

Markov Categories and Entropy

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Utku Boduroğlu, Drew McNeely, and Nico Wittrock

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What Markov categories are not

Definition

A symmetric monoidal category with

objects: commutative comonoids $(X, \text{del}_X, \text{copy}_X)$

morphisms: morphisms of comonoids

What Markov categories are not

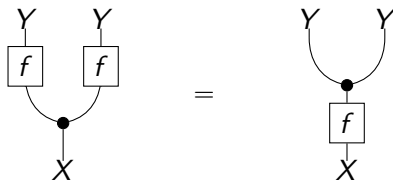
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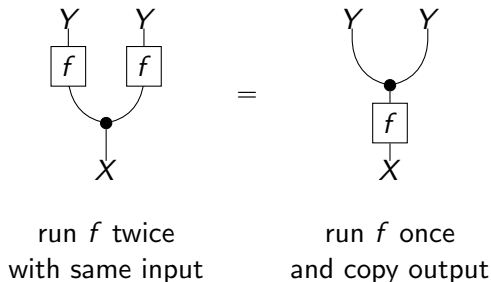
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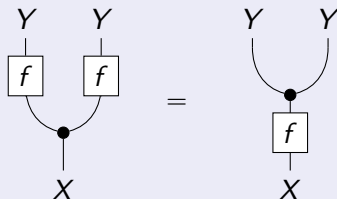
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Determinism in Markov categories

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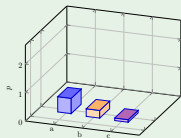
$f : X \rightarrow Y$ is deterministic if



Determinism in FinStoch

Recall: $p : I \rightarrow A := \{a, b, c\}$ is a distribution $\{p(a), p(b), p(c)\}$.

Example

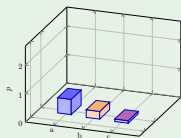


is not deterministic:

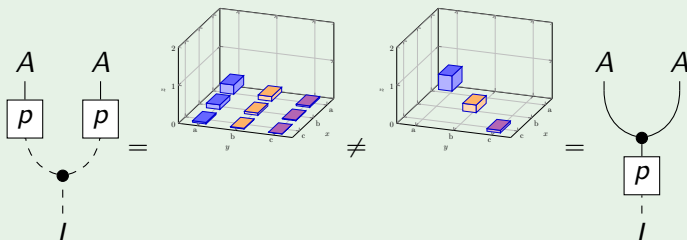
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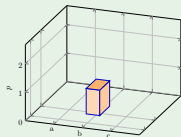
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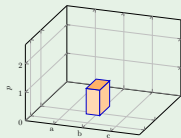
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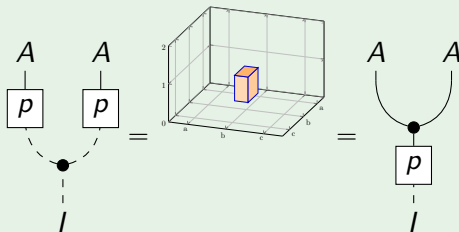
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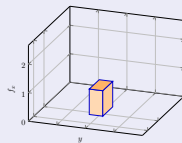


Determinism in FinStoch

Recall: $f : X \rightarrow Y$ has distributions f_x on Y .

Proof.

f is deterministic \Leftrightarrow all f_x of form
 $\Leftrightarrow f$ corresponds to Set-function



Determinism and Products

Theorem

Let \mathcal{C} be a Markov category.

The deterministic morphisms form Markov subcategory $\mathcal{C}_{\text{det}} \subset \mathcal{C}$.

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Example

$$\text{FinStoch} \supset \text{FinStoch}_{\text{det}} \cong \text{FinSet}$$

Products and Determinism

Theorem

Let \mathcal{C} be cartesian monoidal (i.e. \otimes is product functor).

\mathcal{C} is Markov.

All morphisms are deterministic.

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Example (Markov categories without uncertainty)

Set, FinSet, Meas, BorelMeas, ...

