

## CC4

### Chapter II: Data Structures and Address Calculations

1

## Before we start

- **Data Types**
  - Refers to different kinds of data that a variable may hold in a Programming Language
  - Examples:

Data Type	Memory Allocation
Char	1 byte
Int	2 bytes
Float	4 bytes
Double	8 bytes

2

## Data Structures

- Are ways of **storing** and **organizing data** for **efficient utilization**
- Works **together** with **algorithms** to **reduce time and space complexity** while solving **problems**
- Describe how a set of objects are **related** and the set of **operations** that can be applied on the elements
- **Array**
- **Linked-List**
- **Record**
- **Matrix**
- **Stack**
- **Queue**
- **Binary Tree**
- **Binary Search Tree**
- **AVL Tree**
- **Graph**

## Arrays

- Stores **homogenous data** at **contiguous memory addresses** or locations
- Memory space allocated according to **initialized array size**
- Can be of **n-dimensions**, where n is any number from 1 onwards
- Although hidden from us, since it is contiguous, **address calculations** of elements can be determined (later)

# Arrays

- Memory space is not allocated upon declaration, instead upon construction when size is known during instantiation using the keyword *new* in Java.

```
Int [] hardy;  
hardy=new int[5];
```

In single line

```
Int[] hardy=new int[5];
```

Hardy array → 

# Arrays

- One-dimensional Array
- Two-dimensional Arrays
- Multi-dimensional Arrays

## Arrays

- Array Address Calculation Formulae (For Lab Activity#2)

To determine the  $i$ th element of a Single-dimension array:

$$A[i] = \alpha + (i) * \text{esize}$$

where:

$\alpha$  - base or starting address

$i$  - element

esize - element size in bytes

Example: Determine the address of 5<sup>th</sup> element of an integer array A with a starting address of 2000

## Arrays

- Array Address Calculation Formulae (For Lab Activity#2)

To determine the  $i$ th element of a Two-dimension array:

$$A[i][j] = \alpha + [(i) * (UB_2) + (j)] * \text{esize}$$

where:

$UB_2$  - upper bound of the 2<sup>nd</sup> dimension

$\alpha$  - base or starting address

esize - element size in bytes

## Arrays

- Array Address Calculation Formulae (For Lab Activity#2)

To determine the  $i$ th element of a Three-dimension array:

$$A[i][j][k] = \alpha + [(i) * (UB_2) * (UB_3) + (j) * (UB_3) + (k)] * \text{esize}$$

where:

$UB_3$  - upper bound of the 3<sup>rd</sup> dimension

$UB_2$  - upper bound of the 2<sup>nd</sup> dimension

$\alpha$  - base or starting address

esize - element size in bytes

## Arrays

- Exercise:

1. Given  $A[10][3][3][6]$ ,  $\alpha = 2000$ , esize=4 bytes:

- find the formula to represent an element in a 4-dimensional array.
- find the total number of elements
- find the address of  $A[2][2][0][4]$

2. Given  $X[8][3][6][2][3]$ ,  $\alpha = 3000$ , esize=3 bytes:

- find the total number of elements
- find the address of  $X[0][2][5][1][2]$

## Arrays

- **Problem:**

Snow Enterprises gives a fixed salary of P1500 per week to all their newly hired salespeople. They also give an additional 8.5% bonus for each salesperson's Gross Weekly Sales. The company classifies these new employees depending on their Net Weekly Salary as follows:

Employee Classification	Net Weekly Salary
A	P1500 – P1999
B	P2000 – P2499
C	P2500 – P2999
D	P3000 – P3499
E	P3500 – P3999
F	P4000 and above

Write a program that accepts a user's Name and Gross Weekly Sales as inputs. Your program should only have 2 methods outside of the main method:

- Create a *netweekly* method that accepts Gross Weekly Sales as a parameter and returns the Net Weekly Salary as stated in the problem definition.
- Create a *classify* method that accepts Net Weekly Salary as a parameter and returns the Employee Classification as per the table in the problem definition.

**Optional:** If your *classify* method can perform classification using arrays, only 1 loop and 1 condition, additional points will be credited to your next activity.

## Linked Lists

- Consists of 2 values:
  - **Node:** holds **actual element value**
  - **Link:** holds a **reference** to the **next node**, if reference is NULL, end of list
- Memory addresses are allocated **as needed** when **nodes are added**
- Can be **Singly Linked List, Doubly Linked List, Circular Linked List**

## Records

- Stores **heterogenous data**
- Each of these datum is referred to as **fields**
- A combination of records usually create a **database**

13

## Records

Declaration:

```
typedef struct
{
    <data type1> field1;
    <data type2> field2;
    <data type3> field3;
    <data typeN> fieldN;
}RecordType;
```

RecordType

field1
field2
field3
.
.
fieldN

14

# Matrices

- A **matrix** is a rectangular or square grid of numbers arranged into rows and columns.
- Can be generalized to be 2D-arrays you can perform arithmetic calculations on: Addition, Subtraction, Multiplication, and Transposition.

