Information Theory: Part II Applications to Stochastic Processes

 We now consider applying information theory to a long sequence of measurements.

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- In so doing, we will be led to two important quantities
 - Entropy Rate: The irreducible randomness of the system.
 - 2. **Excess Entropy:** A measure of the complexity of the sequence.

Context: Consider a long sequence of discrete random variables. These could be:

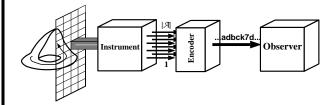
- 1. A long time series of measurements
- 2. A symbolic dynamical system
- 3. A one-dimensional statistical mechanical system

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The Measurement Channel

 Can also picture this long sequence of symbols as resulting from a generalized measurement process:



- On the left is "nature"—some system's state space.
- The act of measurement projects the states down to a lower dimension and discretizes them.
- The measurements may then be encoded (or corrupted by noise).
- They then reach the observer on the right.
- Figure source: Crutchfield, "Knowledge and Meaning ... Chaos and Complexity." In Modeling Complex Systems. L. Lam and H. C. Morris, eds. Springer-Verlag, 1992: 66-10.

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Stochastic Process Notation

- Random variables S_i , $S_i = s \in \mathcal{A}$.
- Infinite sequence of random variables:

$$\overrightarrow{S} = \dots S_{-1} S_0 S_1 S_2 \dots$$

- ullet Block of L consecutive variables: $S^L = S_1, \dots, S_L.$
- $\bullet \Pr(s_i, s_{i+1}, \dots, s_{i+L-1}) = \Pr(s^L)$
- Assume translation invariance or stationarity:

$$\Pr(\,s_{i}, s_{i+1}, \cdots, s_{i+L-1}\,) \,=\, \Pr(\,s_{1}, s_{2}, \cdots, s_{L}\,) \;.$$

- ullet Left half ("past"): $\overset{\leftarrow}{S} \equiv \cdots S_{-3} \, S_{-2} \, S_{-1}$
- ullet Right half ("future"): $\overset{
 ightarrow}{S} \equiv S_0 \, S_1 \, S_2 \cdots$

 $\cdots 11010100101101010101001001010010 \cdots$

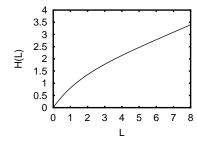
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Entropy Growth

 \bullet Entropy of L-block:

$$H(L) \equiv -\sum_{s^L \in \mathcal{A}^L} \Pr(s^L) \log_2 \Pr(s^L) .$$

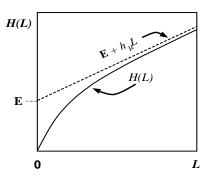
 H(L) = average uncertainty about the outcome of L consecutive variables.



- ullet H(L) increases monotonically and asymptotes to a line
- ullet We can learn a lot from the shape of H(L).

Entropy Rate

• Let's first look at the slope of the line:



- Slope of H(L): $h_{\mu}(L) \equiv H(L) H(L-1)$
- $\bullet\,$ Slope of the line to which H(L) asymptotes is known as the *entropy rate:*

$$h_{\mu} = \lim_{L \to \infty} h_{\mu}(L).$$

- $h_{\mu}(L) = H[S_L|S_1S_1...S_{L-1}]$
- I.e., $h_{\mu}(L)$ is the average uncertainty of the next symbol, given that the previous L symbols have been observed.

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Interpretations of Entropy Rate

- Uncertainty per symbol.
- Irreducible randomness: the randomness that persists even after accounting for correlations over arbitrarily large blocks of variables.
- The randomness that cannot be "explained away".
- Entropy rate is also known as the Entropy Density or the Metric Entropy.
- $h_{\mu} =$ Lyapunov exponent for many classes of 1D maps.
- The entropy rate may also be written: $h_{\mu} \, = \, \lim_{L \to \infty} \frac{H(L)}{L} \; .$
- ullet h_{μ} is equivalent to thermodynamic entropy.
- These limits exist for all stationary processes.

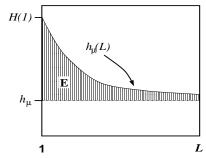
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How does $h_u(L)$ approach h_u ?

 $\bullet \,$ For finite L , $h_{\mu}(L) \geq h_{\mu}.$ Thus, the system appears more random than it is.



- We can learn about the complexity of the system by looking at *how* the entropy density converges to h_{μ} .
- The excess entropy captures the nature of the convergence and is defined as the shaded area above:

$$\mathbf{E} \equiv \sum_{L=1}^{\infty} [h_{\mu}(L) - h_{\mu}] .$$

 E is thus the total amount of randomness that is "explained away" by considering larger blocks of variables. SFI CSSS, Beijing China, July 2005: Information Theory, Part II

Excess Entropy: Other expressions and interpretations Mutual information

 \bullet One can show that E is equal to the mutual information between the "past" and the "future":

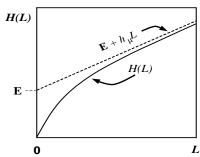
$$\mathbf{E} = I(\overset{\leftarrow}{S};\vec{S}) \equiv \sum_{\overset{\leftarrow}{\{s\}}} \Pr(\overset{\leftrightarrow}{s}) \log_2 \left[\frac{\Pr(\overset{\leftarrow}{S})}{\Pr(\overset{\leftarrow}{s}) \Pr(\vec{s})} \right]$$

- E is thus the amount one half "remembers" about the other, the reduction in uncertainty about the future given knowledge of the past.
- Equivalently, E is the "cost of amnesia:" how much more random the future appears if all historical information is suddenly lost.

