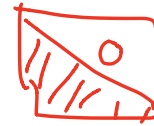


Numerical Linear Algebra

A square matrix A of order n is upper/lower triangular if

$$A_{ij} = 0$$

$$\forall i < j$$



$$A_{ij} = 0$$

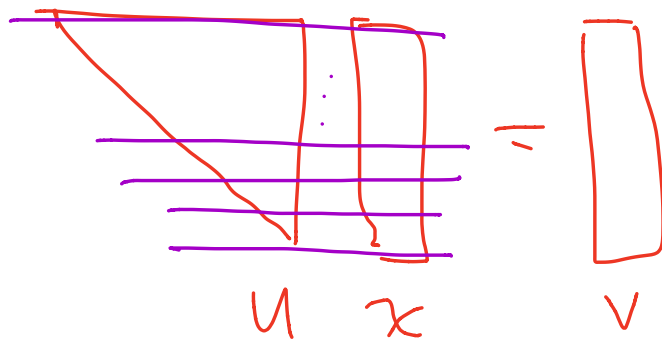
$$\forall i > j$$



We want to compute for a given upper triangular matrix $U \in \mathbb{R}^{n \times n}$

and vector $v \in \mathbb{R}^n$, $U^{-1}v$

Want: $x = U^{-1}v \Rightarrow Ux = v$



Well $u_{nn}x_n = v_n \Rightarrow x_n = \frac{v_n}{u_{nn}}$.

and the second to last row gives

$$u_{n-1,n-1}x_{n-1} + u_{n-1,n}x_n = v_{n-1}$$

$$\Rightarrow x_{n-1} = \frac{1}{u_{n-1,n-1}} (v_{n-1} - u_{n-1,n}x_n)$$

So in general

$$x_i = \frac{1}{u_{ii}} \left(v_i - \sum_{j=i+1}^n u_{ij}x_j \right).$$

This process is called **back solving**

as we solve $n, n-1, \dots, 2, 1$

Prmk: Similar algorithm for lower triangular matrix exists called

forward solving

Rmk: If we wanted $U^T v$?

Well $U^T x = v$ and we could do forward solve with the columns dotted with the rows.

Rmk: In R

$x = U^{-1} v$: $x \leftarrow \text{backsolve}(U, v)$

$x = U^{-T} v$: $x \leftarrow \text{backsolve}(U, v, \text{trans} = \text{TRUE})$

math	R
$A B$	$A \%*\% B$
$A^T B$	$\text{crossprod}(A, B)$
$A B^T$	$\boxed{\text{tcrossprod}(A, B)}$

Some Useful Decompositions

(i) Cholesky

If A is a square symmetric positive definite matrix then we can find an upper triangular matrix C such that

$$A = C^T C$$

C Cholesky factor.

Ex: How to compute $A^{-1}v = x$

$$A^{-1}v = C^{-1} \underbrace{C^{-T}v}_{\text{forward solve}}$$

backward solve.

$$\boxed{\text{R}} \quad C \leftarrow \text{chol}(A)$$

$$y \leftarrow \text{backsolve}(C, v, \text{trans}=\text{TRUE})$$

$$x \leftarrow \text{backsolve}(C, y)$$

□ compute the Cholesky factor.

Ex: Compute the log density

$N(\vec{\mu}, \Sigma)$ at x .

$$\log f(x) = -\frac{n}{2} \log 2\pi - \frac{1}{2} \log |\Sigma|$$

$$- \frac{1}{2} (x - \mu)^T \Sigma^{-1} (x - \mu)$$

Define: $\Sigma = CC^T$ then .

$$\log f(x) = -\frac{n}{2} \log 2\pi - \frac{1}{2} \log |C^T C|$$

$$- \frac{1}{2} (x - \mu)^T C^{-1} C^{-T} (x - \mu)$$

$$= -\frac{n}{2} \log 2\pi - \log |C| - \frac{1}{2} \underbrace{\left[C^{-T} (x - \mu) \right]^T \left[\underbrace{C^{-T} (x - \mu)}_y \right]}_{y^T y}$$

$$= -\frac{n}{2} \log 2\pi - \sum_{i=1}^n \log C_{ii} - \frac{1}{2} y^T y$$

$$\boxed{R} \quad C \leftarrow \text{chol}(\Sigma)$$

$$y \leftarrow \text{bseksolve}(C, x-\mu, \text{trans}=\text{TRANS})$$

$$-\frac{n}{2} \log(2\pi) - \text{Sum}(\log(\text{diag}(C)))$$

$$- \text{Sum}(y^2)$$

(ii) If X is an $n \times p$ matrix
 then we can find an
 $n \times n$ orthogonal matrix Q ,
 $(Q^T Q = Q Q^T = I)$ and $n \times p$
 trapezoidal matrix R .

$$n=p \quad \boxed{} = \boxed{}^Q \begin{array}{c} R \\ \swarrow \end{array}$$

$$n>p \quad \boxed{} = \boxed{\begin{array}{|c|c|} \hline & \\ \hline \end{array}} \begin{array}{c} \boxed{}^R \\ \hline 0 \end{array}$$

$$n<p \quad \boxed{} = \boxed{} \begin{array}{c} \swarrow \\ \end{array}$$

$$\begin{array}{ccc} \text{---} & \text{---} & \text{---} \\ & Q & R \end{array}$$

Take $n > p$ and minimize

$$SS(\beta) = (Y - X\beta)^T (Y - X\beta)$$

$$X = QR = [Q_1, Q_2] \begin{bmatrix} R_1 \\ 0 \end{bmatrix} = \underbrace{QR_1}_{\text{Thin QR}}$$

$$SS(\beta) = \left(\begin{bmatrix} Q_1^T \\ Q_2^T \end{bmatrix} Y - \begin{bmatrix} R_1 \beta \\ 0 \end{bmatrix} \right)^T \begin{pmatrix} Q_1^T Y - R_1 \beta \\ Q_2^T Y \end{pmatrix}$$

$$= \underbrace{(Q_1^T Y - R_1 \beta)^T (Q_1^T Y - R_1 \beta)}_{\text{Minimizer}} + Y^T Q_2 Q_2^T Y$$

Minimizer

$$Q_1^T Y - R_1 \hat{\beta} = 0 \quad \text{or}$$

$$\hat{\beta} = (Q_1^T Q_1)^{-1} Q_1^T Y \quad \text{is the problem}$$

$$\begin{pmatrix} p & - & n_1 & q_1' & y \end{pmatrix} \quad \text{we solve "easy"}$$

So

$$SS(\hat{\beta}) = Y^T Q_2 Q_2^T Y$$

$$\hat{Y} = X\hat{\beta} = \underbrace{Q_1 Q_1^T}_H Y$$

$$\hat{e} = \underbrace{Q_2 Q_2^T}_{1-H} Y$$

- Q partitions the X into the covariate & residual space.