

# UNIVERSIDADE FEDERAL DO CEARÁ

CAMPUS DE RUSSAS

## Algoritmos em Grafos

Aula 17: Fluxo Máximo em Redes(Push-Relabel)

**Professor Pablo Soares** 

2022.1

#### Sumário

- 1. Fluxo Máximo em Redes(última aula);
  - a. Rede Residual;
  - b. Caminho aumentante.
- 2. Intuição;
- 3. Pré-Fluxo;
- 4. Operações.
  - a. Push;
  - b. Relabel;
  - c. Algoritmo Genérico Push-Relabel.
    - i. Exemplo.

### Intuição















### Intuição













#### Push-Relabel(Pré-fluxo)

• Propriedade de conservação de fluxo não é mantida;

- Um Pré-<u>fluxo</u> é uma uma função f:  $Vx V \rightarrow R$ ;
  - a.  $f(V, u) \ge 0$  para todo  $u \in V \{s\}$ 
    - e(u) = f(V, u) (Excesso)
  - b. se  $u \in V \{s, t\} \& e(u) > 0$ 
    - *u* está <u>transbordando</u>
- Também usa-se o conceito de **rede residual**.

- Relabel(u)
  - a. Seja G = (V, E) uma rede com origem s e destino t, e seja f um pré-fluxo em G. Uma função  $h: V \rightarrow N$  é uma função de altura se

■ 
$$h(s) = |V|, h(t) = 0 e h(u) \le h(v) + 1 \forall (u, v) \in E_f$$

- b. Aplica-se quando:
  - e(u) > 0
  - $\forall v \in V, tal que(u, v) \in E_f, h(u) \le h(v)$

Relabel(u)  
1. 
$$h[u] \leftarrow 1 + min\{ h[v]: (u, v) \in E_f \}$$
  
Fim.

$$e(u) > 0$$

3/3

3/3

4/4

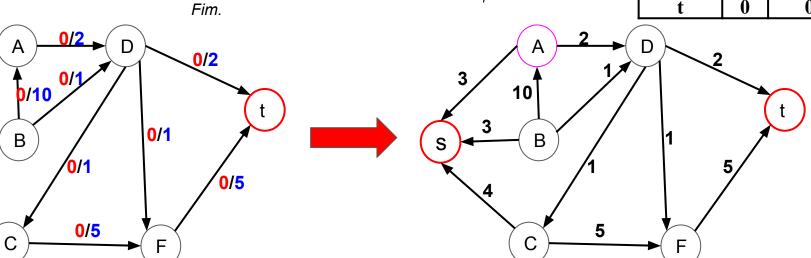
G = (V, E)

$$\forall v \in V, tal que(u, v) \in E_f, h(u) \le h(v)$$

$$\forall v \in V, \ tal \ que \ (u, v) \in E_f, \ h(u) \leq h(v)$$
Relabel(u)

Relabel(u)

1. 
$$h[u] \leftarrow 1 + min\{h[v]: (u, v) \in E_f\}$$



h

0

0

e

 $\infty$ 3

Vértice

B

D

$$\frac{\text{xemplo Relabel}(u)}{\text{xemplo Relabel}(u)}$$

a. Aprica-se quando. 
$$\bullet e(u) > 0$$

0/5

G = (V, E)

C

$$\blacksquare$$
  $\forall v \in$ 

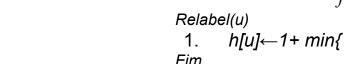
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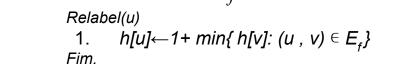
$$\forall v \in V, tal que(u, v) \in E_f, h(u) \leq h(v)$$

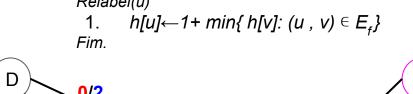
$$\forall V \subseteq V, \ lat \ Que \ (u, V) \subseteq E_f, \ h(u) \le h(V)$$

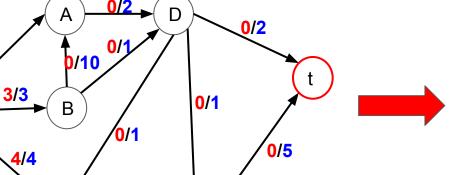
$$Relabel(u)$$

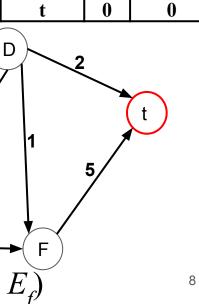
$$1 \qquad h[u] = 1 + \min\{h[v]: (u, v) \in F\}$$











Vértice

B

D

10

В

h

0

e

 $\infty$ 3

Vértice

B

D

h

0

e

 $\infty$ 3

$$\bullet \quad e(u) > 0$$

$$\forall v \in$$

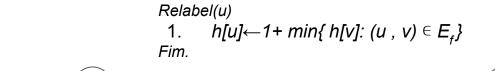
$$\forall v \in V, \ tal \ que \ (u, v) \in E_f, \ h(u) \le h(v)$$

