

Recap

Symmetrized Conductance:

$$\bar{\phi}(s) = \frac{|E(s, \bar{s})|}{|V_d(s) \cdot V_d(\bar{s})|} \cdot \text{Vol}(V)$$

We argued

$$\bar{\phi}_a = \min_{S \subseteq V} \bar{\phi}(s) = \min_{S \subseteq V} \frac{|E(s, \bar{s})|}{|E_{K_a}(s, \bar{s})|}$$

$$= \min_{S \subseteq V} \frac{\mathbf{1}_S^T L_a \mathbf{1}_S}{\mathbf{1}_S^T L_{K_a} \mathbf{1}_S}$$

Relaxing  
cuts  $\rightarrow$  all vectors

$$\geq \min_{x \in \mathbb{R}^n} \frac{x^T L_a x}{x^T L_{K_a} x} = \lambda_2$$

Today:  $\bar{\phi}_a \leq \sqrt{2\lambda_2}$

Metrics metrics over the vertex set.

Def: A metric  $d \in \mathbb{R}^{n \times n}$ :

$$(i) \ d_{ii} = 0 \quad i \in V$$

$$(ii) \ d_{ij} = d_{ji}$$

$$(iii) \ d_{ij} \leq d_{ik} + d_{kj} \quad \forall i, j, k \in V$$

Ex:  $x, y \in \mathbb{R}^k \quad d(x, y) = \|x - y\|_p$

Suppose  $x_1, \dots, x_n \in \mathbb{R}^k$

$$d(x_i, x_j) = \|x_i - x_j\|_2^2 \text{ not}$$

a metric in general.

A semi-metric is a metric when  
(iii) need not hold.

Then if we define

$$\alpha(G) = \sum_{ij \in E} \underbrace{w_{ij} d_{ij}}_{\text{volume of edge}}$$

### Cut Metrics

Def: If  $S \subseteq V$  the cut metric  $\delta^{(S)}$

$$\delta_{ij}^{(S)} = \begin{cases} 0 & \text{Same side of cut} \\ 1 & \text{diff side of cut} \end{cases}$$

### Relation to Conductance

$$\delta^{(S)}(G) = |E(S, \bar{S})|$$

$$\delta^{(S)}(K_n) = \frac{\text{Vol}(S)\text{Vol}(\bar{S})}{\text{Vol}(V)}$$

$$\bar{\phi}(S) = \frac{\delta^{(S)}(G)}{\delta^{(S)}(K_n)}$$

$$\bar{\phi}_n = \min_{S \subseteq V} \frac{\delta(G)}{\delta(K_n)} = \min_{S \subseteq V} \frac{\mathbf{1}_S^T L \mathbf{1}_S}{\mathbf{1}^T L(K_n) \mathbf{1}}$$

Relation to Spectral Gap

$$\lambda_2 = \min_x \frac{\sum_{ij \in E} w_{ij} (x_i - x_j)^2}{\sum_{ij \in V} \frac{d_i d_j}{\text{Vol}(V)} (x_i - x_j)^2}$$

Def: A semi-metric  $g$  is  $\ell_2^2$ -embeddable if there exists some embedding

$v_i \in \mathbb{R}^d$  s.t.

$$g_{ij} = \|v_i - v_j\|_2^2 \quad \forall i, j \in V$$

Claim:  $\lambda_2 = \min_{\substack{g \text{ is } \ell_2^2 \\ \text{embeddable}}} \frac{g(G)}{g(K_n)}$

Lemma  $\frac{\sum a_i}{\sum b_i} \geq \min_i \frac{a_i}{b_i}$

Pf (of claim):

$$\min_{g \in \ell_2^2} \frac{g(G)}{g(K_n)} = \frac{\sum_{i,j \in E} w_{ij} \|v_i - v_j\|^2}{\sum_{i,j \in V} \frac{d_i d_j}{\text{vol}(V)} \|v_i - v_j\|^2}$$

$$w_{ij} \|v_i - v_j\|^2 = \sum_{k=1}^d w_{ij} (v_i - v_j)_k^2$$

$$= \frac{\sum_{k=1}^d \sum_{i,j \in E} w_{ij} (v_i - v_j)_k^2}{\sum_{k=1}^d \sum_{i,j \in V} \frac{d_i d_j}{\text{vol}(V)} (v_i - v_j)_k^2}$$

