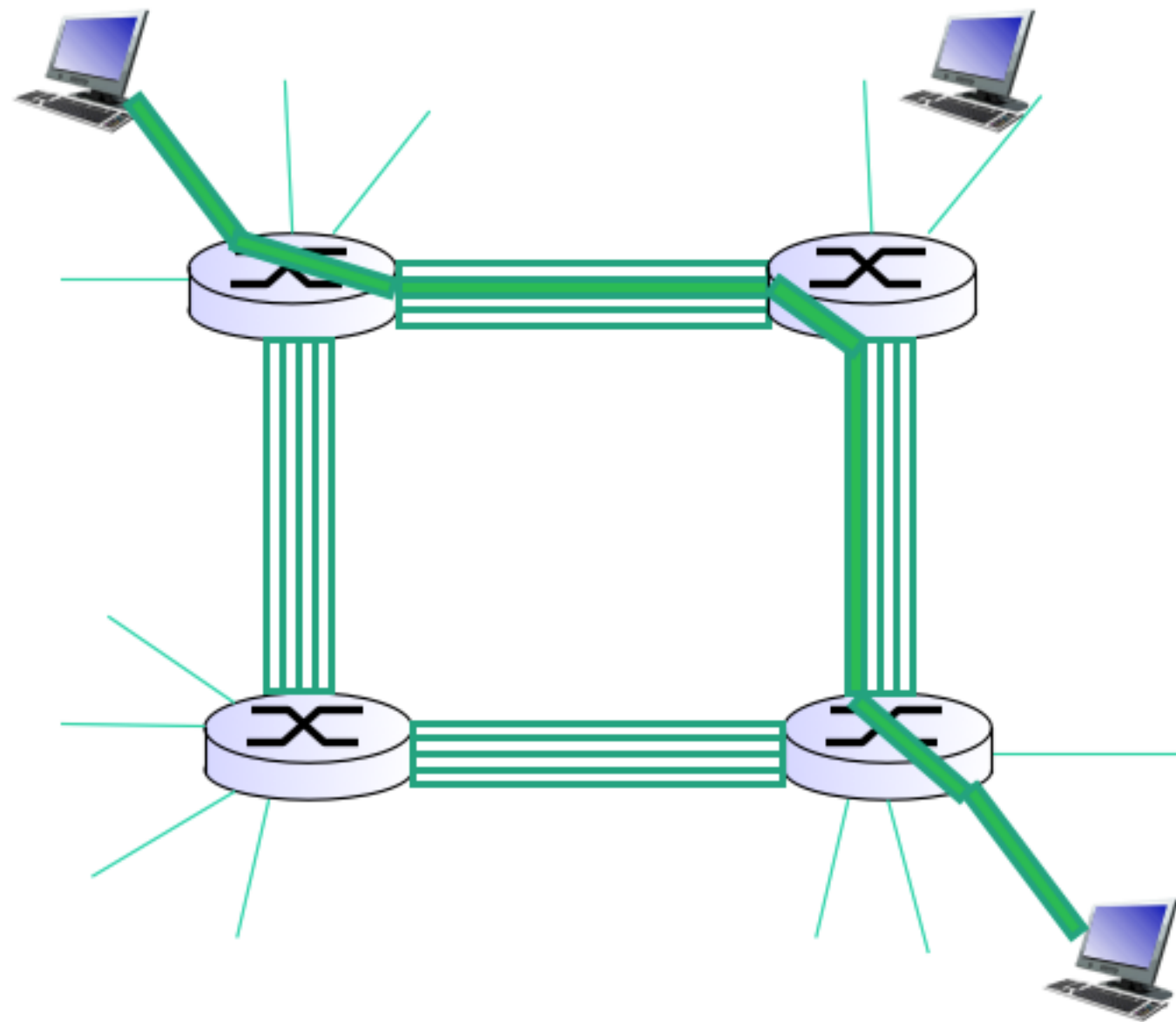
- Determines source-destination route taken by packets
- Utilizes routing algorithms

- **Forwarding**

- Move packets from router's input to appropriate router output

# Alternative Core: Circuit Switching

- **End-to-end resources allocated to, reserved for “call” between source & dest:**
  - In diagram, each link has four circuits
    - Call gets 2<sup>nd</sup> circuit in top link and 1<sup>st</sup> circuit in right link
  - Dedicated resources (i.e. no sharing)
    - Circuit-like (guaranteed) performance
  - Circuit segment idle if not used by a call (no sharing)
  - Commonly used in traditional telephone networks

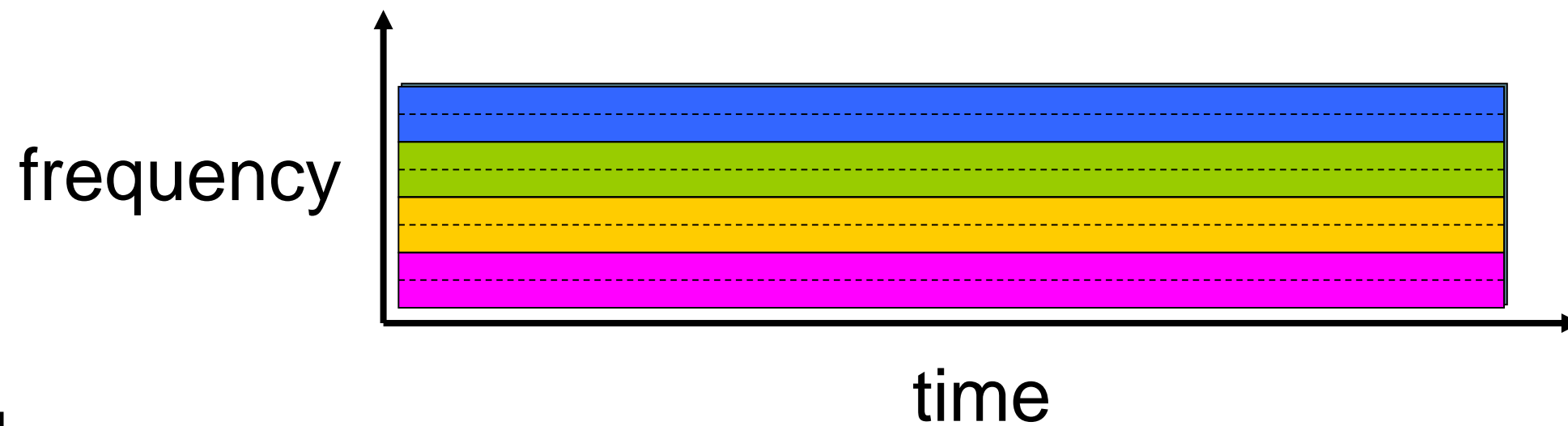


# Circuit switching: FDM versus TDM

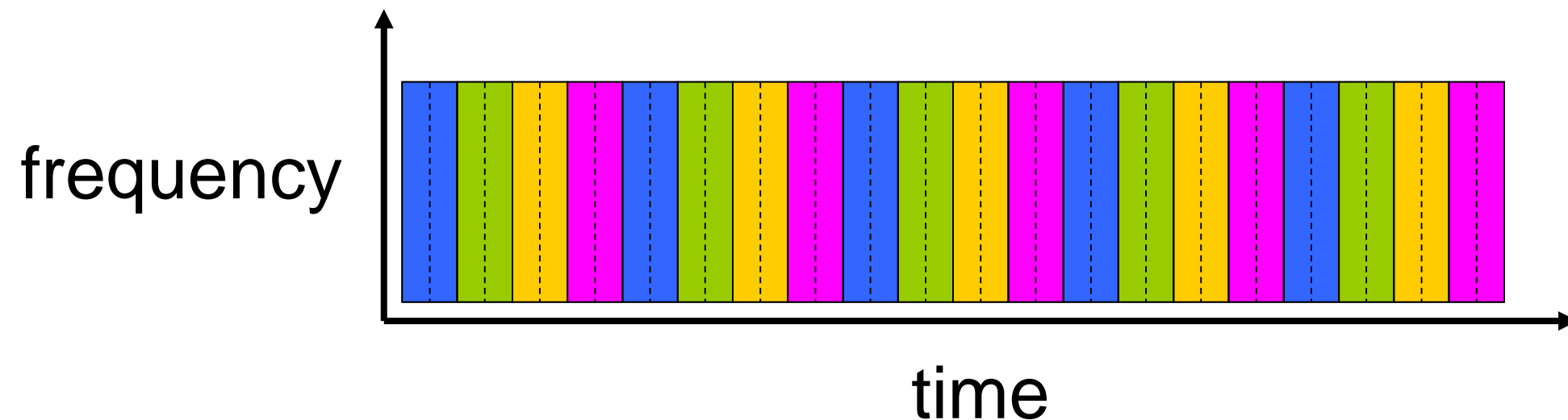
FDM

Example:

4 users



TDM



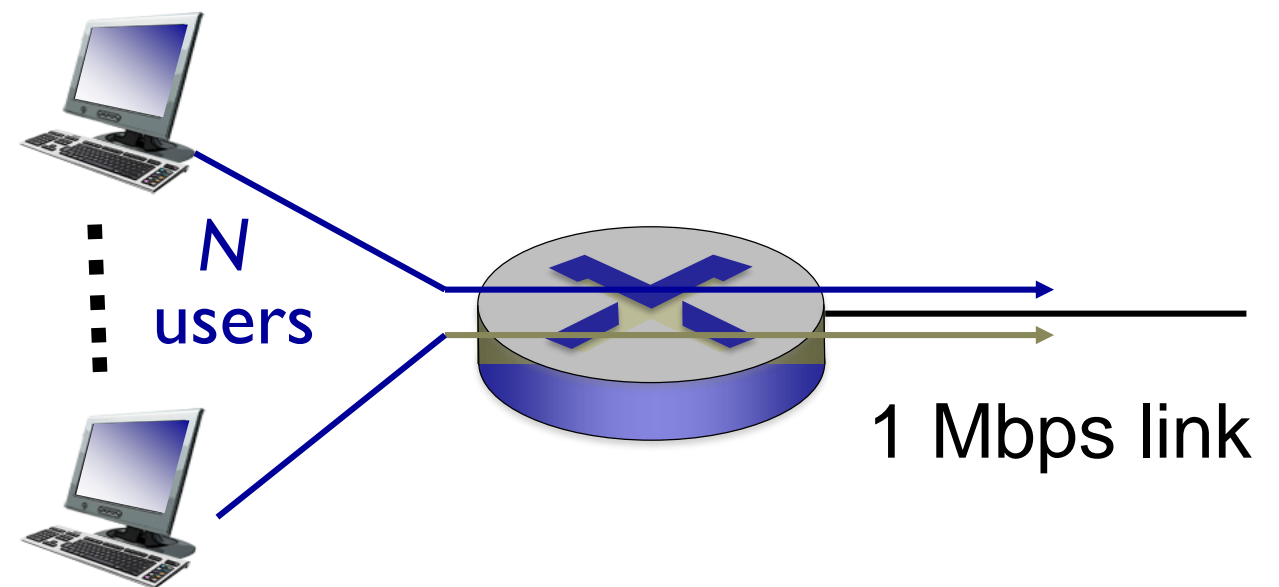


# Packet switching versus circuit switching

***packet switching allows more users to use network!***

**example:**

- **1 Mb/s link**
- **each user:**
  - 100 kb/s when “active”
  - active 10% of time
- ***circuit-switching:***
  - 10 users
- ***packet switching:***
  - with 35 users, probability  $> 10$  active at same time is less than .0004 \*



**Q:** how did we get value 0.0004?

**Q:** what happens if  $> 35$  users ?

# Packet switching versus circuit switching

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**is packet switching a “slam dunk winner?”**

- **great for bursty data**
  - resource sharing
  - simpler, no call setup
- **excessive congestion possible: packet delay and loss**
  - protocols needed for reliable data transfer, congestion control
- **Q: How to provide circuit-like behavior?**
  - bandwidth guarantees needed for audio/video apps
  - still an unsolved problem (chapter 7)
- **Q: Human analogies of reserved resources (circuit switching) versus on-demand allocation (packet-switching)?**

