## **Event-B Course**

## 3. A Mechanical Press Controller

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- 1. Informal presentation of the example
- 2. Presentation of some design patterns
- 3. Writing the requirement document
- 4. Proposing a refinement strategy
- 5. Development of the model using refinements and design patterns

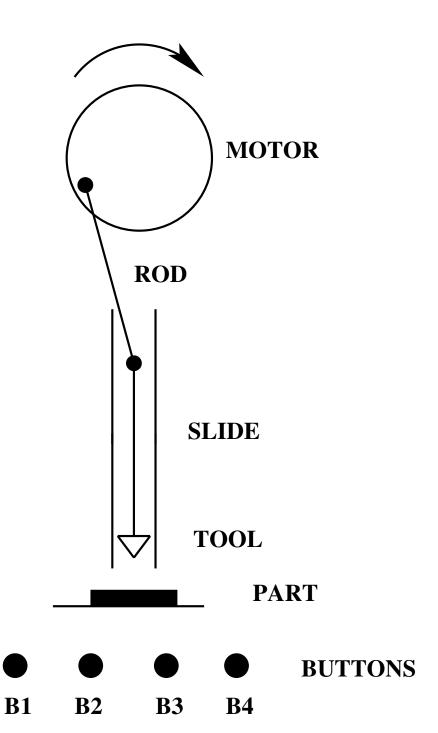
1. Informal Presentation of the Example

- A mechanical press controller

- Adapted from a real system

- The real system is coming from INRST:

Institut National de la Recherche sur la Sécurité du Travail



- A Vertical Slide with a tool at its lower extremity

- An electrical Rotating Motor

- A Rod connecting the motor to the slide.

- A Clutch engaging or disengaging the motor on the rod

- When the clutch is disengaged, the slide stops "immediately"

- Button B1: start motor

- Button B2: stop motor

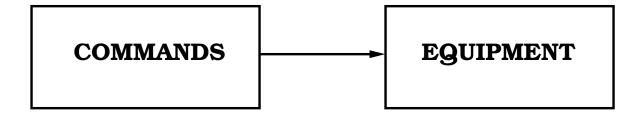
- Button B3: engage clutch

- Button B4: disengage clutch

- Action 1: Change the tool at the lower extremity of the slide

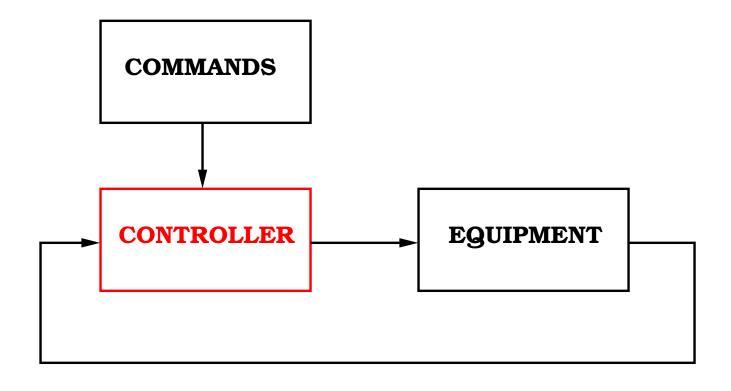
- Action 2: Put a part to be treated under the slide

- Action 3: Remove the part



- 1. start the motor (button B1)
- 2. change the tool (action 1)
- 3. put a part (action 2),
- 4. engage the clutch (button B3): the press now works,
- 5. disengage the clutch (button B4): the press does not work,
- 6. remove the part (action 3),
- 7. repeat zero or more times steps 3 to 6,
- 8. repeat zero or more times steps 2 to 7,
- 9. stop the motor (button B2).

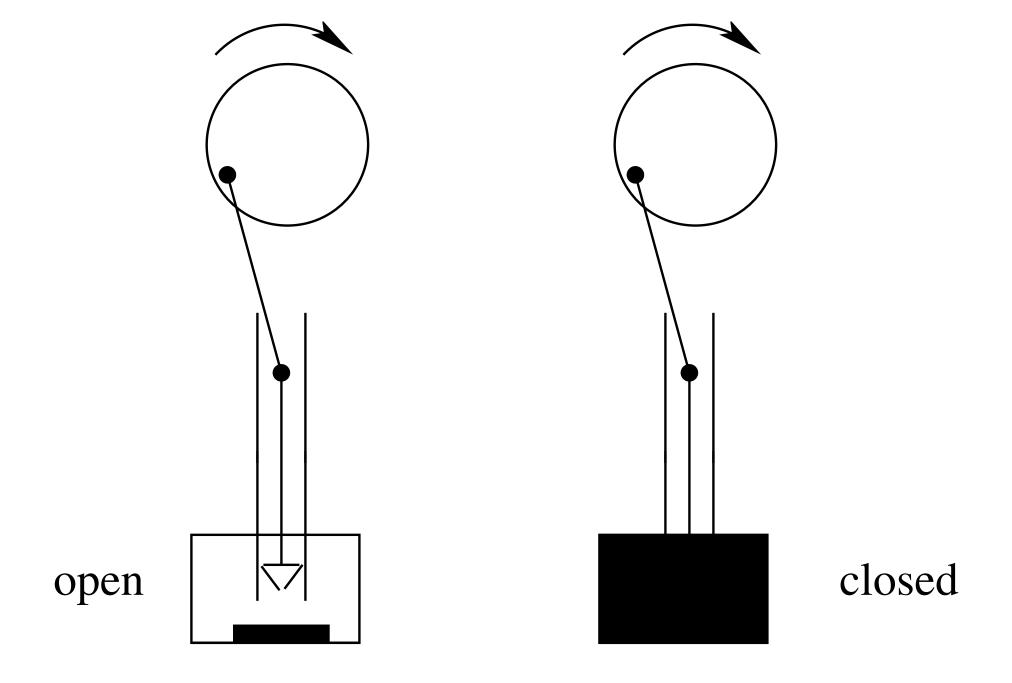
- step 2 (change the tool),
- step 3 (put a part),
- step 6 (remove the part) are all DANGEROUS



- Controlling the way the clutch is engaged or disengaged

- Protection by means of a Front Door

The Front Door



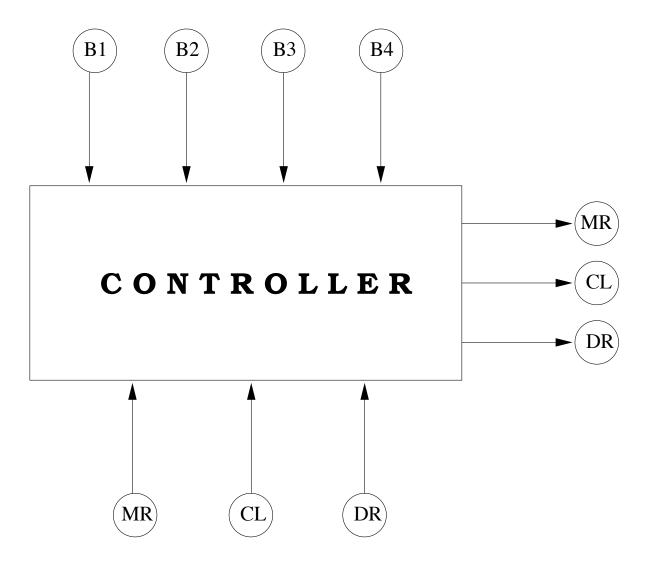
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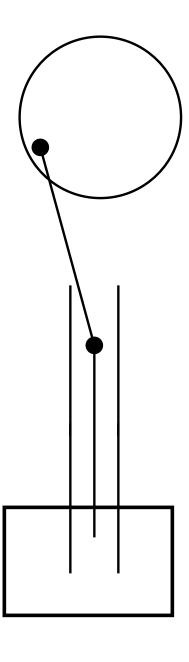
- Initially, the door is open

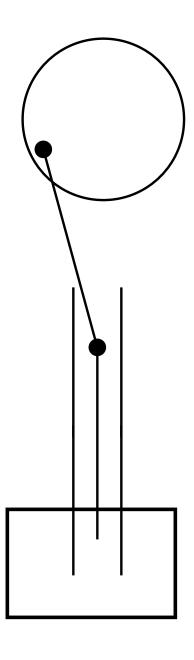
When the user presses button B3 to engage the clutch,
 the door is first closed BEFORE engaging the clutch

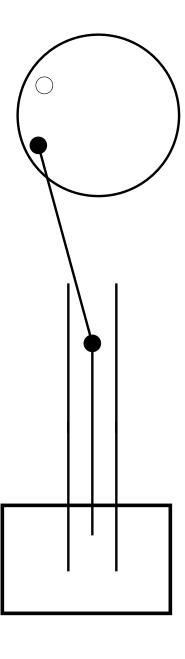
When the user presses button B4 to disengage the clutch,
 the door is opened AFTER disengaging the clutch

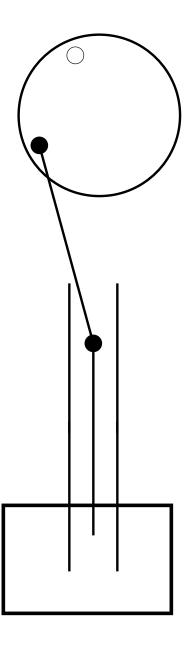
- Notice: The door has no button.

