A Tangram Puzzle Solver in Common Lisp

Michael Wessel

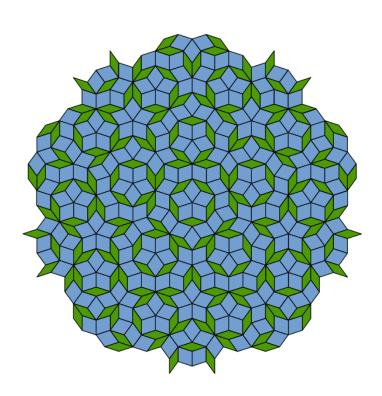


May 3rd 2021 ELS'21

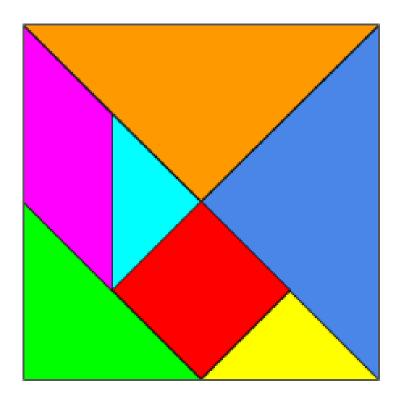
Talk Outline

- Disclaimer: this is a demo talk
 - Video of program in action & some Lisp code
 - No deep theoretical insights, but some practical insights
 - This is a fun recreational spare time project
 - Not associated with my employer in any way
- The program was written in 2003, reactivated now
 - To get back to Lisp Q: Would I still be able to maintain and extent my Lisp code from almost 20 years ago?
 - To **finish what was started** the original program had some deficiencies & bugs
 - To catch up with the Lisp community
 (paper motivation last time I contributed was 15 years ago)
 - Because of a recent book that really fascinated me (see at the end)

Tiling Problems are Fascinating! Penrose Tiling & Tangram Set



Aperiodic Penrose Tiling (Extends Infinitely) Source: Wikipedia CC



Tangram Tiling (Finite) Source: Wikipedia CC







Dog Pattern | Schnauzer | M. C. E... pinterest.com



How to make an Escher tiling static1.squarespace.com



Escher tessellations ... pinterest.com



Tessellation patterns ... pinterest.com



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Tessellation Patterns - From ... widewalls.ch



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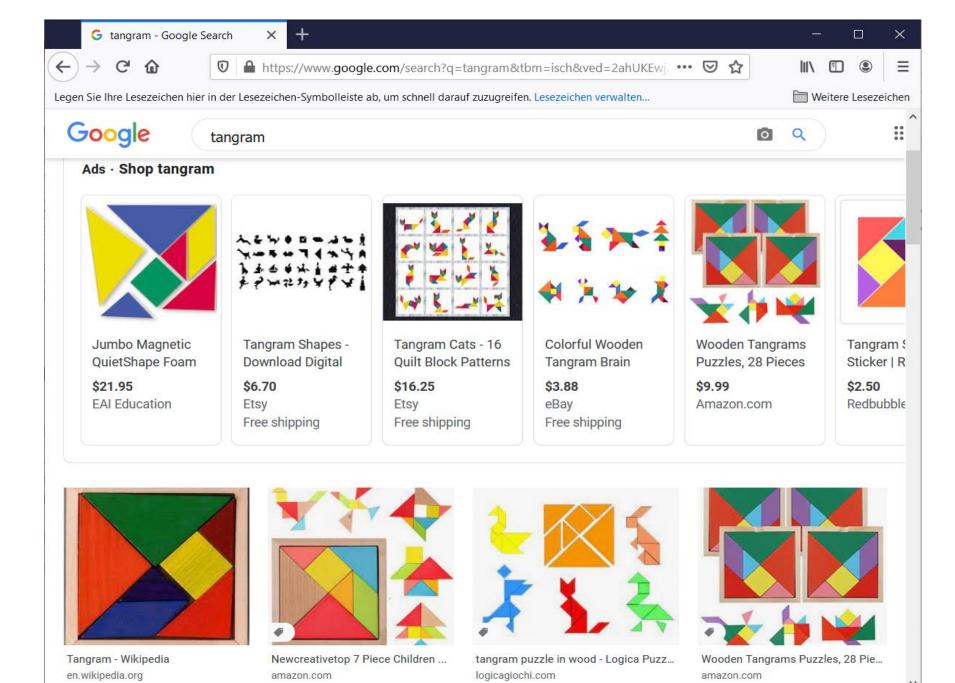
Pin on tessellations pinterest.com



