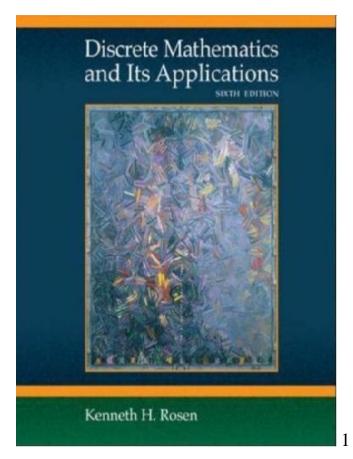


Jiangxi University of Science and Technology

Discrete Mathematics and Its Applications

Lecture012: CHP2:

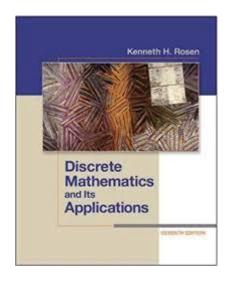
Sequences and Summations





Acknowledgement

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Summation

• The symbol \sum (Greek letter sigma) is used to denote summation.

$$\sum_{i=1}^{n} a_i = a_1 + a_2 + \dots + a_k$$

i is the **index of the summation**, and the choice of letter *i* is arbitrary; the index of the summation runs through all integers, with its **lower limit** 1 and ending **upper limit** k.

• The limit:

$$\sum_{i=1}^{\infty} a_i = \lim_{n \to \infty} \sum_{i=1}^{n} a_i$$



Summation

• The laws for arithmetic apply to summations

$$\sum_{i=1}^{k} (ca_i + b_i) = c \sum_{i=1}^{k} a_i + \sum_{i=1}^{k} b_i$$

Use associativity to separate the b terms from the a terms.

Use distributivity to factor the c's.



• <u>Definition</u>: A <u>sequence</u> is a function from a subset of the natural numbers (usually of the form $\{0, 1, 2, ...\}$ to a set S.

Note: the sets

$$\{0, 1, 2, 3, \dots, k\}$$
 and $\{1, 2, 3, 4, \dots, k\}$

are called *initial segments* of N.

Notation: if f is a function from $\{0, 1, 2, ...\}$ to S we usually denote f(i) by a_i and we write

$${a_0, a_1, a_2, ...} = {a_i}_{i=0}^k = {a_i}_0^k$$

where k is the upper limit (usually ∞).



Examples:

Using zero-origin indexing, if f(i) = 1/(i + 1). then the Sequence

$$f = \{1, 1/2, 1/3, 1/4, \dots\} = \{a_0, a_1, a_2, a_3, \dots\}$$

Using one-origin indexing the sequence f becomes $\{1/2, 1/3, \ldots\} = \{a_1, a_2, a_3, \ldots\}$

• Summation Notation

$$\{a_i\}_0^k$$

Given a sequence we can add together a subset of the sequence by using the summation and function notation

$$a_{g(m)} + a_{g(m+1)} + \dots + a_{g(n)} = \sum_{j=m}^{n} a_{g(j)}$$

or more generally

$$\sum_{j \in S} a_j$$



Examples:
$$r^{0} + r^{1} + r^{2} + ... + r^{n} = \sum_{j=0}^{n} r^{j}$$

$$1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + ... + \sum_{j=1}^{\infty} \frac{1}{i}$$

$$a_{2m} + a_{2(m+1)} + ... + a_{2(n)} = \sum_{j=m}^{n} a_{2j}$$
if $S = \{2, 5, 7, 10\}$ then $\sum_{j \in S} a_{j} = a_{2} + a_{5} + a_{7} + a_{10}$

Similarity for the *product* notation:



$$\prod_{j=m}^{n} a_{j} = a_{m} a_{m+1} \dots a_{n}$$

Definition: A *geometric progression* is a sequence of the form $a, ar, ar^2, ar^3, ar^4, \ldots$

Your book has a proof that

$$\sum_{i=0}^{n} r^{i} = \frac{r^{n+1} - 1}{r - 1} \text{if } r \neq 1$$

(you can figure out what it is if r = 1).

You should also be able to determine the sum

- if the index starts at k vs. 0
- if the index ends at something other than n (e.g., n-1, n+1, etc.).