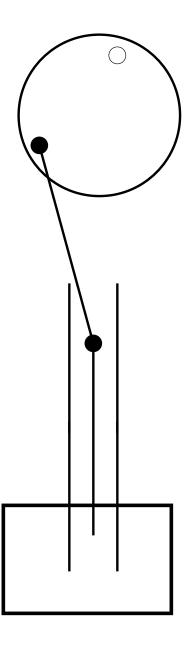


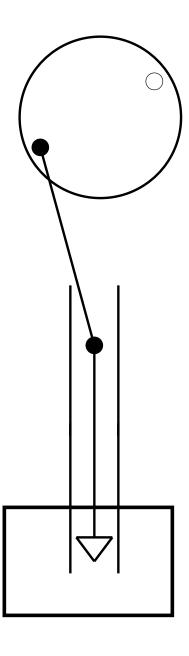


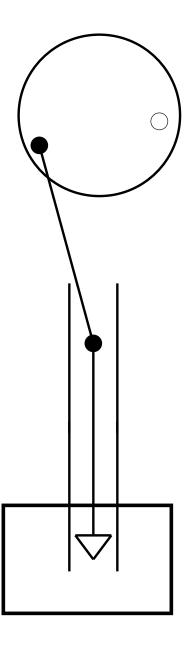


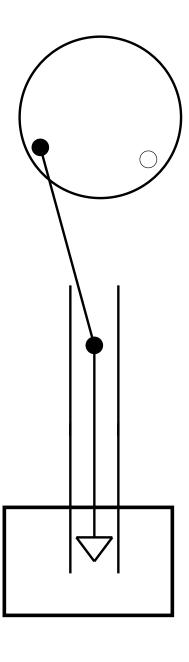


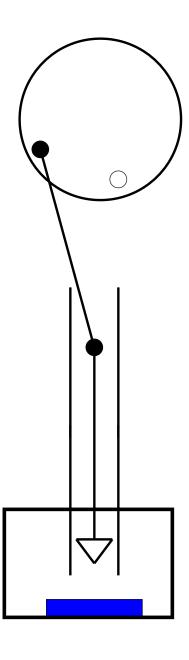


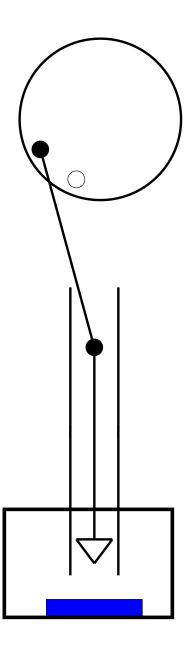


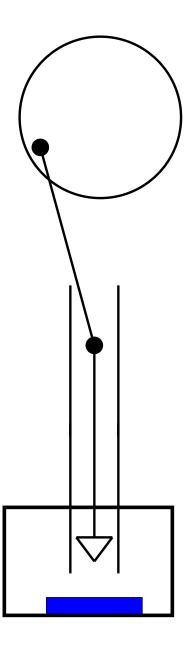


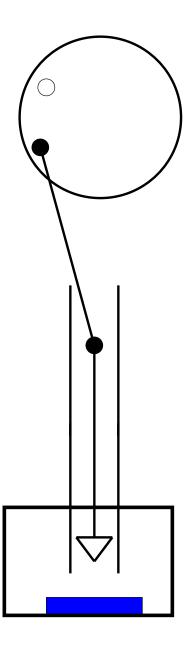


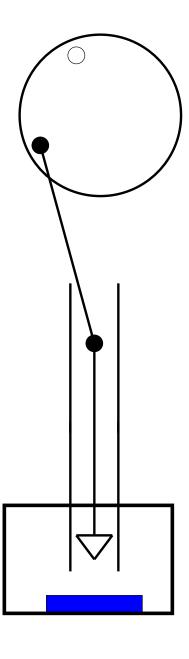


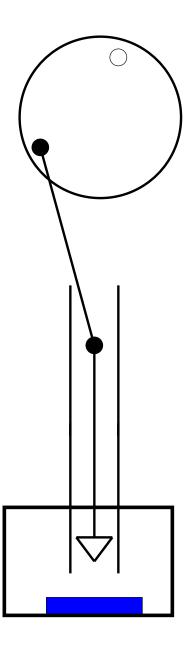


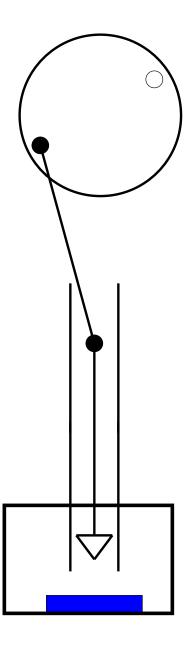


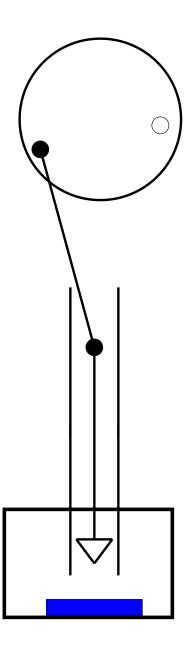


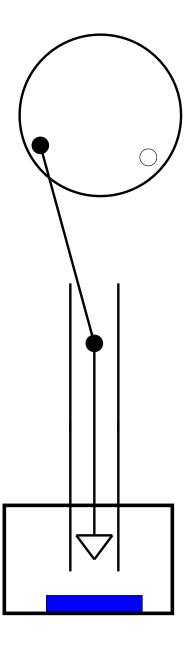


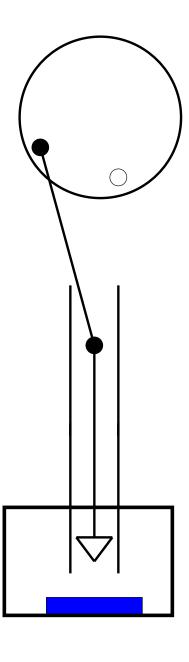


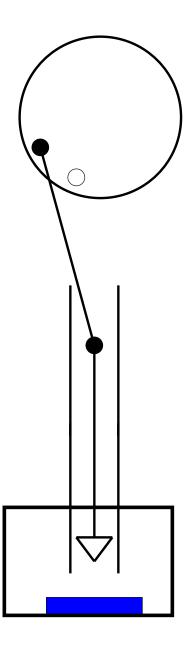


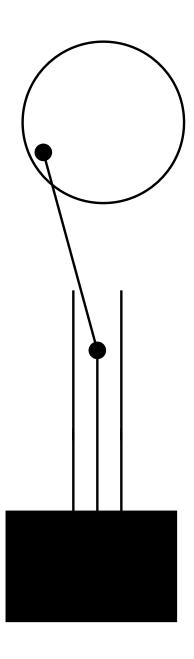


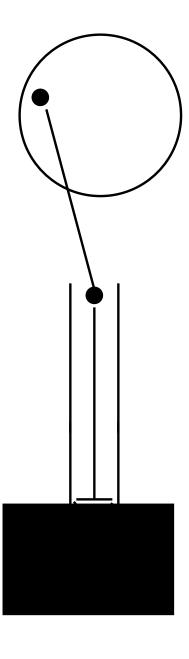


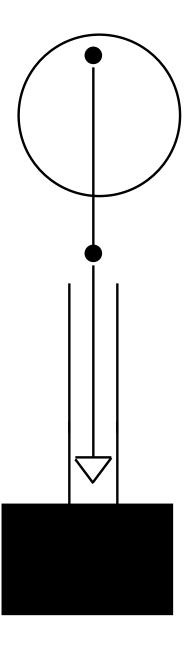


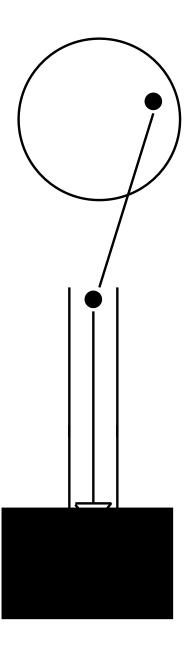


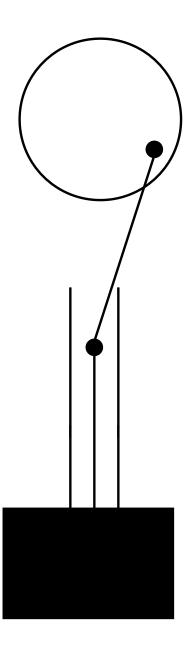


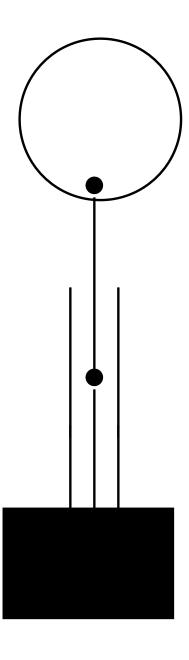


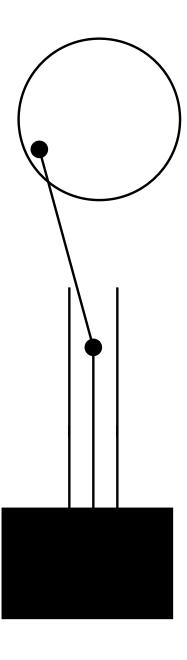


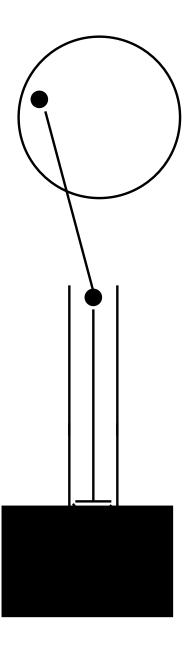


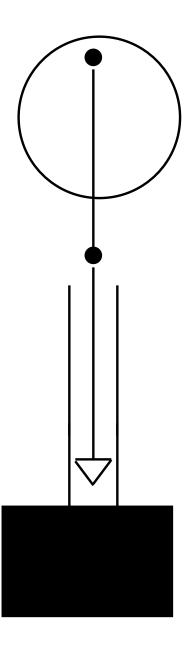


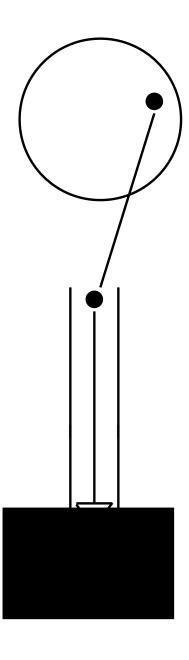


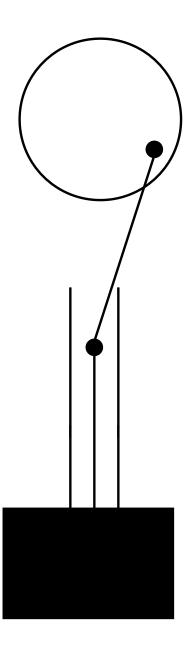


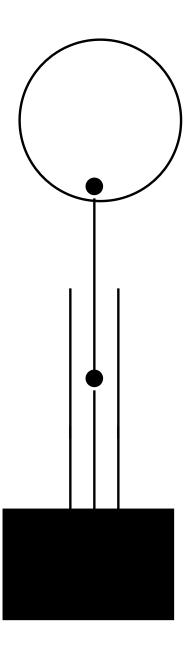


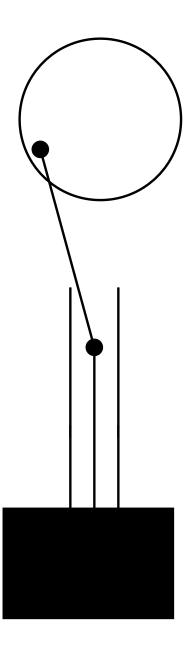


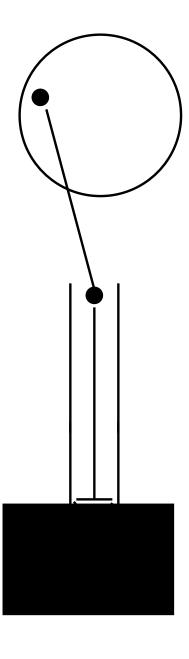


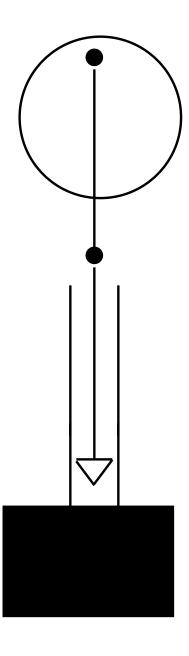


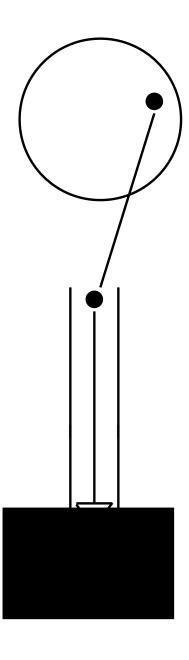


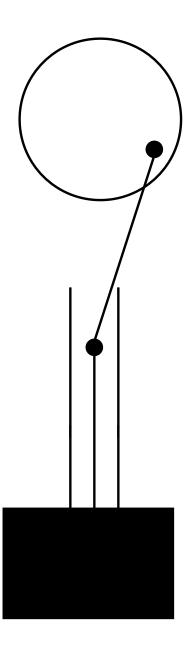


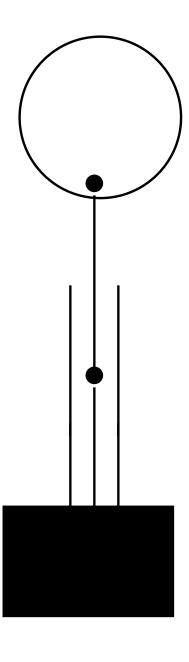


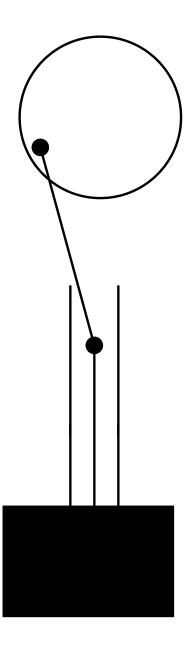


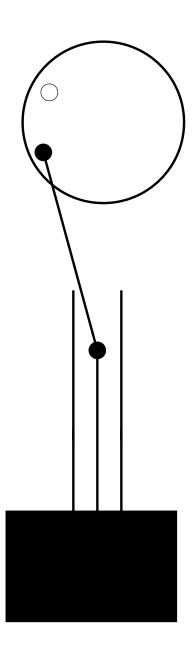


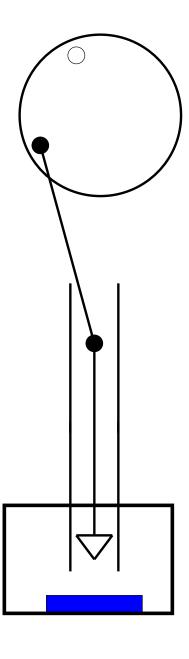


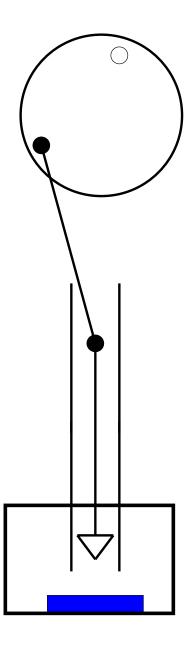


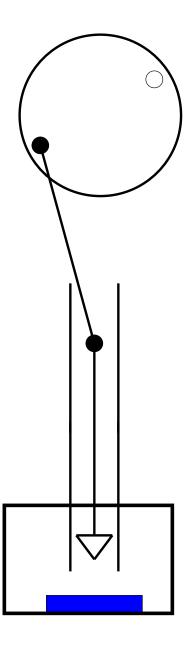


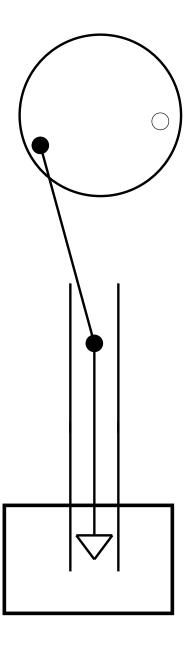


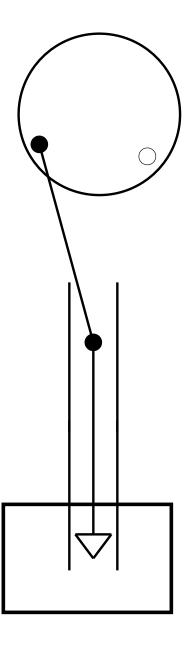


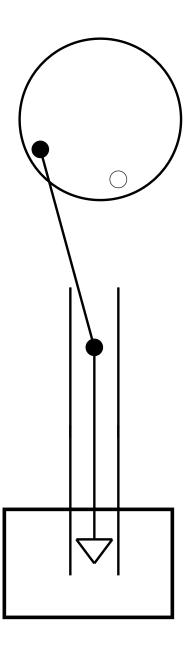


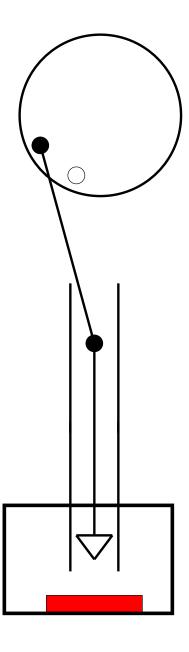


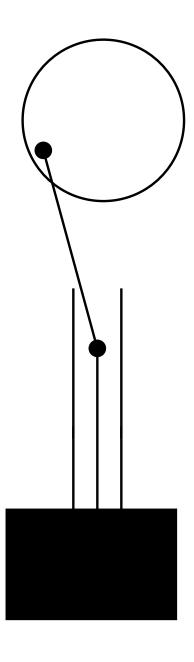


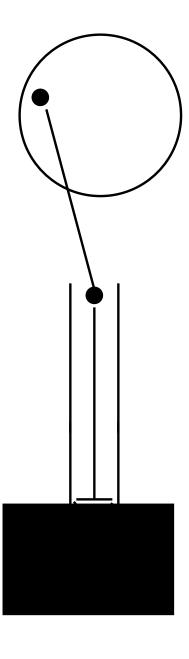


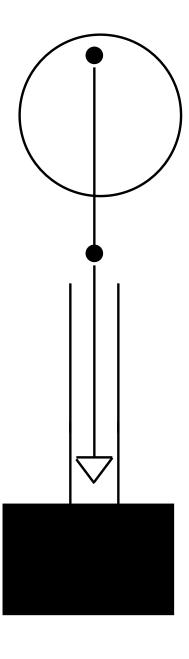


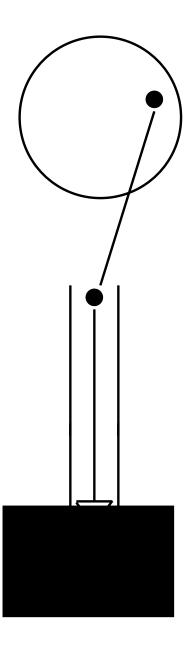


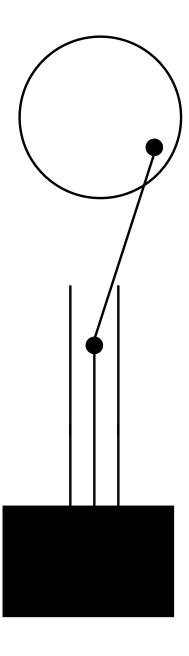


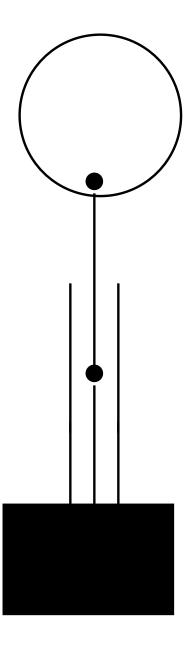


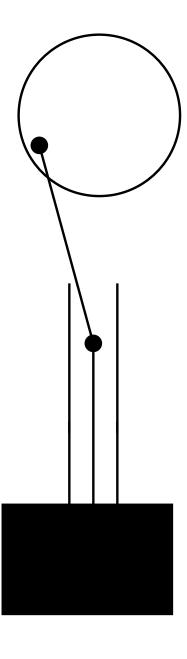


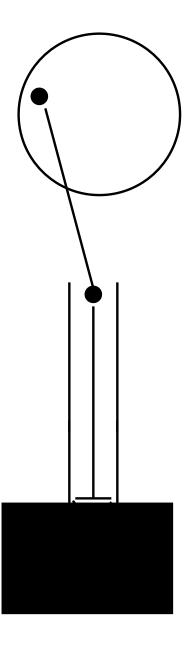


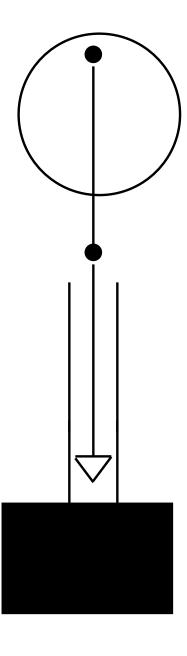


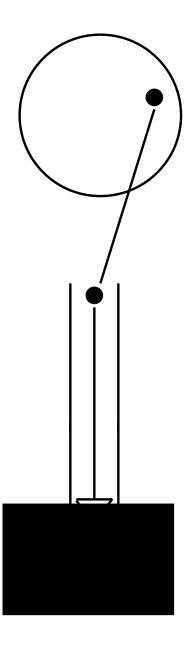


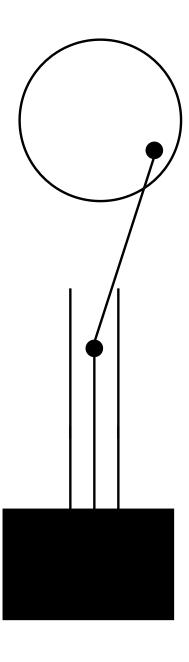


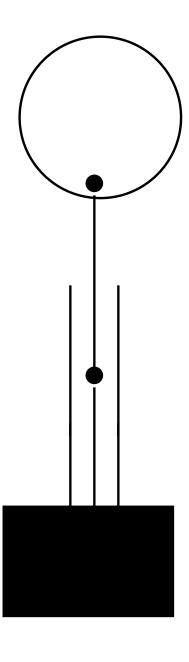


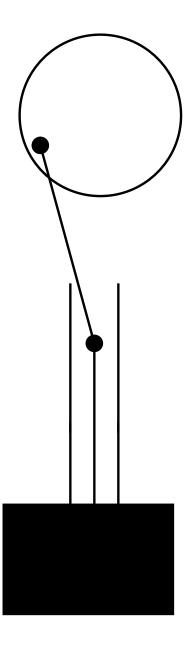


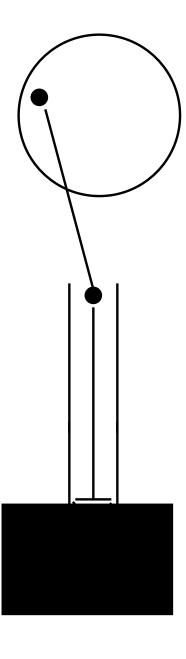


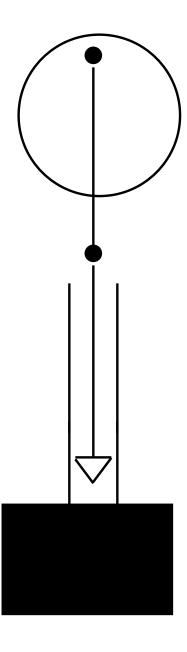


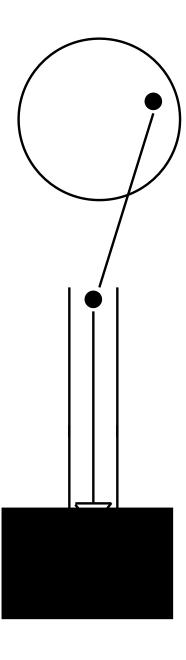


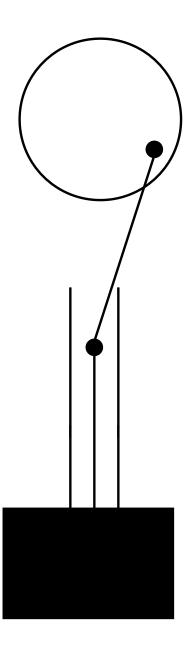


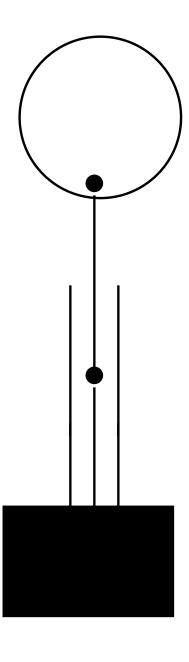


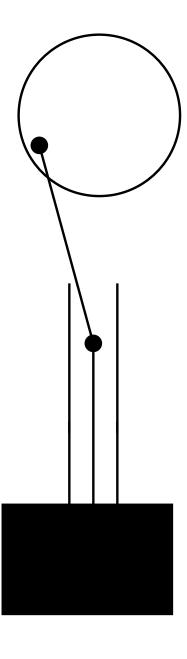


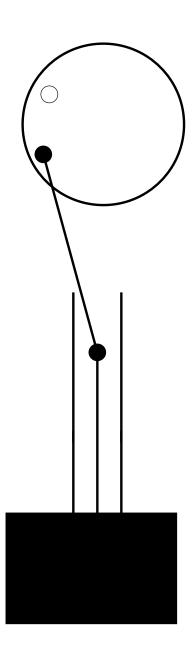


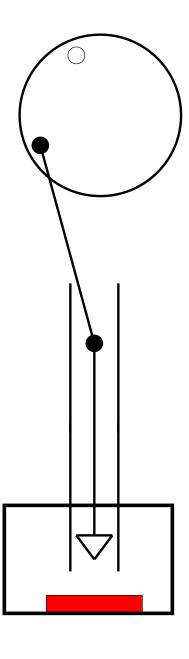


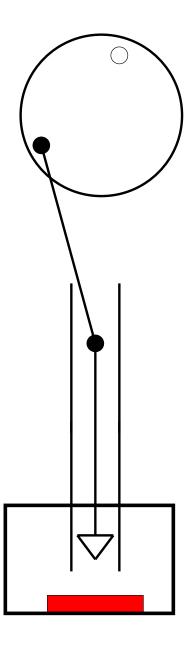


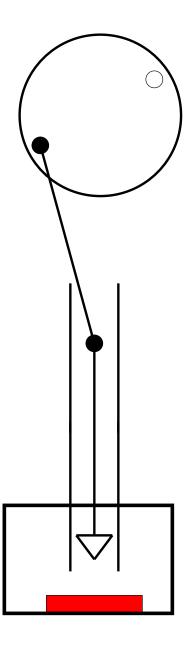


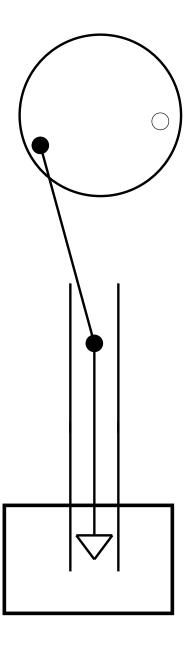


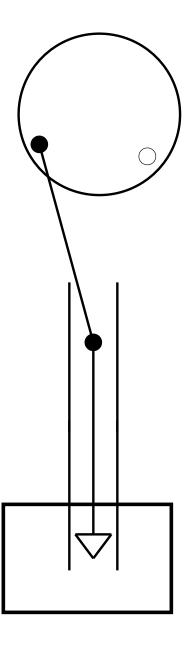


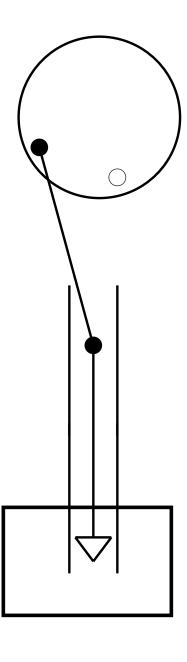


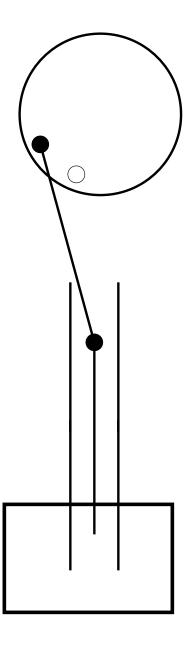


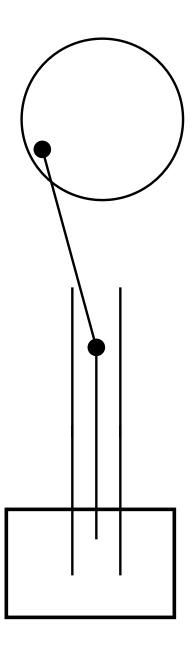


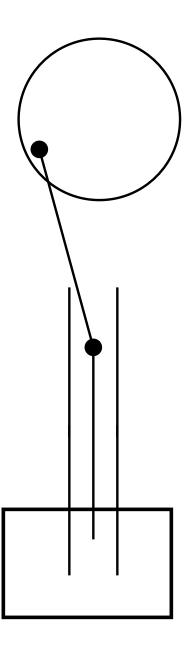


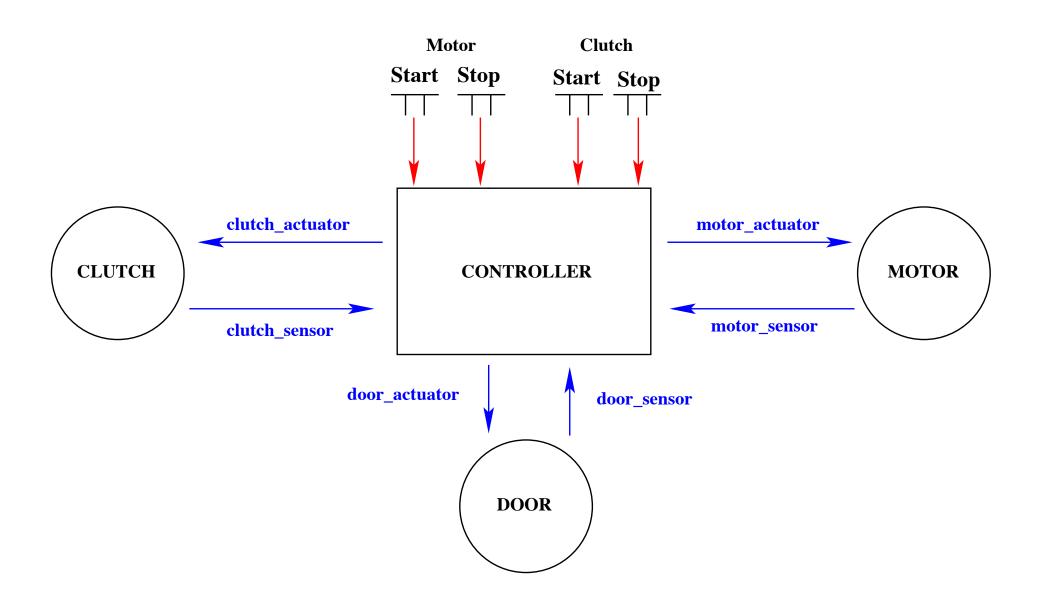












2. Presentation of some Design Patterns

- A number of similar behaviors

- Some complex situations to handle

- A specific action results eventually in having a specific reaction:

- Pushing button B1 results eventually in starting the motor

- Pushing button B4 results eventually in disengaging the clutch

- . . .