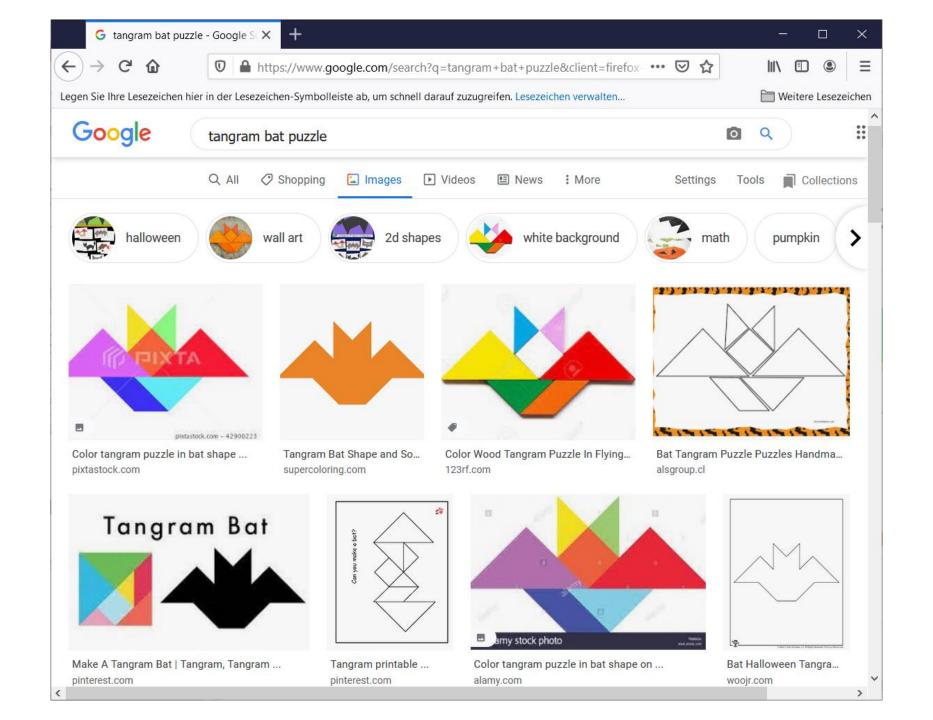
mc escher ... pinterest.com



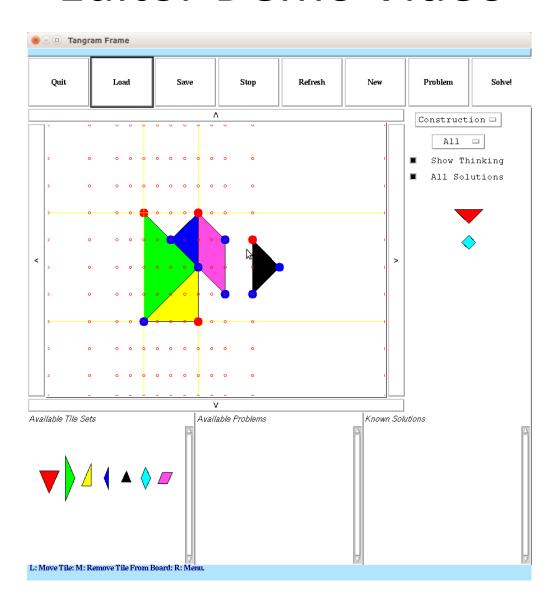
330 Tessellations ideas | art... pinterest.com



