

# CS480

# Translators

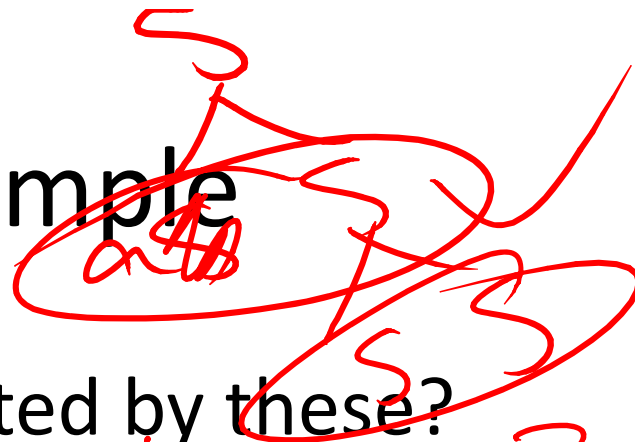
SDT and Intro to Parsing

Chap. 2

~~aaa~~



or



# Class Example

- What language is generated by these?

~~$S \rightarrow 0S1 \mid 01$~~

$L = \{0^m 1^m \mid m \geq 0\}$

~~$S \rightarrow S(S)S \mid \epsilon$~~

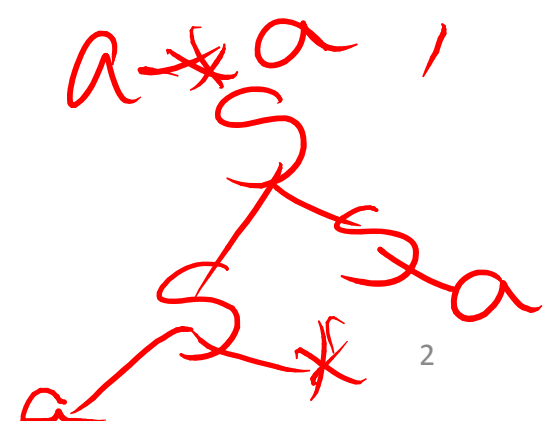
— balanced parens

$S \rightarrow a \mid SS \mid S^*$

$L = \{a^m \mid m \geq 0\}$

- Which are ambiguous?

