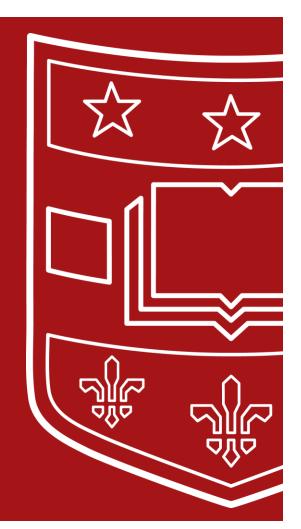
CSE 433S: Introduction to Computer Security

Network Defense Lecture





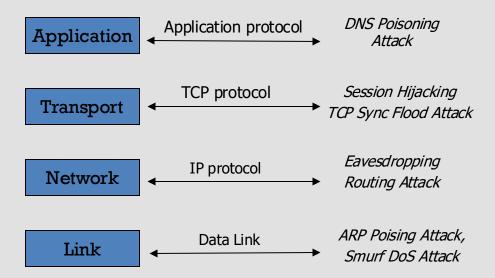
Review Questions

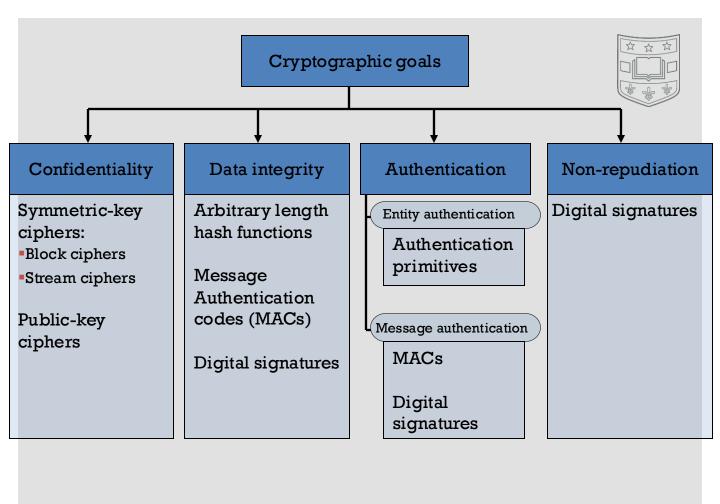


- How many symmetric keys does it take to support secure communications among 4 peers? What are the possible approaches to mitigate this issue?
- What is Diffie-Hellman key exchange, what attack is it vulnerable to, how do you launch that attack?
- What is public key crypto and how is different than symmetric key crypto? What key do Alice use if Alice wants to deliver a secret to Bob.
- In practice, can we use RSA to directly encrypt secret key for communication?
- What is digital signature? What key does Alice use to sign a file that she wants to authenticate and why?

Previously on Introduction to Computer Security – Network Attacks

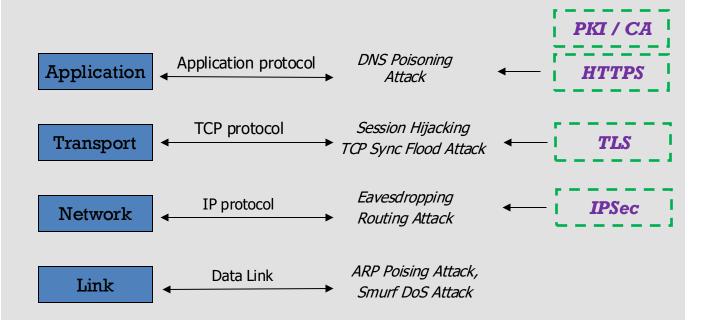






This module - Network defense





Network Defense



- Public key infrastructure (PKI)
 - How do we know we are talking to the right entity on the web
- Certificate authority (CA)
 - Mechanisms to establish the web of trust
- IPSec
 - Enables confidentiality and integrity protection for communication between two network nodes at IP layer
- TLS (Transport Layer Security)
 - Enables security communication link over TCP



PKI/CA

"Web of trust"



- Obtain public keys from friends in person
 - "Key-signing parties" from the PGP time
- Obtain "certificates" on my public key from my friends

- If A knows pk_B, and B issued a certificate for C, then C can send that certificate to A
 - What trust assumptions are being made here?

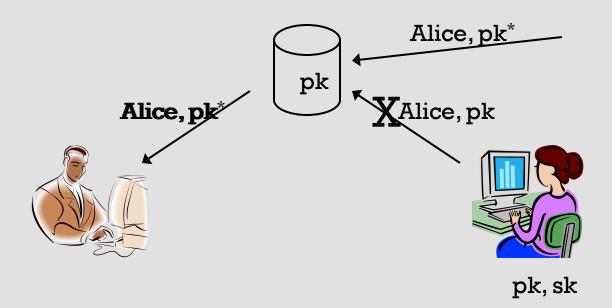
Public-key repository



- Store certificates in a central repository
 - E.g., MIT PGP keyserver
- To find Alice's public key
 - Get all public keys for "Alice," along with certificates on those keys
 - Look for a certificate signed by someone you trust whose public key you already have

