



Basics of Combinatorics

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Introduction

Counting the objects that satisfy some criteria is a very common task in both TopCoder problems and in real-life situations. The myriad ways of counting the number of elements in a set is one of the main tasks in combinatorics, and I'll try to describe some basic aspects of it in this tutorial. These methods are used in a range of applications, from discrete math and probability theory to statistics, physics, biology, and more.

Combinatorial primitives

Let's begin with a quick overview of the basic rules and objects that we will reference later.



The rule of sum

$$A \cap B = \emptyset \Rightarrow |A| + |B| = |A \cup B|$$



The rule of product

$$|A| \cdot |B| = |A \times B|$$

For example, if we have three towns — A, B and C — and there are 3 roads from A to B and 5 roads from B to C, then we can get from A to C through B in $3 \cdot 5 = 15$ different ways.

These rules can be used for a finite collections of sets.

Permutation without repetition

When we choose k objects from n -element set in such a way that the order matters and each object can be chosen only once:

$$(n)_k = P_k^n = \frac{n!}{(n-k)!}$$

For example, suppose we are planning the next 3 challenges and we have a set of 10 easy problems to choose from. We will only use one easy problem in each contest, so we can choose our problems in $(10)_3 = \frac{10!}{(10-3)!}$ different ways.

Permutation (variation) with repetition

The number of possible choices of k objects from a set of n objects when order is important and one object can be chosen more than once:

$$n^k$$

For example, if we have 10 different prizes that need to be divided among 5 people, we can do so in 5^{10} ways.

Permutation with repetition

The number of different permutations of n objects, where there are n_1 indistinguishable objects of type 1, n_2 indistinguishable objects of type 2, ..., and n_k indistinguishable objects of type k ($n_1 + n_2 + \dots + n_k = n$), is:

$$\binom{n}{n_1, n_2, \dots, n_k} = C_n(n_1, n_2, \dots, n_k) = \frac{n!}{n_1! \cdot n_2! \cdot \dots \cdot n_k!}$$

For example, if we have 97 coders and want to assign them to 5 rooms (rooms 1-4 have 20 coders in each, while the 5th room has 17), then there

are $\binom{97}{20, 20, 20, 20, 17} = \frac{97!}{20! \cdot 20! \cdot 20! \cdot 20! \cdot 17!}$ possible ways to do it.

Combinations without repetition

In combinations we choose a set of elements (rather than an arrangement, as in permutations) so the order doesn't matter. The number of different k-element subsets (when each element can be chosen only once) of n-element set is:

$$\binom{n}{k} = C_k^n = \frac{n!}{k!(n-k)!}$$

For example, if we have 7 different colored balls, we can choose any 3 of them in $\binom{7}{3} = \frac{7!}{3!(7-3)!}$ different ways.

Combination with repetition

Let's say we choose k elements from an n-element set, the order doesn't matter and each element can be chosen more than once. In that case, the number of different combinations is:

$$f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}$$

For example, let's say we have 11 identical balls and 3 different pockets, and we need to calculate the number of different divisions of these balls to the pockets. There would be $f_{11}^3 = \binom{3+11-1}{11} = \frac{(3+11-1)!}{11!(3-1)!}$ different combinations.

It is useful to know that f_k^n is also the number of integer solutions to this equation:

$$x_1 + x_2 + \dots + x_n = k, x_i \geq 0 (i \in \overline{1, n})$$

Why? It's easy to prove. Consider a vector $(1, 1, \dots, 1)$ consisting of $(n+k-1)$ ones, in which we want to substitute $n-1$ ones for zeroes in such way that we'll get n groups of ones (some of which may be empty) and the number of ones in the i^{th} group will be the value of x_i :

$$\underbrace{(1, \dots, 1)}_{x_1}, 0, \underbrace{1, \dots, 1}_{x_2}, 0, \dots, 0, \underbrace{1, \dots, 1}_{x_n}$$

The sum of x_i will be k, because k ones are left after substitution.

The Basics

Binary vectors

Some problems, and challenge problems are no exception, can be reformulated in terms of binary vectors. Accordingly, some knowledge of the basic combinatorial properties of binary vectors is rather important. Let's have a look at some simple things associated with them:

1. Number of binary vectors of length n: 2^n .

2. Number of binary vectors of length n and with k '1' is

$$\binom{n}{k}.$$

We just choose k positions for our '1's.

3. The number of ordered pairs (a, b) of binary vectors, such that the distance between them (k) can be calculated as follows: $\binom{n}{k} \cdot 2^n$.

The distance between a and b is the number of components that differs in a and b — for example, the distance between (0, 0, 1, 0) and (1, 0, 1, 1) is 2).