- Correlating two pieces of equipment:
  - When the clutch is engaged then the motor must work
  - When the clutch is engaged then the door must be closed

- Making an action dependent of another one:

- Engaging the clutch implies closing the door first

- Disengaging the clutch means opening the door afterwards

- Here is a sequence of events:
  - (1) User pushes button B1 (start motor)
  - (1') User does not remove his finger from button B1
  - (2) Controller sends the starting command to the motor
  - (3) Motor starts and sends feedback to the controller
  - (4) Controller is aware that the motor works
  - (5) User pushes button B2 (stop motor)
  - (6) Controller sends the stop command to the motor
  - (7) Motor stops and sends feedback to the controller
  - (8) Controller is aware that the motor does not work
  - (9) Controller must not send the starting command to the motor

- Here is a sequence of events:
  - (1) User pushes button B1 (start motor)
  - (2) Controller sends the starting command to the motor
  - (3.1) Motor starts and sends feedback to the controller
  - (3.2) User pushes button B2 (stop motor)

- (3.1) and (3.2) may occur simultaneously
- If controller treats (3.1) before (3.2): motor is stopped
- If controller treats (3.2) before (3.1): motor is not stopped

- We want to build systems which are correct by construction
- We want to have more methods for doing so
- "Design pattern" is an Object Oriented concept
- We would like to borrow this concept for doing formal developments
- A preliminary tentative with reactive system developments

- Advantage: systematic developments and also refinement of proofs

- This is an engineering concept

- It can be used outside OO

- The goal of each DP is to solve a certain category of problems

- But the design pattern has to be adapted to the problem at hand

- Is it compatible with formal developments?

- Let's apply this approach to the design of reactive systems

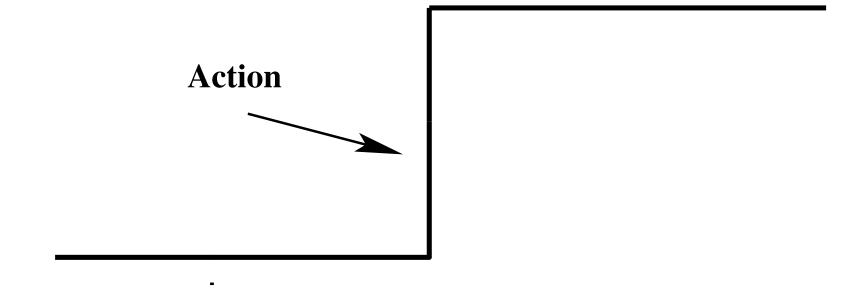
A design pattern isn't a finished design that can be transformed into code

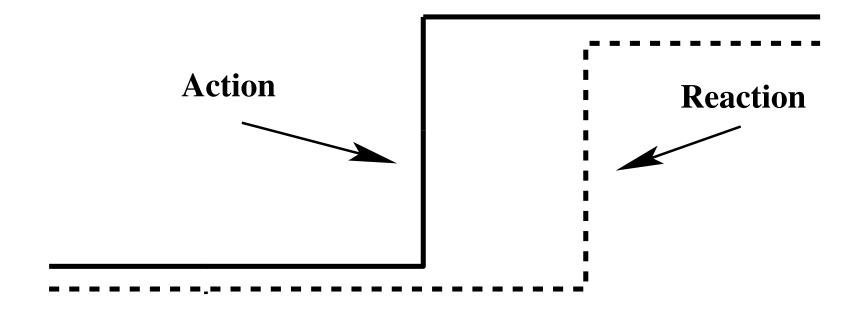
- It is a template for how to solve a problem that can be used in many different situations

Patterns originated as an architectural concept by Christopher Alexander

- "Design Patterns: Elements of Reusable Object-Oriented Software" published in 1994 (Gamma et al)

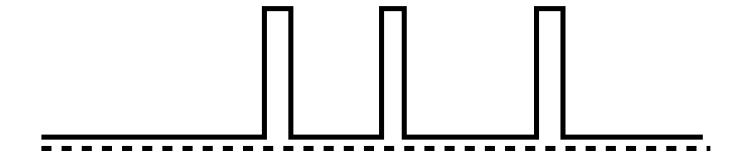
- Design pattern can speed up the development process by providing tested and proven development paradigms
- The documentation for a design pattern should contain enough information about the problem that the pattern addresses, the context in which it is used, and the suggested solution.
- Some feel that the need for patterns results from using computer languages or techniques with insufficient abstraction

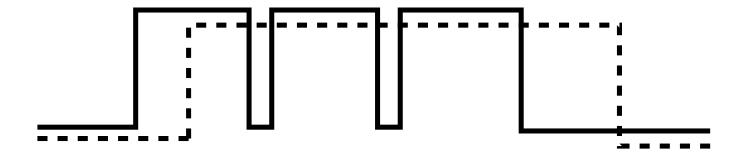




- Sometimes, the reaction has not enough time to react

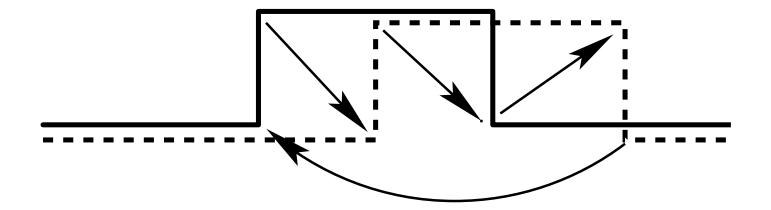
- Because the action moves too quickly





- Sometimes, the reaction always follows the action

- They are both synchronized



- We built first a model of a weak reaction

- The strong reaction will be a refinement of the weak one

variables: a

 $\boldsymbol{r}$ 

pat0\_1:  $a \in \{0,1\}$ 

pat0\_2:  $r \in \{0, 1\}$ 

- a denotes the action
- r denotes the reaction

pat0\_1:  $a \in \{0,1\}$ pat0\_2:  $r \in \{0,1\}$ pat0\_3:  $ca \in \mathbb{N}$ pat0\_4:  $cr \in \mathbb{N}$ 

pat0\_5: cr < ca

- ca and cr denote how many times a and r are set to 1
- pat0\_5 formalizes the weak reaction

```
a_on a=0 when a=0 then a:=1 ca:=ca+1 end
```

$${f a}_{f off}$$
 when  $a=1$  then  $a:=0$  end

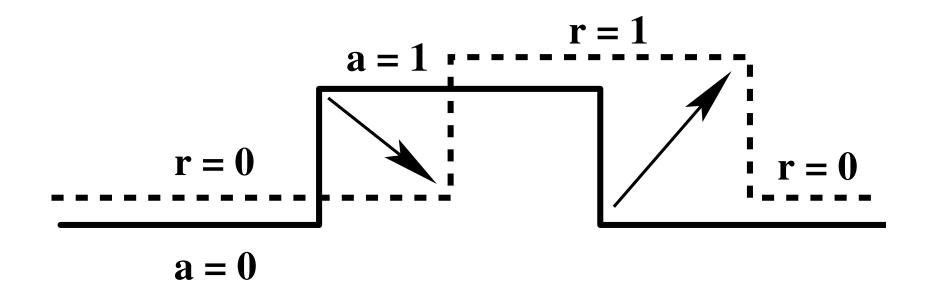
$$a = 1$$

$$a = 0$$

$$\mathbf{a} = \mathbf{0}$$

```
r\_onr=0
a=1
then
r:=1
cr:=cr+1
end
```

$$r\_{
m off}$$
 when  $r=1$   $a=0$  then  $r:=0$  end

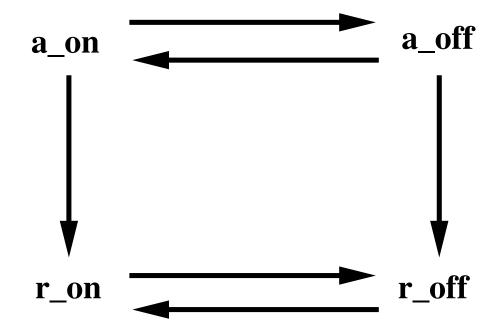


```
a_on a=0 then a:=1 ca:=ca+1 end
```

```
{f a}_{f o} when a=1 then a:=0 end
```

```
r\_{on}
when
r=0
a=1
then
r:=1
cr:=cr+1
end
```

```
r\_{
m off} when r=1 a=0 then r:=0 end
```



 $\begin{array}{cc} \text{variables:} & a, \\ & r, \\ & ca, \\ & cr \end{array}$ 

```
\begin{array}{l} \mathsf{init} \\ a := 0 \\ r := 0 \\ ca := 0 \\ cr := 0 \end{array}
```

```
a_on a=0 then a:=1 ca:=ca+1 end
```

```
{f a}_{f o}{f f} when a=1 then a:=0 end
```

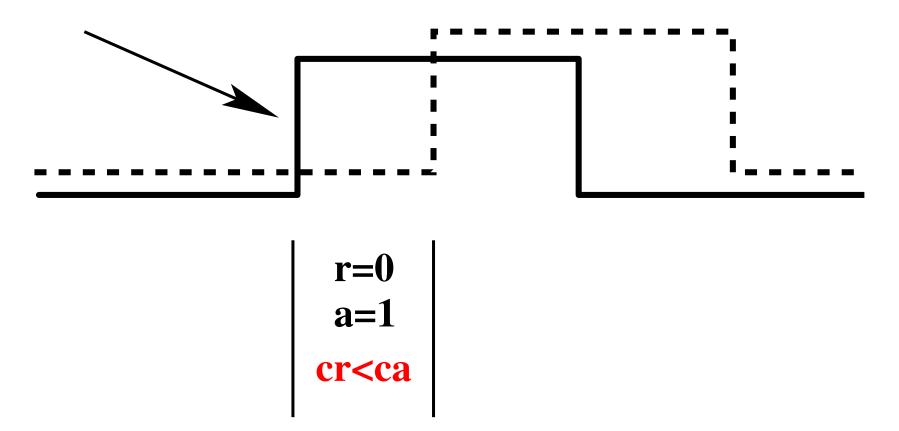
```
r\_{on}
when
r=0
a=1
then
r:=1
cr:=cr+1
end
```

```
r\_{
m off}
{
m when}
r=1
a=0
{
m then}
r:=0
{
m end}
```

## D E M 0 (Showing a Problem and Finding a Solution)

pat0\_6: 
$$r=0 \ \land \ a=1 \ \Rightarrow \ cr < ca$$

## ca is incremented



 $\mathsf{pat0}_{\scriptscriptstyle{-}}\mathsf{1:} \qquad a \ \in \ \{0,1\}$ 

pat0\_2:  $r \in \{0, 1\}$ 

pat0\_3:  $ca \in \mathbb{N}$ 

pat0\_4:  $cr \in \mathbb{N}$ 

pat0\_5:  $cr \leq ca$ 

pat0\_6:  $r = 0 \land a = 1 \Rightarrow cr < ca$ 

The counters have

been removed

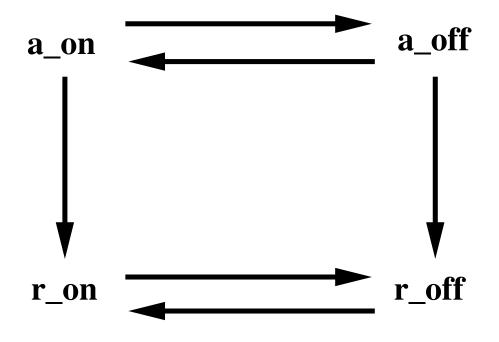
$${f a}_{-}$$
on ${f when} \ a=0 \ {f then} \ a:=1 \ {f end}$ 

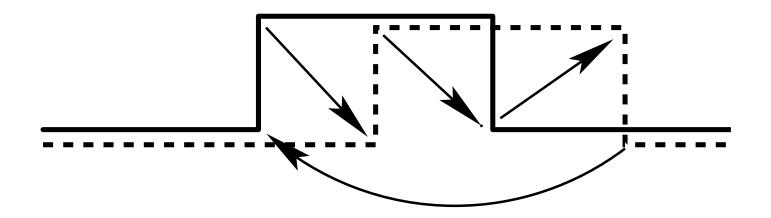
$${f a}_{-}$$
 off when  $a=1$  then  $a:=0$  end

$$egin{aligned} a &:= 0 \ r &:= 0 \end{aligned}$$

$$r\_{
m on}$$
 when  $r=0$   $a=1$  then  $r:=1$  end

$$r\_{
m off}$$
 when  $r=1$   $a=0$  then  $r:=0$  end





- We add the following invariant

pat1\_1: 
$$ca \leq cr + 1$$

- Remember invariant pat0\_5

pat0\_5:  $cr \leq ca$ 

We have thus:  $ca = cr \lor ca = cr + 1$ 

```
pat1_1: ca \leq cr + 1
```

```
a_on when a=0 then a:=1 ca:=ca+1 end
```

```
{f a}_{-}{f off} when a=1 then a:=0 end
```

```
r\_{on}
when
r=0
a=1
then
r:=1
cr:=cr+1
end
```

```
r\_{
m off}
{
m when}
r=1
a=0
{
m then}
r:=0
{
m end}
```

Nothing guarantees that the invariant is preserved

## D E M 0 (Showing Problems and Finding Solutions)

- Putting together these two invariants

pat1\_2: 
$$a=0 \Rightarrow ca=cr$$

pat1\_3: 
$$a=1 \land r=1 \Rightarrow ca=cr$$

- leads to the following

pat1\_4: 
$$a=0 \ \lor \ r=1 \ \Rightarrow \ ca=cr$$

pat0\_5: 
$$cr \leq ca$$

pat0\_6: 
$$a=1 \land r=0 \Rightarrow cr < ca$$

pat1\_1: 
$$ca \leq cr + 1$$

pat1\_4: 
$$a=0 \ \lor \ r=1 \ \Rightarrow \ ca=cr$$

This can be simplified to

pat2\_1: 
$$a=1 \land r=0 \Rightarrow ca=cr+1$$

pat2\_2: 
$$a=0 \ \lor \ r=1 \ \Rightarrow \ ca=cr$$

pat0\_1:  $a \in \{0,1\}$ 

pat0\_2:  $r \in \{0,1\}$ 

pat0\_3:  $ca \in \mathbb{N}$ 

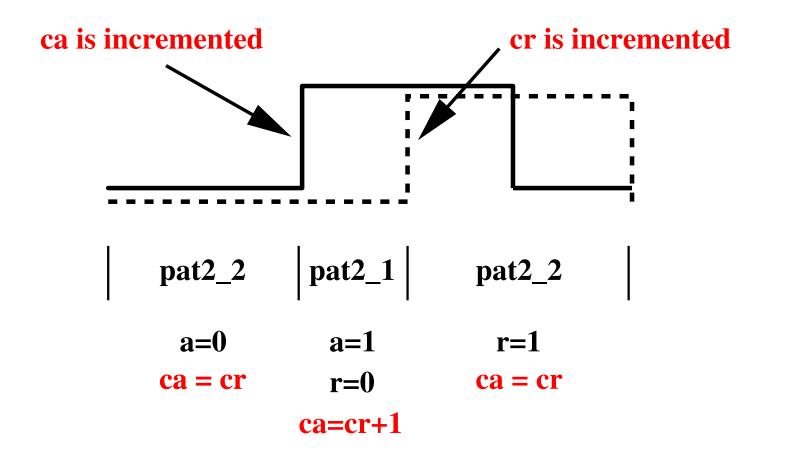
pat0\_4:  $cr \in \mathbb{N}$ 

pat2\_1:  $a=1 \land r=0 \Rightarrow ca=cr+1$ 

pat2\_2:  $a=0 \ \lor \ r=1 \ \Rightarrow \ ca=cr$ 

pat2\_1: 
$$a=1 \land r=0 \Rightarrow ca=cr+1$$

pat2\_2: 
$$a=0 \ \lor \ r=1 \ \Rightarrow \ ca=cr$$



The counters have

been removed

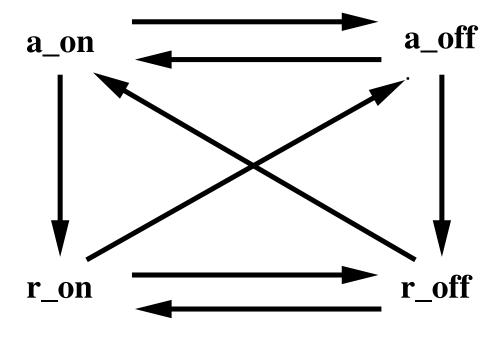
a\_on when 
$$a=0$$
  $r=0$  then  $a:=1$  end

a\_off when 
$$a=1$$
  $r=1$  then  $a:=0$  end

$$\begin{array}{c} \mathsf{init} \\ a := 0 \\ r := 0 \end{array}$$

$$r_{-}$$
on when  $r=0$   $a=1$  then  $r:=1$  end

$$r\_{
m off}$$
 when  $r=1$   $a=0$  then  $r:=0$  end



- Proof failures helped us improving our models
- When an invariant preservation proof fails on an event, there are two solutions:
  - adding a new invariant
  - strengthening the guard
- Modelling considerations helped us choosing one or the other
- At the end, we reached a stable situation (fixpoint)

3. Writing the Requirement Document

The system has got the following pieces of equipment: a Motor, a Clutch, and a Door

EQP<sub>1</sub>

Four Buttons are used to start and stop the motor, and engage and disengage the clutch