Challenges on Public Key Repository





Public Key infrastructure

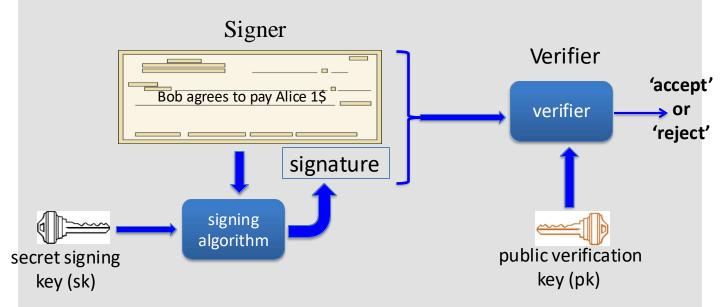


- Problem: How to determine the correct public key of a given entity
 - Binding between IDENTITY and PUBLIC KEY
- Possible attacks
 - Name spoofing: Eve associates Alice's name with Eve's public key
 - Key spoofing: Eve associates Alice's key with Eve's name
 - DoS: Eve associates Alice's name with a nonsensical (bogus) key

Digital signatures



Solution: make signature depend on document





Use signatures for secure key distribution!



- Assume a trusted party with a public key known to everyone
 - CA = certificate authority
 - Public key pk_{CA}
 - Private key sk_{CA}

Use signatures for secure key distribution!



- Alice asks the CA to sign the binding (Alice, pk)
 - $\operatorname{cert}_{CA \to Alice} = \operatorname{Sign}_{sk_{CA}}(Alice, pk)$
- (CA must verify Alice's identity out of band)

Use signatures for secure key distribution!



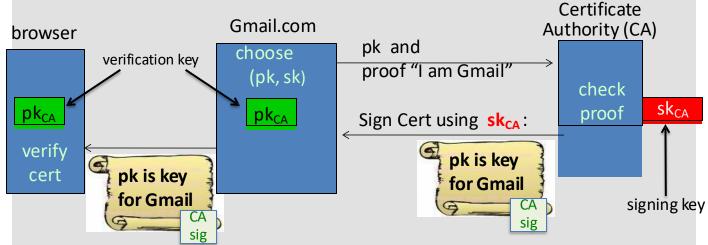
- Bob obtains Alice, pk, and the certificate cert_{CA→Alice} ...
 - ... verifies that $Vrfy_{pK_{CA}}((Alice, pk), cert_{CA \rightarrow Alice}) = 1$
- Bob is then assured that pk is Alice's public key
 - As long as the CA is trustworthy...
 - · Honest, and properly verifies Alice's identity
 - and the CA's private key has not been compromised



Important application: Certificates

Problem: browser needs server's public-key to setup a session key

Solution: server asks trusted 3rd party (CA) to sign its public-key pk



Server uses Cert for an extended period (e.g. one year)

Dan Boneh

Chicken-and-egg problem?



- How does Bob get pk_{CA} in the first place?
- Several possibilities...

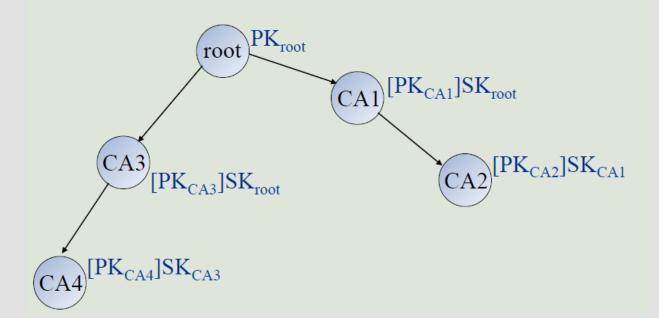
"Roots of trust"



- Bob only needs to securely obtain a small number of CA's public keys
 - Need to ensure secure distribution only for these few, initial public keys
- E.g., distribute as part of an operating system, or web browser
 - Firefox:
 Tools->Options->Privacy & Security->View
 certificates->Authorities

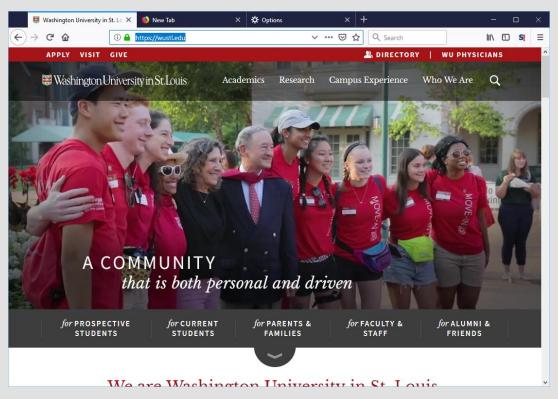
Certificate Hierarchy





Is WUSTL secure?