Let $a = (a_1, a_2, \dots, a_n)$, $b = (b_1, b_2, \dots, b_n)$ and distance between them is k. Next, let's look at the sequence of pairs $(a_1, b_1), (a_2, b_2), \dots, (a_n, b_n)$. There are exactly k indices i in which $a_i \neq b_i$. They can be (0,1) or (1,0), so there are 2 variants, and n-k can be either (0,0) or (1,1), for another 2 variants. To

calculate the answer we can choose k indices in which vectors differs in $\binom{n}{k}$ ways, then we choose components that differs in 2^k ways and components that are equal in 2^{n-k} ways (all of which use the permutation with repetition formula), and in the end we just multiply all these numbers and get $\binom{n}{k} \cdot 2^k \cdot 2^{n-k} = \binom{n}{k} \cdot 2^n$.

Delving deeper

Now let's take a look at a very interesting and useful formula called the inclusion-exclusion principle (also known as the sieve principle):

$$\left| \bigcup_{i=1}^n A_i \right| = \sum_{i=1}^n (-1)^{i-1} \sum_{\substack{I \subseteq \{1, \dots, n\} \\ |I|=i}} \left| \bigcap_{j \in I} A_j \right|$$

This formula is a generalization of:

$$|A \cup B| = |A| + |B| - |A \cap B|$$

There are many different problems that can be solved using the sieve principle, so let's focus our attention on one of them. This problem is best known as "Derangements". A derangement of the finite set X is a [bijection](#) from X into X that doesn't have fixed points. A small example: For set $X = \{1, 2, 3\}$ bijection $\{(1,1), (2,3), (3,2)\}$ is not derangement, because of $(1,1)$, but bijection $\{(1,2), (2,3), (3,1)\}$ is derangement. So let's turn back to the problem, the goal of which is to find the number of derangements of n -element set.

We have $X = \{1, 2, 3, \dots, n\}$. Let:

- A be the set of all bijections from X into X , $|A| = n!$,
- A_0 be the set of all derangements of X ,
- A_i ($i \in X$) be the set of bijections from X into X that have (i,i) ,
- A_I ($I \subseteq X$) be the set of bijections from X into X that have $(i,i) \forall i \in I$, so $A_I = \bigcap_{i \in I} A_i$ and $|A_I| = (n - |A_I|)!$.

In formula we have sums $\sum_{\substack{I \subseteq \{1, \dots, n\} \\ |I|=i}} \left| \bigcap_{j \in I} A_j \right|$, in our case we'll have $\sum_{\substack{I \subseteq \{1, \dots, n\} \\ |I|=i}} \left| \bigcap_{j \in I} A_j \right| = \sum_{\substack{I \subseteq X \\ |I|=i}} |A_I|$, so let's calculate them:

$$\sum_{\substack{I \subseteq X \\ |I|=i}} |A_I| = \binom{n}{i} (n-i)! = \frac{n!}{i!}$$

(because there are exactly $\binom{n}{i}$ i -element subsets of X).

Now we just put that result into the sieve principle's formula:

$$\left| \bigcup_{i=1}^n A_i \right| = \sum_{i=1}^n (-1)^{i-1} \sum_{\substack{I \subseteq X \\ |I|=i}} |A_I| = \sum_{i=1}^n (-1)^{i-1} \frac{n!}{i!} = n! \sum_{i=1}^n \frac{(-1)^{i-1}}{i!}$$

And now the last step, from $A_0 = A \setminus \bigcup_{i=1}^n A_i$ we'll have the answer:

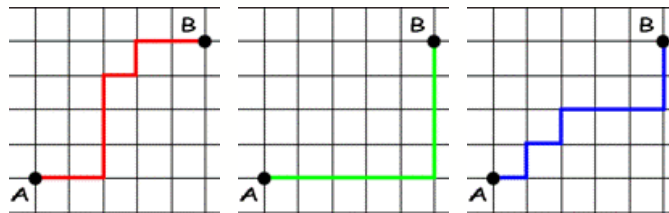
$$|A_0| = |A| - \left| \bigcup_{i=1}^n A_i \right| = n! - n! \sum_{i=1}^n \frac{(-1)^{i-1}}{i!} = n! \sum_{i=0}^n \frac{(-1)^i}{i!}.$$

And the last remark:

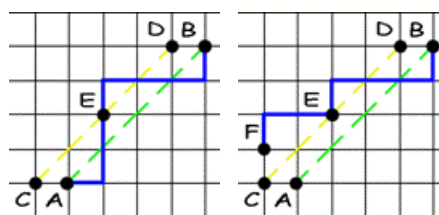
$$n! \sum_{i=0}^n \frac{(-1)^i}{i!} = \left\lfloor \frac{n!}{e} \right\rfloor.$$

This problem may not look very practical for use in TopCoder problems, but the thinking behind it is rather important, and these ideas can be

