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# Securing and Integrating the IoT with a Smart Home Router

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# 1 Introduction

*“The most profound technologies are those that disappear. They weave themselves into the fabric of everyday life until they are indistinguishable from it.”*

– Mark Weiser, *The Computer for the Twenty-First Century*, 1991

The modern home is becoming increasingly filled with a variety of *connected* devices (laptops, tablets, phones, set-top boxes etc.), providing a myriad of different services to users within the home. On top of this, with the advent use of smart phones and introduction of wearable devices, we too are starting to carry around our own personal network of devices everywhere we go, brushing past many others in our daily lives at home, work and on the street. Although all connected to the Internet, these devices are often encapsulated within their own environment and ecosystem, unable to interconnect, creating a fractured and often complex user experience.

Making matters more interesting, the Internet of Things paradigm is once again becoming a field of great interest due to the advent of cheap, low power wireless embedded devices [1]. However, not much consideration has been made for how these Things should be integrated into the existing home network, with many approaches opting to simply bridge the device to the cloud ([6], [8], [19]), with obvious concerns for security, privacy and up-time.

As these devices enter our homes and pockets, bringing with them their own ecosystems, the user is faced with the increasingly difficult burden of managing all of them and the ecosystems [11], [10]. Due to the sheer number and diversity of these devices, many of which will provide overlapping services and functionality, problems arise in how to ensure these devices not only play nicely together but also ensuring the user’s network and information stays secure against new and unanticipated threats.

In order for these multiple layered networks of devices to truly fade away into the fabric of our everyday lives, a platform and relevant protocols need to be engineered to not only support this heterogeneous network securely, but also aid the user in managing both the network and the privacy of their information.

The Homework home router platform was created to these issues. Rather than assume every user is a network administrator, the project investigated the needs and abilities of the average user in order to propose the future of home networking, re-inventing the protocols, models and architectures to truly suit the home environment. This re-invention of the home router allows a user to easily install, manage and use their home network, without the need of a Cisco qualification.

In regards to the Internet of Things development, previous work demonstrated that it was in need a suitable protocol in order to meet the specific needs of a network of Things [16]. Thus, a new protocol was designed and implemented, which could not only run on even the most constrained battery-powered devices (8MHz), but it could also efficiently scale to support hundreds of Things within the same network.

## 2 Statement of Problem

Whilst there have been many attempts, both past and recent, to create the perfect Internet of Things network for the home, many still exhibit several issues. This section will first present a typical user scenario to provide a context for this problem, then discuss several of the issues found in existing work, both our previous work and others', which provide the motivation for this project.

### 2.1 A Scenario

In order to establish the context of the problem, a typical user scenario has been created to demonstrate a typical novice user purchasing a new Thing and attempting to integrate it into their network securely with minimal knowledge and effort.

*Bob buys a new Thing, a motion sensor, from a shop and brings it home. He wants to connect it to his currently existing home network. To do so, Bob turns on the Thing and presses the connect button on both the Thing and his home router. Using his tablet/PC Bob is able to view the newly available Things in the network that he can connect to, in which he can see his new Thing. Bob is able to view various information about the Thing, including its type, manufacturer, functions etc. Bob selects the Thing he wants to connect to the network and within a minute the new thing is now part of the network and he can then customise how he wants to use it in the network.*

### 2.2 The Problem

The Internet of Things protocol created in [16] proved to be a successful proof-of-concept; however, in order for it to be considered for deployment and integration into existing homes, two main issues need to first be addressed.

#### 2.2.1 Security

The initial design of the IoT protocol didn't consider security due to time constraints; however, because such rich and sensitive data can be gathered about the user and their home from Thing's, it becomes extremely important to not only protect the data from unauthorised persons, but also protect the network itself. The IoT protocol needs to be sufficiently secured to prevent eavesdropping of the transmitted data and injection of false events by perpetrators masquerading as sanctioned participants in the network.

Demonstrated in the scenario, a user may want to purchase any kind of compatible Thing from a shop and then want bring it into their home network with minimal effort. The Thing itself would have no prior knowledge of the network and needs a secure mechanism

to enter the network without compromising the rest of the network or exposing its own data to unauthorised users.