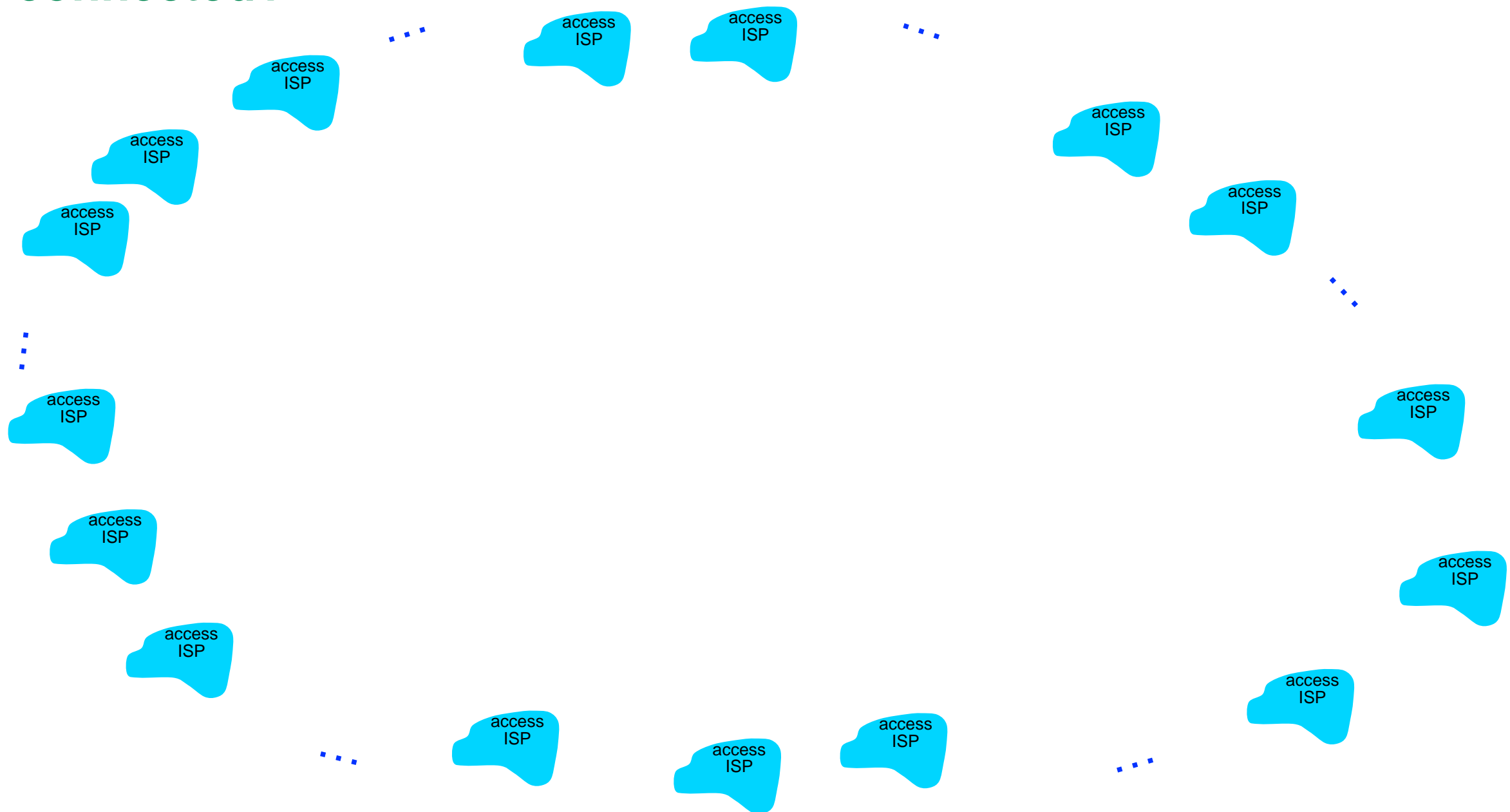
# Internet structure: network of networks

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- End systems connect to Internet via **access ISPs** (Internet Service Providers)
  - residential, company and university ISPs
- Access ISPs in turn must be interconnected.
  - so that any two hosts can send packets to each other
- Resulting network of networks is very complex
  - evolution was driven by **economics and national policies**
- Let's take a stepwise approach to describe current Internet structure

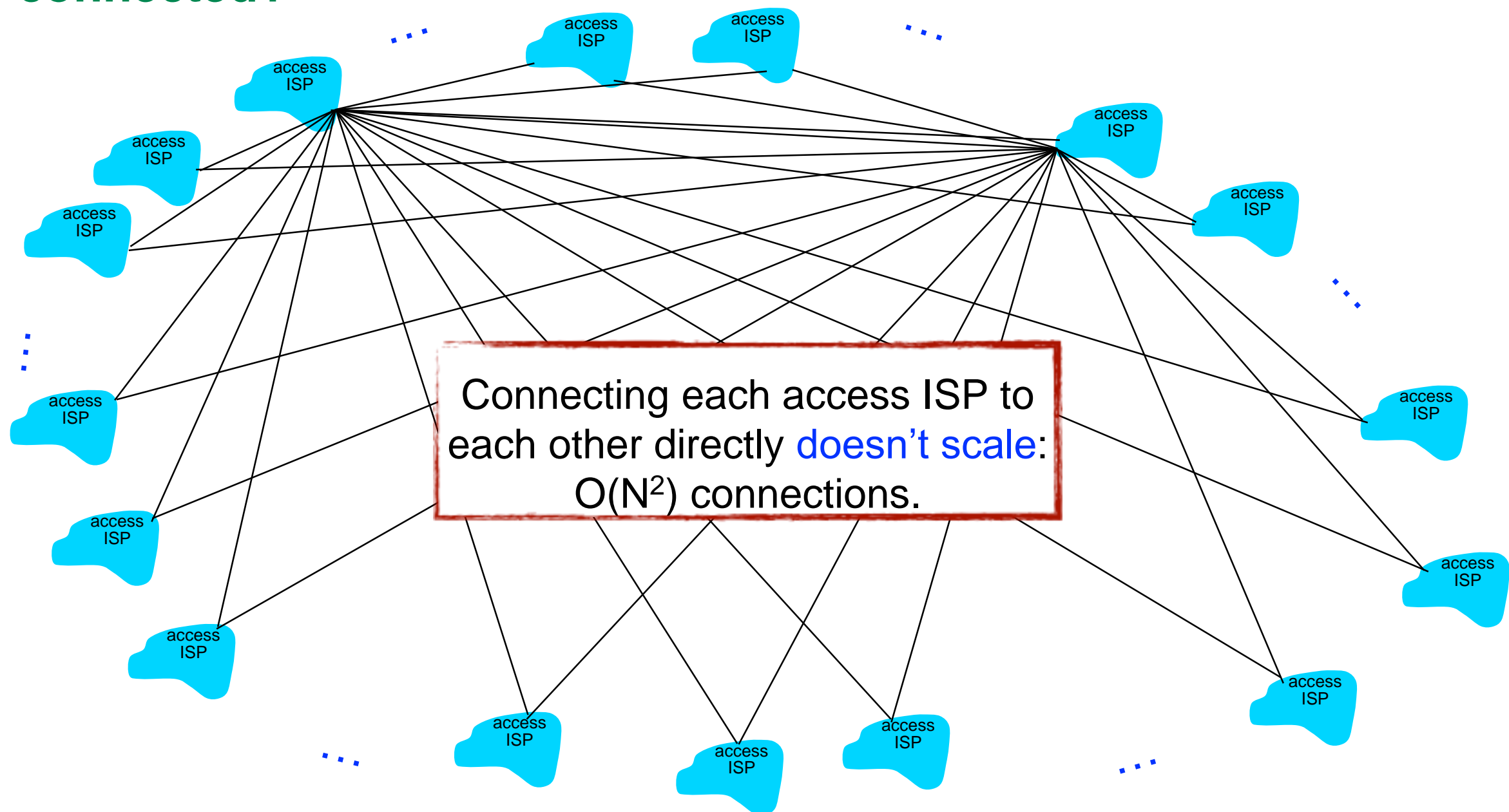
# Internet Structure: Network of Networks

- **Question:** Given millions of access ISPs, how should they be connected?



