

- ①
1. $\frac{\Gamma'}{(Q \vee (R_3 \wedge R_4))}, \gamma_Q \vdash (\frac{\varphi_3 \vee \varphi_4}{Q \vee (R_3 \wedge R_4)}) \text{ (ip)}$
 2. $(Q \vee (R_3 \wedge R_4)), \gamma_Q, Q \vdash Q \text{ (ip)}$
 3. $(Q \vee (R_3 \wedge R_4)), \gamma_Q, Q \vdash \gamma_Q \text{ (ip)}$
 4. $(Q \vee (R_3 \wedge R_4)), \gamma_Q, Q \vdash \perp \quad (\gamma_e, 2, 3)$
- $\dashv_e \frac{\Gamma \vdash \perp}{\Gamma \vdash \Psi}$

$\alpha_1. (Q \vee (R_3 \wedge R_4)), \gamma_Q, Q \vdash (R_3 \wedge R_4) \quad (\perp_e, h)$

$\alpha_2. (Q \vee (R_3 \wedge R_4)), \gamma_Q, (R_3 \wedge R_4) \vdash (R_3 \wedge R_4) \text{ (ip)}$

$\alpha_3. (Q \vee (R_3 \wedge R_4)), \gamma_Q \vdash (R_3 \wedge R_4)^+ \quad (\vee_e, \perp, \alpha_1, \alpha_2)$

$\alpha_4. (Q \vee (R_3 \wedge R_4)) \vdash (\frac{\gamma_Q \rightarrow (R_3 \wedge R_4)}{\varphi_1 \quad \varphi_2}) \quad (\rightarrow_i, \alpha_3)$

$$\frac{\Gamma, \varphi_1 \vdash \varphi_2}{\Gamma \vdash \varphi_1 \rightarrow \varphi_2}$$

$$\vee_e \frac{\frac{\Gamma' \vdash \varphi_3 \vee \varphi_4}{\Gamma' \vdash \varphi}}{\Gamma \vdash \varphi} \quad \frac{\Gamma' \vdash \varphi_3 \vdash \psi}{\Gamma \vdash \varphi} \quad \frac{\Gamma' \vdash \varphi_4 \vdash \psi}{\Gamma \vdash \varphi}$$

$$\forall_e \frac{\Gamma \vdash \forall x. \varphi}{\Gamma \vdash \varphi[x \mapsto t]}$$

$$\forall_i \frac{\Gamma \vdash \varphi[x \mapsto x_0]}{\Gamma \vdash \forall x. \varphi} \quad x_0 \notin \text{vars}(\Gamma, \varphi)$$

$$1. \Gamma \vdash \forall x. \underbrace{P(x, y)}_{\varphi}^+$$

$$2. \Gamma \vdash P(e, y) \quad (\forall_e, 1) \quad i(x)$$

$$1. \Gamma \vdash P(x, y) \quad \text{variable new}$$

$$2. \Gamma \vdash \forall x. \underbrace{P(x, y)}_{\varphi} \quad (\forall_i, 1)$$

3. $\Gamma \vdash \forall x. \exists y. \underbrace{P(x, y)}_{\varphi}$ legt x zu y \rightarrow im P free
in φ

4. $\Gamma \vdash \exists x. P(x, y) \quad (\forall_e, 3)$

$$\exists_i \frac{\Gamma \vdash \varphi[x \mapsto t]}{\Gamma \vdash \exists x. \varphi}$$

$$\exists e \frac{\Gamma \vdash \exists x. \varphi \quad \Gamma, \varphi[x \mapsto x] \vdash \psi}{\Gamma \vdash \psi} \begin{matrix} x_0 \notin \text{vars}(\varphi) \\ \Gamma, \varphi, \psi \end{matrix}$$

$$1. \Gamma \vdash P(i(e), y)$$

$$2. \Gamma \vdash \exists x. P(x, y) (\exists_i, 1)$$

$$3. \Gamma \vdash \exists z. P(i(e), z) (\exists_i, 1)$$

$$1. \Gamma \vdash \exists x. \varphi$$

$$2. \Gamma, P(y) \vdash \varphi (ip) \quad !!$$

$$3. \Gamma \vdash P(y) \quad \begin{matrix} \text{oude variabila,} \\ \text{dar NU } y, x \\ \text{sau van din } \Gamma \end{matrix}$$

$$\vee_e \frac{\Gamma \vdash \varphi_1 \vee \varphi_2 \quad \Gamma, \varphi_1 \vdash \psi \quad \Gamma, \varphi_2 \vdash \psi}{\Gamma \vdash \psi}$$