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Editor Demo Video

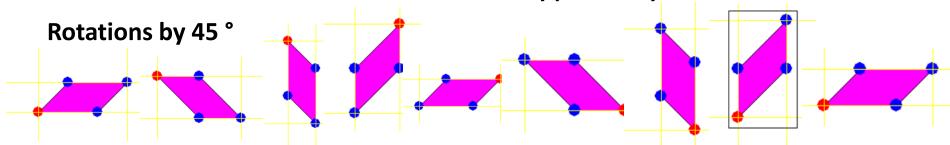


Rules of the Tangram Game

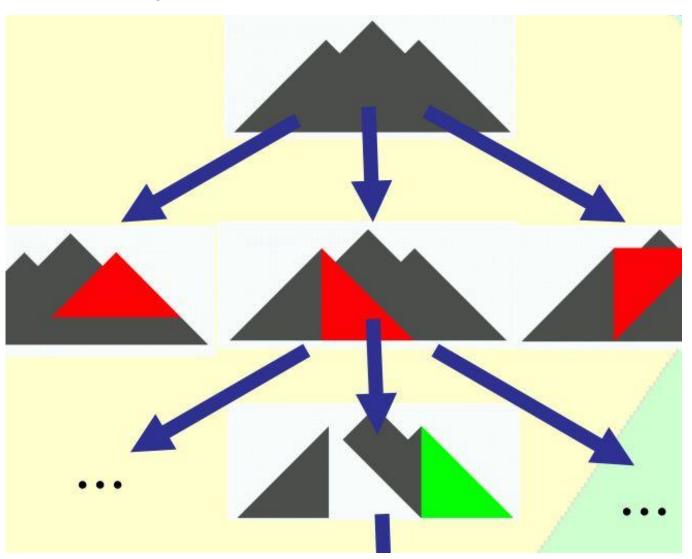
reflected

variant of

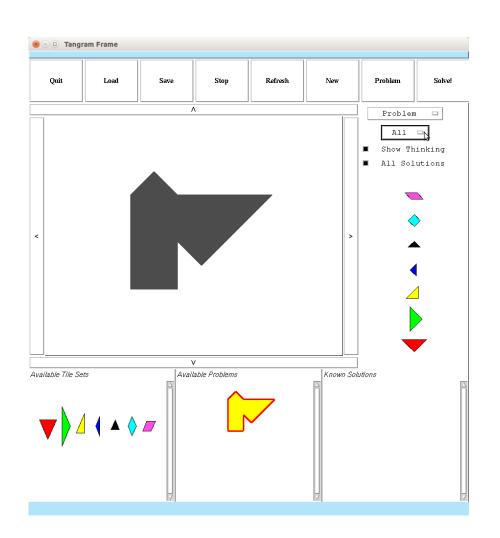
- Each tile must be used EXACTLY once
- The silhoutte / outline polygon must be partioned completely
 - no remaining space
 - no overlaps between tiles
- Tiles can be translated, rotated and reflected (flipped over)
- The reflected variants of each tile (modulo congruency) constitute a **TileType**; only ONE reflected variant of each TileType may be used

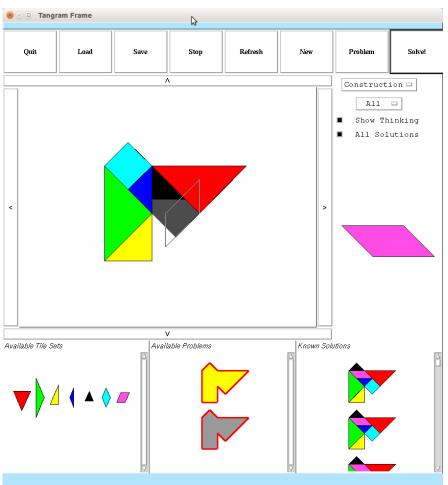


Solving Tiling Problems by Geometric Search



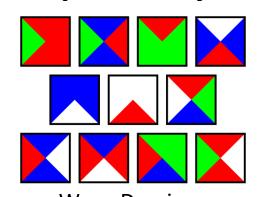
Solver Demo Video



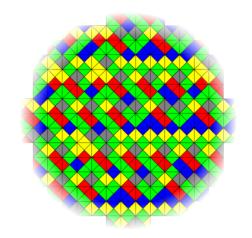


Comments on Problem Complexity

- It is undecidable if an arbirtary set of (Wang domino) tiles tiles the (infinte) plane
 - reduction to Halting Problem of Turing Machines
- Wang 1961:
 - Conjecture: if a set of finite tiles tiles the (infinite) plane, then there also exists a periodic tiling
 - If the Conjecture is true, then it implies the existence of a decision algorithm
- Berger 1966 (Wang's student):
 - That algorithm does not exist hence, the conjecture is wrong
 - What's wrong with it?
 There also exist aperiodic tilings
 - See 11 Wang Dominos (Jeandel and Rao 2015)
 - Penrose tilings
- Tangram is NP-hard
- Certain variants are NP-complete (e.g., symmetric puzzles)