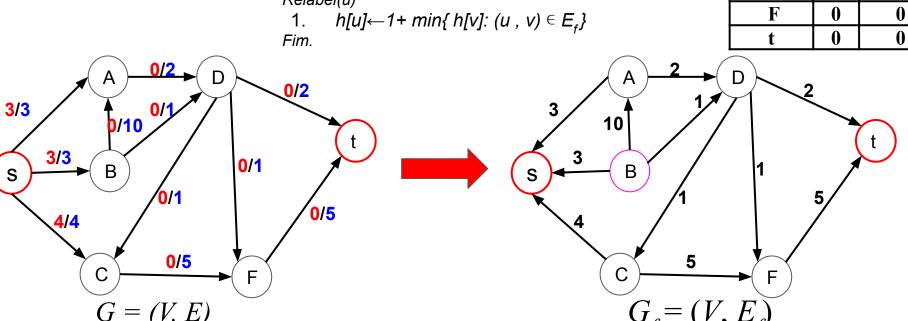
$$\forall v \in V, \ tat \ que \ (u, v) \in E_f, \ n(u) \leq n(v)$$

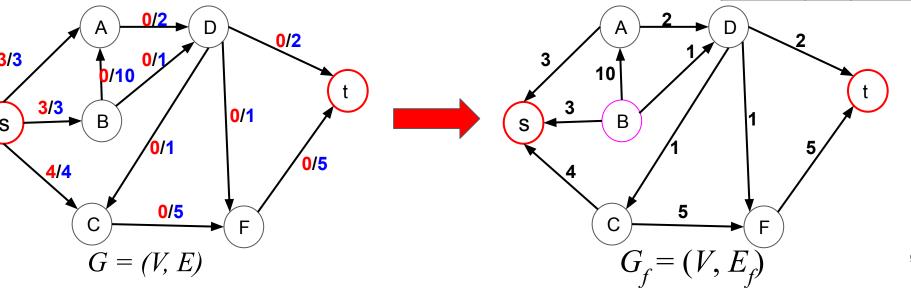
Relabel(u)

Relabel(u)

1. 
$$h[u] \leftarrow 1 + min\{h[v]: (u, v) \in E_f\}$$





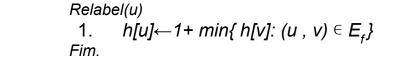


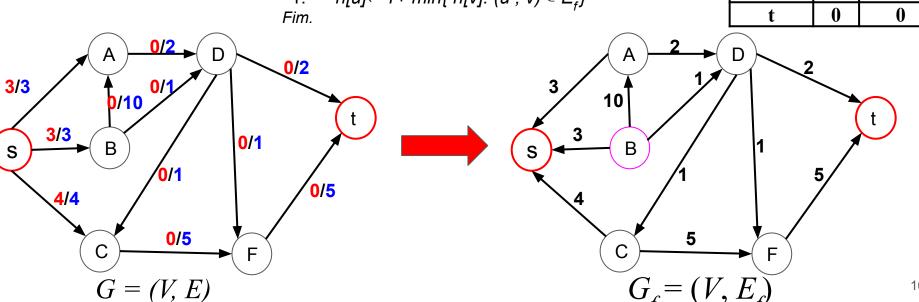
$$\bullet e(u) > 0$$

$$\forall v \in V, tal que(u, v) \in E_f, h(u) \le$$

$$\forall v \in V, \ tal \ que \ (u, v) \in E_f, \ h(u) \le h(v)$$
Relabel(u)

Relabel(u)
$$h[u] \leftarrow 1 + \min\{h[v]: (u, v) \in E_t\}$$





h

0

e

 $\infty$ 3

B

D

- Push(u, v)
  - a. Seja G = (V, E) uma rede com origem s e destino t, e seja f um pré-fluxo em G.
  - b. Aplica-se quando:
    - e(u) > 0

    - $c_f(u, v) > 0$  h(u) = h(v) + 1

- 1.  $d_{f}[u, v] = min(e[u], c_{f}[u, v])$
- 2.  $f[u, v] \leftarrow f[u, v] + d_f[u, v]$
- 3. *f[v, u]←- f[u, v]*
- 4. e[u]←e[u] d<sub>4</sub>[u, v]
- 5.  $e[v] \leftarrow e[v] + d[u, v]$

Fim.