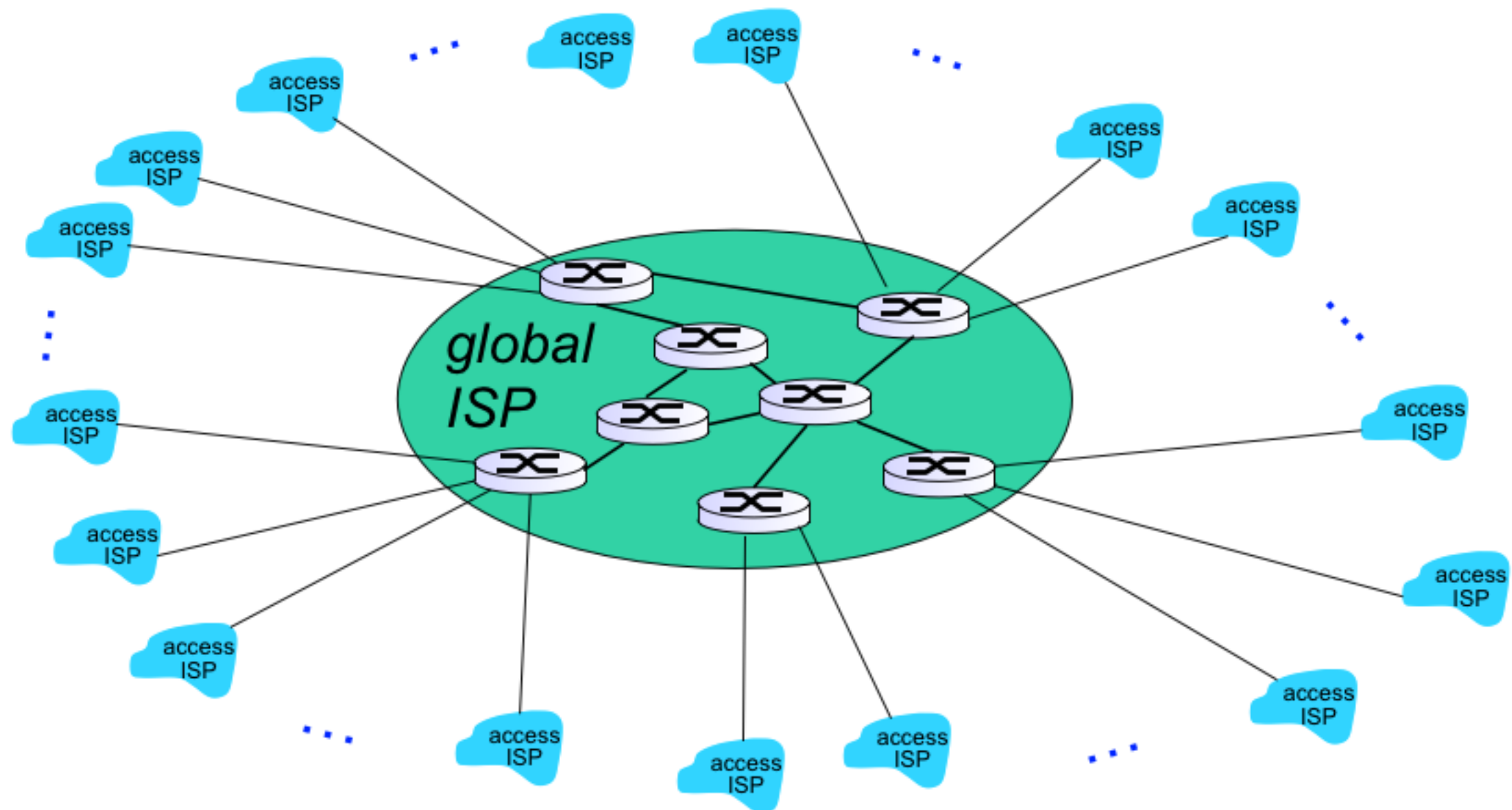
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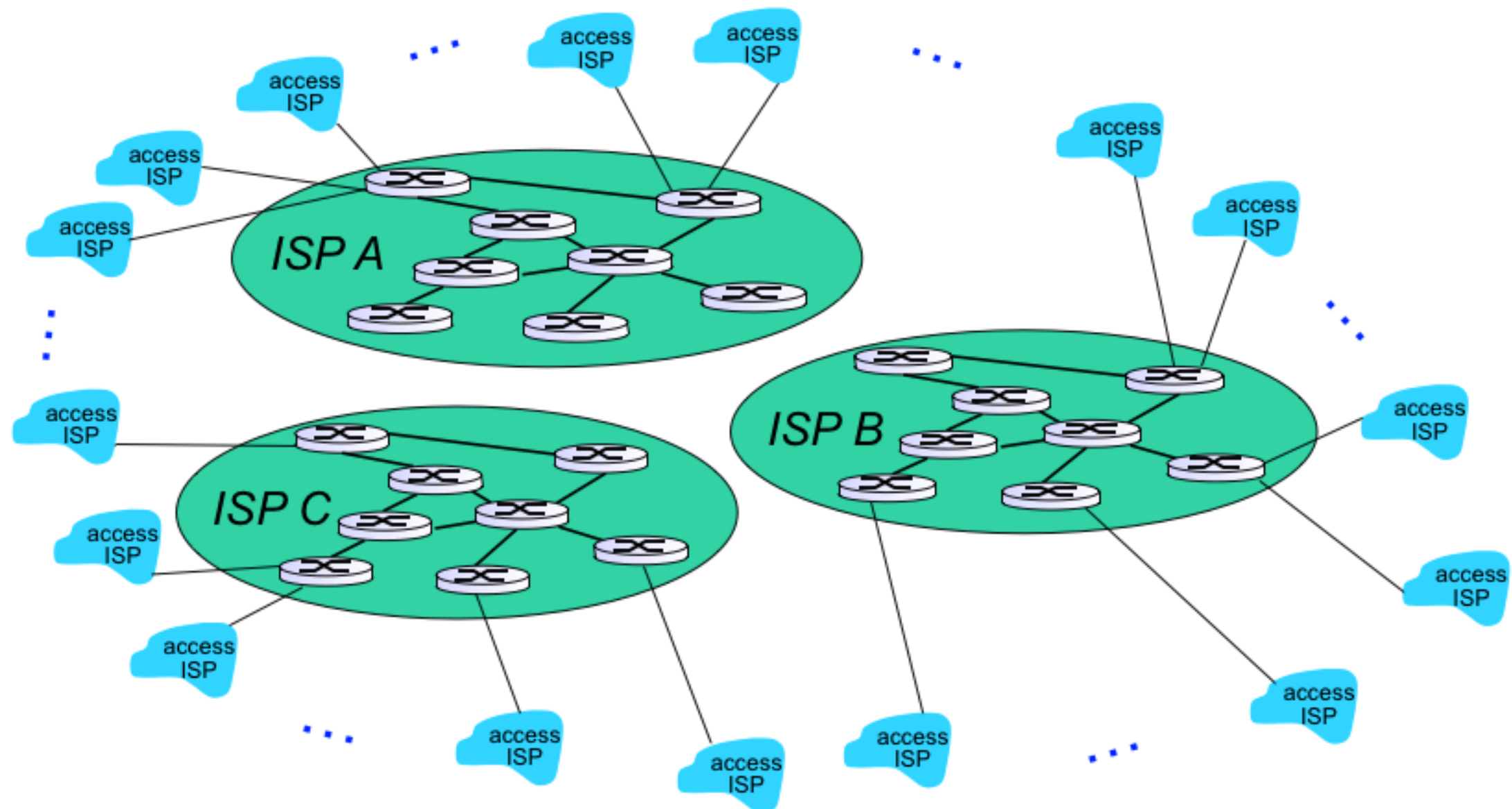
# Internet Structure: Network of Networks

- **Option:** Connect each access ISP to a global transit ISP? Customer and provider ISPs have economic agreement.



# Internet Structure: Network of Networks

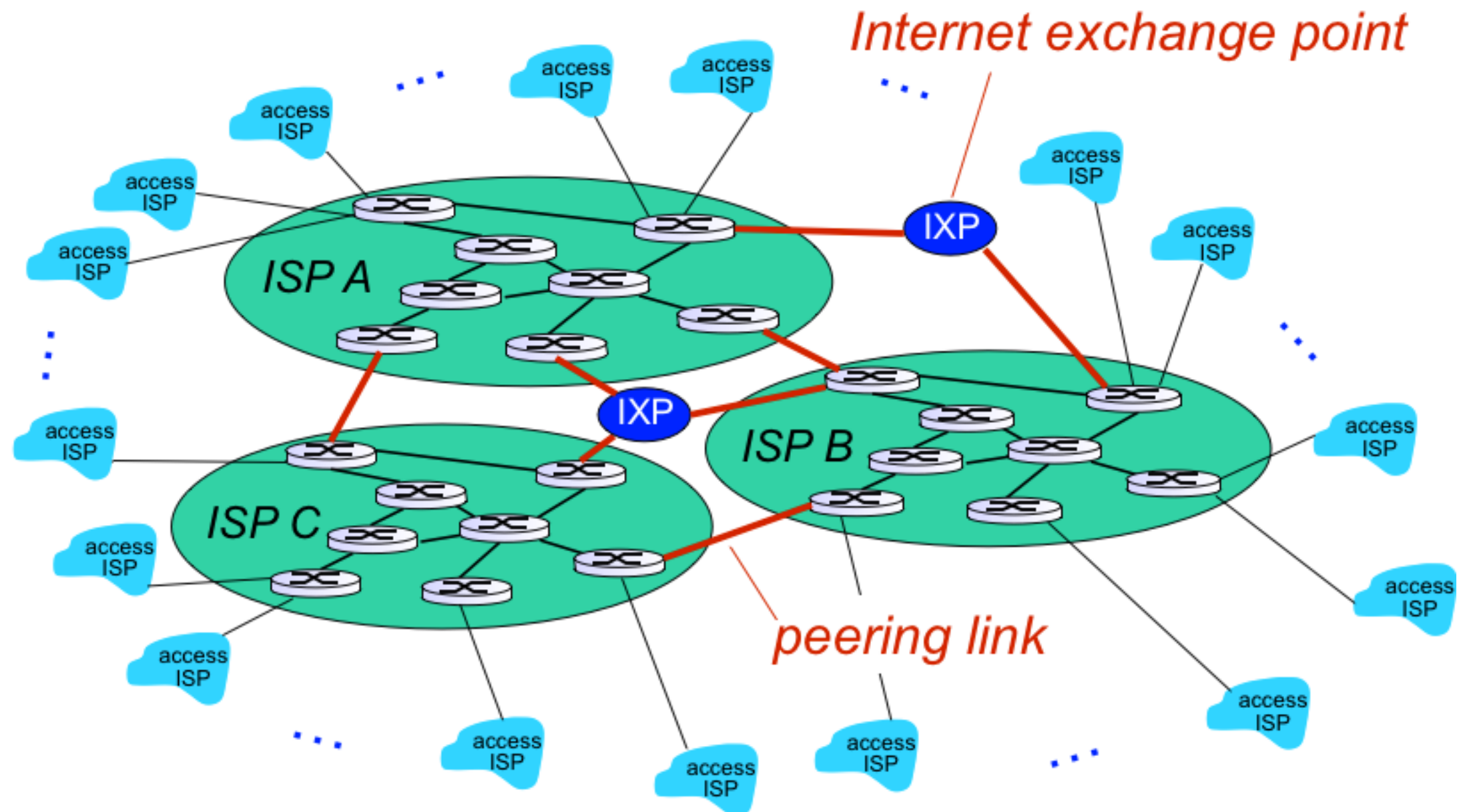
- But if one global ISP is viable business, there will be competitors ...





# Internet Structure: Network of Networks

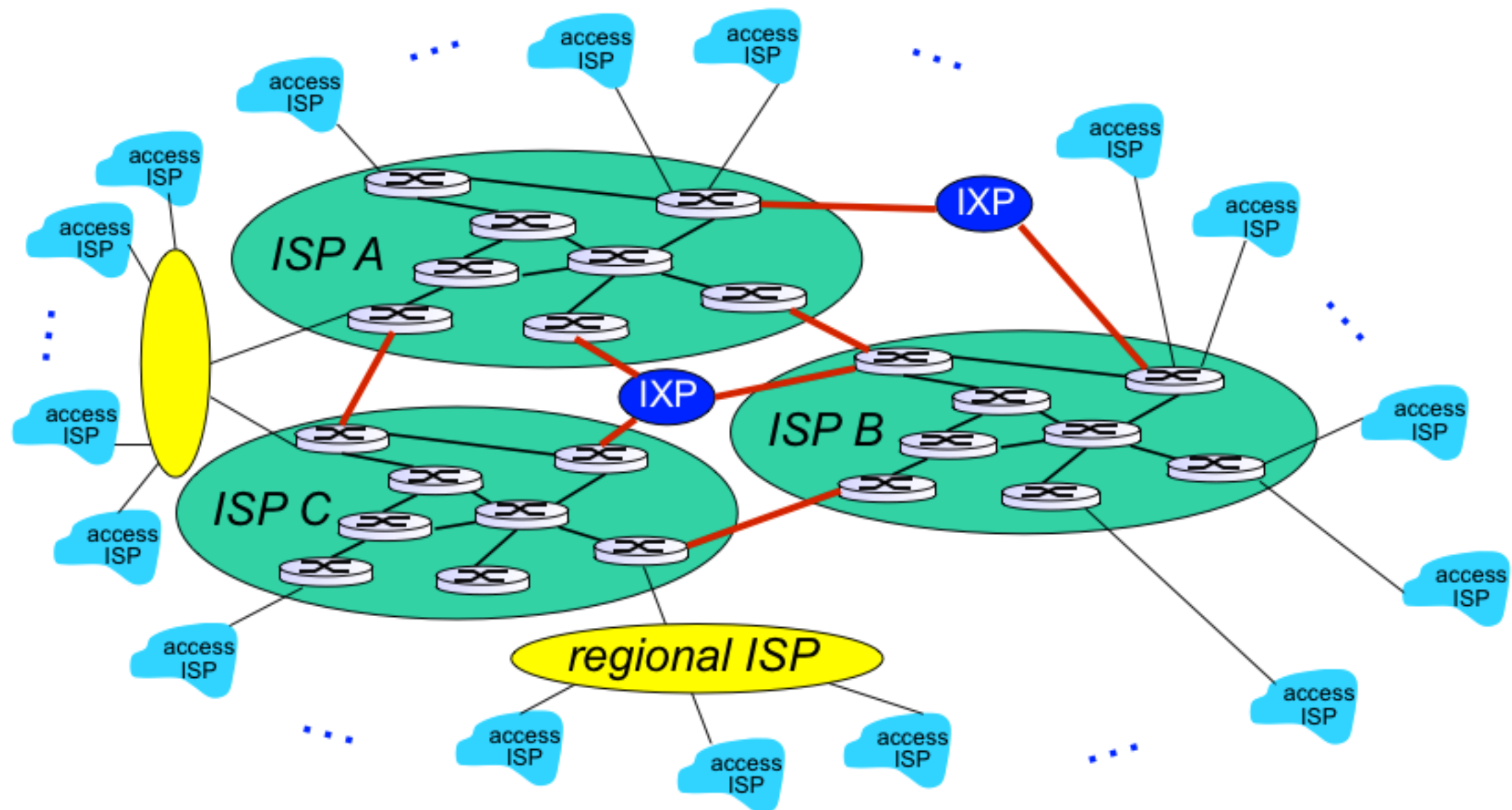
- But if one global ISP is viable business, there will be competitors ... which must be interconnected