### **2.2.2 Integration**

The current implementation exists as a standalone library with several demo applications. Integration of the IoT protocol into a user-friendly platform is necessary to harness the full power of a network of Things. The integrated platform would then be able to discover and connect to available Things, subscribe and log events from the sensors and using user customised rules, use automata to detect if these rules are met and then perform actions by publishing commands to actuators in the network.

## **2.3 Intranet of Things vs Internet of Things**

As described earlier in section 1, many previous deployments of Internet of Things networks have taken a cloud first approach, see [6, 8]. Whilst this yields certain benefits, such as easy external access and integration with other services [4, 9], it also poses several questions regards data security, privacy and up-time. For this project, the focus will be on developing a home first platform, in which all Things communication will be kept local, with no cloud processing involved; thus a more suitable name, the Intranet of Things, will be used.

## **2.4 Project Outcomes**

Outcomes of the project:

- An extended IoT protocol with sufficient security to prevent eavesdropping and un-sanctioned devices.
- An Extended home information platform (Homework), with the IoT controller role implemented, enabling capturing of IoT events from sensors and creation of commands for actuators.
- Use of Homework's automata to implement closed loop control of Things in the home, subscribing to sensor events, processing rules and publishing to actuators to perform actions.

## **2.5 Structure**

The rest of this proposal is structured in the following way. Section 3 surveys previous work on the field of wireless sensor networks security, home networks and Internet of Things

systems and protocols. Section 4 takes what has been learnt from the literature survey and proposes a solution to solve the problem described above. Lastly, section 5 discusses the work plan outlined for the implementation and evaluation part of the project, for the second half of the year.

## 3 Literature Review

This section discusses previous work on Internet of Things protocols, WSN security and home networking which provides the motivation and possible aid in solving the previously mentioned issues.

### 3.1 State-of-the-art IoT Protocols

Over the last 20 years, the Internet of Things has developed and evolved continuously, from 20 years ago where it was the concept of using barcodes, and later RFID[5], to track items in a warehouse, to today where our home appliances are filled with “smart” electronics, able to sense, react and notify the user of events, such as the washing machine finishing a load. Whilst the Internet of Things has been around for quite some time, under various guises, it’s yet to become commercially successful and see wide-spread deployment; many believe that 2013 is the year that the Internet of Things will truly take off and become ubiquitous in our daily lives[1].

The rest of this section will discuss the various different attempts made to create Internet of Things networks in the home.

#### 3.1.1 Smart Things

Launched in 2012, the SmartThings platform aims to allow the user to turn any ordinary object in the home into a “Smart” object by giving it the ability to connect to the Internet. The platform consists of a central hub connected to the Internet, containing a low-power Zigbee radio, from which it connects to an array of SmartThings accessories as well as many pre-existing third party Zigbee devices. Some of these accessories include motion detectors, moisture sensors, vibration detectors, power-plug switches as well as many others. The user can then view and interact with their Smart Things through the cloud service via an Internet-connected PC or a smartphone, allowing them to either control the devices directly or set-up rules/schedules for devices.

As shown in figure 1, Smart Things takes a cloud approach, offloading all processing to the cloud, leaving the hub to just act as a gateway and translator between the Things and

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<sup>1</sup>Image from Smart Things developer support: <https://support.smartthings.com/entries/21603009-Device-Types-Capabilities-Attributes>