Wang Dominos
(Aperiodic)
Source: Wikipedia CC



Wang Tesselation with 13 Tiles
Source: Wikipedia CC

Search Function – 1 / 3

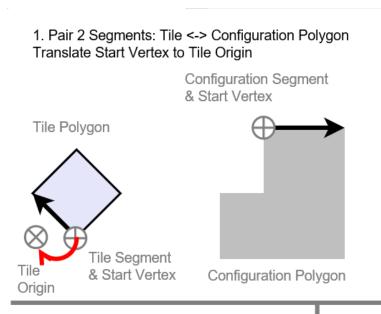
```
(defun find-covering (polygon-set tile-types &key (remove-used-tiles-p t)
                                  frame
                                  all-solutions-p
                                  allow-less-tiles-p
                                  show-thinking-fn)
 (let ((solutions nil))
    (labels ((goal-p (polygons tile-types)
               (and (not polygons)
                    (or allow-less-tiles-p
                        (not tile-types))))
             (do-it (polygons tile-types history)
               (incf *inc-counter*)
               (cond ((goal-p polygons tile-types)
                      (push
                       (list polygons tile-types history)
                       solutions)
                      (when show-thinking-fn
                        (funcall show-thinking-fn polygons tile-types history t))
                      (unless all-solutions-p
                        (return-from find-covering solutions)))
```

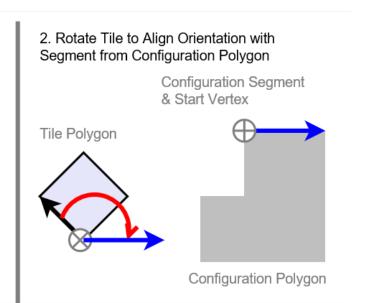
Search Function – 2 / 3

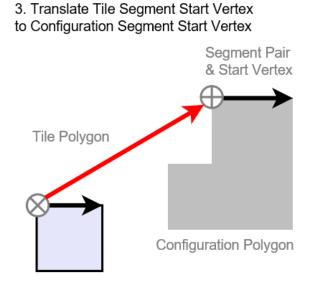
```
(t
(when show-thinking-fn (funcall show-thinking-fn polygons tile-types history nil))
#+:ignore
 (sleep 1)
 (let* ((alignments nil)
        (perfect-alignments nil))
   (dolist (polygon polygons)
     (dolist (tile-type tile-types)
       ;; tile-type is the set of reflected variants of that tile
       (let ((rem-tile-types
              (if remove-used-tiles-p
                  (remove tile-type tile-types)
                tile-types)))
         (dolist (tile tile-type)
           (dolist (succ-conf
                    (find-possible-alignments polygon (shape tile)
                                                   :frame frame
                                                   :color (color tile)
                                                   :score-p t
                                                   :rem-tile-types rem-tile-types))
             ;; succ-conf = (list h aligned-poly succ-conf orig-poly rem-tile-types)
             (if (eg (third succ-conf) :match)
                 (push (list succ-conf tile ) perfect-alignments)
               (push (list succ-conf tile ) alignments)))))))
  (let ((sorted (append (sort perfect-alignments #'tuple-> :kev #'caar)
                          (subseq (sort alignments #'tuple-> :key #'caar)
                                 ;; best n only ?
                                 0 (length alignments)))))
```

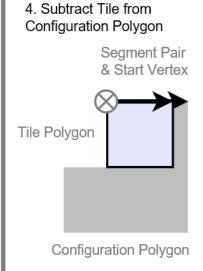
Search Function – 3 / 3

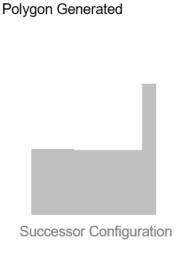
Successor Generator Function









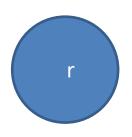


5. New Configuration

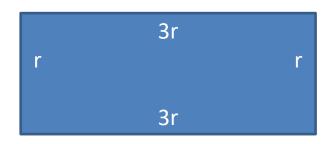
Geometric Heuristics

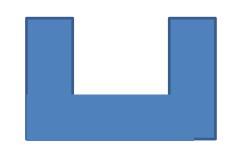
- Tile area (biggest first)
- Compactness of successor polygon P

