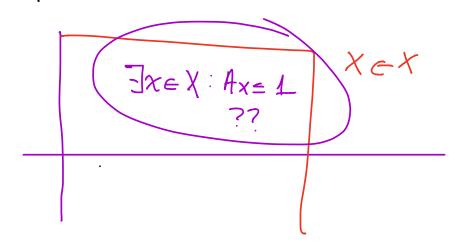
Last time:

Optimize



Came up with a robust way

to answer with $f_{M}(x) = S_{M}(A_{X})$ $f_{M}(x) = S_{M}(A_{X})$ $f(x) = M_{M}(A_{X})$ $f(x) = M_{M}(A_{X})$

Alg: $\overline{X}_0 = 0$ $X_{t+1} = \text{any min} \langle \nabla f_n(\overline{X}_t), \chi \rangle$ $\chi_{\epsilon} \chi$

If turns out we can use
$$X_{t+1} = a_{ny} \{X \in X : \{f_n(X_t), x\} \le 1\}$$

So either X_{t+1} does not exists (no feasible sol.) or ge for $T = O(\frac{p^2 \log n}{2^2})$ iterating then $A \times_T \le 1 + \epsilon$
 $X_T \le X$

Relation to the salle point problem.

min
$$\max_{x \in X} p^T A x - 1 + MH(p)$$

 $f(x)$

+n(x)

Could have also argued through the dual primal plagar relationship.

max (mh pTAx-1)
pe am (xex

Less illustration of the regularization method.

Rock this method extends 1100 with any convex constraints.

For approximate solutions in quadrata programming.

fu(z)=ulog z exp(zi/n)

$$\nabla_{ii} f_{n}(z) = \begin{cases} e^{\frac{2i}{n}} - e^{\frac{2i}{n}} & e^{\frac{2i}{n}} & e^{\frac{2i}{n}} \\ \sum_{e} e^{\frac{2i}{n}} & \sum_{e} e^{\frac{2i}{n}} \\ & & (\sum_{e} e^{\frac{2i}{n}})^{L} \end{cases}$$

$$= \frac{1}{M} \left(\frac{1}{r} - \frac{1}{r^2} \right)$$

$$\nabla_{ij}^{2} f_{n}(z) = \frac{1}{m} \frac{e^{\frac{1}{2}i/m}}{(2ie^{\frac{1}{2}i/m})^{2}}$$

$$\nabla^2 f_n(z) = \frac{1}{n} \left(dling(p) - ppT \right)$$

* Commant generating function.

fn(2+h) = fn(2 + < Dfn(2), h) +

 $\frac{h^{T} n^{2} f(\tilde{z})h}{2} + O(h^{3})$

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So we have 11.1100 someothress.

How does this work for SDB?

∑y. A" > I y ∈ R"

0/11 - /-

$$f_{M}(y) = \lambda_{Max} (2'y' k') - 1$$

$$f_{M}(y) = \max_{X \in \mathbb{Z}_{q'}(k')} - 1$$

$$I \cdot x = 1$$

$$x \neq 0$$

$$H(x) = Tr(X \log X) = H(\lambda_{1}(x))$$

$$\nabla f_{M}(x) = \frac{e^{\sum_{X \in \mathbb{Z}_{q'}(k')}}}{1 \cdot e^{\sum_{X \in \mathbb{Z}_{q'}(k')}}}$$