Exemplo Push(u, v)Vértice Push(u, v) h e

a. Aplica-se quando:
$$e(u) > 0$$

0/2

0/5

0/1

b/10 0/1

0/5

G = (V, E)

В

C

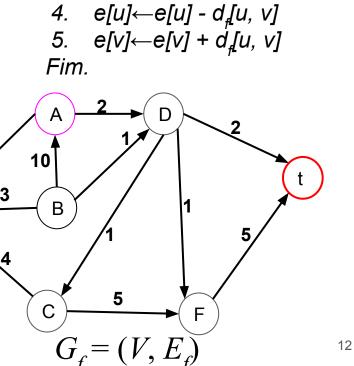
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3/3

4/4

Aprica-se quando.

$$e(u) > 0$$
 $c_f(u, v) > 0$ 
 $h(u) = h(v) + 1$ 
 $c_f(u, v) > 0$ 
 $c_f(u, v) > 0$ 
 $c_f(u, v) > 0$ 
 $c_f(u, v) = 0$ 
 $c_f(u$ 



1.  $d_f[u, v] = min(e[u], c_f[u, v])$ 

2.  $f[u, v] \leftarrow f[u, v] + d_f[u, v]$ 

3. *f*[*v*, *u*]←- *f*[*u*, *v*]

Exemplo Push(u, v)Vértice Push(u, v) h

a. Aplica-se quando:
$$e(u) > 0$$

$$c(u, v) > 0$$
Vértice  $e$ 

$$s$$

$$s$$

$$s$$

$$A$$

$$A$$

$$B$$

$$B$$

$$C$$

0/2

0/5

0/1

b/10 0/1

0/5

G = (V, E)

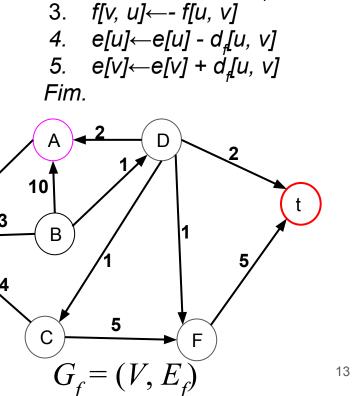
В

C

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3/3

4/4



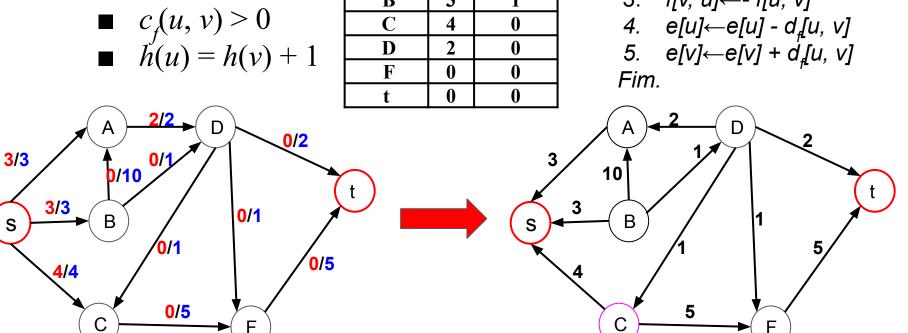
1.  $d_f[u, v] = min(e[u], c_f[u, v])$ 

2.  $f[u, v] \leftarrow f[u, v] + d_f[u, v]$ 

Exemplo Push(u, v)Vértice Push(u, v) h e

a. Aplica-se quando:
$$e(u) > 0$$

G = (V, E)



1.  $d_f[u, v] = min(e[u], c_f[u, v])$ 

2.  $f[u, v] \leftarrow f[u, v] + d_f[u, v]$ 

3. *f*[*v*, *u*]←- *f*[*u*, *v*]

Exemplo Push(u, v)Vértice Push(u, v) h e

0/2

0/5

0/1

b/10 0/1

0/5

G = (V, E)

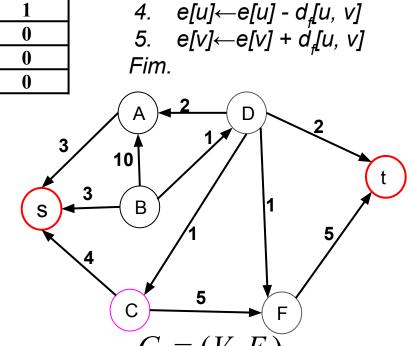
В

C

3/3

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4/4



1.  $d_f[u, v] = min(e[u], c_f[u, v])$ 

15

2.  $f[u, v] \leftarrow f[u, v] + d_f[u, v]$ 

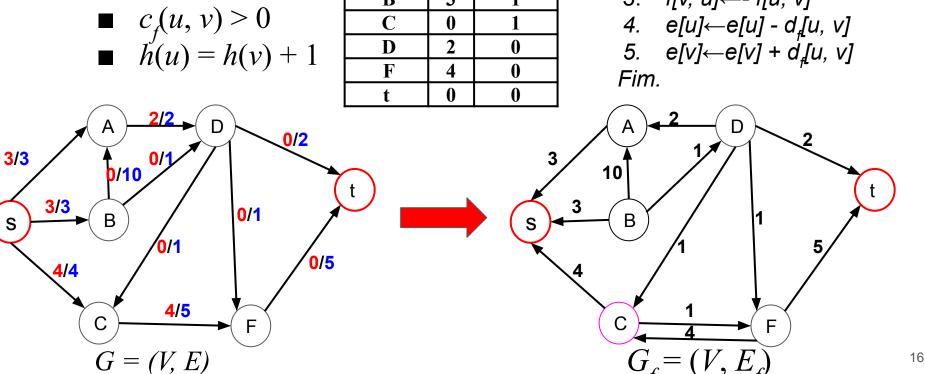
3. *f*[*v*, *u*]←- *f*[*u*, *v*]