$$compactness(P) = \frac{circumference(P)}{\sqrt{area(P)}}$$





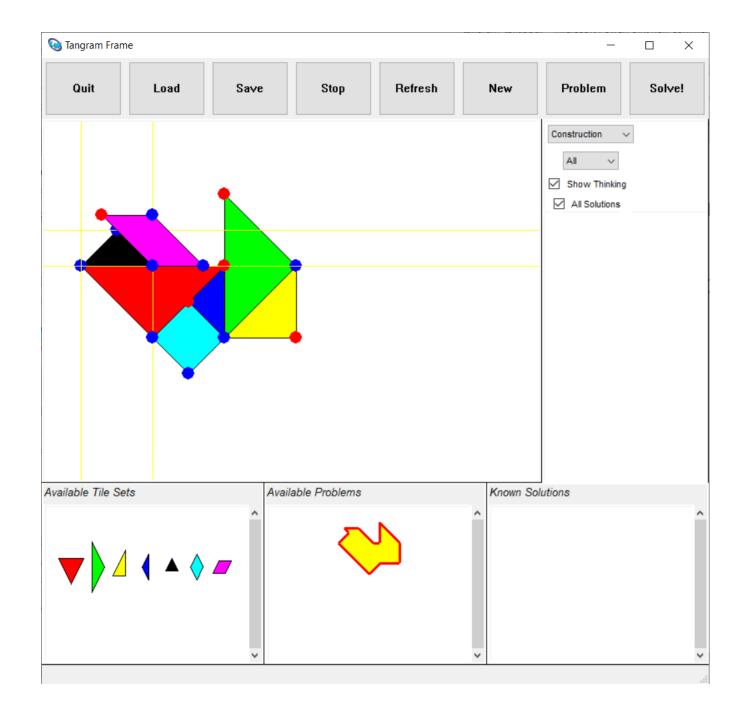


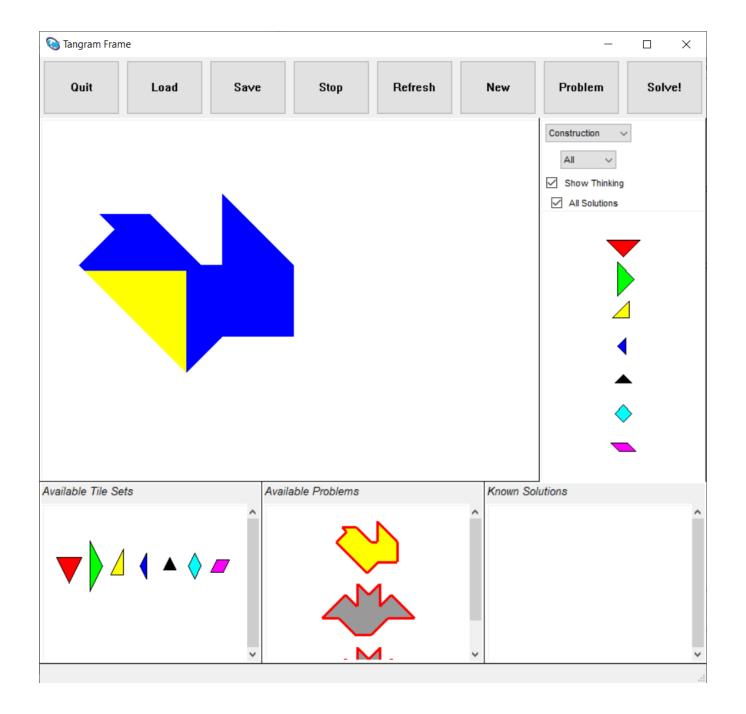


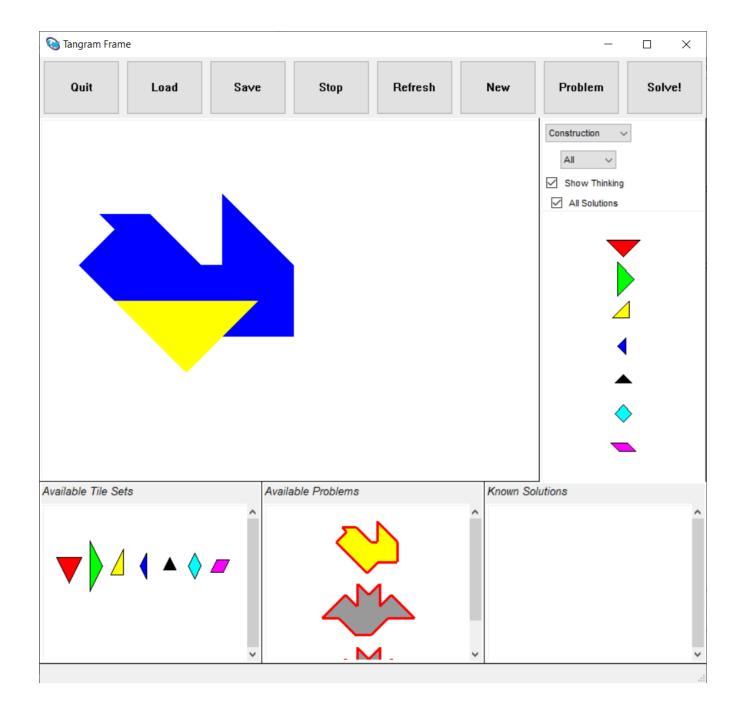
$$\frac{2\pi r}{\sqrt{\pi r^2}} = 2\sqrt{\pi} \sim 3.5 \qquad \frac{4r}{\sqrt{r^2}} = 4$$