~~aaa~~



# Syntax-Directed Translation

- Extend grammar
  - Attributes ✓ rules
  - Translation scheme — actions

# Synthesized Attributes

Opposite  
from inheritance

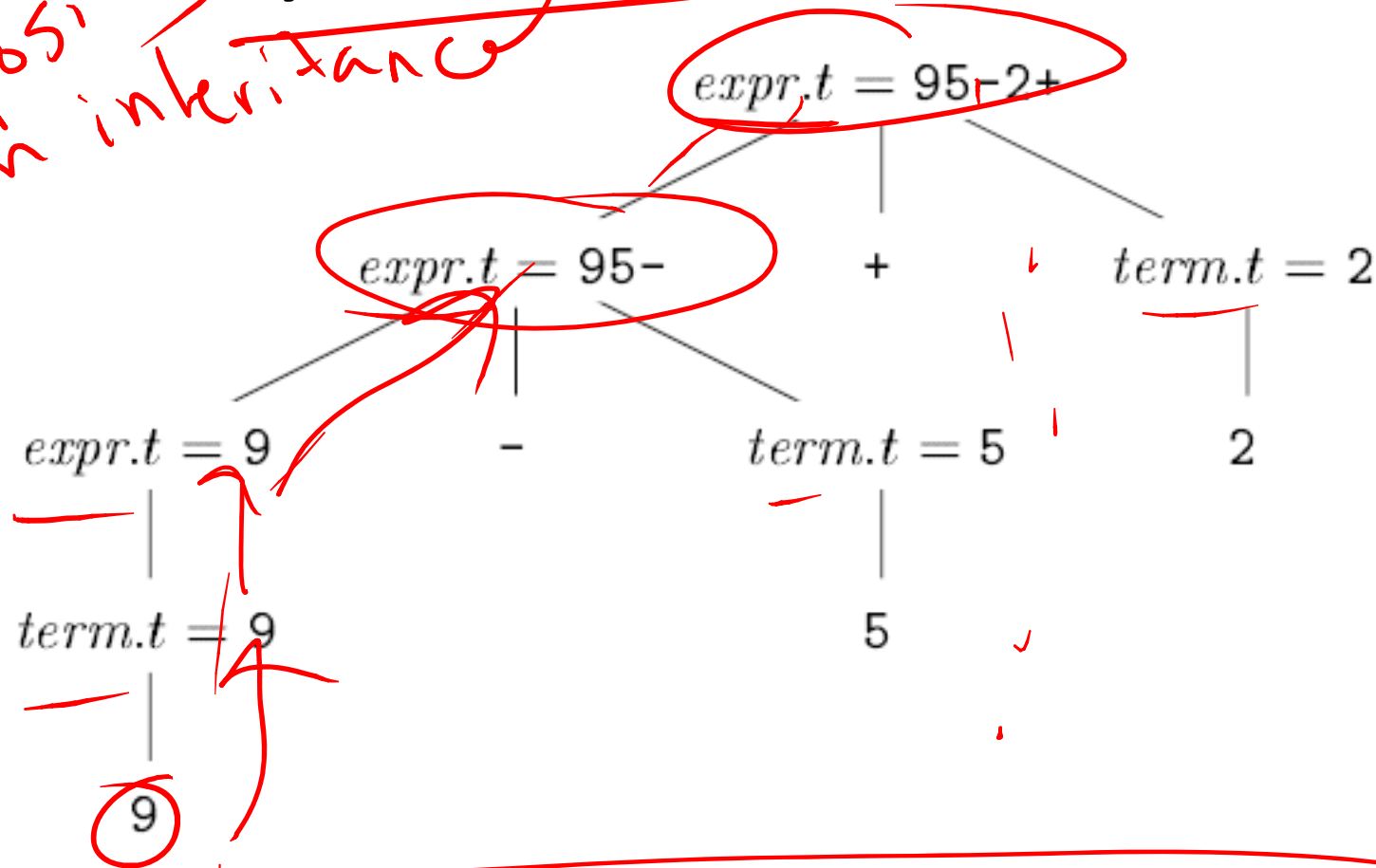


Figure 2.9: Attribute values at nodes in a parse tree

# Syntax-Directed Definition

PRODUCTION	SEMANTIC RULES
$expr \rightarrow expr_1 + term$	$expr.t = \cancel{expr_1.t} \parallel \cancel{term.t} \parallel '+'$
$expr \rightarrow expr_1 - term$	$expr.t = \cancel{expr_1.t} \parallel \cancel{term.t} \parallel '-'$
$expr \rightarrow term$	$expr.t = \underline{term.t}$
$term \rightarrow 0$	$term.t = '0'$
$term \rightarrow 1$	$term.t = '1'$
...	...
$term \rightarrow 9$	$term.t = '9'$

Figure 2.10: Syntax-directed definition for infix to postfix translation

# Syntax-Directed Translation Scheme

$rest \rightarrow + term \{ \text{print}(' + ') \} rest_1$

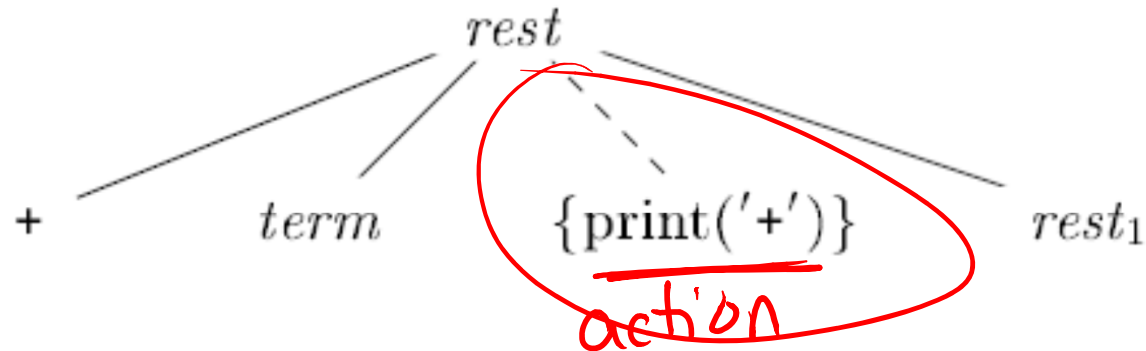


Figure 2.13: An extra leaf is constructed for a semantic action

- How does this differ from synthesized attributes?