• Cardinality **Definition**::

The cardinality of a set A is equal to the cardinality of a set B, denoted |A| = |B|,

if there exists a bijection from A to B.



Definition:

- -If a set has the same cardinality as a subset of the natural numbers N, then the set is called *countable*.
- If |A| = |N|, the set A is *countably infinite*.
- The (transfinite) cardinal number of the set N is $aleph null = \aleph_0$.
- If a set is not countable we say it is *uncountable*.



Examples:

The following sets are uncountable (we show later)

The real numbers in [0, 1]

- P(N), the power set of N
- Note: With infinite sets proper subsets can have the same cardinality. This cannot happen with finite sets.
- Countability carries with it the implication that there is a *listing* of the elements of the set.



- **Definition:** $|A| \le |B|$ if there is an injection from A to B.
 - Note: as you would hope,
- **Theorem:** If $|A| \le |B|$ and $|B| \le |A|$ then |A| = |B|.

This implies

- if there is an injection from A to B
- if there is an injection from B to A then
- there must be a bijection from A to B

This is <u>difficult</u> to prove but is an example of demonstrating existence without construction.

It is often easier to build the injections and then conclude the bijection exists.



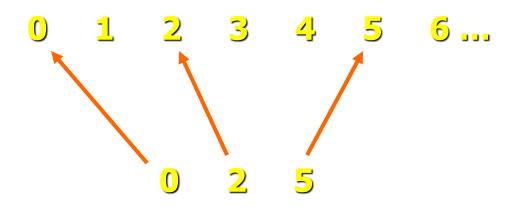
- Example:

Theorem: If A is a subset of B then $|A| \le |B|$.

Proof: the function f(x) = x is an injection from A to B.

Example: $\{0, 2, 5\}$ | ≤ \aleph_0

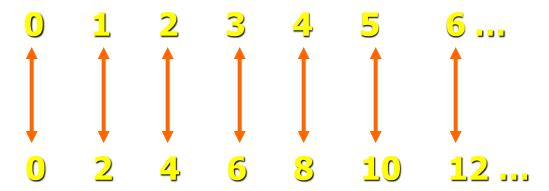
The injection f: $\{0, 2, 5\} \rightarrow N$ defined by f(x) = x is shown below:





- Some Countably Infinite Sets
 - The set of even integers E (0 is considered even) is countably infinite. Note that E is a proper subset of N,

Proof: Let f(x) = 2x. Then f is a bijection from N to E



 $-Z^{+}$, the set of positive integers is countably infinite.



- The set of positive rational numbers Q⁺ is countably infinite.
 - Proof: Z^+ is a subset of Q^+ so $|Z^+| = \aleph_0 \le |Q^+|$.
- Now we have to show that $|Q^+| \le \aleph_0$.
- To do this we show that the positive rational numbers with repetitions, Q_R , is countably infinite.

Then, since Q^+ is a subset of Q_R , it follows that $|Q^+| \le \aleph_0$ and hence $|Q^+| = \aleph_0$.

y X	1	2	3	4	5	6	7	
1	1/1	2/1-	-3/1	4/1	-5/1	6/1	7/1	
2	1/2	2/2	3/2	4/2	5/2	6/2	7/2	
3	1/3	2/3	3/3	4/3	5/3	6/3	7/3	
4	1/4	2/4	3/4	4/4	5/4	6/4	7/4	
5								



- The position on the path (listing) indicates the image of the bijective function f from N to Q_R :

$$f(0) = 1/1$$
, $f(1) = 1/2$, $f(2) = 2/1$, $f(3) = 3/1$, and so forth.

Every rational number appears on the list at least once, some many times (repetitions).

Hence,
$$|N| = |QR| = 80$$
. Q. E. D

- The set of all rational numbers Q, positive and negative, is countably infinite.



 The set of (finite length) strings S over a finite alphabet A is countably infinite.

To show this we assume that

- A is nonvoid
- There is an "alphabetical" ordering of the symbols in A

Proof: List the strings in lexicographic order:

- all the strings of zero length,
- then all the strings of length 1 in alphabetical order,
- then all the strings of length 2 in alphabetical order, etc

This implies a bijection from N to the list of strings and hence it is a countably infinite set.