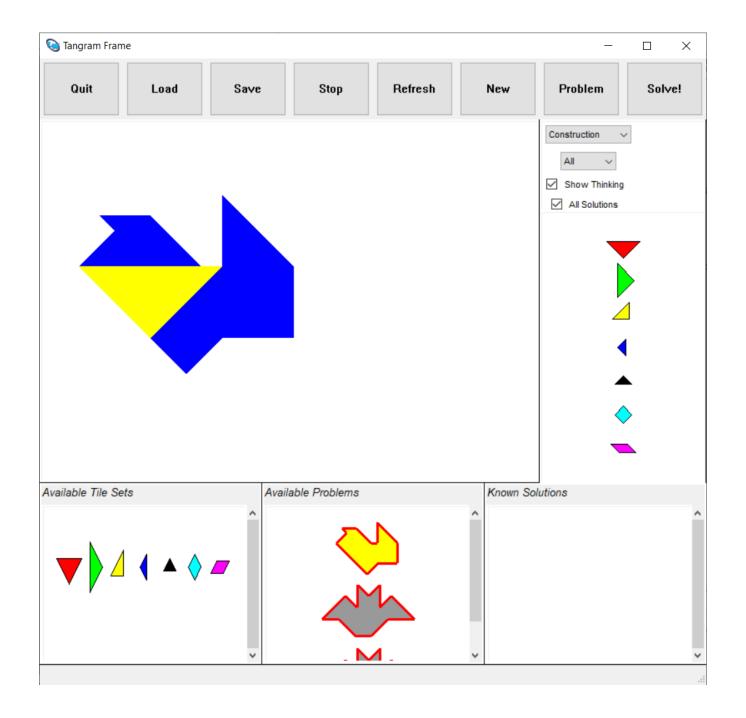
$$\frac{8r}{\sqrt{3r^2}} \sim 4.6$$

$$\frac{20}{\sqrt{14}} \sim 5.3$$









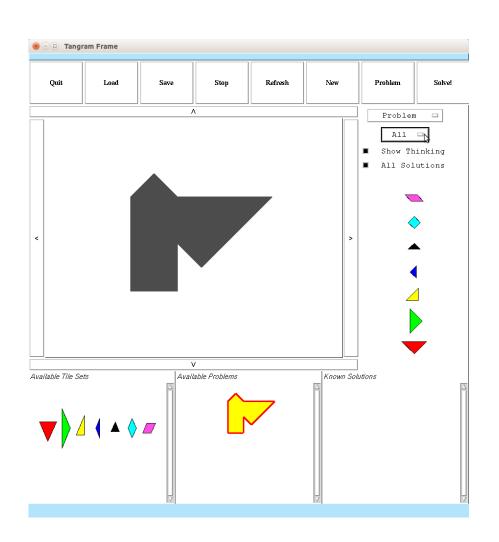
Geometric Heuristics Continued

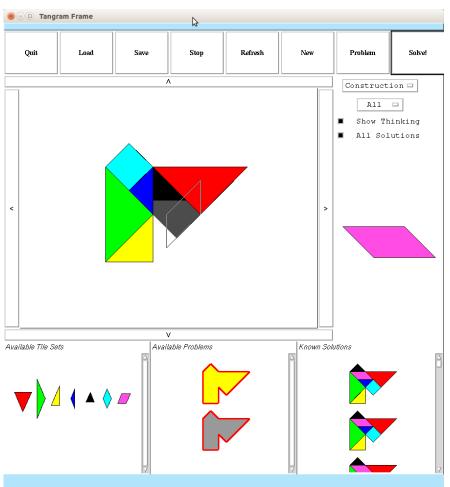
- How to compare two successor configurations wrt. different criteria?
 - Comparing Apples with Pears
 - Combine into one score
 - Weighted sum
 - Difficult, not normalized
 - Use ordered list of criteria instead lexicographic ordering (see tuple->)
 - Score = (tile size, -compactness, #interior vertices, ...)

Scoring Function

```
(defun alignment-result-heuristic-for-polygon (aligned-polygon conf orig-polygon rem-tile-types)
  (cond ((symbolp conf) ; :match!
         (list (tangram-geometry:calculate-area aligned-polygon) 0))
        ((some #'(lambda (poly)
                   (every #'(lambda (tile-type)
                              (let ((tile (first tile-type))); use first reflected variant
                                (< (+ (tangram-geometry:calculate-area poly) 2)</p>
                                   :: account for inaccuracies ...
                                    (tangram-geometry:calculate-area (shape tile)))))
                          rem-tile-types))
               conf)
         : bad )
        (t
        ;; score list for lexicographic tuple->-p sorting
         (list
          (tangram-geometry:calculate-area aligned-polygon)
          (- (count-if #'(lambda (p)
                           (inside-p p orig-polygon))
                       (point-list aligned-polygon)))
          (apply #'min (mapcar #'(lambda (x) (- (compactness x))) conf))
          (- (reduce #'+ (mapcar #'length (mapcar #'segments conf))))))))
```

Demo Video Heuristics





Further Essential Tricks

- Remove congruent successor configurations