1.  $d_f[u, v] = min(e[u], c_f[u, v])$ 

2.  $f[u, v] \leftarrow f[u, v] + d_f[u, v]$ 

3. *f*[*v*, *u*]←- *f*[*u*, *v*]

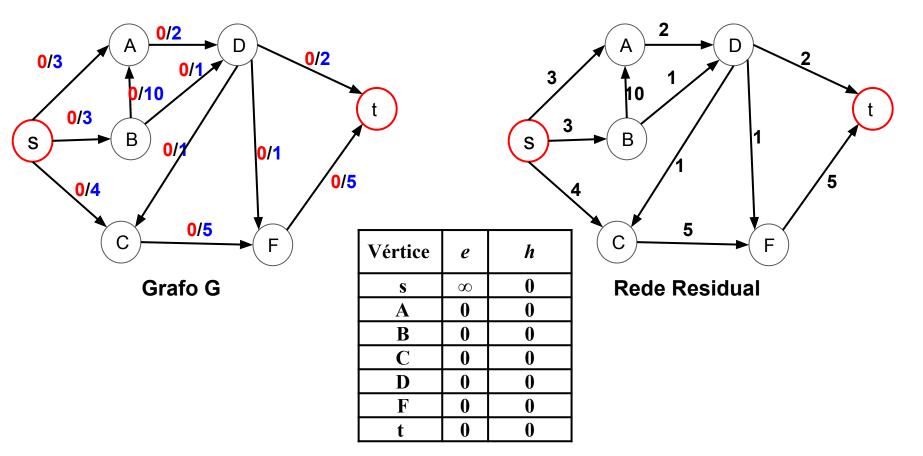
Exemplo Push(u, v)Push(u, v)

