

Compiler Construction

Lexical Analysis

Chapter -3

Convert Finite Automata To Regular Expression

- Arden's Theorem
- State Elimination Method

Arden's Theorem

It states that-

Let P and Q be two regular expressions over Σ .

If **P does not contain a null string**
then

$R = Q + RP$ has a unique solution i.e. $R = QP^*$

Theorem is applicable to only Deterministic Finite Automata

Arden's Theorem (Proof)

Theorem:

Let P and Q be two regular expressions over Σ .

If **P does not contain a null string**
then

$R = Q + RP$ has a unique solution i.e. $R = QP^*$

Proof –

$$\begin{aligned} R &= Q + (Q + RP)P && \text{[After putting the value } R = Q + RP\text{]} \\ &= Q + QP + RPP \end{aligned}$$

When we put the value of **R** recursively again and again, we get the following equation –

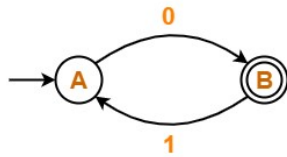
$$R = Q + QP + QP^2 + QP^3 + \dots$$

$$R = Q (\epsilon + P + P^2 + P^3 + \dots)$$

$$R = QP^* \quad \text{[As } P^* \text{ represents } (\epsilon + P + P^2 + P^3 + \dots)]$$

Hence, proved.

Arden's Theorem



For State A:

$$A = \epsilon + B1$$

For State B:

$$B = A0$$

Bring accepting state in the form $R=Q+RP$,

$$B = (\epsilon + B1)0$$

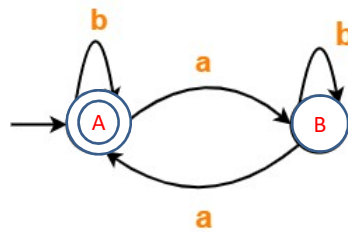
$$B = \epsilon 0 + B10$$

$$B = 0 + B10$$

Here, $B=R$, $Q=0$ and $P=10$

As per $R=QP^*$, **solution will be : $0(10)^*$**

Arden's Theorem



For State A:

$$A = \epsilon + Ab + Ba$$

For State B:

$$B = Aa + Bb$$

Applying Arden's Theorem to State B,

Here, $R=B$, $Q=Aa$ and $P=b$

$$B = Aab^*$$

Bring accepting state in the form $R=Q+RP$,

$$A = \epsilon + Ab + Aab^*a$$

$$A = \epsilon + A(b + ab^*a)$$

Here, $A=R$, $Q=\epsilon$ and $P=b + ab^*a$

As per $R=QP^*$, **solution will be : $(b + ab^*a)^*$**

Arden's Theorem

For State A:

$$A = \epsilon$$

For State B:

$$B = Ab + Ba + Ca$$

For State C:

$$C = Aa$$

Bring accepting state in the form $R = Q + RP$,

$$B = Ab + Ba + Aaa$$

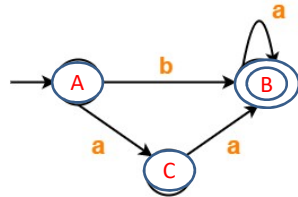
$$B = \epsilon b + Ba + \epsilon aa$$

$$B = b + Ba + aa$$

$$B = b + aa + Ba$$

Here, $B = R$, $Q = b + aa$ and $P = a$

As per $R = QP^*$, **solution will be : $(b + aa)a^*$**



Arden's Theorem

For State A:

$$A = \epsilon + Aa + Ca$$

For State B:

$$B = Ab + Bb + Cb$$

For State C:

$$C = Ba$$

$$B = Ab + Bb + Bab$$

Applying Arden's Theorem

$$B = Ab(b + ab)^*$$

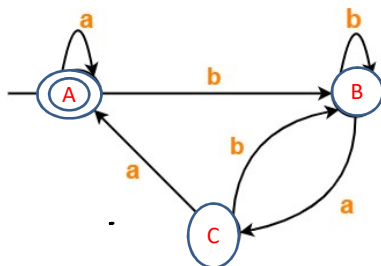
$$C = Ab(b + ab)^*a$$

$$A = \epsilon + Aa + Ab(b + ab)^*aa$$

$$A = \epsilon + A(a + b(b + ab)^*aa)$$

Bring accepting state in the form $R = Q + RP$,

As per $R = QP^*$, **solution will be : $(a + b(b + ab)^*aa)^*$**



Arden's Theorem

For State A:

$$A = \epsilon + A1 + B1$$

For State B:

$$B = A0 \quad B = (1+01)^*0$$

Bring accepting state in the form $R=Q+RP$,

$$A = \epsilon + A1 + A01$$

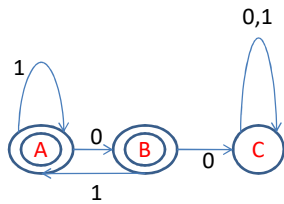
$$A = \epsilon + A(1+01)$$

Here, $A=R$, $Q=\epsilon$ and $P=(1+01)$

As per $R=QP^*$, **$A = (1+01)^*$**

$$\mathbf{B = (1+01)^*0}$$

Final Solution is : $(1+01)^* + (1+01)^*0$



Eliminate dead state directly.

Also remove any incoming arrow to the dead state

Arden's Theorem

For State A:

$$A = \epsilon + A1 + C1$$

For State B:

$$B = A0 + B0 + C0$$

For State C:

$$C = B1$$

Applying Arden's Theorem to state A

$$A = (\epsilon + C1)1^*$$

$$A = 1^* + C11^*$$

$$B = (1^* + C11^*)0 + B0 + C0$$

$$B = 1^*0 + C11^*0 + C0 + B0$$

Applying Arden's Theorem to state B

$$B = (1^*0 + C11^*0 + C0)0^*$$

Bring accepting state in the form $R=Q+RP$,

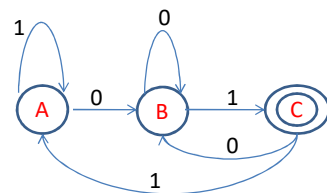
$$C = (1^*0 + C11^*0 + C0)0^*1$$

$$C = 1^*00^*1 + C11^*00^*1 + C00^*1$$

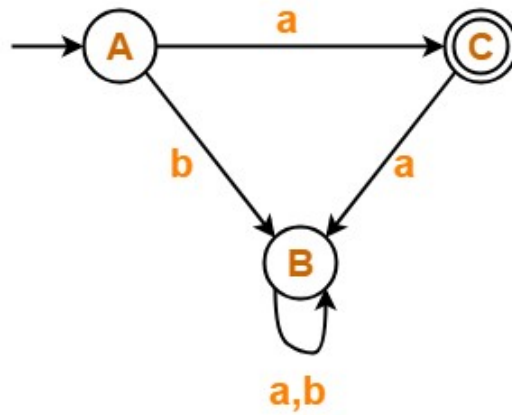
$$C = 1^*00^*1 + C(11^*00^*1 + 00^*1)$$

As per $R=QP^*$,

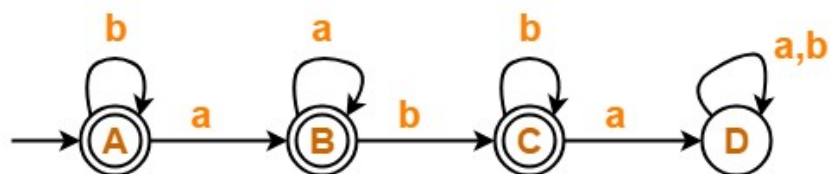
solution will be : $1^*00^*1(11^*00^*1 + 00^*1)^*$



Try Yourself



Try Yourself



Try Yourself