 - Congruent configurations (the same configuration being generated from different input) occur frequently
 - Tiles have rotational symmetries
 - There are congruent tile types
 - Don't explore the same search space regions repeatedly
 - Reduce branching factor / search space
- Remove configurations that have polygons which are smaller than the smallest (remaining) tile
 - Backtrack early
 - Avoid trashing

Performance & Heuristics

Problem	Optimizations A & B	Optimizations A & ! B	Optimizations ! A & ! B
Standard Problem Solution Found at Branching Factor First 3 Levels	~ 3 mins 2705 104, 76, 44	~ 6 mins 4230 104, 76, 44	~ 27 mins 22788 148, 112, 65
Bat Problem	~ 1 min	~ 2.5 mins	~ 45 mins (!)
Other Test Problem	~ 10 seconds		

- Without heuristics, none of these problems can be solved
- ~15 Configurations / Second
- These additional heuristics are hightly effective as well
 Optimization A Remove congruent successor configurations
 Optimization B Remove successor configuration with polygons that are too small for any tile to cover / fit in

About Lisp -1/3

- CLOS is used extensively
 - Geometric classes
 - Compositional hierarchy
 - Segments from Points, Polygons from Segments,

• • •

- Dynamically affected by transformation matrix
- Reader methods for x, y

=> all geometric functions are
 automatically aware of active transformations
 (e.g., calculate-rcc-relation next slide)

About Lisp -2/3

Transformations similar to CLIM

```
(setf (affected-by-matrix-p a) nil)
(setf (affected-by-matrix-p b) t)
(with-translation ( (- xf) (- yf) )
  (with-rotation ( phi )
    (with-translation ( xt yt )
      (when (and frame (show-thinking-p frame))
        (show-polygons (list b)
                       :clear-p nil
                       :inks (list +flipping-ink+)
                       :filled-p nil))
      (let* ((rel (calculate-rcc-relation a b)))
```

