

CC4

Chapter I: Introduction to Algorithms and Data Structures

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Theory of Computation

Early Leanings of Computer Science:

- Use mathematical models to explore **what** computers can do
 - Not restricted to mathematical computations
 - Includes networks, databases, simulations, artificial intelligence, games, security, etc.

Present Leanings:

- Explore **how** to solve computational problems and **how efficient** it can be solved.

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Foundation of Algorithms

- Can a particular task be accomplished by a computing device?
- What is the minimum number of operations for any algorithm to perform a certain function?

Algorithms

- If a computer can solve a particular computational problem, then there must be an **algorithm** for that problem.
- An algorithm is any **well-defined, unambiguous** way/method to solve **all possible instances** of a problem.
- It is a **finite** set of instructions that is **usable by a computer** to accomplish a **particular task**

Algorithms

- Formally an algorithm

“Consists of a set of explicit and unambiguous finite steps, which when carried out for a given set of initial conditions, produce the corresponding output, and terminate in a fixed amount of time.”

- Donald E. Knuth

Guiding Principles of Algorithm Design

Input

- 0 or more quantities which are externally supplied

Output

- At least 1 quantity that is produced

Finiteness

- Must terminate after a finite number of steps; traceability

Definiteness

- Each instruction must be clear and unambiguous

Effectiveness

- Simple but Feasible to implement

Creating Algorithms

- **Requirements:**
 - Information given (**input**) and the results to be produced (**output**) must be **explicitly specified**.
 - **Understand** the problem.
- **Design:**
 - Write an algorithm that will **solve the problem** according to the **requirement**.
- **Analysis:**
 - Trying to come up with **another algorithm** and **compare** it with the previous one.

Creating Algorithms (cont'd)

- **Refinement and Coding:**
 - A **representation** is chosen and a **complete version** of the program is made.
- **Verification:**
 - Consist of 3 aspects:
 - Program
 - Proving
 - Proof by Induction
 - Proof by Contradiction
 - Testing and Debugging

Algorithm Example

Algorithm: Bubble Sort, Ascending

Problem: Given an array of integers, arrange elements from lowest value to highest value.

Input: Array of unsorted integers

Output: Array of sorted integers

Algorithms

Algorithm: Bubble Sort, Ascending

Design:

Let $j=0$, $temp=0$, $array=array$ of integers

- a. Loop through each element pair of the array length
- b. Check the value of the integer at $array[j]$
- c. Check the value of the integer at $array[j+1]$
- d. If $array[j] > array[j+1]$, store the integer in temp**
- e. Place the integer in $array[j+1]$ in $array[j]$**
- f. Place the integer in temp in $array[j+1]$**
- g. Increment j
- h. Repeat steps b to g until length of array, print output

Algorithms

Algorithm: Bubble Sort, Ascending

Refinement and Coding:

```
//Method that accepts an array of integers
static void sort(int arr[]) {
    int n = arr.length;
    //Loop through each element pair given array length n
    for (int i = 0; i < n-1; i++) {
        for (int j = 0; j < n-i-1; j++) {
            //Check and compare values in indices j and j+1
            if (arr[j] > arr[j+1]) { //If the integer in index j is greater
                int temp = arr[j]; //Store it in a temporary variable
                arr[j] = arr[j+1]; //index j is now empty, move the integer in index j+1 here
                arr[j+1] = temp; //index j+1 is now empty, move the integer in temp here
            }
        }
    }
}
```

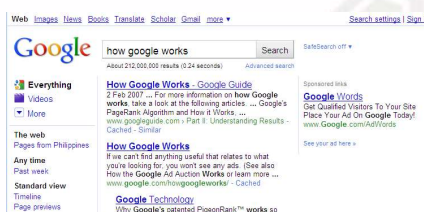
Verification: Testing and debugging



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Importance of Algorithms



The Internet enables people all around the world to quickly access and retrieve (even large amount) information through search algorithms

- Finding good routes from the source host to the destination host uses a clever algorithm
- Search engines that quickly find pages about a particular information, and rank the retrieved results, make use of well-studied algorithms

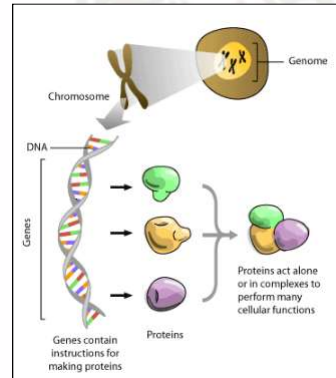


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Importance of Algorithms

- The Human Genome Project involves sophisticated algorithms, enabling it to
 - identify all the (approx. 100,000) genes in human DNA
 - determine the sequences of (approx. 3 billion) chemical base pairs that make up the human DNA
 - store all of the derived information in databases
 - develop tools for data analysis



Importance of Algorithms

Far-reaching impact:

- Business (Stock Market Forecasting, Customer Segmentation, Logistics Set-up,...)
- Hard Science Research (Simulations and visualizations for Protein Folding, Chemical Reactions, Particle Collision,...)
- Security (Encryption/Decryption, Voting Machines, Fraud Detection,...)
- Others (Recommender Systems, File Compression, Cancer Detection, Games,...)