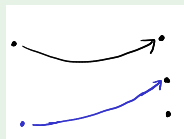
Probability Monads

...introduce uncertainty to cartesian Markov categories

Probability Monads

Example (Powerset Monad)

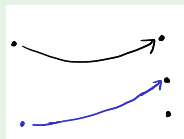
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Probability Monads

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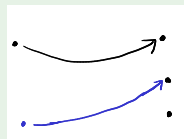


(slightly adjusted) powerset monad $2^\circ : \text{Set} \rightarrow \text{Set}$

Probability Monads

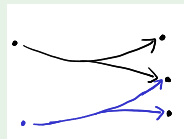
Example (Powerset Monad)

Set is deterministic



(slightly adjusted) powerset monad $2^\circ : \text{Set} \rightarrow \text{Set}$

Kleisli category Set_{2° describes possibility



Probability Monads

Theorem

For a cartesian monoidal category \mathcal{C} with monad $T : \mathcal{C} \rightarrow \mathcal{C}$ which is compatible with I (\rightsquigarrow powerset monad was adjusted) compatible with \otimes its Kleisli category \mathcal{C}_T is Markov.

Probability Monads

Example

Set



2^O



Poss

FinSet



2^O



FinPoss

Meas



Giry monad



Stoch

BorelMeas



Giry monad



BorelStoch

Probability Monads



Probability Monads

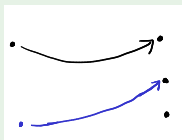


Uncertainty captured by morphisms!

Probability Monads

Example (Powerset Monad)

Set is cartesian monoidal \Rightarrow morphisms are deterministic

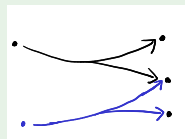


Powerset monad $2^\bullet : \text{Set} \rightarrow \text{Set}$,

$$X \mapsto 2^X := \{M \subseteq X\}$$

Kleisli category Set_{2^\bullet} has morphisms:

$$f : X \rightarrow 2^Y \quad \text{i.e. } f(x) \subseteq Y$$



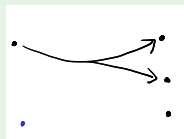
As Pablo stated: describes possibility

Probability Monads

Example (Powerset Monad)

Problem: Kleisli category Set_2^\bullet has morphism:

$$f : X \rightarrow 2^Y \quad \text{with} \quad f(x_i) = \emptyset$$

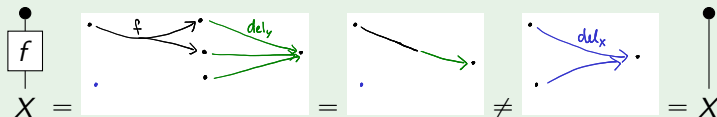
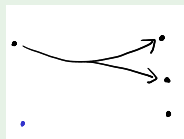


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$\Rightarrow \text{Set}_2^\bullet$ is not Markov!

$$\rightsquigarrow \text{adjust monad} \quad T : \text{Set} \rightarrow \text{Set},$$

$$X \mapsto 2^X - \emptyset$$

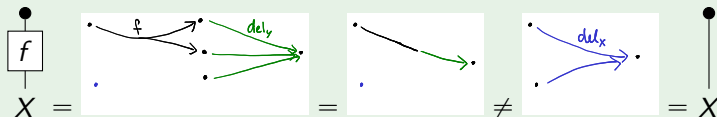
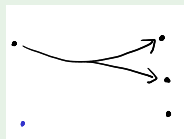
$\Rightarrow \text{Set}_T =: \text{Poss}$ is Markov.

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\rightsquigarrow adjust monad $T : \text{Set} \rightarrow \text{Set}$,

$$X \mapsto 2^X - \emptyset$$

$\Rightarrow \text{Set}_T =: \text{Poss}$ is Markov. With uncertainty.

Outline

This is just me fooling around with Beamer to get a feel for it. None of the contents should be taken seriously at this point.

- 1 Introduction to Markov Categories
- 2 Useful ways to construct Markov categories
 - Randomness in Markov categories
 - Introducing Uncertainty by Probability Monads
- 3 Properties and additional axioms
- 4 Bonus Slides
 - Powerset monad
- 5 Markov Categories
 - Axioms
 - Examples
- 6 Enriched Categories
- 7 Divergences
 - Entropy

Definition of Markov Category

Definition

A Markov category is a semi-Cartesian category in which every object is a commutative comonoid.

Example Markov category

Okay if we create an example, it absolutely needs to involve the cloud cover forecast during the total solar eclipse.

Enriched Categories

So yeah an enriched category is like a category with more structure on the morphisms. Instead of having hom-sets, ya got yourself some hom-objects.

Divergence

Aight, so a divergence is like a metric for probability distributions...
Except it ain't a metric, cause it ain't symmetric.

Entropy

Entropy is a measure of how “random” a probability distribution is, given through the divergence between both sides of the defining equation for deterministic kernels. Neat!!