– For example:

Let
$$A = \{a, b, c\}.$$

Then the lexicographic ordering of A is

 $\{\lambda, a, b, c, aa, ab, ac, ba, bb, bc, ca, cb, cc, aaa, aab, aac, aba,\} = \{f(0), f(1), f(2), f(3), f(4),\}$

- The set of all C programs is countable.
- Proof: Let S be the set of legitimate characters which can appear in a C program.
 - A C compiler will determine if an input program is a syntactically correct C program (the program doesn't have to do anything useful).
 - Use the lexicographic ordering of S and feed the strings into the compiler.
 - If the compiler says YES, this is a syntactically correct C program, we add the program to the list.
 - Else we move on to the next string.
- In this way we construct a list or an implied bijection from N to the set of C programs.
- Hence, the set of C programs is countable.

Q. E. D.



Cantor Diagonalization

An important technique used to construct an object which is not a member of a countable set of objects with (possibly) infinite descriptions

Theorem: The set of real numbers between 0 and 1 is uncountable.

Proof: We assume that it is countable and derive a contradiction.

If it is countable we can list them

(i.e., there is a bijection from a subset of N to the set).

We show that no matter what list you produce we can construct a real number between 0 and 1 which is not in the list.

Hence, there cannot exist a list and therefore the set is not countable



It's actually much bigger than countable. It is said to have the *cardinality of the continuum*, c.

Represent each real number in the list using its decimal expansion.

e.g.,
$$1/3 = .333333333...$$

 $1/2 = .5000000...$
 $= .4999999...$

If there is more than one expansion for a number, it doesn't matter as long as our construction takes this into account.



THE LIST....

$$r_1 = .d_{11}d_{12}d_{13}d_{14}d_{15}d_{16}.....$$

$$r_2 = .d_{21}d_{22}d_{23}d_{24}d_{25}d_{26}....$$

$$r_3 = .d_{31}d_{32}d_{33}d_{34}d_{35}d_{36}....$$

Now construct the number $x = .x_1x_2x_3x_4x_5x_6x_7...$

$$x_i = 3 \text{ if } d_{ii} \neq 3$$

 $x_i = 4 \text{ if } d_{ii} = 3$

(Note: choosing 0 and 9 is not a good idea because of the non uniqueness of decimal expansions.)

Then x is not equal to any number in the list.

Hence, no such list can exist and hence the interval (0,1) is uncountable. Q. E. D.



An extra goody:

Definition: a number x between 0 and 1 is *computable* if there is a C program which when given the input i, will produce the ith digit in the decimal expansion of x.

Example:

The number 1/3 is computable.

The C program which always outputs the digit 3,regardless if the input, computes the number.



Theorem: There is exists a number x between 0 and 1 which is *not* computable.

There *does not exist* a C program (or a program in any other language) which will compute it!

Why? Because there are more numbers between 0 and 1 than there are C programs to compute them.

(in fact there are c such numbers!)

Our second example of the *nonexistence* of programs to compute things!



Useful Summations

Sum	Closed Form			
$\sum_{k=0}^{n} ar^k \ (r \neq 0)$	$\frac{ar^{n+1}-a}{r-1}, r \neq 1$			
$\sum_{k=1}^{n} k$	$\frac{n(n+1)}{2}$			
$\sum_{k=1}^{n} k^2$	$\frac{n(n+1)(2n+1)}{6}$			
$\sum_{k=1}^{n} k^3$	$\frac{n^2(n+1)^2}{4}$			
$\sum_{k=0}^{\infty} x^k, x < 1$	$\frac{1}{1-x}$			
$\sum_{k=1}^{\infty}, kx^{k-1}, x < 1$	$\frac{1}{(1-x)^2}$			



Infinite Cardinality

- How can we extend the notion of cardinality to infinite sets?
- Definition: Two sets **A** and **B** have the same cardinality if and only if there exists a bijection (or a one-to-one correspondence) between them, A ~ B.

