

# Artificial Intelligence

## Inference in First-Order Logic

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# Propositional vs. First-Order Logic

Foo

# Unification and First-Order Inference

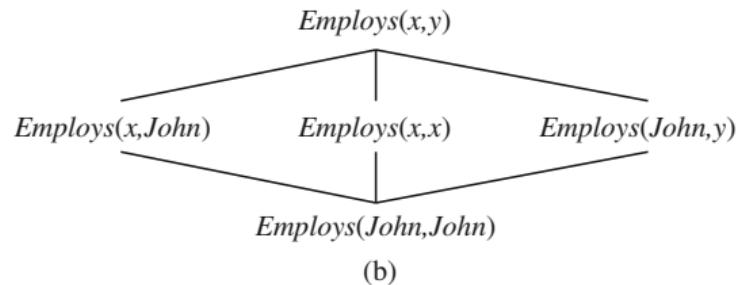
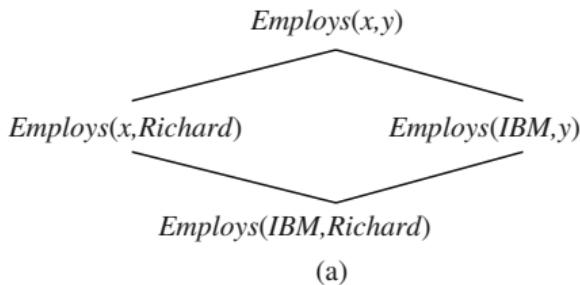
**function** UNIFY( $x, y, \theta = \text{empty}$ ) **returns** a substitution to make  $x$  and  $y$  identical, or *failure*

- if**  $\theta = \text{failure}$  **then return** *failure*
- else if**  $x = y$  **then return**  $\theta$
- else if** VARIABLE?( $x$ ) **then return** UNIFY-VAR( $x, y, \theta$ )
- else if** VARIABLE?( $y$ ) **then return** UNIFY-VAR( $y, x, \theta$ )
- else if** COMPOUND?( $x$ ) **and** COMPOUND?( $y$ ) **then**
  - return** UNIFY(ARGS( $x$ ), ARGS( $y$ ), UNIFY(OP( $x$ ), OP( $y$ ),  $\theta$ ))
- else if** LIST?( $x$ ) **and** LIST?( $y$ ) **then**
  - return** UNIFY(REST( $x$ ), REST( $y$ ), UNIFY(FIRST( $x$ ), FIRST( $y$ ),  $\theta$ ))
- else return** *failure*

**function** UNIFY-VAR( $var, x, \theta$ ) **returns** a substitution

- if**  $\{var/val\} \in \theta$  for some  $val$  **then return** UNIFY( $val, x, \theta$ )
- else if**  $\{x/val\} \in \theta$  for some  $val$  **then return** UNIFY( $var, val, \theta$ )
- else if** OCCUR-CHECK?( $var, x$ ) **then return** *failure*
- else return** add  $\{var/x\}$  to  $\theta$

# Unification and First-Order Inference



# Forward Chaining

**function** FOL-FC-ASK( $KB, \alpha$ ) **returns** a substitution or *false*

**inputs:**  $KB$ , the knowledge base, a set of first-order definite clauses  
 $\alpha$ , the query, an atomic sentence

**while** true **do**

$new \leftarrow \{\}$      *// The set of new sentences inferred on each iteration*

**for each rule in**  $KB$  **do**

$(p_1 \wedge \dots \wedge p_n \Rightarrow q) \leftarrow \text{STANDARDIZE-VARIABLES}(rule)$

**for each**  $\theta$  such that  $\text{SUBST}(\theta, p_1 \wedge \dots \wedge p_n) = \text{SUBST}(\theta, p'_1 \wedge \dots \wedge p'_n)$   
            for some  $p'_1, \dots, p'_n$  in  $KB$

