

Project Report

Self-Organized Criticality in Sandpile Models



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Abstract

blah blah

Acknowledgements

blah blah

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Introduction and Motivations

1.1 to do list

Individual contributions Introduction and Motivations Description of the Model Implementation Simulation Results and Discussion Summary and Outlook Sandpile Model

The Sandpile Model

2.1 Bak-Tang-Wiesenfeld Model

The classical sandpile model represents a cellular automation describing a dynamical system following certain rules that can be described as follows.

The field/lattice, which is chosen to be two-dimensional, represents a sandpile. Each site on the lattice has a certain value z that intuitively represents the height or slope of the sandpile at certain position described with the coordinates x and y. At each time step, a number of grains of sand is placed on top of a random site, which increases its value by a given value, e.g. one. If the value of the site exceeds a critical value z_c (e.g. three), the site collapses/topples and its grains are evenly distributed to its neighbours.

In certain cases some of the adjascent sites will exceed the critical value too and the toppling process will continue until an equilibrium state is again reached. This series of collapsing sites is clasically described as an avalanche. The next grain is not placed until the equilibrium state is reached, meaning that the time scale of the random grain placement and of the development of avalanches are decoupled.

The classical model description can mathematically be represented as follows.

Initially, the lattice is empty:

$$z(x,y) = 0 \quad \forall x, y$$

Then, the value of a random site x, y is increased:

$$z(x,y) \rightarrow z(x_r,y_r) + 1$$

If its value exceeds the critical value $z_c = 3$, then it topples and distributes its grains to its neighbours:

$$z(x,y) \stackrel{?}{>} 3 \Rightarrow z(x,y) \rightarrow z(x,y) - 4$$
$$z(x \pm 1, y) \rightarrow z(x \pm 1, y) + 1$$
$$z(x, y \pm 1) \rightarrow z(x, y \pm 1) + 1$$

Clearly, many variations of the described model can be considered and can produce different results. The classical sandpile model, as originally described by Per Bak, Chao Tang and Kurt Wiesenfeld, represents the starting point of any further investigations considered in this paper.

2.2 Parameters

The behavior of the model is analysed dependent on different parameters such as:

- lattice size
- number of dimensions of lattice
- mass conservation, i.e. if the number of grains removed from a collapsed site is equal to the sum of grains its neighbour sites received
- boundary conditions, see below
- etc.

Different types of boundary conditions can be thought of:

- open: If a site near the border topples, some of its grains leave the system (mass is lost).
- closed: Near-border site does not fully collapse, but keeps the grains that would fall off in an open case.
- periodic: The system has no boundaries, i.e. toppling near the border is "wrapped over".
- mixed: E.g. the lattice is periodic in one dimension and has open boundaries in another dimension.

ASDF ILLUSTRATE DIFFERENT BOUNDARY CONDITIONS

2.3 Abelian Model

One important property which can be used to categorize different sandpile models is whether they behave in a commutative or *abelian* way. In particular, this can be applied to the development of avalanches in the model described above: The question posed here is, whether an equilibrium state resulting from an avalanche depends on the way the avalanche is calculated. More precisely, it can be shown that any avalanche, being a sequence of topplings, always results in the same equilibrium state i.e. does not depend

on the order, in which the topplings occur. The mathematical proof of this hypothesis is nicely presented in [7].

To illustrate this practical but not necessarily obvious fact, a sample 4x4-field with one active site is considered (see figure 2.1). At step (2), two different sites simultaneously become active, therefore creating a "choice", which site to topple first. Depending on such choices, different sequences of topplings occur, but all lead to the same equilibrium state.

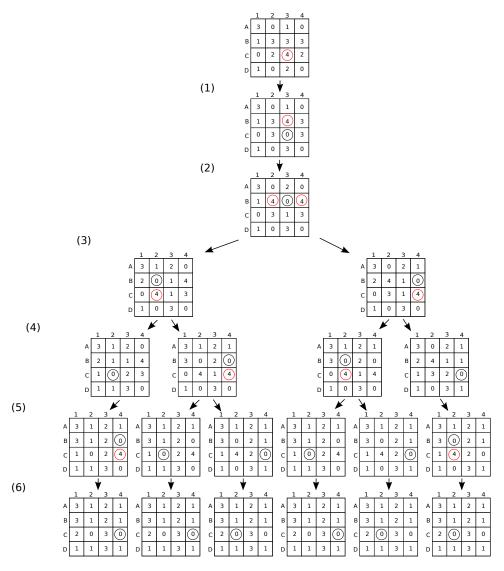


Figure 2.1: Demonstration of the abelian property: six different orders of topplings all lead to the same equilibrium state. The sites circled red indicate sites that have just become active, those circled black have just collapsed. Here, continuous boundary conditions have been used.

