Performance Modeling and Design of Computer Systems- Ch 2 Queueing Theory Terminology

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Introduction

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#### Overview

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### Queueing Theory Theory of Queues

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- The theory behind what happens when you have lots of jobs
- what makes queues appear and how to make them go away
- Queueing theory applies anywhere that queues come up
- Example
  - CPU uses a time-sharing scheduler to serve a queue of jobs waiting for CPU time
  - Router in a network serves a queue of packets waiting to be routed.
- Queueing theory is built on stochastic modeling and analysis
  - Model and analyze service demands of jobs and the interarrival times of jobs as random variables.

#### Goal of Queueing Theory 2 Goals

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- Predicting the system perfor- mance. Ex.
  - predicting mean delay or delay variability in service
  - number of jobs that will be in queue
  - mean number of servers being utilized
- Developing design of improved system
- Example
  - Can we build a better system from 1 slow discs or one faster disc
  - •

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- Consider a system of a single CPU that serves a queue of jobs in First-Come- First-Served (FCFS) order
- Assume any distribution for the arraival
  - $oldsymbol{\bullet}$   $\lambda = \operatorname{arrrival} \operatorname{rate} = \operatorname{No} \operatorname{of} \operatorname{jobs} \operatorname{arrives} \operatorname{per} \operatorname{second}$
  - $\bullet \ \mu = {\rm service} \ {\rm rate} = {\rm No} \ {\rm of} \ {\rm jobs} \ {\rm served} \ {\rm per} \ {\rm second}$
  - response time= time diff. bet. job arrives until it completes service
  - E[T] = mean response time  $\sum (x.P(x))$ . Can you serve customer within old time length?



- **Question**: What if  $\lambda$  is doubled?
  - ullet How the E[T] changes? increases or decreases?
  - If decreases, can using powerful cpu solve the problem?
  - How much powerful CPU is necessary



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• Consider the closed system of figure



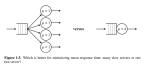
- Replace server 1 with a server that is twice as fast (the new server services jobs at an average rate of 2 jobs every 3 seconds).
  - Does this improvement affect the average response time in the system?
  - Does it affect the throughput?
  - Both cases improvement is no or negligible
- If the system is converted to open system? where arrival times are independent of service com- pletions.
  - Absolutely possible



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- one fast CPU of speed s,or n slow CPUs each of speed s/n (see Figure ).
- Goal is to minimize mean response time.



- one fast CPU of speed s,or n slow CPUs each of speed s/n (see Figure 1.5). Your goal is to minimize mean response time.
  - Choice depends on the variability of the job size
    - Question: when job size variability is high? Answer: we prefer many slow servers because we do not want short jobs getting stuck behind long ones.
    - Question: Which system do you prefer when load is low?
      Answer: When load is low, not all servers will be utilized, so it seems better to go with one fast server.

#### Power of Queueing Theory Design Example 3..cont..

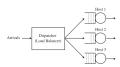
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- If jobs are preemptible; Jobs can be stopped and restarted where they left off.
- Question do you prefer many slow machines as compared to a single fast machine? Answer: If your jobs are preemptible, you could always use a single fast machine to simulate the effect of n slow machines.
- Resources can vary. CPU, GPU, MEMORY etc.
- Complexity and variation grows

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- Assume that all hosts are identical (homogeneous)
- all jobs only use a single resource.
- once jobs are assigned to a host, they are processed there in FCFS order and are non-preemptible.
- Which task assignment policies yields the lowest mean response time?
- Some options Random, Round-Robin,
  Size-Interval-Task-Assignment (SITA), Central- Queue:
- More possible. Answer depends on various parameters
  - If job size variability is low, then the LWL policy is best.
  - If job size variability is high, then it is important to keep short jobs from getting stuck behind long ones, so a SITA-like policy,

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- A single server. Jobs arrive according to a Poisson process.
  Arbitrary Job size
- Question: Which Scheduling policy is best w.r.t. mean response time if non- preemptive service orders?
- Scheduling policies are: First-Come-First-Served (FCFS),
  Non-Preemptive Last-Come- First-Served (LCFS), etc
- Ans: All are same for non-preemptive service orders
- now if Non Preemptive-LCFS policy (PLCFS)-Whenever a new arrival enters the system, it immediately preempts the job in service
- Mean resposne time depends on the variability of the job size distribution.
  - job size distribution is at least moderately variable, then PLCFS will be a huge improvement.

If the job size distribution is hardly variable (basically constant), then PLCFS policy will be up to a factor of 2