We split infinite sets into two groups:

- 1. Sets with the same cardinality as the set of natural numbers
- 2. Sets with different cardinality as the set of natural numbers



Infinite Cardinality

- Definition: A set is **countable** if it is **finite** or has the same **cardinality as the set of positive integers.**
- Definition: A set is uncountable if it is not countable.
- Definition: The cardinality of an infinite set S that is countable is denotes by \aleph_0 (where \aleph is aleph, the first letter of the Hebrew alphabet).
- We write $|S| = \aleph_0$ and say that S has cardinality "aleph null"

Note: Georg Cantor defined the notion of cardinality and was the first to realize that infinite sets can have different cardinalities. \aleph_0 is the cardinality of the natural numbers; the next larger cardinality is aleph-one \aleph_1 , then, \aleph_2 and so on.



Infinite Cardinality: Odd Positive Integers

Example: The set of odd positive integers is a countable set.

Let's define the function f, from Z^+ to the set of odd positive numbers,

$$f(n) = 2 n - 1$$

We have to show that f is both one-to-one and onto.

one-to-one

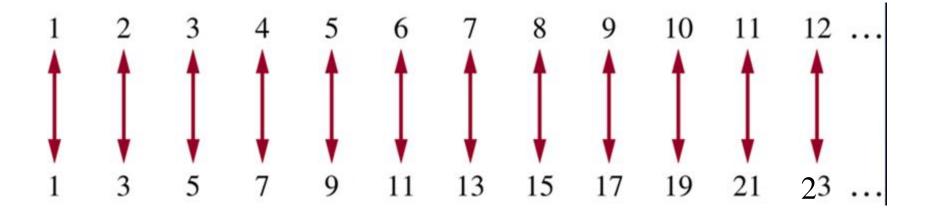
Suppose
$$f(n)=f(m) \rightarrow 2n-1=2m-1 \rightarrow n=m$$

onto

Suppose that t is an odd positive integer. Then t is 1 less than an even integer 2k, where k is a natural number. hence t=2k-1=f(k).



Infinite Cardinality: Odd Positive Integers





Infinite Cardinality:Integers

Example: The set of integers is a countable set.

Lets consider the sequence of all integers, starting with 0: 0,1,-1,2,-2,....

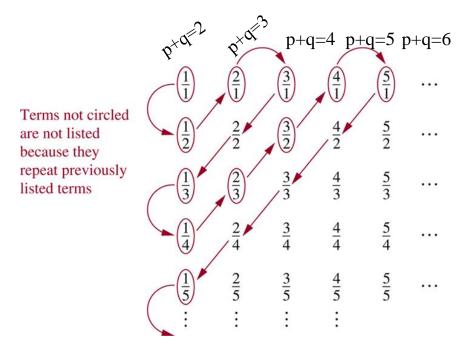
Infinite Cardinality: Rational Numbers

- Example: The set of **positive rational numbers** is a **countable** set.
- Hmm...

Infinite Cardinality: Rational Numbers

Example: The set of **positive rational numbers** is a **countable** set Key aspect to list the rational numbers as a sequence – every positive number is the quotient p/q of two positive integers.

Visualization of the proof.



Since all positive rational numbers are listed once, the set of positive rational numbers is countable.

The set of all infinite sequences of zeros and ones is uncountable. Consider a sequence,

$$a_1, a_2, \dots, a_n, n \rightarrow \infty, a_i = 0 \text{ or } a_i = 1$$

For example:
$$s_1 = (0, 0, 0, 0, 0, 0, 0, ...)$$

 $s_2 = (1, 1, 1, 1, 1, 1, 1, 1, ...)$
 $s_3 = (0, 1, 0, 1, 0, 1, 0, ...)$
 $s_4 = (1, 0, 1, 0, 1, 0, 1, ...)$
 $s_5 = (1, 1, 0, 1, 0, 1, 1, ...)$
 $s_6 = (0, 0, 1, 1, 0, 1, 1, ...)$
 $s_7 = (1, 0, 0, 0, 1, 0, 0, ...)$

So in general we have:

$$s_n = (s_{n,1}, s_{n,2}, s_{n,3}, s_{n,4}, ...)$$

i.e., $s_{n,m}$ is the mth element of the nth sequence on the list.