EQP<sub>2</sub>

A Controller is supposed to manage this equipment

EQP\_3

Buttons and Controller are weakly synchronized

FUN<sub>1</sub>

Controller are Equipment are strongly synchronized

FUN<sub>2</sub>

When the clutch is engaged, the motor must work

SAF<sub>1</sub>

When the clutch is engaged, the door must be closed

SAF<sub>2</sub>

When the clutch is engaged, the door cannot be closed several times, ONLY ONCE

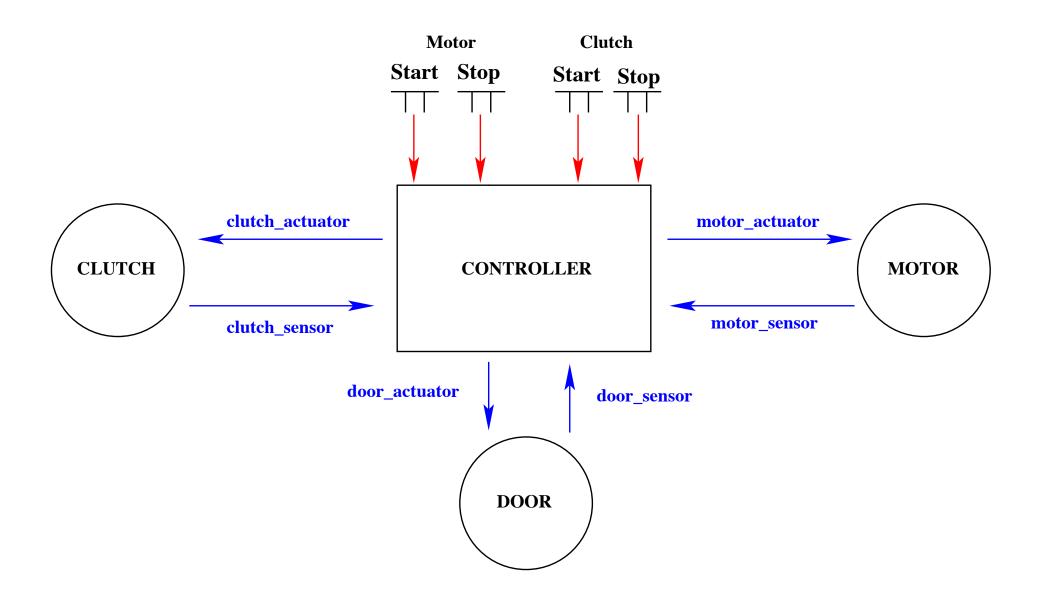
FUN 3

When the door is closed, the clutch cannot be disengaged several times, ONLY ONCE

FUN 4

Opening and closing the door are not independent. It must be synchronized with disengaging and engaging the clutch

FUN<sub>5</sub>



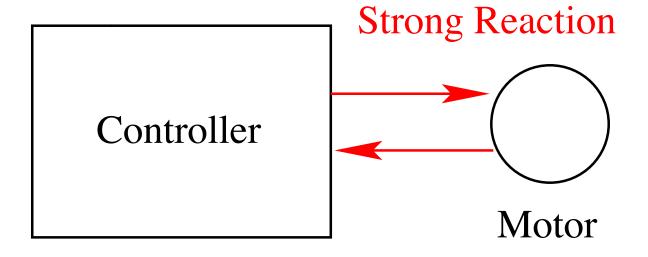
4. Proposing a Refinement Strategy

- Initial model: Connecting the controller to the motor
- 1st refinement: Connecting the motor buttons to the controller
- 2nd refinement: Connecting the controller to the clutch
- 3rd refinement: Constraining the clutch and the motor

- 4th refinement: Connecting the controller to the door
- 5th refinement: Constraining the clutch and the door
- 6th refinement: More constraints between clutch and door

- 7th refinement: Connecting the clutch buttons to the controller

 Development of the Model using Refinements and Design Patterns



Controller are Equipment are strongly synchronized

FUN<sub>2</sub>

The counters have

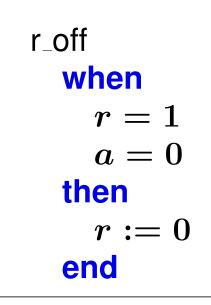
been removed

a\_on when 
$$a=0$$
  $r=0$  then  $a:=1$  end

$${f a}_{-}$$
 off  ${f when}$   $a=1$   $r=1$  then  $a:=0$  end

$$egin{aligned} a &:= 0 \ r &:= 0 \end{aligned}$$

$$r_{-}$$
on $r=0$  $a=1$ then $r:=1$ end



set: STATUS

constants:

 $stopped \\ working$ 

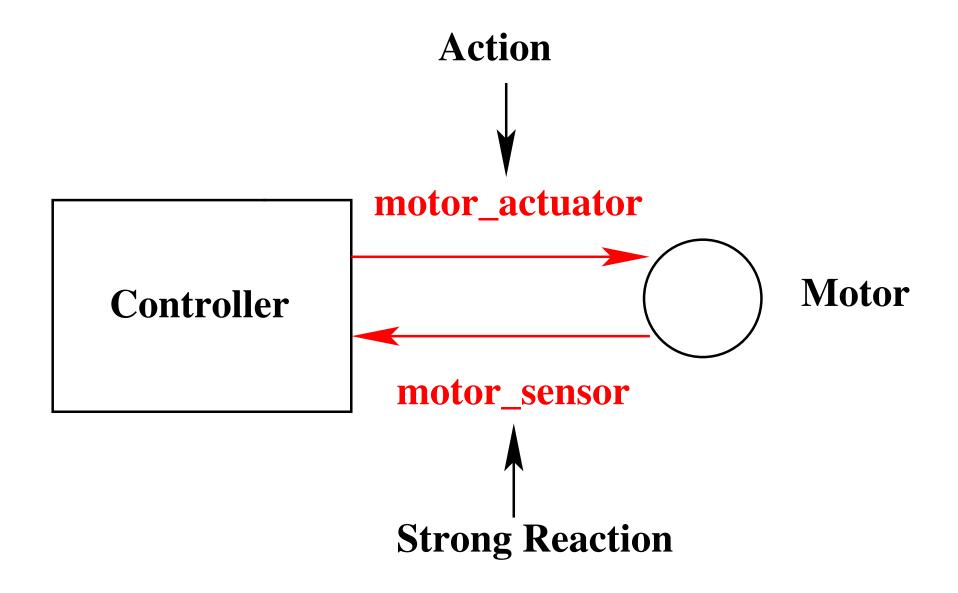
 $axm0_1: STATUS = \{stopped, working\}$ 

axm0\_2:  $stopped \neq working$ 

 $egin{array}{ll} extbf{variables:} & motor\_actuator \\ motor\_sensor \end{array}$ 

inv0\_1:  $motor\_sensor \in STATUS$ 

inv0\_2:  $motor\_actuator \in STATUS$ 



- We instantiate the strong pattern as follows:

 $\boldsymbol{a}$ 

motor actuator

Motor\_stop

- Convention: Controller events start with "treat\_"

 $r_{off} \sim$ 

init

$$a := 0$$

$$r := 0$$

init

 $motor\_actuator := stopped$ 

 $motor\_sensor := stopped$ 

```
{f a}_{-}on{f when} a=0 r=0 then a:=1 end
```

```
egin{array}{ll} {\sf treat\_start\_motor} & {\sf when} \\ & {\it motor\_actuator} = stopped \\ & {\it motor\_sensor} = stopped \\ {\sf then} \\ & {\it motor\_actuator} := working \\ {\sf end} \end{array}
```

```
r\_{
m on} when r=0 a=1 then r:=1 end
```

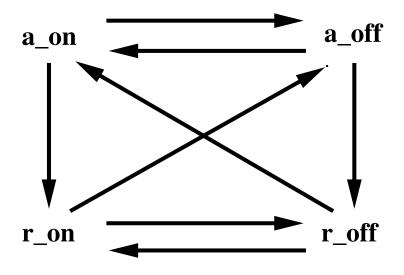
```
egin{array}{ll} {\sf Motor\_start} & {\sf when} \\ & {\it motor\_sensor} = {\it stopped} \\ & {\it motor\_actuator} = {\it working} \\ & {\sf then} \\ & {\it motor\_sensor} := {\it working} \\ & {\sf end} \end{array}
```

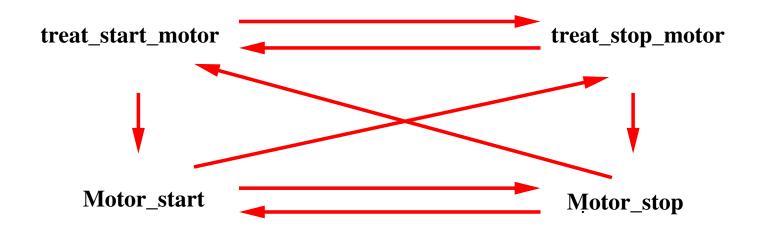
```
{f a}_{-} off {f when} {f a}=1 {f r}=1 then {f a}:=0 end
```

```
egin{array}{ll} {\sf treat\_stop\_motor} & {\sf when} \\ & {\it motor\_actuator} = {\it working} \\ & {\it motor\_sensor} = {\it working} \\ & {\sf then} \\ & {\it motor\_actuator} := {\it stopped} \\ & {\sf end} \\ \end{array}
```

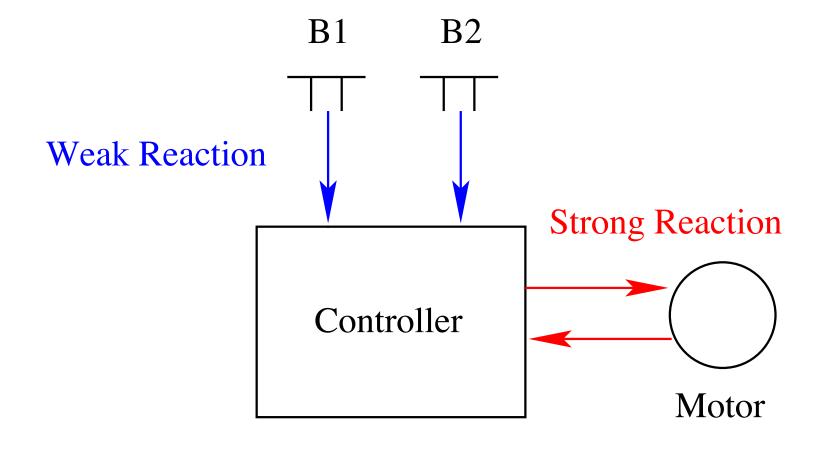
```
r\_{
m off} when r=1 a=0 then r:=0 end
```

```
egin{array}{ll} {\sf Motor\_stop} & {\sf when} \\ & {\it motor\_sensor} = {\it working} \\ & {\it motor\_actuator} = {\it stopped} \\ {\sf then} \\ & {\it motor\_sensor} := {\it stopped} \\ {\sf end} \end{array}
```





- Environment
  - motor\_start
  - motor\_stop
- Controller
  - treat\_start\_motor
  - treat\_stop\_motor



Buttons and Controller are weakly synchronized

FUN<sub>1</sub>

The counters have

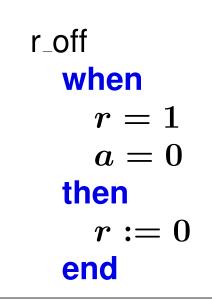
been removed

$${f a}_{-}$$
on ${f when} \ a=0 \ {f then} \ a:=1 \ {f end}$ 

$${f a}_{-}$$
 off when  $a=1$  then  $a:=0$  end

$$\begin{array}{c} \text{init} \\ a := 0 \\ r := 0 \end{array}$$

$$r_{-}$$
on $r=0$  $a=1$ then $r:=1$ end



variables: ...

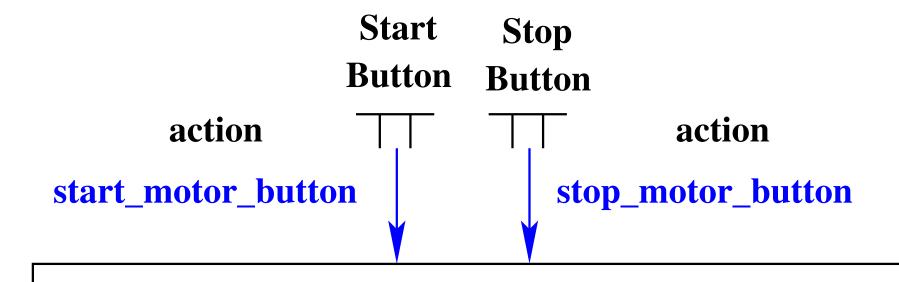
 $start\_motor\_button \\ stop\_motor\_button \\ start\_motor\_impulse \\ stop\_motor\_impulse$ 

inv1\_1:  $stop\_motor\_button \in BOOL$ 

inv1\_2:  $start\_motor\_button \in BOOL$ 

inv1\_3:  $stop\_motor\_impulse \in BOOL$ 

inv1\_4:  $start\_motor\_impulse \in BOOL$ 



start\_motor\_impulse
weak reaction

stop\_motor\_impulse

weak reaction

**CONTROLLER** 

- We instantiate the pattern as follows:

```
egin{array}{lll} a & 
ightsquigarrow & start\_motor\_button \ r & 
ightsquigarrow & start\_motor\_impulse \ 0 & 
ightsquigarrow & {
m FALSE} \ 1 & 
ightsquigarrow & {
m TRUE} \ \end{array}
```

- We rename treat\_start\_motor as treat\_push\_start\_motor\_button

init

a := 0

r := 0

init

 $motor\_actuator := stopped \ motor\_sensor := stopped$ 

 $start\_motor\_button := { t FALSE}$ 

 $start\_motor\_impulse := { t FALSE}$ 

```
{f a}_{-}on when a=0 then a:=1 end
```

```
egin{array}{ll} {\sf push\_start\_motor\_button} & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & \\ & & & \\ & & \\ & & & \\ & & \\ & & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & &
```

```
{f a}_{-} off when a=1 then a:=0 end
```

```
release\_start\_motor\_button
start\_motor\_button = TRUE
then
start\_motor\_button := FALSE
end
```

```
r on
  when
    r = 0
    a = 1
  then
    r := 1
  end
```

```
treat_push_start_motor_button
  refines
    treat_start_motor
  when
    start\_motor\_impulse = {	t FALSE}
    start\_motor\_button = \mathrm{TRUE}
    motor\_actuator = stopped
    motor\_sensor = stopped
  then
    start\_motor\_impulse := \mathrm{TRUE}
    motor\_actuator := working
  end
```

- This is the most important slide of the talk
- We can see how patterns can be superposed

```
a_on a=0 r=0 then a:=1 end
```

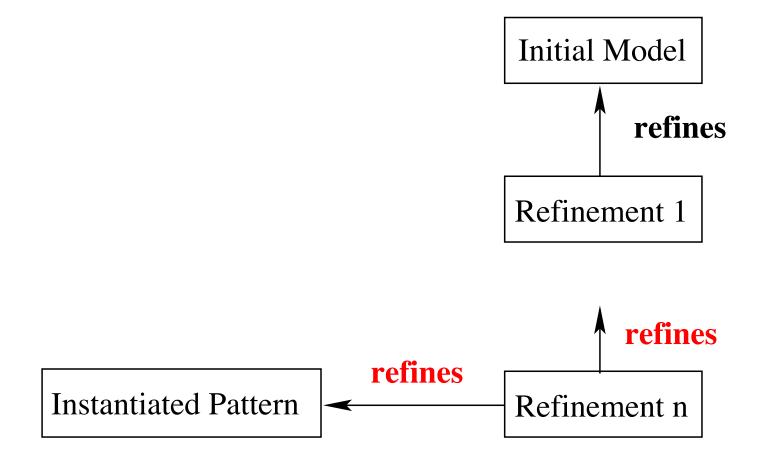
```
treat_start_motor
when

motor_actuator = stopped
motor_sensor = stopped
then

motor_actuator := working
end
```

```
r on
  when
     r = 0
    a = 1
  then
    r := 1
  end
```

```
treat_push_start_motor_button
  when
    start\_motor\_impulse = {
m FALSE}
    start\_motor\_button = \mathrm{TRUE}
    motor\_actuator = stopped
    motor\_sensor = stopped
  then
    start\_motor\_impulse := \mathrm{TRUE}
    motor\_actuator := working
  end
```



```
r\_{
m off} when r=1 a=0 then r:=0 end
```

```
egin{array}{lll} treat_release\_start\_motor\_button & & & & \\ & & & start\_motor\_impulse = TRUE \\ & & start\_motor\_button = FALSE \\ & then & & start\_motor\_impulse := FALSE \\ & end & & \end{array}
```

- We instantiate the pattern as follows:

```
egin{array}{lll} a & 
ightsquigarrow & stop\_motor\_button \ r & 
ightsquigarrow & stop\_motor\_impulse \ 0 & 
ightsquigarrow & {
m FALSE} \ 1 & 
ightsquigarrow & {
m TRUE} \ \end{array}
```

```
a_on  
→  push_stop_motor_button

a_off  
→  release_stop_motor_button

r_on  
→  treat_push_stop_motor_button

r_off  
→  treat_release_stop_motor_button
```

init

a := 0r := 0  $motor\_actuator := stopped \ motor\_sensor := stopped \ start\_motor\_button := FALSE \ start\_motor\_impulse := FALSE \ stop\_motor\_button := FALSE \ stop\_motor\_impulse := FALSE$ 

```
{f a}_{-}on{f when} \ a=0 \ {f then} \ a:=1 \ {f end}
```

```
\begin{array}{l} {\sf push\_stop\_motor\_button} \\ {\it when} \\ {\it stop\_motor\_button} = {\sf FALSE} \\ {\it then} \\ {\it stop\_motor\_button} := {\sf TRUE} \\ {\it end} \end{array}
```

```
{f a}_{-} off when a=1 then a:=0 end
```

```
egin{array}{c} 	ext{release\_stop\_motor\_button} \ 	ext{$stop\_motor\_button} = 	ext{TRUE} \ 	ext{then} \ 	ext{$stop\_motor\_button} := 	ext{FALSE} \ 	ext{end} \end{array}
```

```
r_on
```