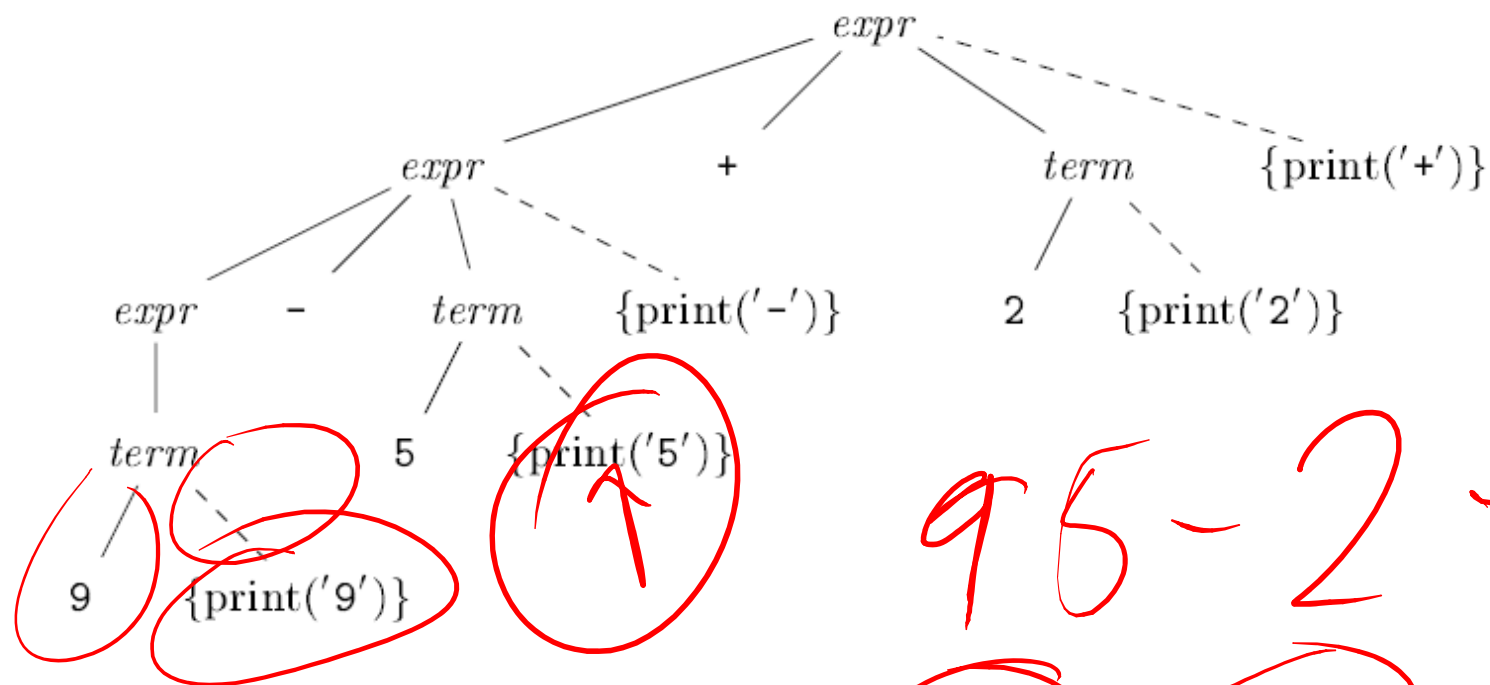


Figure 2.14: Actions translating 9-5+2 into 95-2+

Given

$expr$	$\rightarrow$	$expr_1 + term$	$\{print('+')\}$	-
$expr$	$\rightarrow$	$expr_1 - term$	$\{print('-')\}$	-
$expr$	$\rightarrow$	$term$		
$term$	$\rightarrow$	0	$\{print('0')\}$	-
$term$	$\rightarrow$	1	$\{print('1')\}$	-
		...		
$term$	$\rightarrow$	9	$\{print('9')\}$	-

actions

Figure 2.15: Actions for translating into postfix notation

# What is Parsing?

- Determine how string is generated
- Methods
  - Top-down
  - Bottom-up

*terminals*





# for ( ; expr ; expr ) other

$stmt \rightarrow$  **expr ;**  
          | **if ( expr ) stmt**  
          | **for ( optexpr ; optexpr ; optexpr ) stmt**  
          | **other**