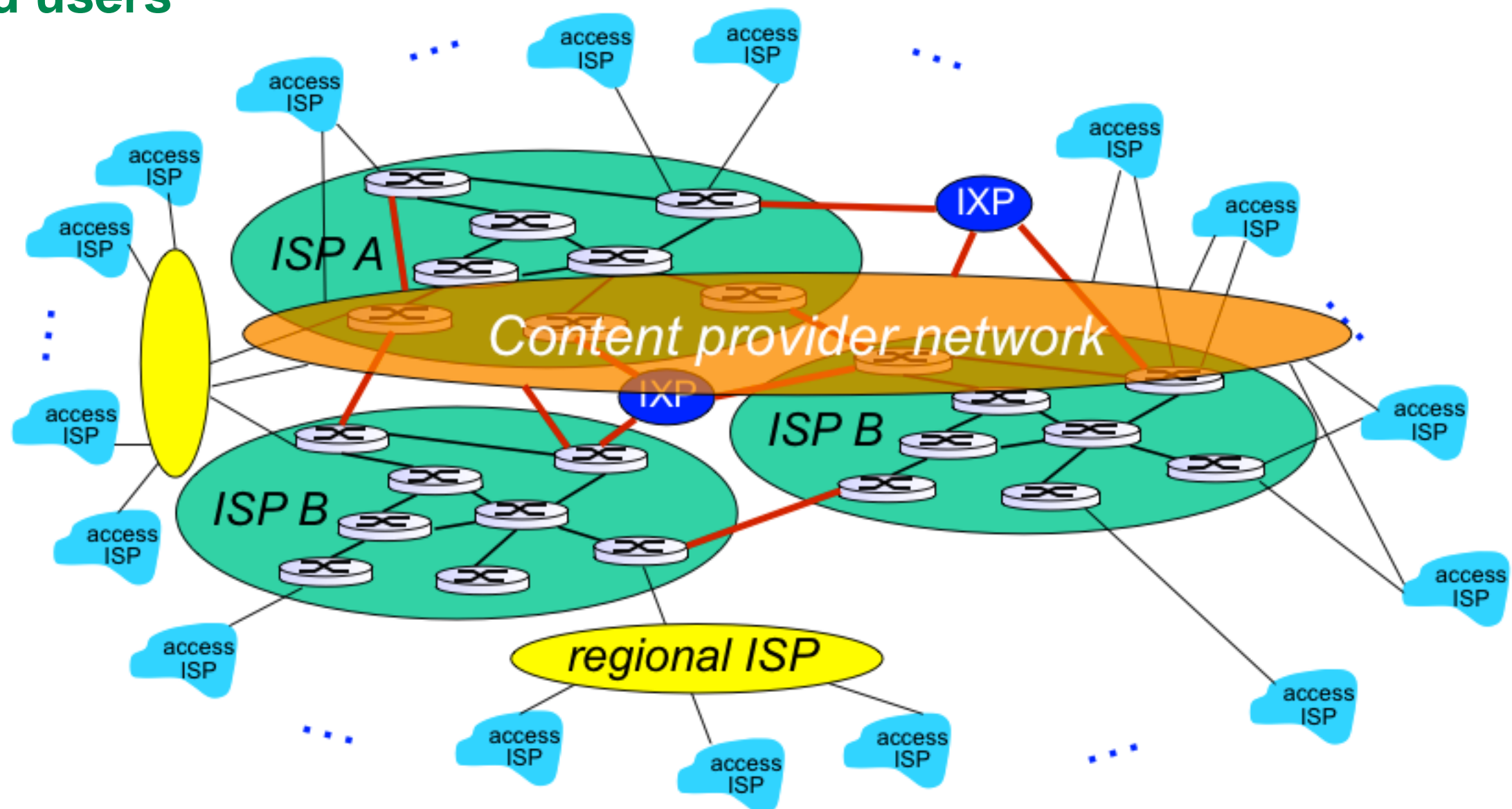
# Internet Structure: Network of Networks

- ... and regional networks may arise to connect access ISPs to ISPs

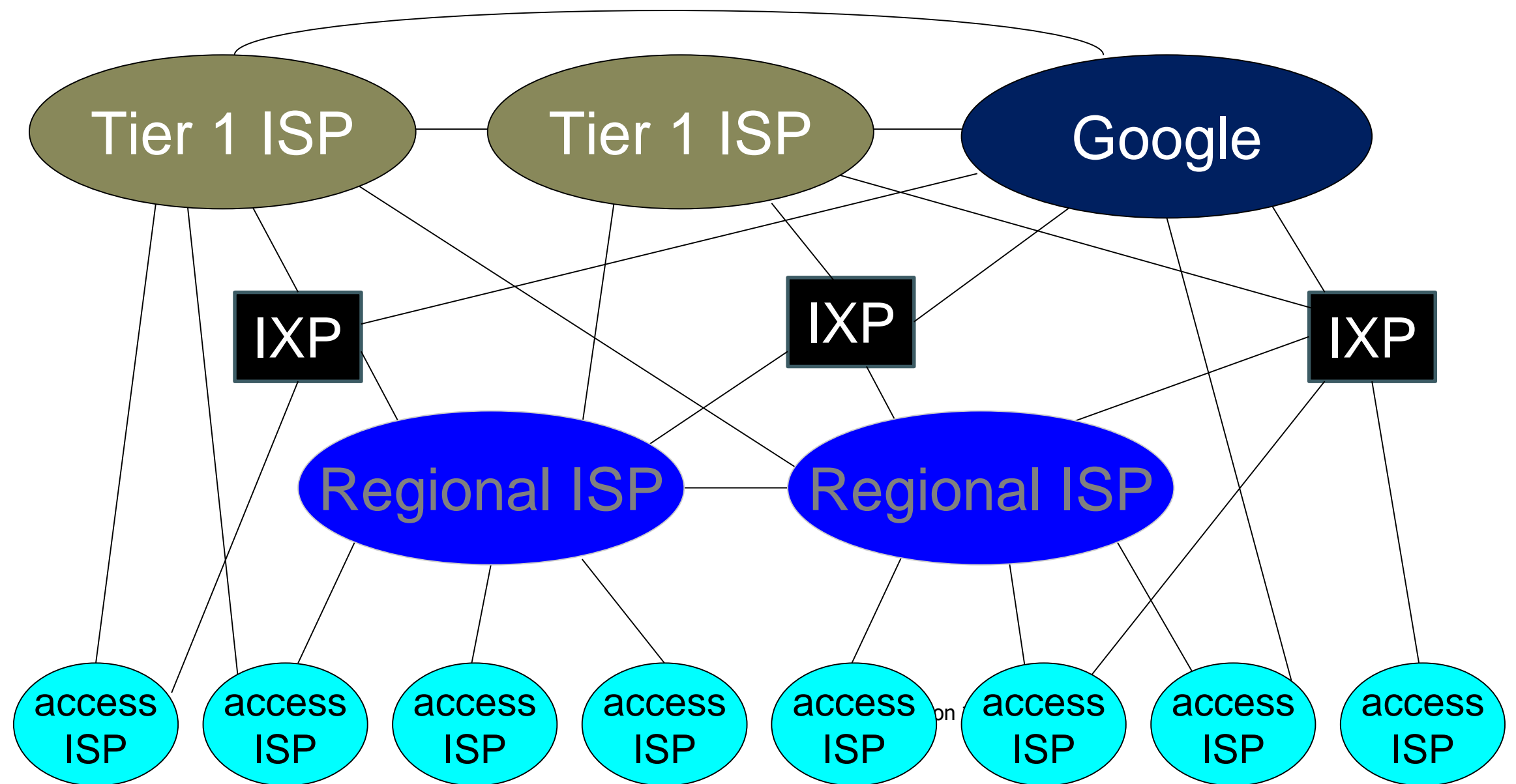


# Internet Structure: Network of Networks

- ... and content provider networks (e.g. Google, Microsoft, Akamai) may run their own network, to bring services and content close to end users



# Internet structure: network of networks



- **Center: small # of well-connected large networks**
  - “tier-1” commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
  - content provider network (e.g., Google): private network that connects its data centers to Internet, often bypassing tier-1, regional ISPs

# Introduction

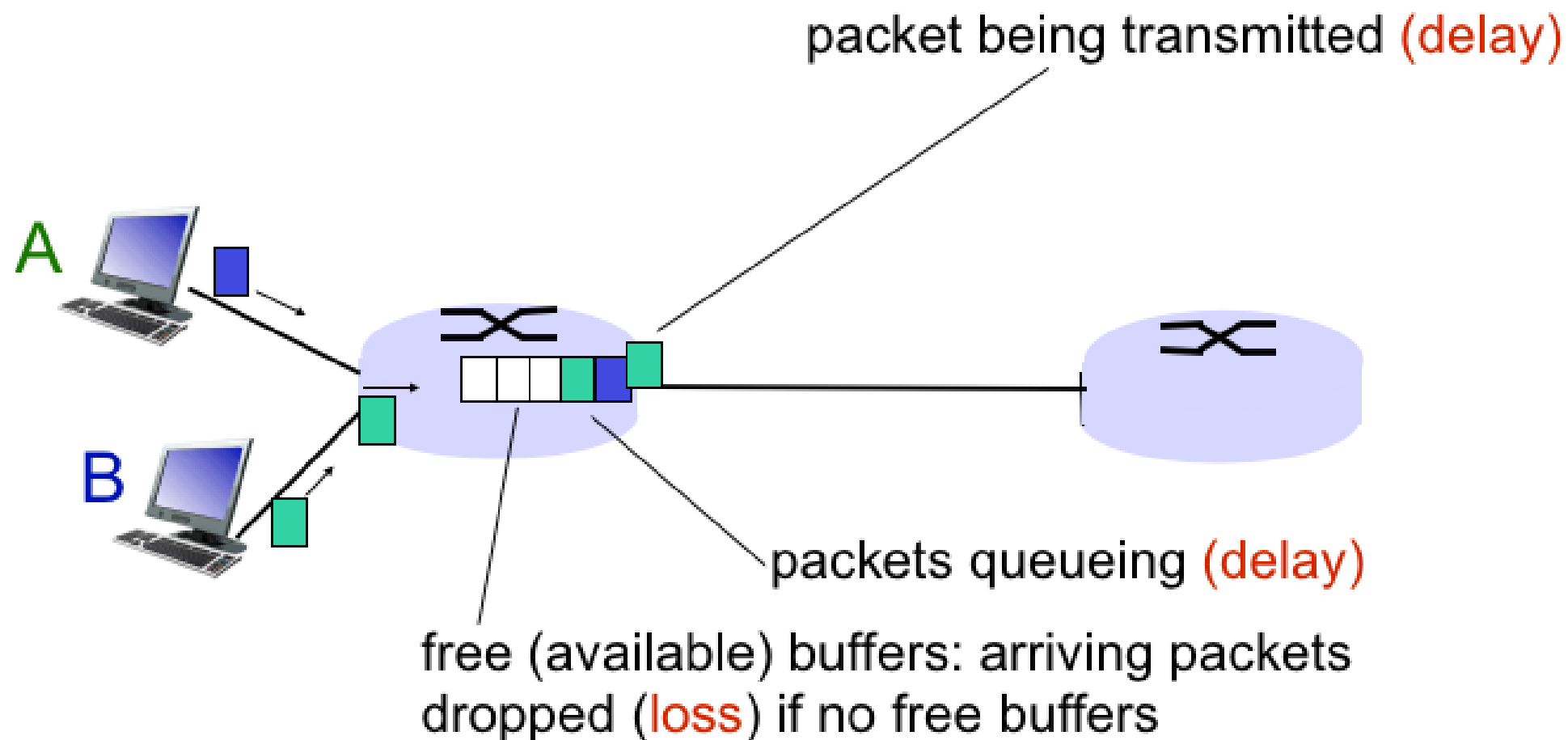
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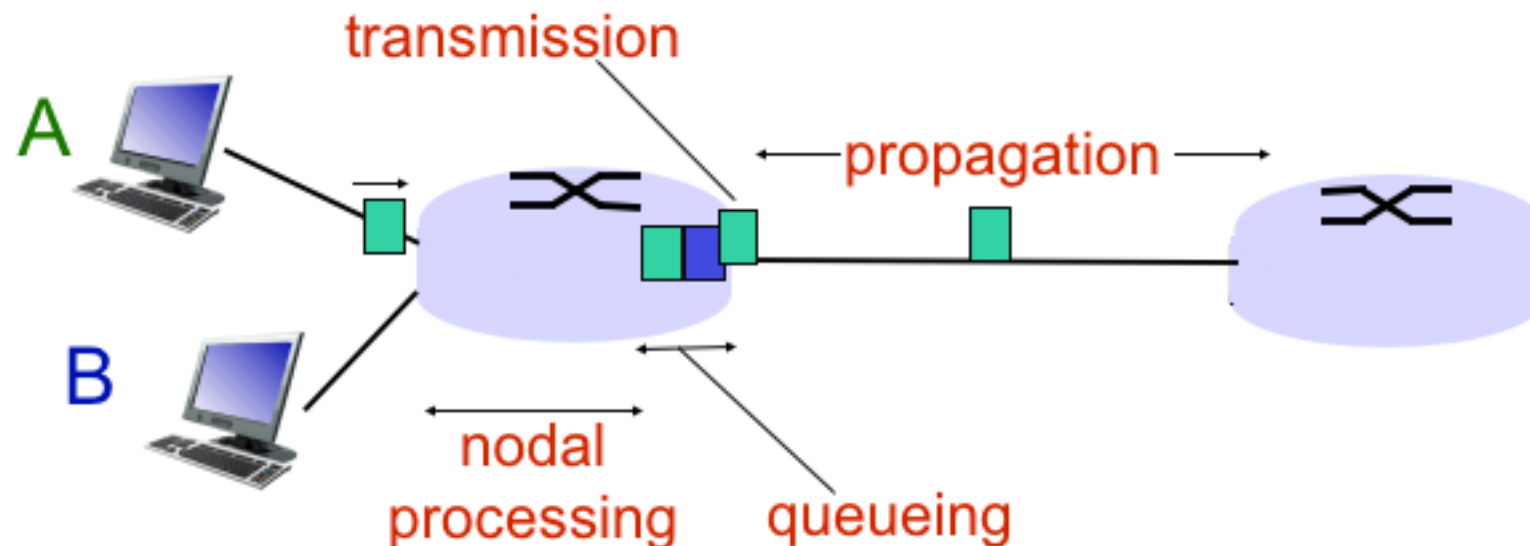
# How Does Packet Delay Occur?