## when

r = 0 a = 1

## then

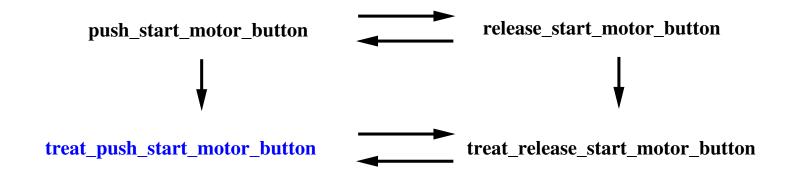
r := 1

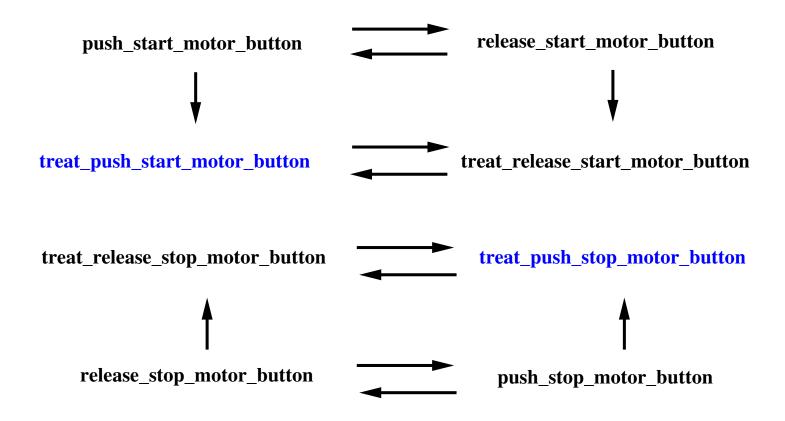
end

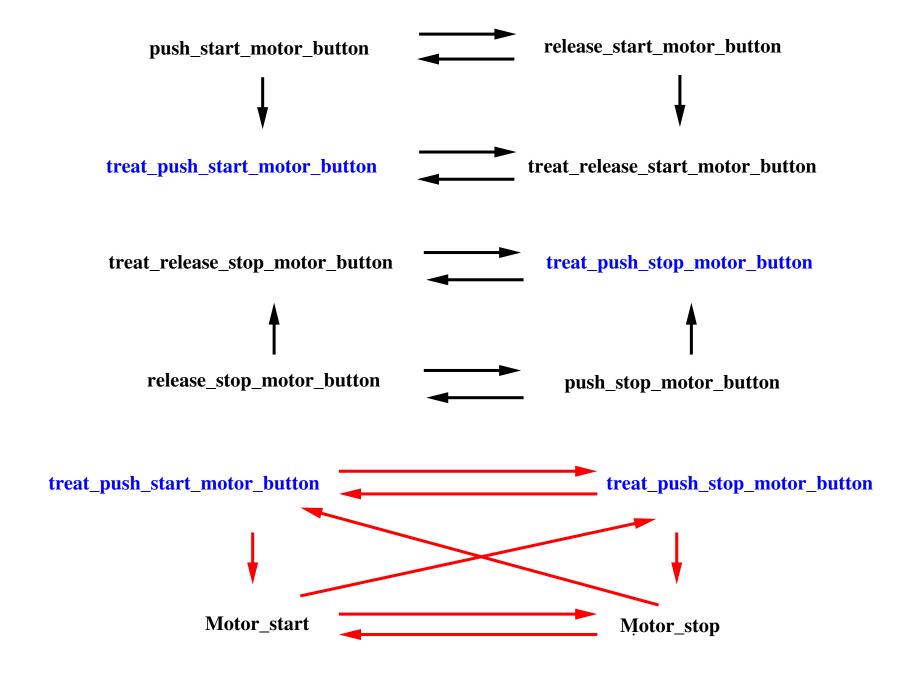
```
treat_push_stop_motor_button
  refines
    treat_stop_motor
  when
    stop\_motor\_impulse = {
m FALSE}
    stop\_motor\_button = \mathrm{TRUE}
    motor\_sensor = working
    motor\_actuator = working
  then
    stop\_motor\_impulse := \mathrm{TRUE}
    motor\_actuator := stopped
  end
```

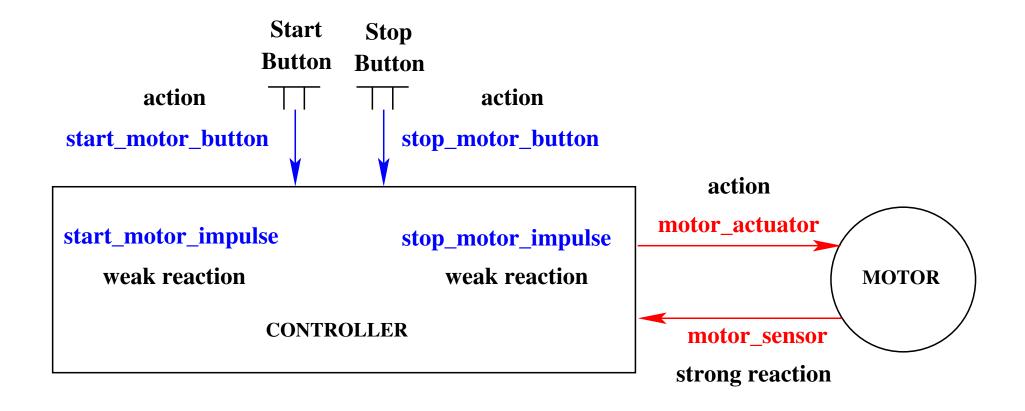
```
r\_{
m off} when r=1 a=0 then r:=0 end
```

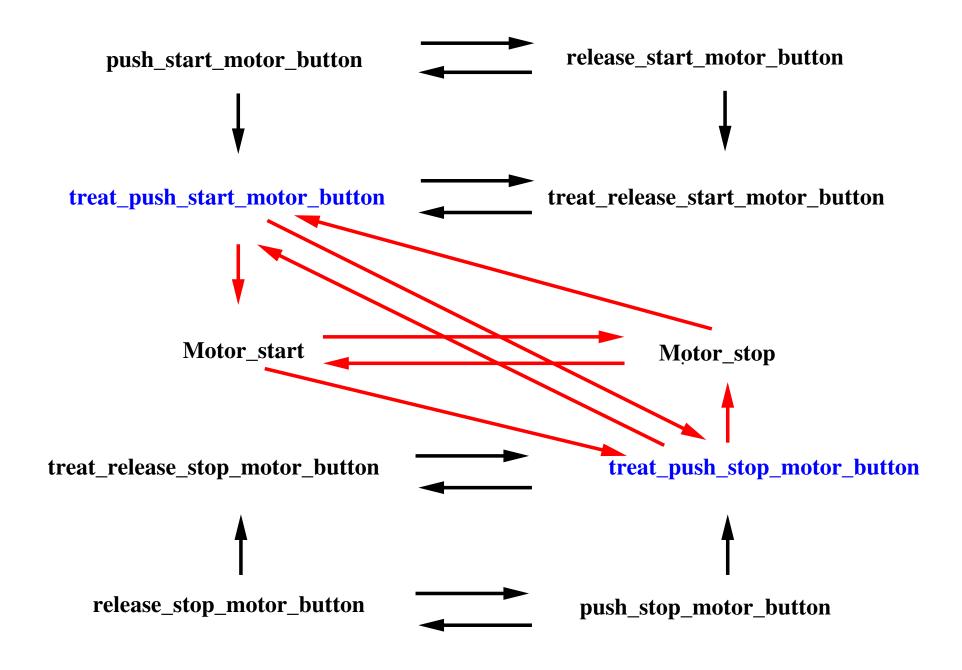
```
egin{array}{lll} treat_release\_stop\_motor\_button & & & & stop\_motor\_impulse = TRUE & & stop\_motor\_button = FALSE & & & then & & stop\_motor\_impulse := FALSE & end & & & & \end{array}
```











```
treat_push_start_motor_button
refines
treat_start_motor
when

start_motor_impulse = FALSE
start_motor_button = TRUE
motor_actuator = stopped
motor_sensor = stopped
then

start_motor_impulse := TRUE
motor_actuator := working
end
```

- What happens when the following hold

```
\neg \left(motor\_actuator = stopped \ \land \ motor\_sensor = stopped \right)
```

- We need another event

```
treat_push_start_motor_button
refines
treat_start_motor
when

start_motor_impulse = FALSE
start_motor_button = TRUE
motor_actuator = stopped
motor_sensor = stopped
then

start_motor_impulse := TRUE
motor_actuator := working
end
```

- In the second case, the button has been pushed but the internal conditions are not met
- However, we need to record that the button has been pushed:

 $start\_motor\_impulse := \text{TRUE}$ 

```
treat_push_stop_motor_button
    refines
    treat_stop_motor
    when

    stop_motor_impulse = FALSE
    stop_motor_button = TRUE
    motor_sensor = working
    motor_actuator = working
    then

    stop_motor_impulse := TRUE
    motor_actuator := stopped
end
```

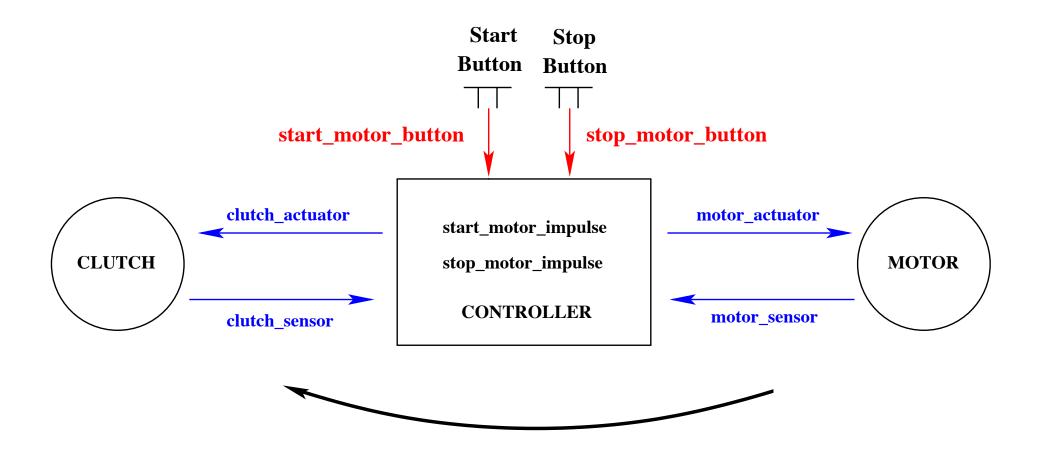
```
when
    stop_motor_impulse = FALSE
    stop_motor_button = TRUE
    ¬(motor_sensor = working \wedge motor_actuator = working)
then
    stop_motor_impulse := TRUE
end
```

- In the second case, the button has been pushed but the internal conditions are not met
- However, we need to record that the button has been pushed:

 $stop\_motor\_impulse := \mathrm{TRUE}$ 

- Environment
  - motor\_start
  - motor\_stop
  - push\_start\_motor\_button
  - release\_start\_motor\_button
  - push\_stop\_motor\_button
  - release\_stop\_motor\_button

- Controller
  - treat\_push\_start\_motor\_button
  - treat\_push\_start\_motor\_button\_false
  - treat\_push\_stop\_motor\_button
  - treat\_push\_stop\_motor\_button\_false
  - treat\_release\_start\_motor\_button
  - treat\_release\_stop\_motor\_button



- We introduce the set in a new context:

$$CLUTCH = \{engaged, disengaged\}$$

- We copy the initial model where we instantiate:

- Environment
  - motor\_start
  - motor\_stop
  - clutch\_start
  - clutch\_stop
  - push\_start\_motor\_button
  - release\_start\_motor\_button
  - push\_stop\_motor\_button
  - release\_stop\_motor\_button

- Controller
  - treat\_push\_start\_motor\_button
  - treat\_push\_start\_motor\_button\_false
  - treat\_push\_stop\_motor\_button
  - treat\_push\_stop\_motor\_button\_false
  - treat\_release\_start\_motor\_button
  - treat\_release\_stop\_motor\_button
  - treat\_start\_clutch
  - treat\_stop\_clutch

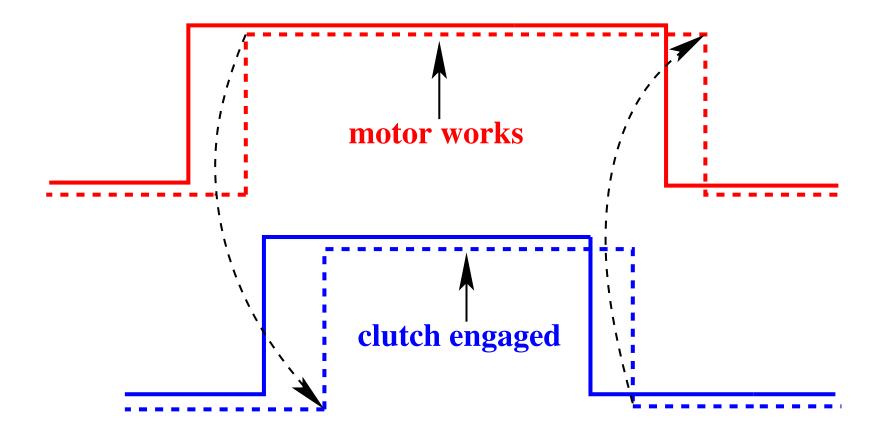
- An additional safety constraint

When the clutch is engaged, the motor must work

SAF<sub>1</sub>

- For this we develop ANOTHER DESIGN PATTERN

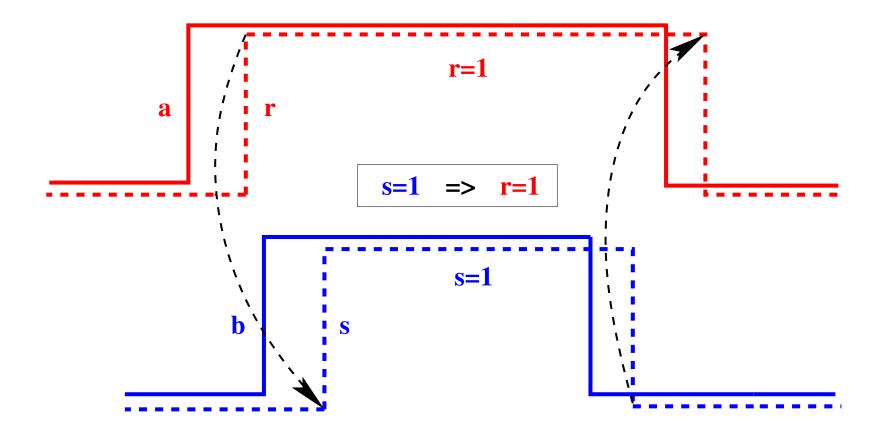
- It is called: Weak synchronization of two Strong Reactions



When the clutch is engaged

then

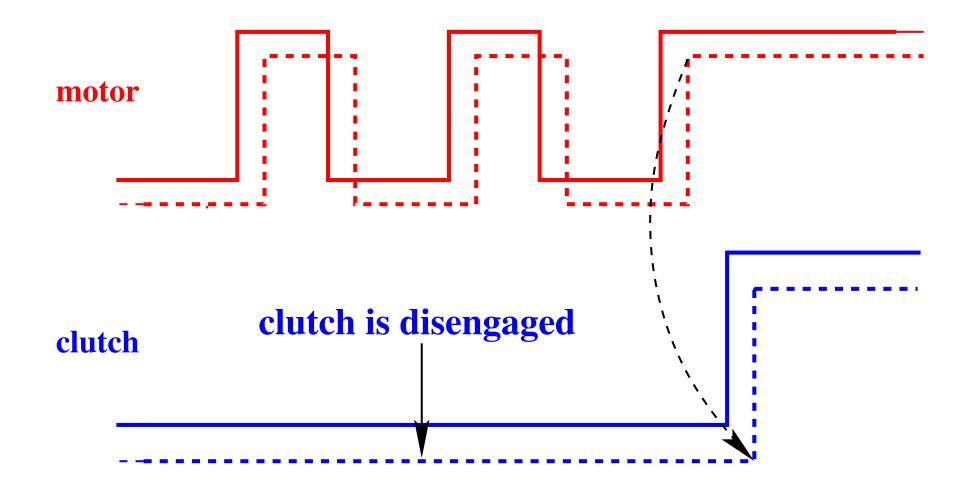
the motor must work



When the clutch is engaged

then

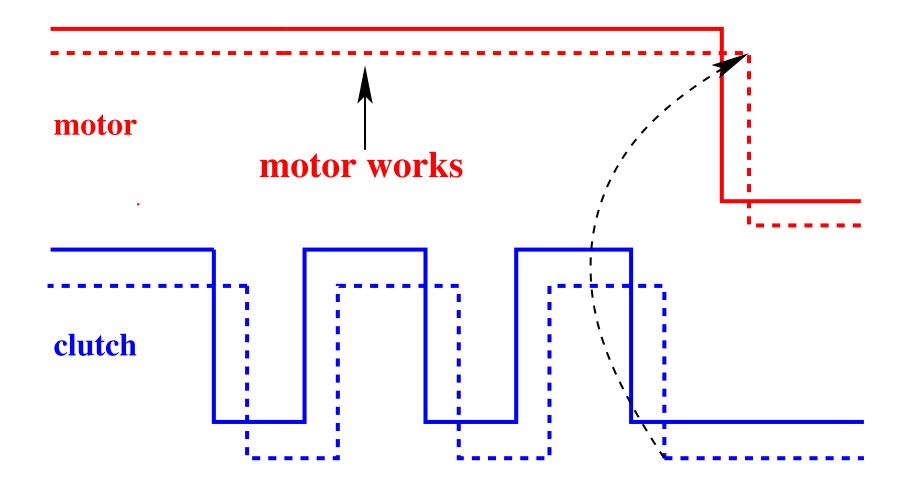
the motor must work



When the clutch is disengaged,

then

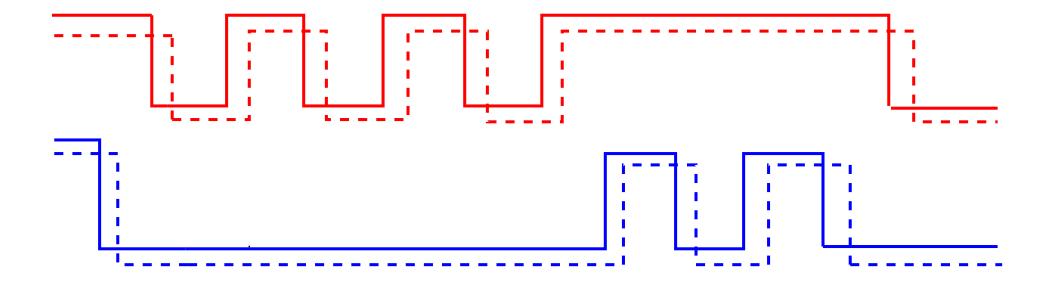
the motor can be started and stopped several times

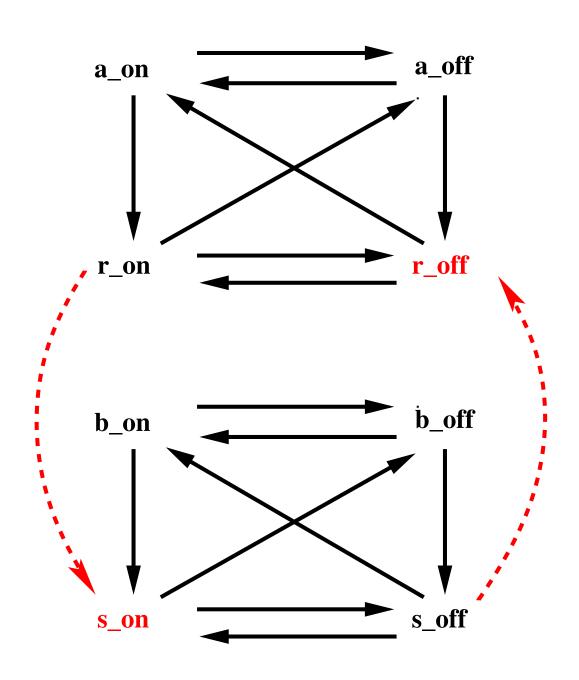


When the motor works,

then

the clutch can be engaged and disengaged several times





```
a~\in~\{0,1\}
dbl0_1:
        r \in \{0,1\}
dbl0_2:
dbl0_3:
        ca \in \mathbb{N}
dbl0_4:
       cr \in \mathbb{N}
dbl0_5: a=1 \land r=0 \Rightarrow ca=cr+1
dbl0_6:
       a=0 \ \lor \ r=1 \ \Rightarrow \ ca=cr
dbl0_7: b \in \{0,1\}
dbl0_8: s \in \{0,1\}
dbl0_9: cb \in \mathbb{N}
dbl0_10: cs \in \mathbb{N}
dbl0_11: b=1 \wedge s=0 \Rightarrow cb=cs+1
dbl0_12: b=0 \lor s=1 \Rightarrow cb=cs
```

```
a_on when a=0 r=0 then a:=1 ca:=ca+1 end
```

```
r\_{on}
when
r=0
a=1
then
r:=1
cr:=cr+1
end
```

```
egin{aligned} {
m a\_off} & {
m when} & a=1 & \ r=1 & {
m then} & \ a:=0 & {
m end} & \end{aligned}
```

$$egin{array}{c} {\sf r\_off} \\ {\sf when} \\ {\it r}=1 \\ {\it a}=0 \\ {\sf then} \\ {\it r}:=0 \\ {\sf end} \\ \end{array}$$

```
b_on when b=0 s=0 then b:=1 cb:=cb+1 end
```

$$s\_on$$
 $s=0$ 
 $b=1$ 
 $then$ 
 $s:=1$ 
 $cs:=cs+1$ 
 $end$ 

```
egin{aligned} & 	extbf{b\_off} \\ & 	extbf{when} \\ & b = 1 \\ & s = 1 \\ & 	extbf{then} \\ & b := 0 \\ & 	extbf{end} \end{aligned}
```

$$s\_off$$
when
 $s=1$ 
 $b=0$ 
then
 $s:=0$ 
end

dbl1\_1: 
$$s=1 \Rightarrow r=1$$

- It seems sufficient to add the following guards

```
egin{array}{ll} 	extsf{s\_on} & when & s = 0 & & & & & & & \\ b = 1 & & & & & & & & & & \\ r = 1 & & & & & & & & & & \\ then & & & & & & & & & & & \\ s := 1 & & & & & & & & & & \\ cs := cs + 1 & & & & & & & \\ end & & & & & & & & & & \\ \end{array}
```

```
r\_{
m off}
{m when}
{m r}=1
{m a}=0
{m s}=0
{m then}
{m r}:=0
{m end}
```

