

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$$

↓
○ ○ ○ ○ ○

$$\binom{5}{3} = \binom{4}{2} + \binom{4}{3}$$

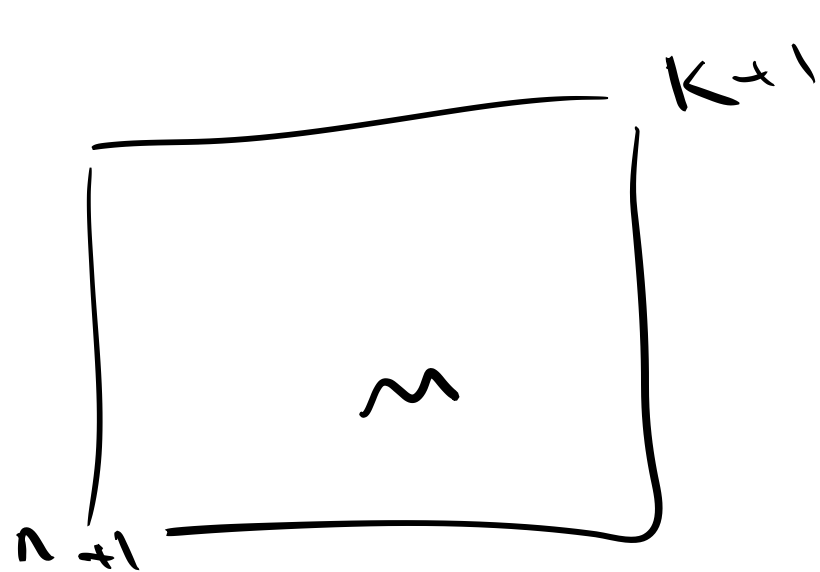
سب دو جگہ

def f(k, n):

if k==0 or k==n:

return 1

return f(k-1, n-1)
+ f(k, n-1)



def $f(n, k)$:

سخت، فزونی

$O(n^k)$

سه خطی $O(nk)$

for $i = 0 \rightarrow n$:

for $j = 0 \rightarrow k$:

if $j = 0$ or $j = k$:

$n[i, j] = 1$

else:

$n[i, j] = n[i-1, j-1] + n[i-1, j]$

return $n[i, j]$

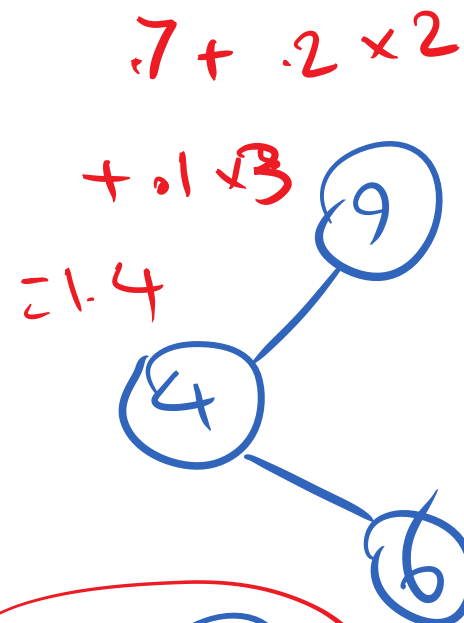
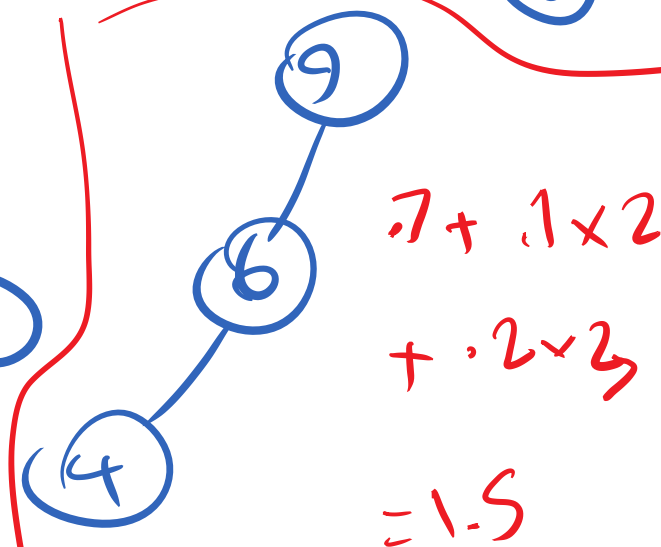
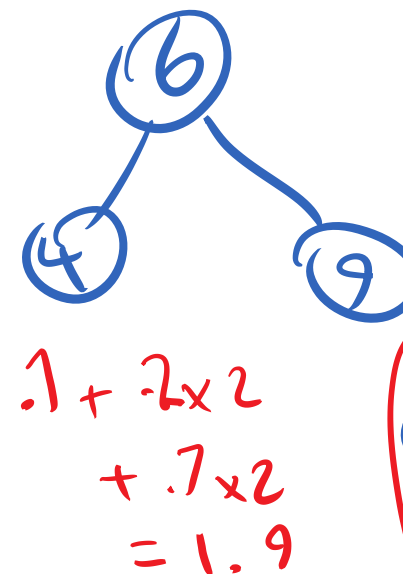
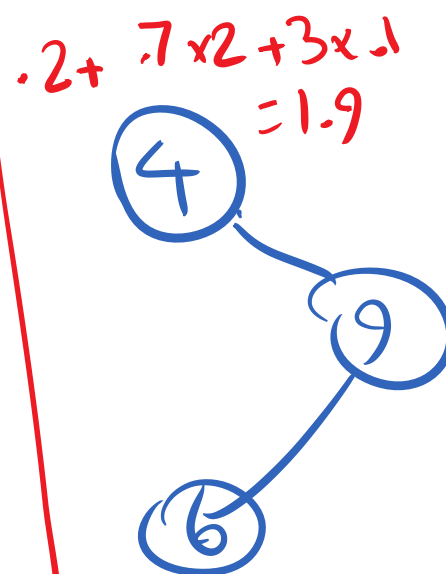
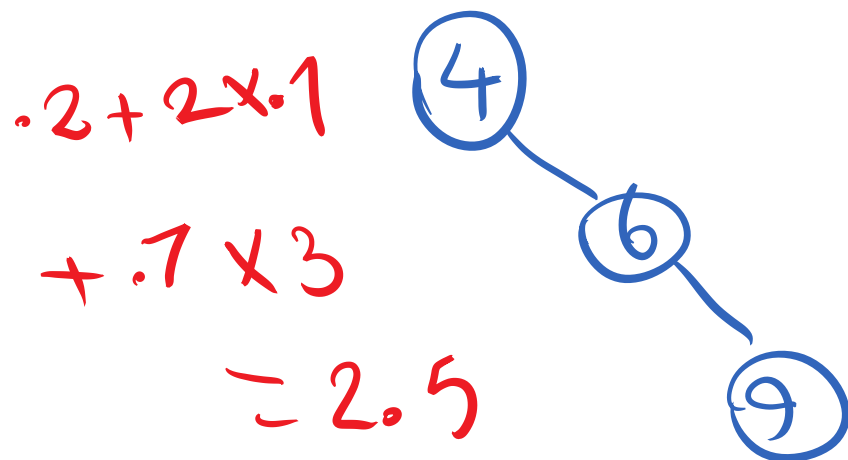
$$\binom{5}{2}$$

n/r	0	1	2
0	1		
1	1	1	
2	1	2	1
3	1	3	3
4	1	4	6
5	1	5	10

Optimal Binary Search Tree

.2 .1 .7
 ↑ ↑ ↑
4, 6, 9

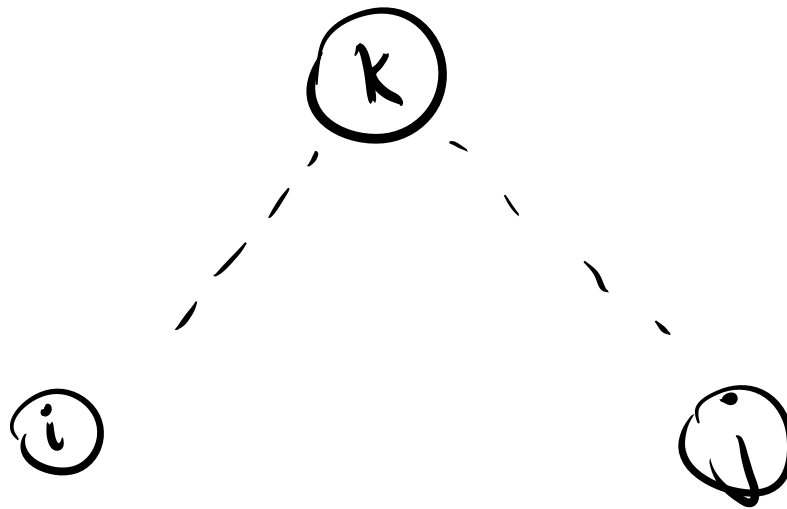
BST



P_1 P_2 P_3
 $.2$ $.1$ $.7$
 \uparrow \uparrow \uparrow
 $\underline{4}$, $\underline{6}$, $\underline{9}$