Certificate Viewer: "wustl.edu"

×



This certificate has been verified for the following uses:

SSL Client Certificate

SSL Server Certificate

Issued To

Common Name (CN) wustl.edu

Organization (O) Washington University in St. Louis

Organizational Unit (OU) WUIT

Serial Number 6A:A4:83:F6:AF:57:FA:A6:0A:08:00:28:0B:AB:3A:9A

Issued By

Common Name (CN) InCommon RSA Server CA

Organization (O) Internet2
Organizational Unit (OU) InCommon

Period of Validity

Begins On Sunday, September 24, 2017
Expires On Thursday, September 24, 2020

Fingerprints

SHA-256 Fingerprint 09:46:35:7D:20:0D:7D:E0:C5:F2:2C:31:0F:EB:86:EA:

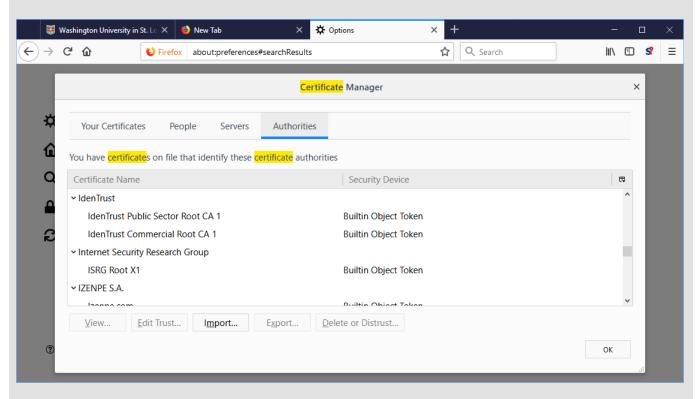
E6:2F:74:B0:1F:11:96:89:3C:72:A0:0F:08:1F:7A:41

SHA1 Fingerprint 0F:5E:04:57:05:2B:93:9A:ED:99:37:1D:E7:A3:09:D0:3D:2B:3C:48



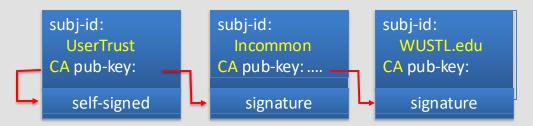
What else are we trusting?

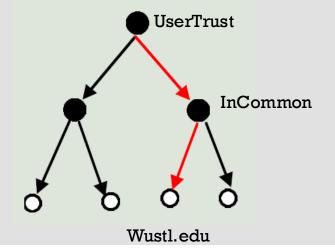




Certificate Path





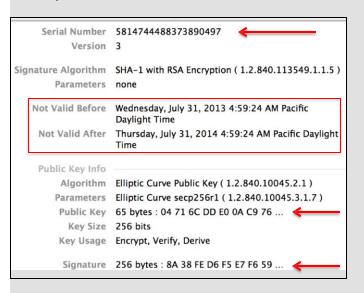


- Verifier has to know the public key of root CA
- Other public keys can be discovered
- All CA on the path has be trusted by the verifier



Certificates: example

Important fields:





PKI in practice... Many problems



- Does not work quite as well as in theory...
 - Proliferation of root CAs
 - Revocation
 - Common Name Issue

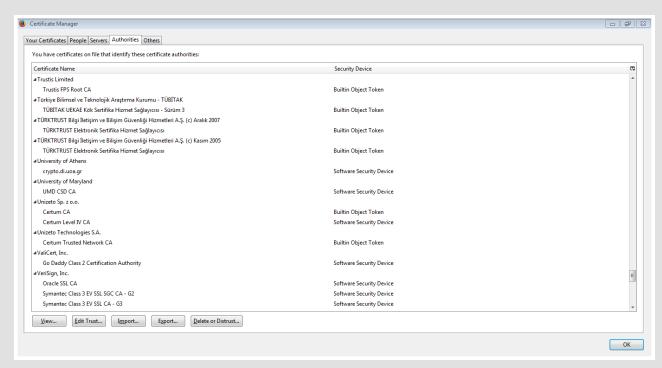
Something you trust



Certificate Viewer: "Verisign Class 1 Public Primary Certification Authority - G3" X General Details Could not verify this certificate because the issuer is unknown. Issued To Common Name (CN) VeriSign Class 1 Public Primary Certification Authority - G3 Organization (O) VeriSign, Inc. Organizational Unit (OU) VeriSign Trust Network 00:8B:5B:75:56:84:54:85:0B:00:CF:AF:38:48:CE:B1:A4 Serial Number Issued By Common Name (CN) VeriSign Class 1 Public Primary Certification Authority - G3 Organization (O) VeriSign, Inc. Organizational Unit (OU) VeriSign Trust Network Period of Validity Thursday, September 30, 1999 Begins On Expires On Wednesday, July 16, 2036 Fingerprints CB:B5:AF:18:5E:94:2A:24:02:F9:EA:CB:C0:ED:5B:B8: SHA-256 Fingerprint 76:EE:A3:C1:22:36:23:D0:04:47:E4:F3:BA:55:4B:65 SHA1 Fingerprint 20:42:85:DC:F7:EB:76:41:95:57:8E:13:6B:D4:B7:D1:E9:8E:46:A5 Close

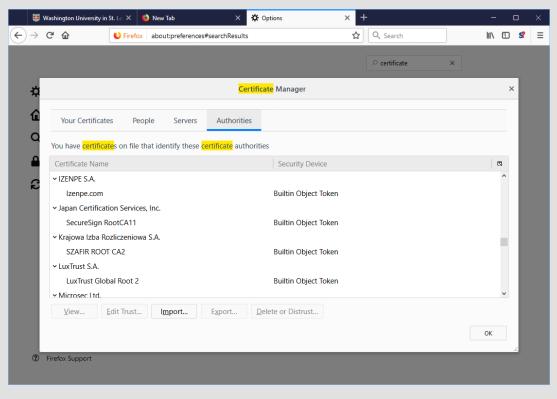
Do you trust them all?