- But we do not want to touch these events

```
s_on
 when
   s = 0
   b=1
           r = 1
 then
   s := 1
   cs := cs + 1
  end
```

```
r_off
 when
    r = 1
    a = 0
            s = 0
  then
    r := 0
  end
```

- We introduce the following additional invariants

dbl1\_2: 
$$b=1 \Rightarrow r=1$$
dbl1\_3:  $a=0 \Rightarrow s=0$ 

dbl1\_2: 
$$b=1 \Rightarrow r=1$$

In order to maintain this invariant, we have to refine bon

 $egin{aligned} & extbf{b\_on} \\ & extbf{when} \\ & extbf{b} = 0 \\ & extbf{s} = 0 \\ & extbf{r} = 1 \\ & extbf{then} \\ & extbf{b} := 1 \\ & extbf{c} b := cb + 1 \\ & extbf{end} \end{aligned}$ 

dbl1\_2: 
$$b=1 \Rightarrow r=1$$
  $(r=0 \Rightarrow b=0)$ 

In order to maintain this invariant, we have to refine r\_off

```
r\_{
m off} when r=1 a=0 then r:=0 end r\_{
m off} when r=1 a=0 b=0 then r:=0
```

- But, again, we do not want to touch this event

```
r\_{
m off} when r=1 a=0 b=0 then r:=0 end
```

- We introduce the following invariant

dbl1\_4: 
$$a = 0 \Rightarrow b = 0$$

dbl1\_3: 
$$a=0 \Rightarrow s=0$$

In order to maintain this invariant, we have to refine a\_off

 $a\_{
m off}$  when a=1 r=1 then a:=0 end

 $a\_{
m off}$ when a=1 r=1 s=0then a:=0end

dbl1\_3: 
$$a=0 \Rightarrow s=0$$
  $(s=1 \Rightarrow a=1)$ 

In order to maintain this invariant, we have to refine son

```
s on
s_on
                           when
 when
                            s = 0
   s = 0
                            b = 1
   b=1
                            a = 1
 then
                           then
   s := 1
                            s := 1
   cs := cs + 1
                            cs := cs + 1
 end
                           end
```

- But, again, we do not want to touch this event

- We have to introduce the following invariant

$$b=1 \Rightarrow a=1$$

- Fortunately, this is **dbl1\_4** ( $a=0 \implies b=0$ ) contraposed

dbl1\_4: 
$$a=0 \Rightarrow b=0$$

In order to maintain this invariant, we have to refine a\_off again

 $a\_{
m off}$ when a=1 r=1 s=0then a:=0end

 $a\_{
m off}$ when a=1 r=1 s=0 b=0then a:=0end

dbl1\_4: 
$$a=0 \Rightarrow b=0$$
  $(b=1 \Rightarrow a=1)$ 

In order to maintain this invariant, we have to refine bon again

```
egin{aligned} & 	extbf{b\_on} \\ & 	extbf{when} \\ & 	extbf{b} = 0 \\ & 	extbf{s} = 0 \\ & 	extbf{r} = 1 \\ & 	extbf{then} \\ & 	extbf{b}, cb := 1, cb + 1 \\ & 	extbf{end} \end{aligned}
```

```
dbl1_1: s=1 \Rightarrow r=1
dbl1_2: b=1 \Rightarrow r=1
dbl1_3: a = 0 \Rightarrow s = 0
dbl1_4: a = 0 \Rightarrow b = 0
```

```
b_on
when
 b = 0
 s = 0
 r = 1
  a = 1
then
 b, cb := 1, cb + 1
end
```

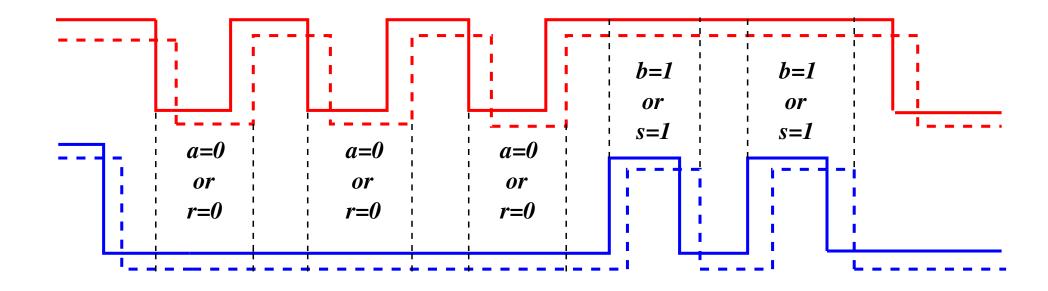
a\_off when a = 1r = 1s = 0b = 0then a := 0end

This can be put into a single invariant

**dbl1\_5:** 
$$b=1$$
  $\vee$   $s=1$   $\Rightarrow$   $a=1$   $\wedge$   $r=1$ 

with the following contraposed form

dbl1\_6: 
$$a=0 \ \lor \ r=0 \ \Rightarrow \ b=0 \ \land \ s=0$$



Reminder: - - - is the motor and - - - is the clutch

**dbl1\_5:** b=1  $\vee$  s=1  $\Rightarrow$  a=1  $\wedge$  r=1

**dbl1\_6:**  $a=0 \lor r=0 \Rightarrow b=0 \land s=0$ 

```
a\_on when a=0 r=0 then a:=1 end
```

```
a_off when a=1 r=1 s=0 b=0 then a:=0 end
```

```
r\_on
when
r=0
a=1
then
r:=1
end
```

```
r\_{off}
when
r=1
a=0
then
r:=0
end
```

```
b_on when b = 0 s = 0 r = 1 a = 1 then b := 1 end
```

```
\begin{array}{c} \text{b\_off}\\ \textbf{when}\\ b=1\\ s=1\\ \textbf{then}\\ b:=0\\ \textbf{end} \end{array}
```

```
s_{-}on when s=0 b=1 then s:=1 end
```

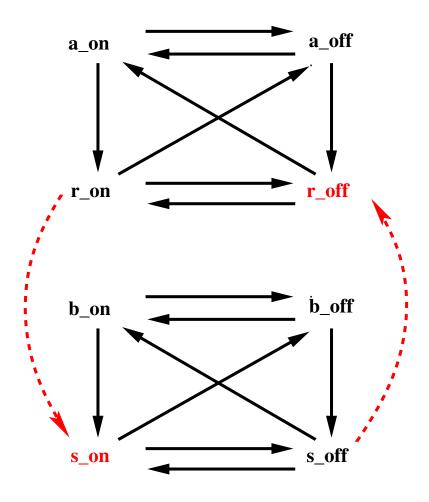
$$s\_off$$
when
 $s=1$ 
 $b=0$ 
then
 $s:=0$ 
end

```
      dbl1_1: s = 1 \Rightarrow r = 1
      (r = 0 \Rightarrow s = 0)

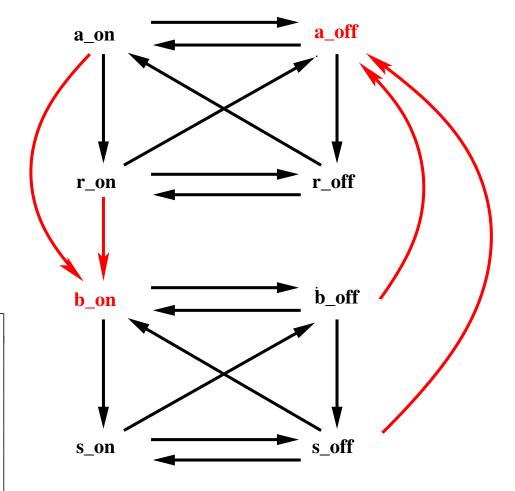
      dbl1_2: b = 1 \Rightarrow r = 1
      (r = 0 \Rightarrow b = 0)

      dbl1_3: a = 0 \Rightarrow s = 0
      (s = 1 \Rightarrow a = 1)

      dbl1_4: a = 0 \Rightarrow b = 0
      (b = 1 \Rightarrow a = 1)
```



dbl1\_1: 
$$s=1 \Rightarrow r=1$$
  $(r=0 \Rightarrow s=0)$ 



 $egin{aligned} & b\_on \\ & when \\ & b=0 \\ & s=0 \\ & r=1 \\ & a=1 \\ & then \\ & b:=1 \\ & end \end{aligned}$ 

 $a\_{
m off}$ when a=1 r=1 s=0 b=0then a:=0end

When the clutch is engaged, the motor must work

SAF\_1

$$egin{array}{ll} ext{inv3\_1:} & clutch\_sensor = engaged \\ \Rightarrow & \\ motor\_sensor = working \end{array}$$

- This is an instance of the previous design pattern

- We instantiate the pattern as follows:

```
treat_push_start_motor_button
                motor actuator
                                                 a_on
                                                             \sim \rightarrow
\boldsymbol{a}
                                                                      treat_push_stop_motor_button
                motor\_sensor
                                                 a_off
       \sim \rightarrow
                                                             \sim \rightarrow
                                                                      Motor start
0
                stopped
                                                 r_on
       \sim \rightarrow
                                                             \sim \rightarrow
                working
                                                                      Motor_stop
                                                 r_off
       \sim \rightarrow
                                                             \sim \rightarrow
         \boldsymbol{b}
                                                         b on
                                                                              treat start clutch
                         clutch\_actuator
                                                                               treat_stop_clutch
                         clutch\_sensor
                                                         b_off
                                                                     \sim \rightarrow
         s
                         disengaged
                                                                               Clutch_start
         0
                                                         s_on
                                                                     \sim \rightarrow
                         engaged
                                                                               Clutch_stop
                                                         s_off
                                                                     \sim \rightarrow
```

dbl1\_1: 
$$s=1 \Rightarrow r=1$$

dbl1\_2: 
$$b=1 \Rightarrow r=1$$

$$clutch\_sensor = engaged$$
inv3\_1:  $\Rightarrow$ 
 $motor\_sensor = working$ 

 $clutch\_actuator = engaged$ inv3\_2:  $\Rightarrow$   $motor\_sensor = working$ 

**dbl1\_3:** 
$$a=0 \Rightarrow s=0$$

dbl1\_4: 
$$a = 0 \Rightarrow b = 0$$

$$motor\_actuator = stopped$$
inv3\_3:  $\Rightarrow$ 
 $clutch\_sensor = disengaged$ 

 $motor\_actuator = stopped$  inv3\_4:  $\Rightarrow$   $clutch\_actuator = disengaged$ 

```
egin{aligned} & b\_on \ & when \ & b=0 \ & s=0 \ & r=1 \ & a=1 \ & then \ & b:=1 \ & end \end{aligned}
```

```
egin{array}{ll} {
m treat\_start\_clutch} & {
m when} \\ & clutch\_actuator = disengaged \\ & clutch\_sensor = disengaged \\ & motor\_sensor = working \\ & motor\_actuator = working \\ {
m then} \\ & clutch\_actuator := engaged \\ {
m end} \\ \end{array}
```

```
a_off when
```

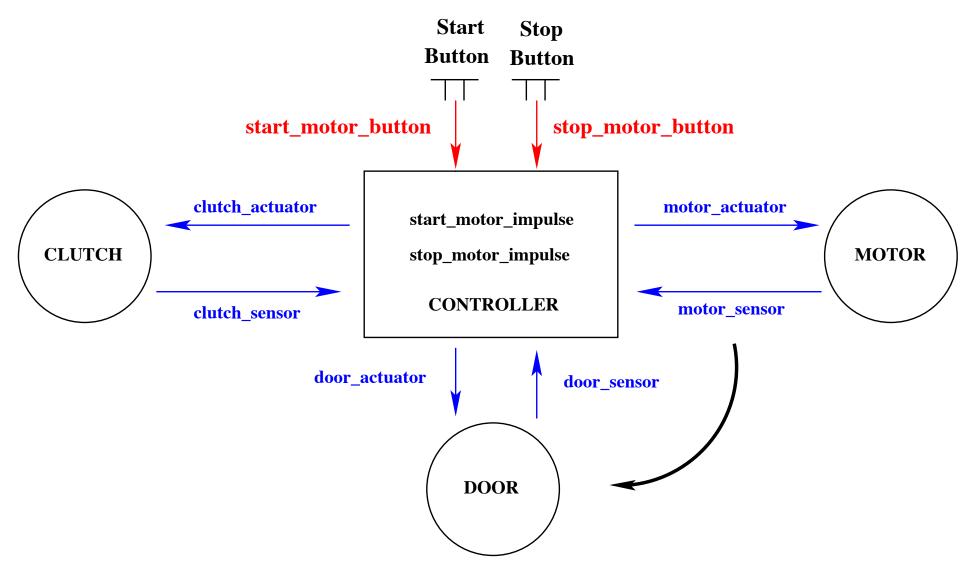
```
a=1 \ r=1 \ s=0 \ b=0 \ 	an a:=0
```

end

```
treat_push_stop_motor_button
 when
    stop\_motor\_impulse = {
m FALSE}
    stop\_motor\_button = \mathrm{TRUE}
    motor\_actuator = working
    motor\_sensor = working
    clutch\_sensor = disengaged
    clutch\_actuator = disengaged
 then
    motor\_actuator := stopped
    stop\_motor\_impulse := \mathrm{TRUE}
  end
```

- Environment (no new events)
  - motor\_start
  - motor\_stop
  - clutch\_start
  - clutch\_stop
  - push\_start\_motor\_button
  - release\_start\_motor\_button
  - push\_stop\_motor\_button
  - release\_stop\_motor\_button

- Controller (no new events)
  - treat\_push\_start\_motor\_button
  - treat\_push\_start\_motor\_button\_false
  - treat\_push\_stop\_motor\_button
  - treat\_push\_stop\_motor\_button\_false
  - treat\_release\_start\_motor\_button
  - treat\_release\_stop\_motor\_button
  - treat\_start\_clutch
  - treat\_stop\_clutch



- We copy (after renaming "motor" to "door") what has been done in the initial model

- We introduce the set in a new context:

$$DOOR = \{open, closed\}$$

- We copy the initial model where we instantiate:

 $egin{array}{lll} motor & 
ightsquare & door \ & STATUS & 
ightsquare & DOOR \ & working & 
ightsquare & closed \ & stopped & 
ightsquare & open \ \end{array}$ 

- Environment
  - motor\_start
  - motor\_stop
  - clutch\_start
  - clutch\_stop
  - door\_close
  - door\_open
  - push\_start\_motor\_button
  - release\_start\_motor\_button
  - push\_stop\_motor\_button
  - release\_stop\_motor\_button

- Controller
  - treat\_push\_start\_motor\_button
  - treat\_push\_start\_motor\_button\_false
  - treat\_push\_stop\_motor\_button
  - treat\_push\_stop\_motor\_button\_false
  - treat\_release\_start\_motor\_button
  - treat\_release\_stop\_motor\_button
  - treat\_start\_clutch
  - treat\_stop\_clutch
  - treat\_close\_door
  - treat\_open\_door

- An additional safety constraint

When the clutch is engaged, the door must be closed

SAF<sub>2</sub>

- We copy (after renaming "motor" to "door") what has been done in the third model:

When the clutch is engaged, the motor must work

SAF\_1

- Can you guess it?

- Can you guess it?

- When the motor is not working, we must allow users:
  - to change the tool
  - to replace the part to be treated

- Can you guess it?

- When the motor is not working, we must allow users:
  - to change the tool
  - to replace the part to be treated

- Hence the following additional requirement (which was forgotten)

When the motor is stopped, the door must be open

SAF\_3

- Can you guess it?
- When the motor is not working, we must allow users:
  - to change the tool
  - to replace the part to be treated
- Hence the following additional requirement (which was forgotten)

When the door is closed, the motor must work

SAF\_3'

- SAF\_3' is the contraposed form of SAF\_3

- Additional safety constraint

When the door is closed, the motor must work

SAF\_3'

- We copy (after renaming "clutch" to "door") what has been done in the third model:

When the clutch is engaged, the motor must work

SAF<sub>1</sub>

When the clutch is engaged, the motor must work

SAF<sub>1</sub>

When the clutch is engaged, the door must be closed

SAF<sub>2</sub>

When the door is closed, the motor must work

SAF 3'

- Requirement SAF\_1 is now redundant: SAF\_2 ∧ SAF\_3' ⇒ SAF\_1

- Initial model: Connecting the controller to the motor
- 1st refinement: Connecting the motor button to the controller
- 2nd refinement: Connecting the controller to the clutch
- 3rd (4th) refinement: Connecting the controller to the door

4th (5th) refinement: Constraining the clutch and the door
 Constraining the motor and the door

- 5th (6th) refinement: More constraints between clutch and door

- 6th (7th) refinement: Connecting the clutch button to the controller

- Environment (no new events)
  - motor\_start
  - motor\_stop
  - clutch\_start
  - clutch\_stop
  - door\_close
  - door\_open
  - push\_start\_motor\_button
  - release\_start\_motor\_button
  - push\_stop\_motor\_button
  - release\_stop\_motor\_button

- Controller (no new events)
  - treat\_push\_start\_motor\_button
  - treat\_push\_start\_motor\_button\_false
  - treat\_push\_stop\_motor\_button
  - treat\_push\_stop\_motor\_button\_false
  - treat\_release\_start\_motor\_button
  - treat\_release\_stop\_motor\_button
  - treat\_start\_clutch
  - treat\_stop\_clutch
  - treat\_close\_door
  - treat\_open\_door

- Adding two functional constraints

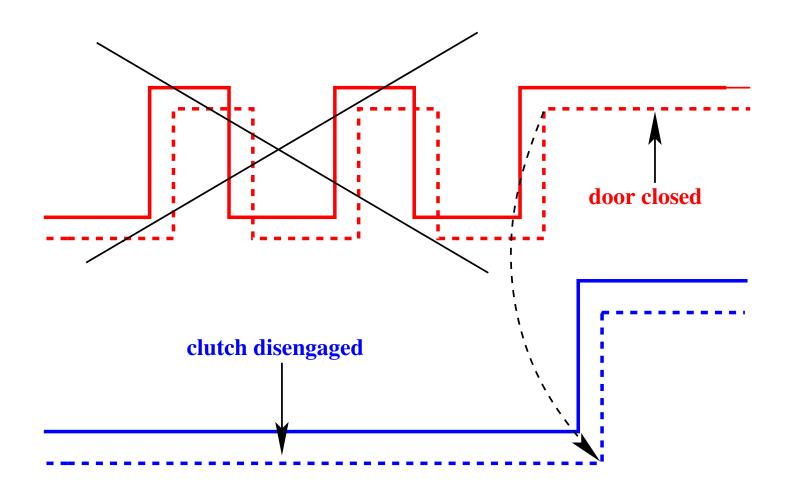
When the clutch is disengaged, the door cannot be closed several times, ONLY ONCE