M_{ij}
 $i \leq k \leq j$

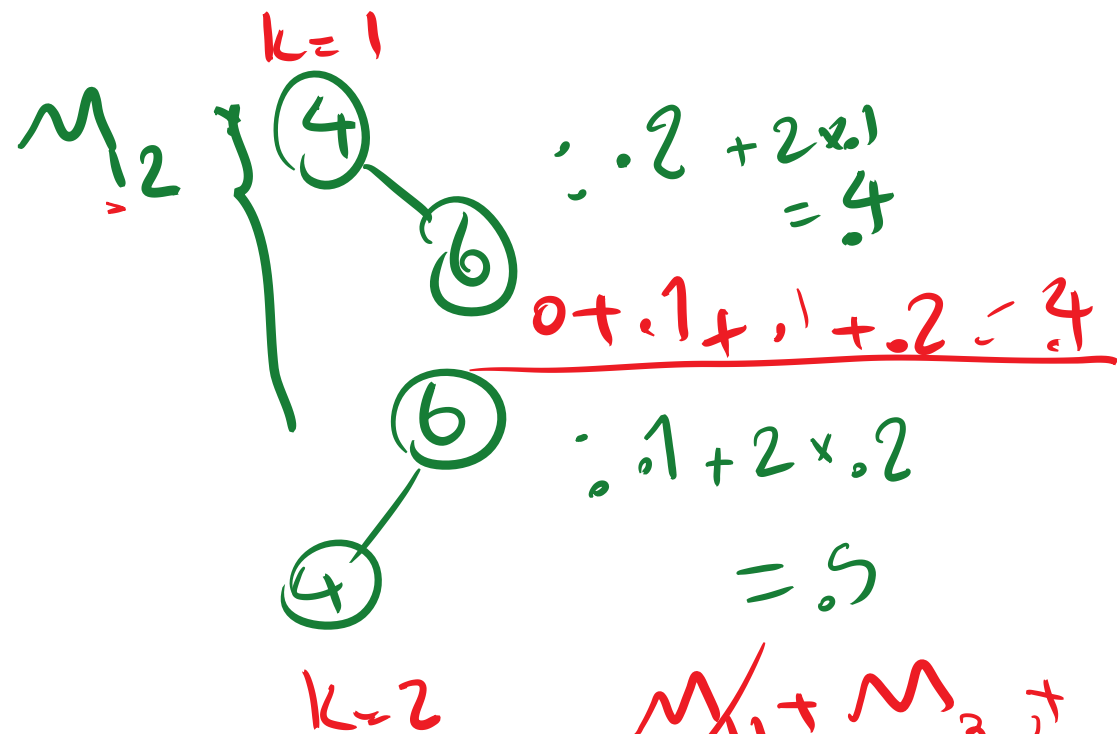
$M_{ii} = P_i$
 $M_{i, i+1}$



P_1 .2 P_2 .1 P_3 .7
 ↑ ↑ ↑
4, 6, 9

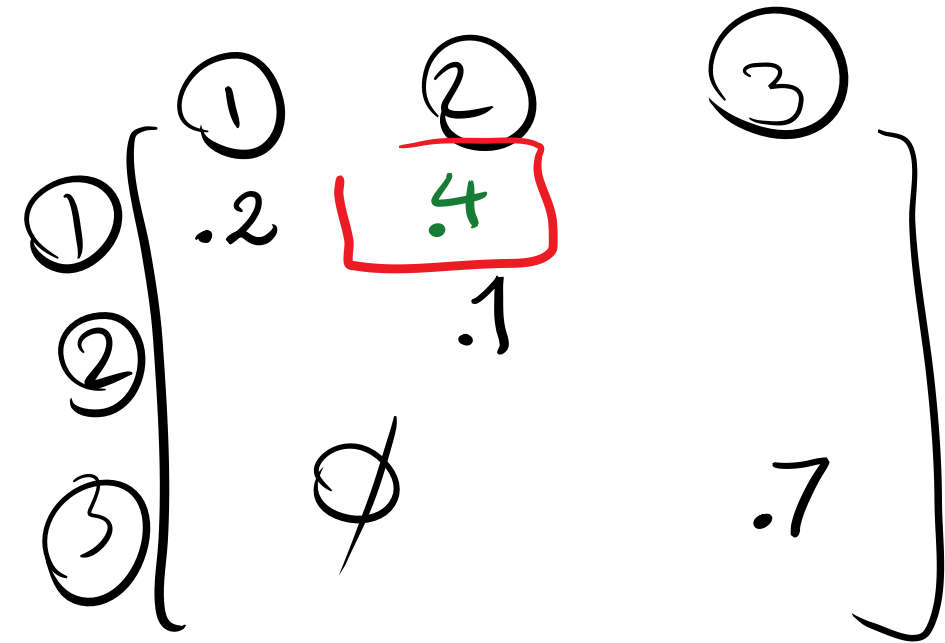
$$M_{ij} : \min_k \{ M_{ik-1} + M_{k+1}j + \sum_{l=i}^j P_l \}$$