A lot of people can sign, and this is just top level





The numbers





Browsers accept certificates from a large number of CAs:

- ■Top level CAs ≈ 60
- Intermediate CAs ≈ 1200





Revocation

Why do we need revocation?



- It could be
 - CA is compromised
 - Bob forget his private key
 - Bob's private key is stolen
 - Bob disclose his private key
- Certificate revocation needs to occur when:
 - certificate holder key compromise/loss
 - CA key compromise
 - end of contract (e.g. certificates for employees)

General Requirement of Revocation



- Timeliness
 - Before using a certificate, must check most recent revocation status
- Efficiency
 - Computation
 - Bandwidth and storage
 - Availability
- Security

Revocation Type



Implicit

- Each certificate is periodically issued
- Alice has a fresh certificate -> Alice not revoked
- No need to distribute/publish revocation info

Explicit

- Only revoked certificates are periodically announced
- Alice's certificate not listed among the revoked ->
 Alice not revoked
- Need to distribute/publish revocation info

Revocation Method



- CRL Certificate Revocation List
 - CRL-DP, indirect CRL, dynamic CRL-DP,
 - delta-CRL, windowed CRL, etc.
 - CRT and other Authenticated Data Structures
- OCSP On-line Certificate Status Protocol

CRS - Certificate Revocation System

Revocation for WUSTL



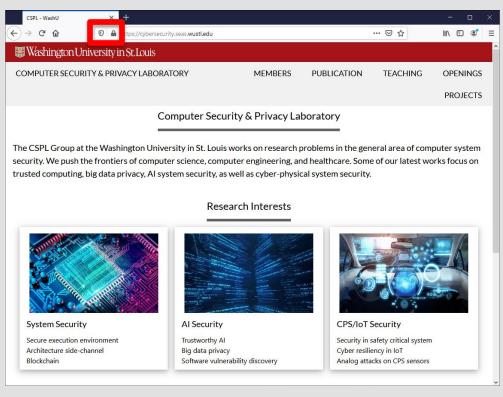
Certificate Viewer: "wustl.edu"	×
General Details	
Certificate <u>H</u> ierarchy	
✓ USERTrust RSA Certification Authority	
✓ InCommon RSA Server CA	
wustl.edu	
Certificate Fields	
Extended Key Usage	^
-Certificate Policies	
CRL Distribution Points	
Authority Information Access	
Certificate Subject Alt Name	
Certificate Signature Algorithm	
Certificate Signature Value	~
Field <u>V</u> alue	
Not Critical URI: http://crl.incommon-rsa.org/InCommonRSAServerCA.crl	
Export	
	<u>C</u> lose

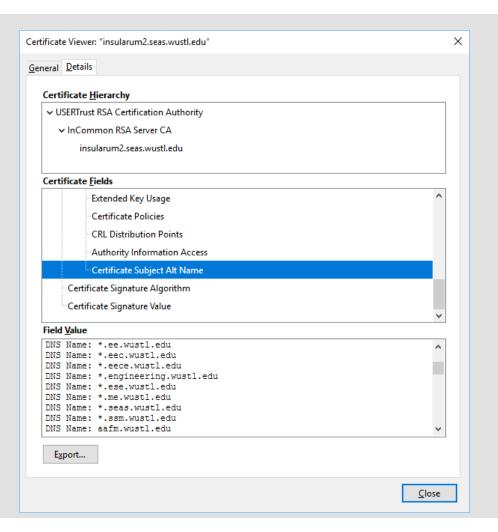


Alternative Names

My website is secure? How?



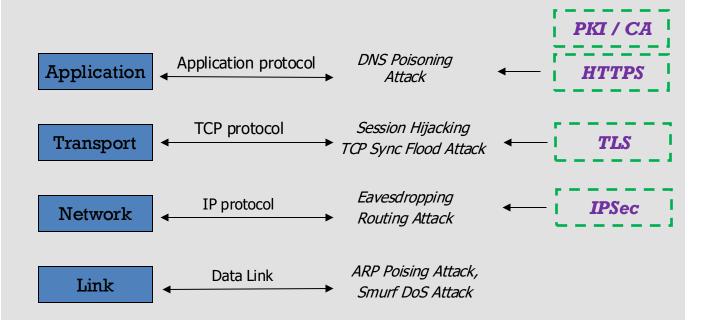






This module - Network defense



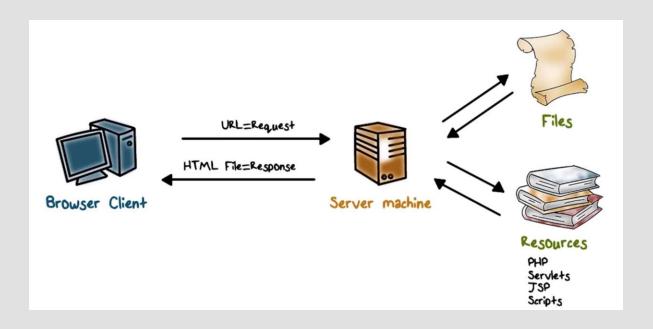




HTTPS/TLS

HTTP





Threat Model







Network Attacker:

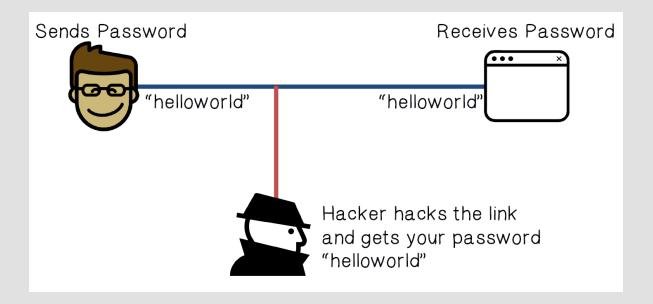
- Controls network infrastructure: Routers, DNS
- Eavesdrops, injects, blocks, and modifies packets

Examples:

- Wireless network at Internet Café
- Internet access at hotels (untrusted ISP)

HTTP





HTTPS



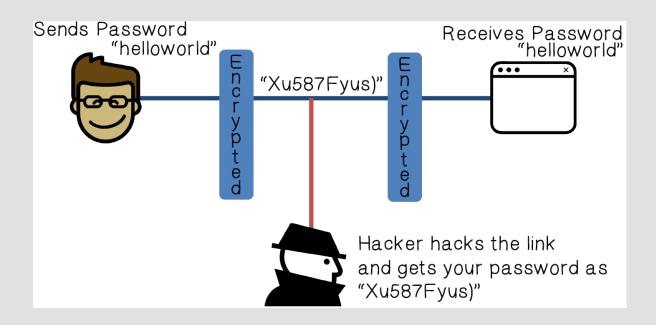


HTTPS: Hyper Text Protocol Secure= all communication between your browser and a website is encrypted.



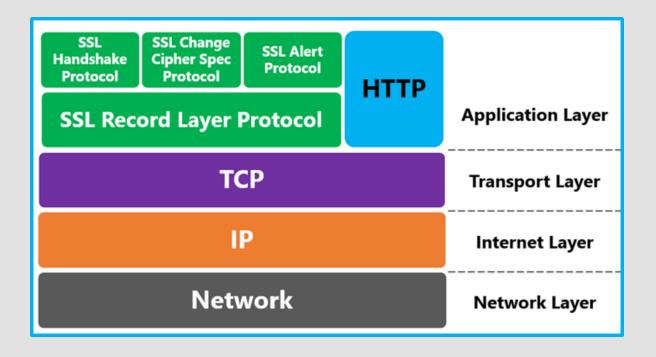
HTTPS





Where is HTTPS







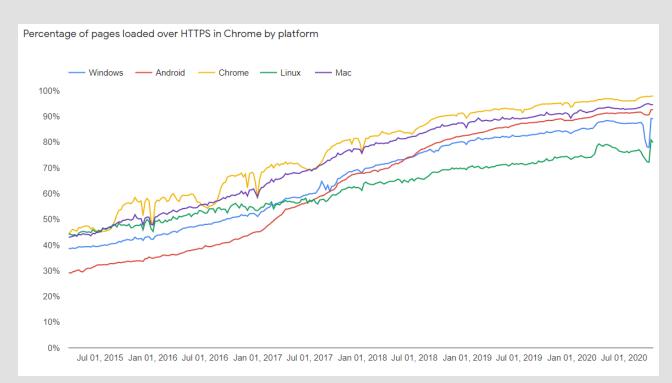


 Creates a secure channel over an insecure network

 Is an effective protection against man-in-themiddle attacks

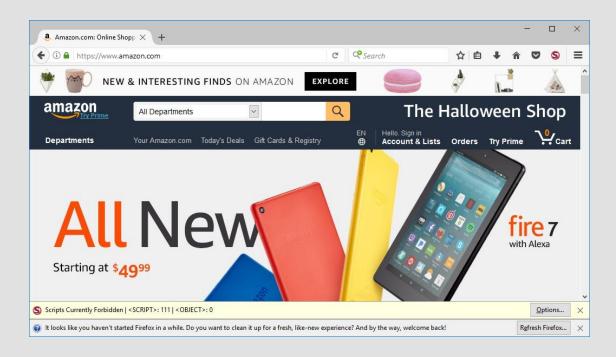
 Can still provide security even when only one side of the communication is secure





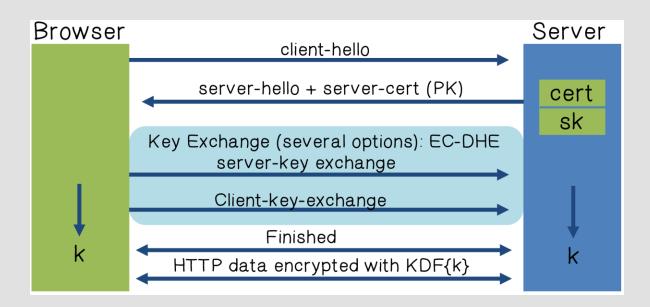
First few hundred millisecs when you go to amazon.com