(2)

$$1. \exists z. R(z), \forall y_3. (R(y_3) \rightarrow Q(y_3)) \vdash y_3 (R(y_3) \rightarrow Q(y_3)) (ip)$$

$$2. \exists z. R(z), \forall y_3. (R(y_3) \rightarrow Q(y_3)) \vdash R(tz) \rightarrow Q(tz) (\forall_e, 1)$$

$$3. \exists z. R(z), \forall y_3. (R(y_3) \rightarrow Q(y_3)) \vdash \exists z. R(z) (ip)$$

$$m. \exists z. R(z), \forall y_3. (R(y_3) \rightarrow Q(y_3)), R(x_0) \vdash R(tz) \quad \begin{matrix} z_0 \notin \text{vars}(\Gamma, \varphi, \psi) \\ z_0 \notin \text{vars}(\Gamma) \checkmark \\ z_0 \notin \text{vars}(\varphi) \checkmark \\ z_0 \notin \text{vars}(\psi) \end{matrix}$$

$$j. \exists z. R(z), \forall y_3. (R(y_3) \rightarrow Q(y_3)) \vdash R(tz) (\exists_e, 3, m) \quad \begin{matrix} tz \text{ nu} \\ \text{contine } z_0 \end{matrix}$$

$$i. \exists z. R(z), \forall y_3. (R(y_3) \rightarrow Q(y_3)) \vdash Q(tz) (\rightarrow_i, 2, j)$$

$$k. \exists z. R(z), \forall y_3. (R(y_3) \rightarrow Q(y_3)) \vdash \exists z. Q(z) (\exists_i, i)$$

$$n. \exists z. R(z), \forall y_3. (R(y_3) \rightarrow Q(y_3)) \vdash \forall z. \exists z. Q(z) (\forall_i, k)$$

1. $\exists z. R(z), \forall y. (R(y) \rightarrow Q(y)) \vdash \exists z. R(z)$ (ip)
2. $\exists z. R(z), \forall y. (R(y) \rightarrow Q(y)), R(x_0) \vdash \forall y. (R(y) \rightarrow Q(y))$ (ip)
3. $\exists z. R(z), \forall y. (R(y) \rightarrow Q(y)), R(x_0) \vdash R(x_0) \rightarrow Q(x_0)$ ($\rightarrow e, 2$)
4. $\exists z. R(z), \forall y. (R(y) \rightarrow Q(y)), R(x_0) \vdash R(x_0)$ (ip)
- j. $\exists z. R(z), \forall y. (R(y) \rightarrow Q(y)), R(x_0) \vdash Q(x_0)$ ($\rightarrow e, 3, h$)
 $x_0 \notin \text{vars}(\Gamma) \checkmark$
 $x_0 \notin \text{vars}(e) \checkmark$
 $x_0 \notin \text{vars}(h) \checkmark$
- m. $\exists z. R(z), \forall y. (R(y) \rightarrow Q(y)), R(x_0) \vdash \exists z. Q(z)$ ($\exists i, j$)
- k. $\exists z. R(z), \forall y. (R(y) \rightarrow Q(y)) \vdash \exists z. Q(z)$ ($\exists e, 1, m$)
- n. $\exists z. R(z), \forall y. (R(y) \rightarrow Q(y)) \vdash \forall z. \exists z. Q(z)$ ($\forall i, k$)

$$③ \varphi = (\exists y. \neg P(a) \vee \neg P(\underline{x}))$$

$$\varphi_a = \neg (\forall y. (\underline{P(b)} \vee \underline{Q(a)}))$$

$$\neg \forall x \varphi_i \equiv \exists x \neg \varphi_i$$

$$\varphi = \varphi_3 \vee \varphi_4 \quad \neg \varphi = \neg \varphi_3 \wedge \neg \varphi_4$$