$q' \leftarrow \text{SUBST}(\theta, q)$

**if**  $q'$  does not unify with some sentence already in  $KB$  or  $new$  **then**

                add  $q'$  to  $new$

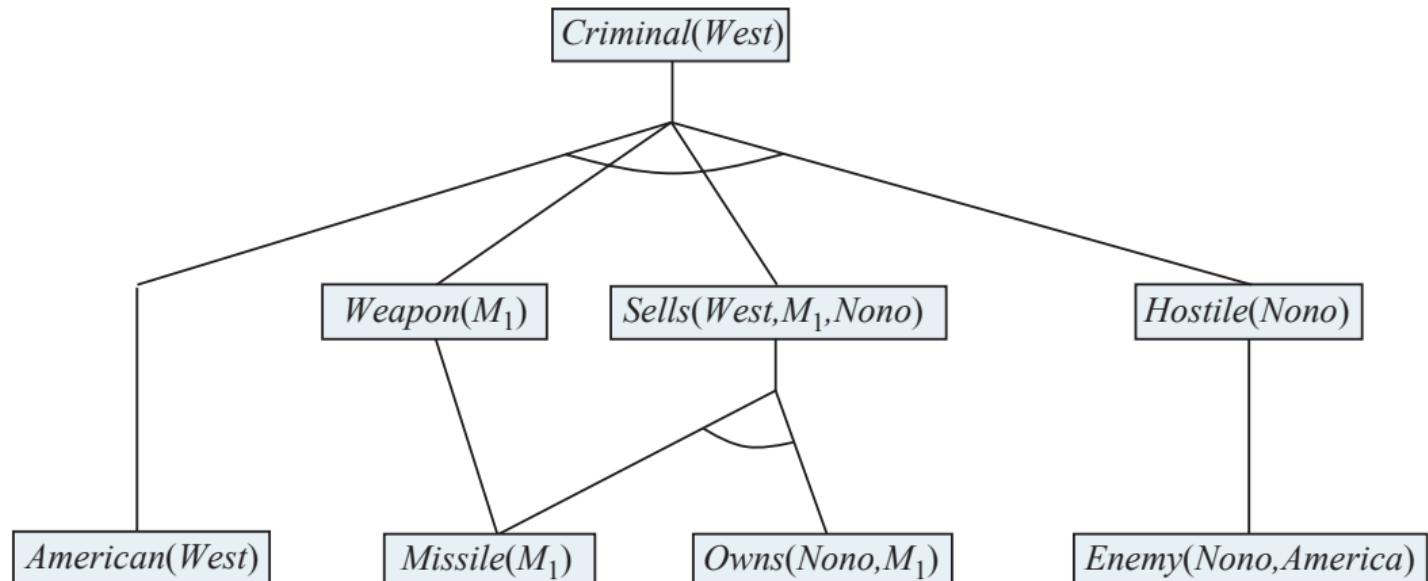
$\phi \leftarrow \text{UNIFY}(q', \alpha)$

**if**  $\phi$  is not failure **then return**  $\phi$

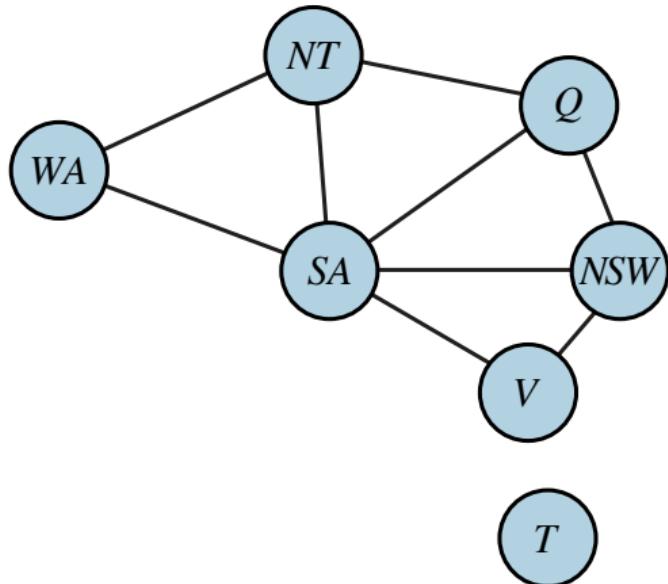
**if**  $new = \{\}$  **then return** false

    add  $new$  to  $KB$

## Forward Chaining



## Forward Chaining



(a)

$Diff(wa, nt) \wedge Diff(wa, sa) \wedge$   
 $Diff(nt, q) \wedge Diff(nt, sa) \wedge$   
 $Diff(q, nsw) \wedge Diff(q, sa) \wedge$   
 $Diff(nsw, v) \wedge Diff(nsw, sa) \wedge$   
 $Diff(v, sa) \Rightarrow Colorable()$

