

Compiler Construction

Lexical Analysis

Chapter -3

Regular Expressions

1. All Strings that start with “tab” or end with “bat”
 $\text{tab}\{A, \dots, Z, a, \dots, z\}^* \mid \{A, \dots, Z, a, \dots, z\}^* \text{bat}$
2. All Strings in Which Digits 1,2,3 exist in ascending numerical order:
 $\{A, \dots, Z\}^* 1 \{A, \dots, Z\}^* 2 \{A, \dots, Z\}^* 3 \{A, \dots, Z\}^*$
3. All strings of lowercase letters in which the letters are in ascending lexicographic order.
 $a^* b^* c^* \dots z^*$

Regular Expression

All strings of lowercase letters that contain the five vowels in order.

want -> $\text{other}^* \text{a} (\text{other}|\text{a})^* \text{e} (\text{other}|\text{e})^* \text{i} (\text{other}|\text{i})^* \text{o} (\text{other}|\text{o})^* \text{u} (\text{other}|\text{u})^*$

other -> $[\text{bcdfghjklmnpqrstvwxyz}]$

All strings of a's and b's that do not contain the **substring** abb.

$\text{b}^*(\text{a}^+\text{b}?)^*$

All strings of a's and b's that do not contain the **subsequence** abb.

$\text{b}^* \mid \text{b}^*\text{a}^+ \mid \text{b}^*\text{a}^+\text{ba}^*$

Regular Expressions

1. All strings of lowercase letters that contain the five vowels in order

$C \rightarrow \text{b}|\text{c}|\text{d}|\text{f}|\dots|\text{z}$

$\text{C}^*\text{a}(\text{C}|\text{a})^*\text{e}(\text{C}|\text{e})^*\text{i}(\text{C}|\text{i})^*\text{o}(\text{C}|\text{o})^*\text{u}(\text{C}|\text{u})^*$

2. All strings of a's and b's with an odd number of a's

$\text{b}^*(\text{ab}^*\text{ab}^*)^*\text{ab}^*$

3. All strings of 0's and 1's in which any two 0's in α are separated by three 1's

$1^*(0111)^*01^* + 1^*$

Find Regular Languages

1. $a(a|b)^*a$
2. $((\epsilon|a)b^*)^*$
3. $(a|b)^*a(a|b)(a|b)$
4. $a^*ba^*ba^*ba^*$
5. $(aa|bb)^*((ab|ba)(aa|bb)^*(ab|ba)(aa|bb)^*)^*$

Regular Definitions

Regular Definitions: Associate names with Regular Expressions

For Example : PASCAL IDs

letter $\rightarrow A | B | C | \dots | Z | a | b | \dots | z$

digit $\rightarrow 0 | 1 | 2 | \dots | 9$

id $\rightarrow \text{letter} (\text{letter} | \text{digit})^*$

Shorthand Notation:

“+” : one or more $r^+ = r^+ \in$ (Kleene) & $r^+ = r r^*$ (Positive)

“?” : zero or one $r? = r | \epsilon$

[range] : set range of characters (replaces “|”)

$[A-Z] = A | B | C | \dots | Z$

Example Using Shorthand : PASCAL IDs

id $\rightarrow [A-Za-z][A-Za-z0-9]^*$

Regular Definitions

Unsigned Number **1240, 39.45, 6.33E15, or 1.578E-41**

digit $\rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$

digits $\rightarrow \text{digit digit}^*$

optional_fraction $\rightarrow . \text{digits} \mid \epsilon$

optional_exponent $\rightarrow (E (+ \mid - \mid \epsilon) \text{digits}) \mid \epsilon$

num $\rightarrow \text{digits optional_fraction optional_exponent}$

Shorthand

digit $\rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$

digits $\rightarrow \text{digit}^+$

optional_fraction $\rightarrow (. \text{digits}) ?$

optional_exponent $\rightarrow (E (+ \mid -) ? \text{digits}) ?$

num $\rightarrow \text{digits optional_fraction optional_exponent}$

Regular Definitions

How can we use concepts developed so far to assist in recognizing tokens of a source language ?

Assume Following Tokens:

{ if, then, else, relop, id, num

What language construct are they used for ?

Given Tokens, What are Patterns ?

if \rightarrow if

then \rightarrow then

else \rightarrow else

relop $\rightarrow < \mid <= \mid > \mid >= \mid = \mid <>$

id $\rightarrow \text{letter (letter } \mid \text{ digit)}^*$

num $\rightarrow \text{digit}^+ (. \text{digit}^+) ? (E (+ \mid -) ? \text{digit}^+) ?$

What does this represent ?