Excess Entropy: Other expressions and interpretations Geometric View

• **E** is the *y*-intercept of the straight line to which H(L)

- asymptotes.
- $\mathbf{E} = \lim_{L \to \infty} \left[H(L) h_{\mu} L \right]$.



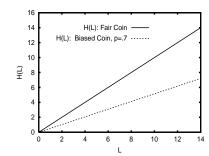
Excess entropy summary:

- Is a structural property of the system measures a feature complementary to entropy.
- Measures memory or spatial structure.
- Lower bound for statistical complexity, minimum amount of information needed for minimal stochastic model of system

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Example I: Fair Coin



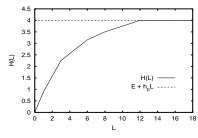
- For fair coin, $h_{\mu} = 1$.
- For the biased coin, $h_{\mu} \approx 0.8831$.
- \bullet For both coins, $\mathbf{E} = 0$.
- Note that two systems with different entropy rates have the same excess entropy.

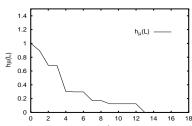
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Example II: Periodic Sequence

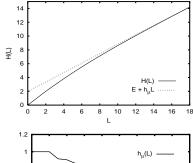


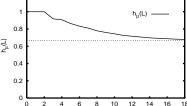


- Sequence: ... 1010111011101110...
- $h_{\mu} pprox 0$; the sequence is perfectly predictable.
- ullet ${f E}=\log_2 16=4$: four bits of phase information
- \bullet For any period-p sequence, $h_{\mu}=0$ and $\mathbf{E}=\log_2 p.$

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Example III: Random, Random, XOR





- Sequence: two random symbols, followed by the XOR of those symbols.
- $h_{\mu} = \frac{2}{3}$; two-thirds of the symbols are unpredictable.
- $\mathbf{E} = \log_2 4 = 2$: two bits of phase information.
- For many more examples, see Crutchfield and Feldman, Chaos, 15: 25-54, 2003.

automata, dynamical systems]

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Excess Entropy: Notes on Terminology

All of the following terms refer to essentially the same quantity.

- Excess Entropy: Crutchfield, Packard, Feldman
- Stored Information: Shaw
- Effective Measure Complexity: Grassberger, Lindgren, Nordahl
- Reduced (Rényi) Information: Szépfalusy, Györgyi, Csordás
- Complexity: Li, Arnold
- Predictive Information: Nemenman, Bialek, Tishby

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Csordás and Szépfalusy, Phys. Rev. A, 39:4767-4777. 1989.
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Shaw, "The Dripping Faucet ...," Aerial Press, 1984. [A

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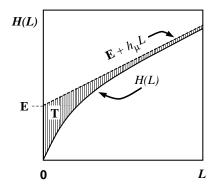
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- Bialek, et al, Neur. Comp., 13:2409-2463. 2001. [Long-range 1D Ising models, machine learning]

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Transient Information T

- $\mathbf{T} \equiv \sum_{L=1}^{\infty} [\mathbf{E} + h_{\mu}L H(L)].$
- T is related to the total uncertainty experienced while synchronizing to a process.



- ullet The shaded area is the transient information T.
- T measures how difficult it is to synchronize to a sequence.

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Some Applications in Agent-Based Modeling Settings

If an agent doesn't have sufficient memory, its
environment will appear more random. In a quantitative
sense, regularities that are missed (as measured by the
excess entropy) are converted into randomness (as
measured by the entropy rate).

Crutchfield and Feldman, Synchronizing to the Environment: Information Theoretic Constraints on Agent Learning. *Advances in Complex Systems*. 4. 251–264. 2001.

2. The average-case difficulty for an agent to synchronize to a periodic environment is measured by the transient information.

Feldman and Crutchfield. Synchronizing to a Periodic Signal: The Transient Information and Synchronization Time of Periodic Sequences. *Advances in Complex Systems*. 7. 329–355. 2004.

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Some Applications in Agent-Based Modeling Settings, continued

- More generally it seems likely that the entropy and mutual information are useful tools for quantifying
 - (a) properties of agents: e.g., how much memory they have
 - (b) the behavior of agents: e.g, how unpredictably they act
 - (c) properties of the environment: e.g., how structured it is

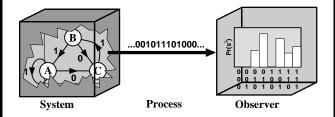
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Estimating Probabilities

• ${\bf E}$ and h_μ can be estimated empirically by observing a process.



• One simply forms histograms of occurrences of particular sequences and uses these to estimate $\Pr(s^L)$, from which $\mathbf E$ and h_μ may be readily calculated.

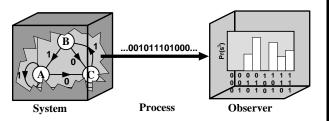
For more sophisticated and accurate ways of inferring h_{μ} , see, e.g.,

- Schürmann and Grassberger. Chaos 6:414-427. 1996.
- Nemenman. http://arXiv.org/physics/0207009.
 2002.

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A look ahead

• Note that the observer sees measurement symbols: 0's and 1's.



- It doesn't see inside the "black box" of the system.
- In particular, it doesn't see the internal, hidden states of the system, A, B, and C.
- Is there a way an observer can infer these hidden states?
- What is the meaning of state?

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