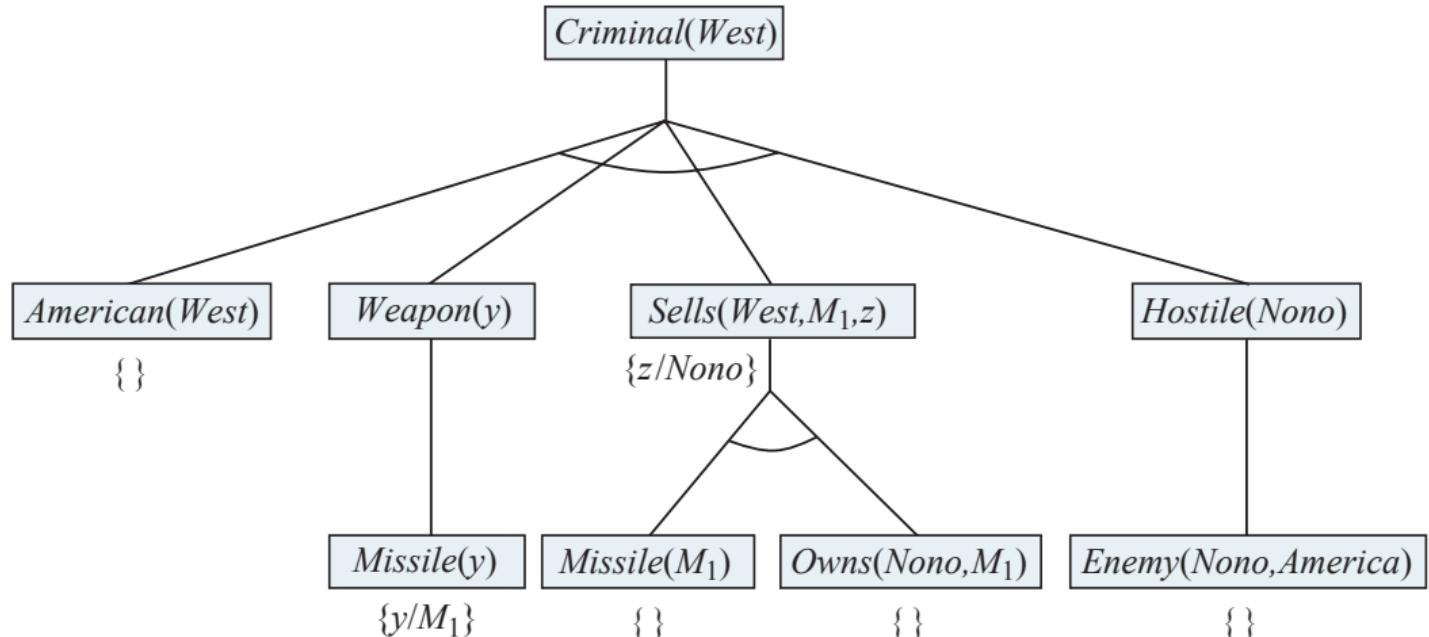
$Diff(Red, Blue) \quad Diff(Red, Green)$   
 $Diff(Green, Red) \quad Diff(Green, Blue)$   
 $Diff(Blue, Red) \quad Diff(Blue, Green)$

(b)

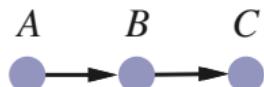
# Backward Chaining

```
function FOL-BC-ASK(KB, query) returns a generator of substitutions
  return FOL-BC-OR(KB, query, {})  
  
function FOL-BC-OR(KB, goal, θ) returns a substitution
  for each rule in FETCH-RULES-FOR-GOAL(KB, goal) do
    (lhs  $\Rightarrow$  rhs)  $\leftarrow$  STANDARDIZE-VARIABLES(rule)
    for each  $\theta'$  in FOL-BC-AND(KB, lhs, UNIFY(rhs, goal, θ)) do
      yield  $\theta'$   
  
function FOL-BC-AND(KB, goals, θ) returns a substitution
  if  $\theta = \text{failure}$  then return
  else if LENGTH(goals) = 0 then yield  $\theta$ 
  else
    first, rest  $\leftarrow$  FIRST(goals), REST(goals)
    for each  $\theta'$  in FOL-BC-OR(KB, SUBST(θ, first), θ) do
      for each  $\theta''$  in FOL-BC-AND(KB, rest, θ') do
        yield  $\theta''$ 
```