FUN 3

When the door is closed, the clutch cannot be disengaged several times, ONLY ONCE

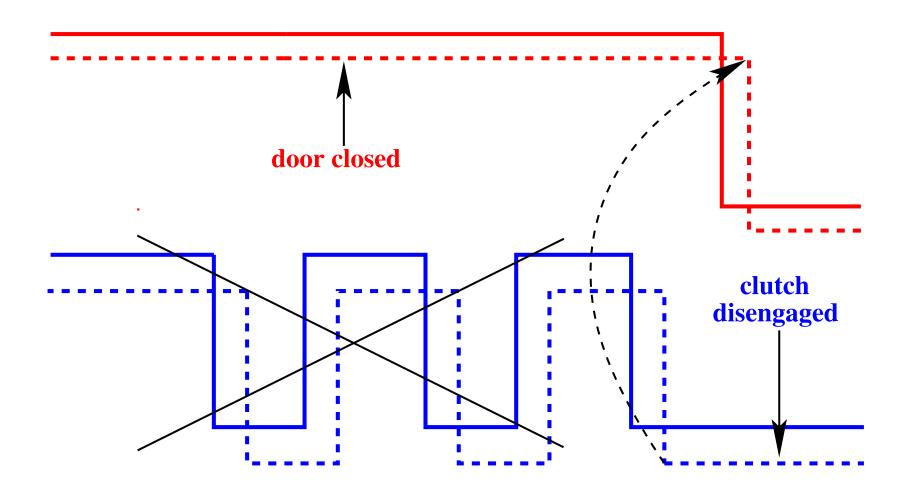
FUN\_4

## Problem with the Weak Synchronization of Strong Reactions 225

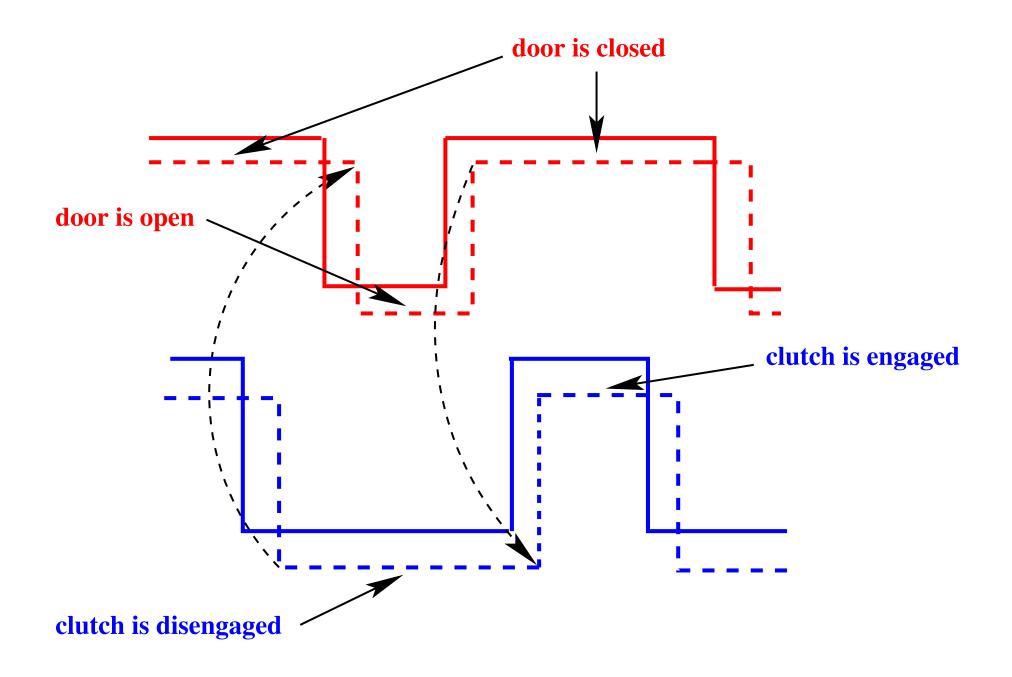


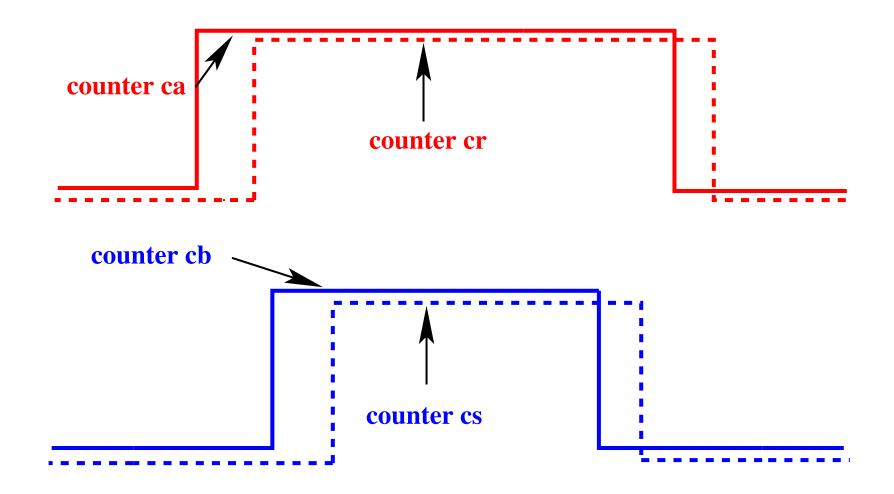
 When the clutch is disengaged, the door cannot be closed several times

## Problem with the Weak Synchronization of Strong Reactions 226



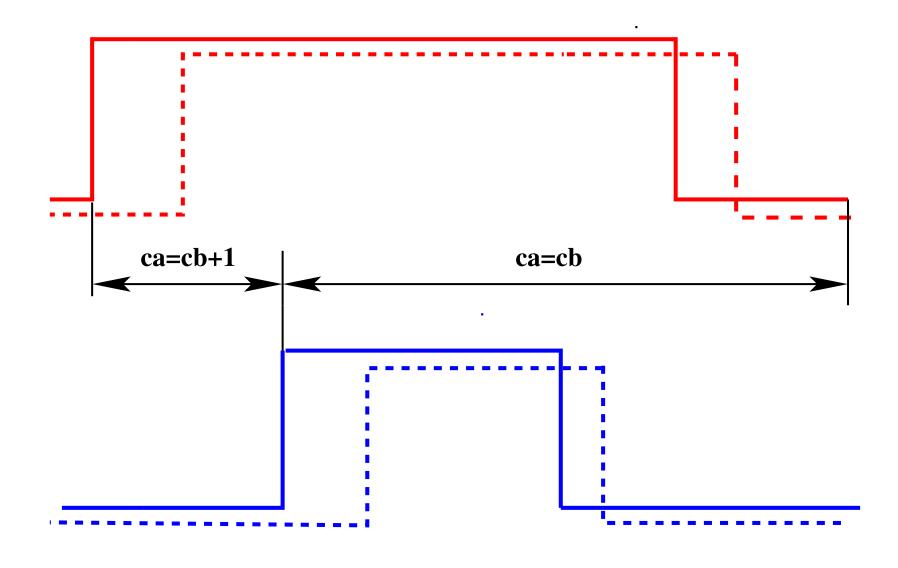
 When the door is closed, the clutch cannot be disengaged several times

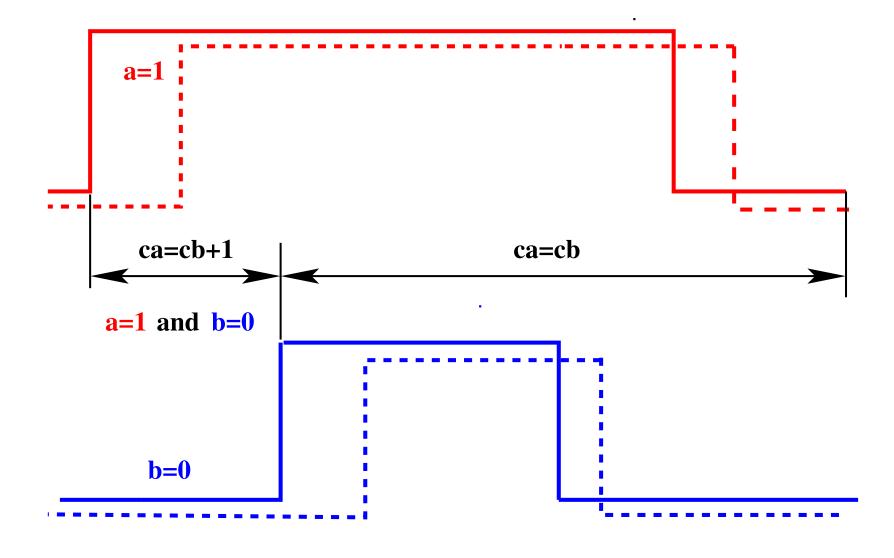


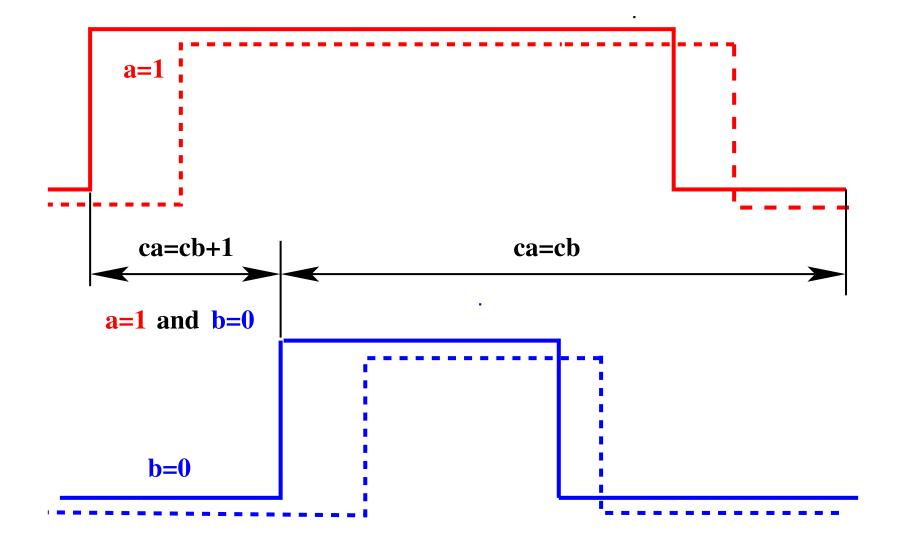


What we want:

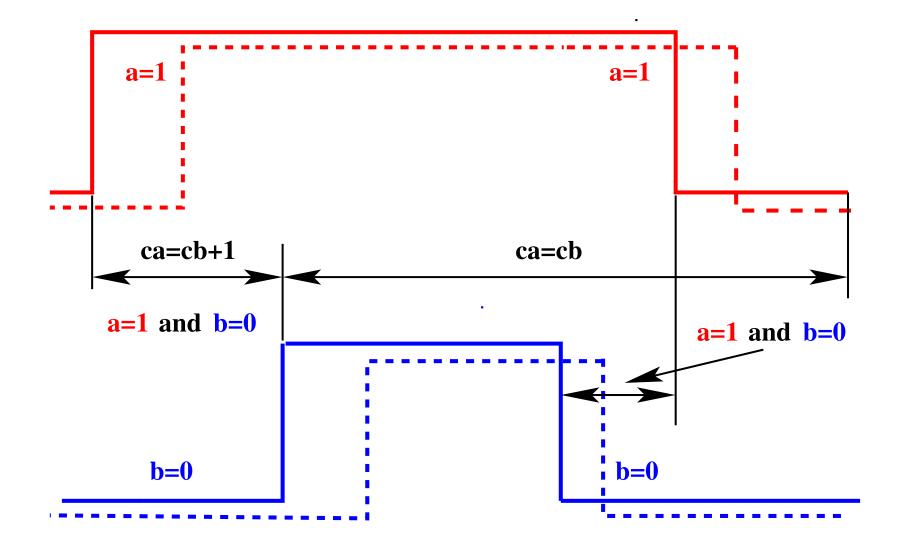
$$egin{aligned} ca &= cb ⅇ & ca &= cb+1 \ & cr &= cs ⅇ & cr &= cs+1 \end{aligned}$$

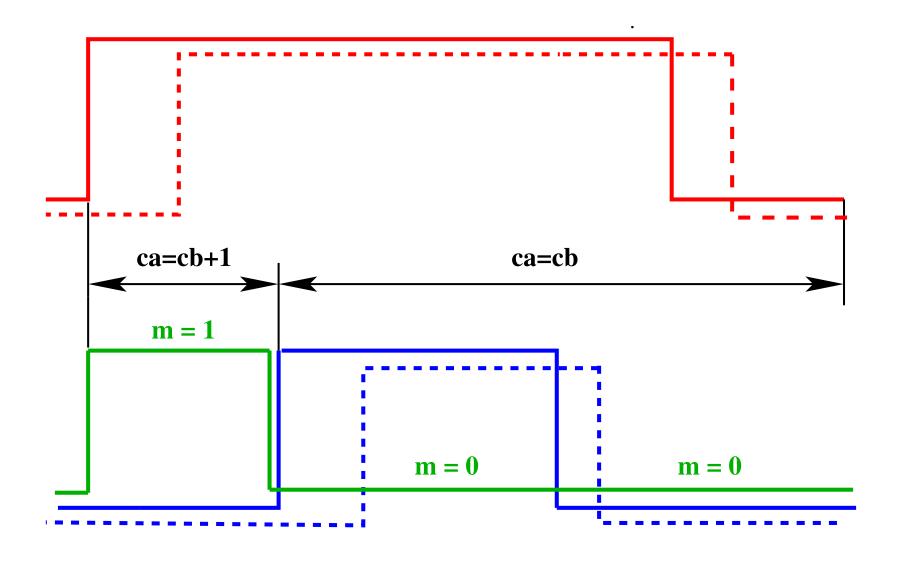




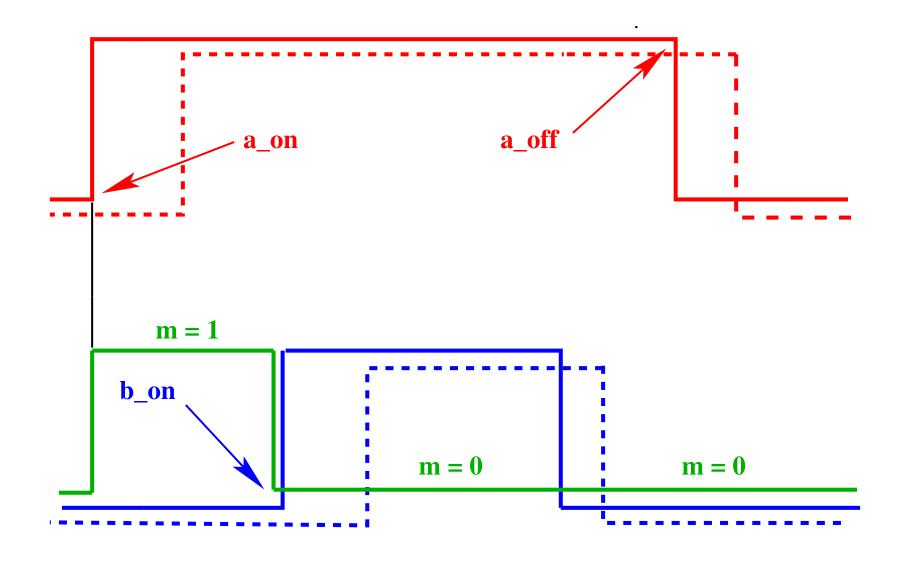


$$a=1 \land b=0 \Rightarrow ca=cb+1$$



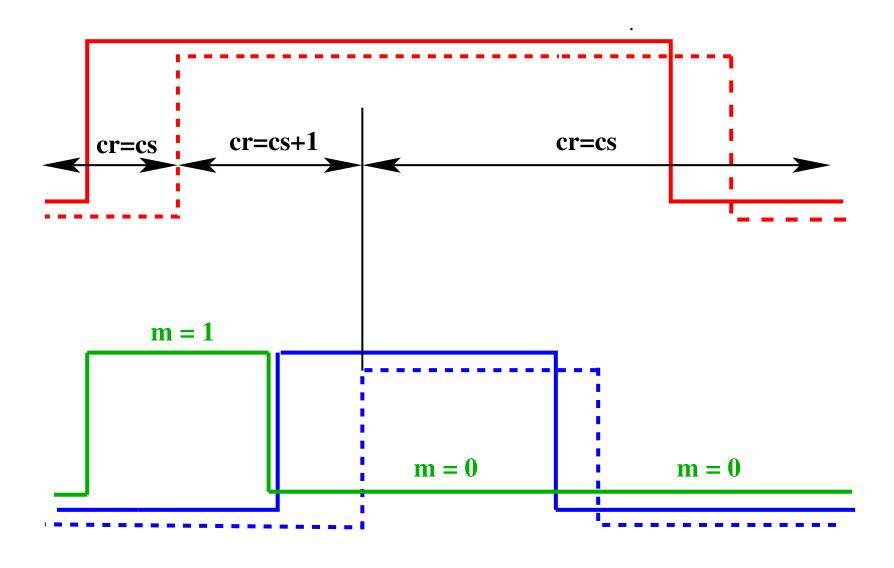


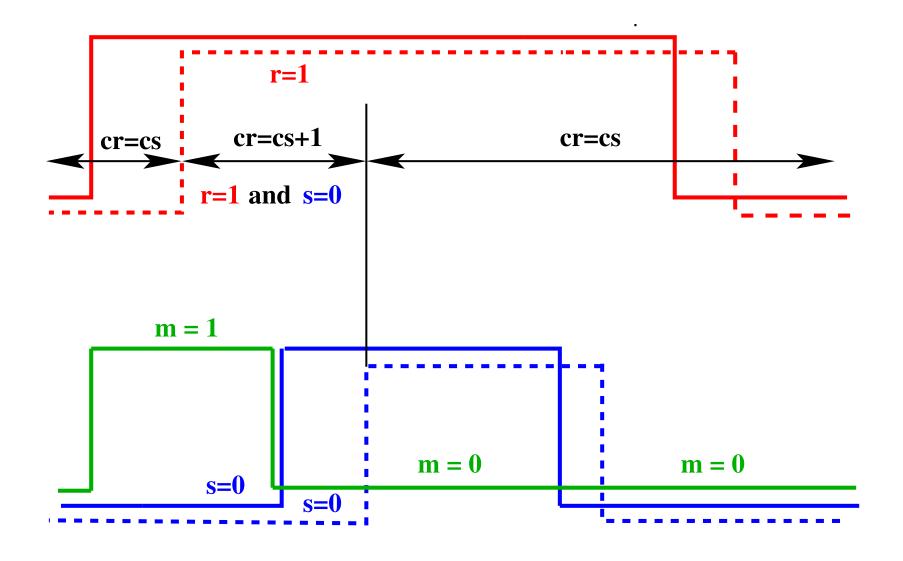
$$m=1 \Rightarrow ca=cb+1 \ m=0 \Rightarrow ca=cb$$