Grammar:

stmt \rightarrow if *expr* then *stmt*

if *expr* then *stmt* else *stmt*

$\mid \epsilon$

expr \rightarrow term relop term \mid term

term \rightarrow id \mid num

Other tasks done by Lexical Analyzer

Scan away *blanks*, new lines, tabs
Can we Define Tokens For These?

blank → **blank**
tab → **tab**
newline → **newline**
delim → **blank** | **tab** | **newline**
ws → **delim**⁺

Ans: No token is returned to parser

Pattern, Token, Attribute-Value

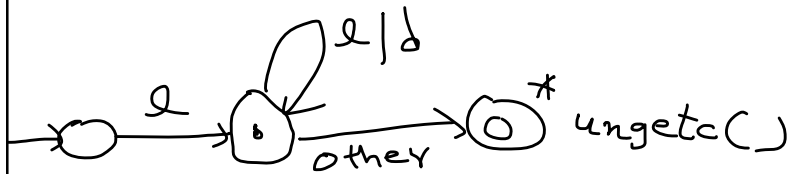
Regular Expression	Token	Attribute-Value
ws	-	-
if	if	-
then	then	-
else	else	-
id	id	pointer to table entry
num	num	Exact value
<	relop	LT
<=	relop	LE
=	relop	EQ
<>	relop	NE
>	relop	GT
>=	relop	GE

Note: Each token has a unique token identifier to define category of lexemes

CDF A)

Transition Diagrams

For Example : PASCAL IDs

letter \rightarrow A | B | C | ... | Z | a | b | ... | zdigit \rightarrow 0 | 1 | 2 | ... | 9✓ id \rightarrow letter (letter | digit)*

Transition Diagrams

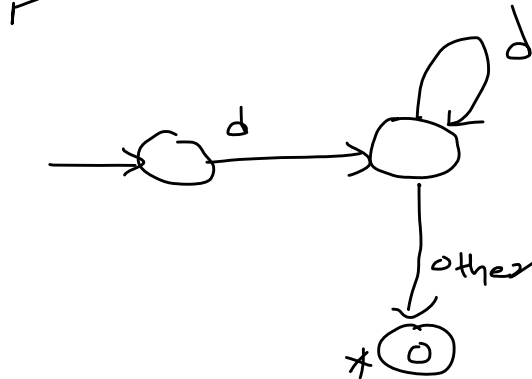
blank \rightarrow blanktab \rightarrow tabnewline \rightarrow newlinedelim \rightarrow blank | tab | newlinews \rightarrow delim +

Transition Diagrams

digit $\rightarrow 0 | 1 | 2 | \dots | 9$

digits $\rightarrow \text{digit } \underline{\text{digit}}^*$

Num $\rightarrow \text{digits}$



Transition Diagrams

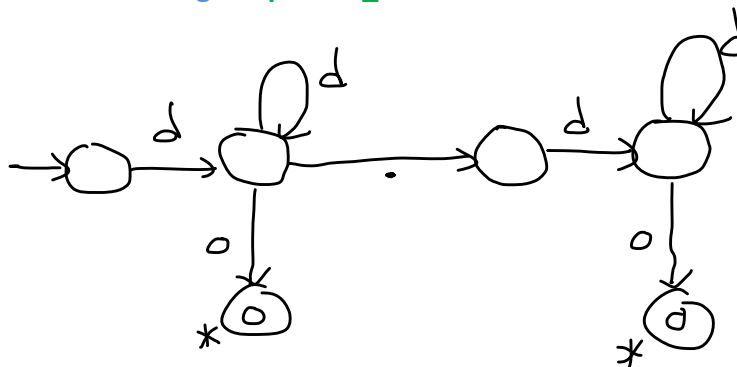
digit $\rightarrow 0 | 1 | 2 | \dots | 9$

digits $\rightarrow \text{digit } \underline{\text{digit}}^*$

optional_fraction $\rightarrow . \text{digits} | \epsilon$

Num $\rightarrow \text{digits } \text{optional_fraction}$

| 2.45



Transition Diagrams

Unsigned Number

1240, 39.45, 6.33E15, or 1.578E-41

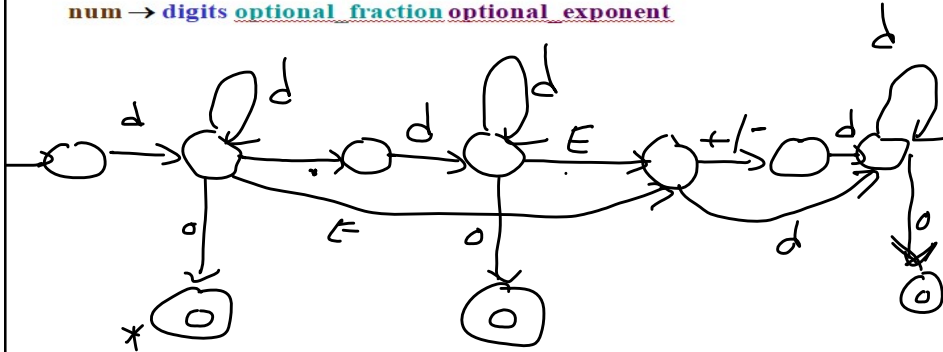
digit $\rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$

