


Modular Exponentiation Master Guide

Problem Statement


Given integers x , n , and M , compute $(x^n \bmod M)$ efficiently.

Why Not Just Use `pow()` or x^n ?

Imagine raising $x = 5$ to power $n = 1e18$ . The number becomes *astronomically huge* and will definitely **overflow** even `long long int`. Hence, we need smarter approaches!

Analogy Time: " Folding Paper "

Imagine you're folding a large paper (exponent) repeatedly:

- Every fold ($\div 2$) makes it smaller!
- Use the folds (binary form) to remember *when* to multiply! 

Brute Force Approach (Not Recommended)

```
long long int PowMod(long long x, long long n, long long M) {  
    long long int ans = 1;  
    for (long long i = 0; i < n; i++) {  
        ans = (ans * x) % M;  
    }  
    return ans;  
}
```

 Time Complexity: $O(n)$

 Space Complexity: $O(1)$


 **Not feasible for large n** like $1e18$. Too slow!

Efficient Approach: **Binary Exponentiation** (a.k.a. Fast Power)

We exploit the fact that:

$$x^n = x^{(n/2)} * x^{(n/2)} \quad \text{if } n \text{ is even}$$
$$x^n = x * x^{(n/2)} * x^{(n/2)} \quad \text{if } n \text{ is odd}$$

Intuition

- Convert n to **binary** 
- Use each bit to decide whether to multiply
- Reduce n by half at each step



Dry Run Example

```
x = 3, n = 5, M = 100
Binary of 5 = 101
Step-by-step:
- ans = 1
- bit 0 (1): ans = ans * 3 = 3
- bit 1 (0): skip multiply
- bit 2 (1): ans = ans * 81 = 243
Final ans = 243 % 100 = 43
```

Final C++ Code

```
class Solution {
public:
    long long int PowMod(long long int x, long long int n, long long int M) {
        long long int ans = 1;
        while (n > 0) {
            if (n & 1) { // if current bit is 1
                ans = (ans * x) % M;
            }
            x = (x * x) % M; // square base
            n = n >> 1;      // move to next bit
        }
        return ans % M;
    }
};
```

Time & Space Complexity


Metric	Value
 Time	$O(\log n)$
 Space	$O(1)$

Recursive Version (Same Logic)




```
long long int PowModRecursive(long long x, long long n, long long M) {
    if (n == 0) return 1;
    long long half = PowModRecursive(x, n / 2, M);
    long long result = (half * half) % M;
    if (n % 2 == 1) result = (result * x) % M;
    return result;
}
```

But watch out for **stack overflow** for huge `n`.

When Do You Use This?

- Cryptography (RSA )
- Hashing
- Competitive Programming
- Large exponent calculations without overflow

Summary Table

Method	Time Complexity	Space	Comment
Brute Force	$O(n)$	$O(1)$	 Not for large <code>n</code>
Iterative Binary	$O(\log n)$	$O(1)$	 Best for large inputs
Recursive Binary	$O(\log n)$	$O(\log n)$	 Risk of stack overflow

Tags

modular exponentiation, binary exponentiation, fast power, math, modulo arithmetic, competitive programming