

## Search element in a sorted matrix

3.4

Given a sorted matrix  $\text{mat}[n][m]$  and an element 'x'. Find position of x in the matrix if it is present, else print -1. Matrix is sorted in a way such that all elements in a row are sorted in increasing order and for row 'i', where  $1 \leq i \leq n-1$ , first element of row 'i' is greater than or equal to the last element of row 'i-1'. The approach should have  $O(\log n + \log m)$  time complexity.

Examples:

```
Input : mat[][] = { {1, 5, 9},
                   {14, 20, 21},
                   {30, 34, 43} }
```

```
      x = 14
```

```
Output : Found at (1, 0)
```

```
Input : mat[][] = { {1, 5, 9, 11},
                   {14, 20, 21, 26},
                   {30, 34, 43, 50} }
```

```
      x = 42
```

```
Output : -1
```

**Recommended:** Please try your approach on [{IDE}](#) first, before moving on to the solution.

Please note that this problem is different from [Search in a row wise and column wise sorted matrix](#). Here matrix is more strictly sorted as first element of a row is greater than last element of previous row.

A **Simple Solution** is to one by one compare x with every element of matrix. If matches, then return position. If we reach end, return -1. Time complexity of this solution is  $O(n \times m)$ .

An **efficient solution** is to typecast given 2D array to 1D array, then apply [binary search](#) on the typecasted array.

**Another efficient approach** that doesn't require typecasting is explained below.

- 1) Perform binary search on the middle column till only two elements are left or till the middle element of some row in the search is the required element 'x'. This search is done to skip the rows that are not required
- 2) The two left elements must be adjacent. Consider the rows of two elements and do following
  - a) check whether the element 'x' equals to the middle element of any one of the 2 rows
  - b) otherwise according to the value of the element 'x' check whether it is present in the 1st half of 1st row, 2nd half of 1st row, 1st half of 2nd row or 2nd half of 2nd row.

Note: This approach works for the matrix  $n \times m$  where  $2 \leq n$ . The algorithm can be modified for matrix  $1 \times m$ , we just need to check whether 2nd row exists or not

### Example:

Consider:     | 1  2  3  4 |  
 x = 3, mat = | 5  6  7  8 |     Middle column:  
               | 9 10 11 12 |     = {2, 6, 10, 14}  
               |13 14 15 16 |     perform binary search on them  
                                   since,  $x < 6$ , discard the  
                                   last 2 rows as 'a' will  
                                   not lie in them(sorted matrix)

Now, only two rows are left

              | 1  2  3  4 |  
 x = 3, mat = | 5  6  7  8 |     Check whether element is present  
                                   on the middle elements of these  
                                   rows = {2, 6}  
                                    $x \neq 2$  or 6

If not, consider the four sub-parts

1st half of 1st row = {1}, 2nd half of 1st row = {3, 4}  
 1st half of 2nd row = {5}, 2nd half of 2nd row = {7, 8}

According the value of 'x' it will be searched in the  
 2nd half of 1st row = {3, 4} and found at (i, j): (0, 2)

```
// C++ implementation to search an element in a
// sorted matrix
#include <bits/stdc++.h>
using namespace std;

const int MAX = 100;

// This function does Binary search for x in i-th
// row. It does the search from mat[i][j_low] to
// mat[i][j_high]
void binarySearch(int mat[][MAX], int i, int j_low,
                  int j_high, int x)
```

```

{
    while (j_low <= j_high)
    {
        int j_mid = (j_low + j_high) / 2;

        // Element found
        if (mat[i][j_mid] == x)
        {
            cout << "Found at (" << i << ", "
                << j_mid << ")";
            return;
        }

        else if (mat[i][j_mid] > x)
            j_high = j_mid - 1;

        else
            j_low = j_mid + 1;
    }

    // element not found
    cout << "Element no found";
}

// Function to perform binary search on the mid
// values of row to get the desired pair of rows
// where the element can be found
void sortedMatrixSearch(int mat[][MAX], int n,
                        int m, int x)
{
    // Single row matrix
    if (n == 1)
    {
        binarySearch(mat, 0, 0, m-1, x);
        return;
    }

    // Do binary search in middle column.
    // Condition to terminate the loop when the
    // 2 desired rows are found
    int i_low = 0;
    int i_high = n-1;
    int j_mid = m/2;
    while ((i_low+1) < i_high)
    {
        int i_mid = (i_low + i_high) / 2;

        // element found
        if (mat[i_mid][j_mid] == x)
        {
            cout << "Found at (" << i_mid << ", "
                << j_mid << ")";
            return;
        }

        else if (mat[i_mid][j_mid] > x)
            i_high = i_mid;

        else
            i_low = i_mid;
    }

    // If element is present on the mid of the
    // two rows
    if (mat[i_low][j_mid] == x)
        cout << "Found at (" << i_low << ", "
            << j_mid << ")";
    else if (mat[i_low+1][j_mid] == x)
        cout << "Found at (" << (i_low+1)
            << ", " << j_mid << ")";
}

```

```

// Ssearch element on 1st half of 1st row
else if (x <= mat[i_low][j_mid-1])
    binarySearch(mat, i_low, 0, j_mid-1, x);

// Search element on 2nd half of 1st row
else if (x >= mat[i_low][j_mid+1] &&
        x <= mat[i_low][m-1])
    binarySearch(mat, i_low, j_mid+1, m-1, x);

// Search element on 1st half of 2nd row
else if (x <= mat[i_low+1][j_mid-1])
    binarySearch(mat, i_low+1, 0, j_mid-1, x);

// search element on 2nd half of 2nd row
else
    binarySearch(mat, i_low+1, j_mid+1, m-1, x);
}

// Driver program to test above
int main()
{
    int n = 4, m = 5, x = 8;
    int mat[][MAX] = {{0, 6, 8, 9, 11},
                      {20, 22, 28, 29, 31},
                      {36, 38, 50, 61, 63},
                      {64, 66, 100, 122, 128}};

    sortedMatrixSearch(mat, n, m, x);
    return 0;
}

```

[Run on IDE](#)

Output:

Found at (2, 1)

Time complexity:  $O(\log n + \log m)$ .  $O(\log n)$  time is required to find the two desired rows. Then  $O(\log m)$  time is required for binary search in one of the four parts with size equal to  $m/2$ .

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