digits $\rightarrow \text{digit digit}^*$

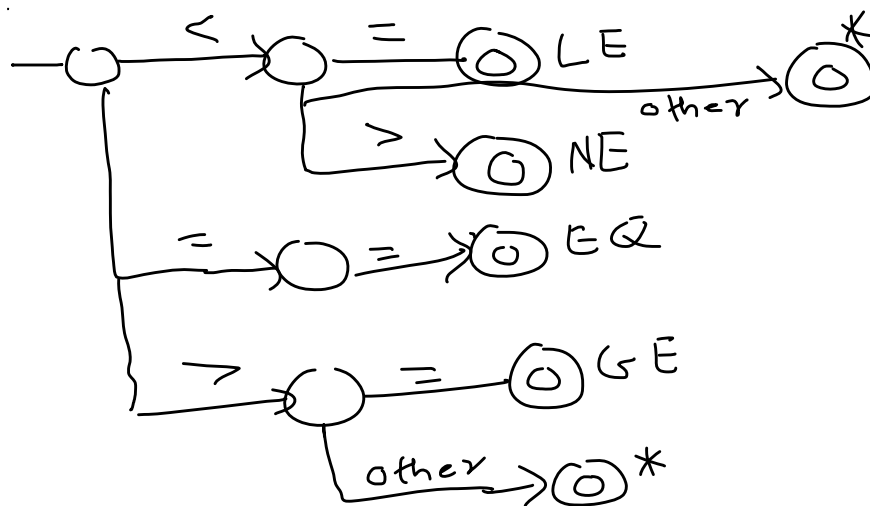
optional_fraction $\rightarrow \cdot \text{digits} \mid \epsilon$

optional_exponent $\rightarrow (E (+ \mid -) \epsilon) \text{ digits} \mid \epsilon$

num $\rightarrow \text{digits optional_fraction optional_exponent}$

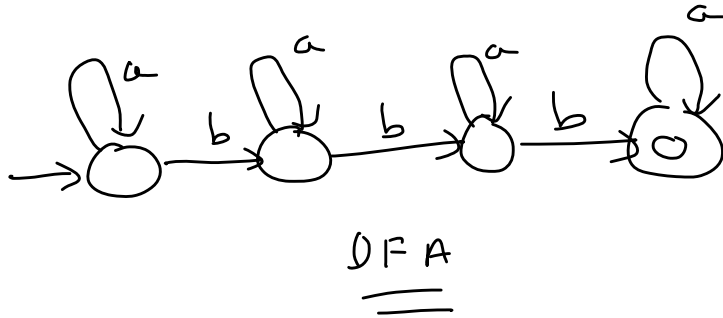


Transition Diagrams (Relational Operators)



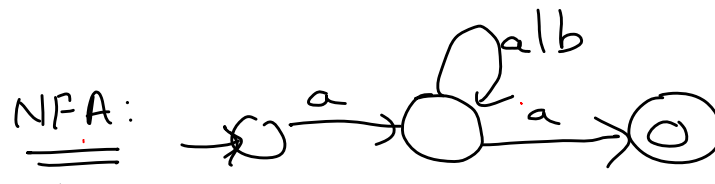
Transition Diagram

$a^*ba^*ba^*ba^*$

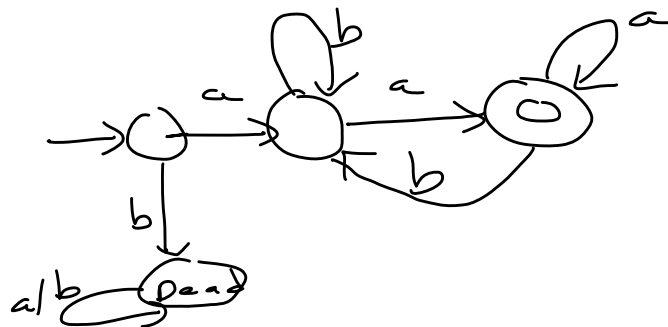


Transition Diagram

$a(a|b)^*a$



DFA



Transition Diagram (DFA)
 $(a|b)^*a(a|b)(a|b)$

Try Yourself

Transition Diagram (DFA)
 $(aa|bb)^*((ab|ba)(aa|bb)^*(ab|ba)(aa|bb)^*)^*$

Try Yourself