

Flow based Graph Partitioning

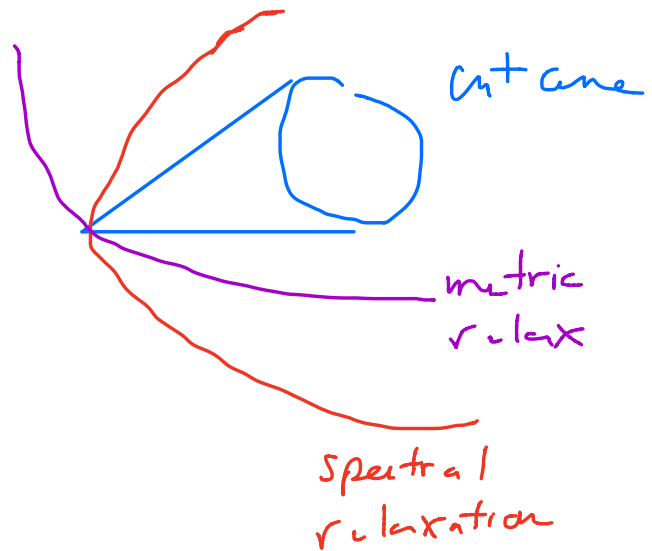
Goal: find the min. cut

$$\overline{\Phi}_G = \min_{S \subseteq V} \overline{\Phi}(S)$$

Approach metric relaxation

Set of
all L_1 -embeddable
metrics
 \cong

(convex(cut-metrics))



Solved via linear programming

$$\min_{e \in G} \sum \delta_e \quad \text{s.t.} \quad \sum_a \frac{d_i d_j}{\text{Vol}(a)} \delta_{ij} = 1$$

$$\forall i, j \quad \forall p \in P_{ij} \quad d_{ij} \leq \sum_{e \in p} d_e \quad f_p$$

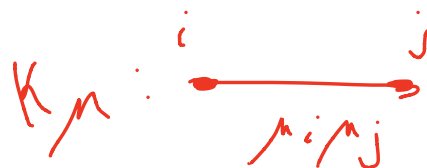
We can generalize this flow prob.
to solve more complex problems

$$\alpha_\mu(S) = \frac{E(S, \bar{S})}{\mu(S)\mu(\bar{S})} \quad \swarrow \text{any measure.}$$

$$\underline{\text{Ex:}} \quad \mu = 1, \quad \alpha(S) = \frac{|E(S, \bar{S})|}{|S||\bar{S}|}$$

In our case we are looking at
the problem

$$\min_{S \subseteq V} \frac{d_S(h)}{d_S(k_\mu)}$$



Dual



$$\max \alpha \quad \text{s.t.} \quad \forall e \in E \quad \sum_{p: e \in p} f_p \leq 1 \quad (\text{Capacity})$$

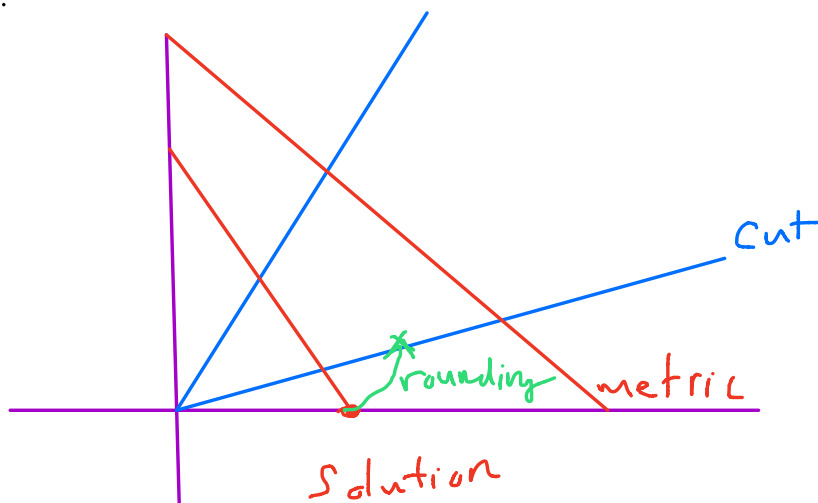
$$\forall_{ij} \quad \sum_{p \in P_{ij}} f_p \geq \alpha r_{ij} \quad (\text{total flow requirements})$$

total flow requirements

$$\forall p \quad f_p \geq 0$$

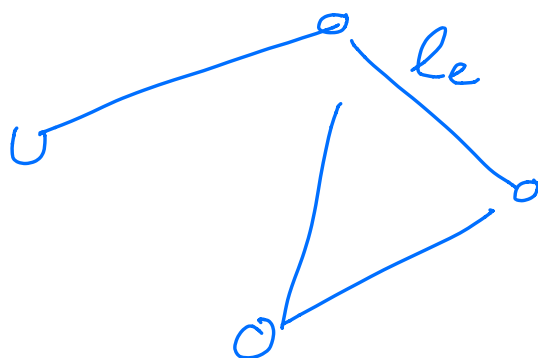
Suppose that we have a max flow.

Now we want to turn this into a cut.



Rounding Sol. to cut center

Suppose we have the sol. to this flow prob.



Lemma: Assume there exists $S \subseteq V$

with $\mu(S) \geq \frac{\mu(V)}{c} \equiv \frac{m}{c}$

s.t. $\forall i, j \in S \quad \delta_{ij} \leq \alpha$

$$\sum_{j \in V \setminus S} \mu_j d(j, S) \quad \left\{ \begin{array}{l} \text{closes distance} \\ \text{from } j \rightarrow c \in S \end{array} \right.$$

$$\leq \sum_j \mu_j \sum_k (\mu_k (d_{jk} + d(k, S)))$$

$$j \in V \setminus S \quad k \in V.$$

$$= \sum_{(j,k) \in (V \setminus S, V)} \mu_j \mu_k d_{jk} + \mu(V \setminus S) \sum_{k \in V} \mu d(k, S)$$

$$\geq \frac{1}{2} (1 - \mu(S))^2 \alpha$$

bc:

$$\sum_{i,j \in S \times S} \mu_i \mu_j d_{ij} \leq (\mu(S))^2 \alpha$$

Fenchel Embedder

$$y_i = d(i, S)$$

$$|y_i - y_j| \leq |d(j, S) - d(i, S)| \leq d_{ij}$$

$$\sum \mu_i \mu_j |y_i - y_j| = \sum \mu_i \mu_j |d(u_i, s) - d(u_j, s)|$$

$$\geq \sum_{i \in V} \mu_i d(i, s)$$

So these ℓ_1 embeddings
are interpreted as restrictions
of the distances.