• It is possible to build a sequence, say s_0 , in such a way that its first element is different from the first element of the first sequence in the list, its second element is different from the second element of the second sequence in the list, and, in general, its nth element is different from the nth element of the nth sequence in the list. In other words, $s_{0,m}$ will be 0 if $s_{m,m}$ is 1, and $s_{0,m}$ will be 1 if $s_{m,m}$ is 0.



$$\begin{split} s_1 &= (\underline{\mathbf{0}}, 0, 0, 0, 0, 0, 0, ...) \\ s_2 &= (1, \underline{\mathbf{1}}, 1, 1, 1, 1, 1, 1, ...) \\ s_3 &= (0, 1, \underline{\mathbf{0}}, 1, 0, 1, 0, ...) \\ s_4 &= (1, 0, 1, \underline{\mathbf{0}}, 1, 0, 1, ...) \\ s_5 &= (1, 1, 0, 1, \underline{\mathbf{0}}, 1, 1, ...) \\ s_6 &= (0, 0, 1, 1, 0, \underline{\mathbf{1}}, 1, ...) \\ s_7 &= (1, 0, 0, 0, 1, 0, \underline{\mathbf{0}}, ...) \\ ... \\ s_0 &= (\underline{\mathbf{1}}, \underline{\mathbf{0}}, \underline{\mathbf{1}}, \underline{\mathbf{1}}, \underline{\mathbf{1}}, \underline{\mathbf{0}}, \underline{\mathbf{1}}, ...) \end{split}$$

Note: the diagonal elements are highlighted, showing why this is called the **diagonal** argument

- The sequence s_0 is distinct from all the sequences in the list. Why?
- Let's say that s_0 is identical to the 100th sequence, therefore, $s_{0,100} = s_{100,100}$.
- In general, if it appeared as the *n*th sequence on the list, we would have $s_{0,n} = s_{n,n}$,
- which, due to the construction of s_0 , is impossible.



From this it follows that the set T, consisting of all infinite sequences of zeros and ones, cannot be put into a list s_1 , s_2 , s_3 , ...

Otherwise, it would be possible by the above process to construct a sequence s_0 which would both be in T (because it is a sequence of 0's and 1's which is by the definition of T in T) and at the same time not in T (because we can deliberately construct it not to be in the list). T, containing all such

sequences, must contain s_0 , which is just such a sequence. But since s_0 does not appear anywhere on the list, T cannot contain s_0 .

Therefore *T* cannot be placed in one-to-one correspondence with the natural numbers. In other words, the set of infinite binary strings is **uncountable**.



Real Numbers

Example;

The set of real numbers is an uncountable set.

Let's assume that the set of real numbers is countable.

Therefore any subset of it is also countable, in particular the interval [0,1].

How many real numbers are in interval [0, 1]?



Real Numbers

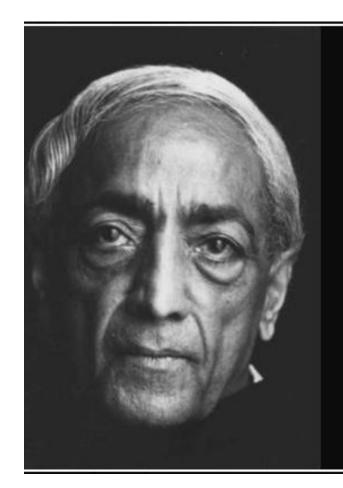
• How many real numbers are in interval [0, 1]?

0(4) 2 9 0 1 3 2 9 8 4 2 0 3 9 ... 0.8 2 3 9 9 1 3 2 7 2 5 8 9 2 5 ... 0.9 2 5 3 9 1 5 9 7 4 5 0 6 2 1 ... "Countably many! There's part of the list!"

Counterexample:
Use diagonalization
to create a new number
that differs in the ith
position of the
ith number
by 1.

"Are you sure they're all there?"

0.5 3 6 ...
So we say the reals are "uncountable."



There is no end to education. It is not that you read a book, pass an examination, and finish with education. The whole of life, from the moment you are born to the moment you die, is a process of learning.

— Jiddu Krishnamurti —

AZ QUOTES