$$i \leq k \leq j$$



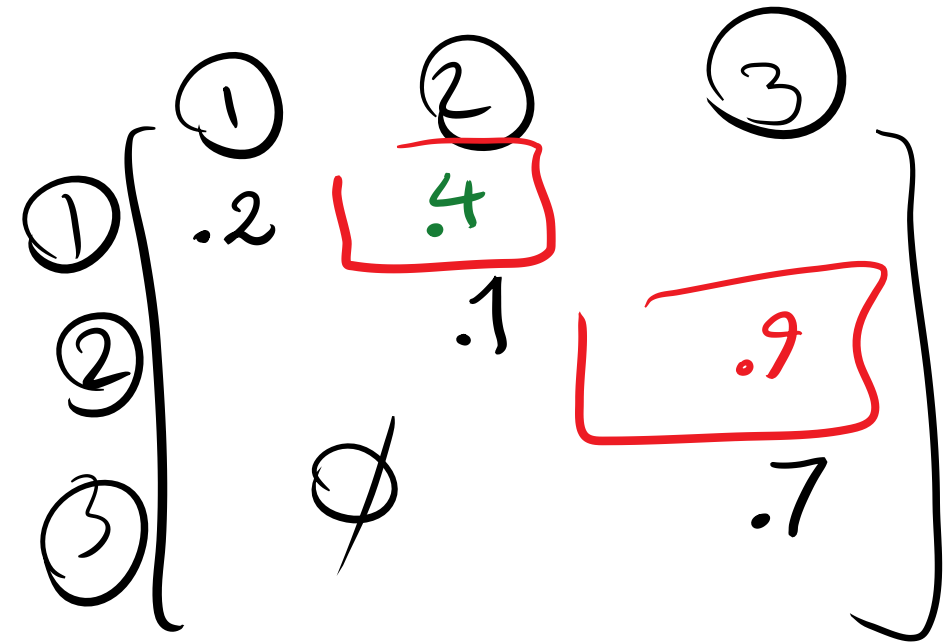
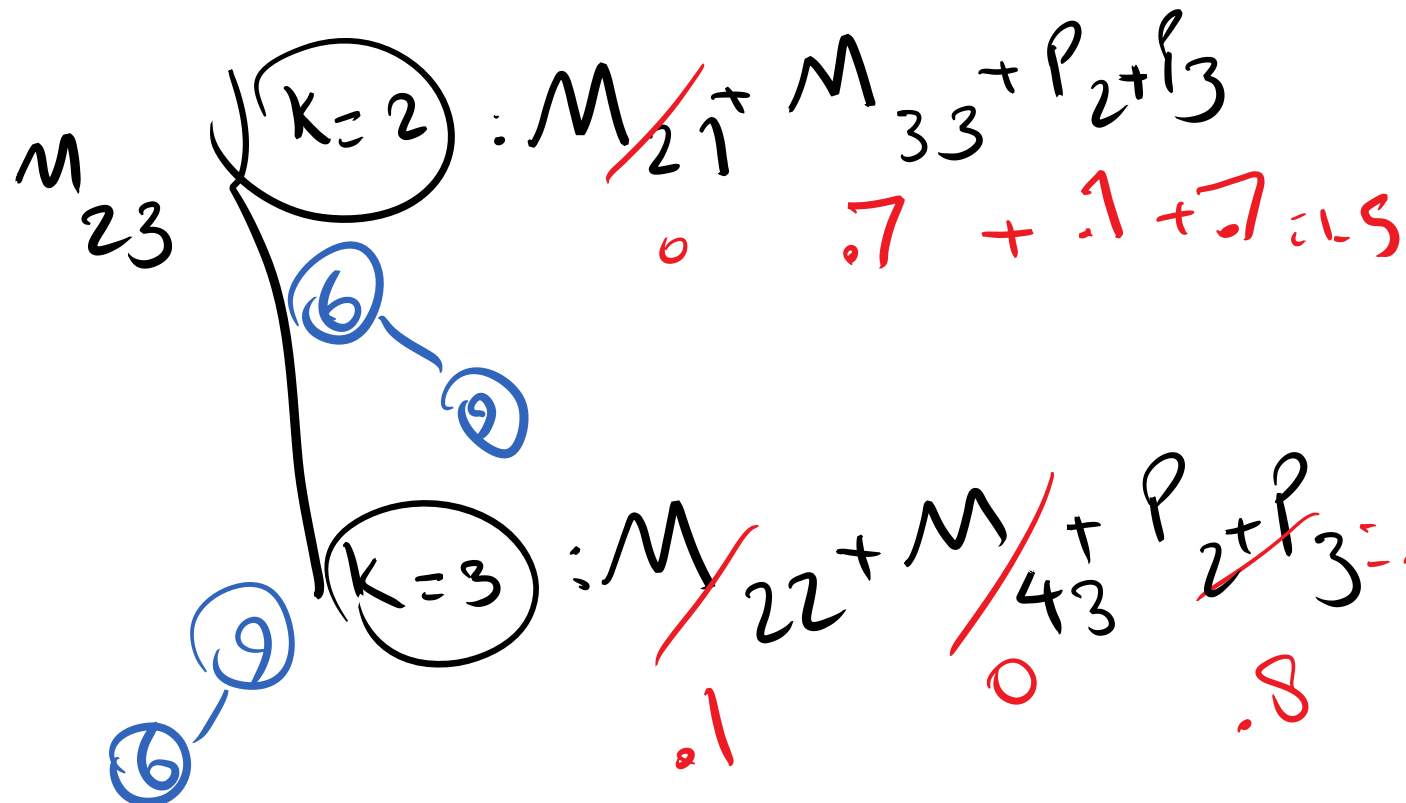
$$M_{1,1} + M_{3,2} + P_1 + P_2$$