Model Implementation in MATLAB/Octave

Note on Code and Programming Sustainability

In order to produce "sustainable" code and to share the spirit of independency coming from the open-source community, the coded routines were tested in MATLAB and in Octave (one of its open-source clones). The source code can be found in Appendix A.

3.1 Basic Sandpile Code

First, a lattice/field is generated using uniformly distributed random numbers from 0 to z_c (critical_state). This is done in order to start with a potentially critical field and not to place single grains of sand until a site gets critical.

```
f = floor(unifrnd(0,critical_state+1,height,width));
```

Another interesting starting point is a *uniform* critical field, where every site is either 0 or z_c .

```
f = floor(unifrnd(0,2,height,width))*critical_state;
```

When the field is ready, a global loop runs through a defined number of timesteps, placing a grain on a random site, checking if the site becomes active and if so, computing the resulting avalanche.

```
end end
```

3.2 Simple Avalanche Code

...ASDF...

3.3 Optimization of Avalanche Code

The simple avalanche code checks the whole field including the fields, that cannot possibly be affected by the avalanche. It can therefore be optimized, for example using a LIFO data structure – a *stack*. The coordinates of very site that needs to be checked are placed on the stack, so that the computation of the avalanche consists of working through the stack and toppling all the active sites in it. During their toppling, their neighbours are again put on the stack, which makes the procedure dynamical and not easily comprehensive. The algorithm is illustrated in figure 3.1.

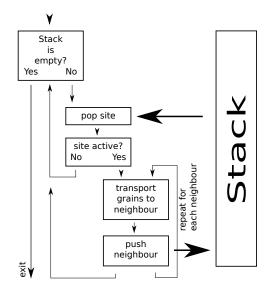


Figure 3.1: using a stack for avalanche calculation

Considering the example from figure 2.1, the stack algorithm results in the following sequence:

- 0. push C3
- 1. pop C3, topple, push its neighbours (C2,C4,B3 and D3) to stack
- 2. pop B3, topple, push B2, B4, A3 and C3 to stack

```
    3. pop B2, ...
    4. pop C2, ...
    5. pop B4, ...
    6. pop C4, ...
```

Figure 3.2 shows the states of the stack after each of these steps. To avoid confusion, only the active sites are shown here.

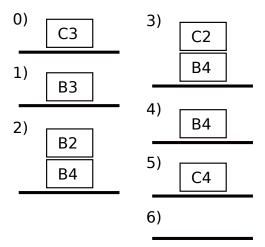


Figure 3.2: A sample sequence of stack states. The algorithm proceeds until the stack is empty.

The main loop including the stack feature looks like this:

```
for t=1:timesteps
    % choose random site % ...
     % place grain
    % push site to stack
stack_n = 1;
stack_x(1) = x;
stack_y(1) = y;
     % avalanche — work through stack
     while (stack_n > 0)
          % pop from stack
x = stack_x(stack_n);
          y = stack_y(stack_n);
          stack_n = stack_n - 1;
          % check if overcritical/active
          if (f(y,x) > critical_state)
               % collapse/topple
               f(y,x) = f(y,x) - neighbours * collapse;
               % look at every neighbour
               for n=1:neighbours
                     % add/transport grain to neighbour
                    f(y+neighbour_offset_y(n),x+neighbour_offset_x(n)) = ...
    f(y+neighbour_offset_y(n),x+neighbour_offset_x(n)) + collapse;
```

8 3.4. Statistics

```
% push neighbour to stack
    stack_n = stack_n + 1;
    stack_x(stack_n) = x + neighbour_offset_x(n);
    stack_y(stack_n) = y + neighbour_offset_y(n);
    end
    end
end
end
```

3.4 Statistics

Many different variables may be of interest for the statistical analysis of sandpile models. The easiest to implement is avalanche size:

```
% check if overcritical/active
if (f(y,x) > critical_state)
% collapse/topple
f(y,x) = f(y,x) - neighbours * collapse;
% record statistics
avalanche_sizes(t) = avalanche_sizes(t) + 1;
```