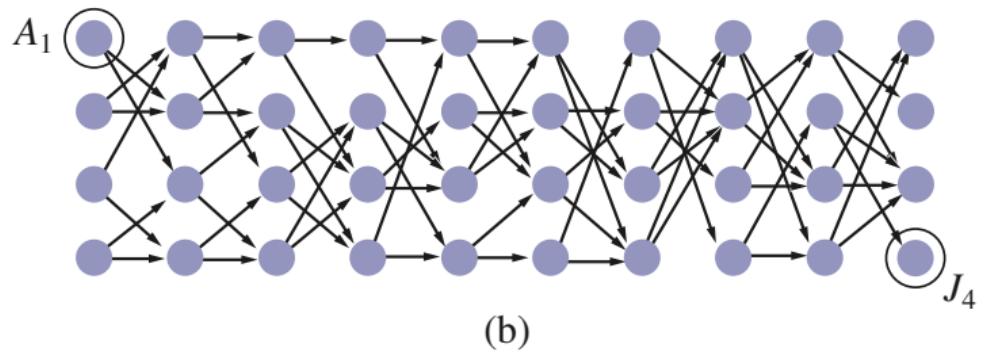
## Backward Chaining



# Logic Programming

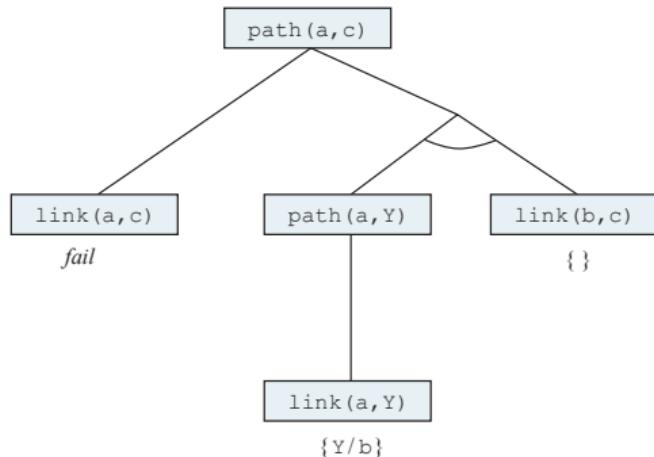


(a)

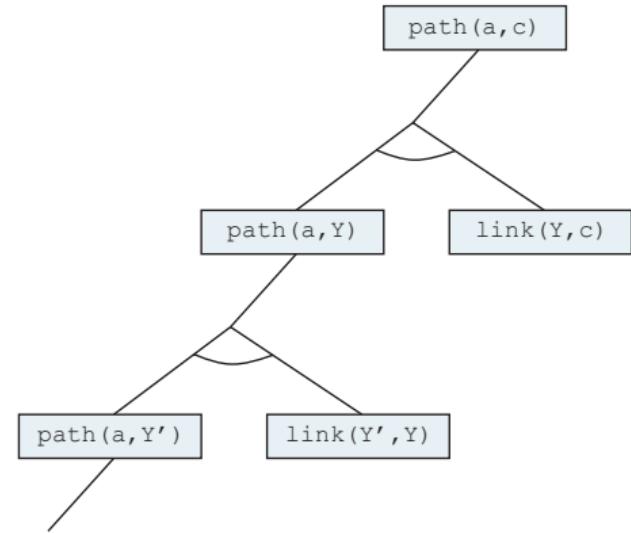


(b)