Executing Algorithms

- **Machines** for executing algorithms involve the study of various **fabrications** and **organizations** for algorithms to be carried out.
- Languages for describing algorithms has 2 phases:
 - Language design
 - Translation

Algorithm Analysis

- Know the behavior of the algorithm. This includes the pattern and performance profile of an algorithm; measured in terms of **execution time** and the **amount of space** consumed.
- Computer-related technology is normally associated with **hardware**
 - **Processor** speed
 - **Memory** size

Which is better? Improving hardware or improving algorithms?

Algorithm Analysis

Making an algorithm more efficient is important:

Illustration 1: Suppose a complicated algorithm for n objects requires n^2 operations

- If each op requires 1 sec, how long will the algorithm run when $n = 1000$?
 - $n = 1000 \Rightarrow n^2 = 1\,000\,000$
 - Run-time = $1\,000\,000 \text{ ops} * (1 \text{ sec} / \text{op}) = 1\,000\,000 \text{ sec}$
 - Approx. 11 $\frac{1}{2}$ days!
- If a (twice-as-) fast processor is developed so that each operation requires 0.5 second only, how long will the algorithm run?
- A “faster” algorithm requiring only $3n$ operations will finish in how much time?
 - $n = 1000 \Rightarrow 3n = 3000$
 - Run-time = $3000 \text{ ops} * (1 \text{ sec} / \text{op}) = 3000 \text{ sec}$ or 50 min only!

Analyzing Programs

1. Priori Estimates

- Obtain a **function** which **bounds** the algorithm's **time complexity**. The amount of time a single execution will take or the number of times a statement is executed.
- However, it is **impossible** to **know** the **exact** amount of **time** to execute any command **unless**:
 - a. the machine we are executing is known;
 - b. machine instruction set
 - c. time required by each machine instruction;
 - d. translation of the compiler from source to machine lang.

Analyzing Programs

2. Posteriori Estimates

- How memory space is consumed by the algorithm when it runs
- Example: use of Arrays vs. Linked-lists

Asymptotic Growth and Notation

- Analyzing the **properties** of a growth **function $f(n)$** as **n increases** in size.
- Asymptotic growth is used to analyze/measure algorithm efficiency by:
 - a) Benchmarking
 - b) **Big-Oh Notation**
 - c) Frequency Count

Benchmarking

- Actually run algorithms and get elapsed (running) time.
- Pitfalls:
 - Requires implementation of algorithm
 - Presence of bugs
 - Inefficient implementation
 - Requires exactly the same machine set-up
 - Interruption from software features
 - Results are valid only for the test cases used

Big-Oh Notation

Big-Oh Notation

Common time complexities:

$$O(1) < O(\log n) < O(n) < O(n \log n) < O(n^2) < O(nk) < O(2n) < O(n!)$$

Where n is the input size.

Big-Oh Notation

$O(1)$

- constant; most instructions are executed once or at most only a few times.

$O(\log n)$

- program slightly slower as N grows; normally in programs that solve a big problem by transforming it into a small problem, cutting the size by some constant factor

$O(n)$

- linear; proportional to the size of N

Big-Oh Notation

$O(n \log n)$

- occurs in algorithms that solve a problem by breaking it up into smaller sub-problems, solve them independently and then combining the solution.

$O(n^2)$

- quadratic; can be seen in algorithms that process all pairs of data items.

$O(n^k)$

- polynomial; algorithms that process polynomial of data items.

Big-Oh Notation

$O(2^n)$

- exponential; brute-force solution

$O(n!)$

- factorial.

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Example

Given 2 algorithms performing the same task on N inputs:

	P1	P2
Actual time complexity:	$10n$	$n^2 / 2$
Big-Oh:	$O(n)$	$O(n^2)$

Which is faster and more efficient?

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Example

Analysis:

N	P1	P2
1		
5		
10		
15		
20		
30		

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Example

Analysis:

N	P1	P2
1	10	0.5
5	50	12.5
10	100	50
15	150	112.5
20	200	200
30	300	450

P2 is faster and more efficient for $N \leq 20$;
BUT P1 proves to be better once $N > 20$

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Efficiency Measure Example

$N = 10,000$

Efficiency	Big-O	Iterations	Estimated Time
Logarithmic	$O(\log n)$	14	Microsecond
Linear	$O(n)$	10,000	Seconds
Linear Logarithmic	$O(n \log n)$	140,000	Seconds
Quadratic	$O(n^2)$	$10,000^2$	Minutes
Polynomial	$O(n^k)$	$10,000^k$	Hours
Exponential	$O(c^n)$	$2^{10,000}$	Intractable
Factorial	$O(n!)$	$10,000!$	Intractable

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Frequency Count

Determine the frequency count and the corresponding Big-Oh Notation.

	<u>Frequency Count</u>
1. $K = 500;$	1
for ($j=1; j \leq K; j++$)	$1 + K + 1 + K$
$x = x + 1;$	K
$n = 200;$	1

Substituting actual value for K .

F.C. = 1504

$O(1)$

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