Another interesting method in combinatorics — and one of my favorites, because of its elegance — is called method of paths (or trajectories). The main idea is to find a geometrical interpretation for the problem in which we should calculate the number of paths of a special type. More precisely, if we have two points A, B on a plane with integer coordinates, then we will operate only with the shortest paths between A and B that pass only through the lines of the integer grid and that can be done only in horizontal or vertical movements with length equal to 1. For example:


$$\binom{n+m}{m} \text{ or } \binom{n+m}{n}.$$

Here's the main idea of the solution: Find the number of paths from A to B that cross segment AB, and call them "incorrect". If path is "incorrect" it has points on the segment CD, where $C = (0, 1)$, $D = (n-1, n)$. Let E be the point nearest to A that belongs to CD (and to the path). Let's symmetrize AE, part of our "incorrect" path with respect to the line CD. After this operation we'll get a path from $F = (-1, 1)$ to B.


$$\binom{2n}{n} - \binom{2n}{n-1} = \frac{(2n)!}{n!n!} - \frac{(2n)!}{(n-1)!(n+1)!} = \frac{\binom{2n}{n}}{n+1}$$

Recurrence relations probably deserves their own separate article, but I should mention that they play a great role in combinatorics. Particularly with regard to TopCoder, most calculation problems seem to require coders to use some recurrence relation and find the solution for the values

of parameters.

If you'd like to learn more, check out these tutorials: [An Introduction to Recursion](#), [Recursion, Part 2](#), and [Dynamic Programming: From novice to advanced](#). Done reading? Let's take a look at some examples.

[ChristmasTree](#) (SRM 331 Division Two – Level Three):

We'll use DP to solve this — it may not be the best way to tackle this problem, but it's the easiest to understand. Let $\text{cnt}[\text{lev}][r][g][b]$ be the number of possible ways to decorate the first lev levels of tree using r red, g green and b blue baubles. To make a recurrent step calculating $\text{cnt}[\text{lev}][r][g][b]$ we consider 3 variants:

“ $\text{cnt}[\text{lev}][r][g][b]=$

1) we fill the last level with one color (red, green or blue), so just:

$$= \text{cnt}[\text{lev}-1][r-\text{lev}][g][b] + \text{cnt}[\text{lev}-1][r][g-\text{lev}][b] + \text{cnt}[\text{lev}-1][r][g][b-\text{lev}] + ;$$

2) if $(\text{lev} \% 2 == 0)$ we fill the last level with two colors (red+green, green+blue or red+blue), then we calculate the number of possible decorations using the *Permutation with repetition* formula. We'll get $\frac{\text{lev}!}{(\text{lev}/2)!(\text{lev}/2)!}$ possible variants for each two colors, so just

$$\frac{\text{lev}!}{(\text{lev}/2)!(\text{lev}/2)!} (\text{cnt}[\text{lev}-1][r-\text{lev}/2][g-\text{lev}/2][b] + \dots + \text{cnt}[\text{lev}-1][r][g-\text{lev}/2][b-\text{lev}/2]) + ;$$

3) if $(\text{lev} \% 3 == 0)$ we fill the last level with three colors and, again using the *Permutation with repetition* formula, we'll get

$$\frac{\text{lev}!}{(\text{lev}/3)!(\text{lev}/3)!(\text{lev}/3)!} \text{possible variants, so we'll get:}$$
$$+ \frac{\text{lev}!}{(\text{lev}/3)!(\text{lev}/3)!(\text{lev}/3)!} \cdot \text{cnt}[\text{lev}-1][r-\text{lev}/3][g-\text{lev}/3][b-\text{lev}/3]$$

(all $\text{cnt}[i][j][k]$ with negative indices are 0).

[DiceGames](#) (SRM 349 Division One – Level Two):

First we should do some preparation for the main part of the solution, by sorting **sides** array in increasing order and calculating only the formations where the numbers on the dice go in non-decreasing order. This preparation saves us from calculating the same formations several times (see [SRM 349 – Problem Set & Analysis](#) for additional explanation). Now we will only need to build the recurrence relation, since the implementation is rather straightforward. Let $\text{ret}[i][j]$ be the number of different formations of the first i dice, with the last dice equal to j (so $i \in \overline{0, n-1}, j \in \overline{0, \text{sides}[i]-1}$, where n is the number of elements in **sides**). Now we can simply write the recurrence relation:

$$\text{ret}[i][j] = \sum_{k=0}^j \text{ret}[i-1][k], \text{ret}[0][i] = 1 (i \in \overline{0, \text{sides}[0]-1})$$

The answer will be $\sum_{i=0}^{\text{sides}[n-1]-1} \text{ret}[n-1][i]$.

Conclusion

As this article was written for novices in combinatorics, it focused mainly on the basic aspects and methods of counting. To learn more, I recommend you check out the reference works listed below, and keep practicing – both in TopCoder SRMs and pure mathematical problems. Good luck!

References:

1. Hall M. [“Combinatorial theory”](#)
2. Cameron P.J. [“Combinatorics: Topics, Techniques, Algorithms”](#)
3. [en.wikipedia.org :-\)](#)

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