$optexpr \rightarrow$   $\epsilon$   
          | **expr**

Figure 2.16: A grammar for some statements in C and Java

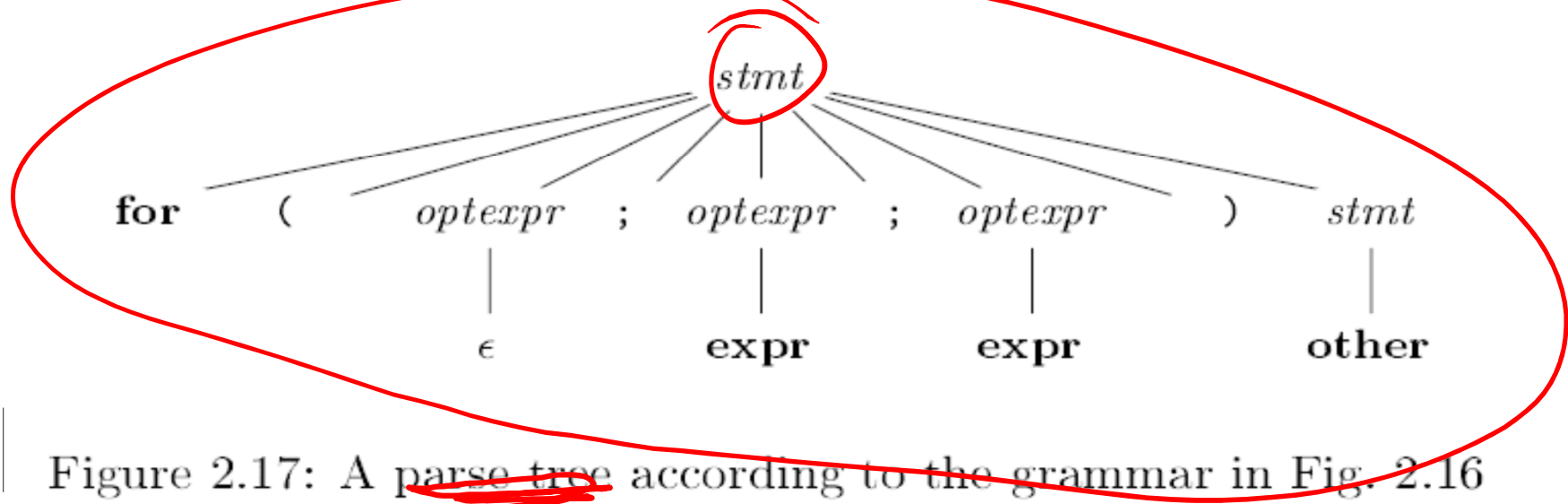


Figure 2.17: A parse tree according to the grammar in Fig. 2.16

# Top-down Parsing

- Recursive Descent

- What is it?
- What do you need?
- Issues?