- **Packets queue in router buffers**

- Packet arrival rate to link (temporarily) exceeds output link capacity
- Packets queue, wait for turn (i.e. they are **delayed**)



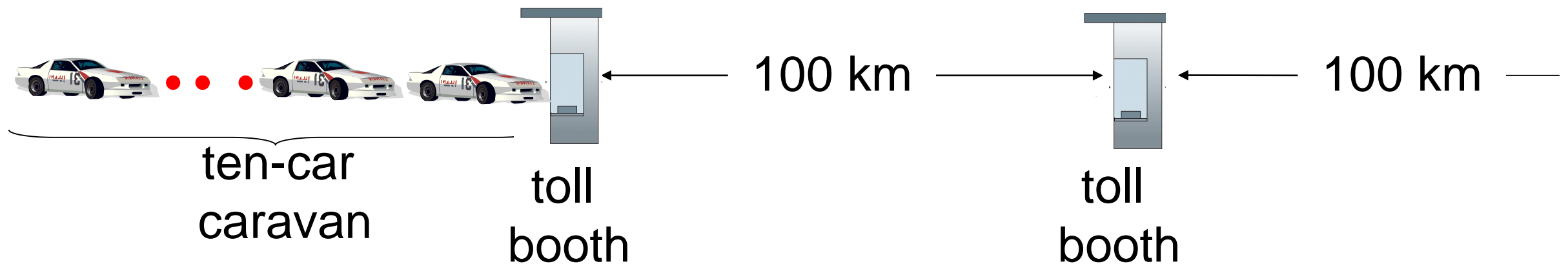
# Four Sources of Packet Delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

- **$d_{\text{proc}}$ : Nodal processing delay** - Check for bit errors, determine output link
- **$d_{\text{queue}}$ : Queueing delay** - Time waiting for output link
- **$d_{\text{trans}}$ : Transmission delay** - (packet length) / (link bandwidth)
- **$d_{\text{prop}}$ : Propagation delay** - (length of physical link) / (propagation speed)

# Caravan analogy



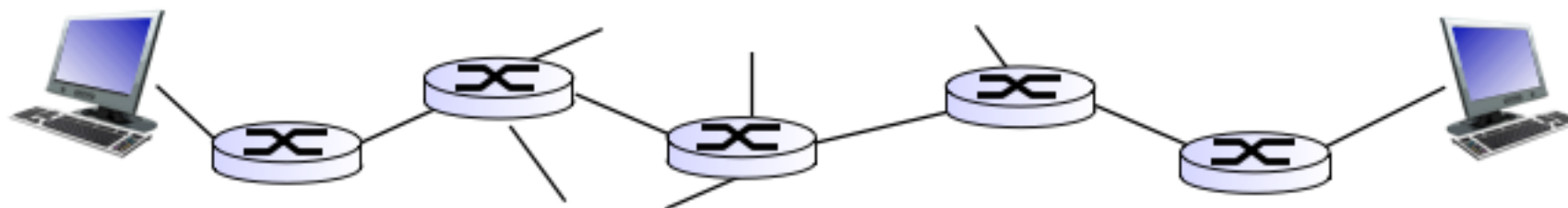
- cars “propagate” at 100 km/hr
- toll booth takes 12 sec to service car (bit transmission time)
- car ~ bit; caravan ~ packet
- **Q: How long until caravan is lined up before 2nd toll booth?**
- time to “push” entire caravan through toll booth onto highway =  $12 * 10 = 120$  sec
- time for last car to propagate from 1st to 2nd toll booth:  $100\text{km} / (100\text{km/hr}) = 1$  hr
- **A: 62 minutes**



# “Real” Internet Delays and Routes

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- What do “real” Internet delay & loss look like?
- Use **traceroute** program
  - Provides delay measurement from source to router along end-to-end Internet path towards destination
  - For all  $i$ :
    - Sends 3 packets that will reach router  $i$  on path towards destination
    - Router  $i$  will return packets to sender
    - Sender times interval between transmission and reply

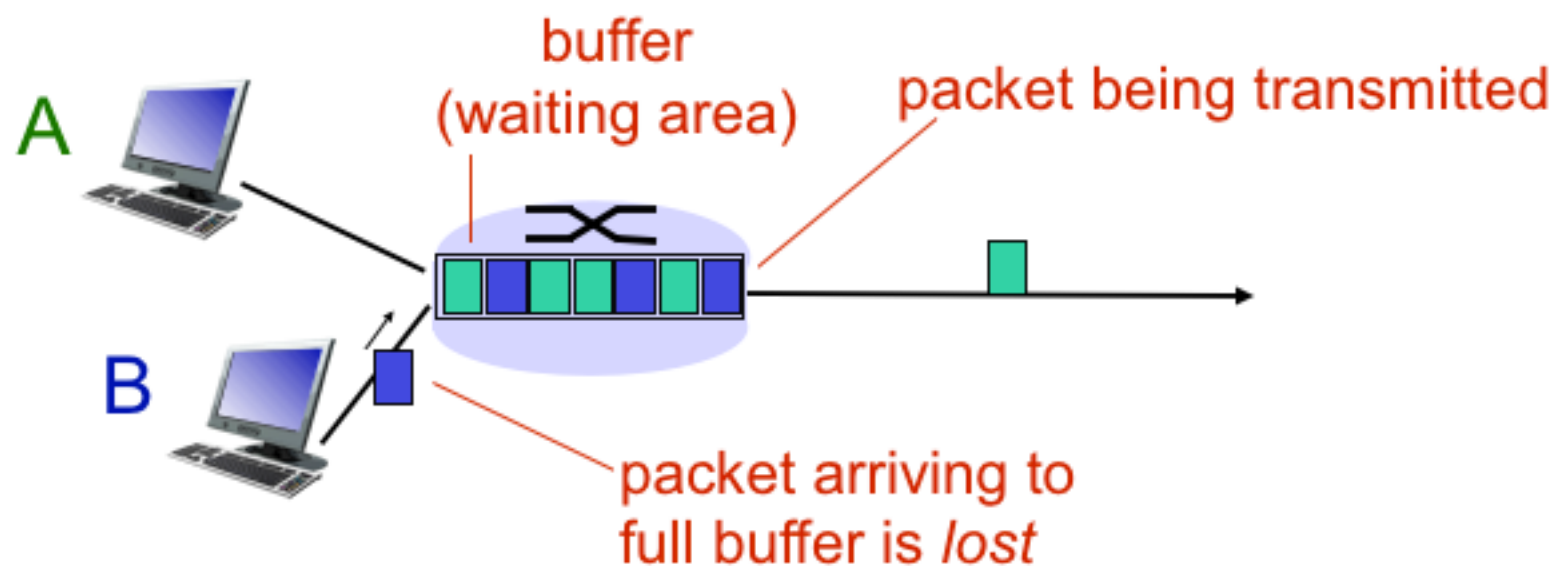




# How Does Packet Loss Occur?

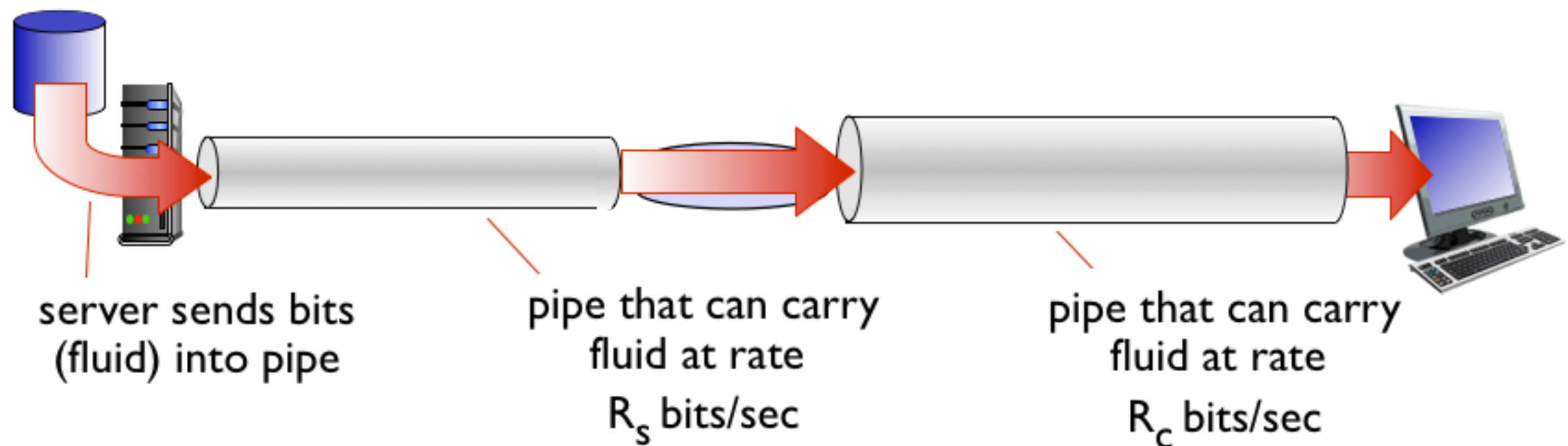
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- Packet queue preceding link has finite capacity
- Packets arriving to a full queue are dropped
- Lost packets may be retransmitted by previous node, by original source, or not at all