Here, the number of avalanches is recorded at every time step by increasing the counter after each toppling that happens during the avalanche. After the main loop, the data is sorted and the distribution is fitted into a power-law distribution given by

$$P(s) = a \cdot s^b$$

where P is the number of avalanches of size s. The coefficients a and b are determined using a simple solver that minimizes $a \cdot s^b - P(s)$.

```
% count avalanche sizes - calculate distribution
for s=1:max(avalanche.sizes)
    avalanche.count(s) = size(avalanche.sizes(avalanche.sizes==s),2);
end
% filter zero values
s = [1:max(avalanche.sizes)];
P = avalanche.count(1:end);
s = s(P>0);
P = P(P>0);
% fit into power—law
[c,fval,info,output]=fsolve(@(c)((c(1).*s.^c(2))-P),[100,1]);
a = c(1);
b = c(2);
```

Simulation Results

Summary and Outlook

12 Bibliography

Bibliography

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- [7] Frank Redig Ronald Meester and Dmitri Znamenski. The abelian sandpile; a mathematical introduction. April 2001.
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Appendix A

MATLAB/Octave-Code

A.1 critical field

make_pictures draw and export all frames or not silent produces no output (except time progress) if true driving_plane_reduction percentage of field close to the boundary not to be affected by driving (putting grains)

A.2 sandpile

14 A.2. sandpile

```
= 0 => use whole field (default)
= 0.2 => put grains at least 0.2*width
and 0.2*height far away from boundary
> 0.5 => invalid [!]
% OUTPUTS
                       avalanche sizes (topplings count) for each timestep size at avalanche—starting—site for eacg t avalanche lifetime for each t final field
     as
     nc
     at
final
% translate parameters
     width = size(f,2);
height = size(f,1);
neighbours = size(neighbour,2);
neighbour_offset_x = neighbour(1,:);
neighbour_offset_y = neighbour(2,:);
collapse = collapse_per_neighbour;
boundary = boundary_type;
                                                           % number ofneighbours to collapse to
% define stack for avalanches
stack_x = 0;
stack_y = 0;
stack_n = 0;
      % avalanche statistics
      avalanche statistics
avalanche.sizes = zeros(1, timesteps);
av_begin.t = zeros(1,timesteps);
avalanche.add = zeros(1,timesteps); % = av_size - av_ltime
% show starting field
if (silent==false)
    disp('starting from this field:');
    disp(f);
      % choose random site y=floor(unifrnd(1,height*(1-2*driving_plane_reduction))+height*driving_plane_reduction); x=floor(unifrnd(1,width*(1-2*driving_plane_reduction))+width*driving_plane_reduction); % unifraction  
           % place grain f(y,x) = f(y,x) + 1;
            % communicate
           if (silent==false)
    disp(['random grain on x' num2str(x) ',y' num2str(y)]);
              save picture of field before collapsing (with active field)
            if (make_pictures)
    draw.field(f,2);
    print(['field' num2str(t) '.png'],'-dpng');
           % push site to stack
stack.n = 1;
stack.x(1) = x;
stack.y(1) = y;
            % save avalanche starting site
           av_begin_x = x;
av_begin_y = y;
av_begin_t(t) = 0; % # topplings at avalanche starting site
            % avalanche - work through stack
            while (stack_n > 0)
                 % pop from stack
                 x = stack_x(stack_n);
y = stack_y(stack_n);
stack_n = stack_n - 1;
                 % display current site
if (silent==false)
    disp(['current site: x ' num2str(x) '; y ' num2str(y)]);
                  end
```

```
% check if overcritical/active if (f(y,x) > critical\_state)
            % communicate collapsing
if (silent==false)
    disp('collapse!');
            end
             % collapse/topple f(y,x) = f(y,x) - neighbours * collapse;
            \$ count future topplings to be caused by this toppling future topplings = 0;
             % look at every neighbour
             for n=1:neighbours
                         % communicate
if (silent==false)
                                      disp(['neighbour ' num2str(n)]);
                         %%%%% check boundary %%%%%%
                                      \mbox{\% 1)} no-boundary conditions (continuous field, pack-man style) if (boundary == 1)
                                                   % modify neighbour offsets if (y+neighbour_offset_y(n) < 1)
                                                               neighbour_offset_y(n) = neighbour_offset_y(n) + height;
                                                  if (y+neighbour_offset_y(n) > height)
   neighbour_offset_y(n) = neighbour_offset_y(n) - height;
                                                   end
                                                   if (x+neighbour_offset_x(n) < 1)
    neighbour_offset_x(n) = neighbour_offset_x(n) + width;</pre>
                                                            (x+neighbour_offset_x(n) > width)
neighbour_offset_x(n) = neighbour_offset_x(n) - width;
                                                   \ \mbox{add/transport grain to neighbour} \ \mbox{f(y+neighbour_offset_x(n)) = f(y+neighbour_offset_x(n)) = f(y+neighbour_offset_x
                                                    % push neighbour to stack
                                                   stack_n = stack_n + 1;
stack_x(stack_n) = x + neighbour_offset_x(n);
stack_y(stack_n) = y + neighbour_offset_y(n);
                                                   % count future topplings to be caused by this toppling
if (f(y+neighbour_offset_y(n),x+neighbour_offset_x(n)) == (critical_stat
    future_topplings = future_topplings + 1;
                                      % 2) energy loss at boundary (table style) elseif (boundary == 2)
                                                   % add/transport grain to neighbour
f(y+neighbour_offset_y(n),x+neighbour_offset_x(n)) = f(y+neighbour_offset_x(n))
                                                               % push neighbour's neighbours to stack
stack_n = stack_n + 1;
stack_x(stack_n) = x + neighbour_offset_x(n);
stack_y(stack_n) = y + neighbour_offset_y(n);
                                                               % count future topplings to be caused by this toppling
if (f(y+neighbour_offset_y(n),x+neighbour_offset_x(n)) == (critical_
    future_topplings = future_topplings + 1;
```

```
% calculate additional topplings caused
if (future_topplings > 0)
    avalanche_add(t) = avalanche_add(t) + future_topplings - 1;
                             % communicate additional topplings to come if (silent==false) disp(['this collapse generates ' num2str(future_topplings - 1) ' additional end
                        else
                              % communicate additional topplings to come
                             if (silent==false)
   disp(['this collapse generates no additional topplings']);
                             end
                        end
            end
            end
            % display field after collapsing
if (silent==false)
                 disp(f);
disp('');
           end
     % return avalanche sizes
as = avalanche_sizes;
     \mbox{\%} return number of topplings at avalanche starting site nc = av_begin_t;
     % return final state final = f;
        return avalanche lifetimes
          = avalanche_sizes - avalanche_add;
end
```