TLS Summary

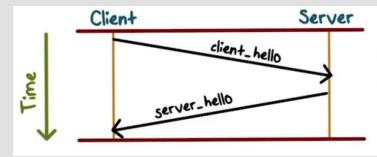




The Handshake Protocol







Phase 1

Establish security capabilities, including protocol version, session ID, cipher suite, compression method, and initial random numbers

The Handshake Protocol



The Parameters:

- Version: the highest TLS version understood by the client
- Random: a 32-bit timestamp and 28 bytes generated by a secure random number generator
- · Session ID: a variable-length session identifier
- CipherSuite: a list containing the combinations of cryptographic algorithms supported by the client
- Compression Method: a list of compression methods supported by the client

HTTPS -> Port 443



```
    Internet Protocol, Src: 172.17.30.63 (172.17.30.63), Dst: 7
    Transmission Control Protocol, Src Port: 50752 (50752), Dst
    Source port: 50752 (50752)
    Destination port https (443)
    Sequence number: 1 (relative sequence number)
    [Next sequence number: 164 (relative sequence number)]
    Acknowledgement number: 1 (relative ack number)
    Header length: 20 bytes
    Flags: 0x18 (PSH, ACK)
    Window size: 64860
```

Client Hello – Version



```
Content Type: Handshake (22)

Version: TLS 1.0 (0x0301)

Length: 158

□ Handshake Protocol: Client Hello

Handshake Type: Client Hello (1)

Length: 154

Version: TLS 1.0 (0x0301)

fd 5c e2 64 00 00 16 03 01 00 9e 01 00 00 9a 03 .\.d...

01 4a 2f 07 ca b9 4f bs 66 7a 06 56 7f ce c9 f7 .J/...o. .z.v....
```

TLSv1 Record Layer: Handshake Protocol: Client Hello

Client Hello - Random



```
■ Random
```

gmt_unix_time: Jun 9, 2009 21:09:30.000000000 random_bytes: B94FB3067A06567FCEC9F737BD5270F7002BB0D6723E551A..

01 4a 2f 07 ca b9 4f b3 06 7a 06 56 7f ce c9 f7 .]/...o. .z.v.... 37 bd 52 70 f7 00 2b b0 d6 72 3e 55 1a 0d 57 d9 7.Rp..+. .r>U..w. 82 00 00 44 c0 0a c0 14 00 88 00 87 00 39 00 38 ...b....9.8

Why do we need this?

Client Hello – Session ID



Session ID Length: 0

82 🔟 00 44 c0 0a c0 14 00 88 00 87 00 39 00 389.8

Client Hello – Cipher Suite



Client Hello – Server Name Extension



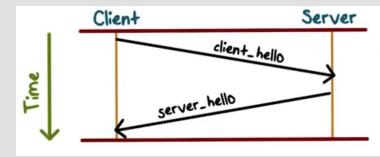
■ Extension: server_name								
							0 00 13	
00 11	00 00	0e 77	77 77	2e 61	6d 61		if 6e 2e	www .amazon.
63 6f	6d 00	0a 00	08 00	06 00	17 00	18 0	0 19 00	COM

Why do we need this?

The Handshake Protocol







Phase 1

Establish security capabilities, including protocol version, session ID, cipher suite, compression method, and initial random numbers

Server Hello



```
Handshake Protocol: Server Hello

Handshake Type: Server Hello (2)

Length: 70

Version: TLS 1.0 (0x0301)

Random

gmt_unix_time: Jun 9, 2009 21:09:30.000000000

random_bytes: 986BE7A3ACDD547D038235895ABA467EED6EB07CD46E5E8A..

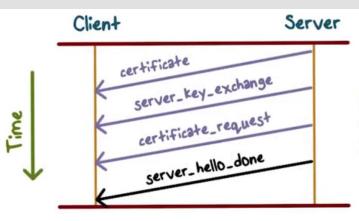
Session ID Length: 32

Session ID: ACD07167996609E3584D15C2D43596D8AB93CFA9E899A82A...

Cipher Suite: TLS_RSA_WITH_RC4_128_MD5 (0x0004)
```



The Handshake Protocol





Phase 2

Server may send certificate, key exchange, and request certificate. Server signals end of hello message phase.

Certificate Message



```
■ Handshake Protocol: Certificate

Handshake Type: Certificate (11)

Length: 2464

Certificates Length: 2461

□ Certificates (2461 bytes)

Certificate Length: 1271

□ Certificate (id-at-commonName=www.amazon.com,id-at-organizationN

□ signedCertificate

version: v3 (2)
```

```
      e8
      99
      a8
      2a
      a1
      65
      2b
      8a
      b6
      67
      a1
      2b
      00
      04
      00
      0b

      00
      09
      a0
      00
      04
      f7
      30
      82
      04
      f3
      30
      82
      03

      db
      a0
      03
      02
      01
      02
      02
      10
      16
      9d
      04
      1c
      31
      30
      be
      3d

      56
      66
      06
      f2
      67
      9b
      a1
      72
      30
      0d
      06
      09
      2a
      86
      48
      86

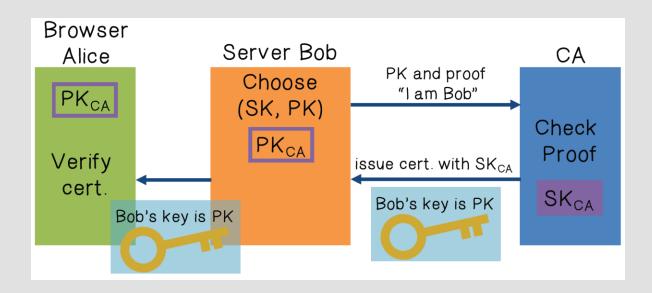
      f7
      0d
      01
      01
      05
      05
      00
      30
      81
      b0
      31
      0b
      30
      09
      06
      03

      55
      04
      06
      13
      02
      55
      53
      31
      17
      30
      15
      06
      03
      55
      04
      0a

      13
      0e
      56
      65
      72
      69
      53
      69
      67
      6e
      2c
      20
      49
```

What is a certificate?





