

Homework 4 - Math 525

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July 2023

P. 1 Consider the simplex method applied to a standard form problem and assume that the rows of the matrix A are linearly independent. For each of the statements that follow, give either a proof or a counterexample.

- a. An iteration of the simplex method may move the feasible solution by a positive distance while leaving the cost unchanged.
- b. A variable that has just left the basis cannot reenter in the very next iteration.
- c. A variable that has just entered the basis cannot leave in the very next iteration.
- d. If there is a nondegenerate optimal basis, then there exists a unique optimal basis.
- e. If x is an optimal solution found by the simplex method, no more than m of its components can be positive, where m is the number of equality constraints.

- a. False. $c'x < c'(x + \theta d)$ given $\theta > 0$ and $d > 0$, values that are always present in any positive direction. Thus the cost cannot be unchanged.
- b. True. When a variable leaves the basis, its corresponding coefficient becomes 0. Since pivoting is done strictly negative values, and the pivoting column corresponds to variables entering the basis, it is not possible for a variable to exit and then enter the next iteration.
- c.
- d. False. We can find a LP problem with multiple, non-degenerate, optimal solutions. Consider the LP problem

$$\begin{aligned} \max \quad & x_1 + x_2 \\ \text{s.t.} \quad & x_1 + x_2 \leq 1 \\ & x \geq 0 \end{aligned}$$

Now every point where $x_1 + x_2 = 1$ and $x \geq 0$ is optimal.

- e. True. By Theorem 2.4(b), $x_i = 0 \forall i \neq B(1), \dots, B(m)$, so there can be only m positive variables.

P. 3 Solve completely (i.e., both Phase I and Phase II) via the simplex method the following problem:

$$\begin{array}{ll}\text{minimize} & 2x_1 + 3x_2 + 3x_3 + x_4 - 2x_5 \\ \text{subject to} & x_1 + 3x_2 + 4x_4 + x_5 = 2 \\ & x_1 + 2x_2 - 3x_4 + x_5 = 2 \\ & -x_1 - 4x_2 + 3x_3 = 1 \\ & x_1, x_2, x_3, x_4, x_5 \geq 0\end{array}$$

We see 4 iterations of Phase I. We construct the first tableau with the introduction of auxiliary variables:

-5	-1	-1	-3	-1	-2	0	0	0
2	1	3	0	4	1	1	0	0
2	1	2	0	-3	1	0	1	0
1	-1	-4	3	0	0	0	0	1

We get 3 as the pivot variable with x_3 entering the basis and x_6 exiting. Performing the necessary operations we get:

-4	-2	-5	0	1	-2	0	1
2	1	3	0	4	1	0	0
2	1	2	0	-3	1	1	0
$\frac{1}{3}$	$-\frac{1}{3}$	$-\frac{4}{3}$	1	0	0	0	$\frac{1}{3}$

Note that we completely removed x_6 from the tableau as it no longer serves any purpose. We now let x_2 enter the basis with x_7 exiting:

$-\frac{2}{3}$	$-\frac{1}{3}$	0	0	$\frac{17}{3}$	$-\frac{1}{3}$	0
$\frac{2}{3}$	$\frac{1}{3}$	1	0	$\frac{4}{3}$	$\frac{1}{3}$	0
$\frac{2}{3}$	$\frac{1}{3}$	0	0	$-\frac{17}{3}$	$\frac{1}{3}$	1
$\frac{11}{9}$	$\frac{1}{9}$	0	1	$\frac{16}{9}$	$\frac{4}{9}$	0

Now we let x_1 enter with x_7 exiting. Thus we get the final tableau:

0	0	0	0	0	0
0	0	1	0	0	0
2	1	0	0	-17	1
1	0	0	1	$\frac{11}{3}$	$\frac{1}{3}$

and the basic feasible solution is $x = (2, 0, 1, 0, 0)$.