A.3 avalanche distribution analysis

```
yy2 = avalanche_count2(1:end);
xx2 = xx2(yy2>0);
yy2 = yy2(yy2>0);
       % plot avalanche count vs size
       figure;
subplot(2,2,1);
plot(xx,yy,'marker','s');
       % fit the curve into power law distribution (f = c1*x^c2) [c, fval, info, output]=fsolve(@(c)((c(1).*xx.^c(2))-yy),[100,1]);
       hold on;
plot(xx,c(1).*xx.^c(2),'r');
xlabel('avalanche size s');
ylabel('avalanche count P(s)');
title(['avalanche distribution and power—law—fit P(s)=' num2str(c(1)) '*s^ ' num2str(c(2))]);
       noid on;
loglog(xx,c(1).*xx.^c(2),'r');
xlabel('avalanche size s');
ylabel('avalanche count P(s)');
title(['avalanche distribution and power-law-fit P(s)=' num2str(c(1)) '*s^ ' num2str(c(2))]);
       % return coefficients
       a = c(1);

b = c(2);
       % plot avalanche count vs lifetime
subplot(2,2,3);
plot(xx2,yy2,'marker','s');
       % fit the curve into power law distribution (f = c1*x^c2) [c,fval,info,output]=fsolve(@(c)((c(1).*xx2.^c(2))-yy2),[100,1]); hold on;
       noid on;
plot(xx2,c(1).*xx2.^c(2),'r');
xlabel('avalanche lifetime t');
ylabel('avalanche count P(t)');
title(['avalanche distribution and power-law-fit P(t)=' num2str(c(1)) '*t^ ' num2str(c(2))]);
       % same on a log-log-scale plot
subplot(2,2,4);
loglog(xx2,yy2,'marker','s');
hold on;
       nold on;
loglog(xx2,c(1).*xx2.^c(2),'r');
xlabel('avalanche lifetime t');
ylabel('avalanche count P(s)');
title(['avalanche distribution and power-law-fit P(t)=' num2str(c(1)) '*t^ ' num2str(c(2))]);
       % return coefficients
a2 = c(1);
b2 = c(2);
end
```

A.4 test sandpile