内存屏障的定义

```
文件 linux-5.6.3/include/linux/compiler-gcc.h。
编译期屏障,使用 memory 标识符。
/* Optimization barrier */
/* The "volatile" is due to gcc bugs */
#define barrier() __asm__ _volatile_("": : :"memory")
文件 linux-5.6.3/include/asm-generic/barrier.h。
#ifdef CONFIG SMP
#ifndef smp_mb
#define smp mb()
                   __smp_mb()
#endif
#ifndef smp_rmb
#define smp_rmb()
                  __smp_rmb()
#endif
#ifndef smp wmb
#define smp wmb()
                   smp wmb()
#endif
文件 linux-5.6.3/arch/x86/include/asm/barrier.h。
编译期屏障、运行期屏障。
#ifdef CONFIG X86 32
#define mb() asm volatile(ALTERNATIVE("lock; addl $0,-4(%%esp)", "mfence", \
                     X86 FEATURE XMM2) ::: "memory", "cc")
#define rmb() asm volatile(ALTERNATIVE("lock; addl $0,-4(%%esp)", "lfence", \
                      X86 FEATURE XMM2) ::: "memory", "cc")
#define wmb() asm volatile(ALTERNATIVE("lock; addl $0,-4(%%esp)", "sfence", \
                      X86 FEATURE XMM2) ::: "memory", "cc")
#else
#define mb()
               asm volatile("mfence":::"memory")
               asm volatile("lfence":::"memory")
#define rmb()
#define wmb()
               asm volatile("sfence" ::: "memory")
#endif
#define dma rmb()
                   barrier()
                   barrier()
#define dma_wmb()
#ifdef CONFIG_X86_32
```

```
#define __smp_mb() asm volatile("lock; addl $0,-4(%esp)" ::: "memory", "cc")
#else
#define __smp_mb() asm volatile("lock; addl $0,-4(%rsp)" ::: "memory", "cc")
#endif
#define __smp_rmb() dma_rmb()
#define __smp_wmb() barrier()
```

内存屏障的使用举例

```
文件 linux-5.6.3/kernel/events/ring_buffer.c。
static void perf_output_put_handle(struct perf_output_handle *handle)
    struct perf_buffer *rb = handle->rb;
   unsigned long head;
   unsigned int nest;
    * If this isn't the outermost nesting, we don't have to update
    * @rb->user page->data head.
    */
    nest = READ_ONCE(rb->nest);
    if (nest > 1) {
        WRITE ONCE(rb->nest, nest - 1);
        goto out;
   }
again:
   /*
    * In order to avoid publishing a head value that goes backwards,
    * we must ensure the load of @rb->head happens after we've
    * incremented @rb->nest.
    * Otherwise we can observe a @rb->head value before one published
    * by an IRQ/NMI happening between the load and the increment.
    */
    barrier();
    head = local_read(&rb->head);
```

```
文件 linux-5.6.3/kernel/exit.c。
使用 smp_mb()保证 2 个函数的变量的相互依赖的顺序。
void rcuwait_wake_up(struct rcuwait *w)
{
   struct task_struct *task;
   rcu_read_lock();
```

```
/*
 * Order condition vs @task, such that everything prior to the load
 * of @task is visible. This is the condition as to why the user called
 * rcuwait_trywake() in the first place. Pairs with set_current_state()
 * barrier (A) in rcuwait wait event().
      WAIT
                          WAKE
      [S] tsk = current [S] cond = true
 *
 *
          MB (A)
                          MB (B)
      [L] cond
                    [L] tsk
 *
 */
smp mb(); /* (B) */
task = rcu_dereference(w->task);
if (task)
    wake_up_process(task);
rcu read unlock();
              struct vcpu_runstate_info *res, unsigned int cpu)
```