Phase II

We copy over the same tableau and calculate the new reduced cost and cost to get :

7	2	3	3	-2	3
2	1	0	0	-17	0
0	0	1	0	7	0
1	0	0	1	$\frac{11}{3}$	$\frac{1}{3}$

We do one more iteration to get:

7	2	$23/7$	3	0	3
2	1		0	0	1
0	0		0	1	0
1	0		1	0	0

We now see that there are no more negative values in c so the optimal solution is $x = (2, 0, 1, 0, 0)$. Note that we can leave some values blank as calculating them does not give any information useful to the problem.

- P. 4** While solving a standard form problem, we arrive at the following tableau, with x_3, x_4 , and x_5 being the basic variables:

-10	δ	-2	0	0	0
4	-1	η	1	0	0
1	α	-4	0	1	0
β	γ	3	0	0	1

The entries $\alpha, \beta, \gamma, \delta, \eta$ in the tableau are unknown parameters. For each one of the following statements, find some parameter values that will make the statement true:

- The current solution is optimal and there are multiple optimal solutions.
- The optimal cost is $-\infty$.
- The current solution is feasible but not optimal.

a.

- b. Note that $\beta \geq 0$ for x to be feasible. We then have to find some u such that no component is positive. This, as seen from section 3.2, ensures that the optimal cost is $-\infty$. Thus, we choose $\delta \leq 0$ and $\alpha, \gamma < 0$. η is free. A possible tableau could be:

-10	-1	-2	0	0	0
4	-1	7	1	0	0
1	-1	-4	0	1	0
1	-1	3	0	0	1

- c. To be feasible $\beta \geq 0$. The rest of the variables are free as long as at least one $\delta, \gamma, \alpha \geq 0$ to avoid an optimal cost of $-\infty$. We will require $\alpha, \beta, \gamma, \delta, \eta > 0$ giving a possible tableau of:

-10	5	-2	0	0	0
4	-1	1	1	0	0
1	1	-4	0	1	0
5	1	3	0	0	1

P. 6 Consider the following linear programming problem with a single constraint:

$$\begin{aligned} & \text{minimize} && \sum_{i=1}^n c_i x_i \\ & \text{subject to} && \sum_{i=1}^n a_i x_i = b \\ & && x_i \geq 0, \quad i = 1, \dots, n. \end{aligned}$$

- a. Derive a simple test for checking the feasibility of this problem. (Hint: Discuss when $b = 0, b > 0$ and $b < 0$)
- b. Assuming that the optimal cost is finite, develop a simple method for obtaining an optimal solution directly.

a. As instructed, we break b into 3 cases:

- * $b = 0$: This is trivial as $x = \mathbf{0}$ is always a solution
- * $b > 0$: We only require some $a_j > 0$ as we can choose x such that $x_i = 0 \quad \forall i \neq j$ and $x_j = \frac{b}{a_j}$.
- * $b < 0$: Similarly, this is feasible if some $a_j < 0$. Again, we can always choose some x such that $x_i = 0 \forall i \neq j$ and $x_j = \frac{b}{a_j}$.

b. Noticing that we only have one constraint, we note that the solution will only have 1 non-zero value. To find this value, consider $\frac{b}{a_i}$. The smallest of these ratios is the optimal cost if we let all other values be 0 and $a = \mathbf{1}$. Because $a \neq \mathbf{1}$ in most cases, we account for this by multiplying c_i to the ratio and considering the smallest. That is, $\frac{b \cdot c_i}{a_i}$. The smallest of this gives an optimal solution. That is, $x_i = \frac{b}{a_j}$ and $x_j = 0 \forall j \neq i$ where i is the smallest of the $\frac{b \cdot c_i}{a_i}$.

P. 8 Consider the following optimization problem (P) : find a vector $x \in \mathbb{R}^n$ that satisfies $Ax = 0$ and $x \geq 0$, and such that the number of positive components of x is maximized. Let (P') be the linear program defined as:

$$\begin{aligned} & \text{maximize} && \sum_{i=1}^n y_i \\ & \text{subject to} && A(z + y) = 0 \\ & && z, y \geq 0 \\ & && y_i \leq 1, \quad i = 1, \dots, n. \end{aligned}$$

Show that (P') can be used to solve (P) . (Hint: You can show that (P) and (P') are equivalent - you must specify how to map a feasible solution of (P) to a feasible solution of (P') with greater or equal cost, and viceversa.)