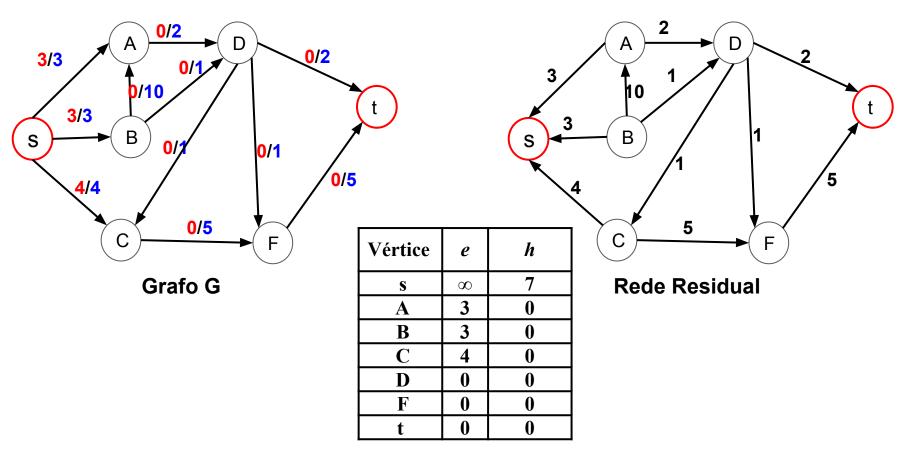


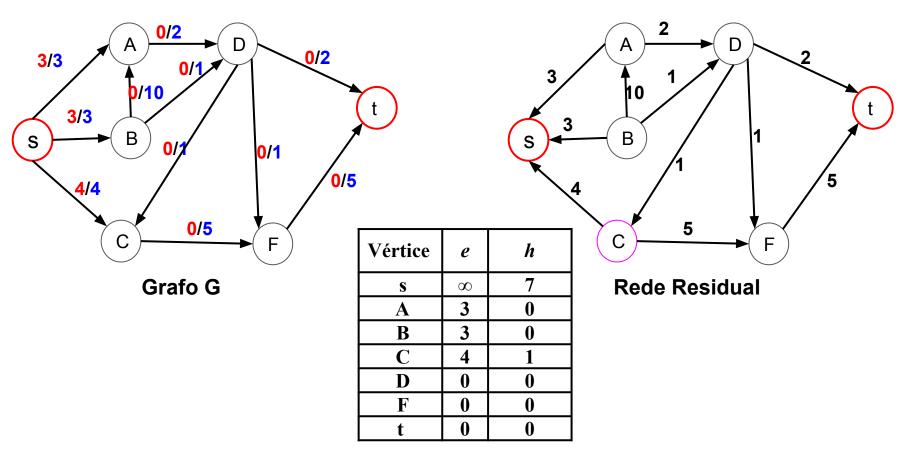
Pseudocódigo Push-Relabel 1. para cada vértice u ← V[G]  $h[u] \leftarrow e[u] \leftarrow 0$ Push(u, v) 3. fimpara 1.  $d_{f}[u, v] = min(e[u], c_{f}[u, v])$ 4. para cada aresta  $(u,v) \leftarrow E[G]$ 2.  $f[u, v] \leftarrow f[u, v] + d_f[u, v]$  $f[u, v] \leftarrow f[v, u] \leftarrow 0$ 3. *f[v, u]←- f[u, v]* 6. fimpara 4. e[u]←e[u] - d<sub>4</sub>[u, v] 7.  $h[s] \leftarrow |V|$ 5.  $e[v] \leftarrow e[v] + d[u, v]$ para cada u ∈ L.adi[s] Fim. f[s, u] $\leftarrow$ c(s, u) Relabel(u) 10.  $f[u, s] \leftarrow -c(s, u)$ 1.  $h[u] \leftarrow 1 + min\{ h[v]: (u, v) \in E_f \}$ 11.  $e[u] \leftarrow c(s, u)$ Fim. 12.  $e[s] \leftarrow e[s] - c(s, u)$ 13. fimpara 14. enquanto existir uma operação de <u>Push</u> ou <u>Relabel</u> faça *15.* selecione uma operação e execute 16. fimenquanto Fim.

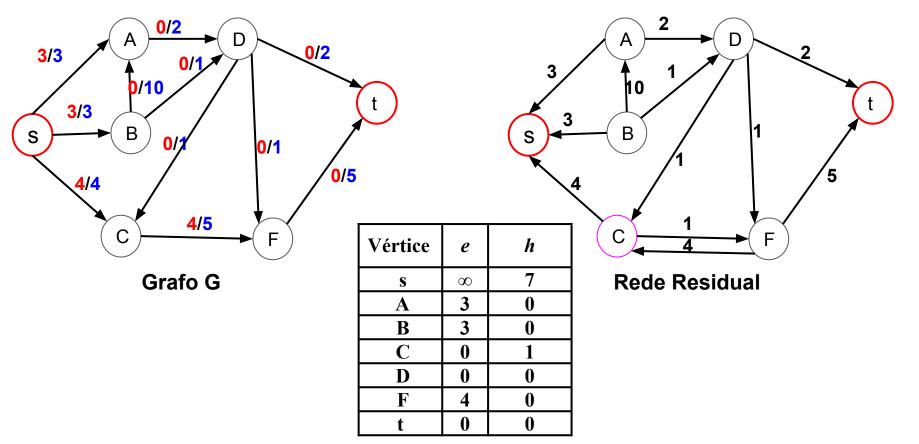
Push-Relabel-Generico(G, s)

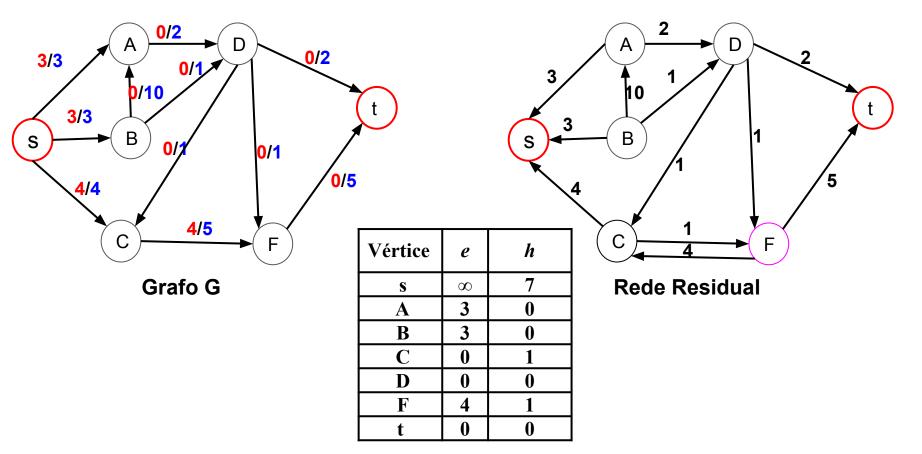
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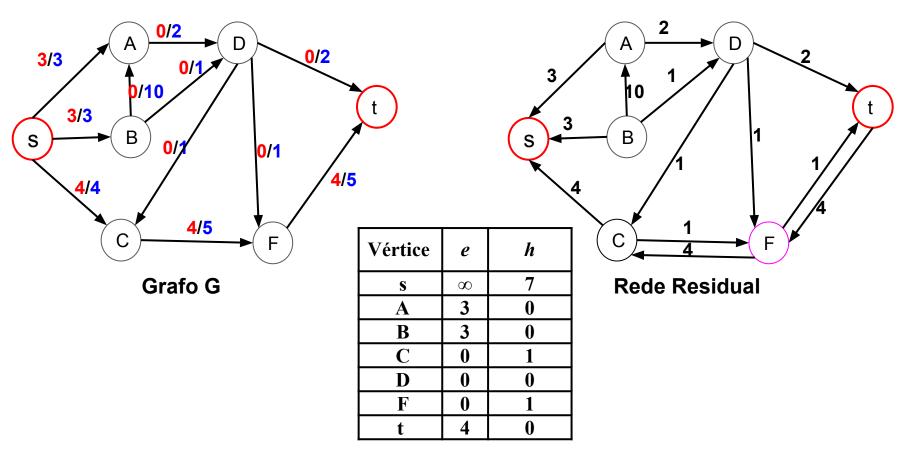