Server Hello Done



```
Handshake Protocol: Server Hello Done
```

Handshake Type: Server Hello Done (14)

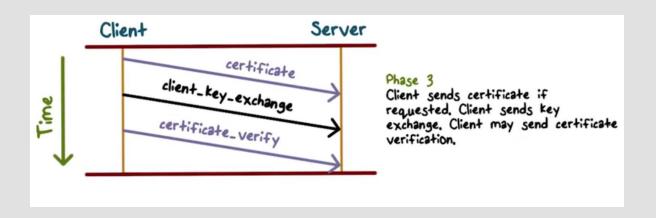
Length: 0

03 90 Oc Oe 00 00 00





The Handshake Protocol



Key Exchange



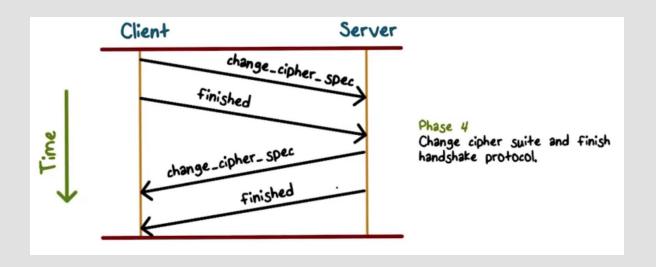
```
Handshake Protocol: Client Key Exchange
```

Handshake Type: Client Key Exchange (16)

Length: 130



The Handshake Protocol



Client Change Cipher Spec



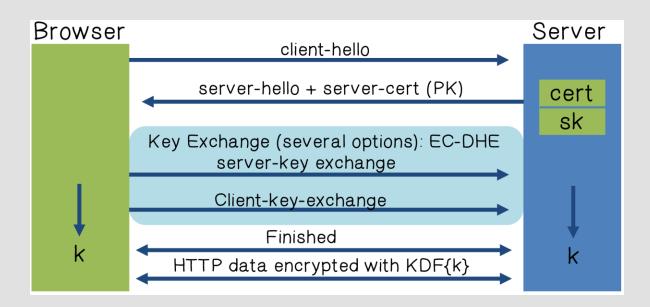
Finally, Encrypted Traffic



```
Secure Socket Layer
  TLSv1 Record Layer: Application Data Protocol: http://dx.doi.org/10.1006/j.jpub.2006/pdf
      Content Type: Application Data (23)
      Version TLS 1.0 (0x0301)
      Length: 466
      Encrypted Application Data: 7FEF3541D25C5F37F46461988729B6EF873D59F3
      00 1a a0 c4 28 48 00 17
                                    87 98 00 08 00 45 00
                                                              ....(H.. .....E.
      01 ff ba 50 40 00 f7 06
                                 e6 00 48 15 cf 41 ac 11
            01 bb c6 40 52
                     d1 f2 57
                                    70 fb 2e 26
               94 01 e9 66 d1
                                                 b6 5d 1b
                                 b0 07 49 f0 0a a6 56 c8
```

TLS Summary





Problems with HTTPS





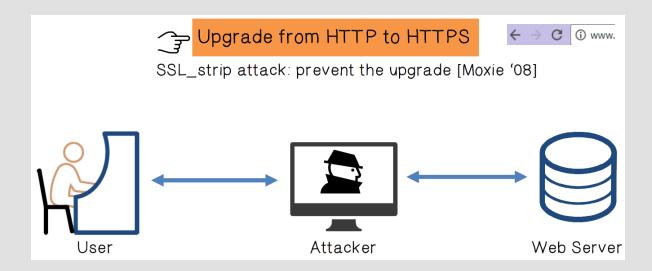
TUpgrade from HTTP to HTTPS

Forged certs

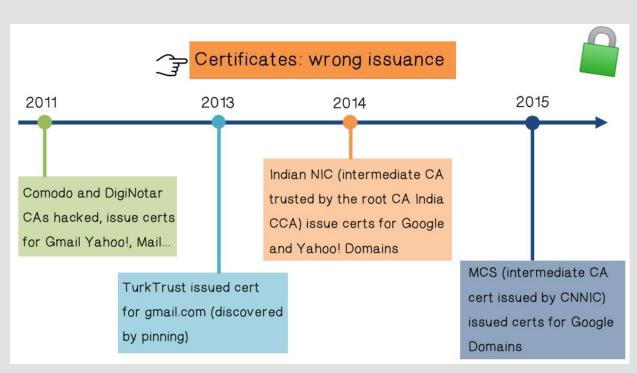
Mixed content: HTTP and HTTPS on the same page