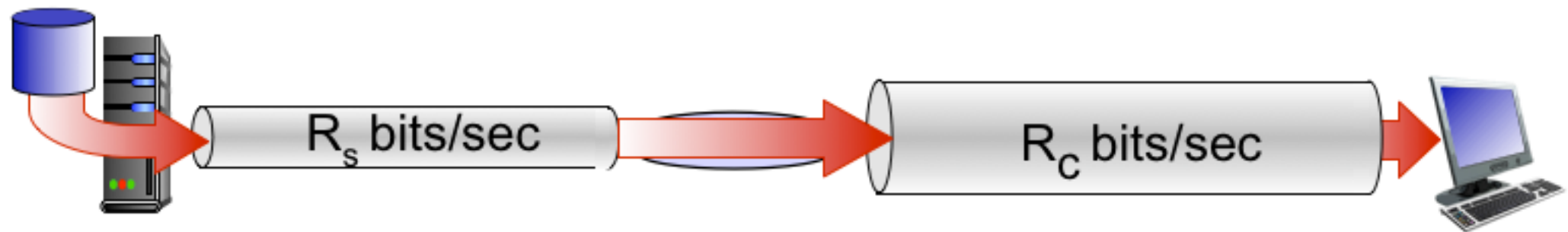
# Throughput

- **Throughput: the rate (bits/time unit) at which bits can be transferred between sender/receiver**
  - Instantaneous throughput: rate at given point in time
  - Average throughput: rate over longer period of time

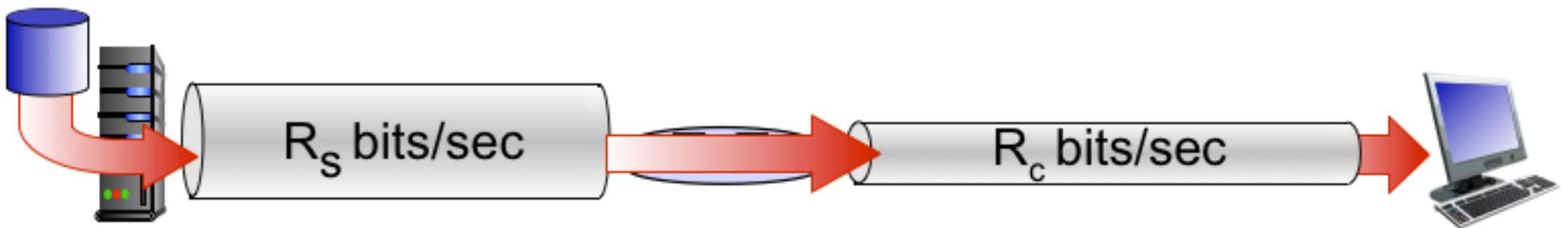


## Throughput (Cont.)

- $R_s < R_c$  : What is average end-to-end throughput?



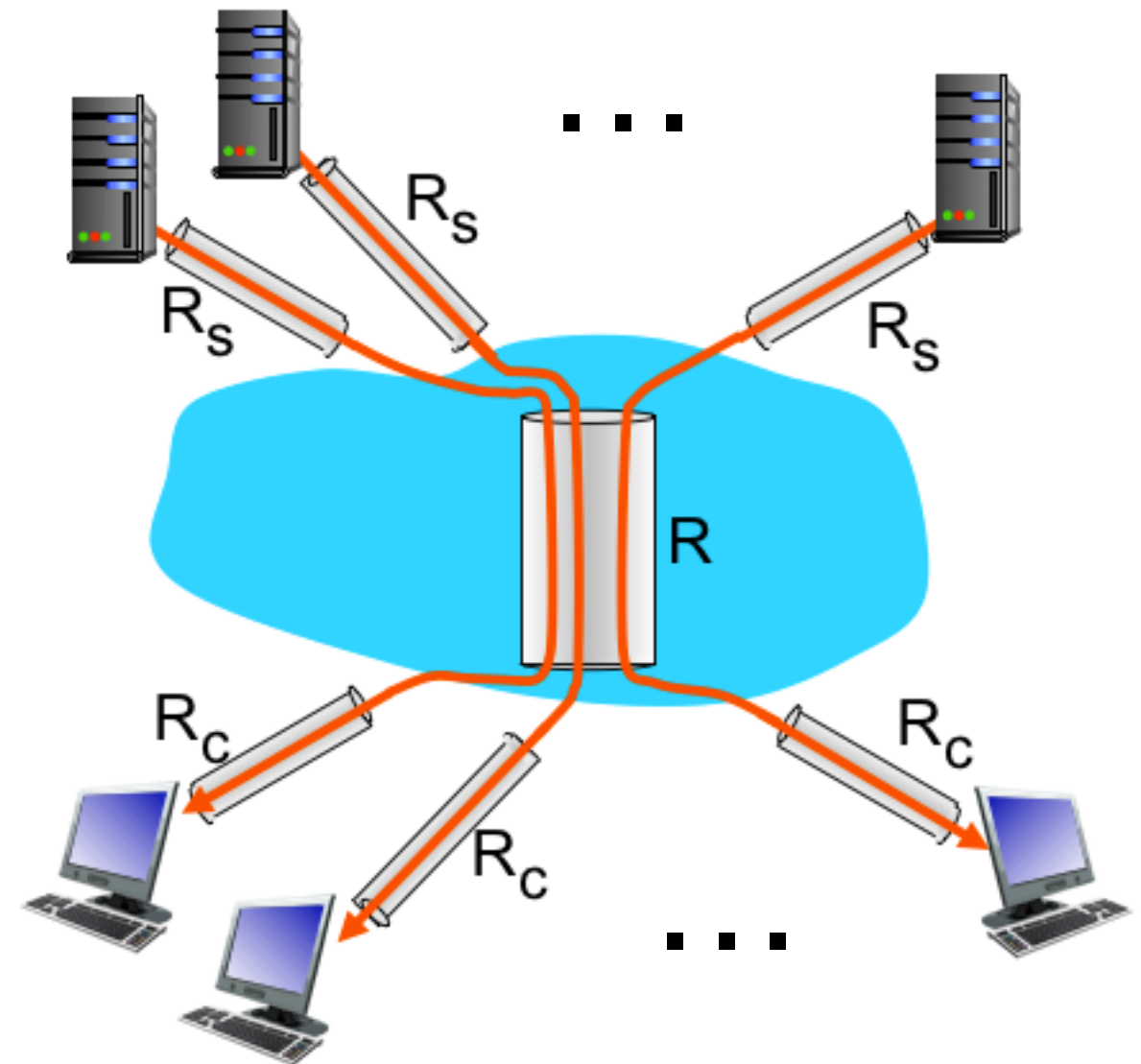
- $R_s > R_c$  : What is average end-to-end throughput?



***Bottleneck Link***: A link on an end-to-end path that constrains the end-end throughput

# Throughput: Internet Scenario

- **Per-connection end-to-end throughput**
  - $\min(R_c, R_s, R/10)$
- **In practice,  $R_c$  or  $R_s$  is often bottleneck**



10 connections (fairly) share  
backbone bottleneck link  $R$  bits/sec

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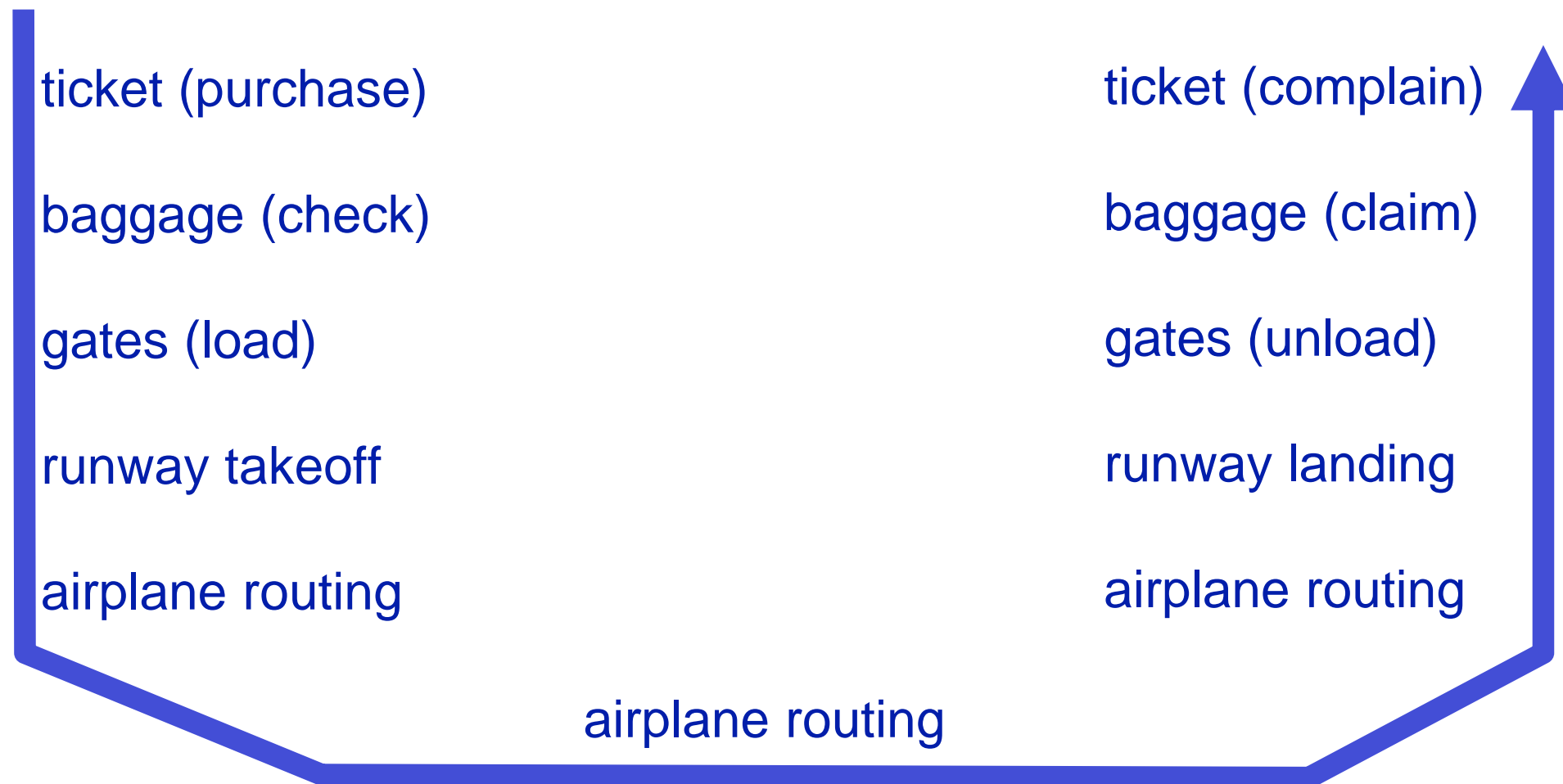
# Protocol “Layers”

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- **Networks are complex, with many components:**
  - Hosts
  - Routers
  - Links of various media
  - Applications
  - Protocols
  - Hardware, software
- **Need some way to organize everything**