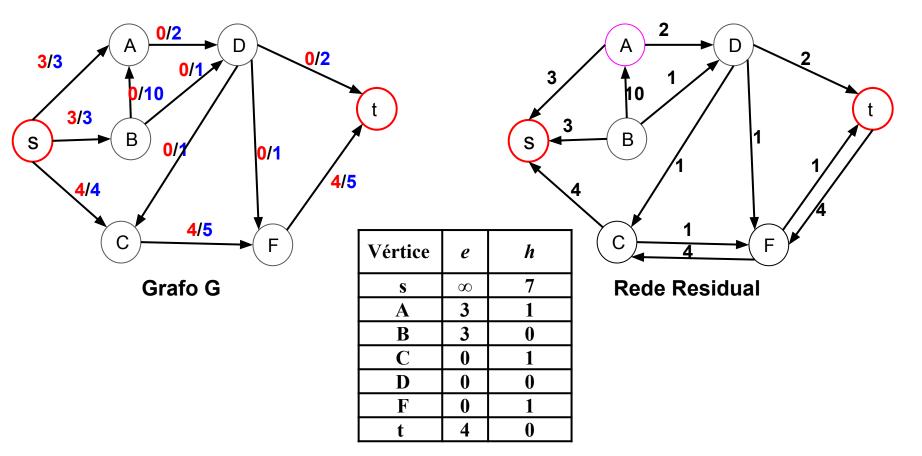


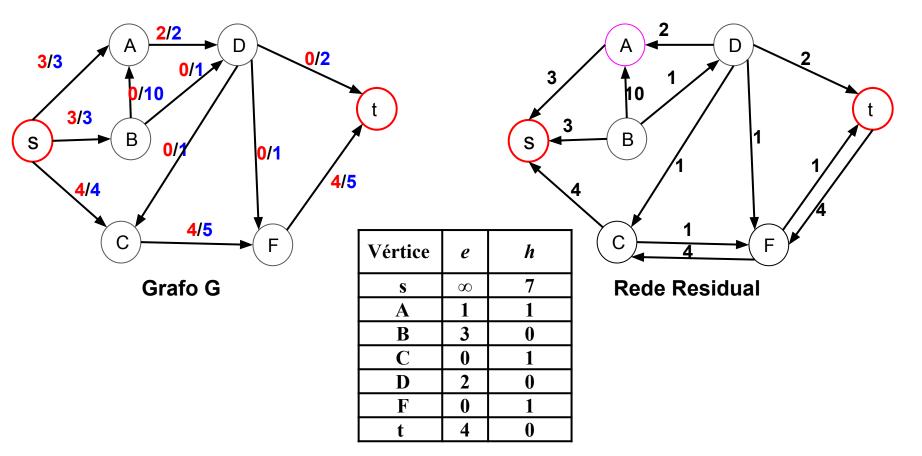


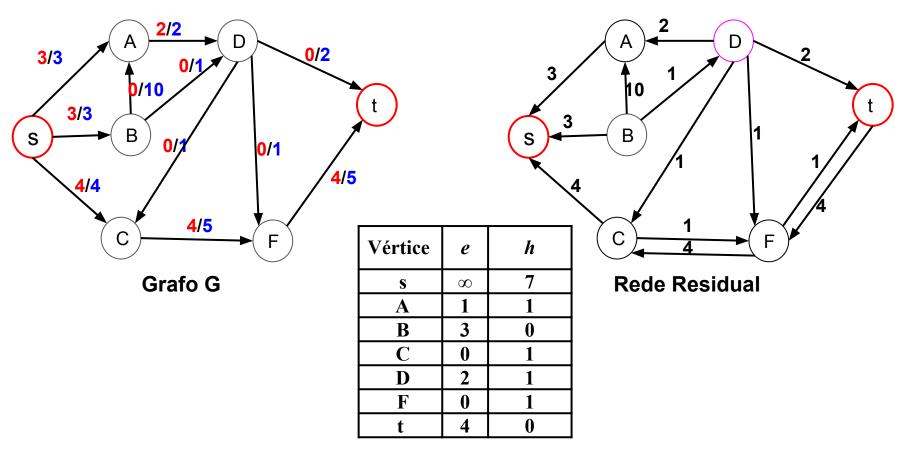


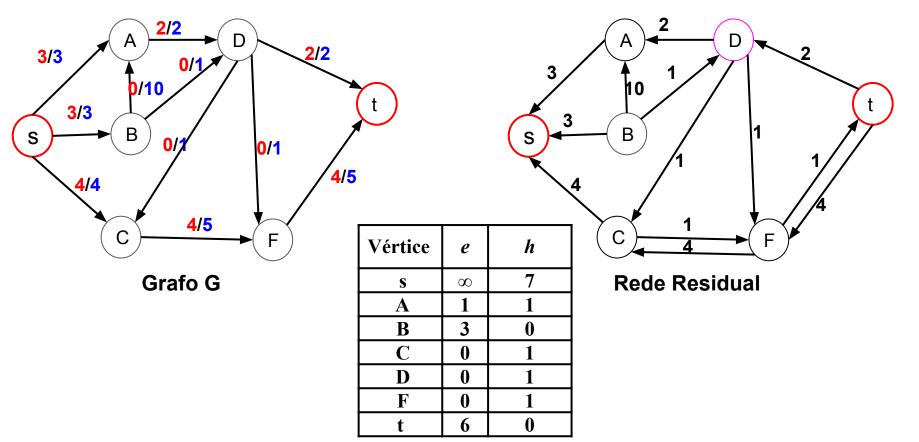