$\varphi \equiv \varphi_a$ ddacă pt orice S și orice S -attribution

$S, \alpha \models \varphi$ ddacă $S, \alpha \models \varphi_a$

$S, \alpha \models \varphi$ ddacă există $u \in D$ a.i. $S, \alpha[y \mapsto u] \models \neg (P(a) \vee \neg P(\underline{x}))$

ddacă există $u \in D$ a.i. Nu avem $\begin{cases} S, \alpha[y \mapsto u] \models \underline{P(a)} \\ \text{ sau } \\ S, \alpha[y \mapsto u] \models \neg P(\underline{x}) \end{cases}$

ddacă există $u \in D$ a.i. Nu avem

$$\begin{cases} P^S(\overline{\alpha[y \mapsto u]}(a)) \end{cases}$$

sau $P^S(\overline{\alpha[y \mapsto u]}(\underline{x}))$ nu are loc

ddacă există un α în Hu are loc $P^S(\alpha^S)$

$S, \alpha \models \varphi$

sau

$P^S(\alpha(\alpha))$ nu are loc

\exists

ddacă Hu nu are loc $P^S(b^S)$

$S, \alpha \models \varphi$

sau

$Q^S(a^S)$ nu are loc

Prin

\exists

$$\varphi_c = \neg (\exists x, P(x)) \wedge \neg (\exists y, \neg P(y))$$

Hu există un α în $P^S(\alpha^S)$ nu are loc $P^S(\alpha(\alpha))$

4) Trec la logica doar dacă invat la logica sau
nu este adevarat ca Terra este plată.

p : Trec la logica

q : Invat la logica

r : este adevarat ca Terra este plată

• $((P \rightarrow q) \vee \neg r)$

• $(P \rightarrow (q \vee \neg r))$

p dacă q $\neg q \rightarrow p$

p doar dacă q $p \rightarrow q$

p dacă nu numai dacă q

$p \leftrightarrow q$

$$\textcircled{5} \quad \left(\underbrace{(q_3 \wedge q_4)}_{\varphi} \rightarrow p \right) \wedge (r \rightarrow p) \vdash \left(\underbrace{(q_3 \wedge q_4)}_{\varphi} \vee r \right) \rightarrow p$$

1. $(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r) \vdash (q \vee r)$ (ip)
2. $(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r), \varphi \vdash (q \rightarrow p) \wedge (r \rightarrow p)$ (ip)
3. $(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r), \varphi \vdash (q \rightarrow p)$ ($\wedge e_1, 2$)
4. $(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r), \varphi \vdash \varphi$ (ip)
- k_1 . $(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r), \underline{\varphi} \vdash p$ ($\rightarrow e, 3, h$)
- x_1 . $(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r), r \vdash (q \rightarrow p) \wedge (r \rightarrow p)$ (ip)
- x_2 . $(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r), r \vdash (r \rightarrow p)$ ($\wedge e_2, x_1$)
- x_3 . $(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r), r \vdash r$ (ip)
- k_2 . $(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r), \underline{r} \vdash p$ ($\rightarrow e, x_2, x_3$)

- k . $(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r) \vdash p$ ($\vee e, 1, k_1, k_2$)
- n . $(q \rightarrow p) \wedge (r \rightarrow p) \vdash \frac{(q \vee r) \rightarrow p}{\varphi} \quad (\rightarrow i, k)$

$$\rightarrow_i \frac{\Gamma, \varphi_1 \vdash \varphi_2}{\Gamma \vdash \varphi_1 \rightarrow \varphi_2} \quad \rightarrow_e \frac{\Gamma \vdash \varphi_3 \rightarrow \varphi_4 \quad \Gamma \vdash \varphi_3}{\Gamma \vdash \varphi_4}$$

$$(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r) \vdash \underline{r \rightarrow p}$$

- x_1 . $(q \rightarrow p) \wedge (r \rightarrow p), (q \vee r) \vdash (q \vee r)$ (ip)

$$x_2 \quad \underline{(q \rightarrow p) \wedge (r \rightarrow p)}, (q \vee r), \underline{r} \vdash \underline{\quad} \quad (????)$$

$$x_3 \quad \underline{(q \rightarrow p) \wedge (r \rightarrow p)}, (q \vee r), \underline{r} \vdash r \quad (\text{ip})$$

$$n. \quad \underline{(q \rightarrow p) \wedge (r \rightarrow p)}, (q \vee r) \vdash r \quad (\vee e, x_1, x_2, x_3)$$