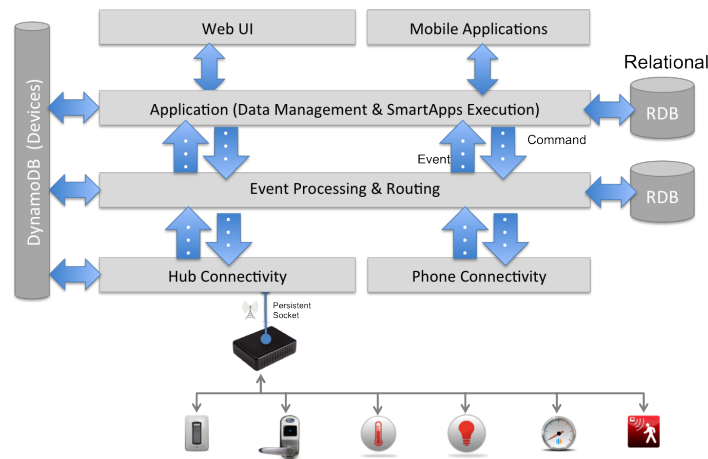


Figure 1: Smart Things - Cloud approach<sup>1</sup>

the cloud. By utilising the cloud in this way, it enables the hub to be a very simple and low-power device, reducing both the purchasing and running costs for the user. It also allows the user to access their network of Things from anywhere, at any time, which would otherwise be difficult to do in a local approach.

However, the cloud approach also has several drawbacks. By offloading the entire processing needs of the network, as show above, the network of Things becomes vulnerable to failure if either the network fails or is unresponsive, or if the cloud service fails or goes down for maintenance. Another issue is the privacy of the user's data; how safe is the data in the cloud and how secure are the services they make available for users to interact with their devices with. There have been many examples of online services that have been attacked, leaking user data and/or shutting down for extended periods of time [7, 2, 3].

Talk about integration of two home networks.

### 3.1.2 CoRE CoAP HIP-DEX

### 3.1.3 KNoT

## 3.2 Home networking

Discuss Homework router and cache db

## 3.3 WSN Security - Symmetric Cryptography

Much research has been carried out to secure WSNs, to ensure data within the network can't be compromised, which could lead to an negative alteration of academic experiment results or cause a leakage sensitive data.



Symmetric key cryptography is a class of cryptographic algorithms which enable two participants to exchange data in secret over a public channel using a shared key, known by both participants. Both encryption and decryption use the same key. One of the key benefits to symmetric algorithms is that it's relatively inexpensive, both in time and space, to encrypt and decrypt data, which makes this ideal for use in a WSN where resources are extremely constrained.

The rest of this section discusses several attempts to secure WSNs using symmetric key cryptography for both the TinyOS and Contiki WSN operating systems, mentioning both the benefits and some apparent drawbacks of each.

### 3.3.1 TinySec

TinySec is a fully functional symmetric security link layer component created for the wireless sensor network operating system, TinyOS. It was the first fully implemented solution for WSNs and was created to address the security worries of running a WSN and transmitting private sensor data in the clear. Unlike conventional security protocol implementations which can afford significant time and space overheads, such as 16-32 bytes for security per packet, WSN typically run on extremely constrained devices with packet sizes of just 30 bytes, making those implementations impossible/extremely expensive to run.

To resolve this, TinySec took a balanced approach making a compromise between the level of security, packet overhead and resource requirements. The end result proved that it's possible to secure a WSN efficiently entirely in software, without the need for additional hardware.

Communication between nodes, not just nodes-to-base-station, in WSNs is often quite important, allowing nodes to not only redirect other's traffic along routes but also consolidate duplicate packets from multiple nodes about the same event, saving the overall network from wasting power receiving and transmitting the extra packets; tinySec chose to engineer in security at the link layer, allowing these mechanisms to perform without alteration. The security goals of TinySec aimed to enable access control, whereby only authorised participants may participate in the network, with unauthorised messages easy to spot and reject; ensure message integrity, so that authorised messages can't be illegally altered by a man-in-the-middle without the receiver noticing; and ensure confidentiality, to ensure information is kept secret from unauthorised eavesdroppers.

The TinySec implementation uses Cipher Block Chaining with an initialisation vector (IV), together these achieve semantic security, therefore ensuring that encrypting the same plain text twice returns a different cipher text each time. So that the receiving end knows how to begin decryption of the data, the IV must be sent in the clear along with the encrypted data. When using an IV, its length needs to be taken into consideration because repeats will occur when the number wraps, causing a security vulnerability. On unconstrained devices an IV is usually 8 or 16 bytes, however due to the packet size limitations of the wireless sensors used, a 8(2 byte counter) byte IV was chosen. In the IV, 6 bytes are made

up of pre-existing fields to conserve space and ensure globally unique IVs in the network e.g. to nodes send the same data event and both happen to have the same counter value, but differ in source (src), so the IV is different, therefore preserving the security.