HTTPS downgrading attack



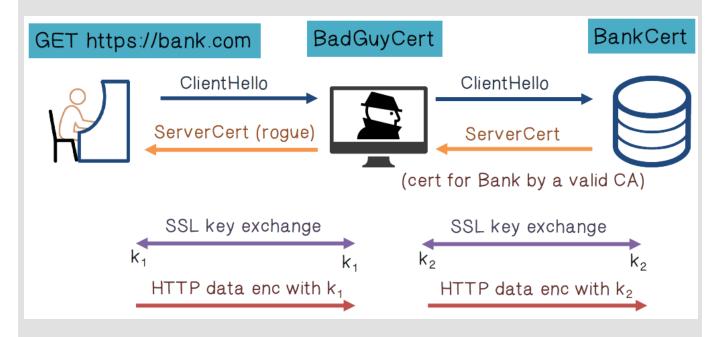






MIMT using rouge cert



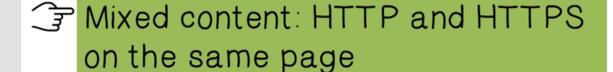


Problems with HTTPS



Upgrade from HTTP to HTTPS

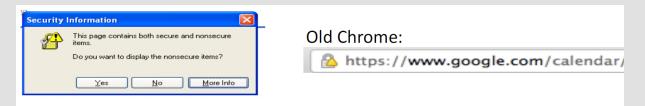
Forged certs



Mixed Contents: HTTP and HTTPS



- Page loads over HTTPS, but contains content over HTTP (e.g. <script src="http://.../script.js>)
- ⇒ Active network attacker can hijack session by modifying script en-route to browser



Mostly ignored by users ...



How do you attack TLS?

Padding Oracle Attack (BEAST)

Export Downgrade (FREAK)

Common Exponent (LogJam, CCS 15 Best Paper)

Side Channel Attack (Chen et al. SP 10)

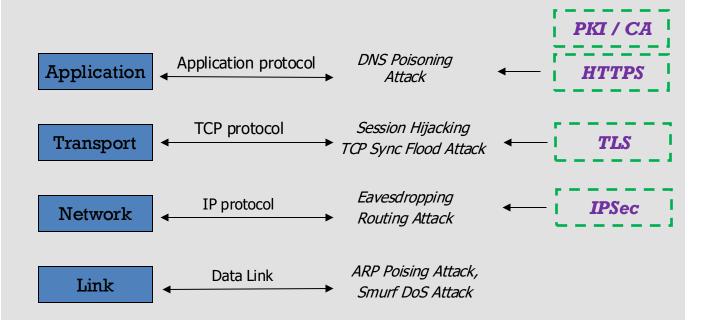
Summary



- TLS is one of the most fundamental methods to secure the internet
 - TLS client hello / server hello
 - Key exchange
- HTTPS attacks
 - Downgrade attack
 - Rouge certificates
- TLS attack
 - Padding oracle (BEAST, Lucky 13)
 - Export downgrade (FREAK)
 - DH component reuse (Logjam)
 - Side channel information leakage

This module - Network defense







IPSec

IP is not Secure!



- IP protocol was designed in the late 70s to early 80s
- Part of DARPA Internet Project
- Very small network
 - All hosts are known!
 - So are the users!
 - Therefore, security was not an issue

Security Issues in IP



- Source spoofing
- Replay packets
- No data integrity or confidentiality

- DOS attacks
- Replay attacks
- MiTM attack
- Interleaving attacks
- Eavesdropping
- and more...

Fundamental Issue: *Networks are not fully secure*

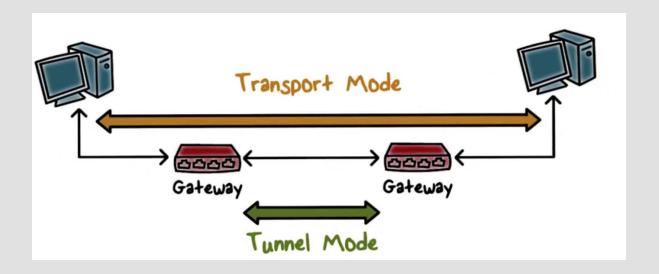
IPSec Architecture



- IPSec provides security in three situations:
 - Host-to-host
 - Host-to-gateway
 - Gateway-to-gateway
- IPSec operates in two modes:
 - Transport mode (for end-to-end, host-to-host)
 - Tunnel mode (for VPN)

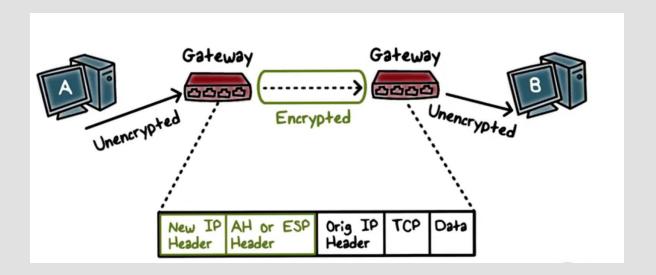
IPSec Modes - Tansport





IPSec Modes - Tunnel Mode





Blockchain



BLOCKCHAIN TECHNOLOGY STACK