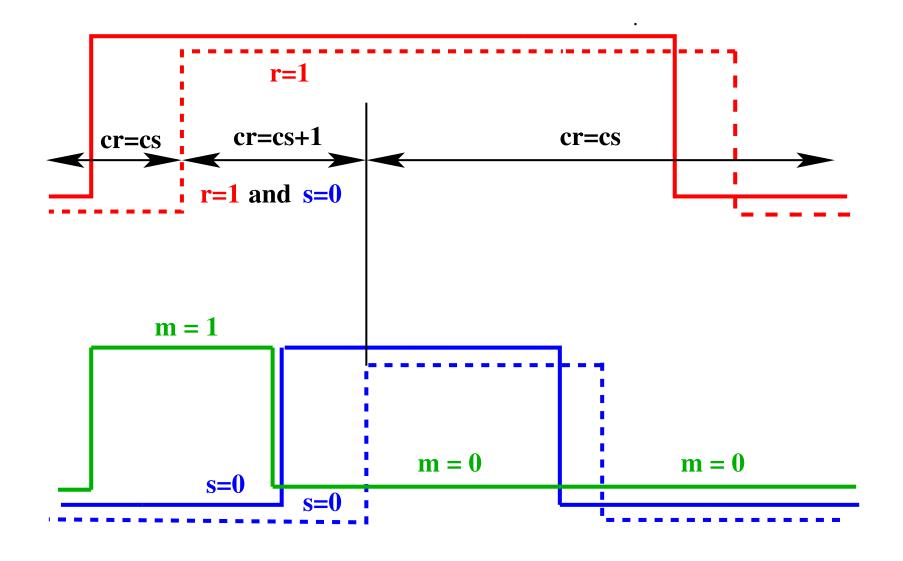


```
a_on when a=0 r=0 then a:=1 ca:=ca+1 m:=1 end
```

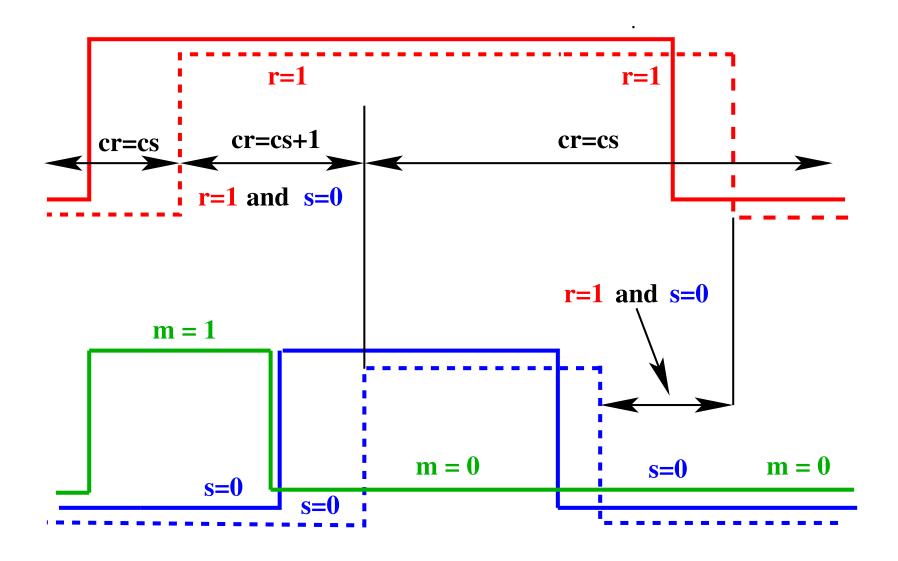
```
b_on
when
 r = 1
 a = 1
 b = 0
 s = 0
 m = 1
then
 b := 1
 cb := cb + 1
 m := 0
end
```

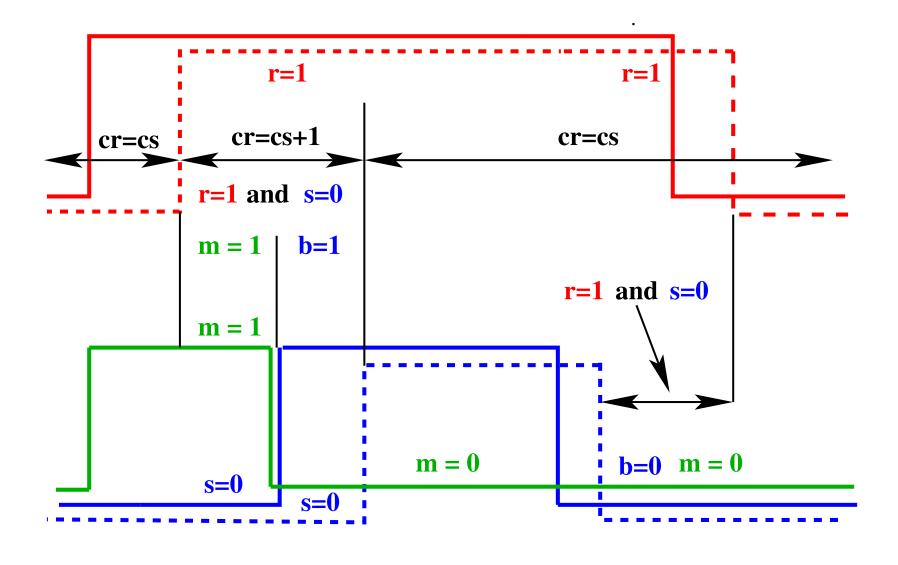




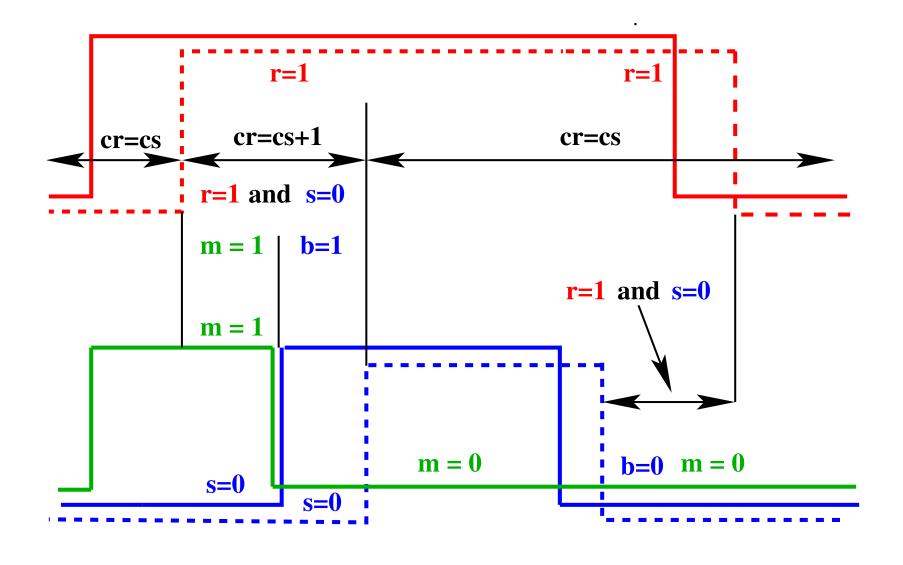


$$r=1 \ \land \ s=0 \ \Rightarrow \ cr=cs+1$$





**The Solution** 



$$egin{aligned} r=1 \; \wedge \; s=0 \; \wedge \; (m=1 \; ee \; b=1) \; \Rightarrow \; cr=cs+1 \ r=0 \; ee \; s=1 \; ee \; (m=0 \; \wedge \; b=0) \; \Rightarrow \; cr=cs \end{aligned}$$

```
dbl2_1: m \in \{0,1\}
```

**dbl2\_2:** 
$$m=1 \Rightarrow ca=cb+1$$

dbl2\_3: 
$$m=0 \Rightarrow ca=cb$$

dbl2\_4: 
$$r=1 \land s=0 \land (m=1 \lor b=1) \Rightarrow cr=cs+1$$

**dbl2\_5:** 
$$r=0 \ \lor \ s=1 \ \lor \ (m=0 \ \land \ b=0) \ \Rightarrow \ cr=cs$$

**dbl2**<sub>-</sub>1: 
$$m \in \{0,1\}$$

dbl2\_2: 
$$m=1 \Rightarrow ca=cb+1$$

dbl2\_3: 
$$m=0 \Rightarrow ca=cb$$

dbl2\_4: 
$$r=1 \land s=0 \land (m=1 \lor b=1) \Rightarrow cr=cs+1$$

**dbl2\_5:** 
$$r=0 \ \lor \ s=1 \ \lor \ (m=0 \ \land \ b=0) \ \Rightarrow \ cr=cs$$

- The following theorems are easy to prove

thm2\_1: 
$$ca = cb \lor ca = cb + 1$$

thm2\_2: 
$$cr = cs \lor cr = cs + 1$$

More Invariant 244

```
dbl2_1: m \in \{0, 1\}

dbl2_2: m = 1 \Rightarrow ca = cb + 1

dbl2_3: m = 0 \Rightarrow ca = cb

dbl2_4: r = 1 \land s = 0 \land (m = 1 \lor b = 1) \Rightarrow cr = cs + 1

dbl2_5: r = 0 \lor s = 1 \lor (m = 0 \land b = 0) \Rightarrow cr = cs

dbl2_6: a = 0 \Rightarrow m = 0
```

- The last new invariants was discovered while doing the proof
- It requires adding the guard m=0 in event a\_off
- The proofs are now completely automatic

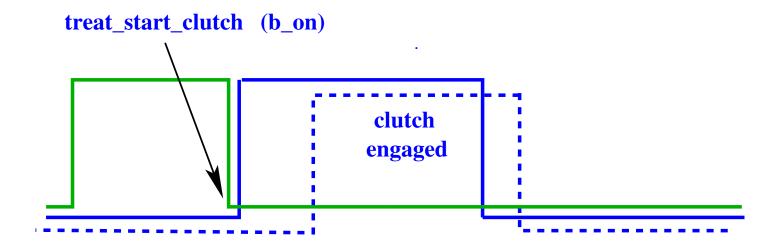
```
a_on when a=0 r=0 then a:=1 ca:=ca+1 m:=1 end
```

```
b_on
when
 r = 1
 a = 1
 b = 0
 s = 0
 m = 1
then
 b := 1
 cb := cb + 1
 m := 0
end
```

```
a_off when a=1 r=1 b=0 s=0 m=0 then a:=0 end
```

Instantiation





- We instantiate the pattern as follows:

```
a_on  

a_off  

treat_close_door  

treat_open_door  

b_on  

treat_start_clutch
```

```
a_on when a=0 r=0
then a:=1 m:=1 end
```

```
egin{array}{ll} {
m treat\_close\_door} & {
m when} & door\_actuator = open & door\_sensor = open & motor\_actuator = working & motor\_sensor = working & {
m then} & door\_actuator := closed & m := 1 & {
m end} & {
m end} & {
m end} & {
m constant} & {
m con
```

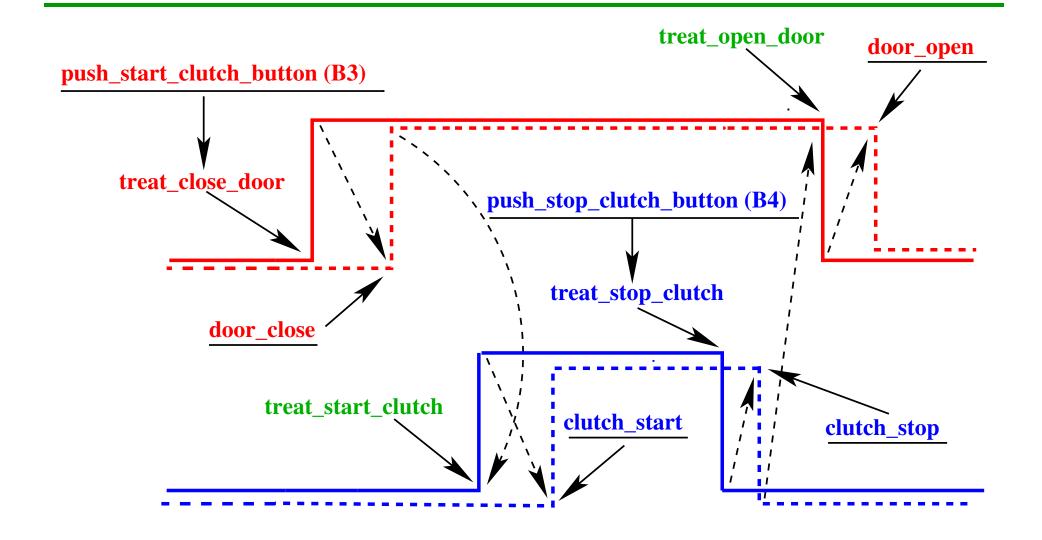
```
b_on
when
```

```
egin{aligned} b &= 0 \ s &= 0 \ r &= 1 \ a &= 1 \ m &= 1 \ then \ b &:= 1 \ m &:= 0 \ end \end{aligned}
```

```
treat_start_clutch
 when
   motor\_actuator = working
   motor\_sensor = working
   clutch\_actuator = disengaged
   clutch\_sensor = disengaged
   door\_sensor = closed
   door\_actuator = closed
   m = 1
 then
   clutch\_actuator := engaged
   m := 0
 end
```

```
a_off when a=1 r=1 s=0 b=0 m=0 then a:=0 end
```

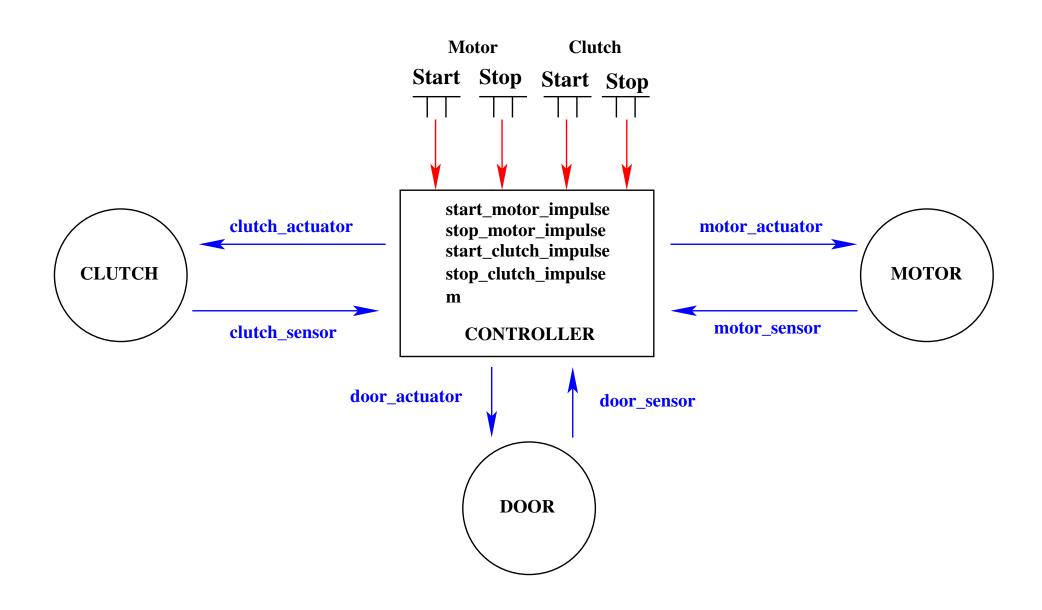
```
egin{array}{ll} {\sf treat\_open\_door} & {\sf when} \\ & door\_actuator = closed \\ & door\_sensor = closed \\ & clutch\_sensor = disengaged \\ & clutch\_actuator = disengaged \\ & m = 0 \\ {\sf then} \\ & door\_actuator := open \\ {\sf end} \\ \end{array}
```



- treat\_close\_door is the result of depressing button B3
- treat\_stop\_clutch is the result of depressing button B4
- treat\_start\_clutch and treat\_open\_door are automatic

- Environment (no new events)
  - motor\_start
  - motor\_stop
  - clutch\_start
  - clutch\_stop
  - door\_close
  - door\_open
  - push\_start\_motor\_button
  - release\_start\_motor\_button
  - push\_stop\_motor\_button
  - release\_stop\_motor\_button

- Controller (no new events)
  - treat\_push\_start\_motor\_button
  - treat\_push\_start\_motor\_button\_false
  - treat\_push\_stop\_motor\_button
  - treat\_push\_stop\_motor\_button\_false
  - treat\_release\_start\_motor\_button
  - treat\_release\_stop\_motor\_button
  - treat\_start\_clutch
  - treat\_stop\_clutch
  - treat\_close\_door
  - treat\_open\_door



- There are no door buttons

- The door must be closed before engaging the clutch

- The door must be opened after disengaging the clutch

- It is sufficient to connect:
  - button B3 to the door (closing the door)
  - button B4 to the clutch (disengaging the clutch)

- motor\_start
- motor\_stop
- clutch\_start
- clutch\_stop
- door\_close
- door\_open
- push\_start\_motor\_button
- release\_start\_motor\_button
- push\_stop\_motor\_button
- release\_stop\_motor\_button
- push\_start\_clutch\_button
- release\_start\_clutch\_button
- push\_stop\_clutch\_button
- release\_stop\_clutch\_button

- treat\_push\_start\_motor\_button
- treat\_push\_start\_motor\_button\_false
- treat\_push\_stop\_motor\_button
- treat\_push\_stop\_motor\_button\_false
- treat\_release\_start\_motor\_button
- treat\_release\_stop\_motor\_button
- treat\_start\_clutch
- treat\_stop\_clutch
- treat\_close\_door
- treat\_open\_door
- treat\_close\_door\_false
- treat\_stop\_clutch\_false
- treat\_release\_start\_clutch\_button
- treat\_release\_stop\_clutch\_button

- The environment events

- The environment variables modified by environment events

- The sensor variables modified by environment events

- The actuator variables read by environment events

- The controller variables not seen by environment events

- No environment variables in this model

- The controller events

- The controller variables modified by controller events

- The sensor variables read by controller events

- The actuator variables modified by controller events

- The environment variables not seen by controller events

- No environment variables in this model

- 7 sensor variables:
  - $motor\_sensor$
  - $clutch\_sensor$
  - $door\_sensor$
  - $start\_motor\_button$
  - $-stop\_motor\_button$
  - $start\_clutch\_button$
  - $-stop\_clutch\_button$

- 3 actuator variables:
  - $motor\_actuator$
  - $clutch\_actuator$
  - $door\_actuator$
- 5 controller variables (without the counter variables):
  - $start\_motor\_impulse$
  - $stop\_motor\_impulse$
  - $-start\_clutch\_impulse$
  - $-stop\_clutch\_impulse$
  - *m*

- 14 environment events,

- 14 controller events,

- 130 lines for environment events,

- 180 lines for controller events.

- 4 weak reactions: 4 buttons (B1, B2, B3, B4)

- 3 strong reactions: 3 devices (motor, clutch, door)

- 3 strong-weak reactions: motor-clutch, clutch-door, motor-door

- 1 strong-strong reaction: clutch-door

- Weak reaction: 6
- Strong reaction: 3
- Strong-weak reaction: 16
- Strong-strong reaction: 6
- Total: 31

- Press (typing): 15
- Total: 15

- Weak reaction: 18
- Strong reaction: 12
- Strong-weak reaction: 61
- Strong-strong reaction: 35
- Total: 126

- Press: 0

- PO saving: 4x18 + 3x12 + 3x61 + 35 = 326

- Design patterns: 126 (all automatic)

- Press: 0

- This design pattern approach is very fruitful

- It results in a very systematic formal development

- Many other patterns have to be developed

- More automation has to be provided (plug-in)