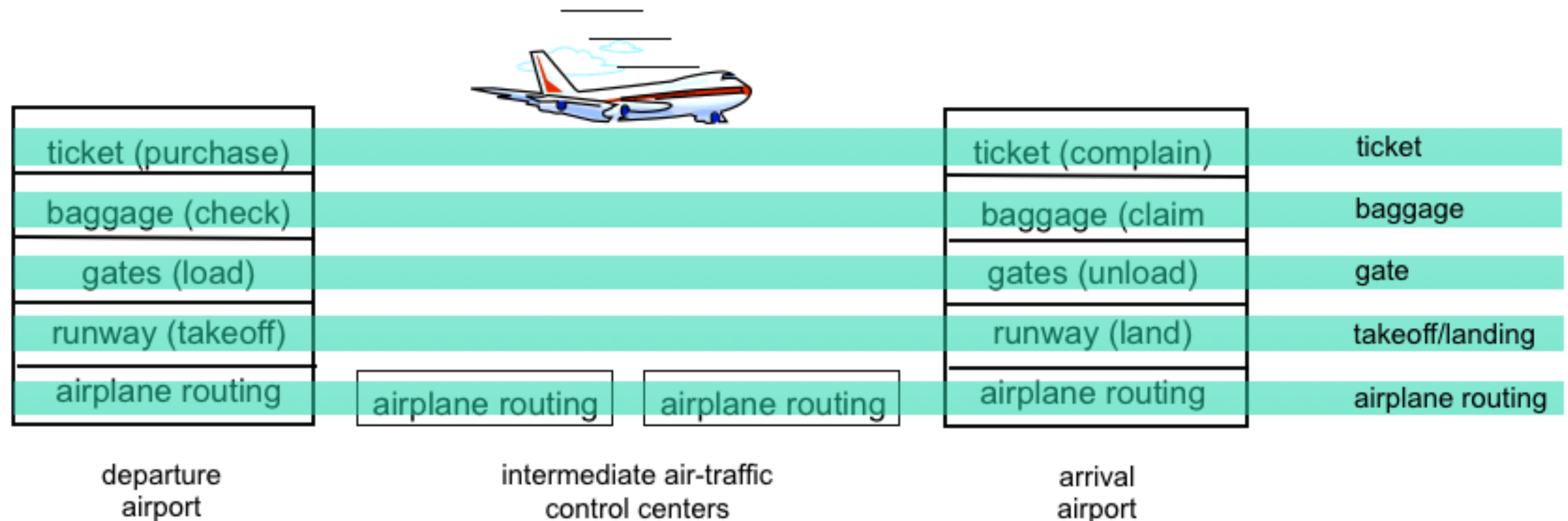
# Example: Organization of Air Travel

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- **Each layer implements a service**
  - Uses its own internal-layer actions
  - Relies on services provided by layer below

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# Why Layering?

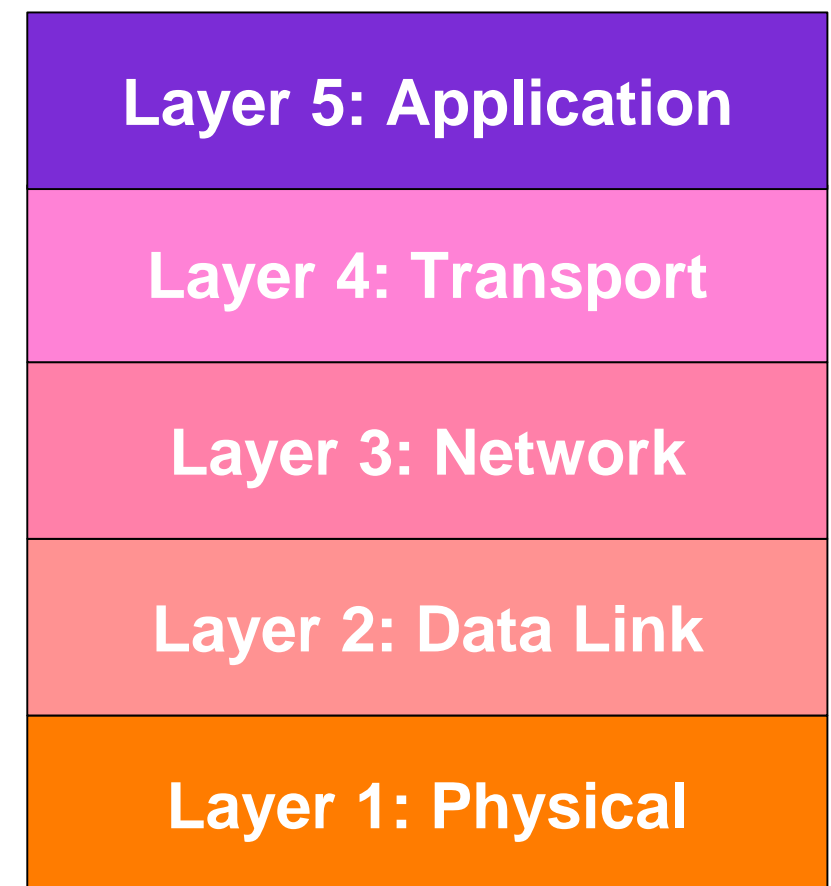
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- **Systems can be complex**
- **Explicit structure allows clear identification and relationship of complex system components**
  - Layered reference model for discussion
- **Modularization eases maintenance, updating of system**
- **Changing implementation of a single layer's service is transparent to rest of system**
  - e.g. change in gate procedure doesn't affect rest of system

# Five-Layer Internet Protocol Stack (TCP/IP Model)

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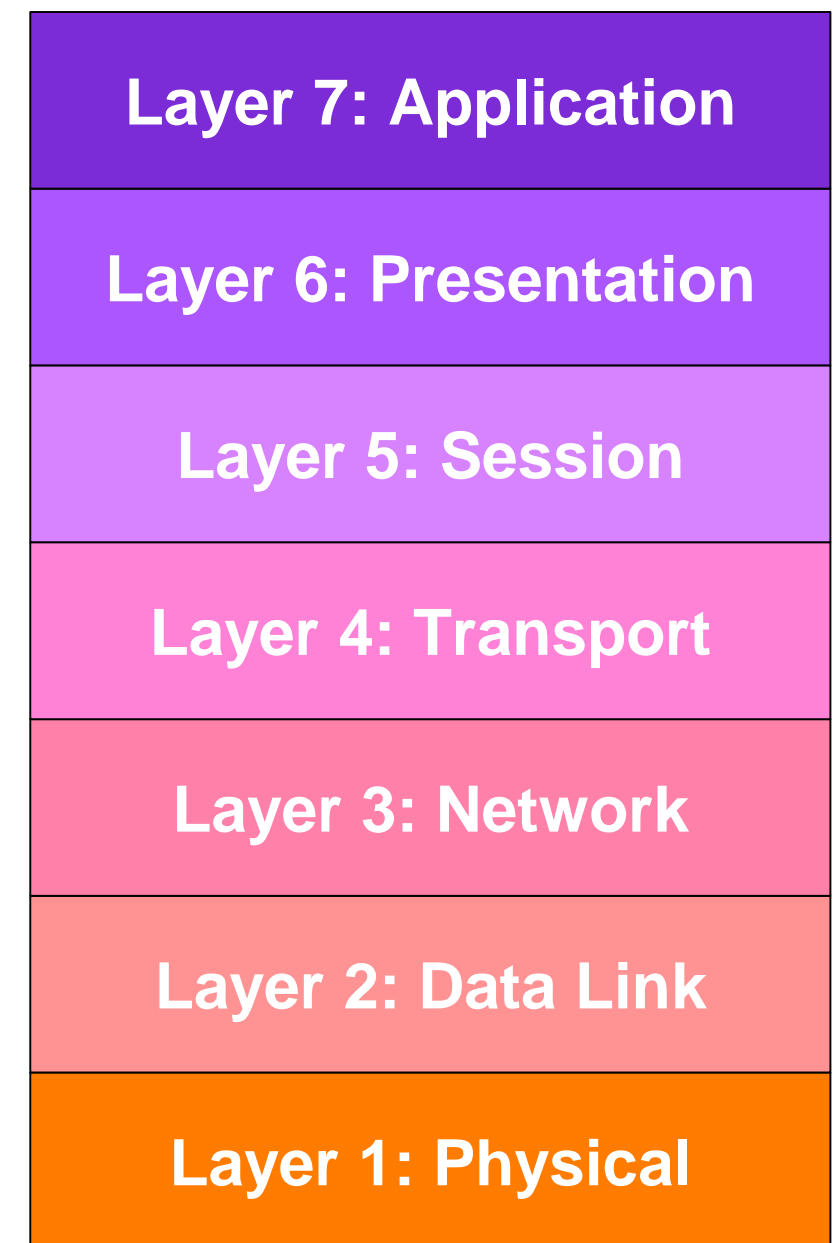
- **Application:** supporting network applications
  - FTP, SMTP, HTTP
- **Transport:** process-to-process data transfer
  - TCP, UDP
- **Network:** routing of datagrams from source to destination
  - IP, routing protocols
- **Link:** data transfer between neighboring network elements
  - Ethernet, 802.11 (WiFi), PPP
- **Physical:** bits on the wire



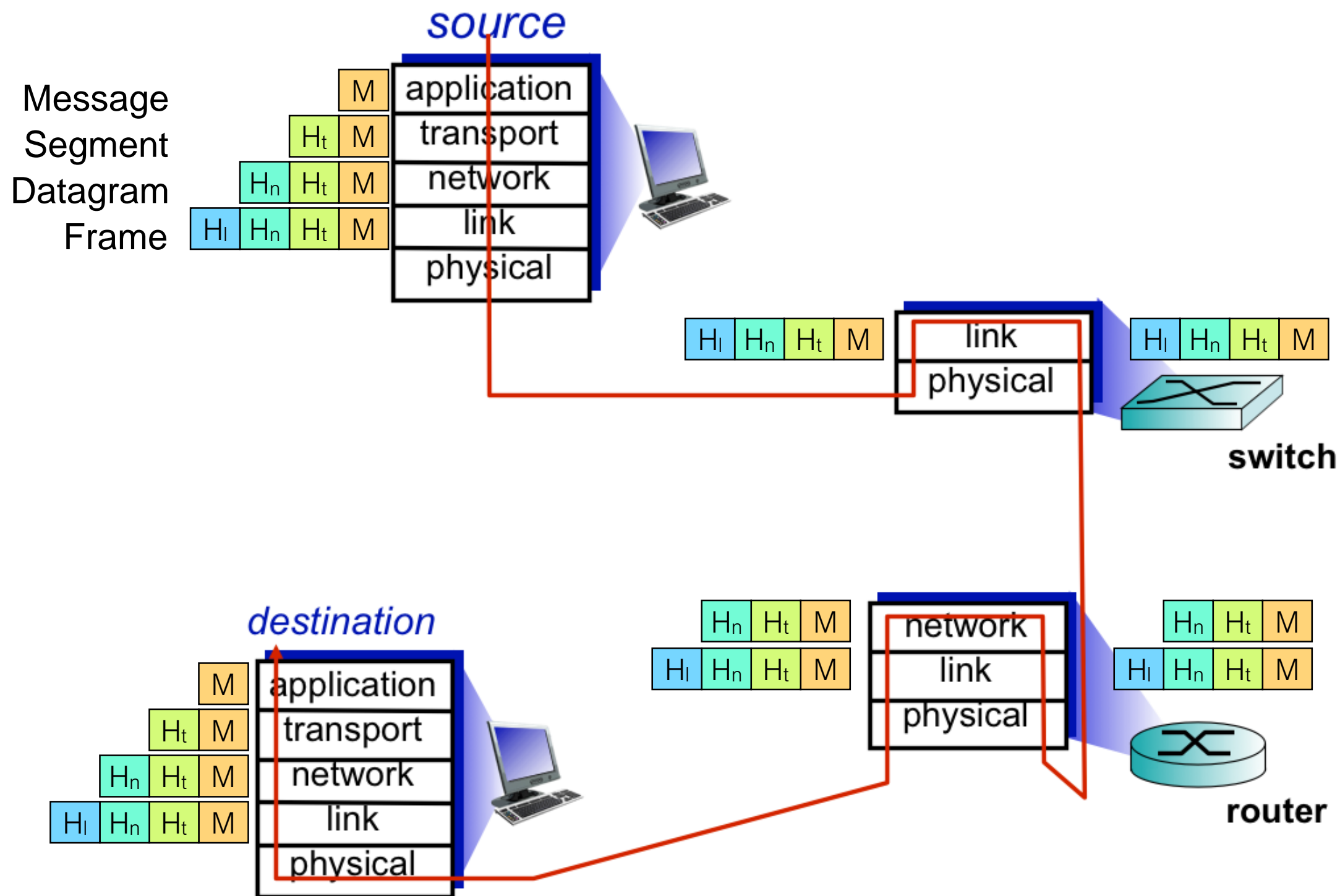
# Seven-Layer ISO/OSI Network Model

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- **Presentation:** allows applications to interpret meaning of data
  - e.g. encryption, compression, machine-specific conventions
- **Session:** synchronization, checkpointing, recovery of data exchange
- Internet stack is missing these layers
- These services, if needed, must be implemented in the Application Layer



# Encapsulation



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# Network security

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- **Field of network security:**
  - how bad guys can attack computer networks
  - how we can defend networks against attacks
  - how to design architectures that are immune to attacks
- **Internet not originally designed with (much) security in mind**
  - *original vision*: “a group of mutually trusting users attached to a transparent network” 😊
  - Internet protocol designers playing “catch-up”
  - security considerations in all layers!