Matrix M

$$M = \begin{bmatrix} 4 & 1 & 5 \\ 3 & 6 & 2 \end{bmatrix}$$

# Matrix Operations

## Matrix Addition

- **RULES:**
  1. Matrices must have the **same dimensions**
  2. Add the **values** of each element at the **same position**

$$\begin{bmatrix} 8 & 4 & 2 \\ 6 & 1 & 5 \end{bmatrix} + \begin{bmatrix} 3 & 10 & 4 \\ 5 & 6 & 1 \end{bmatrix} = \begin{bmatrix} 11 & 14 & 6 \\ 11 & 7 & 6 \end{bmatrix}$$



# Matrix Operations

## Matrix Subtraction

- RULES:**

1. Matrices must have the **same dimensions**
2. Subtract the **values** of each element at the **same position**

$$\begin{bmatrix} 8 & 4 & 2 \\ 6 & 1 & 5 \end{bmatrix} - \begin{bmatrix} 3 & 10 & 4 \\ 5 & 6 & 1 \end{bmatrix} = \begin{bmatrix} 5 & -6 & -2 \\ 1 & -5 & 4 \end{bmatrix}$$

# Matrix Operations

## Matrix Multiplication (with a Constant)

- RULES:**

1. Multiply each element with the constant value

$$2 \times \begin{bmatrix} 4 & 1 \\ 3 & 6 \end{bmatrix} = \begin{bmatrix} 8 & 2 \\ 6 & 12 \end{bmatrix}$$

$$-1 \times \begin{bmatrix} -4 & 1 \\ -3 & 5 \\ 6 & -2 \end{bmatrix} = \begin{bmatrix} 4 & -1 \\ 3 & -5 \\ -6 & 2 \end{bmatrix}$$

# Matrix Operations

## Matrix Multiplication

- RULES:**

- Matrices can **only** be multiplied if the **# of columns** in the **first matrix** is **equal** to the **# of rows** in the **second matrix**.

$$A = \begin{bmatrix} 4 & -1 & 5 \\ 3 & 6 & -2 \end{bmatrix}$$

2 x 3

$$B = \begin{bmatrix} 3 \\ 6 \\ -2 \end{bmatrix}$$

3 x 1

# Matrix Operations

## Matrix Multiplication

- Therefore, multiplying matrices is **not commutative**

$$B = \begin{bmatrix} 3 \\ 6 \\ -2 \end{bmatrix}$$

3 x 1

$$A = \begin{bmatrix} 4 & -1 & 5 \\ 3 & 6 & -2 \end{bmatrix}$$

2 x 3

# Matrix Operations

## Matrix Multiplication

- RULES:**

2. **Sum of the products** of elements are computed from **rows** of **first matrix** and **columns** of **second matrix**.

Let matrix  $A = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$  and  $B = \begin{bmatrix} e & f \\ g & h \end{bmatrix}$

$$A(B) = \begin{bmatrix} a & b \\ c & d \end{bmatrix} \begin{bmatrix} e & f \\ g & h \end{bmatrix} = \begin{bmatrix} ae + bg & af + bh \\ ce + dg & cf + dh \end{bmatrix}$$

# Matrix Operations

## Matrix Multiplication

$$A = \begin{bmatrix} 5 & 1 & 6 \\ 0 & 8 & -2 \end{bmatrix}$$

**2 x 3**

$$B = \begin{bmatrix} 2 \\ 3 \\ 4 \end{bmatrix}$$

**3 x 1**

$$A(B) = \begin{bmatrix} 5(2) + 1(3) + 6(4) \\ 0(2) + 8(3) + (-2)(4) \end{bmatrix} = \begin{bmatrix} 37 \\ 16 \end{bmatrix}$$

- Notice A(B) is now a **2 x 1** matrix.

# Matrix Operations

## Matrix Transposition

- Swap the rows for the columns

$$\begin{bmatrix} 4 & 1 & 5 \\ 3 & 6 & 2 \end{bmatrix}^T = \begin{bmatrix} 4 & 3 \\ 1 & 6 \\ 5 & 2 \end{bmatrix}$$

# Matrix Operations

## Exercises

### Given:

$$A = \begin{bmatrix} 3 & 1 & 5 \\ 6 & 2 & 0 \end{bmatrix} \quad B = \begin{bmatrix} 6 \\ 4 \\ -1 \end{bmatrix} \quad C = \begin{bmatrix} 2 & 4 \\ 3 & 6 \\ -1 & 2 \end{bmatrix} \quad D = \begin{bmatrix} 5 & 2 \\ 3 & 1 \end{bmatrix} \quad E = \begin{bmatrix} 3 & -2 \\ 1 & 4 \end{bmatrix}$$

$$F = \begin{bmatrix} 2 & 1 & 3 \\ 5 & 7 & -2 \end{bmatrix}$$

1.  $A + F$
2.  $E - D$
3.  $C + B$
4.  $C(D)$
5.  $A(F)$
6.  $C^T$
7.  $F^T(E)$