$$.2 + 0 + .2 + .1$$



P_1 .2
 P_2 .1
 P_3 .7
 ↑ ↑ ↑
4, 6, 9

$M_{ij} : \min_k \{ M_{ik-1} + M_{k+1}j + \sum_{l=i}^j P_l \}$
 $i \leq k \leq j$



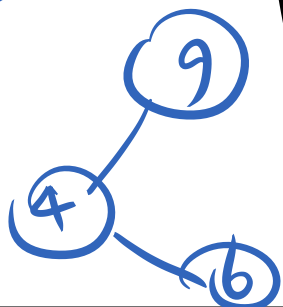
P_1 P_2 P_3
 .2 .1 .7
 ↑ ↑ ↑
4, 6, 9

$$M_{ij} = \min_k \left\{ M_{ik-1} + M_{k+1} + \sum_{l=i}^j p_l \right\}$$

$i \leq k \leq j$

$$M_{13} \left\{ \begin{array}{l} (k=1) : M_{10} + M_{23} + 1 = 1.9 \end{array} \right.$$

$$\boxed{K=2} \quad \cancel{M_{11}}_{-2} + \cancel{M_{33}}_{.7} + 1 = 1.9$$



$k=3$: $M^4_{12} + M^1_{43} = 1.4$

Handwritten diagram illustrating the insertion sort process on the list $[1, 2, 3]$:

- Initial list: $[1, 2, 3]$
- Step 1: Compare 1 and 2. Since $1 < 2$, no shift occurs. The list remains $[1, 2, 3]$.
- Step 2: Compare 1 and 3. Since $1 < 3$, no shift occurs. The list remains $[1, 2, 3]$.
- Step 3: Compare 1 and 2. Since $1 < 2$, no shift occurs. The list remains $[1, 2, 3]$.
- Final list: $[1, 2, 3]$

def OBST ($P[1..n]$):

$M = \text{Zeros}(n, n)$

for $i: 1 \rightarrow n$:

$M_{ii} = P_i$

for $s = 1$ to $n-1$:

for $i = 1$ to $n-s$:

$j = i+s$

$M_{ij} = \min_{i \leq k \leq j} \{ M_{ik-1} + M_{k+1,j} \} + \sum_{l=i}^j P_l$

M_{ij} → ϕ
 i, j

$T(n) = O(n^3)$: بهترین زمان

$O(n^2)$: کمتر حافظه