For ensuring authenticity and integrity of messages, TinySec uses Cipher Block Chaining Message Authentication Codes (CBC-MAC) of 4 bytes in length. Similar to a CRC, CBC-MAC runs over the data and produces a 4 byte MAC which is appended to the packet. If a message was to be altered, the attacker has a 1 in  $2^{32}$  of blindly forging a valid MAC. In a WSN with a limited send rate of 19.2Kb/s it would take over 20 months to send enough packets to possibly succeed in forging a MAC. In the case of attack, a receive heuristic could be used to detect multiple failed MAC transmissions at a nearby node, triggering an alert to the rest of the network.

Whilst TinySec can secure a WSN against eavesdropping and forged messages, it has two significant drawbacks. Firstly, in regards to key distribution, encryption and authentication keys need to be loaded to the nodes prior to deployment. This can cause issues when the shared keys need to be changed, such as when they are compromised, as all nodes in the network will need the new key. This can be especially difficult post-deployment, simply due to the number of nodes and often embedded and/or difficult to reach locations. Secondly, if a node in the network is compromised, because the authentication key is a network wide shared key, the illegitimate node can pretend to be any other node in the network, making it difficult to protect, never mind counter against.

### 3.3.2 MiniSec

MiniSec was created to tackle several problems apparent in the then current WSN security protocols, TinySec and Zigbee. The pre-cursor to MiniSec, TinySec, received much attention and use due to its power and resource efficient security implementation, but because of limited authentication and no replay prevention the overall security of it was deemed insufficient to protect a WSN. A commercial alternative, Zigbee, exhibits significantly higher security, but does so at the cost of high energy consumption. MiniSec was designed to find the middle-ground between the two, increasing security whilst remaining energy efficient.

In contrast to TinySec, for its encryption mode of operation, MiniSec uses Offset Code Book. Unlike cipher block chaining (CBC) which requires two passes to provide both encryption and authentication, OCB provides both in only one pass over the data. This one pass also performs faster than CBC's two passes and only requires one key for operation, thus making it more appropriate for a constrained device in terms of power and storage.

MiniSec also differs from TinySec by reducing the size of the IV counter sent in a packet, yet managing to increase the size of the IV so that it wraps less often. This is achieved by storing some state about the IV counter locally and only transmitting the last  $n$  bits of it to the receiving node. This also requires some logic on both sides to manage the counter in the event of packet loss larger than the range of values stored in the last  $n$  bits sent, i.e. when  $> 2^n$  are lost. In the paper, the authors chose 3 as  $n$ . Because of this, only 3 extra

bits needed to be sent with a packet, instead of the 2 extra bytes in TinySec. However, this overhead could be removed altogether. The maximum default packet size in TinyOS is 29 bytes, therefore the 3 most significant bits in the packet-size byte aren't used and can instead be used to store the IV counter. Therefore removing the need for the extra byte to store those bits.<sup>2</sup>

Another improvement, is the use of a synchronised counter as a nonce to prevent replay attacks. In every packet a monotonically increasing nonce is encrypted with the payload, when the recipient receives a packet with a nonce lower than its counter, it ignores the packet, therefore replayed packets would be dropped.

### 3.3.3 ContikiSec

[12] Discuss port of security to contiki, comparisons of different encryption algorithms, use of OCB.

## 3.4 Key Distribution Problem

Whilst previous work has demonstrated it is feasible to secure a wireless sensor against eavesdropping, unauthorised participants and replay attacks, the issue of key distribution still remains. In the typical academic scenario, this may be a non-issue due to the expertise of the users. However, as described in the scenario in section 2, the typical home user will only have a very basic knowledge if any at all of how to use and operate the devices, let alone configure and distribute network encryption keys on all of their Things. Therefore, it becomes necessary to investigate another method in which the keys for symmetric cryptography can be distributed, with minimal effort and knowledge required from the user.