About Lisp -3/3

Multi-Methods are nice for Case Analysis

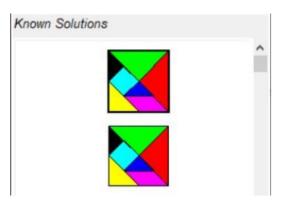
```
(defmethod touches-p ((line geom-line) (poly geom-polygon))
  (and (intersects-p line poly)
       (or (lies-on-p line poly)
           (not (one-part-is-inside-p line poly)))))
(defmethod touches-p ((poly geom-polygon) (line geom-line))
 (touches-p line poly))
(defmethod touches-p ((chain1 geom-chain) (chain2 geom-chain))
 ;;; (error "TOUCHES-P ~A ~A: Not well-defined!" chain1 chain2))
(defmethod-memo touches-p ((obj geom-chain) (poly geom-polygon))
    ((obj poly (trafo-id *matrix*)))
 (and (intersects-p obj poly)
       (every #'(lambda (segment)
                  (=> (intersects-p segment poly)
                      (touches-p segment poly)))
              (segments obi))))
(defmethod touches-p ((poly geom-polygon) (obj geom-chain))
  (touches-p obj polv))
(defmethod-memo touches-p ((poly1 geom-polygon) (poly2 geom-polygon))
    ((poly1 poly2 (trafo-id *matrix*)))
 (and (not (eq poly1 poly2))
       (intersects-p poly1 poly2)
       (every #'(lambda (segment)
                  (=> (intersects-p segment poly2)
                      (touches-p segment poly2)))
              (segments poly1))))
```

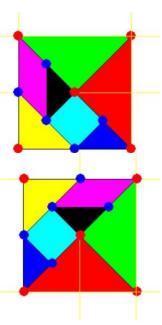
About CLIM

- CLIM is great for building prototypes quickly
- CLOS classes are presentations
 - Active objects that are alive on the screen (highlighting)
 - Automatically generated context menus etc.
 - You get A LOT of functionallity with minimal effort

Observations

- Frequently finds interesting solutions
- Is it **superhuman**? I guess it depends:
 - Search frequently is, not only since DNNs (e.g., Checkers, Chess / DeepBlue, ...)
 - Compared to my tangram puzzle abilities, yes!
 - Not sure how fast people are
- Are other programs faster?
 - There are few available online... and they appeared only recently
 - Most of them are special-purpose and Tangram-specific
 - Whereas ours is generic and can (in principle) solve arbirtary tiling problems
 - Our program is complete (probably not the case for some of the others)
 - Sometimes, completeness has to be sacrificied for performance
- **Show Thinking** checkbox is useful for checking the progress
 - However, 60 % slower
 - The heuristics were devised by watching the search program and giving the program some "common sense" and human insight





Wrap Up & Future Work

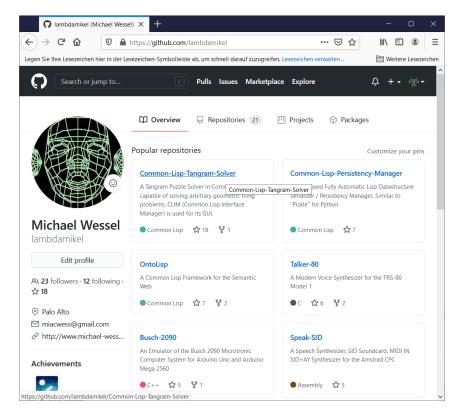
- Q: Would I still be able to maintain and extent my Lisp code from almost 20 years ago?
- A: YES
 - Lisp was fun to work with again after many years of absence
 - Lisp has good code maintainability
 - Concise, precise, semi-formal, small computational units (functions & methods)
 - The program is running much better than in 2003;
 I feel that, for now, I finished what was started
- The search space is very large
 - Getting the initial positions of the first two tiles right is essential
 - Else backtracking to them won't happen for a long time
 - Without heuristics no solutions can be found within 24 hours
- Future Work
 - SBCL / McCLIM, Penrose Tilings, ...
 - More optimizations (Random Restarts, ...)

Some Resources — Parens for You!

Github repositories

https://github.com/lambdamikel

- Geometric substrate
- Persistence module
- Some reusable functions there for you
- Sourcecode and executables for Tangram
 - Try tweaking it, change heuristics!
 It's easy and inutitive
- Uses Lispworks (6.1) + CLIM
 - Linux, Windows
 - M1 Silicon Mac –Lispworks 7.1
 - Thanks to Rainer Joswig (Lispm)!



"A rollercoaster adventure.... A nail-biting narrative [of] the thrill of scientific discovery." —NATURE

THE SECOND KIND OF IMPOSSIBLE

THE EXTRAORDINARY QUEST FOR A NEW FORM OF MATTER



STEINHARDT

Thank You!

Questions?