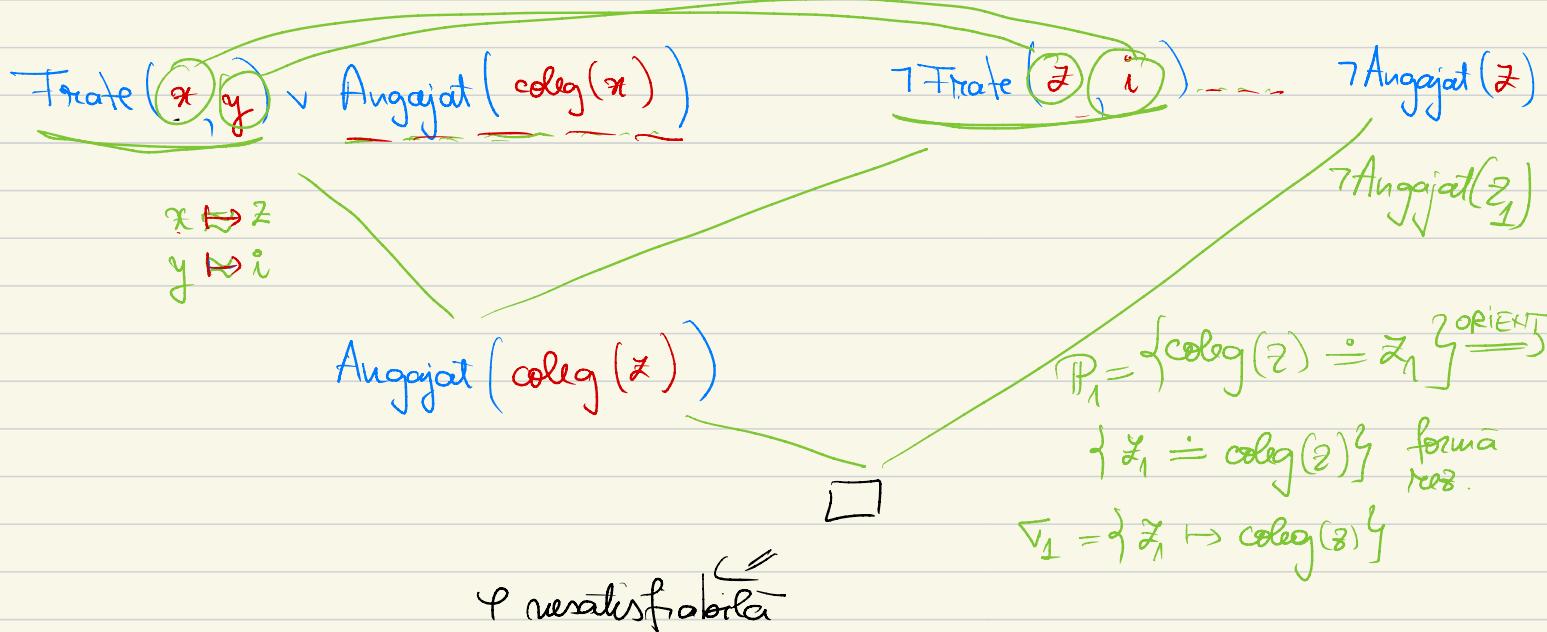
$$k. \quad \underline{(q \rightarrow p) \wedge (r \rightarrow p)}, (q \vee r) \vdash p \quad (\vee e, 1, k_1, k_2)$$

Res. $\varphi = \text{FHSC}$

$$\text{RB} \frac{\mathcal{P}(t_1 \dots t_n) \vee C_1 \quad \neg \mathcal{P}(t'_1 \dots t'_n) \vee C_2}{\nabla^b(C_1 \vee C_2)} \quad V_1 \cap V_2 = \emptyset$$

$$\tau = \text{mgu}(\{t_1 \doteq t'_1, \dots, t_n \doteq t'_n\})$$

$$\varphi = \forall x. \forall y. \forall z. ((\text{Frate}(x, y) \vee \text{Angajat}(\text{coleg}(x)) \wedge (\neg \text{Frate}(z, i)) \wedge \neg \text{Angajat}(z))$$



$$\mathcal{P}(e) \quad \neg \mathcal{P}(i(x))$$

$$P_2 = \{ e \doteq i(x) \} \xrightarrow{\text{conflict}} \perp \Rightarrow \text{HV aplicam RB}$$

$$\text{FP} \frac{\mathcal{P}(t_1 \dots t_n) \vee \mathcal{P}(t'_1 \dots t'_n) \vee C}{\nabla^b(\mathcal{P}(t_1 \dots t_n) \vee C)} \quad \tau = \text{mgu}(\{t_1 \doteq t'_1, \dots, t_n \doteq t'_n\})$$

$$\forall x. (\varphi_1 \wedge \varphi_2) \equiv (\forall x. \varphi_1) \wedge (\forall x. \varphi_2) \equiv (\forall x. \varphi_1)_1 (\forall x_1. \varphi_2[x \mapsto x_1])$$

LR

$$\equiv \forall x. \forall x_1. (\varphi_1)_1 \varphi_2[x \mapsto x_1]$$