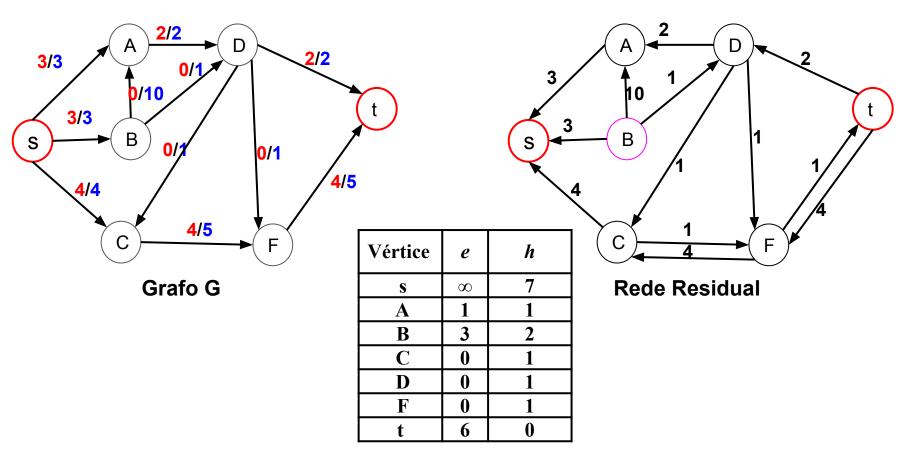


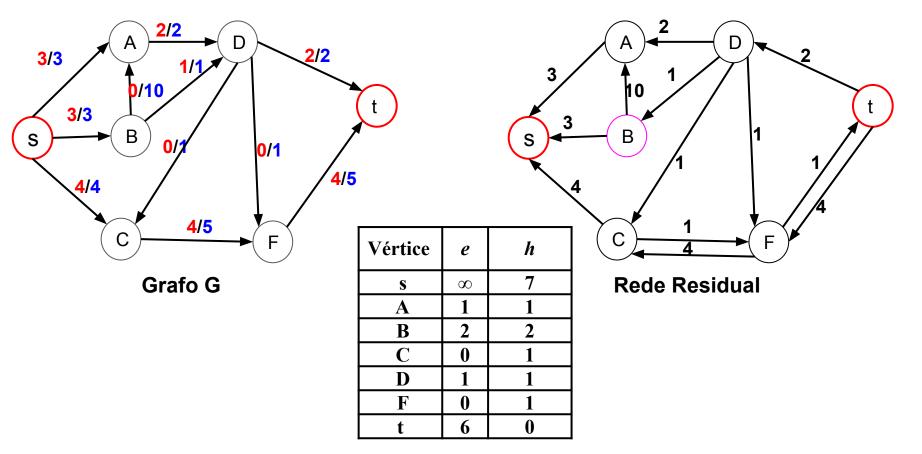


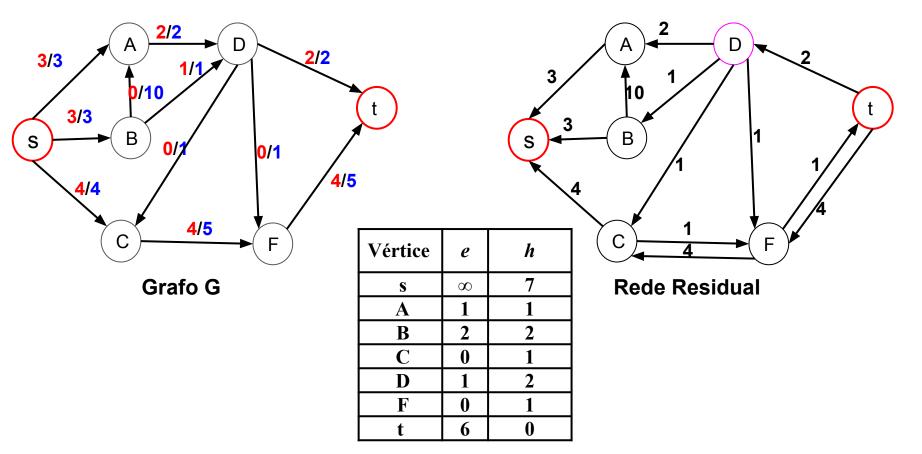


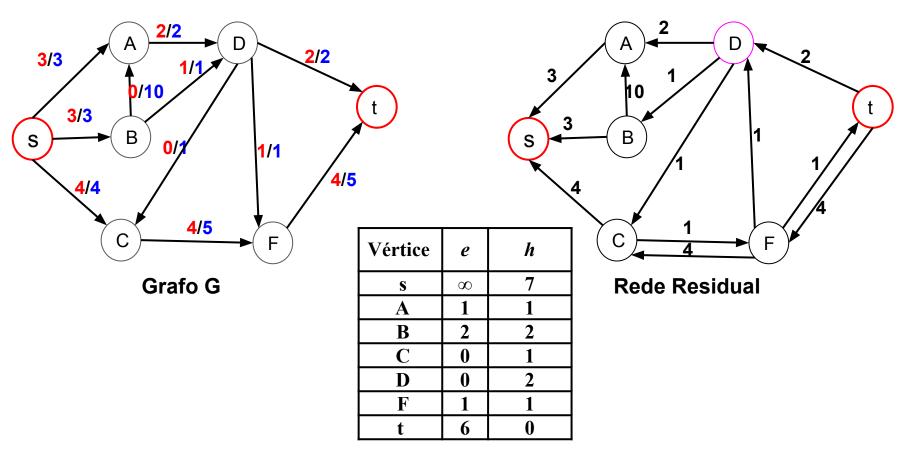


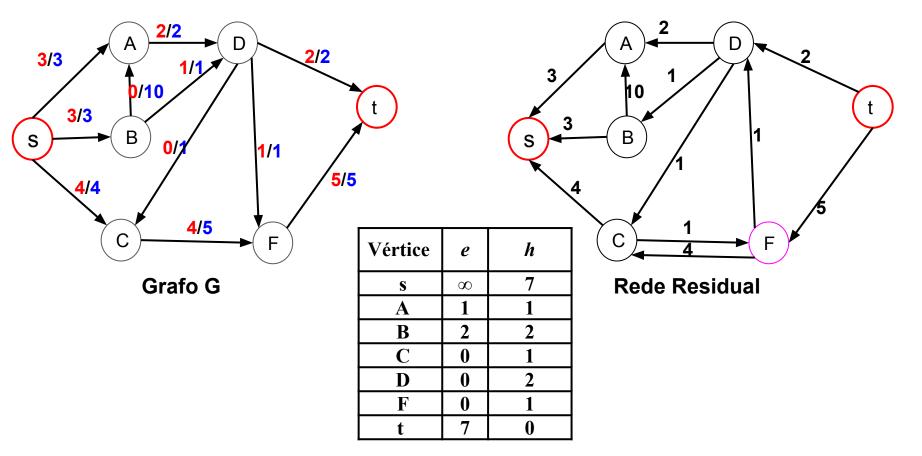






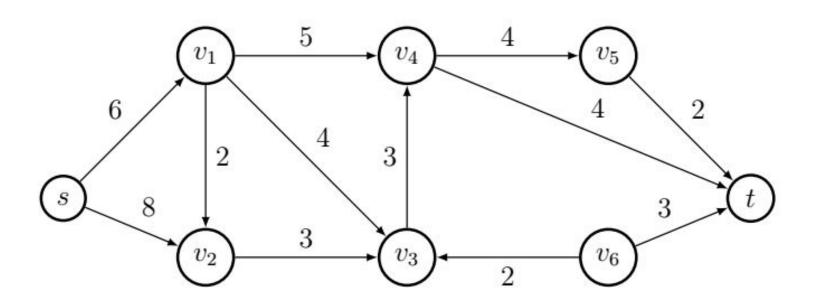






#### Exercício de Fixação

Determine o fluxo máximo na rede abaixo.





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