### 3.4.1 USB interface connectivity

A simple method of distributing keys could be to force the user to plug the Thing into their PC or router, which would then automatically program the correct key into the device with minimal user interaction. Whilst this seems attractive at first, the main issue with this method is that it requires every Thing to have a USB port, plus the appropriate circuitry and logic to perform the programming. This not only increases the cost for the Things, as they now require additional components, but it can also increase the physical size and complexity of using them.

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<sup>2</sup>The publicly available source code for MiniSec (<https://sparrow.ece.cmu.edu/group/minisec.html>) does not make use of this technique, instead it appends an additional byte, thus reducing the benefit of the proposed reduced state based counter.

### 3.4.2 Message in a Bottle

Tried to solve using small Faraday cages to encapsulate the keying devices...[15]

### 3.4.3 Asymmetric Cryptography

Another possible alternative for distributing the symmetric shared keys, would be to use asymmetric cryptography to establish an expensive albeit short-term secure channel between the two participants over an insecure channel and then exchange the key over this. The following section describes this in more detail, as well as some previous work related to WSN and asymmetric cryptography.

## 3.5 WSN Security - Asymmetric Cryptography

Asymmetric cryptography, also known as public-key cryptography, is a class of cryptographic algorithms which enable two or more participants to establish a secure channel over an insecure channel with no prior knowledge or shared secrets beforehand. Participants generate a pair of keys before starting to communicate, known as their public and private keys. The private key must be kept secret, whereas the public key can be distributed freely.

To securely send data to another participant, both must exchange one another's public keys, this can be done on an insecure channel. The newly received public key can then be used to encrypt a packet to send to the other participant. Once a packet has been encrypted with a participant's public key, only they can decrypt it using their private key and read the secret inside, an eavesdropper would only have the public key which can't be used to decrypt the packet, hence the term asymmetric.

Alternatively, the private key can also be used to sign some data, which other participants can authenticate using the public key. This only assures the recipient that the data was sent by the signee, the only one with the private key; however, because anyone with the public key can view the data, it is therefore not secret.

Whilst asymmetric cryptography seems to solve the problems previously discussed in symmetric cryptography, it also has some issues of its own. Firstly, in terms of size and speed, asymmetric algorithms tend to be several orders of magnitude slower than symmetric algorithms, due to the more complex algorithms used to provide the asymmetry, therefore use of them continuously, especially on a WSN where processing power is severely limited, can reduce the throughput and efficiency of the device. Secondly, whilst it's possible to distribute public keys in the open, a man-in-the-middle attack can compromise the security, in which a third party could intercept the public keys and switch them out for their own; after which, the participants would believe they are communicating securely, but are instead passing traffic through the third party.

The rest of this section will discuss two different implementations of asymmetric cryptography algorithms for WSN, as well as a possible solution to the authentication problem.

### 3.5.1 Certificates and Public Key Infrastructures

By itself, asymmetric cryptography only allows two participants to create a secure link over an unsecured channel, preventing eavesdroppers. It does not allow authentication of either participant, meaning that an active eavesdropper could perform a man-in-the-middle attack.

To counter this, a public key infrastructure can be used to pass the authentication up to a trusted body. This trusted body is known as a certificate authority. The CA's job is to issue public key certificates for agents that want to be trusted by others. An agent can obtain a PKC by presenting the CA with its details and public key, which the CA approves and signs the pair using its own private key; it then returns the signed certificate to the agent along with the CA's public key. Now, when the agent carries out an asymmetric key exchange with another party, it can send its certificate to the other party, which can then verify it using its own copy of the CA's public key. This allows two parties to now authenticate each other's public key, ensuring that an active eavesdropper hasn't performed a man-in-the-middle attack and switched out the keys for their own. However, now both parties must trust the CA and obtain a personal PKC and CA private key from it; if the CA were to be compromised, all trust is lost.