$$\geq \min_{1 \leq k \leq d} \frac{\sum_{i,j \in E} w_{ij} (v_i - v_j)_k^2}{\sum_{i,j \in V} \frac{d_i d_j}{\text{vol}(V)} (v_i - v_j)_k^2} \geq \lambda_2$$

$$1 = \frac{1}{n} \sum \frac{d(i,j)}{d(V)} (v_i - v_j)^2$$

Plugging in  $\lambda_2$  then

$$\min_{g \in \mathbb{R}^2} \frac{g(G)}{g(K_n)} = \lambda_2$$

Rmk: A cut metric is  $\mathbb{R}^2$ -embeddable.

$$\begin{array}{c} 0 \qquad \qquad 1 \\ | \qquad \qquad | \\ \hline s \qquad \qquad \overline{s} \end{array} \quad v_i = \begin{cases} 0 & i \in S \\ 1 & i \in \overline{S} \end{cases}$$

$$f_{ij}^{(1)} = |v_i - v_j|^2$$

Where we are

cut metrics - discrete

$\mathbb{R}^2$ -embeddable - continuous.

## Cut Cone

Def: A set  $A \in \mathbb{R}^k$  is a conv if

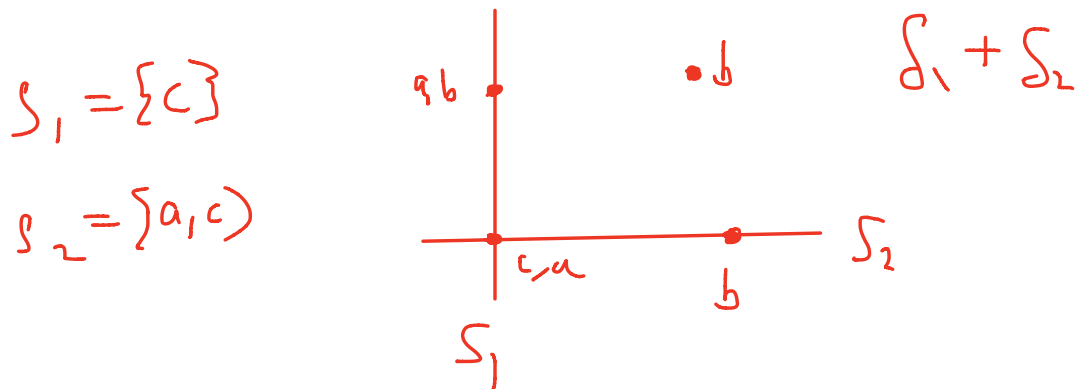
$$x \in A \Rightarrow \alpha x \in A \quad \forall \alpha > 0$$

Def A convex cone generated by a set of points  $B = \{v_i \in \mathbb{R}^n\}$

$$\text{Con}(B) = \{x \in \mathbb{R}^n : x = \sum \alpha_i v_i, \alpha_i \geq 0\}$$

Def: The Cut Cone is  $\text{Cone}(\{d\}_{s \leq v})$

$\Sigma x$  :  $a, b, c$



1.  $\text{C}_2\text{H}_5\text{OH} + \text{H}^+ \rightarrow \text{C}_2\text{H}_5\text{OH}_2^+$

Defn. CUT - METRICS THAT ARE EMBEDDABLE  
IN  $\ell_1$ .

Def  $\ell_1$  embeddability: A metric  $d$

is  $\ell_1$  embeddable if  $\exists \{v_i \in \mathbb{R}^K\}$

s.t.  $d_{ij} = \|v_i - v_j\|_1$

Thm: The set of  $\ell_1$ -embeddable  
metrics are in CUT

(1) Cut Metrics  $\delta^{(s)}$  in  $\ell_1$  embeddable

(2) Any  $\ell_1$  embeddable metric  $d$

is a conic combination of cut metrics

$$d_{ij} = \|v_i - v_j\|_1 = \sum_s \alpha_s \delta^{(s)}(i, j)$$

Claim It suffices to show this  
for one dimension





sweep  
cut  $S_t = \{1, 2, \dots, t\}$

$$d = \sum_{t=1}^{j-1} |v_{t+1} - v_t| \delta^{(S_t)}$$

$$d_{ij} = \sum_{k=i}^{j-1} |v_{k+1} - v_k| = v_i - v_j$$