# Logic Programming

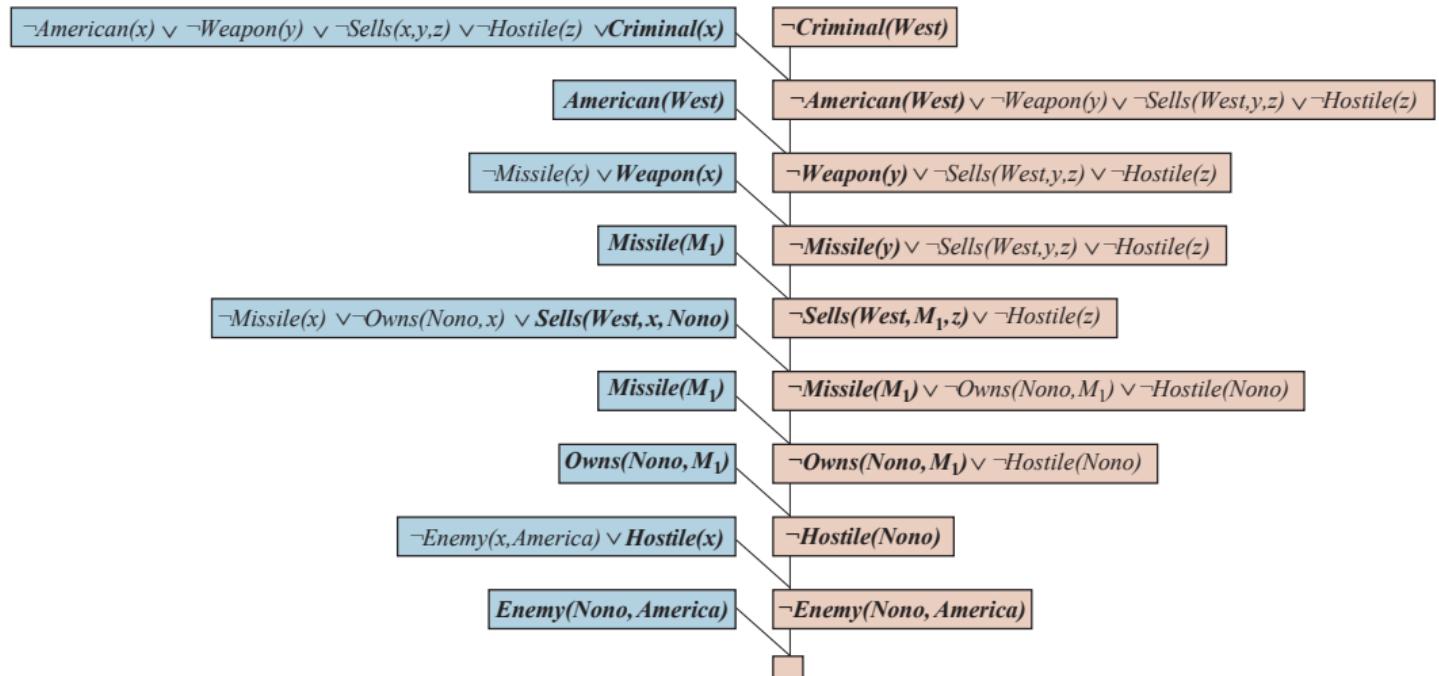


(a)

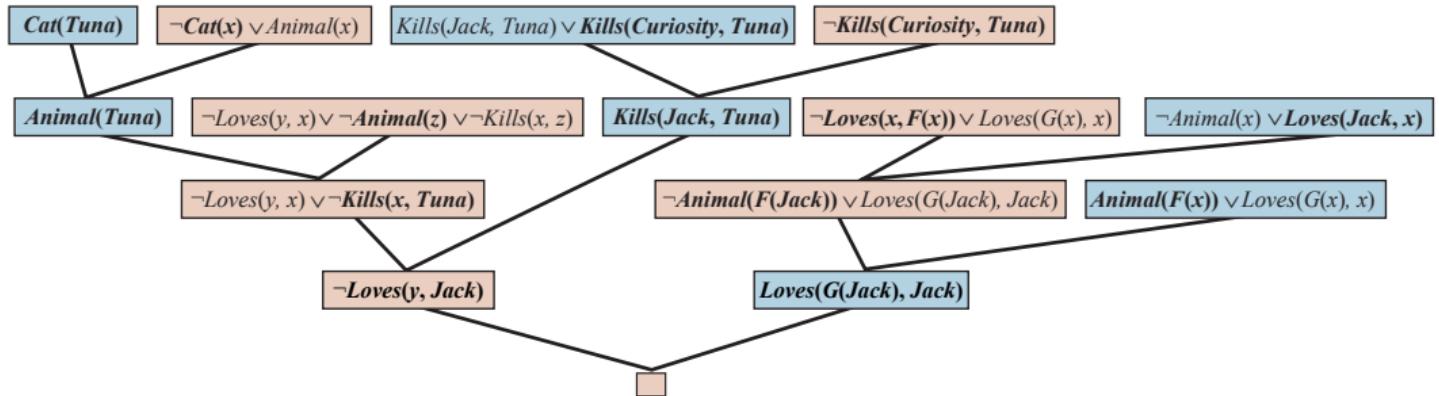


(b)

# Resolution



# Resolution



# Completeness

Any set of sentences  $S$  is representable in clausal form



Assume  $S$  is unsatisfiable, and in clausal form



Herbrand's theorem

Some set  $S'$  of ground instances is unsatisfiable



Ground resolution theorem

Resolution can find a contradiction in  $S'$



Lifting lemma

There is a resolution proof for the contradiction in  $S'$

# Gödel's Incompleteness Theorem

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