### 3.5.2 TinyPK

TinyPK is an implementation of public-key protocols using the Rivest Shamir Adelman (RSA) algorithm for use on TinyOS WSNs[20]. The purpose of the paper was to demonstrate that it's viable to selectively use public key protocols to securely distribute shared session keys for later use in symmetric cryptography.

In the context the paper presents, there are two entities, a sensor network and a third party that wishes to connect to the network. Also apparent is a public key infrastructure, which contains a trusted certificate authority. This CA assists with authenticating all communicating members. Prior to deployment, public and private key pairs are generated for all members (static keys). The CA signs each member's public key with its (CA) private key and then gives back the signed public key and the its own public key to the respective owner. This enables any member to authenticate a public key received from another member, assuring it that a man-in-the-middle attack can't have occurred.

One of the issues TinyPK tries to solve, is the high time cost associated with performing private key functions on the motes, stating that they could take tens of minutes to perform. To avoid this, it is assumed that the third party that wishes to connect to the WSN has more resources at hand and can compute the private functions within reasonable time bounds. This in effect means that any signing required on the sensor side needs to either

be avoided or be precomputed by the CA using its private key. An example in this context, is a mini certificate used to identify a sensor, containing meta-data about the sensor such as ID, date of construction and type. Because the sensor doesn't have the computational power to sign it, the CA can sign it using the CA's private key, therefore any other device with the public key can read it and trust the CA.

The security protocol is itself a simple challenge-response protocol. First, the third party decides to connect to the sensor network, sends it's CA signed public key and sets a challenge in the form of a nonce to protect against replay attacks. If the receiving sensor is legitimate it will contain the CA's public key which can be used to authenticate the public key, using this the device can then un-sign the challenge nonce. Now that the sensor can prove the third party is legitimate, it can then choose to reply to it using the valid nonce it received along with the session key. This time the packet is encrypted using the third parties received public key, ensuring no-one else can eavesdrop on it. Upon reception the third party verifies the nonce and then stores the session key for future use.

**DISCUSS issue of mote authentication happening after key is distributed..is it safe?**

### 3.5.3 TinyECC

[17] A public key crypto implementation of Elliptic Curve cryptography, supporting higher effective security with smaller keys (than RSA). **DISCUSS efficient implementation, sizes of keys and why it's perfect for WSNs**

## 3.6 Other Works

### 3.6.1 MQTT

### 3.6.2 IETF Work

[13, 19]

## 3.7 Problems with previous work

Typical wireless sensor network security has considered and used symmetric cryptography[14, 18, 12] for securing networks in the wild, however key distribution has been bootstrapped (burned to ROM) prior to deployment. This would not only increase the difficulty of bringing a new Thing into a home network, but would also require each Thing have additional hardware in order to interface with the key distributor i.e. PC.

Alternatively, asymmetric cryptography has also been proved viable on typical WSN platforms thanks to elliptic curve cryptography which provides significant performance benefits of previously existing methods[17]; therefore allowing pairs of nodes in a network to dynamically create shared keys. Whilst there is some improvement in performance, it still considerably slower and consumes far more resources than symmetric cryptography, as shown in the table below.

	Symmetric(MiniSec)	Asymmetric(TinyECC)
ROM size	3KB	15.65KB
RAM size	300B	1.8KB
Encrypt	<30ms/pkt	3.2s/pkt
Decrypt	<30ms/pkt	2.1s/pkt

Table 1: Cryptographic implementation sizes and performances on the TelosB mote

**Discuss issue of segmentation of home network and IoT network** Little has been done to present a fully deployable and dynamic secure WSN that is suitable for the home.

## 4 Proposed Approach

### 4.1 Solving the Problem

### 4.2 Security Architecture

#### 4.2.1 Symmetric Key Cryptography

#### 4.2.2 Asymmetric Key Cryptography

### 4.3 Implementation of IoT Protocol on TinyOS

### 4.4 Integration of IoT with Smart Home Router

## 5 Work Plan

- Extend IoT protocol to include security
- Port Secure IoT protocol on TinyOS
- Create host proxy for Secure IoT Protocol
- Create automata for Homework

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