# Bad guys: put malware into hosts via Internet

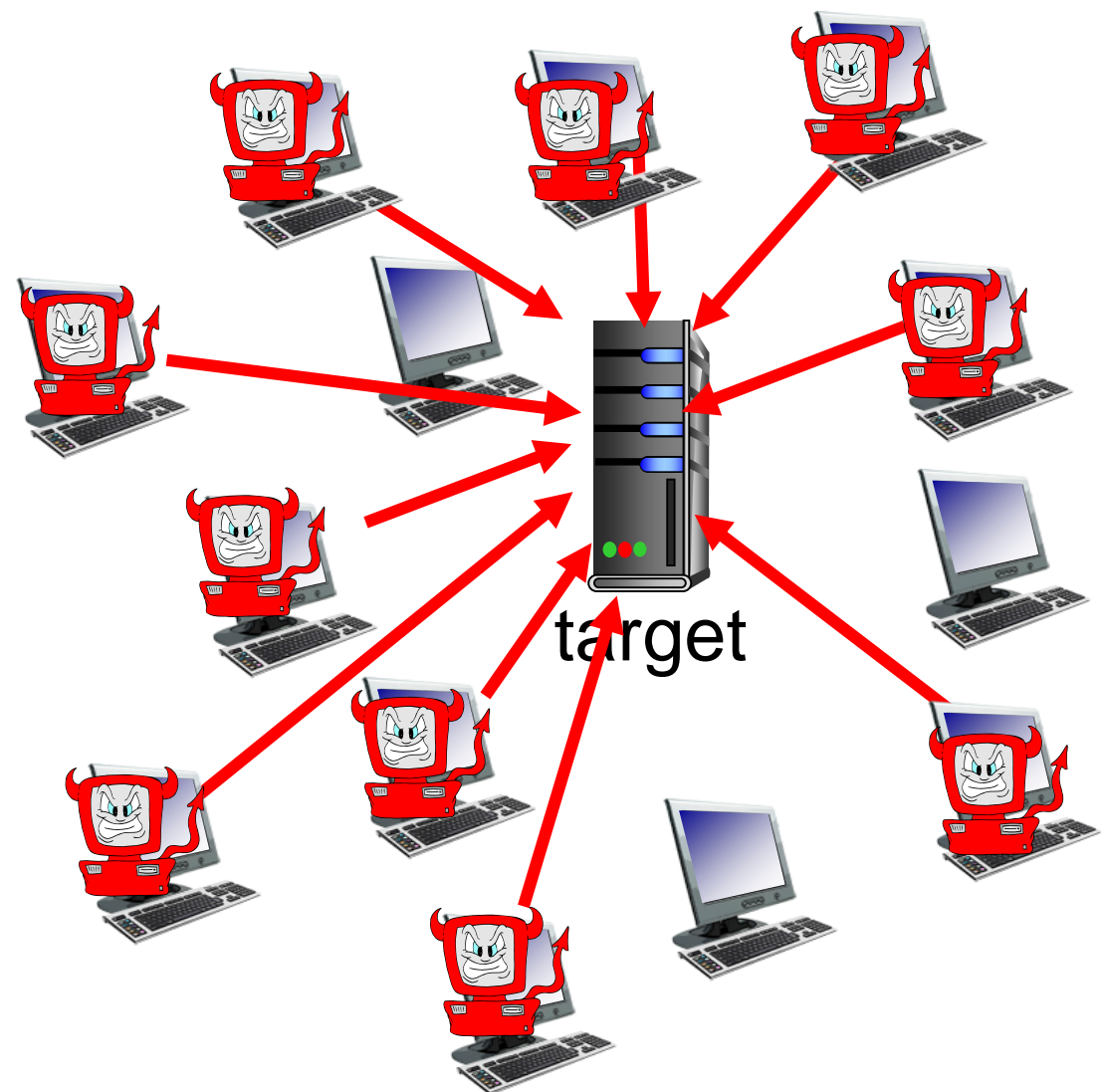
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- **malware can get in host from:**
  - *virus*: self-replicating infection by receiving/executing object (e.g., e-mail attachment)
  - *worm*: self-replicating infection by passively receiving object that gets itself executed
- **spyware malware can record keystrokes, web sites visited, upload info to collection site**
- **infected host can be enrolled in botnet, used for spam. DDoS attacks**

# Bad guys: attack server, network infrastructure

**Denial of Service (DoS):** attackers make resources (server, bandwidth) unavailable to legitimate traffic by overwhelming resource with bogus traffic

1. select target
2. break into hosts around the network (see botnet)
3. send packets to target from compromised hosts

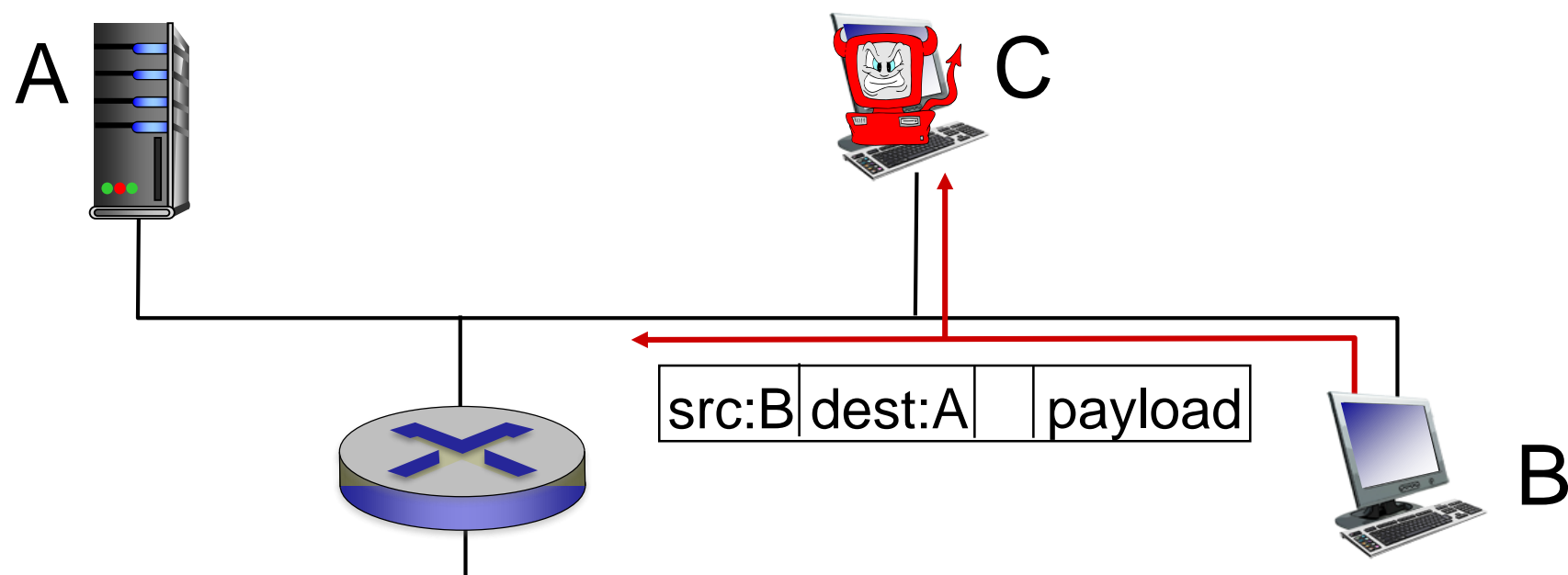




# Bad guys can sniff packets

## ***packet “sniffing”:***

- broadcast media (shared Ethernet, wireless)
- promiscuous network interface reads/records all packets (e.g., including passwords!) passing by

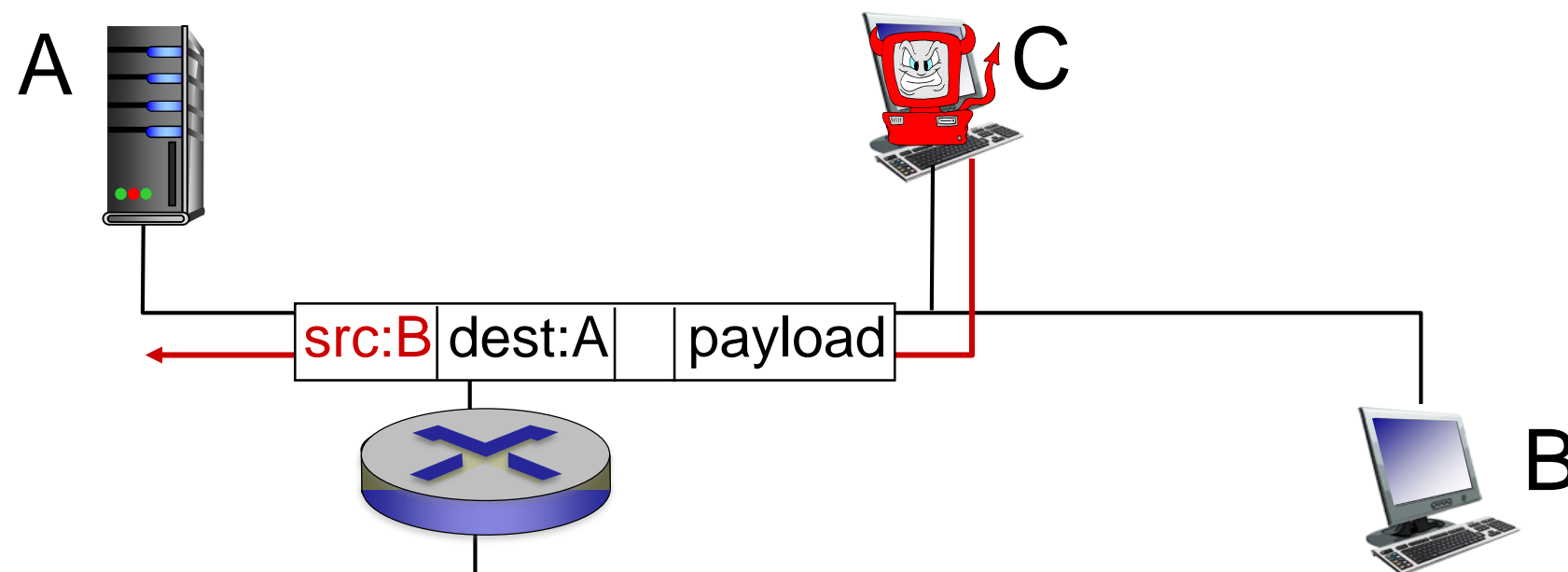


wireshark software used for end-of-chapter  
labs is a (free) packet-sniffer

# Bad guys can use fake addresses

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***IP spoofing:*** send packet with false source address



*... lots more on security (throughout, Chapter 8)*

# Introduction: summary

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***covered a “ton” of material!***

- Internet overview
- what’s a protocol?
- network edge, core, access network
  - packet-switching versus circuit-switching
  - Internet structure
- performance: loss, delay, throughput
- layering, service models
- security
- history

***you now have:***

- context, overview, “feel” of networking
- more depth, detail *to follow!*