- Predictive Parsing

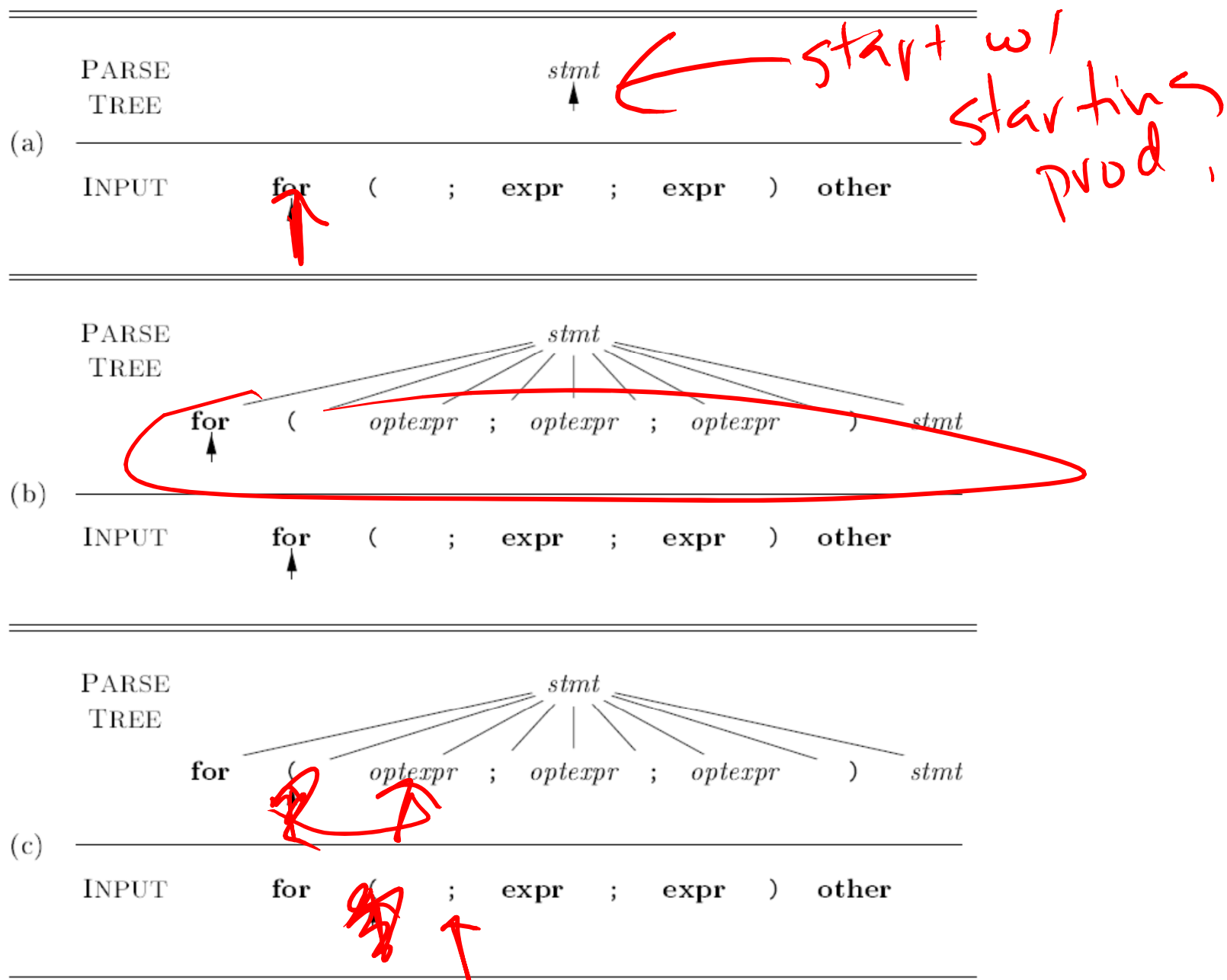
left to right  
scan over

the terms  
leftmost der.

backtracking  
doesn't work  
for

left lookahead (S)S  
recursive  
grammars  
S(S)S

LL(1)  
no backtracking



```

void stmt() {
    switch ( lookahead ) {
    case expr:
        match(expr); match(';'); break;
    case if:
        match(if); match('('); match(expr); match(')'); stmt();
        break;
    case for:
        match(for); match('(');
        optexpr(); match(';'); optexpr(); match(';'); optexpr();
        match(')'); stmt(); break;
    case other:
        match(other); break;
    default:
        report("syntax error");
    }
}

void optexpr() {
    if ( lookahead == expr ) match(expr);
}

void match(terminal t) {
    if ( lookahead == t ) lookahead = nextTerminal;
    else report("syntax error");
}

```

no back tracking

# Predictive Parsing

- Relies on:
  - $\alpha$  is string in grammar
  - $\text{FIRST}(\alpha)$  is set of terminals that appear first
  - If  $\alpha$  generates  $\epsilon$ , then  $\epsilon$  is in  $\text{FIRST}(\alpha)$
  - If  $A \rightarrow \alpha \mid \beta$ , then  $\text{FIRST}(\alpha)$  and  $\text{FIRST}(\beta)$  disjoint
- Example:
  - $\text{FIRST}(\text{stmt}) = \{\text{expr}, \text{if}, \text{for}, \text{other}\}$
  - $\text{FIRST}(\text{expr};) = \{\text{expr}\}$

# Looping Forever

- Left Recursion

- Rewrite:

$- \text{expr} \rightarrow \text{expr} + \text{term} \mid \text{term}$

$- A \rightarrow A\alpha \mid \beta$

Vs.

$- A \rightarrow \beta R$

$- R \rightarrow \alpha R \mid \varepsilon$

- What is the A,  $\alpha$ , and  $\beta$ ?

# Infix to Postfix/What's wrong?

$expr$	$\rightarrow$	$expr + term$	{ print('+' ) }
		$expr - term$	{ print('-' ) }
		$term$	
$term$	$\rightarrow$	0	{ print('0' ) }
		1	{ print('1' ) }
		...	
		9	{ print('9' ) }

Figure 2.21: Actions for translating into postfix notation

# How is this different/same?

$expr \rightarrow term\ rest$

$rest \rightarrow +\ term\ \{ \text{print}(' + ') \}\ rest$

$\quad \quad \quad -\ term\ \{ \text{print}(' - ') \}\ rest$

$\quad \quad \quad \epsilon$

$term \rightarrow 0\ \{ \text{print}(' 0 ') \}$

$\quad \quad \quad 1\ \{ \text{print}(' 1 ') \}$

$\quad \quad \quad \dots$

$\quad \quad \quad 9\ \{ \text{print}(' 9 ') \}$

Figure 2.23: Translation scheme after left-recursion elimination



# Reading/Assignments

- Continue Milestone 1
- Finish Chap. 2